# Cell Broadband Engine

# Programming Handbook

*Including the PowerXCell 8i Processor*

## Version 1.11

May 12, 2008


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# Preface

This handbook describes the extensive programming facilities of the Cell Broadband Engine (Cell/B.E.) and IBM PowerXCell™ 8i processors, which are collectively called Cell Broadband Engine Architecture processors (CBEA processors). The Cell Broadband Engine Architecture (CBEA) was developed jointly by Sony, Toshiba, and IBM; it extends the 64-bit *PowerPC Architecture*™. The Cell/B.E. processor is a result of that collaboration between Sony, Toshiba, and IBM. It is the first implementation of a multiprocessor family conforming to the CBEA. The PowerXCell 8i processor also conforms to the CBEA. It provides a double data rate 2 (DDR2) memory interface and improved double-precision, floating-point performance and additional double-precision instructions. Both CBEA processors comprise eight Synergistic Processor Elements (SPEs) and one PowerPC Processor Element (PPE).

The handbook contains extended content targeted for programmers who are interested in developing user applications, libraries, drivers, middleware, compilers, or operating systems for CBEA processors. The sections describe all of the facilities unique to the CBEA processors that are needed to develop such programs. In general, the sections contain the following topics:

- General hardware and programming-environment overview: Sections 1 through 3.
- Additional hardware overview plus privilege-state (supervisor) programming: Sections 4 through 16. A few user-programming topics, such as decrementers, are also located in these sections.
- Problem-state (user) programming: Sections 17 through 24, and appendixes.

The handbook assumes that the reader is an experienced C/C++ programmer and familiar with basic concepts of single-instruction, multiple-data (SIMD) vector instruction sets such as the IBM PowerPC® Architecture vector/SIMD multimedia extensions, AltiVec, Intel® MMX, SSE, 3DNOW!, x86-64, or VIS instruction sets. The handbook is system-independent; it makes no assumptions about development-tool or operating-system environments.

The handbook provides adequate detail for general software development. Additional specific implementation details might be available under nondisclosure agreement from either Sony, Toshiba, or IBM. System-specific details should be available from the specific system supplier.

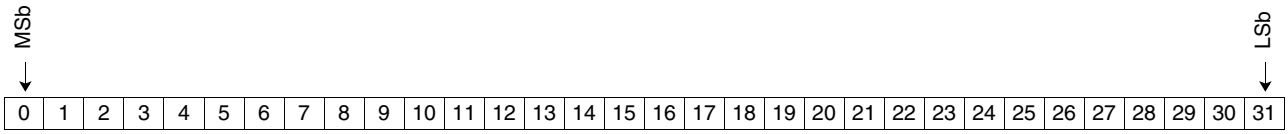
## Related Publications

| Title | Version | Date |
|---|---|---|
| *Cell Broadband Engine Architecture* | 1.02 | October 2007 |
| *PowerPC User Instruction Set Architecture, Book I* | 2.02 | January 28, 2005 |
| *PowerPC Virtual Environment Architecture, Book II* | 2.02 | January 28, 2005 |
| *PowerPC Operating Environment Architecture, Book III* | 2.02 | January 28, 2005 |
| *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors* | 2.0 | June 10, 2003 |
| *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual* | 2.07c | October 2006 |
| *Synergistic Processor Unit Instruction Set Architecture* | 1.2 | January 2007 |
| *C/C++ Language Extensions for Cell Broadband Engine Architecture* | 2.5 | February 2008 |

| Title | Version | Date |
|---|---|---|
| *SPU Application Binary Interface Specification* | 1.8 | September 2007 |
| *SPU Assembly Language Specification* | 1.6 | September 2007 |
| *Cell Broadband Engine Registers* | 1.6 | June 2007 |
| *PowerPC Compiler Writer's Guide* | 1.0 | 1996 |
| *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide* | 1.01 | June 2007 |
| *Linux™ Standard Base Core Specification for PPC 3.0* (http://www.linuxbase.org/spec) | 3.0 | 2004 |
| *64-bit PowerPC ELF Application Binary Interface Supplement* | 1.7.1 | July 2004 |
| *Cell Broadband Engine Linux Reference Implementation Application Binary Interface Specification*. | 1.0 | November 2005 |
| *AltiVec Technology Programming Interface Manual* | 0 | June 1999 |
| *IBM Research Hypervisor (rHype)* (This is available in source form with a general public license (GPL) at the following URL: http://www.research.ibm.com/hypervisor/) | 1.3 | March 2005 |
| *Tool Interface Standard (TIS) Executable and Linking Format (ELF)* specification. TIS Committee. | 1.2 | May 1995 |
| *SYSTEM V APPLICATION BINARY INTERFACE PowerPC Processor Supplement* (this document specifies the 32-bit application binary interface [ABI]). | | September 1995 |
| *Advanced Encryption Standard (AES)*, Federal Information Processing Standards Publications, FIPS PUB 197. | | 2001 |
| *Logical Partitions on IBM PowerPC: A Guide to Working with LPAR on Power5 i5 Servers* (This paper can be downloaded from the following Web site: http://www.redbooks.ibm.com/redbooks.nsf/RedpieceAbstracts/sg248000.html?Open) | | February 2005 |
| *Partitioning Implementations for IBM eServer™ p5 Servers* (This paper can be downloaded from the following Web site: http://www.redbooks.ibm.com/abstracts/sg247039.html) | | February 2005 |
| *Advanced Power Virtualization on IBM eServer p5 Servers* (This paper can be downloaded from the following Web site: http://www.redbooks.ibm.com/abstracts/sg247940.html?Open) | | December 2005 |

## Conventions and Notation

In this document, standard IBM notation is used, meaning that bits and bytes are numbered in ascending order from left to right. Thus, for a 4-byte word, bit 0 is the most-significant (high) bit, and bit 31 is the least-significant (low) bit.




Notation for bit encoding is as follows:

- Hexadecimal values are preceded by x and enclosed in single quotation marks. For example: x'0A00'.

- Binary values in sentences appear in single quotation marks. For example: ‘1010’.

The following software documentation conventions are used in this manual:

- Command (or instruction) names are written in **bold** type. For example: **put**, **get, rdch, wrch, rchcnt**.
- Variables are written in italic type. Required parameters are enclosed in angle brackets. Optional parameters are enclosed in brackets. For example: **get*<f,b>[s]***.
- The notation <***f,b***> indicates that either the tag-specific fence or tag-specific barrier form is available for a referenced memory flow controller (MFC) command.

The following symbols are used in this document:

| | |
|---|---|
| & | bitwise AND |
| \| | bitwise OR |
| % | modulus |
| = | equal to |
| ! = | not equal to |
| x ≥ | greater than or equal to |
| x ≤ | less than or equal to |
| x >> y | shift to the right; for example, 6 >> 2 = 1; least-significant y-bits are dropped |
| x << y | shift to the left; for example, 3 << 2 = 12; least-significant y-bits are replaced by zeros |

## Referencing Registers, Fields, and Bit Ranges

Registers are referred to by their full name or by their short name (also called the register mnemonic). Fields within registers are referred to by their full field name or by their field name. The field name or names are enclosed in brackets [ ].The following table describes how registers, fields, and bit ranges are referred to in this document and provides examples of the references.

| Type of Reference | Format | Example |
|---|---|---|
| Reference to a specific register and a specific field using the register short name and the field names, bit numbers, or bit range. | Register_Short_Name[FieldName] | MSR[FE0] |
| | Register_Short_Name[Bit_Number] | MSR[52] |
| | Register_Short_Name[Field_Name1, Field_Name2] | MSR[FE0, FE1] |
| | Register_Short_Name[Bit_Number, Bit_Number] | MSR[52, 55] |
| | Register_Short_Name[Starting_Bit_Number:Ending_Bit_Number] | MSR[39:44] |
| **Note:** The register short name is also called the register mnemonic. | | |

| Type of Reference | Format | Example |
|---|---|---|
| Reference to a specific register and the setting for a specific field, bit, or range of bits using the register short name and the field names, bit numbers, or bit range, which is followed by an equal sign (=) and a value. | Register_Short_Name[Field_Name] = '*n*'<br>(where *n* is a binary value for the bit or bit range) | MSR[FE0] = '1' |
| | Register_Short_Name[Field_Name] = x'*n*'<br>(where *n* is a hexadecimal value for the bit or bit range) | MSR[FE] = x'F' |
| | Register_Short_Name[Bit_Number] = '*n*'<br>(where *n* is a binary value for the bit or bit range) | MSR[52] = '0' |
| | Register_Short_Name[Bit_Number] = x'*n*'<br>(where *n* is a hexadecimal value for the bit or bit range) | MSR[52] = x'F' |
| | Register_Short_Name[Starting_Bit_Number:Ending_Bit_Number] = '*n*'<br>(where *n* is the binary value for the bit or bit range) | MSR[39:43] = '10010' |
| | Register_Short_Name[Starting_Bit_Number:Ending_Bit_Number] = '*n*'<br>(where *n* is the hexadecimal value for the field or bits) | MSR[39:43] = x'11' |
| **Note:** The register short name is also called the register mnemonic. | | |

## Terminology

In this handbook, a "memory-mapped I/O (MMIO)" register is any internal or external register that is accessed through the main-storage space with load and store instructions, whether or not the register is associated with an I/O controller or device.

## Reserved Regions of Memory and Registers

Reserved areas of the MMIO-register memory map that are not assigned to any functional unit should not be read from or written to. Doing so will cause serious errors in software as follows. For reads or writes generated from the Synergistic Processor Element (SPE) to unassigned reserved spaces, at least one of the following MFC_FIR[46, 53, 56, 58, 61] bits will be set and in most cases will cause a checkstop. For reads or writes generated from the PowerPC Processor Element (PPE) to unassigned reserved spaces, at least one of the CIU_FIR[7,8] will be set, and a checkstop will occur. For reads or writes generated from the I/O interface controller (IOC) to unassigned reserved spaces, the IOC will respond back to the I/O Interface (IOIF) device that sourced the address request with an error (ERR) response. No IOC Fault Isolation Register (FIR) bits are set.

# Revision Log

Each release of this document supersedes all previously released versions. The revision log lists all significant changes made to the document since its initial release. In the rest of the document, change bars in the margin indicate that the adjacent text was significantly modified from the previous release of this document.

| Revision Date | Version | Contents of Modification |
|---|---|---|
| May 12, 2008 | 1.11 | • Revised the *Preface* on page 29 to indicate that the handbook includes information about the PowerXCell 8i processor.<br>• Updated *Related Publications* on page 29.<br>• Revised *Section 1 Overview of CBEA Processors* on page 39 to indicate that the handbook includes information about the PowerXCell 8i processor.<br>• Revised *Figure 1-1 Overview of CBEA Processors* on page 40 to illustrate the double data rate 2 (DDR2) synchronous dynamic random access memory (SDRAM) interface provided by the PowerXCell 8i processor.<br>• Revised *Section 1.2.3 Memory Interface Controller* on page 45 to include information about the PowerXCell 8i processor.<br>• Revised *Section 3.1.4 Floating-Point Support* on page 70 to include information about the PowerXCell 8i processor.<br>• Expanded the description of the IOC_IOST_Origin[IOST Size] field (see *Section 7.3.2.1 Register Layout and Fields* on page 166).<br>• Removed extreme data rate (XDR) from *Figure 8-2 Managed Resources and Requesters* on page 205.<br>• Revised *Section 8.5.5.1 Memory Banks* on page 216 to include PowerXCell 8i processors as well as Cell/B.E. processors.<br>• Changed `IOC_BaseAddr1` to `IOC_BaseAddrMask1` (see *Section 8.7.3 Changing a Requester's Address Map* on page 235).<br>• Removed XDR from *Figure 9-1 Organization of Interrupt Handling* on page 243.<br>• Revised *Section 9.7.1 Mediated External Interrupt Architecture* on page 276 and *Section 9.7.2 Mediated External Interrupt Implementation* on page 279.<br>• Changed "LPCR[RMI]" to "LPCR[LPES] bit 0" in the SRR1 and HSRR1 rows of *Table 9-27 Registers Altered by an External Interrupt* on page 279.<br>• Revised some descriptions to make it clear that a PowerPC Processor Element (PPE) is a dual-instruction dispatch machine (see *Section 10.5.3.1 Low Priority Combined with Low Priority* on page 316, *Section 10.5.3.2 Low Priority Combined with Medium Priority* on page 316, *Section 10.5.3.3 Low Priority Combined with High Priority* on page 317, *Section 10.5.3.4 Medium Priority Combined with Medium Priority* on page 318, and *Section 10.7.2.4 Microcoded Instruction* on page 323).<br>• Revised *Section 20.1.6.2 SPE-to-PPE Communications* on page 585 to correctly describe how to ensure ordering, and deleted *Section 20.1.6.3 Implementation-Specific Ordering*.<br>• Revised *Section 21.4.4 Multiple Execution Units* on page 618 to include PowerXCell 8i processors.<br>• Deleted a sentence fragment (see *Table 23-6 Unique SPU-to-Vector/SIMD Multimedia Extension Data-Type Mappings* on page 689).<br>• Revised *Section 24.2 SPU Pipelines and Dual-Issue Rules* on page 698 and *Table 24-1 Pipeline 0 Instructions and Latencies* on page 698 to include information about the PowerXCell 8i processor.<br>• Corrected *Figure A-1 Dual-Issue Combinations* on page 761.<br>• Revised *Appendix B.1.2 Instructions* on page 771 and *Table B-1 SPU Instructions* on page 772 to include information about the PowerXCell 8i processor and double-precision floating-point instructions.<br>• Revised *Section B.1.3.2 Issue* on page 780. Renamed *Table B-5 Cell/B.E. SPU Execution Pipelines and Result Latency* on page 782 and added the double-precision (DP) instruction to stage S/7. Added *Table B-6 PowerXCell 8i SPU Execution Pipelines and Result Latency* on page 783 and an explanation of latency on the PowerXCell 8i processor.<br>• Added the vector test special value intrinsic to *Table B-13 SPU Intrinsics* on page 788.<br>• Added or revised the following terms: CBEA, CBEA processor, Cell/B.E., Cell/B.E. processor, Cell Broadband Engine Architecture, DDR2, PowerXCell 8i processor, X2D (see the *Glossary* on page 835). |

| Revision Date | Version | Contents of Modification |
|---|---|---|
| May 12, 2008 (continued) | 1.11 | • Changed “local store” to “local storage” throughout. Therefore, changed MFC Local Store Address Compare Register to MFC Local Storage Address Compare Register and MFC Local Store Compare Results Register to MFC Local Storage Compare Results Register (see *Section 4.3.7.2 Address-Compare Exceptions* on page 120, *Section 9.8.3.1 Storage-Protection Errors* on page 286, and *Section 11.3.4.5 Debug and Performance Monitoring* on page 350).<br>• Changed “CBE processor” to “CBEA processor” or “Cell/B.E. processor”, as appropriate, throughout. Changed other instances of “CBE” to “Cell/B.E.” |
| April 24, 2007 | 1.1 | • Updated the list of related publications (see *Related Publications* on page 29).<br>• Corrected the example of the shift to the left symbol (see *Conventions and Notation* on page 30).<br>• Moved the revision log to the front of the book and added text introducing the table. Updated a section title to reflect the change (see *Revision Log* on page 33 and *Appendixes, Glossary, and Index* on page 721).<br>• Updated *Figure 2-2 PPE Functional Units* on page 52.<br>• Corrected the description of address translation (see *Section 3.1.1.1 Addressing and Address Aliasing* on page 66).<br>• Removed *Table 2-2 PowerPC Instruction Summary* (see *Section 2.4.3 Instructions* on page 58) and referred the reader to *Table A-1 PowerPC Instructions by Execution Unit* on page 723, which is more complete.<br>• Changed byte zero of VT from “A.1” to “A.0” (see *Figure 2-6 Byte-Shuffle (Permute) Operation* on page 60).<br>• Removed *Table 2-3 Vector/SIMD Multimedia Extension Instruction Summary* (see *Section 2.5.5 Instructions* on page 62) and referred the reader to *Table A-8 Vector/SIMD Multimedia Extension Instructions* on page 748, which is more complete.<br>• Deleted a redundant list of vector register values (see *Section 2.6.1 Vector Data Types* on page 62).<br>• Replaced *Table 3-2 SPE Floating-Point Support* with *Table 3-2 Single-Precision (Extended-Range Mode) Minimum and Maximum Values* on page 71, *Table 3-3 Double-Precision (IEEE Mode) Minimum and Maximum Values* on page 72, and *Table 3-4 Single-Precision (IEEE Mode) Minimum and Maximum Values* on page 72.<br>• Removed *Table 3-3 SPU Instruction Summary* (see *Section 3.3.2 Instructions* on page 77) and referred the reader to *Table B-1 SPU Instructions* on page 772, which is more complete.<br>• Corrected the granularity of the real-mode address boundary (see *Section 4.3.6 Real Addressing Mode* on page 117).<br>• Deleted an incorrect sentence (see *Section 5.1.3 Reserved Regions of Memory* on page 126).<br>• Corrected the offsets in *Table 5-3 PPE Privileged Memory-Mapped Register Groups* on page 126.<br>• Changed the number of branches needed to clear the ICache from 512 to 256 (see *Section 6.1.3.6 ICache Invalidation* on page 138).<br>• Indicated that the PPE can send multiple overlapping independent (rather than dependent) loads to the L2 (see *Section 6.1.3.7 DCache Load Misses* on page 139).<br>• Indicated that HID4[enb_force_ci] and HID4[dis_force_ci] work only for load/store instructions and affect only the L1 and L2 data caches (see *Section 6.1.3.13 L1 Bypass (Inhibit) Configuration* on page 140). |

| Revision Date | Version | Contents of Modification |
|---|---|---|
| April 24, 2007 (continued) | 1.1 | • Revised the description of the **dcbt** and **dcbtst** instructions (see *Section 6.1.6.2 Data Cache Block Touch* on page 147).<br>• Updated the memory-allocation and IOIF-allocation percentage equations (see *Section 8.5.10.1 Memory Allocation* on page 225 and *Section 8.5.10.2 IOIF Allocation* on page 225).<br>• Reworded the description of the MC_COMP_EN configuration-ring register. Added a cross reference to *Table 8-8 IOIF1_COMP_EN and BE_MMIO_Base Settings* on page 233 to the description of the IOIF1_COMP_EN configuration-ring register (see *Table 8-6 SPE Address-Range Registers* on page 232).<br>• Indicated that the TKM registers can only be changed when there are no outstanding requests in the token manager and that the Allocation registers must be written to zero before the allocation rate is changed (see *Section 8.7.6 Changing TKM Registers* on page 236).<br>• Corrected a footnote describing PPE floating-point exceptions (see *Section 9.2 Summary of Interrupt Architecture* on page 240).<br>• Changed an instance of HID4[en_dcset] to HID4[en_dcway]. Removed a statement indicating that, if the DABR match feature is enabled, any non-VMU cacheable (I = ‘0’) load or store instruction crossing an 8-byte boundary causes an alignment interrupt. Indicated that the PPE does not modify the DSISR when an alignment interrupt is taken (see *Section 9.5.8 Alignment Interrupt (x‘00..00000600’)* on page 255).<br>• Described software requirements for using the interrupt status and masking registers (see *Section 9.6.3.1 Status and Masking* on page 271).<br>• Changed “Vector Multimedia Registers” to “Vector Registers” and “VMR” to “VR” (see *Section 10.2.1 Registers* on page 301).<br>• Explained how stall conditions are handled when the instruction that caused the stall is dependent on a caching-inhibited load instruction (see *Section 10.2.3.1 Pipeline Stall Points* on page 304).<br>• Added a note indicating that MSR[ME] can only be modified by the **rfid** and **hrfid** instructions while in hypervisor mode (see *Section 10.4.1 Machine State Register (MSR)* on page 307).<br>• Added information about the maximum time-base frequency limit (see *Section 13.2.3 Time-Base Frequency* on page 383).<br>• Indicated that the lock-line reservation lost event will not be raised when the reservation is reset by a matching **putllc**, **putlluc**, or **putqlluc** operation, as opposed to a **put** operation, which will cause the lock-line reservation event. Added a note on how to avoid the starvation of writes and low-priority reads (see *Section 18.6.4 Procedure for Handling the Lock-Line Reservation Lost Event* on page 485).<br>• Corrected certain first level interrupt handler (FLIH) code samples to enable proper handling of the Link Register (see *Section 18.7.2 FLIH Design* on page 496, *Section 18.8.2 FLIH Design for Nested Interrupts* on page 502, and *Section 18.9 Using a Dedicated Interrupt Stack* on page 504).<br>• Indicated that, for naturally aligned 1, 2, 4, and 8-byte transfers, the source and destination addresses must have the same 4 least significant bits (see *Section 19.2.1 DMA Commands* on page 516).<br>• Corrected the **putlb** row in *Table 19-4 MFC DMA Put Commands* on page 516. The PPE cannot initiate the **putlb** command.<br>• Added a footnote indicating that 64-bit access to an address range that includes a 32-bit MMIO register is not allowed unless otherwise explicitly specified (see *Section 19.3.1 MFC Command Issue* on page 523).<br>• Corrected the description of how the Cell Broadband Engine handles DMA list element transfers that cross the 4 GB area defined by the EAH (see *Section 19.4.4.2 Initiating the Transfers Specified in the List* on page 537).<br>• Removed extraneous **eieio** instructions from three code examples (see *Section 19.6.5.4 PPE Side* on page 546, *Section 19.6.6.2 PPE Side* on page 547, and *Section 19.7.6.1 From the PPE* on page 553).<br>• Corrected the effect of **sync**, **dsync**, and **syncc** instructions issued by an SPU on the LS domain (see *Table 20-1 Effects of Synchronization on Address and Communication Domains1* on page 563).<br>• Deleted an obsolete paragraph indicating that data transfer only occurs on the first **getllar** instruction (see *Section 20.3.1.1 The Get Lock-Line and Reserve Command—getllar* on page 598). |

| Revision Date | Version | Contents of Modification |
|---|---|---|
| April 24, 2007 (continued) | 1.1 | • Removed a statement indicating that an SPE external event can be used to determine the status of an atomic command (see *Section 20.3.2 The MFC Read Atomic Command Status Channel* on page 599).<br>• Provided a code example that illustrates the use of sleeping locks (see *Section 20.3.3.1 Lock-Line Reservation Lost Event* on page 600).<br>• Clarified the third method used for triangle subdivision (see *Method 3: Evaluate Four Vertices at a Time Using SOA* on page 633).<br>• Corrected a code example (see *Section 22.2.1.7 SIMDization Interaction with Other Optimizations* on page 650).<br>• Removed stores and branch hints from *Table 24-2 Pipeline 1 Instructions and Latencies* on page 699.<br>• Corrected the statement, “The HBR instruction must be placed within 64 instructions of the branch instruction.” The branch instruction is encoded using a split 9-bit field, which is interpreted with 2 implied zero bits. Therefore, the instruction must be within -256 to +255 instructions (see *Section 24.3.3.5 Rules for Using Branch Hints* on page 704).<br>• Removed an incorrect bullet describing pipelined instructions (see *Appendix A.1.2 PPE Instructions* on page 723).<br>• Updated *Table A-1 PowerPC Instructions by Execution Unit* on page 723.<br>• Updated *Table A-2 Storage Alignment for PowerPC Instructions* on page 733.<br>• Changed the code sequence for **mtmsr** and **mtmsrd** (see *Table A-3 Unconditionally Microcoded Instructions (Except Loads and Stores)* on page 735).<br>• Added **tlbsync** to the lists of optional PowerPC instructions implemented in the PPE (see *Appendix A.2.3.2 Book III Optional Instructions Implemented* on page 746).<br>• Changed the title of a section from “Book II Optional Instructions Implemented” to “Book III Optional Instructions Implemented.” In that section, removed the **slbie** and **slbia** instructions because they are not optional (see *Section A.2.3.2 Book III Optional Instructions Implemented* on page 746).<br>• Added the **fre(.), frsqrtes(.)**, and **popcntb** instructions to the list of instructions not implemented in the PPE (see *Appendix A.2.4.1 Book I Unimplemented Instructions* on page 747).<br>• Removed **bccbr** from the list of obsolete user-mode instructions that are not implemented (this instruction is not documented in the latest version of *PowerPC User Instruction Set Architecture, Book I*) (see *Appendix A.2.4.1 Book I Unimplemented Instructions* on page 747).<br>• For graphics rounding mode, corrected the value of the logarithm base-2 estimate output when the input is +0 (see *Section A.3.3.5 Logarithm Base-2 Estimate* on page 753).<br>• Corrected *Figure A-1 Dual-Issue Combinations* on page 761.<br>• Added several recommendations that can be used to help improve performance (see *Section A.7.4 General Recommendations* on page 770).<br>• Updated *Table B-1 SPU Instructions* on page 772.<br>• Changed “JSRE-compliant compiler” to “compiler compliant with the *C/C++ Language Extensions for Cell Broadband Engine Architecture* specification” (see *Section B.2.1 Vector Data Types* on page 784).<br>• Updated *Appendix C-2* begining on page 794:<br>– Added a table note that addresses the functionality restrictions of signal bits 6 and 25 in *Appendix C.2.1 PPU Instruction Unit* on page 797.<br>– Added a comment regarding the correct use of the shared L2 trace bus in *Appendix C.3.2 PPSS L2 Cache Controller - Group 1 (NClk/2)* on page 800, *Appendix C.3.3 PPSS L2 Cache Controller - Group 2 (NClk/2)* on page 801, and *Appendix C.3.4 PPSS L2 Cache Controller - Group 3 (NClk/2)* on page 802.<br>– Replaced obsolete terms in *Appendix C.4.1* begining on page 804.<br>– Added bit 63 in *Appendix C.4.3 SPU Event (NClk)* on page 807.<br>– Added bit 22 in *Appendix C.6.2 EIB Address Concentrator 1 (NClk/2)* on page 813.<br>– Provided clearer meanings to several bit descriptions in *Appendix C.6.3 EIB Data Ring Arbiter - Group 1 (NClk/2)* on page 814 and *Appendix C.6.4 EIB Data Ring Arbiter - Group 2 (NClk/2)* on page 815. |

| Revision Date | Version | Contents of Modification |
|---|---|---|
| April 24, 2007 (continued) | 1.1 | • Added “architecture,” “Cell Broadband Engine Architecture,” “JSRE”, and “island” to the glossary (see *Glossary* on page 835).<br>• Changed “*SPU C/C++ Language Extensions*” to “*C/C++ Language Extensions for Cell Broadband Engine Architecture*” throughout.<br>• Changed “L2_RMT_Setup” to “L2_RMT_Data” throughout.<br>• Changed “SPU_RdEventStatMask” to “SPU_RdEventMask” throughout.<br>• Made other changes to improve the clarity and consistency of the document. |
| April 19, 2006 | 1.0 | Initial release |

# 1. Overview of CBEA Processors

This handbook describes the extensive programming facilities of the Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors, which are collectively called Cell Broadband Engine Architecture processors (CBEA processors). The Cell Broadband Engine Architecture (CBEA) was developed jointly by Sony, Toshiba, and IBM; it extends the 64-bit PowerPC Architecture™. The Cell/B.E. processor is a result of that collaboration between Sony, Toshiba, and IBM. It is the first implementation of a multiprocessor family conforming to the CBEA. The IBM PowerXCell 8i processor also conforms to the CBEA. It provides a double data rate 2 (DDR2) memory interface and improved double-precision, floating-point performance and additional double-precision instructions. Both CBEA processors comprise eight Synergistic Processor Elements (SPEs) and one PowerPC Processor Element (PPE).

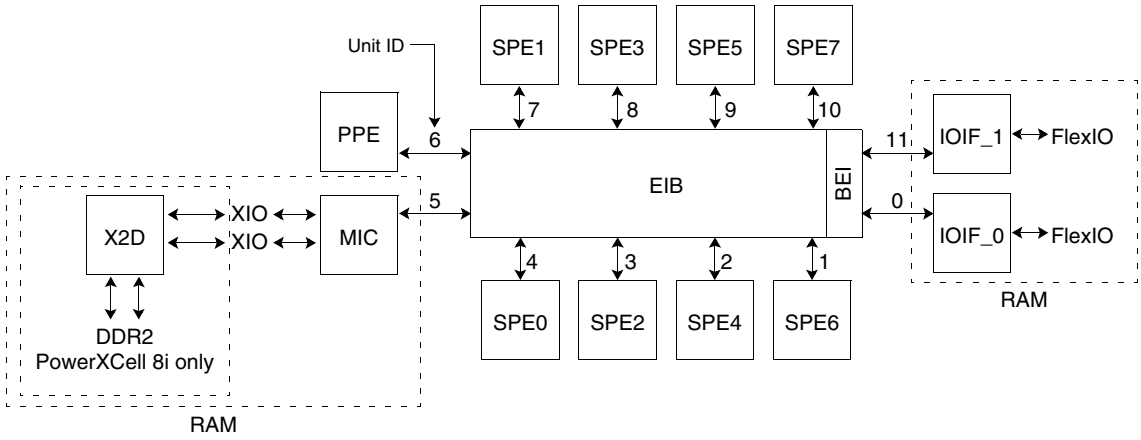
Although the Cell/B.E. processor is initially intended for applications in media-rich consumer-electronics devices such as game consoles and high-definition televisions, the architecture has been designed to enable fundamental advances in processor performance. These advances are expected to support a broad range of applications in both commercial and scientific fields.

This handbook is written for the complete range of programmers, including those developing applications (user programs), libraries, device drivers, middleware, compilers, and operating systems. It assumes that the reader is an experienced C/C++ programmer. It describes and presents examples of both basic and advanced programming concepts for single-instruction, multiple-data (SIMD) vector applications and the system software that supports such applications.

The handbook is system-independent, making no assumptions about development-tool or operating-system environments, other than the C/C++ language environment. The examples are chosen to highlight the general principals required to program the CBEA processors, such that an experienced programmer can apply this knowledge to their particular system environment.

*Figure 1-1* on page 40 shows a block diagram of the CBEA-processor hardware. This figure is referred to later in this section and in subsequent sections.

*Figure 1-1. Overview of CBEA Processors*




| | |
|---|---|
| BEI | Cell Broadband Engine Interface |
| DDR2 | Double-Data-Rate 2 Synchronous Dynamic Random Access Memory (SDRAM) Interface |
| EIB | Element Interconnect Bus |
| FlexIO | Rambus FlexIO Bus |
| IOIF | I/O Interface |
| MIC | Memory Interface Controller |
| PPE | PowerPC Processor Element |
| RAM | Resource Allocation Management |
| SPE | Synergistic Processor Element |
| X2D | XIO to DDR2 logic |
| XIO | Rambus XDR I/O (XIO) Cell |

## 1.1 Background

### 1.1.1 Motivation

The Cell Broadband Engine Architecture has been designed to support a very broad range of applications. The CBEA processors described in this handbook are single-chip multiprocessors with nine processor elements operating on a shared, coherent memory, as shown in *Figure 1-1*. In this respect, CBEA processors extend current trends in PC and server processors. The most distinguishing feature of the CBEA processors is that, although all processor elements share memory, their function is specialized into two types: the PowerPC Processor Element (PPE) and the Synergistic Processor Element (SPE). The CBEA processors have one PPE and eight SPEs.

The first type of processor element, the PPE, contains a 64-bit PowerPC Architecture core. It complies with the 64-bit PowerPC Architecture and can run 32-bit and 64-bit operating systems and applications. The second type of processor element, the SPE, is optimized for running compute-intensive SIMD applications; it is not optimized for running an operating system. The SPEs are independent processor elements, each running their own individual application programs or threads. Each SPE has full access to coherent shared memory, including the memory-mapped I/O space. There is a mutual dependence between the PPE and the SPEs. The

SPEs depend on the PPE to run the operating system, and, in many cases, the top-level thread control for an application. The PPE depends on the SPEs to provide the bulk of the application performance.

The SPEs are designed to be programmed in high-level languages, such as (but certainly not limited to) C/C++. They support a rich instruction set that includes extensive SIMD functionality. However, like conventional processors with SIMD extensions, use of SIMD data types is preferred, not mandatory. For programming convenience, the PPE also supports the standard PowerPC Architecture instructions and the vector/SIMD multimedia extensions.

To an application programmer, the CBEA processors look like 9-way coherent multiprocessors. The PPE is more adept than the SPEs at control-intensive tasks and quicker at task switching. The SPEs are more adept at compute-intensive tasks and slower than the PPE at task switching. However, either processor element is capable of both types of functions. This specialization is a significant factor accounting for the order-of-magnitude improvement in peak computational performance and chip-area-and-power efficiency that the CBEA processors achieve over conventional PC processors.

The more significant difference between the SPE and PPE lies in how they access memory. The PPE accesses main storage (the effective-address space) with load and store instructions that move data between main storage and a private register file, the contents of which may be cached. The SPEs, in contrast, access main storage with direct memory access (DMA) commands that move data and instructions between main storage and a private local memory, called *local storage* (LS). An SPE's instruction-fetches and load and store instructions access its private LS rather than shared main storage, and the LS has no associated cache. This 3-level organization of storage (register file, LS, main storage), with asynchronous DMA transfers between LS and main storage, is a radical break from conventional architecture and programming models, because it explicitly parallelizes computation with the transfers of data and instructions that feed computation and store the results of computation in main storage.

One of the enabling accomplishments for this radical change is that memory latency, measured in processor cycles, has gone up several hundredfold from about the years 1980 to 2000. The result is that application performance is, in most cases, limited by memory latency rather than peak compute capability or peak bandwidth. When a sequential program on a conventional architecture performs a load instruction that misses in the caches, program execution comes to a halt for several hundred cycles. Compared to this penalty, the few cycles it takes to set up a DMA transfer for an SPE are a much better trade-off, especially considering the fact that each of the eight SPE's DMA controller can have up to 16 DMA transfer in flight simultaneously. Conventional processors, even with deep and costly speculation, manage to get, at best, a handful of independent memory accesses in flight.

One of the SPE's DMA transfer methods supports a list (such as a scatter-gather list) of DMA transfers that is constructed in an SPE's local storage, so that the SPE's DMA controller can process the list asynchronously while the SPE operates on previously transferred data. In several cases, this approach to accessing memory has led to application performance exceeding that of conventional processors by almost two orders of magnitude—significantly more than one would expect from the peak performance ratio (approximately 10x) between the CBEA processors and conventional PC processors. The DMA transfers can be set up and controlled by the SPE that is sourcing or receiving the data, or by the PPE or another SPE.

### 1.1.2 Power, Memory, and Frequency

The CBEA processors overcome three important limitations of contemporary microprocessor performance—power use, memory use, and clock frequency.

Microprocessor performance is approaching limits of power dissipation rather than integrated-circuit resources (transistors and wires). The only way to significantly increase the performance of microprocessors in this environment is to improve power efficiency at approximately the same rate as the performance increase. The CBEA processors do this by differentiating between the PPE, optimized to run an operating system and control-intensive code, and the eight SPEs, optimized to run compute-intensive applications. The control-plane PPE leaves the eight data-plane SPEs free to compute data-rich applications.

On today's symmetric multiprocessors—even those with integrated memory controllers—latency to DRAM memory is approaching 1000 cycles. As a result, program performance is dominated by moving data between main storage and the processor. Compilers and application writers must manage this data movement explicitly, even though the hardware cache mechanisms are supposed to relieve them of this task. In contrast, the CBEA processor mechanisms for dealing with memory latencies—the 3-level SPE memory structure (main storage, local storages, and large register files), and asynchronous DMA transfers—enable programmers to schedule simultaneous data and code transfers to cover long memory latencies. At 16 simultaneous transfers per SPE, the CBEA processors can support up to 128 simultaneous transfers between the SPE local storages and main storage. This surpasses the bandwidth of conventional processors by a factor of almost twenty.

Conventional processors require increasingly deeper instruction pipelines to achieve higher operating frequencies. This technique has reached a point of diminishing returns—and even negative returns if power is taken into account. By specializing the PPE for control-intensive tasks and the SPEs for compute-intensive tasks, these processing elements run at high frequencies without excessive overhead. The PPE achieves efficiency by executing two threads simultaneously rather than by optimizing single-thread performance. Each SPE achieves efficiency by using a large register file that supports many simultaneous in-flight instructions without the overhead of register-renaming or out-of-order processing, and asynchronous DMA transfers, that support many concurrent memory operations without the overhead of speculation.

By distinguishing and separately optimizing control-plane and data-plane processor elements, the CBEA processors mitigate the problems posed by power, memory, and frequency limitations. The net result is a multiprocessor that, at the power budget of a conventional PC processor, can provide approximately ten-fold the peak performance of a conventional processor. Of course, actual application performance varies. Some applications may benefit little from the SPEs, whereas others show a performance increase well in excess of ten-fold. In general, compute-intensive applications that use 32-bit or smaller data formats (such as single-precision floating-point and integer) are excellent candidates for the CBEA processors.

### 1.1.3 Scope of this Handbook

As mentioned at the beginning, this handbook is written for the complete range of programmers, including those developing applications, libraries, device drivers, middleware, compilers, and operating systems.

The remainder of this section summarizes highlights of the hardware and programming environments. The two sections that follow expand on these topics for the CBEA processor programming targets, the PPE and the SPEs.

This is followed in Sections 4 through 16 with additional hardware overviews and details primarily relating to privilege-state (supervisor) programming. A few user programming topics are also presented in these sections. The topics covered include virtual-memory management, cache management, I/O, resource-allocation management (the management of memory and I/O resources), interrupts, multithreading, hypervisors (the highest privilege state), SPE context-switching, the CBEA processor clocks (time base) and decrementers, object-file formats, and power and performance monitoring.

Problem-state (user) programming topics are covered primarily in Sections 17 through 24, and the appendixes. The topics include the SPE channel interface, SPE events, DMA transfers and interprocessor communications, shared-storage synchronization, parallel programming techniques, SIMD programming techniques, PPE versus SPE support for vector operations, general SPE programming tips, and summaries of the PPE and SPE instruction sets and C/C++ intrinsics. A few supervisor-programming and compiler-writer topics are also covered in these sections.

There are many programming examples throughout the handbook, and active hyperlinks provide cross-references to relevant topics in various sections. The *Glossary* on page 835 defines terms that have special meanings in the context of the CBEA processors.

## 1.2 Hardware Environment

### 1.2.1 The Processor Elements

*Figure 1-1* on page 40 shows a high-level block diagram of the CBEA processor hardware. There is one PPE, and there are eight identical SPEs. All processor elements are connected to each other and to the on-chip memory and I/O controllers by the memory-coherent element interconnect bus (EIB).

The PPE contains a 64-bit, dual-thread PowerPC Architecture RISC core and supports a PowerPC virtual-memory subsystem. It has 32 KB level-1 (L1) instruction and data caches and a 512 KB level-2 (L2) unified (instruction and data) cache. It is intended primarily for control processing, running operating systems, managing system resources, and managing SPE threads. It can run existing PowerPC Architecture software and is well-suited to executing system-control code. The instruction set for the PPE is an extended version of the PowerPC instruction set. It includes the vector/SIMD multimedia extensions and associated C/C++ intrinsic extensions.

The eight identical SPEs are single-instruction, multiple-data (SIMD) processor elements that are optimized for data-rich operations allocated to them by the PPE. Each SPE contains a RISC core, 256 KB software-controlled LS for instructions and data, and a 128-bit, 128-entry unified register file. The SPEs support a special SIMD instruction set—the *Synergistic Processor Unit Instruction Set Architecture*—and a unique set of commands for managing DMA transfers and interprocessor messaging and control. SPE DMA transfers access main storage using PowerPC effective addresses. As in the PPE, SPE address translation is governed by PowerPC Architecture segment and page tables, which are loaded into the SPEs by privileged software running on the PPE. The SPEs are not intended to run an operating system.

An SPE controls DMA transfers and communicates with the system by means of *channels* that are implemented in and managed by the SPE's memory flow controller (MFC). The channels are unidirectional message-passing interfaces. The PPE and other devices in the system, including other SPEs, can also access this MFC state through the MFC's memory-mapped I/O (MMIO) registers and queues, which are visible to software in the main-storage address space.

For more information about the PPE, see *Section 2* on page 51. For more information about SPEs, see *Section 3* on page 65.

### 1.2.2 Element Interconnect Bus

The element interconnect bus (EIB) is the communication path for commands and data between all processor elements on the CBEA processors and the on-chip controllers for memory and I/O. The EIB supports full memory-coherent and symmetric multiprocessor (SMP) operations. Thus, the CBEA processors are designed to be ganged coherently with other CBEA processors to produce a cluster.

The EIB consists of four 16-byte-wide data rings. Each ring transfers 128 bytes (one PPE cache line) at a time. Each processor element has one on-ramp and one off-ramp. Processor elements can drive and receive data simultaneously. *Figure 1-1* on page 40 shows the unit ID numbers of each element and the order in which the elements are connected to the EIB. The connection order is important to programmers seeking to minimize the latency of transfers on the EIB:

latency is a function of the number of connection hops, such that transfers between adjacent elements have the shortest latencies and transfers between elements separated by six hops have the longest latencies.

The EIB's internal maximum bandwidth is 96 bytes per processor-clock cycle. Multiple transfers can be in-process concurrently on each ring, including more than 100 outstanding DMA memory requests between main storage and the SPEs. The EIB does not support any particular quality-of-service (QoS) behavior other than to guarantee forward progress. However, a resource allocation management (RAM) facility, shown in *Figure 1-1* on page 40, resides in the EIB. Privileged software can use it to regulate the rate at which resource requesters (the PPE, SPEs, and I/O devices) can use memory and I/O resources.

### 1.2.3 Memory Interface Controller

The on-chip memory interface controller (MIC) provides the interface between the EIB and physical memory. It supports one or two Rambus extreme data rate (XDR) memory interfaces, which together support between 64 MB and 64 GB of XDR DRAM memory. The PowerXCell 8i processor supports one or two 128-bit DDR2 memory channels.

Memory accesses on each interface are 1 to 8, 16, 32, 64, or 128 bytes, with coherent memory-ordering. Up to 64 reads and 64 writes can be queued. The resource-allocation token manager provides feedback about queue levels.

The MIC has multiple software-controlled modes, including fast-path mode (for improved latency when command queues are empty), high-priority read (for prioritizing SPE reads in front of all other reads), early read (for starting a read before a previous write completes), speculative read, and slow mode (for power management). The MIC implements a closed-page controller (bank rows are closed after being read, written, or refreshed), memory initialization, and memory scrubbing.

The XDR DRAM memory is ECC-protected, with multi-bit error detection and optional single-bit error correction. It also supports write-masking, initial and periodic timing calibration, dynamic width control, sub-page activation, dynamic clock gating, and 4, 8, or 16 banks.

The PowerXCell 8i processor incorporates an X2D macro that converts XDR packets to DDR2 commands. The processor can be configured with one or two DDR2 memory channels, providing 1 GB to 32 GB of error correction code (ECC) corrected memory.

### 1.2.4 Cell Broadband Engine Interface Unit

The on-chip Cell Broadband Engine interface (BEI) unit supports I/O interfacing. It includes a bus interface controller (BIC), I/O controller (IOC), and internal interrupt controller (IIC), as defined in the *Cell Broadband Engine Architecture* document. It manages data transfers between the EIB and I/O devices and provides I/O address translation and command processing.

The BEI supports two Rambus FlexIO interfaces. One of the two interfaces (IOIF1) supports only a noncoherent I/O interface (IOIF) protocol, which is suitable for I/O devices. The other interface (IOIF0, also called BIF/IOIF0) is software-selectable between the noncoherent IOIF protocol and the memory-coherent Cell Broadband Engine interface (BIF) protocol. The BIF protocol is the EIB's internal protocol. It can be used to coherently extend the EIB, through IOIF0, to another memory-coherent device, that can be another CBEA processor.

## 1.3 Programming Environment

CBEA-processor software development in the C/C++ language is supported by a rich set of language extensions that define C/C++ data types for SIMD operations and map C/C++ intrinsics (which are commands, in the form of function calls) to one or more assembly instructions. These language extensions, summarized in *Appendix A* on page 723 and *Appendix B* on page 771, give C/C++ programmers great control over code performance, without the need for assembly-language programming.

### 1.3.1 Instruction Sets

The instruction set for the PPE is an extended version of the PowerPC Architecture instruction set. The extensions consist of the vector/SIMD multimedia extensions, a few additions and changes to PowerPC Architecture instructions, and C/C++ intrinsics for the vector/SIMD multimedia extensions.

The instruction set for the SPEs is a new SIMD instruction set, the *Synergistic Processor Unit Instruction Set Architecture*, with accompanying C/C++ intrinsics, and a unique set of commands for managing DMA transfer, external events, interprocessor messaging, and other functions. The instruction set for the SPEs is similar to that of the PPE's vector/SIMD multimedia extensions, in the sense that they operate on SIMD vectors. However, the two vector instruction sets are distinct, and programs for the PPE and SPEs are often compiled by different compilers.

Although most coding for the CBEA processors will probably be done in a high-level language like C or C++, an understanding of the PPE and SPE machine instructions adds considerably to a software developer's ability to produce efficient, optimized code. This is particularly true because most of the C/C++ intrinsics have a mnemonic that relates directly to the underlying assembly-language mnemonic.

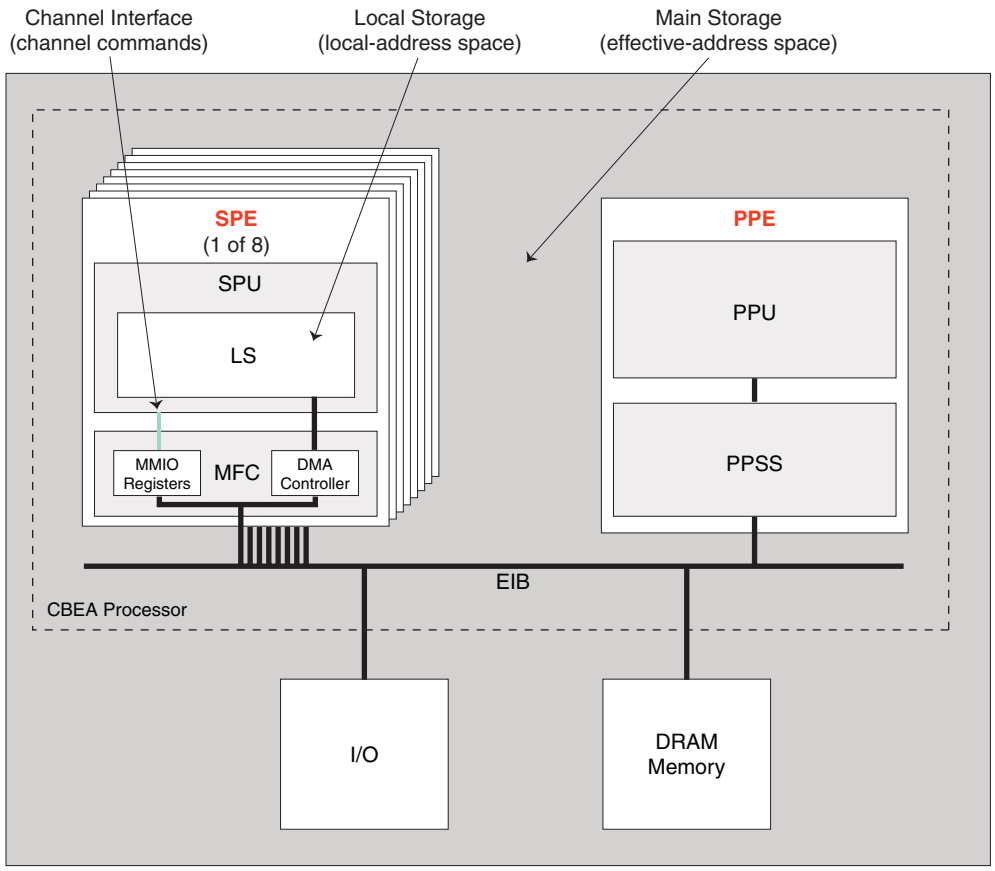
### 1.3.2 Storage Domains and Interfaces

The CBEA processors have two types of storage domains—*one main-storage domain* and *eight SPE local-storage (LS) domains*, as shown in *Figure 1-2* on page 47. In addition, each SPE has a *channel interface* for communication between its synergistic processor unit (SPU) and its MFC. The main-storage domain, which is the entire effective-address space, can be configured by PPE privileged software to be shared by all processor elements and memory-mapped devices in the system[1]. An MFC's state is accessed by its associated SPU through the channel interface, and this state can also be accessed by the PPE and other devices (including other SPEs) by means of the MFC's MMIO registers in the main-storage space. An SPU's LS can also be accessed by the PPE and other devices through the main-storage space. The PPE accesses the main-storage space through its PowerPC processor storage subsystem (PPSS).

The address-translation mechanisms used in the main-storage domain are described in *Section 4 Virtual Storage Environment* on page 79. The channel domain is described in *Section 19 DMA Transfers and Interprocessor Communication* on page 513.

1. In the PowerPC Architecture, all I/O is memory-mapped.

*Figure 1-2. Storage and Domains and Interfaces*




| | | | |
|---|---|---|---|
| DMA | Direct Memory Access | PPE | PowerPC Processor Element |
| EIB | Element Interconnect Bus | PPSS | PowerPC Processor Storage Subsystem |
| LS | Local Storage | PPU | PowerPC Processor Unit |
| MFC | Memory Flow Controller | SPE | Synergistic Processor Element |
| MMIO | Memory-Mapped I/O | SPU | Synergistic Processor Unit |

An SPE's SPU can fetch instructions only from its own LS, and load and store instructions executed by the SPU can only access the LS. SPU software uses LS addresses (not main-storage effective addresses) to do this. Each SPE's MFC contains a DMA controller. DMA-transfer requests contain both an LS address and an effective address, thereby facilitating transfers between the domains.

Data transfers between an SPE's LS and main storage are performed by the associated SPE, or by the PPE or another SPE, using the DMA controller in the MFC associated with the LS. Software running on the associated SPE interacts with its own MFC through its channel interface. The channels support enqueueing of DMA commands and other facilities, such as mailbox and signal-notification messages. Software running on the PPE or another SPE can interact with an MFC through MMIO registers, which are associated with the channels and visible in the main-storage space.

Each MFC maintains and processes two independent command queues for DMA and other commands—one queue for its associated SPU, and another queue for other devices accessing the SPE through the main-storage space. Each MFC can process multiple in-progress DMA commands. Each MFC can also autonomously manage a sequence of DMA transfers in response to a DMA list command from its associated SPU (but not from the PPE or other SPEs). Each DMA command is tagged with a tag group ID that allows software to check or wait on the completion of commands in a particular tag group.

The MFCs support naturally aligned DMA transfer sizes of 1, 2, 4, or 8 bytes, and multiples of 16 bytes, with a maximum transfer size of 16 KB per DMA transfer. DMA list commands can initiate up to 2048 such DMA transfers. Peak transfer performance is achieved if both the effective addresses and the LS addresses are 128-byte aligned and the size of the transfer is an even multiple of 128 bytes.

Each MFC has a synergistic memory management (SMM) unit that processes address-translation and access-permission information supplied by the PPE operating system. To process an effective address provided by a DMA command, the SMM uses essentially the same address-translation and protection mechanism used by the memory management unit (MMU) in the PPE's PowerPC processor storage subsystem (PPSS)[2]. Thus, DMA transfers are coherent with respect to system storage, and the attributes of system storage are governed by the page and segment tables of the *PowerPC Architecture*.

### 1.3.3 Byte Ordering and Bit Numbering

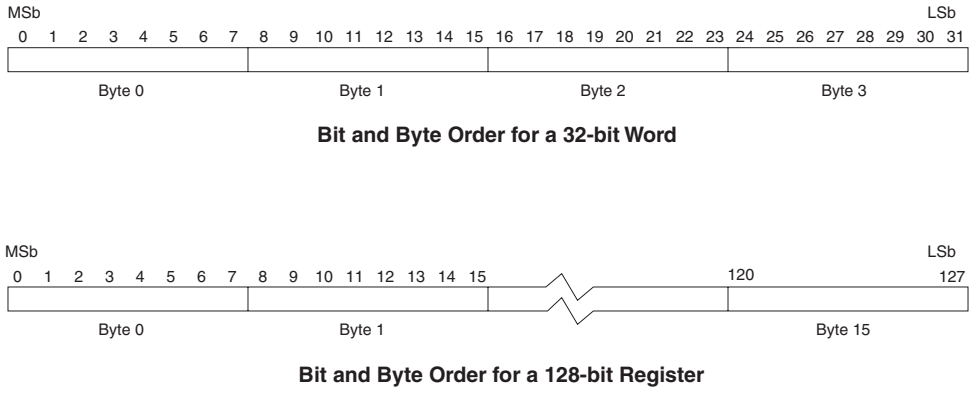
Storage of data and instructions in the CBEA processors uses big-endian ordering, which has the following characteristics:

- Most-significant byte stored at the lowest address, and least-significant byte stored at the highest address.
- Bit numbering within a byte goes from most-significant bit (bit 0) to least-significant bit (bit n). This differs from some other big-endian processors.

A summary of the byte-ordering and bit-ordering in memory and the bit-numbering conventions is shown in *Figure 1-3* on page 49.

2. See *Section 4.3.6* on page 117 for the different way in which the SMM translates real addresses, as compared with the MMU.

*Figure 1-3. Big-Endian Byte and Bit Ordering*




Neither the PPE nor the SPEs, including their MFCs, support little-endian byte-ordering. The MFC's DMA transfers are simply byte moves, without regard to the numeric significance of any byte. Thus, the big-endian or little-endian issue becomes irrelevant to the actual movement of a block of data. The byte-order mapping only becomes significant when data is loaded or interpreted by a processor element or an MFC.

### 1.3.4 Runtime Environment

The PPE runs PowerPC Architecture applications and operating systems, which can include both PowerPC Architecture instructions and vector/SIMD multimedia extension instructions. To use all of the CBEA processor features, the PPE requires an operating system that supports these features, such as multiprocessing with the SPEs, access to the PPE vector/SIMD multimedia extension operations, the interrupt controller, and all the other functions provided by the CBEA processors. This handbook is system-independent, in the sense of making no assumptions about specific operating-system or development-tool environments.

It is common to run a main program on the PPE that allocates threads to the SPEs. In such an application, the *main thread* is said to spawn one or more *tasks*. A task has one or more main threads associated with it, along with some number of *SPE threads*. An SPE thread is a thread that is spawned to run on an available SPE. These terms are defined in *Table 1-1*. The software threads are unrelated to the hardware multithreading capability of the PPE.

*Table 1-1. Definition of Threads and Tasks*

| Term | Definition |
|---|---|
| Main Thread | A thread running on the PPE. |
| Task | A task running on the PPE and SPE. Each such task:<br>• Has one or more main threads and some number of SPE threads.<br>• All the main threads within the task share the task's resources, including access to the SPE threads. |
| SPE Thread | A thread running on an SPE. Each such thread:<br>• Has its own 128 × 128-bit register file, program counter and MFC-DMA command queues.<br>• Can communicate with other execution units (or with main storage through the MFC channel interface). |

A main thread can interact directly with an SPE thread through the SPE's LS. It can interact indirectly through the main-storage space. A thread can poll or sleep, waiting for SPE threads.

The operating system defines the mechanism and policy for selecting an available SPE. It must prioritize among all the applications in the system, and it must schedule SPE execution independently from regular main threads. The operating system is also responsible for runtime loading, passing parameters to SPE programs, notification of SPE events and errors, and debugger support.

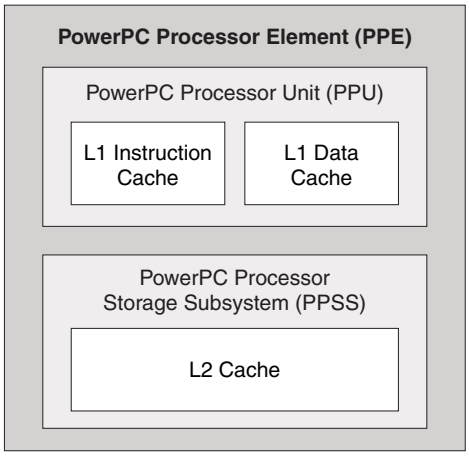
The next two sections provide an expanded overview of the hardware and programming environments for the PPE (*Section 2* on page 51) and the SPEs (*Section 3* on page 65).

# 2. PowerPC Processor Element

The PowerPC Processor Element (PPE) is a general-purpose, dual-threaded, 64-bit RISC processor that conforms to the *PowerPC Architecture*, version 2.02, with the vector/SIMD multimedia extensions. Programs written for the PowerPC 970 processor, for example, should run without modification on the Cell Broadband Engine Architecture (CBEA) processors.[1]

The PPE is responsible for overall control of a system. It runs the operating systems for all applications running on the PPE and Synergistic Processor Elements (SPEs). The PPE consists of two main units, the PowerPC processor unit (PPU) and the PowerPC processor storage subsystem (PPSS), shown in *Figure 2-1*.

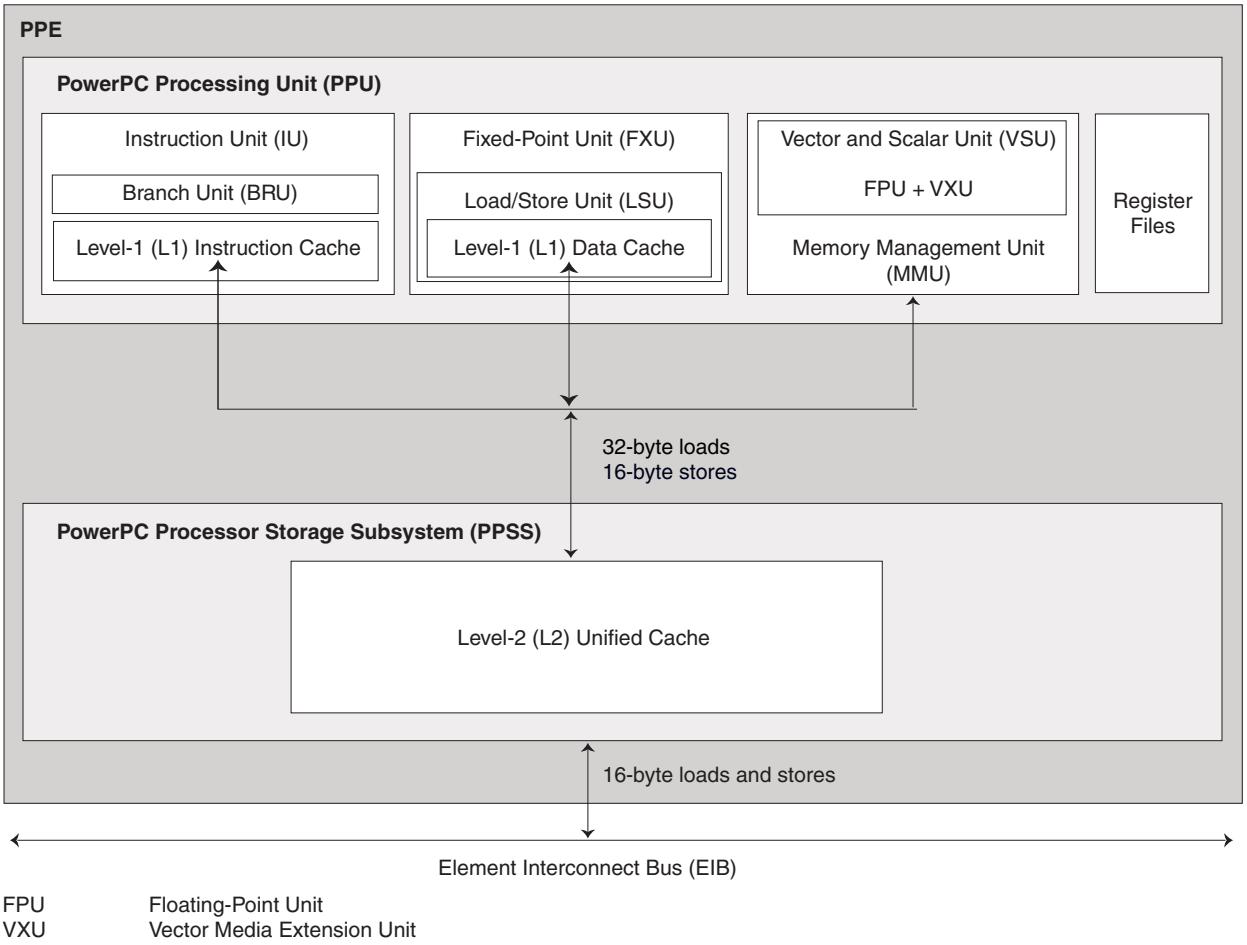
*Figure 2-1. PPE Block Diagram*




The PPU performs instruction execution. Its has a level-1 (L1) instruction cache and data cache and six execution units. It can load 32 bytes and store 16 bytes, independently and memory-coherently, per processor cycle. The PPSS handles memory requests from the PPU and external requests to the PPE from SPEs or I/O devices. It has a unified level-2 (L2) instruction and data cache. The PPU and the PPSS and their functional units are shown in *Figure 2-2* on page 52. The sections that follow describe their functions.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

*Figure 2-2. PPE Functional Units*




## 2.1 PowerPC Processor Unit

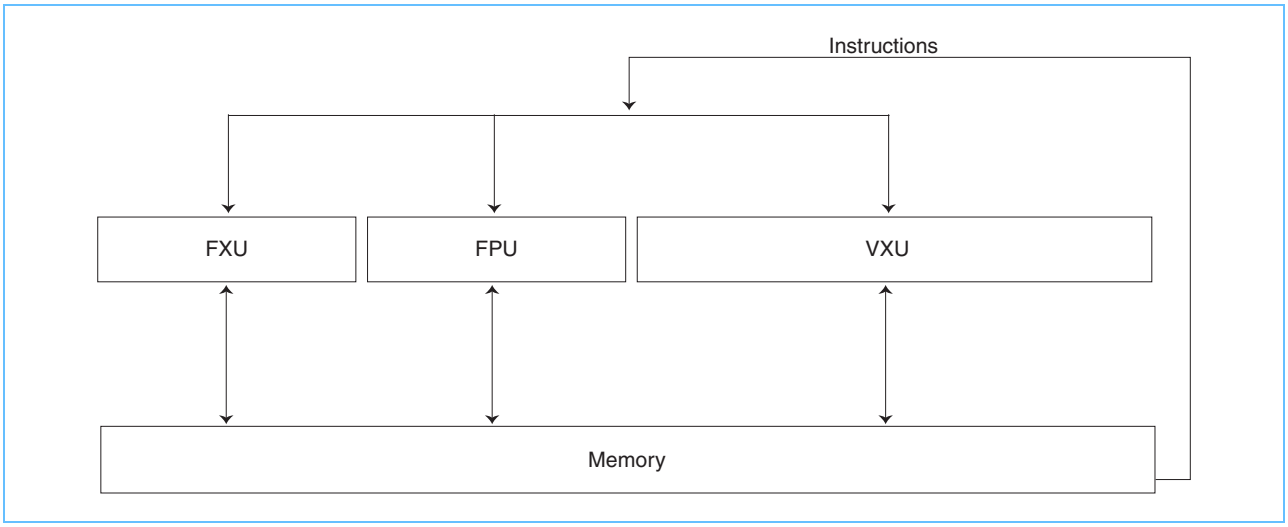
The PPU executes the PowerPC Architecture instruction set and the vector/SIMD multimedia extension instructions. It has duplicate sets of the PowerPC and vector user-state register files (one set for each thread) plus one set of the following functional units:

- *Instruction Unit (IU)*—The IU performs the instruction-fetch, decode, dispatch, issue, branch, and completion portions of execution. It contains the L1 instruction cache, which is 32 KB, 2-way set-associative, reload-on-error, and parity protected. The cache-line size is 128 bytes.
- *Load and Store Unit (LSU)*—The LSU performs all data accesses, including execution of load and store instructions. It contains the L1 data cache, which is 32 KB, 4-way set-associative, write-through, and parity protected. The cache-line size is 128 bytes.
- *Vector/Scalar Unit (VSU)*—The VSU includes a floating-point unit (FPU) and a 128-bit vector/SIMD multimedia extension unit (VXU), which together execute floating-point and vector/SIMD multimedia extension instructions.

- *Fixed-Point Unit (FXU)*—The FXU executes fixed-point (integer) operations, including add, multiply, divide, compare, shift, rotate, and logical instructions.
- *Memory Management Unit (MMU)*—The MMU manages address translation for all memory accesses. It has a 64-entry segment lookaside buffer (SLB) and 1024-entry, unified, parity-protected translation lookaside buffer (TLB). It supports three simultaneous page sizes—4 KB, plus two sizes selectable from 64 KB, 1 MB, or 16 MB.

The 128-bit VXU operates concurrently with the FXU and the FPU, as shown *Figure 2-3* and subject to limitations described in *Appendix A.5* on page 760. All vector/SIMD multimedia extension instructions are designed to be easily pipelined. Parallel execution with the fixed-point and floating-point instructions is simplified by the fact that vector/SIMD multimedia extension instructions do not generate exceptions (other than data-storage interrupts on loads and stores), do not support complex functions, and share few resources or communication paths with the other PPE execution units. A description of the PPE instructions, the units on which they execute, and their latencies is given in *Table A-1 PowerPC Instructions by Execution Unit* on page 723 and *Table A-8 Vector/SIMD Multimedia Extension Instructions* on page 748.

*Figure 2-3. Concurrent Execution of Fixed-Point, Floating-Point, and Vector Extension Units*




The MMU manages virtual-memory address translation and memory protection in compliance with the 64-bit PowerPC Architecture. The MMU supports the following address-space sizes:

- *Real Address (RA) Space*—$2^{42}$ bytes. These are the addresses, in real storage or on an I/O device, whose physical addresses have been mapped by privileged PPE software to the RA space. Real storage can include on-chip SPE local storage (LS) memory and off-chip memory and I/O devices.
- *Effective Address (EA) Space*—$2^{64}$ bytes. These are the addresses generated by programs. *Figure 1-2 Storage and Domains and Interfaces* on page 47 illustrates the effective address space.
- *Virtual Address (VA) Space*—$2^{65}$ bytes. These are the addresses used by the MMU in the PPE and in the memory flow controller (MFC) of each SPE to translate between EAs and RAs.

## 2.2 PowerPC Processor Storage Subsystem

The PPSS handles all memory accesses by the PPU and memory-coherence (snooping) operation from the element interconnect bus (EIB). The PPSS has a unified, 512 KB, 8-way set-associative, write-back L2 cache with error-correction code (ECC). Like the L1 caches, the cache-line size for the L2 is 128 bytes. The cache has a single-port read/write interface to main storage that supports eight software-managed data-prefetch streams. It includes the contents of the L1 data cache but is not guaranteed to contain the contents of the L1 instruction cache, and it provides fully coherent symmetric multiprocessor (SMP) support.

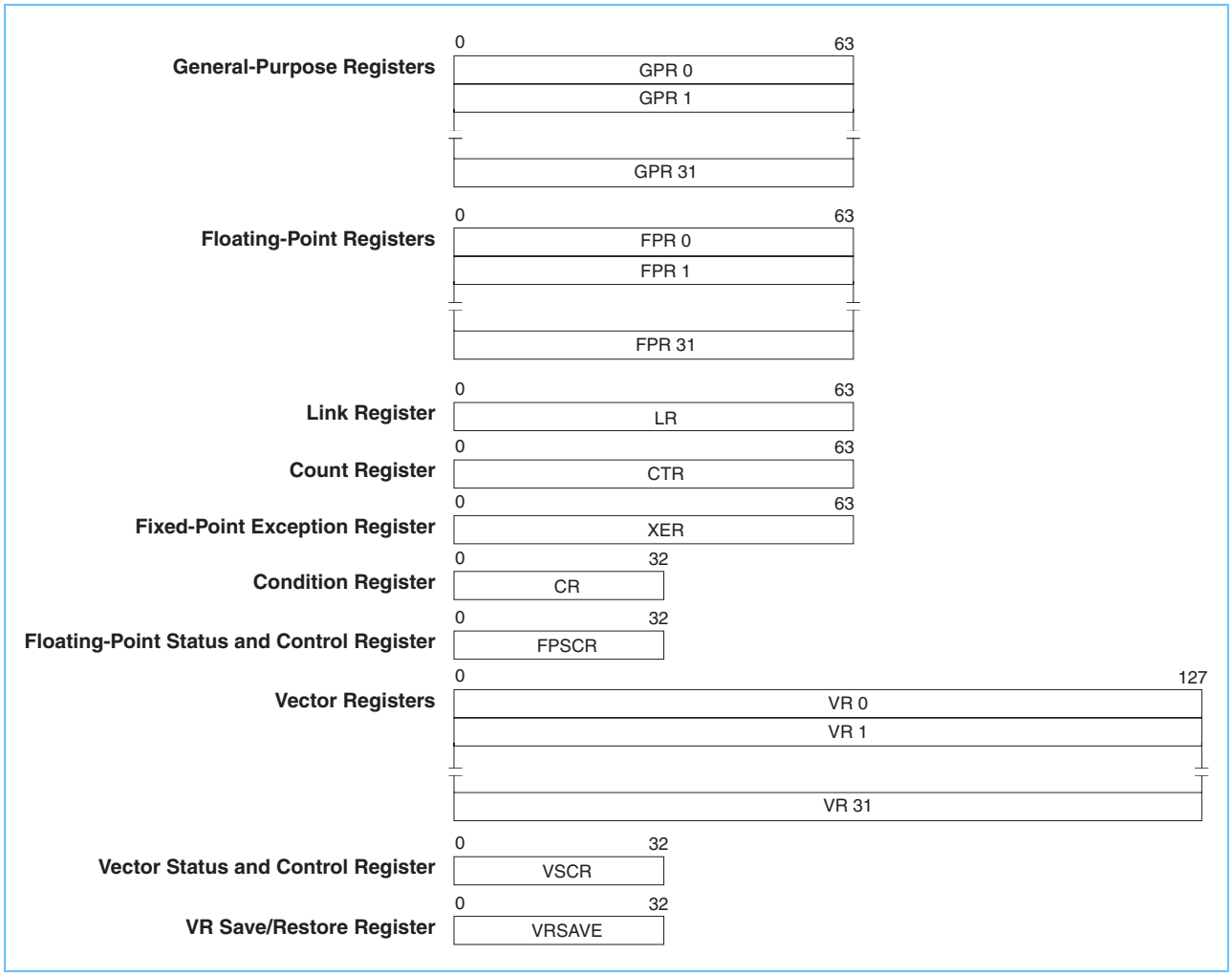
The PPSS performs data-prefetch for the PPU and bus arbitration and pacing onto the EIB. Traffic between the PPU and PPSS is supported by a 32-byte load port (shared by MMU, L1 instruction cache, and L1 data cache requests), and a 16-byte store port (shared by MMU and L1 data cache requests).

The interface between the PPSS and EIB supports 16-byte load and 16-byte store buses. One storage access occurs at a time, and all accesses appear to occur in program order. The interface supports resource allocation management (*Section 8* on page 203), which allows privileged software to control the amount of time allocated to various resource groups. The L2 cache and the TLB use replacement management tables (*Section 6.3* on page 154), which allow privileged software to control the use of the L2 and TLB. This software control over cache and TLB resources is especially useful for real-time programming.

## 2.3 PPE Registers

The PPE problem-state (user) registers are shown in *Figure 2-4* on page 55. All computational instructions operate on registers; no computational instructions modify main storage. To use a storage operand in a computation and then modify the same or another storage location, the contents of the storage operand must be loaded into a register, modified, and then stored back to the target location.

*Figure 2-4. PPE User Register Set*




The PPE registers include:

- *General-Purpose Registers (GPRs)*—The 32 GPRs are 64 bits wide. Fixed-point instructions operate on the full 64 bits. These instructions are mode-independent, except that in 32-bit mode[2], the processor uses only the low-order 32 bits for determining a memory address and the carry, overflow, and record status bits.
- *Floating-Point Registers (FPRs)*—The 32 FPRs are 64 bits wide. The internal format of floating-point data is the IEEE 754 double-precision format. Single-precision results are maintained internally in the double-precision format.
- *Link Register (LR)*—The 64-bit LR can be used to hold the effective address of a branch target. Branch instructions with the link bit (LK) set to '1' (that is, subroutine-call instructions) copy the next instruction address into the LR. The contents of a GPR can be copied to the LR using a move to special-purpose register instruction or from the LR using a move from special-purpose register instruction.

2. Machine State Register `MSR[SF]` bit = '0'.

- *Count Register (CTR)*—The 64-bit CTR can be used to hold either a loop counter or the effective address of a branch target. Some conditional-branch instruction forms decrement the CTR and test it for a '0' value. The contents of a GPR can be copied to the CTR using a move to special-purpose register instruction or from the CTR using a move from special-purpose register instruction.
- *Fixed-Point Exception Register (XER)*—The 64-bit XER contains the carry and overflow bits and the byte count for the move-assist instructions. Most arithmetic operations have instruction forms for setting the carry bit and overflow bit.
- *Condition Register (CR)*—Conditional comparisons are performed by first setting a condition code in the 32-bit CR with a compare instruction or with a recording instruction. The condition code is then available as a value or can be tested by a branch instruction to control program flow. The CR consists of eight independent 4-bit fields grouped together for convenient save or restore during a context switch. Each field can hold status information from a comparison, arithmetic, or logical operation. The compiler can schedule CR fields to avoid data hazards in the same way that it schedules use of GPRs. Writes to the CR occur only for instructions that explicitly request them; most operations have recording and nonrecording instruction forms.
- *Floating-Point Status and Control Register (FPSCR)*—The processor updates the 32-bit FPSCR after every PowerPC (but not vector/SIMD multimedia extension) floating-point operation to record information about the result and any associated exceptions. The status information required by IEEE 754 is included, plus some additional information for exception handling.
- *Vector Registers (VRs)*—There are 32 128-bit-wide VRs. They serve as source and destination registers for all vector instructions.
- *Vector Status and Control Register (VSCR)*—The 32-bit VSCR is read and written in a manner similar to the FPSCR. It has 2 defined bits, a non-Java™ mode bit and a saturation bit; the remaining bits are reserved. Special instructions are provided to move the VSCR from and to a VR.
- *Vector Save Register (VRSAVE)*—The 32-bit VRSAVE register assists user and privileged software in saving and restoring the architectural state across context switches.

The PPE hardware supports two simultaneous threads of execution. All architected and special-purpose registers are duplicated, except those that deal with system-level resources such as logical partitions (*Section 11* on page 331), memory, and thread-control (*Section 10* on page 299). Most nonarchitected resources, such as caches and queues, are shared by both threads, except in cases where the resource is small or offers a critical performance improvement to multithreaded applications. Because of this duplication of state, the PPE can be viewed as a 2-way multiprocessor with shared dataflow. The two hardware-thread environments appear to software as two independent processors.

For details on the complete set of PPE registers—both problem-state (user) and privilege-state (supervisor)—see the following documents:

- *Cell Broadband Engine Registers*
- *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors*
- *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual*

## 2.4 PowerPC Instructions

The PPE's instruction set is compatible with version 2.0.2 of the *PowerPC Architecture* instruction set, with vector/SIMD multimedia extensions, a few new instructions, new meaning for a few existing *PowerPC Architecture* instructions, and implementation of various optional *PowerPC Architecture* instructions. C/C++ intrinsics for the vector/SIMD multimedia extension instructions are also provided.

This section gives a brief overview of the *PowerPC Architecture* instructions. *Section 2.5* on page 59 gives an overview of the vector/SIMD multimedia extension instructions, and *Section 2.6* on page 62 gives an overview of the C/C++ intrinsics. *Appendix A PPE Instruction Set and Intrinsics* on page 723 gives a more complete summary of the instructions and intrinsics, including details about new instructions, compatibility with *PowerPC Architecture* instructions, issue rules, pipeline stages, and compiler optimizations.

### 2.4.1 Data Types

The PowerPC instruction set defines the data types shown in *Table 2-1*. Additional 128-bit vector and scalar data types defined in the vector/SIMD multimedia extensions are described in *Section 2.5.2* on page 61. For details about alignment and performance characteristics for misaligned loads, see *Table A-2 Storage Alignment for PowerPC Instructions* on page 733.

*Table 2-1. PowerPC Data Types*

| Data Types | Width (Bits) |
|---|---|
| **Fixed-Point** | |
| Byte (signed and unsigned) | 8 |
| Halfword (signed and unsigned) | 16 |
| Word (signed and unsigned) | 32 |
| Doubleword (signed and unsigned) | 64 |
| **Floating-Point** | |
| Single-precision | 32 |
| Double-precision | 64 |

### 2.4.2 Addressing Modes

Whenever instruction addresses are presented to the processor, the low-order 2 bits are ignored. Similarly, whenever the processor develops an instruction address, the low-order 2 bits are '0'. The address of either an instruction or a multiple-byte data value is its lowest-numbered byte. This address points to the most-significant byte (the big-endian convention).

All instructions, except branches, generate addresses by incrementing a program counter. For load and store instructions that specify a base register, the effective address in memory for a data value is calculated relative to the base register in one of three ways:

- *Register + Displacement*—The displacement (D) forms of the load and store instructions form the sum of a displacement specified by the sign-extended 16-bit immediate field of the instruction plus the contents of the base register.

- *Register + Register*—The indexed (X) forms of the load and store instructions form the sum of the contents of the index register plus the contents of the base register.
- *Register*—The load string immediate and store string immediate instructions use the unmodified contents of the base register.

Loads and stores can specify an update form that reloads the base register with the computed address. Loads and stores also have byte-reverse forms.

Branches are the only instructions that explicitly specify the address of the next instruction. A branch instruction specifies the effective address of the branch target in one of the following ways:

- *Branch Not Taken*—The byte address of the next instruction is the byte address of the current instruction, plus 4.
- *Absolute*—Branch instructions to absolute addresses transfer control to the word address given in an immediate field of the branch instruction.
- *Relative*—Branch instructions to relative addresses transfer control to the word address given by the sum of the immediate field of the branch instruction and the word address of the branch instruction itself.
- *Link Register* or *Count Register*—The branch conditional to link register instruction transfers control to the effective byte address of the branch target specified in the Link Register; the branch conditional to count register instruction transfers control to the effective byte address of the branch target specified in the Count Register.

### 2.4.3 Instructions

All PowerPC instructions are 4 bytes long and aligned on word (4-byte) boundaries. Most instructions can have up to three operands. Most computational instructions specify two source operands and one destination operand. Signed integers are represented in twos-complement form.

The instructions include the following types:

- *Load and Store Instructions*—These include fixed-point and floating-point load and store instructions, with byte-reverse, multiple, and string options for the fixed-point loads and stores. The fixed-point loads and stores support byte, halfword, word, and doubleword operand accesses between storage and the 32 general-purpose registers (GPRs). The floating-point loads and stores support word and doubleword operand accesses between storage and the 32 floating-point registers (FPRs). The byte-reverse forms have the effect of loading and storing data in little-endian order, although the CBEA processors do not otherwise support little-endian order.
- *Fixed-Point Instructions*—These include arithmetic, compare, logical, and rotate/shift instructions. They operate on byte, halfword, word, and doubleword operands.
- *Floating-Point Instructions*—These include floating-point arithmetic, multiply-add, compare, and move instructions, as well as instructions that affect the Floating-Point Status and Control Register (FPSCR). Floating-point instructions operate on single-precision and double-precision floating-point operands.
- *Memory Synchronization Instructions*—These instructions control the order in which memory operations are completed with respect to asynchronous events, and the order in which memory operations are seen by other processors or memory-access mechanisms. The instruction

types include load and store with reservation, synchronization, and enforce in-order execution of I/O. They are especially useful for multiprocessing. For details about synchronization, see *Section 20 Shared-Storage Synchronization* on page 561.

- *Flow Control Instructions*—These include branch, Condition-Register logical, trap, and other instructions that affect the instruction flow.
- *Processor Control Instructions*—These instructions are used for synchronizing memory accesses and managing caches, TLBs, segment registers, and other privileged processor state. They include move-to/from special-purpose register instructions.
- *Memory and Cache Control Instructions*—These instructions control caches, TLBs, and segment registers.

*Table A-1 PowerPC Instructions by Execution Unit* on page 723 lists the PowerPC instructions, with their latencies, throughputs, and execution units. For details of the instruction pipeline stages, see *Appendix A.6 Pipeline Stages* on page 762. For a complete description of the PowerPC instructions, see the *PowerPC Microprocessor Family: Programming Environments Manual for 64-Bit Microprocessors*.

## 2.5 Vector/SIMD Multimedia Extension Instructions

Vector/SIMD multimedia extension instructions can be freely mixed with PowerPC instructions in a single program. The 128-bit vector/SIMD multimedia extension unit (VXU) operates concurrently with the PPU's 32-bit (or 64-bit in 64-bit mode) fixed-point unit (FXU) and 64-bit floating-point unit (FPU). Like PowerPC instructions, the vector/SIMD multimedia extension instructions are four bytes long and word-aligned. The vector/SIMD multimedia extension instructions support simultaneous execution on multiple elements that make up the 128-bit vector operands.

Vector/SIMD multimedia extension instructions do not generate exceptions (other than data-storage interrupt exceptions on loads and stores), do not support complex functions, and share few resources or communication paths with the other PPE execution units.

The sections that follow briefly summarize features of the extension. For a more detailed list of instructions, see *Appendix A.3* on page 748. For a complete description of the vector/SIMD instruction set, see *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual*.

### 2.5.1 SIMD Vectorization

A *vector* is an instruction operand containing a set of data elements packed into a one-dimensional array. The elements can be fixed-point or floating-point values. Most vector/SIMD multimedia extension and synergistic processor unit (SPU) instructions operate on vector operands. Vectors are also called single-instruction, multiple-data (*SIMD) operands*, or *packed operands*.

SIMD processing exploits data-level parallelism. Data-level parallelism means that the operations required to transform a set of vector elements can be performed on all elements of the vector at the same time. That is, a single instruction can be applied to multiple data elements in parallel.

Support for SIMD operations is pervasive in the CBEA processors. In the PPE, they are supported by the vector/SIMD multimedia extension instructions, described in the sections that follow. In the SPEs, they are supported by the SPU instruction set, described in *Section 3.3* on page 76.

In both the PPE and SPEs, vector registers hold multiple data elements as a single vector. The data paths and registers supporting SIMD operations are 128 bits wide, corresponding to four full 32-bit words. This means that four 32-bit words can be loaded into a single register, and, for example, added to four other words in a different register in a single operation. *Figure 2-5* shows such an operation. Similar operations can be performed on vector operands containing 16 bytes, 8 halfwords, or 2 doublewords.

*Figure 2-5. Four Concurrent Add Operations*

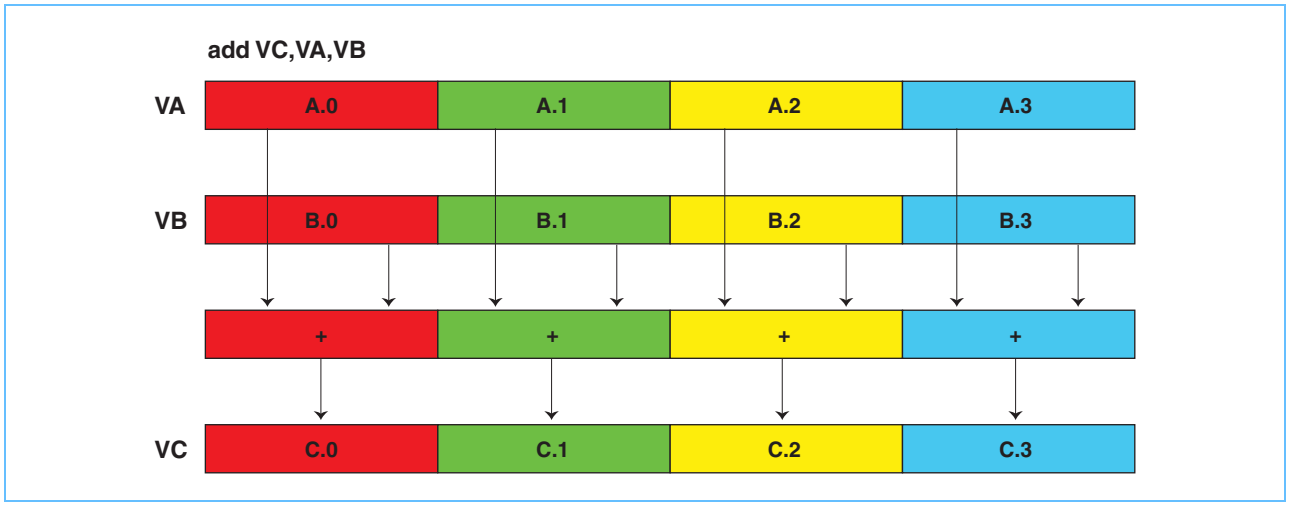
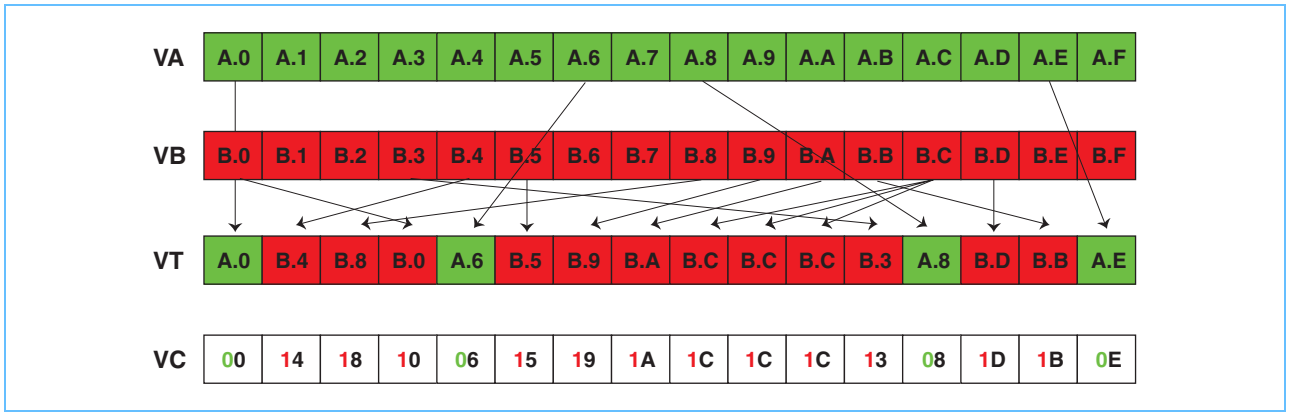



*Figure 2-6* shows another example of a SIMD operation—in this case, a byte-shuffle (permute) operation. Here, the bytes selected for the shuffle from the source registers, VA and VB, are based on byte entries in the control vector, VC, in which a 0 specifies VA and a 1 specifies VB. The result of the shuffle is placed in register VT. The shuffle operation is extremely powerful and finds its way into many applications in which data reordering, selection, or merging is required.

*Figure 2-6. Byte-Shuffle (Permute) Operation*



The process of preparing a program for use on a vector processor is called *vectorization* or *SIMDization*. It can be done manually by the programmer, or it can be done by a compiler that does *auto-vectorization*. See *Section 22 SIMD Programming* on page 629 for details.

### 2.5.2 Data Types

The operands for most vector/SIMD multimedia extension instructions are quadword (128-bit) vectors. These vectors can contain elements that are (a) fixed-point bytes, halfwords, or words, or (b) floating-point words.

The vector registers are 128 bits wide and can contain:

- Sixteen 8-bit values, signed or unsigned
- Eight 16-bit values, signed or unsigned
- Four 32-bit values, signed or unsigned
- Four single-precision IEEE-754 floating point values

The C/C++ intrinsics for the vector/SIMD multimedia extension define C/C++ language data types for vectors, as described in *Section 2.6* on page 62.

### 2.5.3 Addressing Modes

The PPE supports not only basic load and store operations but also load and store vector left or right indexed forms. All vector/SIMD multimedia extension load and store operations use the register + register indexed addressing mode, which forms the sum of the contents of an index GPR plus the contents of a base-address GPR. This addressing mode is very useful for accessing arrays.

In addition to the load and store operations, the vector/SIMD multimedia extension instructions provide a powerful set of element-manipulation instructions—for example, permute (similar to the SPEs' shuffle, *Figure 2-6* on page 60), rotate, and shift—to manipulate vector elements into the required alignment and arrangement after the vectors have been loaded into vector registers.

### 2.5.4 Instruction Types

Most vector/SIMD multimedia extension instructions have three or four 128-bit vector operands—two or three source operands and one result. Also, most instructions are SIMD in nature. The instructions have been chosen for their utility in digital signal processing (DSP) algorithms, including 3D graphics.

The vector/SIMD multimedia extension instructions include the following types:

- *Vector Fixed-Point Instructions*—Vector arithmetic (including saturating arithmetic), compare, logical, rotate, and shift instructions. They operate on byte, halfword, and word vector elements.
- *Vector Floating-Point Instructions*—Floating-point, multiply/add, rounding and conversion, compare, and estimate instructions. They operate on single-precision floating-point vector elements.

- *Vector Load and Store Instructions*—Basic fixed-point and floating-point load and store instructions. No update forms of the load and store instruction are provided. They operate on 128-bit vectors.
- *Vector Permutation and Formatting Instructions*—Vector pack, unpack, merge, splat, permute, select, and shift instructions.
- *Processor Control Instructions*—Instructions that read and write the Vector Status and Control Register (VSCR).
- *Memory Control Instructions*—Instructions for managing caches (user-level and supervisor-level). These instructions are no-ops on the CBEA processors.

### 2.5.5 Instructions

*Table A-8 Vector/SIMD Multimedia Extension Instructions* on page 748 lists the vector/SIMD multimedia extension instructions, with their pipelines, latencies, and throughputs. For a complete description of the instructions, see *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual*.

### 2.5.6 Graphics Rounding Mode

The CBEA processors implement a graphics rounding mode that allows programs written with vector/SIMD multimedia extension instructions to produce floating-point results that are equivalent in precision to those written in the SPU instruction set. In this vector/SIMD multimedia extension mode, as in the SPU environment, the default rounding mode is *round-to-zero*, denormals are treated as zero, and there are no infinities or NaNs.

For details about the graphics rounding mode, see *Section A.3.3* on page 752.

## 2.6 Vector/SIMD Multimedia Extension C/C++ Language Intrinsics

Both the vector/SIMD multimedia extension and SPU instructions are supported by C/C++ language extensions that define a set of vector data types and vector intrinsics. The intrinsics are essentially inline assembly-language instructions in the form of C-language function calls. The intrinsics provide explicit control of the underlying assembly instructions without directly managing registers and scheduling instructions, as assembly-language programming requires. A compiler that supports the C/C++ language extensions for the vector/SIMD multimedia extension instructions will emit code optimized for the Vector/SIMD Multimedia Extension Architecture.

*Section A.4* on page 754 and the *C/C++ Language Extensions for Cell Broadband Engine Architecture* specification provide details about the intrinsics, including a complete list of the intrinsics. Although the intrinsics provided by the PPE and SPE instruction sets are similar in function, their naming conventions and function-call forms are different. The vector/SIMD multimedia extension intrinsics are summarized in this section. The SPE intrinsics are summarized in *Section 3.4* on page 77.

### 2.6.1 Vector Data Types

The vector/SIMD multimedia extension C/C++ language extensions define a set of fundamental data types, called *vector types*. These are summarized in *Table A-10* on page 754.

Introducing fundamental vector data types permits a compiler to provide stronger type-checking and supports overloaded operations on vector types. Because the token, *vector*, is a keyword in the vector/SIMD multimedia extension data types, it is recommended that the term not be used elsewhere in the program as, for example, a variable name.

### 2.6.2 Vector Literals

The vector/SIMD multimedia extension C/C++ language extensions define a set of vector literal formats. These are summarized in *Table A-11* on page 755. These formats consist of a parenthesized vector type followed by a parenthesized set of constant expressions.

### 2.6.3 Intrinsics

The vector/SIMD multimedia extension intrinsics are summarized in *Table A-13* on page 757.

#### 2.6.3.1 *Classes*

The vector/SIMD multimedia extension intrinsics map directly to one or more vector/SIMD multimedia extension assembly-language instructions. The intrinsics are grouped in the following three classes:

- *Specific Intrinsics*—Intrinsics that have a one-to-one mapping with a single assembly-language instruction. The specific intrinsics are rarely used.
- *Generic Intrinsics*—Intrinsics that map to one or more assembly-language instructions as a function of the type of input parameters. In general, the generic class of intrinsics contains the most useful intrinsics.
- *Predicate Intrinsics*—Intrinsics that compare values and return an integer that can be used directly as a value or as a condition for branching.

The specific and generic intrinsics use the prefix, “`vec_`” in front of an assembly-language or operation mnemonic; predicate intrinsics use the prefixes “`vec_all`” and “`vec_any`”. When complied, the intrinsics generate one or more vector/SIMD multimedia extension assembly-language instructions.

Vector/SIMD multimedia extension intrinsics can be used anywhere in a C/C++ language program. There is no need for setup or to enter a special mode.

#### 2.6.3.2 *Example*

The following code example contains two versions of a function that sums an input array of 16 byte values. In such an array-summing function, one could unroll the scalar code to slightly improve the performance or one could use the vector/SIMD multimedia extension intrinsics to significantly improve its performance and eliminate the loop entirely.

The first version follows; it performs 16 iterations of the loop. The second version uses vector/SIMD multimedia extension intrinsics to eliminate the loop entirely.

```
// 16 iterations of a loop
int rolled_sum(unsigned char bytes[16])
{
```

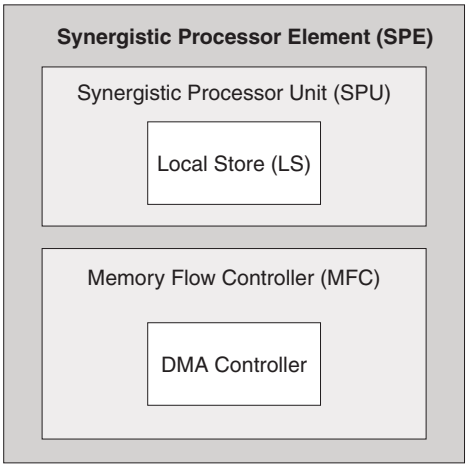
```
    int i;
    int sum = 0;
    for (i = 0; i < 16; ++i)
    {
        sum += bytes[i];
    }
    return sum;
}


// Vectorized for vector/SIMD multimedia extension
int vectorized_sum(unsigned char bytes[16])
{
  vector unsigned char vbytes;
  union {
    int i[4];
    vector signed int v;
  } sum;
  vector unsigned int zero = (vector unsigned int){0};

  // Perform a misaligned vector load of the 16 bytes.
  vbytes = vec_perm(vec_ld(0, bytes), vec_ld(16, bytes), vec_lvsl(0, bytes));

  // Sum the 16 bytes of the vector
  sum.v = vec_sums((vector signed int)vec_sum4s(vbytes, zero),
    (vector signed int)zero);

  // Extract the sum and return the result.
  return (sum.i[3]);
}
```

# 3. Synergistic Processor Elements

The eight Synergistic Processor Elements (SPEs) execute a new single instruction, multiple data (SIMD) instruction set—the *Synergistic Processor Unit Instruction Set Architecture*. Each SPE is a 128-bit RISC processor specialized for data-rich, compute-intensive SIMD and scalar applications. It consists of two main units, the synergistic processor unit (SPU) and the memory flow controller (MFC), shown in *Figure 3-1*.

In this document, the term “SPU” refers to the instruction set or the unit that executes the instruction set, and the term “SPE” refers generally to functionality of any part of the SPE processor element, including the MFC.

*Figure 3-1. SPE Block Diagram*




The SPEs provide a deterministic operating environment. They do not have caches, so cache misses are not a factor in their performance. Pipeline-scheduling rules are simple, so it is easy to statically determine the performance of code and for software to generate high-quality, static schedules.

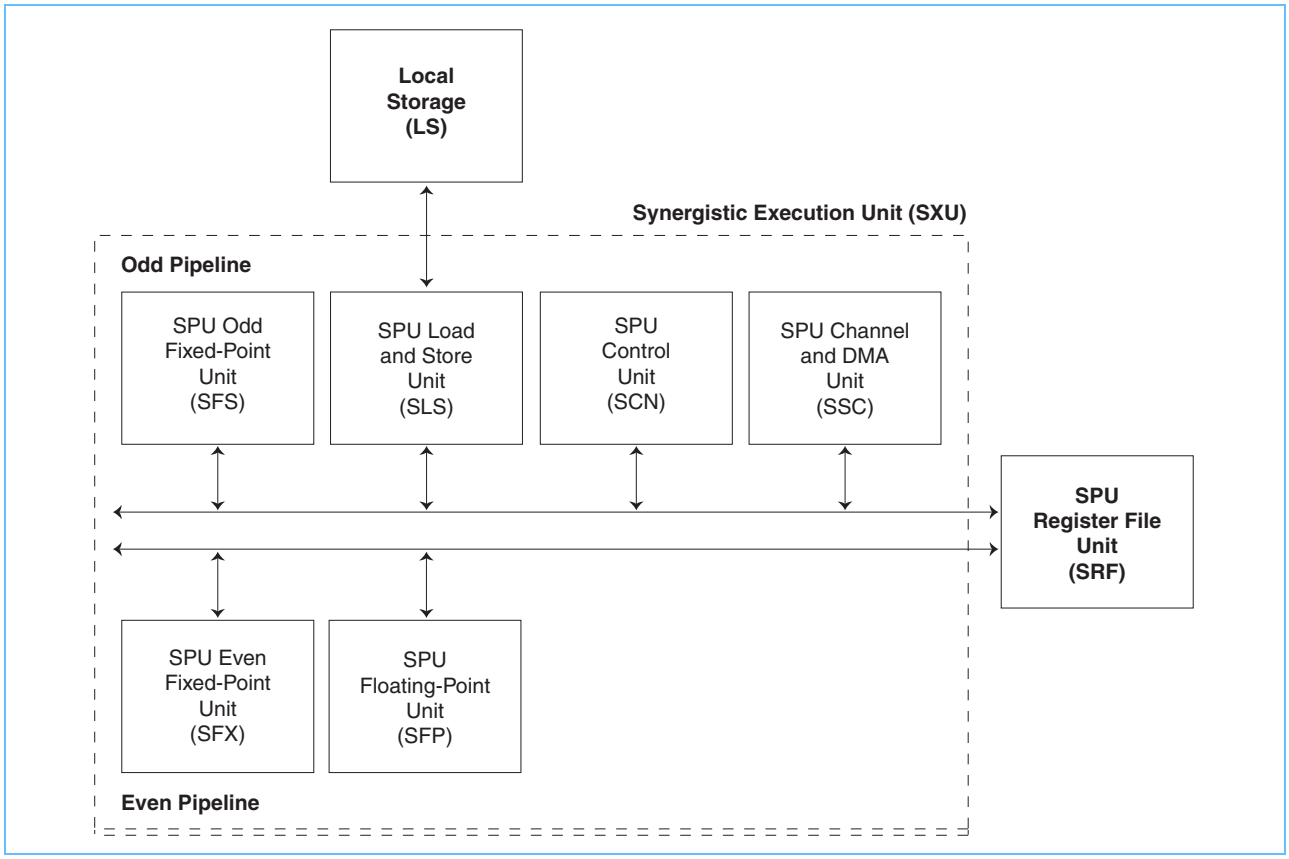
## 3.1 Synergistic Processor Unit

The SPU fetches instructions from its unified (instructions and data) 256 KB local storage (LS), and it loads and stores data between its LS and its single register file for all data types, which has 128 registers, each 128 bits wide. The SPU has four execution units, a DMA interface, and a channel interface for communicating with its MFC, the PowerPC Processor Element (PPE) and other devices (including other SPEs).

Each SPU is an independent processor element with its own program counter, optimized to run SPU programs. The SPU fills its LS by requesting DMA transfers from its MFC, which implements the DMA transfers using its DMA controller. Then, the SPU fetches and executes instructions from its LS, and it loads and stores data to and from its LS.

The main SPU functional units are shown in *Figure 3-2*. These include the synergistic execution unit (SXU), the LS, and the SPU register file unit (SRF). The SXU contains six execution units, which are described in *Section 3.1.3* on page 70.

*Figure 3-2. SPU Functional Units*




The SPU can issue and complete up to two instructions per cycle, one on each of the two (odd and even) execution pipelines. Whether an instruction goes to the odd or even pipeline depends on the instruction type. The instruction type is also related to the execution unit that performs the function, as summarized in *Section 3.3.2* on page 77.

## 3.1.1 Local Storage

The LS is a 256 KB, error-correcting code (ECC)-protected, single-ported, noncaching memory. It stores all instructions and data used by the SPU. It supports one access per cycle from either SPE software or DMA transfers. SPU instruction prefetches are 128 bytes per cycle. SPU data-access bandwidth is 16 bytes per cycle, quadword aligned. DMA-access bandwidth is 128 bytes per cycle. DMA transfers perform a read-modify-write of LS for writes less than a quadword.

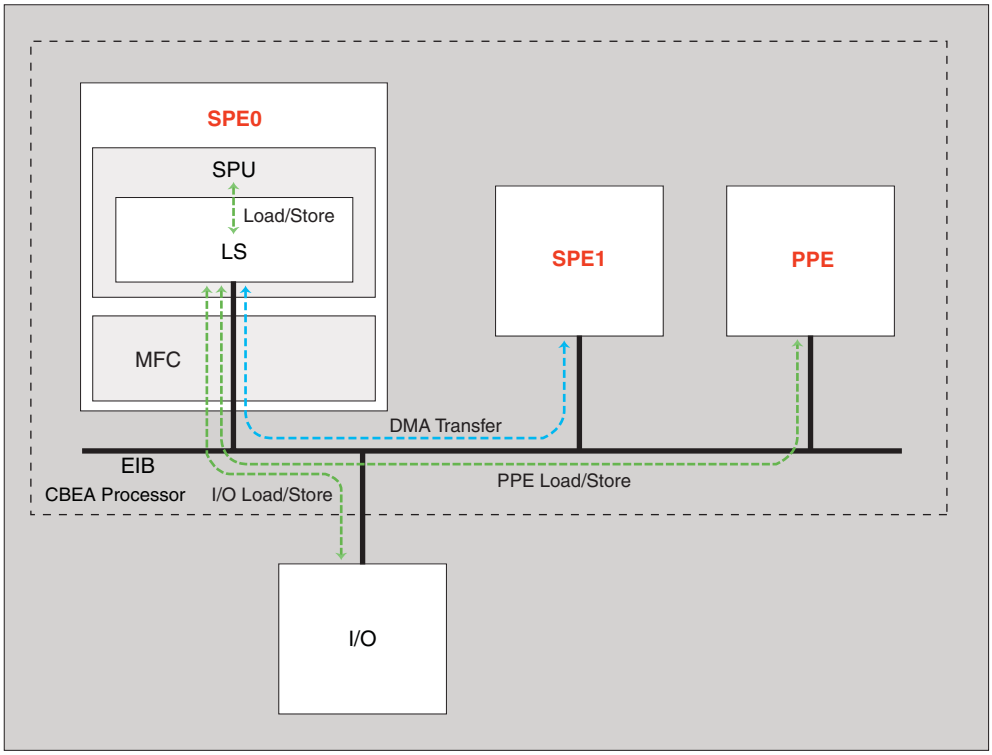
### 3.1.1.1 *Addressing and Address Aliasing*

The SPU accesses its LS with load and store instructions, and it performs no address translation for such accesses. Privileged software on the PPE can assign effective-address aliases to an LS. This enables the PPE and other SPEs to access the LS in the main-storage domain. The PPE

performs such accesses with load and store instructions, without the need for DMA transfers. However, other SPEs must use DMA transfers to access the LS in the main-storage domain. When aliasing is set up by privileged software on the PPE, the SPE that is initiating the request performs address translation, as described in *Section 3.2 Memory Flow Controller* on page 72.

*Figure 3-3* illustrates the methods by which an SPU, the PPE, other SPEs, and I/O devices access the SPU's associated LS, when the LS has been aliased to the main storage domain. See *Figure 1-2 Storage and Domains and Interfaces* on page 47 for an overview of the storage domains. For details about how the PPE does this aliasing, see *Section 4 Virtual Storage Environment* on page 79.

*Figure 3-3. LS Access Methods*




| | | | |
|---|---|---|---|
| DMA | Direct Memory Access | MFC | Memory Flow Controller |
| EIB | Element Interconnect Bus | PPE | PowerPC Processor Element |
| LS | Local Storage | SPE | Synergistic Processor Element |

### 3.1.1.2 *Coherence and Synchronization*

DMA transfers between the LS and main storage are coherent in the system. A pointer to a data structure created on the PPE can be passed by the PPE, through the main-storage space, to an SPU, and the SPU can use this pointer to issue a DMA command to bring the data structure into its LS. Memory-mapped mailboxes or atomic MFC synchronization commands can be used for synchronization and mutual exclusion. See *Section 19* on page 513 and *Section 20* on page 561 for details.

#### 3.1.1.3 *LS Access Priority*

Although the LS is shared between DMA transfers, load and store operations by the associated SPU, and instruction prefetches by the associated SPU, DMA operations are buffered and can only access the LS at most one of every eight cycles. Instruction prefetches deliver at least 17 instructions sequentially from the branch target. Thus, the impact of DMA operations on loads and stores and program-execution times is, by design, limited.

When there is competition for access to the LS by loads, stores, DMA reads, DMA writes, and instruction fetches, the SPU arbitrates access to the LS according the following priorities (highest priority first):

1. DMA reads and writes
2. SPU loads and stores
3. Instruction prefetch

*Table 3-1* summarizes the complete LS-arbitration priorities and transfer sizes. Memory-mapped I/O (MMIO) accesses and DMA transfers always have highest priority. These operations occupy, at most, one of every eight cycles (one of sixteen for DMA reads, and one of sixteen for DMA writes) to the LS. Thus, except for highly optimized code, the impact of DMA reads and writes on LS availability for loads, stores, and instruction fetches can be ignored.

*Table 3-1. LS-Access Arbitration Priority and Transfer Size*

| Transaction | Transfer Size (Bytes) | Priority | Maximum LS Occupancy (SPU Cycle) | Access Path |
|---|---|---|---|---|
| MMIO[1] | ≤ 16 | 1-Highest | 1/8 | Line Interface |
| DMA | ≤ 128 | 1 | | |
| DMA List Transfer-Element Fetch | 128 | 1 | 1/4 | Quadword Interface |
| ECC Scrub[2] | 16 | 2 | 1/10 | |
| SPU Load or Store | 16 | 3 | 1 | |
| Hint Fetch[3] | 128 | 3 | 1 | Line Interface |
| Inline Fetch | 128 | 4-Lowest | 1/16 for inline code | |

1. Access to LS contents by the PPE and other devices (including other SPEs) that can execute loads and stores in the main-storage space.
2. The DMA logic has an ECC-scrub state machine. It is initialized when an ECC error is detected during a load instruction or DMA read, and it accesses every quadword in LS to find and correct every ECC error.
3. See *Section 24.3.3 Branch Hints* on page 701.

After MMIO and DMA accesses, the next-highest priorities are given to ECC scrubbing. This is followed by data accesses by SPU load and store instructions and by hint fetches. The rationale for giving SPU loads and stores higher priority than inline instruction fetches (sequential instruction prefetch) is that loads and stores typically help a program's progress, whereas instruction fetches are often speculative. The SPE supports only 16-byte load and store operations that are 16-byte-aligned. The SPE uses a second instruction (byte shuffle) to place bytes in a different order if, for example, the program requires only a 4-byte quantity or a quantity with a different data alignment. To store data that is not aligned, use a read-modify-write operation.

The lowest priority for LS access is given to inline instruction fetches. Instruction fetches load 32 instructions per SPU request by accessing all banks of the LS simultaneously. Because the LS is single-ported, it is important that DMA and instruction-fetch activity transfer as much useful data as possible in each LS request.

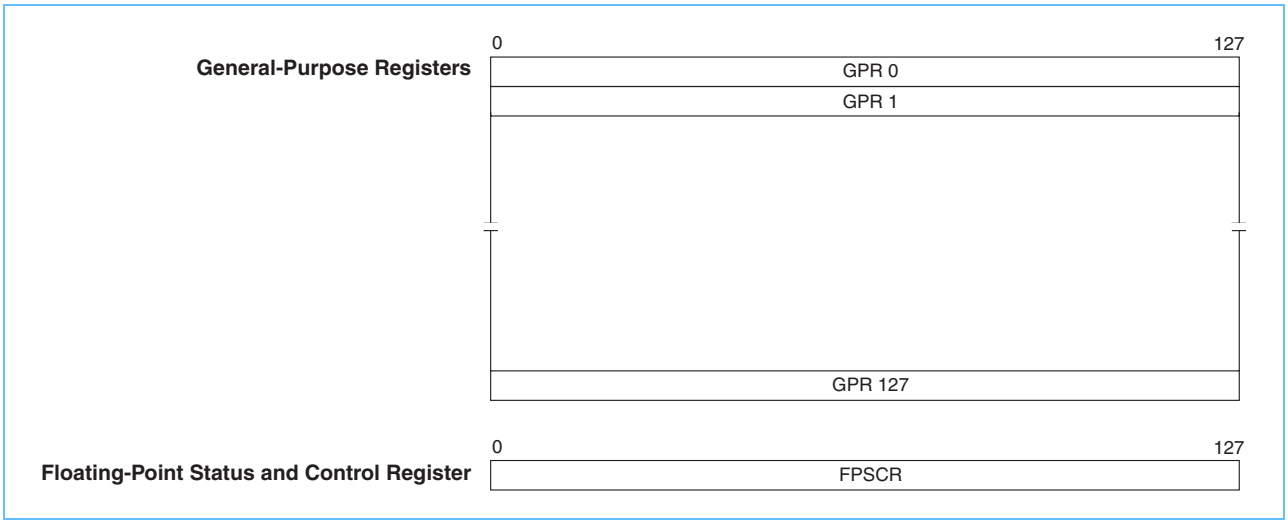
### 3.1.2 Register File

The SPU's 128-entry, 128-bit register file stores all data types—integer, single-precision and double-precision floating-point, scalars, vectors, logicals, bytes, and others. It also stores return addresses, results of comparisons, and so forth. All computational instructions operate on registers—there are no computational instructions that modify storage. The large size of the register file and its unified architecture achieve high performance without the use of expensive hardware techniques such as out-of-order processing or deep speculation.

The complete set of SPE problem-state (user) registers is shown in *Figure 3-4*. These registers include:

- *General-Purpose Registers (GPRs)*—All data types can be stored in the 128-bit GPRs, of which there are 128.
- *Floating-Point Status and Control Register (FPSCR)*—The processor updates the 128-bit FPSCR after every floating-point operation to record information about the result and any associated exceptions.

*Figure 3-4. SPE Problem-State (User) Register Set*




In addition to the problem-state registers shown in *Figure 3-4*, the SPU also supports privilege-state (supervisor) registers that are accessible by privileged software running on the PPE. The SPU does not have any special-purpose registers (SPRs). The MFC, independently of the SPU, supports many MMIO registers and queues. For a summary, see *Section 3.2* on page 72. For details of all SPU and MFC registers, see the *Cell Broadband Engine Registers* specification.

### 3.1.3 Execution Units

The SPU supports dual-issue of instructions to its two execution pipelines, shown in *Figure 3-2* on page 66. The pipelines are referred to as even (pipeline 0) and odd (pipeline 1). The units execute the following types of operations:

- *SPU Odd Fixed-Point Unit (SFS)*—Executes byte granularity shift, rotate mask, and shuffle operations on quadwords.
- *SPU Even Fixed-Point Unit (SFX)*—Executes arithmetic instructions, logical instructions, word SIMD shifts and rotates, floating-point compares, and floating-point reciprocal and reciprocal square-root estimates.
- *SPU Floating-Point Unit (SFP)*—Executes single-precision and double-precision floating-point instructions, integer multiplies and conversions, and byte operations. The SPU supports only 16-bit multiplies for integers, so 32-bit multiplies are implemented in software using 16-bit multiplies. For details on floating-point operations, see *Section 3.1.4*.
- *SPU Load and Store Unit (SLS)*—Executes load and store instructions and hint for branch (HBR) instructions. It also handles DMA requests to the LS.
- *SPU Control Unit (SCN)*—Fetches and issues instructions to the two pipelines, executes branch instructions, arbitrates access to the LS and register file, and performs other control functions.
- *SPU Channel and DMA Unit (SSC)*—Enables communication, data transfer, and control into and out of the SPU. The functions of SSC, and those of the associated DMA controller in the MFC, are described in *Section 3.2* on page 72.

The SPU issues all instructions in program order according to the pipeline assignment. Each instruction is part of a doubleword-aligned instruction pair called a *fetch group*. For details about the rules for fetching and issuing instructions, see *Section B.1.3* on page 779.

### 3.1.4 Floating-Point Support

The SPU supports both single-precision and double-precision floating-point operations. Single-precision instructions are performed in 4-way SIMD fashion. The data formats for single-precision and double-precision instructions are those defined by IEEE Standard 754, but the results calculated by single-precision instructions depart from the IEEE Standard 754 by placing emphasis on real-time graphics requirements that are typical of multimedia processing.

For single-precision operations, the range of normalized numbers is extended beyond the IEEE standard. The representable, positive, nonzero numbers range from Smin = $2^{-126}$ to Smax = $(2 - 2^{-23}) \times 2^{128}$. If the exact result overflows (that is, if it is larger in magnitude than Smax), the rounded result is set to Smax with the appropriate sign. If the exact result underflows (that is, if it is smaller in magnitude than Smin), the rounded result is forced to zero. A zero result is always a positive zero.

Single-precision floating-point operations implement IEEE 754 arithmetic with the following extensions and differences:

- Only one rounding mode is supported: *round towards zero*, also known as *truncation*.
- Denormal operands are treated as zero, and denormal results are forced to zero.
- Numbers with an exponent of all ones are interpreted as normalized numbers and not as infinity or not-a-number (NaN).

Double-precision operations do not support the IEEE precise trap (exception) mode. If a double-precision denormal or not-a-number (NaN) result does not conform to IEEE Standard 754, then the deviation is recorded in a sticky bit in the FPSCR register, which can be accessed using the **fscrrd** and **fscrwr** instructions or the **spu_mffpscr** and **spu_mtfpscr** intrinsics.

Double-precision instructions are performed as two double-precision operations in 2-way SIMD fashion. However, the SPU on the Cell/B.E. processor is capable of performing only one double-precision operation per cycle. Thus, the Cell/B.E. SPU executes double-precision instructions by breaking up the SIMD operands and executing the two operations in consecutive instruction slots in the pipeline. Although double-precision instructions have 13-clock-cycle latencies, on the Cell/B.E. processor, only the final seven cycles are pipelined. No other instructions are dual-issued with double-precision instructions, and no instructions of any kind are issued for six cycles after a double-precision instruction is issued.

The PowerXCell 8i processor is capable of performing two double-precision operations per cycle with a 9-clock-cycle latency. The double-precision instructions are dual-issued with pipe 1 instructions.

*Table 3-2*, *Table 3-3*, and *Table 3-4* summarize the SPE floating-point support.

*Table 3-2. Single-Precision (Extended-Range Mode) Minimum and Maximum Values*

| Number Format | Minimum Positive Magnitude (Smin) | | | Maximum Positive Magnitude (Smax) | | | Notes |
|---|---|---|---|---|---|---|---|
| Register Value | x‘00800000’ | | | x‘7FFFFFFF’ | | | |
| Bit Fields | Sign | 8-Bit Biased Exponent | Fraction (implied [1] and 23 bits) | Sign | 8-Bit Biased Exponent | Fraction (implied [1] and 23 bits) | 1 |
| | 0 | 00000001 | [1.]000...000 | 0 | 11111111 | [1.]111...111 | |
| Value in Powers of 2 | + | $2^{(1 - 127)}$ | 1 | + | $2^{(255 - 127)}$ | $2 - 2^{-23}$ | 2 |
| Combined Exponent and Fraction | $2^{-126} \times (+1)$ | | | $2^{128} \times (+[2 - 2^{-23}])$ | | | |
| Value of Register in Decimal | $1.2 \times 10^{-38}$ | | | $6.8 \times 10^{38}$ | | | |

**Notes:**

1. The exponent field is biased by +127.
2. The value $2 - 2^{-23}$ is one least significant bit (LSb) less than 2.

*Table 3-3. Double-Precision (IEEE Mode) Minimum and Maximum Values*

| Number Format | Minimum Positive Denormalized Magnitude (Dmin) | | | Maximum Positive Normalized Magnitude (Dmax) | | | Notes |
|---|---|---|---|---|---|---|---|
| Register Value | x'0000000000000001' | | | x'7FEFFFFFFFFFFFFF' | | | |
| Bit Fields | Sign | 11-Bit Biased Exponent | Fraction (implied [0] and 52 bits for denormalized number) | Sign | 11-Bit Biased Exponent | Fraction (implied [1] and 52 bits for normalized number) | 1 |
| | 0 | 00000000000 | [0.]000...001 | 0 | 11111111110 | [1.]111...111 | 2 |
| Value in Powers of 2 | + | $2^{(0+1-1023)}$ | $2^{-52}$ | + | $2^{(2046-1023)}$ | $2 - 2^{-52}$ | 3, 4 |
| Combined Exponent and Fraction | $2^{-1022} \times (+2^{-52})$ | | | $2^{1023} \times (+[2 - 2^{-52}])$ | | | |
| Value of Register in Decimal | $4.9 \times 10^{-324}$ | | | $1.8 \times 10^{308}$ | | | |

**Notes:**

1. The exponent is biased by +1023.
2. An exponent field of all ones is reserved for not-a-number (NaN) and infinity.
3. The value $2 - 2^{-52}$ is one LSb less than 2.
4. An extra 1 is added to the exponent for denormalized numbers.

*Table 3-4. Single-Precision (IEEE Mode) Minimum and Maximum Values*

| Number Format | Minimum Positive Denormalized Magnitude (Smin) | | | Maximum Positive Magnitude (Smax) | | | Notes |
|---|---|---|---|---|---|---|---|
| Register Value | x'00000001' | | | x'7F7FFFFF' | | | |
| Bit Fields | Sign | 8-Bit Biased Exponent | Fraction (implied [0] and 23 bits) | Sign | 8-Bit Biased Exponent | Fraction (implied [1] and 23 bits) | 1 |
| | 0 | 00000000 | [0.]000..001 | 0 | 11111110 | [1.]111...111 | |
| Value in Powers of 2 | + | $2^{(0+1-127)}$ | $2^{-23}$ | + | $2^{(254-127)}$ | $2 - 2^{-23}$ | 2 |
| Combined Exponent and Fraction | $2^{-126} \times 2^{-23}$ | | | $2^{127} \times (2 - 2^{-23})$ | | | |
| Value of Register in Decimal | $1.4 \times 10^{-45}$ | | | $3.4 \times 10^{38}$ | | | |

**Notes:**

1. The exponent field is biased by +127.
2. The value $2 - 2^{-23}$ is 1 LSb less than 2.

## 3.2 Memory Flow Controller

Each SPU has its own MFC. The MFC serves as the SPU's interface, by means of the element interconnect bus (EIB), to main-storage (*Figure 1-2* on page 47) and other processor elements and system devices. The MFC's primary role is to interface its LS-storage domain with the main-storage domain. It does this by means of a DMA controller that moves instructions and data between its LS and main storage. The MFC also supports storage protection on the main-storage side of its DMA transfers, synchronization between main storage and the LS, and communication functions (such as mailbox and signal-notification messaging) with the PPE and other SPEs and devices.

*Figure 3-5* shows a block diagram of the MFC.

*Figure 3-5. MFC Block Diagram*

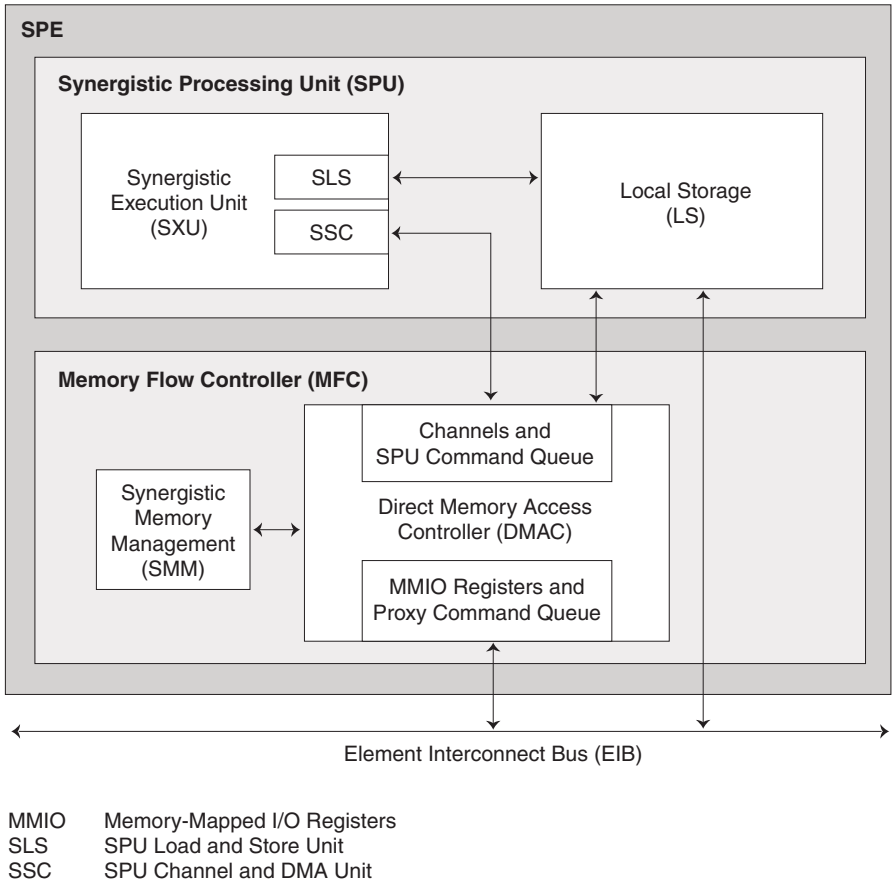



*Figure 3-5* shows the MFC's primary special-function interfaces to the SPU and EIB:

- *LS Read and Write Interface*. This links the LS with the EIB. It is used for DMA transfers (*Section 3.2.4.1* on page 75), atomic transfers (*Section 20.3* on page 597), and snoop requests from the EIB.
- *LS DMA List Element Read Interface*. This links the LS with the MFC's channel and SPU command queue. It is used to implement efficient transfers of DMA lists (*Section 3.2.4.2* on page 75).
- *Channel Interface*. This links the SXU's SPU channel and DMA unit (SSC) with the MFC's channel and SPU command queue. See *Section 3.2.1* on page 74 for an overview of channels.
- *EIB Command and Data Interfaces*. These link the MFC with the EIB. The data interface consists of a 128-bit data-in bus and a 128-bit data-out bus. As an EIB master, the MFC can send up to 16 outstanding MFC commands, and it supports atomic requests. As an EIB subordinate, the MFC supports snoop requests (read and write) of the SPU's LS memory and the MFC's MMIO registers, and atomic requests.

### 3.2.1 Channels

SPE software communicates with main storage, the PPE, and other SPEs and devices through its channels, shown in *Figure 3-5* on page 73 and in *Figure 1-2* on page 47. Channels are unidirectional message-passing interfaces that support 32-bit messages and commands. Each SPE has its own set of channels. SPE software accesses channels with special channel-read and channel-write instructions that enqueue MFC commands (*Section 3.2.3*).

Software on the PPE and other SPEs and devices can gain access to most of an SPE's channel-interface functions by accessing associated problem-state (user) MMIO registers in the main-storage space.

For details about the channel facilities, see *Section 17 SPE Channel and Related MMIO Interface* on page 447.

### 3.2.2 Mailboxes and Signalling

The channel interface in each SPE supports two mailboxes for sending messages from the SPE to the PPE. One mailbox is provided for sending messages from the PPE to the SPE. The channel interface also supports two signal-notification channels (*signals*, for short) for inbound messages to the SPE.

The PPE is often used as an application controller, managing and distributing work to the SPEs. A large part of this task might involve loading main storage with data to be processed, then notifying an SPE by means of a mailbox or signal message. The SPE can also use its outbound mailboxes to inform the PPE that is has finished with a task.

### 3.2.3 MFC Commands and Command Queues

Software on the SPE, the PPE, and other SPEs and devices use MFC commands to initiate DMA transfers, query DMA status, perform MFC synchronization, perform interprocessor-communication via mailboxes and signal-notification, and so forth. The MFC commands that implement DMA transfers between the LS and main storage are called DMA commands.

The MFC maintains two separate command queues, shown in *Figure 3-5* on page 73—an SPU command queue for commands from the MFC's associated SPU, and a proxy command queue for commands from the PPE and other SPEs and devices. The two queues are independent.

The MFC supports out-of-order execution of DMA commands. DMA commands can be tagged with one of 32 tag-group IDs. By doing so, software can determine the status of the entire group of commands. Software can, for example, use the tags to wait on the completion of queued commands in a group. Commands within a tag group can be synchronized with a fence or barrier option. The SPE can also use atomic-update commands that implement uninterrupted updates (read followed by write) of a 128-byte cache-line-size storage location.

For details about MFC commands and command queues, see *Section 19.2 MFC Commands* on page 514.

### 3.2.4 Direct Memory Access Controller

The MFC's DMA controller (DMAC) implements DMA transfers of instructions and data between the SPU's LS and main storage. Programs running on the associated SPU, the PPE, or another SPE or device, can issue the DMA commands. (The PPE can also access an LS through the main-storage domain using load and store instructions.)

The MFC executes DMA commands autonomously, which allows the SPU to continue execution in parallel with the DMA transfers. Each DMAC can initiate up to 16 independent DMA transfers to or from its LS.

#### 3.2.4.1 *DMA Transfers*

To initiate a DMA transfer, software on an SPE uses a channel instruction to write the transfer parameters to the MFC command-queue channels. Software on the PPE and other SPEs and devices performs the comparable function by writing transfer parameters to the main-storage addresses of the MMIO registers that correspond to those channels.

Six parameters need to be written, but a single C/C++ intrinsic can accomplish all of them. Special techniques, such as multibuffering, can be used to overlap DMA transfers with SPE computation.

DMA transfers are coherent with respect to main storage. Virtual-memory address-translation information is provided to each MFC by the PPE operating system. Attributes of system storage (address translation and protection) are governed by the page and segment tables of the *PowerPC Architecture*. Although privileged software on the PPE can map LS addresses and certain MFC resources to the main-storage address space, enabling the PPE or other SPUs in the system to access these resources, this aliased memory is not coherent in the system.

For details about the DMA transfers, including code examples, see *Section 19.2.1 DMA Commands* on page 516, *Section 19.3 PPE-Initiated DMA Transfers* on page 523, and *Section 19.4 SPE-Initiated DMA Transfers* on page 529. For examples of using double-buffering techniques to overlap DMA transfers with computation on an SPU, see *Section 24.1.2* on page 692.

#### 3.2.4.2 *DMA List Transfers*

A DMA list is a sequence of *list elements* that, together with an initiating DMA list command, specifies a sequence of DMA transfers between a single area of LS and possibly discontinuous areas in main storage. Such lists are stored in an SPE's LS, and the sequence of transfers is initiated with a DMA list command. DMA list commands can only be issued by programs running on the associated SPE, but the PPE or other devices (including other SPEs) can create and store the lists in an SPE's LS. DMA lists can be used to implement scatter-gather functions between main storage and the LS.

Each list element in the list contains a transfer size, the low half of an effective address, and a stall-and-notify bit that can be used to suspend list execution after transferring a list element whose stall-and-notify bit is set. Each DMA transfer specified in a list can transfer up to 16 KB of data, and the list can have up to 2048 (2 K) list elements.

For details about the DMA list transfers, including examples, see *Section 19.2.2 DMA List Commands* on page 518 and *Section 19.4.4* on page 536.

### 3.2.5 Synergistic Memory Management Unit

The MFC's synergistic memory management (SMM) unit, shown in *Figure 3-5* on page 73, provides the address-translation and protection features defined by *PowerPC Architecture* for accesses to main storage. The SMM does this based on address-translation information provided to it by privileged PPE software.

*Figure 3-3* on page 67 illustrates the manner in which the PPE, other SPEs, and I/O devices access the SPU's associated LS when the LS has been aliased to main storage. The SMM supports address translations for DMA accesses by the associated SPU to main storage.

The SMM has an 8-entry segment lookaside buffer (SLB) and a 256-entry translation lookaside buffer (TLB) that supports 4 KB, 64 KB, 1 MB, and 16 MB page sizes. Neither the SMM nor the SPU has direct access to system control facilities, such as page-table entries. For details about address translation, see *Section 4 Virtual Storage Environment* on page 79.

The following sections introduce the SPU instruction set and its C/C++ intrinsics. These instructions and intrinsics, together with the MFC's own command set (described in *Section 19.2* on page 514), are used to control DMA transfers and other essential aspects of the MFC.

## 3.3 SPU Instruction Set

The SPU supports an instruction set, as specified by the *Synergistic Processor Unit Instruction Set Architecture* (ISA), and a set of C/C++ intrinsics, as specified by the *C/C++ Language Extensions for Cell Broadband Engine Architecture* document. Although most of the coding for Cell Broadband Engine Architecture (CBEA) processors[1] will be in a high-level language like C or C++, an understanding the SPU instruction set adds considerably to a programmer's ability to produce efficient code. This is particularly true because most of the intrinsics have mnemonics that relate directly to the underlying assembly-language mnemonics.

The SPU ISA operates primarily on SIMD vector operands, both fixed-point and floating-point, with support for some scalar operands. The PPE and the SPE both execute SIMD instructions, but the two processors execute different instruction sets, and programs for the PPE and SPEs are often compiled by different compilers. *Section 2.5.1* on page 59 introduces and illustrates the concept of SIMD operations.

### 3.3.1 Data Types

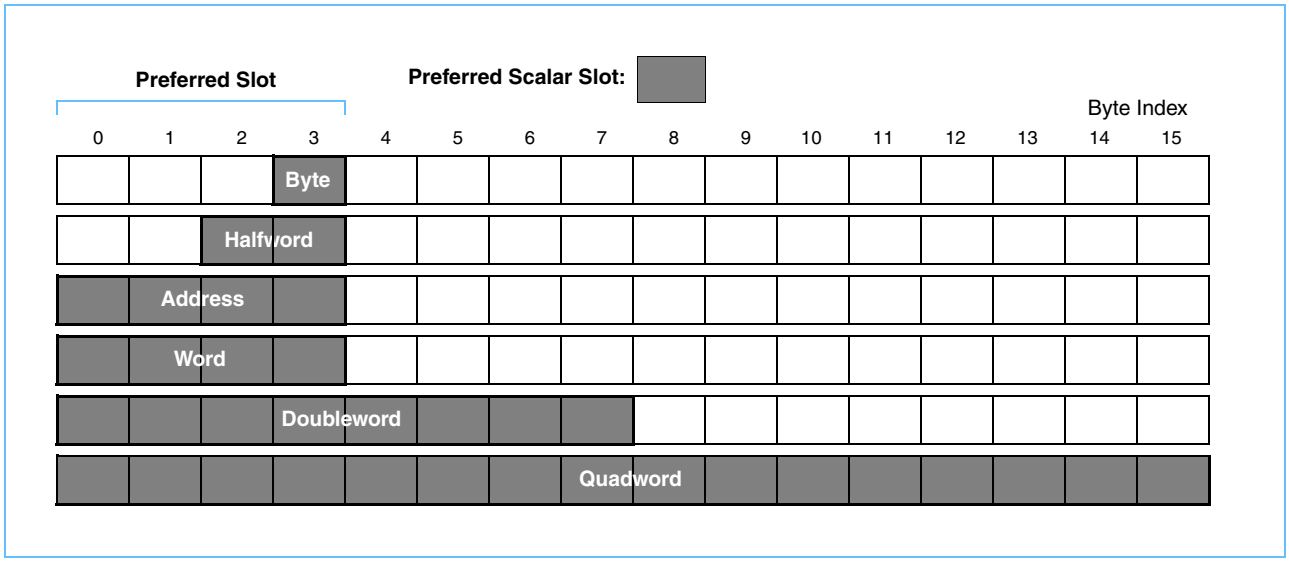
The SPU hardware supports the following data types:

- Byte—8 bits
- Halfword—16 bits
- Word—32 bits
- Doubleword—64 bits
- Quadword—128 bits

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

These data types are indicated by shading in *Figure 3-6* on page 77. The left-most word (bytes 0, 1, 2, and 3) of a register is called the *preferred slot* in the *Synergistic Processor Unit Instruction Set Architecture*, although they are called the *preferred scalar slot* in the *SPU Application Binary Interface Specification*. When instructions use or produce scalar operands or addresses, the values are in the preferred scalar slot. Because the SPU accesses its LS a quadword at a time, there is a set of store-assist instructions for insertion of bytes, halfwords, words, and double-words into a quadword for a subsequent store.

*Figure 3-6. Register Layout of Data Types and Preferred Scalar Slot*




### 3.3.2 Instructions

The SPU has two pipelines into which it can issue and complete up to two instructions per cycle, one in each of the pipelines. If two instructions with latency *n* and *m* are issued to the same pipeline, *n-m* cycles apart, the instructions are retired in order, with the later one stalling for one cycle. Whether an instruction goes to the even or odd pipeline depends on its instruction type.

*Table B-1 SPU Instructions* on page 772 lists the SPU instructions, with their latencies and their pipeline assignments. For details about the SPU instruction pipeline stages, see *Appendix B.1.3 Fetch and Issue Rules* on page 779.

## 3.4 SPU C/C++ Language Intrinsics

A set of C-language extensions, including data types and intrinsics, makes the underlying SPU instruction set and hardware features conveniently available to C programmers. The intrinsics can be used in place of assembly-language code when writing in the C or C++ languages.

The intrinsics represent in-line assembly-language instructions in the form of C-language function calls that are built into a compiler. They provide the programmer with explicit control of the SPE SIMD instructions without directly managing registers. A compiler that supports these intrinsics will emit efficient code for the SPE Architecture. The techniques used by compilers to generate efficient code should typically include:

- Register coloring
- Instruction scheduling (dual-issue optimization)
- Data loads and stores
- Loop blocking, fusion, unrolling
- Correct up-stream placement of branch hints
- Literal vector construction

For example, a C or C++ compiler should generate a floating-point add instruction (**fa rt,ra,rb**) for the SPU intrinsic **`t = spu_add(a,b)`**, assuming **`t`**, **`a`**, and **`b`** are vector float variables. The system header file, `spu_intrinsics.h`, available from IBM, defines the SPU language extensions. The PPE and SPU instruction sets have similar, but distinct, intrinsics.

### 3.4.1 Vector Data Types

The C/C++ language extensions' vector data types are described in *Appendix B.2.1* on page 784.

### 3.4.2 Vector Literals

The C/C++ language extensions' vector literals are described in *Appendix B.2.2* on page 786.

### 3.4.3 Intrinsics

The C/C++ language extension intrinsics are grouped into the following three classes:

- *Specific Intrinsics*—Intrinsics that have a one-to-one mapping with a single assembly-language instruction. Programmers rarely need the specific intrinsics for implementing inline assembly code because most instructions are accessible through generic intrinsics.
- *Generic Intrinsics and Built-Ins*—Intrinsics that map to one of several assembly-language instructions or instruction sequences, depending on the type of operands.
- *Composite Intrinsics*—Convenience functions that map to assembly-language instruction sequences. A composite intrinsic can be expressed as a sequence of generic intrinsics.

Intrinsics are not provided for all assembly-language instructions. Some assembly-language instructions (for example, branches, branch hints, and interrupt return) are naturally accessible through the C/C++ language semantics.

*Section B.2.3* on page 787 lists the most useful intrinsics. The intrinsics are defined fully in the *C/C++ Language Extensions for Cell Broadband Engine Architecture* document.

# 4. Virtual Storage Environment

## 4.1 Introduction

The Cell Broadband Engine Architecture (CBEA) processors[1] implement a virtual-storage environment that is compatible with that defined by the *PowerPC Operating Environment Architecture, Book III*. When virtual addressing is enabled, the combination of CBEA-processor hardware and operating-system software supports a virtual-address space that is larger than either the effective-address (main-storage) space referenced by programs or the real-address space supported by the CBEA processors.

Programs access memory by means of effective addresses. Programs on the PowerPC Processor Element (PPE) do this by executing instructions. Programs on a Synergistic Processor Element (SPE) do it by generating memory flow controller (MFC) direct memory access (DMA) commands, which contain pairs of addresses—a local storage (LS) address and an effective address (EA). The real-address (RA) space is the set of all addressable bytes in physical memory and on devices whose physical addresses have been mapped to the RA space—such as an SPE's on-chip LS or an I/O device's off-chip memory-mapped I/O (MMIO) register or queue. The EA space is called the main-storage space and is illustrated in *Figure 1-2* on page 47.

The operating system can enable virtual-address translation (also called *relocation*) independently for instructions and data. When it does so, the address-translation mechanism allows each program running on the PPE or addressing main storage by means of an SPE's DMA-transfer facility to access up to $2^{64}$ bytes of EA space. Each program's EA space is a subset of the larger $2^{65}$ bytes of virtual-address (VA) space. The memory-management hardware in the PPE and in the MFC of each SPE translates EAs to VAs and then to RAs, to access $2^{42}$ bytes of RA space.

EAs are translated to VAs by a *segmentation* mechanism. *Segments* are protected, nonoverlapping areas of virtual memory that contain 256 MB of contiguous addresses. VAs are translated to RAs by a *paging* mechanism. *Pages* are protected, nonoverlapping, relocatable areas of real memory that contain between 4 KB and 16 MB of contiguous addresses.

Because the physical memory available in a system is often not large enough to map all virtual pages used by all currently executing programs, the operating system can transfer pages from disk storage to physical memory as demanded by the programs—a technique called *demand-paging*. The PPE performs address translation and demand-paging dynamically, during the execution of instructions. The operating system can revise the mapping information to relocate programs on the fly, during each demand-paged transfer from disk into memory. The PPE operating system manages address-mapping by setting up segment-buffer entries and page-table entries used by the hardware to locate addresses in the virtual-addressing environment. The PPE operating system also provides each MFC with the segment-buffer and page-table entries needed by the MFC's associated SPE for its DMA transfers to and from main storage.

In PowerPC Architecture paging, page-table entries perform an *inverted* mapping—from the small physical-address space of installed memory into the large ($2^{65}$ bytes) virtual-address space. In contrast, conventional methods of paging map from a large virtual-address space to a small physical-address space. The inverted method of paging uses an *inverted page table,* which

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

can be much smaller than conventional page tables. Whereas the size of conventional page tables is commensurate with the size of the virtual-address space allocated by the operating system, the size of a *PowerPC Architecture* inverted page table is commensurate with the size of installed memory. Lookup time for the inverted page table is reduced by using a hash algorithm.

In addition to supporting demand-paging, the virtual-address mechanism allows the operating system to impose restrictions on access to each segment and page[2], such as privilege-level and access-type restrictions, and that impose attributes (such as cacheability and memory coherence) on the storage being accessed.

If the operating system does not enable virtual addressing, the system operates in *real addressing mode*. In this mode (also called *real mode*), the EAs generated by programs are used as RAs[3], there is no page table, and there is no demand-paging. The only addressable memory is the physically installed memory.

This section provides an overview of, and several details about, the Cell/B.E. and PowerXCell 8i implementation of its virtual-memory environment and the means and techniques for loading and maintaining the related hardware resources. For more detail, including that related to the PowerPC Architecture in general, see the *PowerPC Architecture* books, the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors*, and the *Cell Broadband Engine Architecture*.

## 4.2 PPE Memory Management

The PPE's memory management unit (MMU) controls access by the PPE to main storage, including the manner in which EAs are translated to VAs and RAs.

Address translation is managed in the MMU by the following hardware units:

- *Effective-to-Real-Address Translation (ERAT) Buffers*—Two 64-entry caches, one for instructions and one for data. Each ERAT entry contains recent EA-to-RA translations for a 4 KB block of main storage, even if this block of storage is translated using a large page (*Section 4.2.6.3* on page 88).
- *Segment Lookaside Buffer (SLB)*—Two unified (instruction and data), 64-entry caches, one per PPE thread, that provide EA-to-VA translations. The PPE supports up to $2^{37}$ segments of 256 MB each. Segments are protected areas of virtual memory.
- *Page Table*—A page table is a hardware-accessed data structure in main storage that is maintained by the operating system. Page-table entries (PTEs) provide VA-to-RA translations. Pages are protected areas of real memory. There is one page table per logical partition (see *Section 11 Logical Partitions and a Hypervisor* on page 331).
- *Translation Lookaside Buffer (TLB)*—The TLB is a unified, 1024-entry cache that stores recently accessed PTEs. A page table in memory is only accessed if a translation cannot be found in the TLB (called a *TLB miss)*.

2. See *Section 4.2.5.1* on page 85 for segment protections and *Section 4.2.6.2* on page 88 for page protections.
3. For the PPE, but not the SPEs, RAs can be offset by a base address in real mode, as described in *Section 4.2.8* on page 100.

### 4.2.1 Memory Management Unit

*Table 4-1* summarizes the features of the PPE's memory management unit (MMU).

*Table 4-1. PPE Memory Management Unit Features*

| Parameter | | PPE Value or Feature |
|---|---|---|
| Addresses | Effective Address (EA) Size | 64 bits. |
| | Virtual Address (VA) Size | 65 bits[1]. This value is designated "n" in the *PowerPC Architecture*. |
| | Real Address (RA) Size | 42 bits[2]. This value is designated "m" in the *PowerPC Architecture*. |
| ERAT | D-ERAT Entries | 64, shared by both threads (2 way × 32). |
| | I-ERAT Entries | 64, shared by both threads (2 way × 32). |
| SLB | SLB Entries | 64 per thread. |
| | Segment Size | 256 MB |
| | Number of Segments | $2^{37}$ |
| | Segment Protection | Segments selectable as execute or no-execute. |
| | SLB Maintenance | Special instructions for maintaining the SLB. |
| Page Table | Page Sizes | 4 KB, plus two of the following large-page sizes ($2^p$)[3]:<br>• 64 KB ($p = 16$)<br>• 1 MB ($p = 20$)<br>• 16 MB ($p = 24$) |
| | Number of Page Tables | • With hypervisor[4]: one page table per logical partition<br>• Without hypervisor: one page table |
| | Page Table Size (number of pages) | Determined by SDR1 register. |
| | Table Structure | Hashed page tables in memory. Inverted mapping, from small real-address space to large virtual-address space. |
| | Page Protection | Page Protection (PP) bits selectable as problem (user) or privileged (supervisor) state, and read-write or read-only. |
| | Page History | Referenced (R) and Changed (C) bits maintained. |
| TLB | TLB Entries | 1024, shared by both threads. |
| | TLB Maintenance | • Hardware-managed or software-managed TLB update<br>• Special instructions for maintaining TLB<br>• Pseudo least recently used (pseudo-LRU) replacement policy<br>• Replacement management table (RMT) for TLB replacement |

1. High-order bits above 65 bits in the 80-bit virtual address (VA[0:14]) are not implemented. The hardware always treats these bits as '0'. Software must not set these bits to any other value than '0' or the results are undefined in the PPE.
2. High-order bits above 42 bits in the 64-bit real address (RA[2:21]) are not implemented. The hardware always treats these bits as '0'. Software must not set these bits to any other value than '0' or the results are undefined in the PPE. The two most-significant bits (RA[0:1]) are architecturally reserved and must be '0'.
3. The value "p" is a power-of-2 variable in the *PowerPC Architecture* representing the size of a page.
4. See *Section 11 Logical Partitions and a Hypervisor* on page 331.

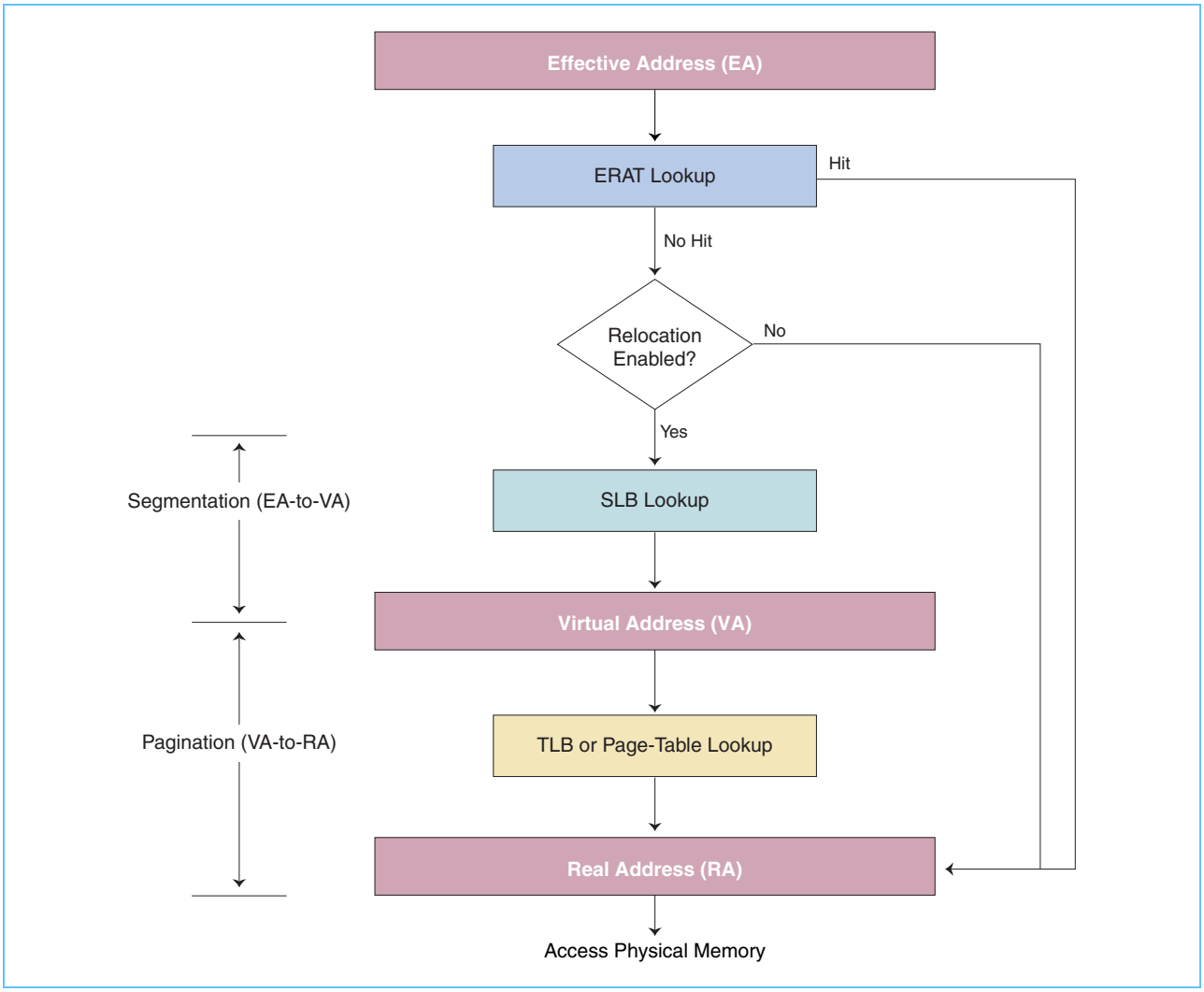
### 4.2.2 Address-Translation Sequence

*Figure 4-1* gives an overview of the address-translation sequence. Two basic steps are used to convert an EA to an RA:

1. *Segmentation*—Convert the EA to a VA, using the SLB.
2. *Pagination*—Convert the VA to an RA, using a page table.

The SLB (*Section 4.2.5.1* on page 85) is used for the first conversion. The TLB (*Section 4.2.7 on page 93*) caches recently used page-table entries and is used to enhance the performance of the second conversion. All EAs are first looked up in one of the two ERAT arrays. This ERAT lookup takes place regardless of whether the PPE is operating in virtual or real mode. If the address is not found in the ERAT, and address relocation is enabled, the address is sent to the segmentation and paging functions. If relocation is disabled, the real-mode translation facility is used.

*Figure 4-1. PPE Address-Translation Overview*



### 4.2.3 Enabling Address Translation

Virtual-address translation is enabled in the PPE by setting either or both of the Relocation bits in the Machine State Register (MSR) to '1':

- Instruction Relocation (MSR[IR])
- Data Relocation (MSR[DR])

If either MSR[IR] = '0' or MSR[DR] = '0', PPE accesses to instructions or data (respectively) in main storage operate in real mode, as described in *Section 4.2.8 Real Addressing Mode* on page 100.

The following sections describe MMU address translation. Additional details are given in the Privileged-Mode Environment section of the *Cell Broadband Engine Architecture*, *Cell Broadband Engine Registers*, and the *PowerPC Operating Environment Architecture, Book III*.

### 4.2.4 Effective-to-Real-Address Translation

Recently used EA-to-RA translations are stored in the effective-to-real-address translation (ERAT) arrays. There are two such arrays—the instruction ERAT (I-ERAT) array and the data ERAT (D-ERAT).

#### 4.2.4.1 *ERAT Functions*

The I-ERAT and D-ERAT are both 64-entry, 2-way set-associative arrays. They are not part of the normal *PowerPC Architecture* MMU. They provide faster address-translation information than the MMU. All main-storage accesses, regardless of the MSR[IR, DR] settings, pass through one of the two ERATs.

Each ERAT entry holds the EA-to-RA translation for an aligned 4 KB area of memory. When using a 4 KB page size, each ERAT entry holds the information for exactly one page. When using large pages, each ERAT entry contains a 4 KB section of the page, meaning that large pages can occupy several ERAT entries. All EA-to-RA mappings are kept in the ERAT including both real-mode and virtual-mode addresses (that is, addresses accessed with MSR[IR] equal to '0' or '1'). The ERATs identify each translation entry with some combination of the MSR[SF, IR, DR, PR, and HV] bits, depending on whether the entry is in the I-ERAT or D-ERAT. This allows the ERATs to distinguish between translations that are valid for the various modes of operation.

The ERATs are shared by both threads, but all entries are identified by the thread ID that created the entry. Each thread maintains its own entries in the ERATs and cannot use the entries created for the other thread.

On instruction fetches, the least-significant bits of the EA are used to simultaneously index into the L1 ICache, the L1 ICache directory, the I-ERAT, and the branch history table (BHT); see *Section 6.1 PPE Caches* on page 133 for descriptions of the caches and the BHT. Several conditions need to be satisfied to consider the fetch a hit:

- The EA of the instruction must match the EA contained in the I-ERAT entry being indexed, and the I-ERAT entry must also be valid.
- The Instruction Relocate (IR) and Hypervisor State (HV) bits from the MSR must match the corresponding bits in the I-ERAT, which were set at the time the I-ERAT entry was loaded.
- The thread performing the fetch must match the thread ID from the I-ERAT entry.

- The RA from the L1 ICache directory must match the RA provided by the I-ERAT.
- The page-protection bits must allow the access to occur.

If all of these conditions are met, and if no L1 ICache parity errors have occurred, the fetch is a hit. If an instruction misses in the I-ERAT, it is held in an ERAT miss queue (there is one miss queue for instructions and another for data) while a request is sent to the MMU for translation. This invokes an 11-cycle penalty to translate the address in the MMU. The instruction following the ERAT miss is flushed, refetched, and held at dispatch while the other thread is given all of the dispatch slots. If this translation hits in the TLB, the ERAT entry is reloaded and the fetch is attempted again.

The operation of the D-ERAT is very similar to the I-ERAT. The D-ERAT is accessed in parallel with the L1 DCache tags, and a comparison is performed at the end to determine if the access was a hit or not.

#### 4.2.4.2 *ERAT Maintenance*

One MMIO register field affects use of the ERATs: the ERAT Cache Inhibit (`dis_force_ci`) bit in Hardware Implementation Register 4 (`HID4[dis_force_ci]`). This bit should be set for normal use of the ERATs.

The replacement policy used by the ERATs is a simple 1-bit least recently used (LRU) policy. Each ERAT entry is set to the invalid state after power-on reset (POR). Normal ERAT operation is maintained by hardware. Any load or store that misses the appropriate ERAT causes all subsequent instructions to flush while that ERAT attempts to reload the translation from the MMU. However, in a hypervisor-controlled system, the ERAT entries are unique to each logical partition (LPAR) such that when a partition is switched by the hypervisor, the ERAT entries must be invalidated.

Each ERAT entry is obtained from a page-table search based on the contents of an SLB entry. To maintain consistency with the SLB, the following instructions cause all entries in an ERAT to be invalidated:

- **slbia** (invalidates entries belonging to the same thread only)
- **tlbie(l)**[4] to a large page only (invalidates all entries regardless of thread)

The **slbie** instruction causes a thread-based invalidation of the ERAT that is sensitive to the SLB class (see *Section 4.2.5.1* on page 85). That is, it invalidates any entries that have a thread-ID and Class bit match with the thread that issued and the segment class indicated by the **slbie** instruction. In addition, the execution of **tlbie(l)** to a small page (L = ‘0’), or the detection of a snooped-**tlbie** operation from another processor element, causes an index-based invalidate to occur in the ERAT. All ERAT entries that have EA bits `[47:51]` matching the `RB[47:51]` bits provided by the instruction are invalidated.

When operating in real mode (*Section 4.2.8 Real Addressing Mode* on page 100), the following sequence of instructions must be used to clear the ERAT without affecting the TLB:

4. The notation, **tlbie(l)**, indicates either the **tlbie** or the **tlbiel** instruction.

```
mtspr RMOR/HRMOR/LPCR/HID6
isync
tlbie L = 1, IS = 00, VPN(38:79-p) = nonmapped virtual address
sync
```

This sequence should be executed at the end of all real-mode phases, before enabling virtual-address translation, because the ERAT is always referenced, be it real or virtual address.

### 4.2.5 Segmentation

Segmentation is implemented with the segment lookaside buffers (SLBs). The SLBs translate EAs to VAs. The PPE contains two unified (instruction and data), 64-entry, fully associative SLBs, one per thread. Segments are protected areas of virtual memory. Each 256 MB segment can hold 64K pages of 4 KB each (or fewer of larger page sizes). The pages that make up a segment can overlap, but they must do so completely, not partially. The operating system can concatenate several contiguous segments (segments referenced in contiguous segment registers) to create very large address spaces. The first byte in a segment is also the first byte in the first page of that segment, and the last byte in the segment is the last byte in the last page.

The *small* or *flat* segmentation model is the simplest. It uses a single segment for all instructions, data, and stacks. In this model, virtual-address translation is, in effect, implemented entirely with paging. Segmentation gives the operating system flexibility in allocating protected address spaces to specialized areas of memory.

Multiple segments can be allocated for use by the instructions, private data, and stack of each process. They can also be allocated for shared libraries, shared memory-mapped I/O ports, and so forth. The operating system can also map file images from disk into shared segments, allowing applications to access file elements by moving a pointer or array index rather than by performing reads and writes through an I/O buffer. Two processes can thereby open the same file from different program-generated addresses and the accesses can map to the same area in the virtual address space. Access to segments can be protected by access-type and privilege attributes, so processes and data can be shared without interference from unauthorized programs.

#### 4.2.5.1 *Segment Lookaside Buffer*

*Figure 4-2* on page 86 shows the format for an SLB entry (SLBE). Each entry specifies the mapping between an effective segment ID (ESID) and a virtual segment ID (VSID). The ESID contains the most-significant bits of the EA; the VSID contains the most-significant bits of the VA. The PPE supports up to $2^{37}$ segments, as specified in the bits 52:88 of the VSID field. The PPE supports a 65-bit virtual address, so bits [37:51] of an SLB entry (`VSID[0:14]`) are always ‘0’.

When hardware searches the SLB, all entries in the SLB are tested for a VA match with the EA. If the SLB search fails, an instruction segment or data segment exception occurs, depending on whether the effective address is for an instruction fetch or for a data access.

*Figure 4-2. SLB Entry Format*


The other fields in the SLB entry include the Valid (V), Supervisor Key (Ks), Problem Key (Kp), No-execute (N), Large Page Indicator (L), Class (C), and Large Page Selector (LP) fields. See the *Cell Broadband Engine Architecture* for definitions of these fields. The PPE adds the Large Page (LP) field to the *PowerPC Architecture* SLB entry (see *Section 4.2.6.3 Large Pages* on page 88). The LP field is shown in *Figure 4-2* to be 3 bits wide, as specified in the *Cell Broadband Engine Architecture*. However, the PPE implements only the low-order bit of the LP field. For this reason, other LP references in this document describe the LP field as a single bit.

#### 4.2.5.2 *SLB Maintenance*

Software must ensure that the SLB contains at most one entry that translates a given EA (that is, a given ESID must be contained in no more than one SLB entry). Software maintains the SLB using the **slbmte**, **slbmfee**, **slbmfev**, **slbie**, and **slbia** instructions. All other optional *PowerPC Architecture* instructions (**mtsr**, **mtsrin**, **mfsr**, and **mfsrin**) provided for 32-bit operating systems and for changing the contents of the SLB are not implemented in the PPE.

The SLB-management instructions are implemented as move to or move from special-purpose register (SPR) instructions in microcode. Because these are microcoded, process performance is degraded when using these instructions. Software should attempt to minimize the number and frequency of these instructions.

For all SLB-management instructions, the first 15 bits [0:14] of the VSID field and the first six bits [52:57] of the Index field are not implemented. Writing them is ignored and reading them returns ‘0’.

The PPE adds the LP field to the RS register definition of the **slbmte** instruction and to the RT register definition of the **slbmfev** instruction. In both cases, this field is added as bits [57:59] of the corresponding register. Bits [57:58] are not implemented, so the LP field is treated as a single bit [59].

The **slbie** and **slbia** instructions clear the Valid bit of the specified SLB entries to ‘0’. The remaining contents of the SLB entry are unchanged and can be read for purposes of debugging. Therefore, when the hypervisor initiates a process or logical-partition context-switch, the hypervisor must set any used SLB entries to ‘0’ to prevent establishing a covert communication channel between processes or partitions.

For the effect these instructions have on the ERATs, see *Section 4.2.4 Effective-to-Real-Address Translation* on page 83.

Information derived from the SLB can be cached in the I-ERAT or the D-ERAT along with information from the TLB. As a result, many of the SLB-management instructions have effects on the ERATs as well as on the SLB itself. Because the SLB is managed by the operating system, multiple entries might be incorrectly set up to provide translations for the same EA. This will result in undefined behavior and can result in a checkstop.

#### 4.2.5.3 *SLB After POR*

All SLB entries and TLB entries are invalidated after power-on reset (POR).

### 4.2.6 Paging

Paging is set up by the operating system and can be actively managed by hardware or software, as described in *Section 4.2.7.1 Enabling Hardware or Software TLB Management* on page 93. The operating system creates page-table entries (PTEs), and hardware or software loads recently accessed PTEs into the translation-lookaside buffer (TLB). A high percentage of paging accesses—typically in excess of 95%—hit in the TLB.

On systems that support disk peripherals, only recently accessed pages are kept in physical memory; the remaining pages are stored on disk. If a program attempts to access a page containing instructions or data that are not currently in physical memory, a *page fault* occurs and the operating system must load the page from disk and update the page table with the physical and virtual addresses of that page.

#### 4.2.6.1 *Page Table*

The operating system must provide PTEs in the page table for all of real memory. There is one page table per logical partition (see *Section 11 Logical Partitions and a Hypervisor* on page 331), or one page table per system when logical partitions are not used. The operating system places PTEs in the page table by using a hashing algorithm. The hashed page table (HTAB) is a variable-sized data structure. Each PTE in the HTAB is a 128-bit data structure that maps one virtual page number (VPN) to its currently assigned real page number (RPN), thus specifying a mapping between VA and RA. The PTE also provides protection inputs to the memory-protection and cache mechanisms.

Each process can have its own segmented address space in the virtual space, but there is only one page table that is shared globally by all processes in a logical partition. During context switches, segments are changed but the TLB is not necessarily purged and the page table remains unchanged. To change a segment, the operating system simply loads a new value into the SLB, which is faster than reloading the TLB. A hashing function applied to a VA provides an index into the page table. The entire page table needed to map, for example, 1 GB of real memory would be a minimum of 2 MB in size (1 GB, times 8 bytes per PTE, divided by 4 KB per page). This architectural economy of paging overhead accounts, in part, for the superior performance of memory management in *PowerPC Architecture* processors, as compared with some other processor architectures.

#### 4.2.6.2 *Page Table Entries*

*Table 4-2* summarizes the implemented bits of a PTE. Software should set all unimplemented bits in the page table to '0'. *Table 4-2* views the PTE as a quadword when giving bit definitions. For more information, see the *PowerPC Operating Environment Architecture, Book III* and the *Cell Broadband Engine Architecture*.

*Table 4-2. PTE Fields*

| Bits | Field Name | Description |
|---|---|---|
| 0:14 | Reserved | Software must set these bits to '0' or results are undefined. |
| 15:56 | AVPN | Abbreviated virtual page number. |
| 57:60 | SW | Available for software use. |
| 61 | L | Large page indicator. |
| 62 | H | Hash function identifier. |
| 63 | V | Valid. |
| 64:85 | Reserved | Software must set these bits to '0' or results are undefined. |
| 86:**115** | RPN | Real page number (but see bit 115, in the next row). |
| **115** | LP | Large page selector (last bit of RPN if L = '1'). |
| 116:117 | Reserved | Software must set these bits to '0' or results are undefined. |
| 118 | AC | Address compare bit. |
| 119 | R | Reference bit. |
| 120 | C | Change bit. |
| 121 | W | Write-through. Hardware always treats this as '0'. See *Table 4-4* on page 92. |
| 122 | I | Caching-inhibited bit. |
| 123 | M | Memory-coherence. Hardware always treats this as '1' if I '0', or '0' if I = '1'. See *Table 4-4* on page 92. |
| 124 | G | Guarded bit. |
| 125 | N | No execute bit. |
| 126:127 | PP[0:1] | Page-protection bits 1:2. |

#### 4.2.6.3 *Large Pages*

The page size can be 4 KB or any two of the three Large Page sizes: 64 KB, 1 MB, or 16 MB. Large pages reduce page-table size and increase TLB efficiency. They are useful in situations that might suffer from demand-paged thrashing, because they minimize the overhead of filling real memory if single large pages can be operated on for extended periods of time.

The selection between the two large page sizes is controlled by the lower bit of the Large Page Selector (`LP`) field of the PTE when the Large Page Mode (`L`) bit of the PPE Translation Lookaside Buffer Virtual-Page Number Register (`PPE_TLB_VPN[L]`) is set to '1'.

The address of the real page (that is, the real page number, or RPN) is formed by concatenating the PTE's abbreviated RPN (ARPN) with the upper seven bits of the LP field and a '0'. Because the RPN must be on a natural page boundary, some bits of the ARPN and LP fields might be required to be '0'. If the L bit is not set in the PPE_TLB_VPN register, the RPN is formed by concatenating the ARPN with the whole LP field, and the page size is 4 KB.

#### 4.2.6.4 *Page-Table Organization and Lookups*

The *Cell Broadband Engine Architecture* adds a bit to the *PowerPC Architecture* definition of the Logical Partition Control Register (LPCR)—the TLB Load (TL) bit—that specifies whether TLB misses are handled by hardware or software. In the hardware-managed TLB mode, when a virtual address is presented to the TLB and no matching TLB entry is found, the hardware initiates a reload of the TLB from the page table. In the software-managed TLB mode, a page fault is generated immediately following any TLB miss, and the MMU will not attempt to reload the TLB. For details on software accessibility and maintenance of the TLB, see *Section 4.2.7.4* on page 95.

Each logical partition can set LPCR[TL] differently. Changes to its value (see *Section 4.2.7.1* on page 93 and *Section 11.2.3* on page 337) must follow all the same synchronization rules as are required when changing the logical partition ID (LPID). See *Section 20 Shared-Storage Synchronization* on page 561 and the *PowerPC Operating Environment Architecture, Book III* for the required synchronization instructions and sequences that must be used by the hypervisor. Failure to properly synchronize these changes can lead to inconsistent state and failure of the application or hypervisor.

In general, the MMU does not do speculative table lookups (tablewalks). The only situation in which the MMU performs a speculative table lookup is a caching-inhibited store that takes an alignment interrupt. Otherwise, the MMU waits to perform the table lookup until the instruction is past the point at which any exception can occur.

The PPE always performs page-table lookups and PTE writes in real mode (see *Section 4.2.8* on page 100). Because the PPE only supports a 42-bit RA, the following restrictions regarding the Storage Description Register 1 (SDR1) apply to hardware table lookups:

- Maximum Hash Table Size (HTABSIZE) allowed in SDR1 is 24 ('11000').
- Maximum Hash Table Origin (HTABORG) allowed in SDR1 is x'FFFFFF'.

The 16-byte PTEs are organized in memory as groups of eight entries, called PTE groups (PTEGs), each one a full 128-byte cache line. A hardware table lookup consists of searching a primary PTEG and then, if necessary, searching a secondary PTEG to find the correct PTE to be reloaded into the TLB. The L2 provides the PTEG data two 16-byte PTEs at a time, in a burst of 32 bytes. These 32 bytes are then run through a 2:1 multiplexer, producing 64 bytes—half of a primary or secondary PTEG. Searching of the primary or secondary PTEG is done as follows: first, the 64-byte even half (PTEs 0, 2, 4, and 6) is searched; if the PTE is not found here, the 64-byte odd half (PTEs 1, 3, 5, and 7) is searched. To accomplish two such searches, the data must be re-requested from the L2 for the odd half of the PTEG, and the multiplexer must be set to look at the appropriate half of the PTEG.

To summarize, hardware searches PTEGs in the following order:

1. Request the even primary PTEG entries.
2. Search PTE[0], PTE[2], PTE[4], and PTE[6].

3. Request the odd primary PTEG entries.
4. Search PTE[1], PTE[3], PTE[5], and PTE[7].
5. Repeat steps 1 through 4 with the secondary PTE.
6. If no match occurs, raise a data storage exception.

For best performance, the page table should be constructed to exploit this search order.

#### 4.2.6.5 *Page-Table Match Criteria*

The page-table match criteria are described in the Page Table Search section of *PowerPC Operating Environment Architecture, Book III*. The PPE also requires that PTE[LP] = SLBE[LP] whenever PTE[L] = '1'. In other words, the PTE page size must match the SLBE page size exactly for a page-table match to occur. If more than one match is found in the page table, then the matching entries must all have the same PTE[LP] value, or the results are undefined, as explained later in this section.

To conserve bits in the PTE, the least-significant bit, bit 115, of the PTE[RPN] field doubles as the PTE[LP] bit. The PTE[LP] bit exists in the page table only if PTE[L] = '1'. When this is true, PTE[LP] is no longer part of the PTE[RPN] field and instead serves as the PTE[LP] bit. See *Section 4.2.6.6* on page 90 for details.

If multiple PTEs match yet differ in any field other than SW, H, AC, R, or C, then the results are undefined. The PPE will either return the first matching entry encountered or it will logically OR the multiple matching entries together before storing them in the TLB. This can lead to unpredictable address-translation results, and the system might checkstop as a result.

The PPE does not implement PTE[0:15]. Software must zero these bits in all page table entries. The PPE will ignore these bits if set, and a match will occur as if PTE[0:15] = x'0000'.

#### 4.2.6.6 *Decoding for Large Page Sizes*

The PPE supports two large page sizes concurrently in the system, in addition to the small 4 KB page size. Each page can be a different size, but each segment must consist of pages that are all of the same size (see *Section 4.2.6.5*).

Two large page sizes can be chosen from a selection of three sizes (64 KB, 1 MB, and 16 MB). To accomplish this, the PPE defines the LP field in the PTE to override the low-order bit of the RPN in the PTE and the low-order bit of the VPN in the **tlbie(l)** instruction if the large page (L) bit is set to '1' for the PTE or **tlbie(l)** instruction. Whenever L = '1', RPN[51], and VPN[67] are implicitly '0'. This is why the PPE can override these bits to be a large page selector (LP) whenever L equals '1'.

There are two possible values of the LP field ('0' and '1'), which allow two different large page sizes to be concurrently active in the system. To determine the actual page size, four bits are kept in HID6[LB], bits [16:19]. *Table 4-3* on page 91 shows how the page size is determined given L, LP, and LB, where "x" means that any value is allowed.

*Table 4-3. Decoding for Large Page Sizes*

| L | LP | LB[0:1] | LB[2:3] | Page Size |
|---|---|---|---|---|
| 0 | x | x | x | 4 KB |
| 1 | 0 | 11 | x | Reserved |
| 1 | 0 | 10 | x | 64 KB |
| 1 | 0 | 01 | x | 1 MB |
| 1 | 0 | 00 | x | 16 MB |
| 1 | 1 | x | 11 | Reserved |
| 1 | 1 | x | 10 | 64 KB |
| 1 | 1 | x | 01 | 1 MB |
| 1 | 1 | x | 00 | 16 MB |

Each logical partition can have a different LB value. Thus, support for large page sizes can be selected on a per-partition basis. However, the hypervisor is only permitted to change the LB field in HID6 if it is switching to or creating a partition. The partition is not allowed to dynamically change this value. The following sequence (or a similar one) must be issued to clean up translation after changing LB:

1. Issue **slbia** for thread 0.
2. Issue **slbia** for thread 1.
3. For (i = ‘0’; i < 256; i++)

   **tlbiel** is = 11 rb[44:51] = i; // invalidate the entire TLB by congruence class.
4. **sync** L = ‘0’.
5. **mtspr** hid6.
6. **sync** L = ‘0’.
7. **isync**.

***Note:*** *The preceding procedure can impact performance negatively and should only be used when switching between partitions that use a different set of large page sizes.*

#### 4.2.6.7 *WIMG-Bit Storage Control Attributes*

The PPE supports a *weakly consistent* shared-storage model, as described in *Section 20 Shared-Storage Synchronization* on page 561. Storage consistency between processing elements and I/O devices is controlled by privileged software using the following four storage-control attributes and their associated bits (called the *WIMG* bits):

- Write-through (W bit)
- Caching-inhibited (I bit)
- Memory coherency (M bit)
- Guarded (G bit)

All PPE instruction accesses and all PPE and SPE data accesses to main storage are performed under the control of these storage attributes. The attributes are maintained in each PTE. The W and I bits control how the processor element performing an access uses its own cache (in the case of the CBEA processors, this applies to the PPE L1 and L2 caches and the SPE atomic caches). The M bit ensures that coherency is maintained for all copies of the addressed memory location. When an access requires coherency, the processor element performing the access informs the coherency mechanisms throughout the system of this requirement. The G bit prevents out-of-order loading and prefetching from the addressed location.

Like many PowerPC processors, the PPE supports only three WIMG-bit combinations: '0010', '0100', and '0101'. All other settings are implemented as shown in *Table 4-4*. In this table, a dash (–) represents either a '0' or a '1'. If software writes an unimplemented WIMG value to a page-table or TLB entry, the W and M bits might be implicitly overwritten in that entry, to the value shown in *Table 4-4*, if an update to a Reference or Change bit occurs (see *Section 4.2.6.8*).

*Table 4-4. Summary of Implemented WIMG Settings*

| WIMG Setting | PPE WIMG Implementation |
|---|---|
| – 0 – – | 0 0 1 0 |
| – 1 – 0 | 0 1 0 0 |
| – 1 – 1 | 0 1 0 1 |

As *Table 4-4* shows, if a storage access is cacheable, then it is also write-back, memory-coherent, and nonguarded. See *Section 4.2.8.2* on page 101 for details about how the storage-control attributes are set in real mode. For further details about storage control attributes in real mode, including those for hypervisor accesses, see the Real Addressing Mode section of the *PowerPC Operating Environment Architecture, Book III*.

The storage attributes are relevant only when an EA is translated. The attributes are not saved along with data in the cache (for cacheable accesses), nor are they associated with subsequent accesses made by other processor elements. Memory-mapped I/O operations address main storage and are therefore subject to the same WIMG attributes as all other main-storage accesses.

#### 4.2.6.8 *Updates to Reference and Change Bits*

If a hardware table lookup is performed and a matching PTE entry is found with the Reference (R) bit set to '0', the MMU performs a PTE store update, with the R bit set to '1'. Therefore, because all table lookups are nonspeculative (with the exception of a caching-inhibited misaligned store[5]), the R bit is always updated for any page that has a load or store request, regardless of any storage interrupts that might occur.

If a hardware table lookup is the result of a store operation, a matching PTE is tested to see if the Change (C) bit is set to '0'. If so, the MMU performs a PTE store update, with the C bit set to '1'. The only case in which this is not true is a storage exception. The MMU does not update the C bit if the storage access (load or store) results in a storage exception, as described in the Storage Protection section of *PowerPC Operating Environment Architecture, Book III*. The PPE performs all store operations, including PTE updates, in-order.

5. Misaligned, caching-inhibited stores are speculative because the instruction is not yet committed; an alignment interrupt has been taken, and the table lookup has not yet been performed at the time of the store.

A store operation to a page whose corresponding PTE's C bit is '0' in the TLB entry causes a page fault when in software-managed TLB mode (LPCR[TL] = '1'). To avoid this, software can set C = '1' on the initial load of a new TLB entry, so that an additional page fault is not incurred during the subsequent store operation.

### 4.2.7 Translation Lookaside Buffer

The PPE's translation lookaside buffer (TLB) is a unified (combined both instruction and data), 1024-entry, 4-way set associative cache that stores recently accessed PTEs. It is shared by both PPE threads. A page table in memory is only accessed if a translation cannot be found in the TLB (called a TLB miss).

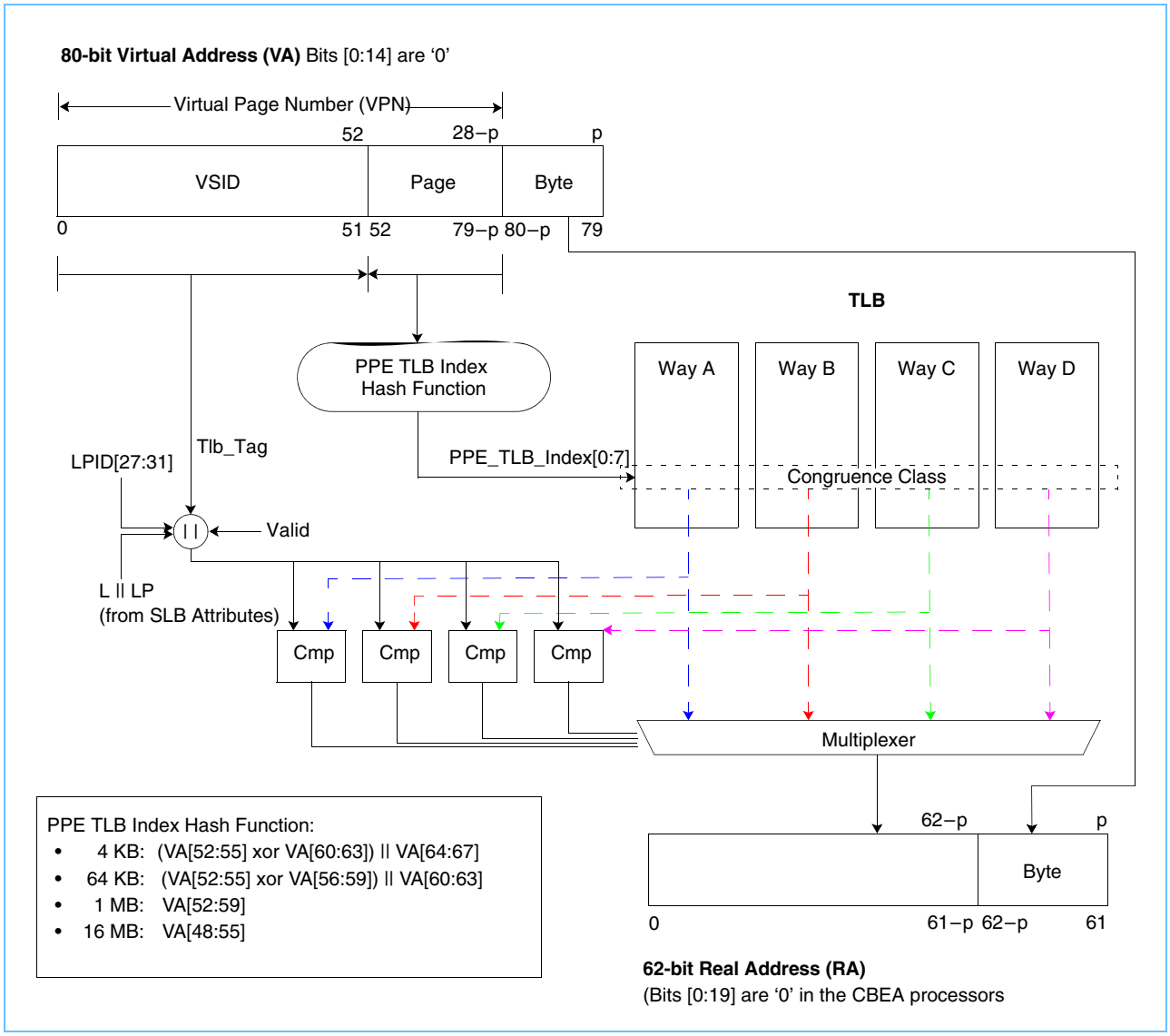
#### 4.2.7.1 *Enabling Hardware or Software TLB Management*

Translation lookaside buffer (TLB) entries and reloads can be managed either by hardware or by hypervisor software. The hardware-managed TLB mode is enabled when the TLB Load (TL) bit of the Logical Partition Control Register is cleared to '0' (LPCR[TL] = '0') (this bit is not defined as such in the *PowerPC Architecture*). The software-managed TLB mode is enabled when LPCR[TL] = '1'. In the software mode, the hypervisor uses MMIO-register accesses to manage the TLB as described in *Section 4.2.7.4 Software Management of the TLB on page 95*.

#### 4.2.7.2 *TLB Lookup*

The result of an SLB lookup is a VA with segment attributes. This information is passed to the TLB to be converted to an RA with page attributes. *Figure 4-3* on page 94 shows the VA-to-RA conversion performed by a TLB entry in the hardware-managed TLB mode. The VA is split to form the Page field, which is input to an index hash function, and the VSID field, which is used as a compare tag (Tlb_Tag). The result of the index hash function produces the index (PPE_TLB_Index) that chooses among the 256 congruence classes (sets, or rows) and the four ways (columns) of the TLB.

*Figure 4-3. Mapping a Virtual Address to a Real Address*




The compare tag (Tlb_Tag) also contains the LPID, Valid (V, set to '1'), Large Page (L), and Large Page Selector (LP) fields of the PTE. Both L and LP come from the SLB segment attributes. The compare tag is compared against the data stored in each way of the indexed congruence class. The width of the compare tag changes for each page size, as shown in *Table 4-5* on page 95.

*Table 4-5. TLB Compare-Tag Width*

| Page Size | Compare Tag (Tlb_Tag) |
|---|---|
| 4 KB | VA[0:59] \|\| Valid \|\| L \|\| LP \|\| LPID[27:31] |
| 64 KB | VA[0:55] \|\| Valid \|\| L \|\| LP \|\| LPID[27:31] |
| 1 MB | VA[0:51] \|\| Valid \|\| L \|\| LP \|\| LPID[27:31] |
| 16 MB | VA[0:47] \|\| Valid \|\| L \|\| LP \|\| LPID[27:31] |

If a hit results from the compare, the entry for the TLB way (see *Figure 4-3* on page 94) is read out through the multiplexer. This data contains a previously cached PTE corresponding to the VA presented to the TLB, which holds the RPN and page attributes.

If this is an instruction-address translation, the page attributes are checked to ensure that the page is not marked Guarded (G = '1') or No-execute (N = '1'), plus all other architected checks to ensure that access is allowed. If any condition is violated, an instruction storage interrupt is generated. Otherwise, if this is a data transaction or if the conditions are not violated, the RPN and page attributes are returned to the PPE's ERAT.

#### 4.2.7.3 *TLB Replacement Policy*

If a TLB miss results from the preceding compare operation, and hardware replacement is enabled (*Section 4.2.7.1* on page 93), a pseudo-LRU algorithm is used to replace the TLB entry. The replacement policy is a 3-bit pseudo-LRU policy in binary-tree format. It has 256 entries, each representing a single congruence class in the TLB. The pseudo-LRU algorithm used for PPE TLB replacement is identical to that used for SPE TLB replacement, as described in *Section 4.3.5.2 Hardware TLB Replacement* on page 110.

The LRU uses the following basic rules:

- The LRU is updated every time an entry in the TLB is hit by a translation. The entry that was referenced becomes the most-recently used, and the LRU bits are updated accordingly. However, the LRU is not updated when software reads the TLB for debug purposes using **mfspr** `PPE_TLB_VPN` or **mfspr** `PPE_TLB_RPN`.
- On a TLB miss, the update algorithm determines which entry to replace, using the following criteria:
  - *Invalid Entry*—Any invalid entry is replaced. If more than one entry is invalid, the left-most (lowest-numbered) way is chosen as invalid.
  - *Valid Entry*—The entry that the replacement management table (RMT) determines to be both least-recently used and eligible for replacement is replaced.

For a description of the replacement management table facility, see *Section 4.2.7.7 TLB Replacement Management Table* on page 99.

#### 4.2.7.4 *Software Management of the TLB*

TLB reloading can operate in either a hardware-managed or software-managed TLB mode (see *Section 4.2.7.1* on page 93). In the hardware-managed mode, described in the preceding sections, software need not keep the TLB consistent with the hardware-accessed page table in

main storage. In the software-managed mode, TLB management instructions and MMIO registers directly specify VA-to-RA translations, without the use of a hardware-accessed page table in main storage.

Software management of the TLB is restricted to the hypervisor, because memory is a system resource. Any attempt to execute TLB-management instructions when `MSR[PR, HV]` are not equal to '01' causes an illegal instruction interrupt.

The PPE implements the following MMIO registers, defined in the *Cell Broadband Engine Architecture*, for software management of the TLB:

- PPE Translation Lookaside Buffer Index Hint Register (`PPE_TLB_Index_Hint`)
- PPE Translation Lookaside Buffer Index Register (`PPE_TLB_Index`)
- PPE Translation Lookaside Buffer Virtual-Page Number Register (`PPE_TLB_VPN`)
- PPE Translation Lookaside Buffer Real-Page Number Register (`PPE_TLB_RPN`)

The PPE does not implement the following registers:

- TLB Invalidate Entry Register (`TLB_Invalidate_Entry`)
- TLB Invalidate All Register (`TLB_Invalidate_All`)

Software management of the TLB allows the system-software designer to forego the requirement of a hardware-accessible page table and use any page-table format. The format and size of a software-managed page table are not restricted by hardware. Software interprets the page table to load the TLB in response to a data storage interrupt caused by a TLB miss. Software TLB management can also be used in combination with hardware management to preload translations into the TLB.

One difference between a hardware and software TLB reload is the point at which the data storage interrupt (DSI) or SPE interrupt is presented to the PPE. For a hardware TLB reload, the interrupt is generated only after the page table is searched by hardware and a translation is not found. For a software TLB reload, the interrupt is generated when a translation is not found in the TLB.

In the software-managed TLB mode, all TLB misses result in a page fault leading to a storage interrupt. Software can then write a new entry to the TLB using the code sequences and MMIO registers shown in *Table 4-6* on page 97. Software can also use this facility to preload the TLB with translation entries that are anticipated in future code execution. This is useful for avoiding TLB misses altogether.

Software should use the **tlbie** instruction to invalidate entries in the TLB, as described in *Table 4-6* and *Section 4.2.7.6 Software TLB Invalidation* on page 98. In *Table 4-6*, the command names in parentheses are recommended compiler mnemonics.

*Table 4-6. TLB Software Management*

| TLB Operation[1] | Code Sequence |
|---|---|
| Write an entry (tlbmte) | `mtspr PPE_TLB_Index, RB;`<br>`mtspr PPE_TLB_RPN, RB;`<br>`mtspr PPE_TLB_VPN, RB;` |
| Read from Index Hint (tlbmfih) | `mfspr RT, PPE_TLB_Index_Hint;` |
| Set the index for writing (tlbmti) | `mtspr PPE_TLB_Index, RB;` |
| Debug: Read from entry (tlbmfe) | `mtspr PPE_TLB_Index, RB;`<br>`mfspr RT1, PPE_TLB_VPN;`<br>`mfspr RT2, PPE_TLB_RPN;` |
| Invalidate an entry | `tlbie(1)`[2] |
| Invalidate all entries | See *Section 4.2.7.6 Software TLB Invalidation* on page 98. |

1. Recommended compiler mnemonics are shown in parentheses.
2. See *Section 4.2.7.6 Software TLB Invalidation* on page 98.

Writing the PPE_TLB_VPN register causes an atomic write of the TLB with the last-written contents of PPE_TLB_RPN and the Lower Virtual Page Number (LVPN) field of PPE_TLB_Index. The TLB entry that is written is the entry pointed to the last-written value of PPE_TLB_Index.

The TLB tags each entry with the PPE logical partition ID (LPID) value at the time the entry is written. The hypervisor should set the LPID register to the value it wants to have the TLB entry tagged with at the time the entry is written. See *Section 4.2.7.8 Logical-Partition ID Tagging of TLB Entries* on page 99 for details. Software must set bit 56 of the PPE_TLB_VPN register (the low-order bit of the AVPN field) to '0' when writing a 16 MB TLB entry.

Reading the PPE_TLB_VPN or PPE_TLB_RPN register returns the PTE data associated with the TLB entry pointed to by the last-written value of PPE_TLB_Index. Reading the PPE_TLB_Index_Hint register returns the location that hardware would have chosen to replace when the last TLB miss occurred as a result of a translation for the thread performing the read operation (PPE_TLB_Index_Hint is duplicated per thread). If no miss has occurred, the register returns its POR value (typically all zeros).

The ability to read the TLB is provided for hypervisor debugging purposes only. The LPID and LVPN fields of the TLB entry are not accessible to the hypervisor, because they are not defined in the PTE that the PPE_TLB_VPN and PPE_TLB_RPN registers reflect. This makes it impossible to save and restore the TLB directly.

Bit 56 of the PPE_TLB_VPN register (the low-order bit of the AVPN field) is undefined when the TLB entry is a 1 MB page.

#### 4.2.7.5 *Rules for Software TLB Management*

When reading or writing the TLB, software must follow these rules:

1. Both threads should not attempt to write the TLB at the same time. Doing so can lead to boundedly undefined results, because there is a nonatomic sequence of register updates required for TLB writes (even though the write to the TLB itself is atomic). For example, the following situation should be avoided: thread 0 writes the PPE_TLB_Index and PPE_TLB_RPN registers; then thread 1 writes over the PPE_TLB_Index register; then thread 0 writes the

PPE_TLB_VPN register. The result is that the wrong TLB entry is written because thread 1 caused the PPE_TLB_Index register to point to the wrong location in the TLB. An easy way to solve this problem is for the hypervisor to reserve a lock when writing the TLB.

2. Both threads should not attempt to read the TLB at the same time, because reading the TLB using the **mtspr** and **mfspr** instructions requires first writing the PPE_TLB_Index register, so rule 1 applies. In other words, the same hypervisor lock that is used to write the TLB should also be acquired to read the TLB using the **mtspr** and **mfspr** instructions. This prevents one thread from corrupting the PPE_TLB_Index pointer that the other thread is using to read the TLB. Failure to acquire a lock can lead to boundedly undefined results.

   Also, because reading the TLB using the **mtspr** and **mfspr** instructions is a nonatomic sequence, hardware updates of the TLB by the other thread should be prevented. For example, the following situation should be avoided: thread 0 reads the PPE_TLB_RPN register; then thread 1 performs a hardware table-lookup update of the same TLB entry; then thread 0 reads the PPE_TLB_VPN register. The PPE_TLB_RPN and PPE_TLB_VPN results will not correspond in this case. Because reading of the TLB using the **mtspr** and **mfspr** instructions is intended for debug use only, the hypervisor should stop the other thread from accessing the TLB while performing this operation.

3. Locking is not required to translate in the TLB. Because writes to the TLB itself are atomic, the thread translating in the TLB sees either the previous or the new TLB entry, but never an intermediate result. This is analogous to hardware updates of the TLB.

4. The TLB is a noncoherent cache for the page table. If the TLB is completely software-managed, then it is software's responsibility to ensure that all TLB contents are maintained properly and that coherency with the page table is never violated.

   Invalidating a TLB entry requires use of the **tlbie** or **tlbiel** instruction (see *Section 4.2.7.6*). This is required so that the ERATs maintain coherency with the TLB. This means, for example, that simply writing the Valid bit to ‘0’ for a TLB entry is not sufficient for removing a translation. Only the **tlbie(l)** instruction properly snoop-invalidates the ERATs.

5. Software must maintain Reference (R) and Change (C) bits in the page table. See *Section 4.2.6.8* on page 92 for R and C bit maintenance by the hardware table lookup. If software writes a TLB entry with C = ‘0’, a subsequent store operation in the software-managed TLB mode takes a page fault. The R bit is ignored by the TLB, so writing R = ‘0’ is ignored and treated like R = ‘1’.

Failure of the hypervisor to properly synchronize TLB read and write operations can lead to undefined results throughout the entire system, requiring a system reset.

#### 4.2.7.6 *Software TLB Invalidation*

Software must use the **tlbie** or **tlbiel** instructions to invalidate entries in the TLB. The **tlbie** instruction is broadcast to all processor elements in the system. The architecture requires that only one processor element per logical partition can issue a **tlbie** at a given time.

The PPE supports bits [22:(63 – *p*)], of the RB source register for **tlbie** and **tlbiel**, which results in a selective invalidation in the TLB based on VPN[38:(79 – *p*)] and the page size, *p*[6]. In other words, any entry in the TLB that matches the VPN[38:(79 – *p*)] and the page size is invalidated.

6. The value “p” is a power-of-2 variable in the *PowerPC Architecture* representing the size of a page.

The PPE adds new fields to the RB register of the **tlbiel** instruction that are not currently defined in the *PowerPC Architecture*. These include the Large Page Selector (LP) and Invalidation Selector (IS) fields. The IS field is provided in RB[52:53] of the **tlbiel** instruction. *Table 4-7* gives details of the IS field. Bit 1 of the IS field is ignored.

*Table 4-7. Supported Invalidation Selector (IS) Values in the **tlbiel** Instruction*

| IS Value | Behavior |
|---|---|
| '00' | The TLB is as selective as possible when invalidating TLB entries. The invalidation match criteria is VPN[38:79-p], L, LP, and LPID. |
| '01' | Reserved. Implemented the same as IS = '00'. |
| '10' | Reserved. Implemented the same as IS = '11'. |
| '11' | The TLB does a congruence-class (index-based) invalidate. All entries in the TLB matching the index of the virtual page number (VPN) supplied will be invalidated. |

The **tlbia** instruction is not supported by the PPE. To remove the entire contents of the TLB, the hypervisor should perform 256 **tlbiel** operations with IS set to '11' and RB[44:51] set to increment through each TLB congruence class.

#### 4.2.7.7 *TLB Replacement Management Table*

The PPE provides a method of controlling the TLB replacement policy based on a replacement-class ID (RclassID). Software can use the replacement management table (RMT) stored in the PPE Translation Lookaside Buffer RMT Register (`PPE_TLB_RMT`) to lock translation entries into the TLB by reserving a particular way of the TLB to a specific program EA range. For details, see *Section 6.3 Replacement Management Tables* on page 154.

#### 4.2.7.8 *Logical-Partition ID Tagging of TLB Entries*

The TLB supports entries for multiple logical partitions simultaneously. Each entry in the TLB is tagged with a 5-bit logical partition ID (`LPID`), which is specified in the Logical Partition Identity Register (`LPIDR`). This tag is assigned when the TLB entry is updated by either hardware or software.

The `LPID` tag on each TLB entry is used by hardware for the following purposes:

- To identify a TLB entry as belonging to a given partition. The TLB does not need to be saved, invalidated, and restored on a logical-partition switch, thus supporting high-performance partition-switching.
- To determine if a TLB entry is a match for a translation request. To be considered a match, the current value of the processor element's `LPID` register and the tag in the TLB must be equivalent. If not, a TLB miss occurs.
- To determine if a **tlbie** should invalidate a given TLB entry. Each processor appends the current `LPID` information to every **tlbie** sent out to the bus. When a **tlbie** is received by the TLB, the `LPID` information can be used to selectively invalidate only the TLB entries that correspond to the same partition as the **tlbie**.

It is possible for one partition to invalidate or discard the TLB entries of another partition by issuing a **tlbie** instruction with IS set to '11'. See *Section 4.2.7.6* on page 98 for details.

Software can read TLB entries, but software cannot read the value of a TLB entry's LPID tag. Theoretically, this allows one partition to read the TLB entries of another partition, but it does not allow the partition to know which partition owns the TLB entry. However, reading from the TLB is restricted to the hypervisor, so the hypervisor should not provide a service that allows operating systems or applications to read the TLB-entry values. Reading TLB entries is not atomic, is low performance, and can cause a security risk. It is intended for hypervisor software debugging only.

Processor elements ignore **tlbie** instructions that do not have an LPID value matching the processor element's current LPIDR value. Because of this, when moving a partition from one processor element to another, software must invalidate all TLB entries on the current processor element before changing the LPIDR. The LPID tagging only benefits partition-switching on a single processor element, not switching a partition between multiple processor elements. The Invalidation Selector (IS) field of the **tlbiel** instruction can be used to override this behavior. If IS = '11', then the TLB performs a congruence-class invalidation of the TLB without regard to the LPID value supplied.

The LPID is ignored by the TLB replacement logic when performing a hardware table lookup.

### 4.2.8 Real Addressing Mode

If the PPE operating system does not enable virtual addressing, as described in *Section 4.2.3 Enabling Address Translation* on page 83, the system operates in *real addressing mode* (also called *real mode*). In this mode, EAs of application programs running on the PPE are used as RAs (possibly offset by an Real Mode Offset Register [RMOR] or Hypervisor Real Mode Offset Register [HRMOR] base address) and there is no page table or paging. In contrast to this PPE behavior, real-mode addresses generated by an SPE's MFC are never offset by RMOR or HRMOR base-address registers and they do not use the MSR[HV] and LPCR[LPES] fields and the RMLR register, as described in *Section 4.3.6* on page 117.

All real-mode accesses are treated as if they belong to a 4 KB page. If a PPE system has a smaller physical-address range than a program's EA range, the extra high-order bits of the EA are ignored (treated as if they were zeros) in the generation of the RA.

#### 4.2.8.1 *Types of Real Addressing Modes*

The contents of the following registers affect the real addressing mode, including real-mode accesses for page-table lookups:

- Real Mode Offset Register (RMOR)
- Hypervisor Real Mode Offset Register (HRMOR)
- The Real Mode Limit Select (RMLS), Real-Mode Caching Inhibited (RMI), and Logical Partitioning Environment Selector (LPES) (bit 1) fields of the Logical Partition Control Register (LPCR)
- The Real-Mode Storage Control (RMSC) field of the Hardware Implementation Register 6 (HID6)

*Table 4-8* on page 101 summarizes how the RA is determined from the EA in real mode. An address-translation fault causes an instruction storage or data storage interrupt. For further details, see the Real Addressing Mode section of the *PowerPC Operating Environment Architecture, Book III*.

*Table 4-8. Summary of Real Addressing Modes*

| Mode Name | Mode Bits | Real-Address Calculation |
|---|---|---|
| Hypervisor Offset Mode | MSR[HV] = '1'<br>EA(0) = '0' | RA = (EA[22:43] \| HRMOR[22:43]) \|\| EA[44:63] |
| Hypervisor Real Mode | MSR[HV] = '1'<br>EA(0) = '1' | RA = EA[22:63] |
| Real Offset Mode | MSR[HV] = '0'<br>LPCR[LPES] bit 1 = '1' | RA = (EA[22:43] \| RMOR[22:43]) \|\| EA[44:63] |
| Mode Fault | MSR[HV] = '0'<br>LPCR[LPES] bit 1 = '0' | LPAR Error (Interrupt) |

For details on how the WIMG bits are set in real mode, see *Section 4.2.6.7 WIMG-Bit Storage Control Attributes* on page 91. Other attribute bits are set in real mode as shown in *Table 4-9*.

*Table 4-9. Summary of Real-Mode Attributes*

| Attribute Bit | Value | Attribute Description |
|---|---|---|
| L | 0 | Large Page |
| LP | 0 | Large Page Selector |
| Ks | 0 | Supervisor State Storage Key |
| Kp | 0 | Problem State Protection Key |
| CL | 0 | Class |
| AC | 0 | Address Compare |
| PP[0:1] | 10 | Page Protection |
| C | 1 | Change |
| N | 0 | No Execute |

#### 4.2.8.2 *Real-Mode Storage Control Facility*

The PPE supports a real-mode storage control (RMSC) facility. This facility allows for altering the storage-control attributes in real mode based on a boundary in the real-address space. The value in the 4-bit RMSC field stored in HID6[26:29] determines the boundary as shown in *Table 4-10*.

*Table 4-10. Real-Mode Storage Control Values (Sheet 1 of 2)*

| Value in HID6[RMSC] Register Field | Real-Address Boundary |
|---|---|
| 0000 | 0 |
| 0001 | 256 MB |
| 0010 | 512 MB |
| 0011 | 1 GB |
| 0100 | 2 GB |

*Table 4-10. Real-Mode Storage Control Values (Sheet 2 of 2)*

| Value in HID6[RMSC] Register Field | Real-Address Boundary |
|---|---|
| 0101 | 4 GB |
| 0110 | 8 GB |
| 0111 | 16 GB |
| 1000 | 32 GB |
| 1001 | 64 GB |
| 1010 | 128 GB |
| 1011 | 256 GB |
| 1100 | 512 GB |
| 1101 | 1 TB |
| 1110 | 2 TB |
| 1111 | 4 TB |

*Figure 4-4* demonstrates how the I (caching-inhibited) and G (guarded) real-mode storage-attribute bits are set according to the location of the Real Address of the storage access relative to the boundary. “RMI” is the Real-Mode Caching Inhibited bit in `LPCR[62]`.

*Figure 4-4. Real-Mode I and G Bit Settings*


#### 4.2.8.3 *Real-Mode Maintenance*

Real-mode translations are kept in the ERAT, so software must issue the following sequence when changing any of the `RMOR`, `HRMOR`, `LPCR`, or `HID6` registers:

- **mtspr** for `RMOR`, or `HRMOR`, or `LPCR[RMLS, RMI, LPES bit 1]`, or `HID6[RMSC]`
- **isync**
- **tlbiel** L = ‘1’, IS = ‘00’, RB = ‘0’
- **sync** L = ‘1’

This form of **tlbie** causes all ERAT entries to be invalidated without affecting the TLB.

Changing any of these registers changes the context of all real-mode translations, so care must be taken to ensure that any previous real-mode translations are completed from a storage perspective and are invalidated from the ERAT. The preceding sequence should also be followed

when multiple registers are to be changed at once—for example, when switching between logical partitions, the RMOR, RMLS, and LPIDR register need to by changed at once. When changing any of these register values, software must not cause an implicit branch, as described in the Implicit Branch section of *PowerPC Operating Environment Architecture, Book III*.

### 4.2.9 Effective Addresses in 32-Bit Mode

The PPE supports both the 64-bit and 32-bit memory management models defined by the *PowerPC Operating Environment Architecture, Book III*. The *PowerPC Architecture* also defines a 32-bit mode of operation for 64-bit implementations, which the PPE supports.

In this 32-bit mode (MSR[SF] = '0'), the 64-bit EA is first calculated as usual, and then the high-order 32 bits of the EA are treated as '0' for the purposes of addressing memory. This occurs for both instruction and data accesses, and it occurs independently from the setting of the MSR[IR]bit, which enables instruction address translation, and the MSR[DR] bit, which enables data address translation. The truncation of the EA is the only way in which memory accesses are affected by the 32-bit mode of operation.

## 4.3 SPE Memory Management

All information in main storage is addressable by EAs generated by programs running on the PPE, SPEs, and I/O devices. An SPE program accesses main storage by issuing a DMA command, with the appropriate EA and LS addresses. The EA part of the DMA-transfer address-pair references main storage. Each SPE's MFC has two command queues with associated control and status registers. One queue—the MFC synergistic processor unit (SPU) command queue—can only be used by its associated SPE. The other queue—the MFC proxy command queue—can be mapped to the EA address space so that the PPE and other SPEs and devices can initiate DMA operations involving the LS of the associated SPE.

When virtual addressing is enabled by privileged software on the PPE, the MFC of an SPE uses its synergistic memory management (SMM) unit to translate EAs from an SPE program into RAs for access to main storage. The SMM's functions are similar to those of the PPE's MMU described in *Section 4.2 PPE Memory Management* on page 80, except that the PPE provides an SMM with segment-buffer and page-table information needed for address translation. This section describes the differences between the SPE's SMM and the PPE's MMU.

### 4.3.1 Synergistic Memory Management Unit

Address translation is managed in the SMM by the following hardware units:

- *Segment Lookaside Buffer (SLB)*—A unified (instruction and data), 8-entry array that provides EA-to-VA translations. The SLB is mapped to main storage as MMIO registers. The SPE supports up to $2^{37}$ segments of 256 MB each.
- *Translation Lookaside Buffer (TLB)*—The TLB is a 256-entry, 4-way set-associative cache that stores recently accessed PTEs supplied by the PPE operating system.

*Table 4-11* on page 104 summarizes the features of the SMM unit.

*Table 4-11. SPE Synergistic Memory Management Unit Features*

| Parameter | | SPE Implemented Value or Feature |
|---|---|---|
| Addresses | Effective Address (EA) Size | Same as PPE. |
| | Virtual Address (VA) Size | |
| | Real Address (RA) Size | |
| SLB | SLB Entries | 8 |
| | Segment Size | 256 MB |
| | Number of Segments | $2^{37}$ |
| | Segment Protection | None. The Class (C) and No-Execute (N) bits are stored in the SLB but are not used within the SMM. |
| | SLB Maintenance | Maintained by accesses to MMIO registers. |
| Page Table | Number of Page Tables | The SMM uses page-table entries supplied by the PPE. |
| | Table Structure | |
| | Page Protection | |
| | Page History | |
| | Page Sizes | 4 KB, plus two of the following Large Page Sizes:<br>• 64 KB (p = 16)<br>• 1 MB (p = 20)<br>• 16 MB (p = 24) |
| | Number of Pages | Determined by MFC_SDR[HTABSIZE] register field. Page sizes specified in this field must agree with the PPE SDR1[HTABSIZE] register field for the same page table. |
| TLB | TLB Entries | 256 |
| | TLB Maintenance | • Hardware-managed or software-managed TLB update<br>• Maintained by accesses to MMIO registers<br>• Pseudo-LRU replacement policy<br>• Replacement management table (RMT) for TLB replacement |

### 4.3.2 Enabling Address Translation

Virtual-address translation is enabled in an SPE by setting the Relocate (R) bit in the MFC State Register 1 (`MFC_SR1[R]`) to '1'. When this is done, the SLB, TLB, and page table facilities in the MFC are used to translate EAs in MFC commands to VAs, and then to a RAs.

If `MFC_SR1[R]` = '0', the main-storage side of MFC accesses operates in real mode. In this mode, EAs are used as RAs, but unlike PowerPC real addresses, the MFC real addresses are never offset by an RMOR or HRMOR base address. In addition, a page table is not used in real mode, and there is no demand-paging, as described in *Section 4.2.8 Real Addressing Mode* on page 100.

The following sections describe details of SMM address translation. Additional details are given in the Privileged-Mode Environment section of the *Cell Broadband Engine Architecture*, *Cell Broadband Engine Registers*, and the *PowerPC Operating Environment Architecture, Book III*.

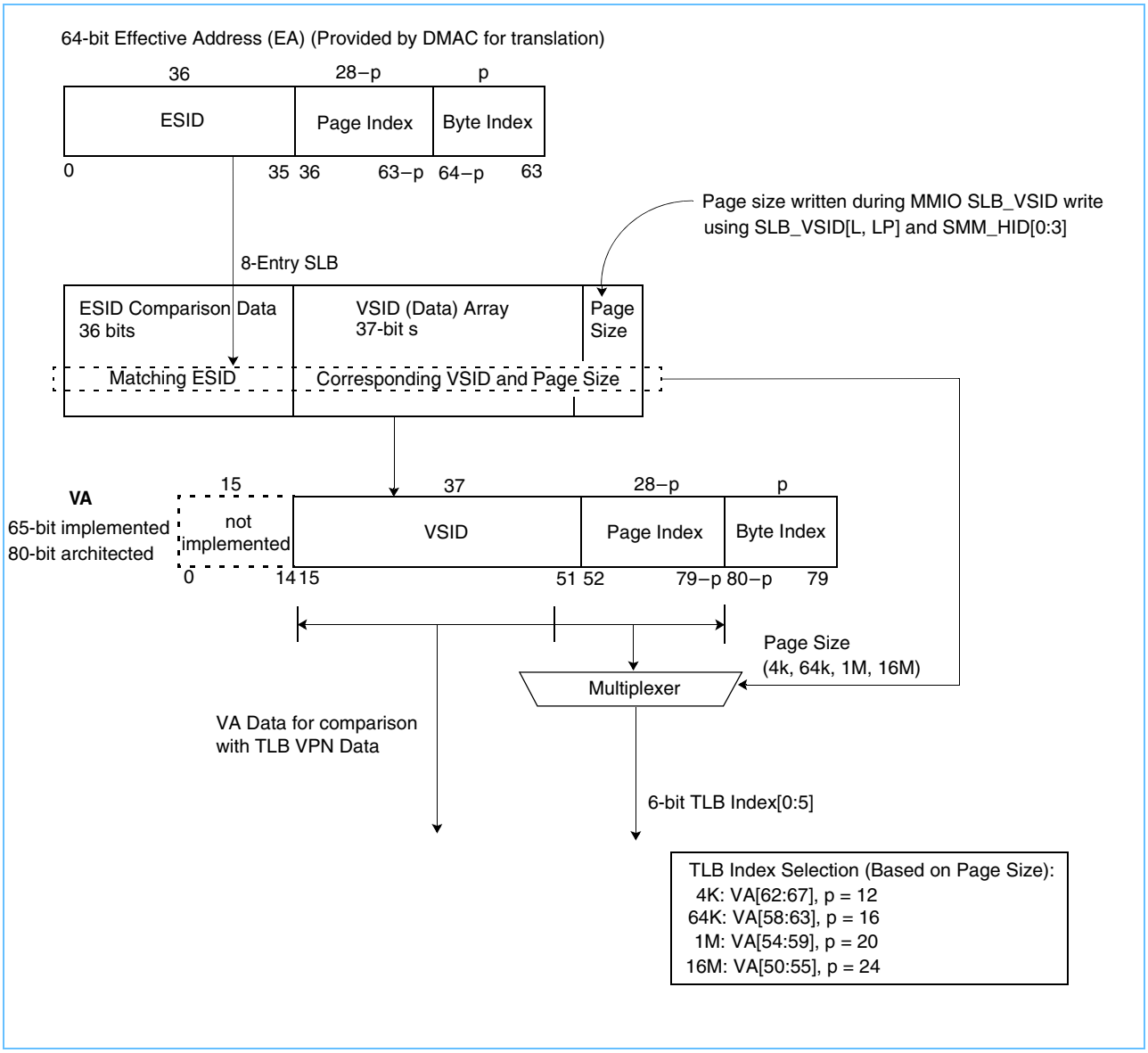
### 4.3.3 Segmentation

The SMM implements an 8-entry SLB. The SLB translates effective segment IDs (ESIDs) to virtual segment IDs (VSIDs)—which also translates EAs to VAs. The SMM has a set of MMIO registers for management of its SLB. These registers provide the same functionality for the SMM as do PowerPC instructions for the PPE's MMU, described in *Section 4.2.5 Segmentation* on page 85.

The SLB's Virtual Segment ID (`VSID`) field implements 37 bits of the architecture-defined 52-bit `VSID`. The Class (`C`) and No-execute (`N`) bits are stored in the SLB array but are not used within the SMM. All portions of the SLB, except the ESID and Valid bit, are protected by parity. Parity generation and checking can be disabled in the `SMM_HID` register.

Unlike the PPE's MMU, the SMM uses only a direct-mapped TLB—thus, all pages of the same size use the same virtual-address bits to obtain the TLB index for TLB lookups. Due to timing constraints, TLB-index selection occurs within the SLB array. Page-size encoding is saved in a separate 4-bit field of the SLB. The SMM uses six SLBs of the virtual page number (VPN) selected by the encoded page size to select a TLB entry. The page size is set and stored in the array on a write to the SLB Virtual Segment ID Register (`SLB_VSID`), and it is determined by the `L` and `LP` bits in the `SLB_VSID` data and the page-size encodings in the SMM Hardware Implementation Dependent Register (`SMM_HID`). See *Table 4-13* on page 113 for information about setting the page size. *Figure 4-5* on page 106 illustrates the EA-to-VA mapping and TLB index selection.

*Figure 4-5. SPE Effective- to Virtual-Address Mapping and TLB-Index Selection*




#### 4.3.3.1 *SLB Registers*

The SMM has a set of MMIO registers for SLB management. These registers provide the same functionality for the SMM that is provided by the *PowerPC Architecture* instructions (**slbie**, **slbia**, **slbmte**, **slbmfev**, and **slbmfee**) for managing the PPE's SLB. The SLB-management registers include:

- SLB Index Register (`SLB_Index`)
- SLB Effective Segment ID Register (`SLB_ESID`)
- SLB Virtual Segment ID Register (`SLB_VSID`)

- SLB Invalidate Entry Register (`SLB_Invalidate_Entry`)
- SLB Invalidate All Register (`SLB_Invalidate_All`)

*Section 4.3.3.2* provides details. See the *Cell Broadband Engine Registers* document for the implemented bits of these *Cell Broadband Engine Architecture* registers.

#### 4.3.3.2 *SLB-Entry Replacement*

The SLB requires a write sequence, similar to the one required for the TLB, to properly load the SLB with data. The SLB MMIO registers must be accessed by privileged PPE software with storage attributes of caching-inhibited and guarded.

First, the index must be written to specify the entry for loading. The `VSID` and `ESID` fields must be written independently, unlike the TLB writes, but the `SLB_VSID` write should follow the index. The `SLB_ESID` data is written last because it contains the Valid bit, and the entry should not be valid until all data is loaded. The `SLB_Index` register is write-only. The `SLB_ESID` and `SLB_VSID` registers are both read-write; however, MMIO reads retrieve data from the array instead of the architected register.

The sequence to replace an SLB entry is:

1. For each entry to be replaced:
   a. Set the index of the SLB entry to be replaced.
   b. Use the `SLB_Invalidate_Entry` register to invalidate the SLB entry.
2. Set the new contents for the `VSID` portion of the SLB entry.
3. Set the new contents for the `ESID` portion of the SLB entry along with the valid bit.

The contents of an SLB entry are accessed by using the `SLB_ESID` and `SLB_VSID` registers. The `SLB_Index` register points to the SLB entry to be accessed by the `SLB_ESID` and `SLB_VSID` registers.

#### 4.3.3.3 *SLB Mappings and Interrupts*

Software must maintain a one-to-one mapping between ESID and VSID entries, because multiple hit-detection is not supported by the SLB. The bits in the VSID are ANDed together if there are multiple matching ESIDs, and the results are unpredictable. If an SLB lookup fails due to an EA mismatch or cleared Valid bit, the SMM asserts a segment interrupt. All other DMA translation-request interrupts on a single translation are masked by a segment interrupt.

The SMM has a nonblocking, hit-under-miss pipeline architecture. The SMM can handle further DMA-translation hits under a pre-existing interrupt. Any subsequent interrupts are masked, and that translation is returned as a miss to the direct memory access controller (DMAC) until the original interrupt is cleared. After the original interrupt is cleared, the subsequent interrupt can be seen.

#### 4.3.3.4 *SLB Invalidation*

Software must use the SLB_Invalidate_Entry register to maintain the SLB and update entries. However, SLB invalidations occur by writing an EA in the SLB_Invalidate_Entry register that matches the ESID data in the entry to be replaced. Entries are invalidated by clearing the Valid bit in the ESID field. The remainder of the ESID and VSID data is untouched for register-readout debugging purposes.

SLB_Invalidate_Entry EA data is written into the SLB_ESID register because they share EA format data. The SLB_Invalidate_Entry register is write-only, and the SLB_Invalidate_Entry data cannot be retrieved because SLB_ESID reads return data from the SLB array and not the architected register. Reads from this register return '0' data.

No data is retained for writes to the SLB_Invalidate_All register. The write action prompts the SLB array to clear all eight Valid bits in the array and, therefore, invalidate all registers. This is a write-only register. Reads from this register return '0' data.

#### 4.3.3.5 *SLB after POR*

All SLB entries and TLB entries are invalidated after power-on reset (POR).

### 4.3.4 Paging

#### 4.3.4.1 *Page Table*

An SPE has no direct access to system-control facilities, such as the page table or page table entries (PTEs) for a logical partition. This restriction is enforced by PPE privileged software. The SPE's atomic (ATO) unit provides PPE-generated PTE data to the SMM for hardware table lookups, as needed.

#### 4.3.4.2 *Large Pages*

The page size can be 4 KB or any two of the three Large Page Sizes: 64 KB, 1 MB, or 16 MB. The selection between the two large page sizes is controlled by:

- Page Size Decode (Pg_Sz_Dcd) field of the SMM Hardware Implementation Dependent Register (SMM_HID). See *SMM_HID Register* on page 113.
- Large Page (L) bit of the MFC TLB Virtual-Page Number Register (MFC_TLB_VPN[L]).
- Large Page (LP) field of the MFC TLB Real Page Number Register (MFC_TLB_RPN[LP]).
- Large Page (L) bit and Large Page-Size (LP) bit of the SLB Effective Segment ID Register (SLB_ESID[L, LP]).

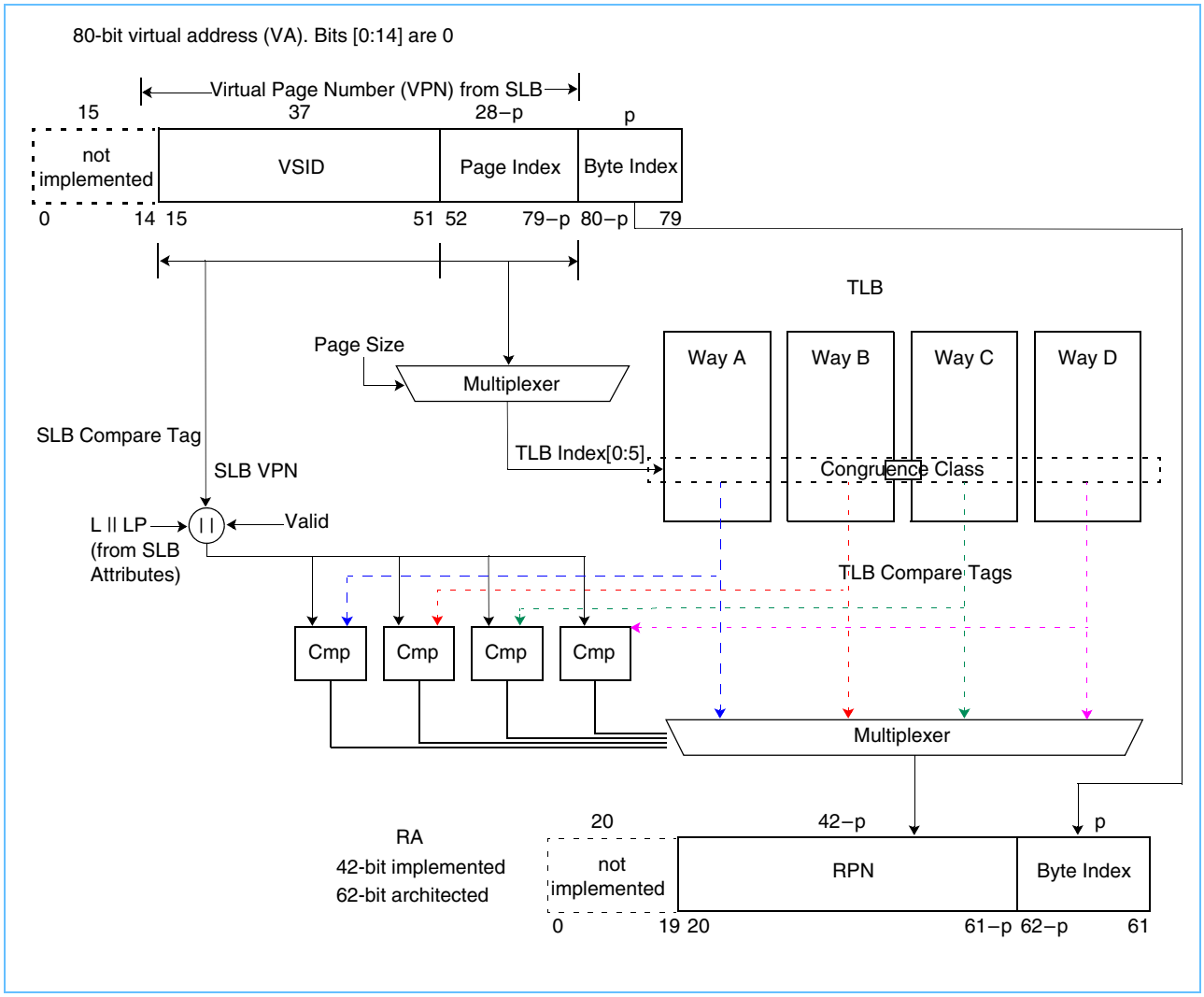
See *Table 4-13* on page 113 for large-page-size decoding.

### 4.3.5 Translation Lookaside Buffer

The TLB is a 256-entry, 4-way, set-associative cache that holds most-recently used PTEs that are provided by the PPE. There are 64 congruence classes (sets), each requiring six bits for the TLB index. When a translation is not found in the TLB, the SMM permits either a hardware or

software TLB reload, as specified in the MFC State Register 1 (MFC_SR1). The TLB array is parity protected. Parity generation and checking can be disabled in the SMM_HID register. *Figure 4-6* illustrates the VA-to-RA mapping.

*Figure 4-6. SMM Virtual- to Real-Address Mapping*




The 6-bit TLB index selects four PTE entries (ways) in a congruence class. The four VPN fields (the TLB Compare Tags) of the selected entries are compared to the VPN field obtained from SLB lookup (the SLB Compare Tag). *Table 4-12* on page 110 shows the SLB and TLB Compare Tag definitions for each page size. The TLB VPN data is a concatenation of MFC_TLB_VPN and MFC_TLB_Index register data during MMIO TLB writes. The TLB RPN data is a concatenation of the MFC_TLB_RPN[ARPN] and MFC_TLB_RPN[LP] register fields.

For large pages in which the L bit is set to '1', the RPN field concatenates the ARPN range with the upper seven bits of the LP field and a value of '0'. The lower bit of the LP field is used to select between the two large pages indicated by the SMM_HID register. The lower bit of the LP field is also

used as part of the TLB compare tag, as shown in *Table 4-12*. If the L bit is not set to '1', the RPN field is formed by concatenating the ARPN field with the whole LP field. See *MFC_TLB_Index, MFC_TLB_RPN, and MFC_TLB_VPN Registers* on page 114 for more information.

*Table 4-12. SLB and TLB Compare Tags*

| Page Size | SLB Compare Tag | TLB Compare Tag |
|---|---|---|
| 4 K | VA[15:61] \|\| Valid \|\| L \|\| LP | VPN[15:61] \|\| Valid \|\| L \|\| LP |
| 64 K | VA[15:57] \|\| Valid \|\| L \|\| LP | VPN[15:57] \|\| Valid \|\| L \|\| LP |
| 1 M | VA[15:53] \|\| Valid \|\| L \|\| LP | VPN[15:53] \|\| Valid \|\| L \|\| LP |
| 16 M | VA[15:49] \|\| Valid \|\| L \|\| LP | VPN[15:49] \|\| Valid \|\| L \|\| LP |

#### 4.3.5.1 *Enabling Hardware or Software TLB Replacement*

Hardware or software can be used to replace the TLB entry. The mode is selected by the TL bit of MFC State Register 1 (MFC_SR1).

#### 4.3.5.2 *Hardware TLB Replacement*

When there is a TLB-lookup miss, the SMM determines which TLB entry to replace by using a pseudo-LRU algorithm and a replacement management table (RMT), as described in *TLB Replacement Policy* on page 111. That data is saved in the MFC_TLB_Index_Hint register.

*Replacement Method*

If hardware replacement is selected, the SMM performs a hardware table lookup to find the replacement data and write the TLB at the location indicated by the MFC_TLB_Index_Hint register. A PTE request from the SMM is basically a store request. See *Section 6.2 SPE Caches* on page 151 for background information about an SPE's TLB and atomic cache.

If the cache line containing the PTE is in the MFC's atomic cache, which allocates a maximum of one line for PTEs, and the line is in the modified or exclusive cache state, the PTE is sent to the SMM. If the SMM requests an update for Reference (R) or Change (C) bits in the PTE, the atomic unit executes the store and updates those bits.

If the cache line containing the PTE is in the MFC's atomic cache, and in a state other than modified or exclusive, a bus request is sent to invalidate copies of the data in other processor elements, and the PTE is sent to the SMM. If the SMM requests an update for R or C bits in the PTE, the atomic unit executes the store and updates those bits.

If the cache line containing the PTE is not in the atomic unit, a bus request is sent to invalidate copies of the data in other caches and to read the data. If the atomic cache is full, the same castout scenario applies here as with the **putllc** or **putlluc** requests (see *Section 6.2.2 Atomic Unit and Cache* on page 151 for more information). That is, the PTE is sent to the SMM. If the SMM requests an update for R or C bits in the PTE, the atomic unit executes the store and updates those bits.

It is a software error if the PTE is used by the semaphore atomic unit operation as a locked line.

*TLB Replacement Policy*

The SMM uses a 3-bit, binary-tree pseudo-LRU algorithm to select a TLB entry for replacement after a TLB miss. There are 64 TLB congruence classes, each containing four (ways), thus 256 TLB entries as shown in *Figure 4-6* on page 109. A replacement management table (RMT) can be used to lock entries and prevent their replacement, as described in *Section 6.3.3 SPE TLB Replacement Management Table* on page 158.

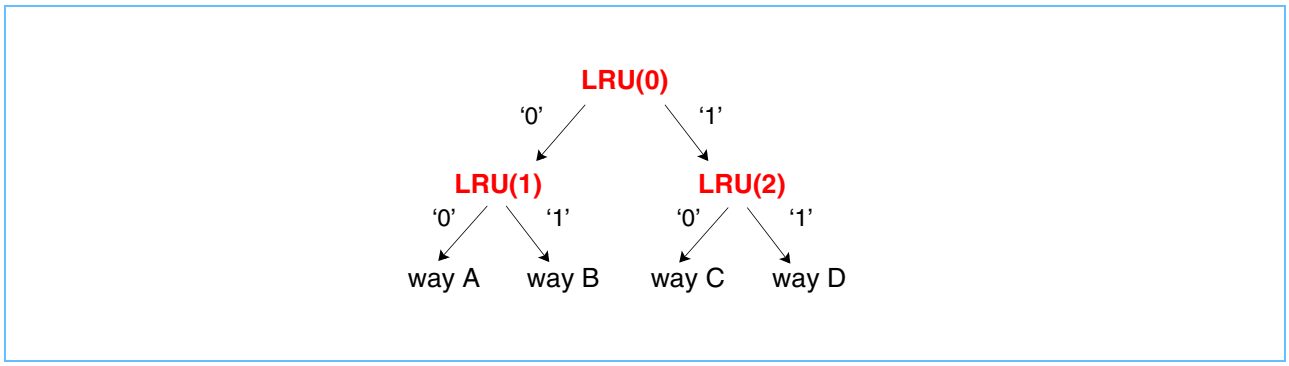
The following rules govern LRU use and TLB-entry replacement:

1. The LRU is updated every time an entry in the TLB is hit by a translation operation. The entry that was referenced becomes the most-recently used, and the LRU bits are updated accordingly. The LRU is not updated if that TLB entry is locked by the RMT.
2. On a TLB miss, the update algorithm determines which entry to replace, using the following priority:
   a. Replace the entry that hit in the TLB, but was a **put** command and the Change (C) bit was ‘0’.
   b. Replace the first entry (way) in the congruence class that was invalid (Valid bit was ‘0’).
   c. Replace the least-recently used entry, indicated by the pseudo-LRU algorithm, that is not locked by the RMT.

After a TLB entry is chosen for replacement using these rules, the TLB replacement index, which specifies TLB entries as a congruence class and way, is saved in the `MFC_TLB_Index_Hint` register (see *MFC_TLB_Index_Hint Register* on page 114). The last four bits of the TLB_Index_Hint register are a one-hot vector to indicate the entry to be replaced. The bits represent the four TLB ways within a congruence class like the four bits on an RMT entry, described in *Section 6.3.3 SPE TLB Replacement Management Table* on page 158. The `MFC_TLB_Index_Hint` vector can only indicate the congruence-class way that is selected for replacement, whereas the RMT entry indicates all ways in the congruence class that are available for replacement.

*Figure 4-7* shows the 3-bit pseudo-LRU binary tree. For each pseudo-LRU bit, a ‘0’ value moves to the left and a ‘1’ moves to the right. Only two bits are needed to determine one of the four congruence-class ways.

*Figure 4-7. Pseudo-LRU Binary Tree for SPE TLB*




Take, for example, an address translation with `RClassID` equal to x‘2’ that misses in the TLB at index x‘34’. The LRU is then read at index x‘34’ to retrieve the three pseudo-LRU bits. If LRU(0) equals ‘1’, we follow the tree to the right, and if LRU(2) equals ‘0’, we follow the tree to the left and way C is chosen for replacement.

The RMT must be checked to see if way C has been locked by software. The RClassID value selects RMT entry 2. If the RMT entry equals ‘1110’, then way C is unlocked and can be replaced. If the RMT entry equals ‘1100’, however, LRU(0) and LRU(2) are overridden by the RMT, and LRU(1) chooses either way A or way B for replacement. If LRU(1) equals ‘1’, way B is chosen for replacement, the MFC_TLB_Index_Hint register is written with index x‘34’, and way B is written as follows: MFC_TLB_Index_Hint[54:63] equals ‘11_0100_0100’.

*Interrupt Conditions*

The SMM contains logic to detect and report interrupt conditions associated with TLB data, including atomic access to caching-inhibited pages (real mode and virtual mode), page-protection (PP) violations, data address compares (DACs), TLB page faults (hardware-managed and software-managed modes), and parity errors during DMA and MMIO-register accesses. Real-mode translation caching-inhibited errors are detected by the DMAC and reported to the SMM for data storage interrupt (DSI) assertion. These errors are summarized in *Section 4.3.7 Exception Handling and Storage Protection* on page 118.

#### 4.3.5.3 *Software TLB Replacement*

As an alternative to the hardware-managed TLB mode (*Section 4.3.5.1 Enabling Hardware or Software TLB Replacement* on page 110), most SMM maintenance can be performed by PPE software accessing MMIO registers. If software replacement is enabled, privileged PPE software searches the PPE’s page table and provides a replacement using access to MMIO registers in the SMM. Software can use the MFC_TLB_Index_Hint data or choose another location within the TLB.

The SMM contains only privilege 1 (hypervisor) and privilege 2 (operating system) registers.

The TLB-management registers include:

- SMM Hardware Implementation Dependent Register (SMM_HID)
- MFC TLB Invalidate Entry Register (MFC_TLB_Invalidate_Entry)
- MFC TLB Replacement Management Table Data Register (MFC_TLB_RMT_Data)
- MFC Storage Description Register (MFC_SDR)
- MFC Address Compare Control Register (MFC_ACCR)
- MFC TLB Index Hint Register (MFC_TLB_Index_Hint)
- MFC Data-Storage Interrupt Status Register (MFC_DSISR)
- MFC TLB Index Register (MFC_TLB_Index)
- MFC TLB Real Page Number Register (MFC_TLB_RPN)
- MFC TLB Virtual-Page Number Register (MFC_TLB_VPN)
- MFC Data Address Register (MFC_DAR)

Descriptions of these registers are given in the following sections. See the *Cell Broadband Engine Registers* document for complete descriptions of the implemented portions of the MFC TLB MMIO registers defined in the *Cell Broadband Engine Architecture.* Writes to invalid register-offset addresses in the assigned SMM range are ignored. Reads to invalid addresses return ‘0’ data. The same applies to nonimplemented bits.

The synergistic bus interface (SBI) unit retrieves all register requests from the element interconnect bus (EIB) and formats the requests for the MFC. All valid SMM requests are formatted for 64-bit data. The MFC_DSISR register can be addressed as a 32-bit register, but the SBI converts it to a 64-bit format.

*SMM_HID Register*

The SMM Hardware Implementation Dependent Register (SMM_HID) contains configuration bits that control many essential functions of the SMM. This register should be written before the SMM performs any translations and before the arrays are loaded with data.

The Pg_Sz_Dcd bits are the most important field in this register. These bits, along with the LP bit in the SLB_VSID register, determine the supported large-page sizes for each array entry when the L bit in SLB_VSID is '1', enabling large-page mode. The large-page (LP) sizes are determined during a write to the SLB_VSID register. See *Table 4-13* for large-page size decoding. By default, the SMM_HID register is cleared to '0', so 16 MB pages are decoded. There are two sets of 2-bit HID values to support the two concurrent large-page sizes. This portion of the SMM_HID register should be written before any SMM activity occurs, and it can only change with a context switch.

*Table 4-13. SMM_HID Page-Size Decoding*

| L | SLB[LP] or TLBIE[LS] | SMM_HID[0:1] | SMM_HID[2:3] | Page Size |
|---|---|---|---|---|
| 0 | — | — | — | 4 KB |
| 1 | 0 | 00 | — | 16 MB |
| 1 | 0 | 01 | — | 1 MB |
| 1 | 0 | 10 | — | 64 KB |
| 1 | 0 | 11 | — | Reserved |
| 1 | 1 | — | 00 | 16 MB |
| 1 | 1 | — | 01 | 1 MB |
| 1 | 1 | — | 10 | 64 KB |
| 1 | 1 | — | 11 | Reserved |

*Table 4-13* is applicable for MFC_TLB_Invalidate_Entry purposes as well. The MFC_TLB_Invalidate_Entry[L, LS] bits are used with SMM_HID[0:3] to select an MFC_TLB_Index for invalidation. See the MFC TLB Invalidate Entry MMIO register in *Cell Broadband Engine Registers* for more information. SMM page sizes must agree with PPE page sizes for the same page table.

Parity bits in the SLB are generated during a write to the SLB_VSID register. Parity-generation can be disabled by setting SMM_HID[SPGEN_Dsbl] to '1'. The parity bits are checked during address translation and reads from the SLB_VSID register. Parity-checking can be disabled by setting SMM_HID[SPCHK_Dsbl] to '1'.

Similarly, parity bits are generated for the TLB entries during writes to the TLB and upon entry reloads by the table-lookup mechanism. TLB parity-generation can be disabled by setting SMM_HID[TPGEN_Dsbl] to '1'. Parity-checking occurs during address translation and reads from the MFC_TLB_VPN or MFC_TLB_RPN registers. The checking can be disabled by setting

SMM_HID[TPCHK_Dsbl] to ‘1’. Enabling of parity-generation and disabling parity-checking causes data error (DERR) interrupts; see *Table 4-15 Translation, MMIO, and Address Match-Generated Exceptions* on page 118 for more about DERR interrupts.

RMT functionality in the SMM can be disabled by setting SMM_HID[RMT_Dsbl] to ‘1’.

With SMM_HID[Stall_DMA_Issue] equal to '0' (the default setting), a pending TLBIE condition stalls the DMAC from issuing new requests to the atomic and SBI units. When SMM_HID[Stall_DMA_Issue] is set to '1', a pending TLBIE condition does not cause the DMAC to stall. With SMM_HID[ATO_TLBIE_Comp] equal to '0' (the default setting) and one of three livelock conditions (set by read and claim [RC], castout, or PTE fetch) detected, the atomic unit does not force a pending TLBIE to complete. When SMM_HID[ATO_TLBIE_Comp] is set to '1', a detected live-lock condition forces a TLBIE completion.

### *MFC_SDR Register*

The MFC Storage Description Register (MFC_SDR) contains the page-table origin and size. The function of this register is identical to that of the PPE_SDR1 register.

### *MFC_TLB_Index_Hint Register*

The function of the MFC_TLB_Index_Hint register is identical to that of the PPE_TLB_Index_Hint register, except that the SMM version has only one thread and implements a 6-bit index in bits[53:59].

### *MFC_TLB_Index, MFC_TLB_RPN, and MFC_TLB_VPN Registers*

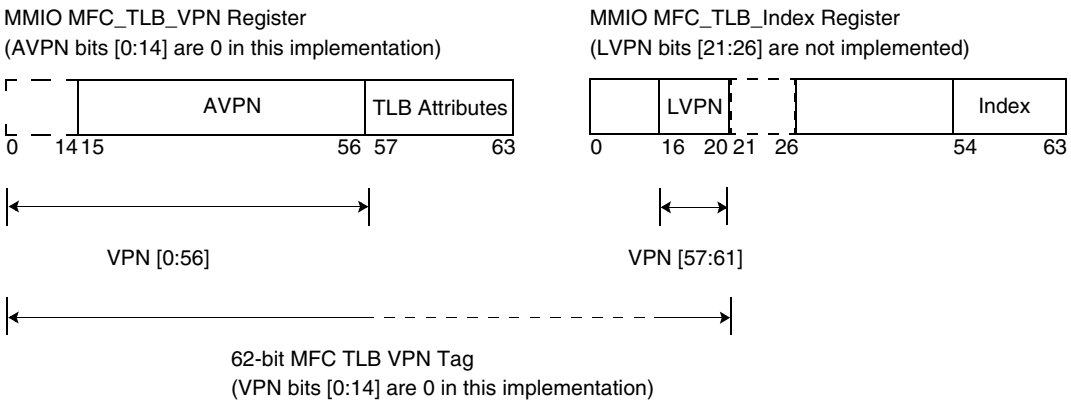
To write a new entry into the TLB, software must write the following registers, in the following order:

- MFC TLB Index (MFC_TLB_Index)
- MFC TLB Real-Page Number (MFC_TLB_RPN)
- MFC TLB Virtual-Page Number (MFC_TLB_VPN)

After this is done, hardware writes the entire TLB entry. Because an MMIO write to the MFC_TLB_VPN register sets off an automatic hardware write to the MFC_TLB_VPN and MFC_TLB_RPN register fields, whether the Index and RPN were previously written or not, the RPN data must be written after the index, as in the preceding list.

The complete MFC TLB VPN tag is created by concatenating the MFC_TLB_VPN[AVPN] and MFC_TLB_Index[LVPN] fields, as shown in *Figure 4-8* on page 115. This is the complete tag stored in the TLB during MMIO writes and used for comparison with the VA during TLB lookup.

*Figure 4-8. MFC TLB VPN Tag*




For TLB reads, the Index should be written first, as with the writes, but only bits [54:63]. MMIO VPN reads must occur in two separate MMIO reads. Reading the MFC_TLB_VPN register returns the register data just as it was written, and it also updates the LVPN field in the MFC_TLB_Index register. A separate MFC_TLB_Index register read must be done to retrieve the LVPN data. MFC_TLB_RPN register reads can occur independently and use the Index field of the MFC_TLB_Index register to select the entry to read.

The W bit of the MFC_TLB_RPN register is not implemented by the Cell/B.E. and PowerXCell 8i processors. The R and M bits are not retained in the TLB array, but they are returned as '1' during an MMIO read of the MFC_TLB_RPN register.

### *MFC_TLB_Invalidate_Entry Register*

Writes to the MFC_TLB_Invalidate_Entry register invalidate all four ways within a TLB congruence class. To invalidate the entire array, this MMIO request must be executed 64 times for each index position.

To invalidate an entry, a 6-bit index must be selected from the VPN and LP fields. The correct range to select is determined from the L and LS bits in conjunction with the SMM_HID[0:3] bits. If L equals '0', then the small-page (4 KB) index range is selected. When L equals '1', the page size is determined as shown in *Table 4-14*.

*Table 4-14. MFC_TLB_Invalidate_Entry Index Selection*

| Page Size | MFC_TLB_Invalidate_Entry Index Range |
|---|---|
| 4 KB | MFC_TLB_Invalidate_Entry[56:61] |
| 64 KB | MFC_TLB_Invalidate_Entry[52:57] |
| 1 MB | MFC_TLB_Invalidate_Entry[48:53] |
| 16 MB | MFC_TLB_Invalidate_Entry[44:49] |

The MFC_TLB_Invalidate_Entry data is actually written into the MFC_TLB_VPN register because these two registers share VPN format data. There is no conflict in using the MFC_TLB_VPN register for both functions. The MFC_TLB_Invalidate_Entry register is write-only. The MFC_TLB_VPN register reads data retrieved from the array and not the MFC_TLB_VPN register. Reads from the MFC_TLB_Invalidate_Entry MMIO register return '0' data.

### *MFC_TLB_Invalidate_All Register*

The MFC_TLB_Invalidate_All register defined by the *Cell Broadband Engine Architecture* is not implemented in the Cell/B.E. and PowerXCell 8i processors.

### *MFC_ACCR Register*

The MFC Address Compare Control Register (MFC_ACCR) implements a data address compare (DAC) mechanism that allows the detection of a DMAC access to either a virtual page with the Address Compare (AC) bit in the PTE set to '1', or to a range of addresses within the LS. See the *Cell Broadband Engine Architecture* for the DAC-detection algorithm.

### *MFC_DSISR Register*

The MFC Data-Storage Interrupt Status Register (MFC_DSISR) contains status bits relating to the last data-storage interrupt (DSI) generated by the SMM. As DSIs are generated by the SMM, the status bits are written to this register in the middle of a translation request. If software attempts to write this register at the same time that the SMM is writing the register, the SMM hardware data overrides the software-written data (this makes the interrupt information available for software debug).

The MFC_DSISR register is used in this way, whether the SMM is operating in virtual addressing mode or real addressing mode. The SMM only writes to this register when a DSI is asserted. The register is locked until the Restart bit, MFC_CNTL[R], is set to '1'. The register is never cleared; it is only overwritten with new DSI information.

### *MFC_DAR Register*

The MFC Data Address Register (MFC_DAR) contains the 64-bit EA from a DMA command. The MFC_DAR register is writable both by the SMM hardware and by software. During translation, the SMM stores the EA in this register for every DMA command until a potential error is discovered. Then, the register becomes locked until the potential error passes or a previously reported error is cleared by setting the Restart bit, MFC_CNTL[R], to '1'.

If SMM hardware and software attempt to write the MFC_DAR register simultaneously, the SMM hardware data overrides the software-written data (this makes the interrupt information available for software debug). This register is used for both virtual and real addressing modes. The EA is stored in this register for all SMM errors and interrupts, except parity errors on MMIO array reads.

#### 4.3.5.4 *Storage Coherency*

The TLB maintains storage coherency by means of snooped TLBIE requests. The TLBIE requests are broadcast over the EIB and collected by the SBI unit. The SBI unit checks the logical partition ID (LPID) for the request to see if the TLBIE is targeted for its partition. If so, the request is sent to the atomic unit in the MFC. The atomic unit generates the TLBIE request for the SMM.

#### 4.3.5.5 *TLB After POR*

All TLB entries and SLB entries are invalidated after power-on reset (POR).

### 4.3.6 Real Addressing Mode

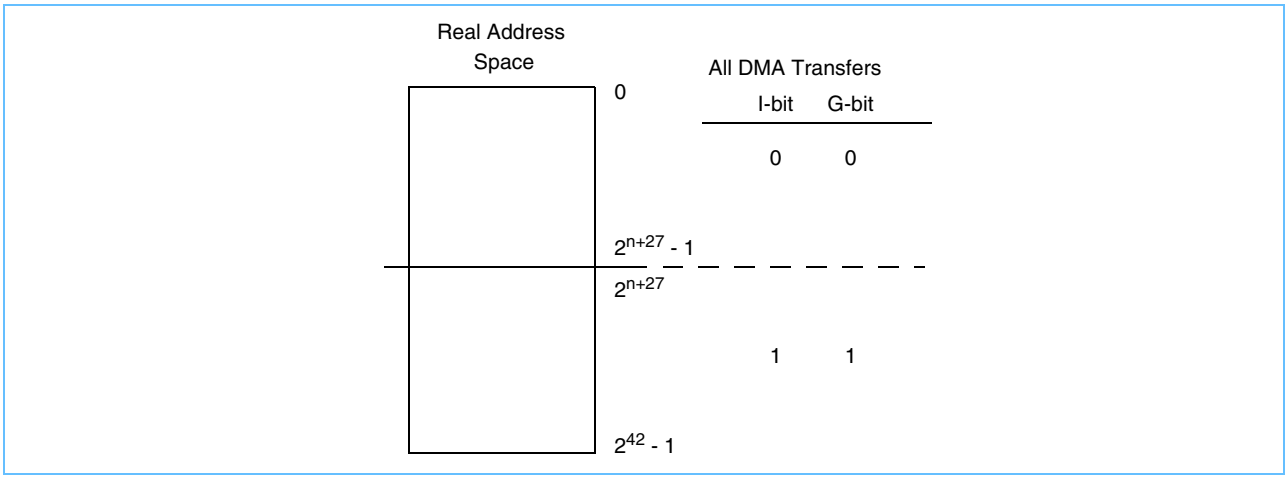
If PPE privileged software has set MFC_SR1[R] = '0', the main-storage side of MFC accesses operates in real-addressing mode. In this mode, EAs provided to the MFC by programs are used as RAs. However, unlike the PPE's real-addressing mode, described in *Section 4.2.8* on page 100, the SPE's real-addressing mode does not use the RMOR or HRMOR base-address registers, or the MSR[HV] and LPCR[LPES] fields and the RMLR register.

The MFC supports a real-mode address boundary facility, specified by the 4-bit MFC_RMAB[RMB] register field. The RMB field is used to control caching-inhibited (I) and guarded (G) storage-control attributes when the MFC is running in real mode. The field sets a boundary, illustrated in *Figure 4-9* on page 118, between storage that is considered well-behaved and cacheable and storage that is not. The granularity of this boundary is a power of 2 from 256 MB up to 4 TB for $1 \le n \le 15$.

Specifically, when in real-addressing mode, if MFC_RMAB[RMB] is set to a value of 'n' (for $1 \le n \le 15$), only those accesses within the first $2^{n+27}$ bytes of the real-address space are considered cacheable and well-behaved (that is, load-load or load-store ordering can be assumed). All accesses outside the first $2^{n+27}$ bytes are considered neither cacheable nor well-behaved. If MFC_RMAB[RMB] = x'0', no accesses are considered well-behaved and cacheable.

The RMB field has no effect when the MFC is running with translation on (MFC_SR1[R] = '1'). PPE privileged software must suspend all MFC operations before modifying the contents of MFC_RMAB[RMB] field.

*Figure 4-9. Real-Mode Storage Boundary*




### 4.3.7 Exception Handling and Storage Protection

The SMM checks for storage-protection errors during both virtual-mode and real-mode address translation, although real-mode checking is very limited. The DMAC also checks for LS-address compares. *Table 4-15* describes translation, address-compare, and parity errors detected and reported by the SMM and DMAC. These exceptions are reported in the following MMIO registers:

- Class 1 Interrupt Status Register (`INT_Stat_class1`)
- MFC Data-Storage Interrupt Status Register (`MFC_DSISR`)
- MFC Fault Isolation Register (`MFC_FIR`)

See the *Cell Broadband Engine Registers* document for a description of the seven Fault Isolation Registers (FIRs), including MFC Fault Isolation Register (`MFC_FIR`), MFC Fault Isolation Register Error Mask (`MFC_FIR_Err`), and MFC Fault Isolation Register Checkstop Enable (`MFC_FIR_ChkStpEnbl`) registers.

*Table 4-15. Translation, MMIO, and Address Match-Generated Exceptions* (Sheet 1 of 2)

| Exception Name | INT_Stat_class1 Register | | | MFC_DSISR Register | DMAC Suspend | Data Error (DERR) | MFC_FIR Register |
|---|---|---|---|---|---|---|---|
| | LP or LG Bits[60:61] | MF (DSI) Bit[62] | SF Bit[63] | | | | |
| SLB Segment Fault | | | X | | | | |
| Atomic access to caching-inhibited page when address relocation is enabled | | X | | A bit[37] | X | | |
| Atomic access to caching-inhibited page when address relocation is disabled (real mode) | | X | | A bit[37] | X | | |
| Page protection violation (PP) | | X | | P bit[36] | | | |

*Table 4-15. Translation, MMIO, and Address Match-Generated Exceptions* (Sheet 2 of 2)

| Exception Name | INT_Stat_class1 Register | | | MFC_DSISR Register | DMAC Suspend | Data Error (DERR) | MFC_FIR Register |
|---|---|---|---|---|---|---|---|
| | LP or LG Bits[60:61] | MF (DSI) Bit[62] | SF Bit[63] | | | | |
| Local-storage address compare (reported by DMAC and not reported as SMM exception) | X | | | | X | | |
| Data address compare (DAC) | | X | | C bit[41] | X | | |
| TLB page fault (hardware or software) | | X | | M bit[33] | | | |
| Parity error detected during SLB/TLB translation operation | | | | | X | | SLB or TLB bits [43:44} |
| Parity error detected during MMIO read to SLB or TLB array (only reported through DERR bit attached to the MMIO data returned to the requester) | | | | | | X | |

#### 4.3.7.1 *Address-Translation Exceptions*

For all translation exceptions reported by the SMM, except parity errors detected on SLB or TLB MMIO reads, the 64-bit EA is saved to the MFC_DAR register. If a DSI exception occurs, information about which kind of DSI error and other related information is written into the MFC_DSISR register, and the MFC Command Queue Error ID Register (MFC_CMDQ_Err_ID) entry index is saved to the MFC Data-Storage Interrupt Pointer Register (MFC_DSIPR). The MFC_DAR, MFC_DSIPR, and MFC_DSISR registers are locked to the value of the current outstanding exception or miss.

The SMM supports only one outstanding translation miss or exception, excluding MMIO-read parity errors. The SMM has the ability to serve hits under one outstanding translation miss or exception. After the first outstanding miss or exception is detected by the SMM, any subsequent exception or miss of an MFC-command translation is reported back to the DMAC as a miss. The DMAC sets the dependency bit associated with the MFC command and stalls the command until the dependency bit is cleared.

When the outstanding miss or exception condition is cleared, all 24 dependency bits are cleared (16 for the MFC SPU command queue and 8 for the MFC proxy command queue), and the DMAC resends the oldest MFC command to the SMM for translation. Exceptions are cleared by setting the Restart bit, MFC_CNTL[R], to ‘1’. Setting MFC_CNTL[R] to ‘1’ also unlocks the MFC_DAR and MFC_DSISR registers and allows the SMM to report a new exception. The MFC_CNTL[Sc] bit should be cleared to ‘0’ to unsuspend the MFC command queues only after all exception conditions have been cleared. The MFC_DSIPR register is unlocked by reading the register or purging the MFC command queues.

#### 4.3.7.2 *Address-Compare Exceptions*

LS-address-compare exceptions occur when the LS is accessed within the address range specified in the MFC Local Storage Address Compare Register (MFC_LSACR). The DMAC writes the MFC Local Storage Compare Results Register (MFC_LSCRR) with the LS address that met the compare conditions set in the MFC_LSACR register, along with the MFC command-queue index and the DMA command type. This exception causes the DMAC to suspend. To clear the error, the LP or LG bits in the INT_Stat_class1 register should be cleared to '0', and the MFC_CNTL[Sc] bit should be cleared to '0' to unsuspend the MFC command queues. See the MFC Local Storage Compare Results (MFC_LSCRR) MMIO register and the MFC Local Storage Address Compare (MFC_LSACR) MMIO register in *Cell Broadband Engine Registers* for more information.

#### 4.3.7.3 *Segment Exceptions*

SLB segment faults occur when there is no matching ESID in the SLB for the EA in the translation request, or the Valid bit is not set in the SLB for the matching ESID entry. When this fault occurs, the MFC_DAR register contains the 64-bit EA, as described in *Section 4.3.7.1 Address-Translation Exceptions* on page 119. The MFC_DAR register and the SLB can be read to determine which entry to replace in the SLB. Then, the entry should be invalidated with a write to the SLB_Invalidate_Entry or SLB_Invalidate_All register, followed by writes to the SLB_Index, SLB_ESID, and SLB_VSID registers to put the proper entry into the SLB. The INT_Stat_class1[SF] should then be cleared to '0', and the MFC_CNTL[R] bit should be set to '1' to resume normal SMM and DMAC operation.

#### 4.3.7.4 *Data-Storage Exceptions*

All DSI exceptions are defined in the *PowerPC Operating Environment Architecture, Book III*. When these exceptions occur, the SMM writes the 64-bit EA to the MFC_DAR register and records the DSI exception type in the MFC_DSISR register, as shown in *Table 4-15* on page 118. That table also shows which DSI exceptions cause the DMAC to suspend. MFC_DSISR[S] is also set to '1' if the DMAC translation request was a **put[rlfs]**, **putll[u]c**, or **sdcrz** operation. The SMM reports the DSI to the SBI to set the INT_Stat_class1[MF] bit to '1'. Software can read the MFC_DAR and MFC_DSISR registers and the SLB and TLB arrays to determine the type of DSI fault and the method for fixing it. After the error condition has been fixed, the INT_Stat_class1[MF] bit should be cleared to '0', and the MFC_CNTL[R] bit should be set to '1'.

#### 4.3.7.5 *Parity Exceptions*

Software might be able to recover from an SLB or TLB parity errors by writing over the faulty entry. Parity errors detected during address translation set bits in the MFC_FIR register as shown in *Table 4-15* on page 118. Only parity errors detected on an address translation cause the MFC_DAR to be locked with the EA value captured in it. After the error has been fixed, the MFC_FIR bits should be cleared and the MFC_DSIPR register written to restart the SMM and unsuspend the DMAC. Parity errors detected during MMIO reads on the SLB or TLB cause the DERR bit to be set on the EIB when the MMIO data is returned to the requester. The SMM can support an unlimited number of MMIO-read parity errors because they do not lock the SMM or DMAC interrupt registers or affect address translation.

# 5. Memory Map

*Section 4 Virtual Storage Environment* on page 79 describes how memory addresses are translated (mapped) between virtual-address, effective-address, and real-address forms. This section describes how real addresses are mapped to the real-address space, which includes physical memory, memory-mapped registers, and the Synergistic Processor Element (SPE) local storage (LS) areas.

## 5.1 Introduction

*Figure 5-1* on page 122 shows the real-address space memory map of the Cell Broadband Engine Architecture (CBEA) processors[1]. The locations in the map are indicated by the sum of two values—an explicit offset added to a base real address. For example, "`x'1E 0000' + SPE3 BE_MMIO_Base`" means that the offset `x'1E 0000'` is added to the real address contained in the `SPE3 BE_MMIO_Base` base-address register. The region marked "Available to Software" can be used by software for any purpose if it is known to be populated with memory chips or external memory-mapped I/O (MMIO) registers.

Although the *Cell Broadband Engine Architecture (CBEA)* defines only one base register (`BP_Base`) for relocating all of this address space, the CBEA processors implement multiple base registers for this relocation, each associated with one of the main logic groups—PowerPC Processor Element (PPE), Synergistic Processor Elements (SPEs), Cell Broadband Engine interface (BEI) unit, memory interface controller (MIC), and pervasive logic (PRV). The base-register values are initialized from the configuration ring during power-on reset (POR), as described in the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*. System developers should use the same address for all eight SPE base addresses. *Table 5-2* on page 124 provides details about the allocated regions of memory.

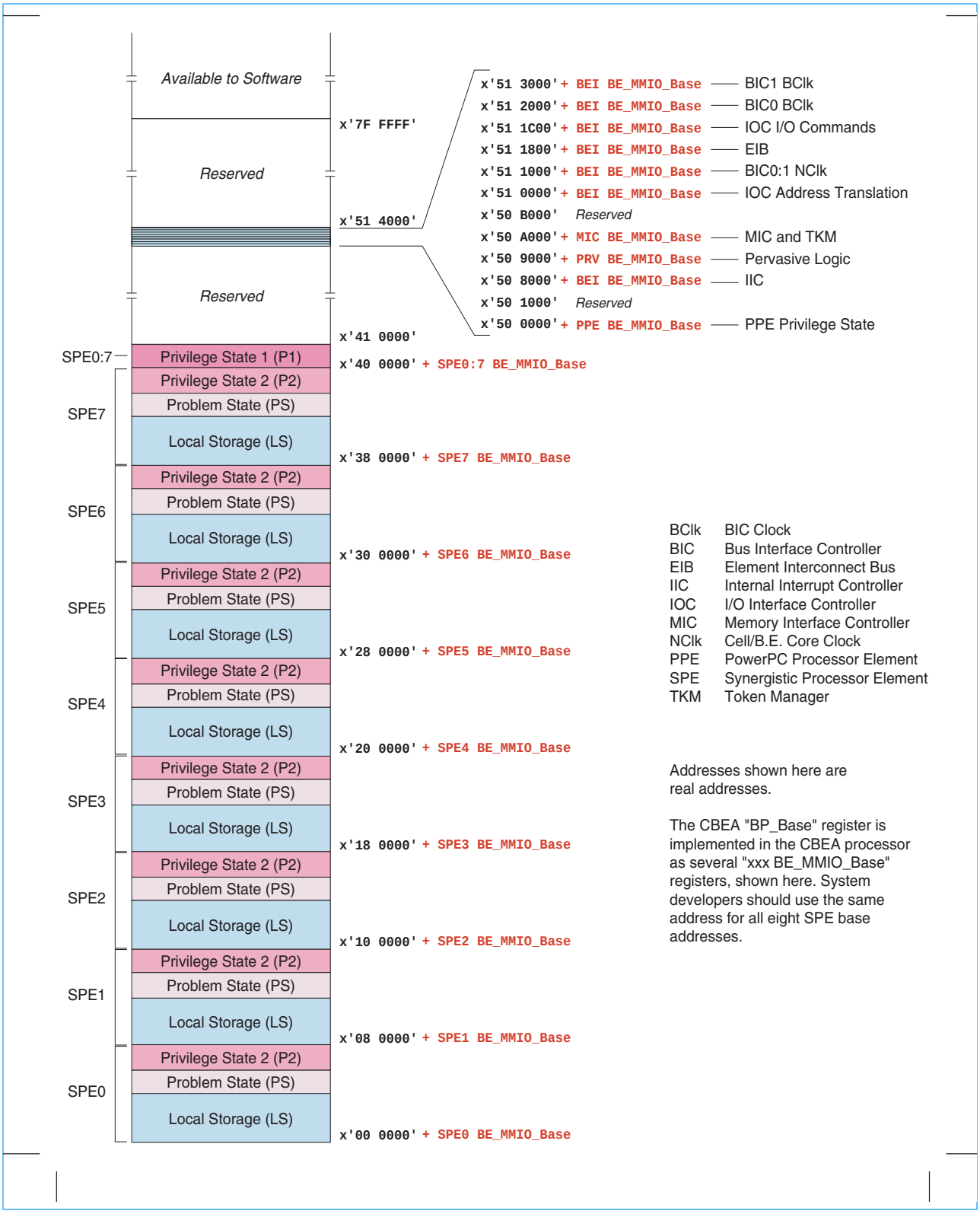
Areas of memory that are marked as *reserved* in *Figure 5-1* on page 122 are not assigned to any functional unit and should not be read from or written to: doing so will cause serious errors in software.[2] Areas of memory that are not identified as allocated in *Figure 5-1*, and that are known to be populated with memory devices, are available to software for any purpose.

When accessing memory that is shared by the PPE, SPEs, or I/O devices, it might be necessary to synchronize storage accesses. The facilities for doing so are described in *Section 20 Shared-Storage Synchronization* on page 561. Loading cooperative PPE and SPE programs into memory requires loading the PPE and SPE programs into distinct memory spaces. The procedures for doing so are described in *Section 14 Objects, Executables, and SPE Loading* on page 397. Additional information regarding stack frames and other software conventions for the use of memory are given in the *SPU Application Binary Interface Specification*.

The documentation for the CBEA processors defines an MMIO register as any internal or external register that is accessed through the main-storage space with load and store instructions, whether or not the register is associated with an I/O device.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.
2. See *Reserved Regions of Memory and Registers* on page 32 for further details.

*Figure 5-1. CBEA Processor Memory Map*



### 5.1.1 Configuration-Ring Initialization

*Table 5-1* shows the memory-base registers for each of the main logic groups. These base registers are initialized from the configuration ring during power-on reset (POR), as described in the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

*Table 5-1. Memory-Base Registers Loaded by Configuration Ring*

| Memory-Base Register | Logic Group | Number of Bits in Configuration-Ring Field | Section of Configuration Ring |
|---|---|---|---|
| SPE0 BE_MMIO_BASE<br>SPE1 BE_MMIO_BASE<br>SPE2 BE_MMIO_BASE<br>SPE3 BE_MMIO_BASE<br>SPE4 BE_MMIO_BASE<br>SPE5 BE_MMIO_BASE<br>SPE6 BE_MMIO_BASE<br>SPE7 BE_MMIO_BASE | • SPE0:7<br>– Local Storage (LS)<br>– Problem State (PS)<br>– Privilege State 2 (P2)<br>– Privilege State 1 (P1) | 19<br>for each SPE | MFC |
| PPE BE_MMIO_BASE | • PPE | 30 | PowerPC Processor Storage Subsystem (PPSS) |
| BEI BE_MMIO_BASE | • Cell Broadband Engine Interface (BEI) Unit:<br>– Element Interconnect Bus (EIB)<br>– Internal Interrupt Controller (IIC)<br>– I/O Controller (IOC) I/O-Address Translation<br>– I/O Controller (IOC) I/O Commands<br>– Bus Interface Controller (BIC) | 22 | IOC |
| MIC BE_MMIO_BASE | • Memory Interface Controller (MIC) | 30 | MIC |
| PRV BE_MMIO_BASE | • Pervasive Logic (PRV)[1] | 30 | MIC |

1. Pervasive logic includes power management, thermal management, clock control, software-performance monitoring, trace analysis, and so forth.

The bit-fields for the base registers in the configuration ring correspond to the most-significant bits of the 42-bit real address implemented in the CBEA processors. The most-significant 19 bits of all these configuration-ring fields should be set to the same value. If a configuration-ring field has more than 19 bits, the additional bits should be set to a value consistent with the settings in *Table 5-2* on page 124 for the starting address of that area of memory. Although each SPE has its own `BE_MMIO_Base` register field in the configuration ring, all such fields should be initialized with the same value.

### 5.1.2 Allocated Regions of Memory

*Table 5-2* on page 124 provides details about the allocated regions of memory.

*Table 5-2. CBEA Processor Memory Map (Sheet 1 of 2)*

| Base Register from Configuration Ring | Offset From Base Register | | Area[1] | | Size in Hexadecimal |
|---|---|---|---|---|---|
| | Start | End | | | |
| SPE0 BE_MMIO_Base | x‘00 0000’ | x‘03 FFFF’ | SPE0 | Local Storage | x‘4 0000’ |
| | x‘04 0000’ | x‘05 FFFF’ | | Problem State | x‘2 0000’ |
| | x‘06 0000’ | x‘07 FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE1 BE_MMIO_Base | x‘08 0000’ | x‘0B FFFF’ | SPE1 | Local Storage | x‘4 0000’ |
| | x‘0C 0000’ | x‘0D FFFF’ | | Problem State | x‘2 0000’ |
| | x‘0E 0000’ | x‘0F FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE2 BE_MMIO_Base | x‘10 0000’ | x‘13 FFFF’ | SPE2 | Local Storage | x‘4 0000’ |
| | x‘14 0000’ | x‘15 FFFF’ | | Problem State | x‘2 0000’ |
| | x‘16 0000’ | x‘17 FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE3 BE_MMIO_Base | x‘18 0000’ | x‘1B FFFF’ | SPE3 | Local Storage | x‘4 0000’ |
| | x‘1C 0000’ | x‘1D FFFF’ | | Problem State | x‘2 0000’ |
| | x‘1E 0000’ | x‘1F FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE4 BE_MMIO_Base | x‘20 0000’ | x‘23 FFFF’ | SPE4 | Local Storage | x‘4 0000’ |
| | x‘24 0000’ | x‘25 FFFF’ | | Problem State | x‘2 0000’ |
| | x‘26 0000’ | x‘27 FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE5 BE_MMIO_Base | x‘28 0000’ | x‘2B FFFF’ | SPE5 | Local Storage | x‘4 0000’ |
| | x‘2C 0000’ | x‘2D FFFF’ | | Problem State | x‘2 0000’ |
| | x‘2E 0000’ | x‘2F FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE6 BE_MMIO_Base | x‘30 0000’ | x‘33FFFF’ | SPE6 | Local Storage | x‘4 0000’ |
| | x‘34 0000’ | x‘35 FFFF’ | | Problem State | x‘2 0000’ |
| | x‘36 0000’ | x‘37 FFFF’ | | Privilege 2 Area | x‘2 0000’ |
| SPE7 BE_MMIO_Base | x‘38 0000’ | x‘3B FFFF’ | SPE7 | Local Storage | x‘4 0000’ |
| | x‘3C 0000’ | x‘3D FFFF’ | | Problem State | x‘2 0000’ |
| | x‘3E 0000’ | x‘3F FFFF’ | | Privilege 2 Area | x‘2 0000’ |

BClk: BIC Clock
BIC: Bus Interface Controller
EIB: Element Interconnect Bus
IIC: Internal Interrupt Controller
IOC: I/O Interface Controller
NClk: Core Clock
PRV: Pervasive Logic

1. See *Section 11.3.1 Access Privilege* on page 346 for definitions of SPE privilege 1 and privilege 2.
2. Areas of memory that are marked as reserved are not assigned to any functional unit and should not be read from or written to: doing so will cause serious errors in software.
3. Areas of memory that are not identified as allocated in *Figure 5-1* on page 122, and that are known to be populated with memory chips or external MMIO registers, are available to software for any purpose.

*Table 5-2. CBEA Processor Memory Map (Sheet 2 of 2)*

| Base Register from Configuration Ring | Offset From Base Register | | Area[1] | | Size in Hexadecimal |
|---|---|---|---|---|---|
| | Start | End | | | |
| SPE0 BE_MMIO_Base | x‘40 0000’ | x‘40 1FFF’ | SPE0 | Privilege 1 Area | x‘2000’ |
| SPE1 BE_MMIO_Base | x‘40 2000’ | x‘40 3FFF’ | SPE1 | | x‘2000’ |
| SPE2 BE_MMIO_Base | x‘40 4000’ | x‘40 5FFF’ | SPE2 | | x‘2000’ |
| SPE3 BE_MMIO_Base | x‘40 6000’ | x‘40 7FFF’ | SPE3 | | x‘2000’ |
| SPE4 BE_MMIO_Base | x‘40 8000’ | x‘40 9FFF’ | SPE4 | | x‘2000’ |
| SPE5 BE_MMIO_Base | x‘40 A000’ | x‘40 BFFF’ | SPE5 | | x‘2000’ |
| SPE6 BE_MMIO_Base | x‘40 C000’ | x‘40 DFFF’ | SPE6 | | x‘2000’ |
| SPE7 BE_MMIO_Base | x‘40 E000’ | x‘40 FFFF’ | SPE7 | | x‘2000’ |
| | x‘41 1000’ | x‘4F FFFF’ | *Reserved*[2] | | |
| PPE BE_MMIO_Base | x‘50 0000’ | x‘50 0FFF’ | PPE | Privilege Area | x‘1000’ |
| | x‘50 1000’ | x‘50 7FFF’ | *Reserved*[2] | | |
| BEI BE_MMIO_Base | x‘50 8000’ | x‘50 8FFF’ | IIC | | x‘1000’ |
| | x‘50 9000’ | x‘50 93FF’ | *Reserved*[2] | | |
| PRV BE_MMIO_Base | x‘50 9400’ | x‘50 97FF’ | PRV | Performance Monitor | x‘0400’ |
| | x‘50 9800’ | x‘50 9BFF’ | | Thermal and Power Management | x‘0400’ |
| | x‘50 9C00’ | x‘50 9FFF’ | | Reliability, Availability, Serviceability (RAS) | x‘0400’ |
| MIC BE_MMIO_Base | x‘50 A000’ | x‘50 AFFF’ | MIC and TKM | | x‘1000’ |
| | x‘50 B000’ | x‘50 FFFF’ | *Reserved*[2] | | |
| BEI BE_MMIO_Base | x‘51 0000’ | x‘51 0FFF’ | IOC | I/O Address Translation | x‘1000’ |
| | x‘51 1000’ | x‘51 13FF’ | BIC | BIC0 NClk | x‘0400’ |
| | x‘51 1400’ | x‘51 17FF’ | | BIC1 NClk | x‘0400’ |
| | x‘51 1800’ | x‘51 1BFF’ | EIB | | x‘0400’ |
| | x‘51 1C00’ | x‘51 1FFF’ | IOC | I/O Commands | x‘0400’ |
| | x‘51 2000’ | x‘51 2FFF’ | BIC | BIC0 BClk | x‘1000’ |
| | x‘51 3000’ | x‘51 3FFF’ | | BIC1 BClk | x‘1000’ |
| | x‘51 4000’ | x‘7F FFFF | *Reserved*[2] | | |
| | x‘80 0000’ | System Maximum | *Available to Software*[3] | | |

BClk: BIC Clock
BIC: Bus Interface Controller
EIB: Element Interconnect Bus

IIC: Internal Interrupt Controller
IOC: I/O Interface Controller
NClk: Core Clock
PRV: Pervasive Logic

1. See *Section 11.3.1 Access Privilege* on page 346 for definitions of SPE privilege 1 and privilege 2.
2. Areas of memory that are marked as reserved are not assigned to any functional unit and should not be read from or written to: doing so will cause serious errors in software.
3. Areas of memory that are not identified as allocated in *Figure 5-1* on page 122, and that are known to be populated with memory chips or external MMIO registers, are available to software for any purpose.

### 5.1.3 Reserved Regions of Memory

Reserved areas are not assigned to any unit and should not be read from or written to; doing so will cause serious errors in software as follows. For reads or writes generated from the SPE to unassigned reserved spaces, at least one of the following `MFC_FIR[46,53,56,58,61]` bits will be set and in most cases will cause a checkstop. For reads or writes generated from the PPE to unassigned reserved spaces, at least one of the `CIU_FIR[7,8]` will be set and a checkstop will occur. For reads or writes generated from the IOC to unassigned reserved spaces, the IOC will respond back to the I/O interface (IOIF) device that sourced the address request with an error (ERR) response. No IOC Fault Isolation Register (FIR) bits are set.

### 5.1.4 The Guarded Attribute

If some areas of real memory are not fully populated with memory devices, privileged software might want to give the pages that map these areas the guarded attribute, as described in *Section 4 Virtual Storage Environment* on page 79. The guarded attribute can protect the system from undesired accesses caused by out-of-order load operations or instruction prefetches that might lead to the generation of a machine check exception. Also, the guarded attribute can be used to prevent out-of-order (speculative) load operations or prefetches from occurring to I/O devices that produce undesired results when accessed in this way. However, the guarded attribute can result in lower performance than is the case with unguarded storage.

## 5.2 PPE Memory Map

### 5.2.1 PPE Memory-Mapped Registers

*Figure 5-1* on page 122 and *Table 5-1* on page 123 show the PPE memory area and its offset from the value in the `PPE BE_MMIO_Base` register that is filled by the configuration ring at POR. All locations within this area that are not populated as described in this section are reserved.

*Table 5-3* lists the groups of PPE MMIO registers in this memory area. All registers in this area can only be accessed by privileged software. For details about the registers in these groups, see the *Cell Broadband Engine Registers* document and the *PowerPC Operating Environment Architecture, Book III*.

*Table 5-3. PPE Privileged Memory-Mapped Register Groups*

| x‘50 0000’ + Offset From Base Register | | Register Group |
|---|---|---|
| Start | End | |
| x‘0310’ | x‘0870’ | Level 2 (L2) MMIO Registers |
| x‘0938’ | x‘0948’ | Core Interface Unit (CIU) MMIO Registers |
| x‘0A48’ | x‘0A48’ | Noncacheable Unit (NCU) MMIO Registers |
| x‘0B48’ | x‘0B78’ | Bus Interface Unit (BIU) MMIO Registers |

### 5.2.2 Predefined Real-Address Locations

There are four areas of the PPE real-address space[3] that have predefined uses, as shown in *Table 5-4*. The first 256 bytes are reserved for operating-system use. The remainder of the 4 KB page beginning at address x‘0000’ is used for exception vectors. The two contiguous 4 KB pages beginning at real address x‘1000’ are reserved for implementation-specific purposes. A contiguous sequence of real pages beginning at the real address specified by Storage Description Register 1 (SDR1) contains the page table. See the *PowerPC Operating Environment Architecture, Book III* for information about the SDR1 register.

*Table 5-4. Predefined Real-Address Locations*

| Real Address | | Predefined Use |
|---|---|---|
| Start | End | |
| x‘0000’ | x‘00FF’ | Operating system |
| x‘0100’ | x‘0FFF’ | Exception vectors[1] |
| x‘1000’ | x‘2FFF’ | Implementation-specific use |
| Software-specified in Storage Description Register 1 (SDR1) | | Page table |

1. See *Section 9 PPE Interrupts* on page 239 for the assignment of the exception-vector offsets.

## 5.3 SPE Memory Map

Each SPE is allocated an area of memory for the following purposes:

- LS, if aliased to main-storage space (see *Section 5.3.1* on page 128)
- Problem State MMIO Registers
- Privilege State 2 MMIO Registers
- Privilege State 1 MMIO Registers

*Figure 1-2* on page 47 shows the relationship between an SPE’s LS and main storage. SPE software accesses its LS directly, by fetching instructions and loading and storing data using an LS address (LSA). SPE software has no direct access to main storage; instead, it must use the memory flow controller (MFC) DMA facilities to move data between the LS and the effective-address (EA) space.

Access, by an SPE, to its channels that are associated with its MMIO registers can be done through the use of read channel (**rdch**) and write channel (**wrch**) instructions. See *Section 17 SPE Channel and Related MMIO Interface* on page 447 and *Section 19 DMA Transfers and Interprocessor Communication* on page 513 for details.

Access to an SPE’s MMIO registers by the PPE and other SPEs and devices that can generate EAs is done through the EA space. As indicated, the SPE’s MMIO registers are divided into three groups, based on privilege:

3. In the real addressing mode for the PPE, accesses to main storage are performed in a manner that depends on the contents of the MSR[HV] bit, the LPCR[LPES] field, and the HRMOR, RMLR, and RMOR registers. For details, see *Section 4.2.8 Real Addressing Mode* on page 100, *Section 11 Logical Partitions and a Hypervisor* on page 331, and *PowerPC Operating Environment Architecture, Book III*.

- *Privilege State 1 (most privileged)*—Used by a hypervisor (*Section 11* on page 331) to manage an SPE on behalf of a logical partition or, if there is no hypervisor, used by the operating system.
- *Privilege State 2*—Used by the operating system in a logical partition to manage the SPE within the partition or, if there is no hypervisor, used by the operating system.
- *Problem State (Privilege 3, least privileged*)—Used by Problem State (user, or application) software, if direct access to the SPE from user space is supported by the operating system.

### 5.3.1 SPE Local-Storage Memory Map

The LS area for each SPE is 256 KB in size. If these LS areas are aliased into the EA space, they begin at the real address (RA) offsets shown in *Table 5-2* on page 124.

#### 5.3.1.1 *LS Aliasing to Main Storage*

Privileged software on the PPE can alias an SPE's LS address space to the main-storage EA space. The aliasing is done by setting the LS Real-Address Decode (D) bit of the MFC State Register One (`MFC_SR1`) to '1'. This aliasing allows the PPE, other SPEs, and external I/O devices to physically address an LS through a translation method supported by the PPE or SPE, or, for I/O devices, through a translation method supported by the EIB (*Section 7.4* on page 176).

When aliasing is implemented, each LS address (LSA) is assigned an RA within the effective-address (EA) space. Privileged software can, for example, map an LSA space into the EA space of an application to allow DMA transfers between the LS of one synergistic processor unit (SPU) and the LS of another SPU. Data transfers that use the LS area aliased in main storage should do so as caching-inhibited storage, because these accesses are not coherent with the SPU LS accesses (that is, SPU load, store, instruction fetch) in its LS domain.

Each SPE has its own `MFC_SR1` register, located in the privilege 1 memory-map area (*Figure 5-1* on page 122) for that SPE. When an SPE's `MFC_SR1[D]` bit is set to '1', the MFC associated with that SPE decodes RAs into LSAs.

As an alternative to aliasing, privileged software can use the `MFC_SR1[D]` bit to prevent the LS from being addressed using RAs. This can be done to facilitate a more isolated SPU program environment.

The value of the `MFC_SR1[D]` bit has no effect on LS access by MFC DMA commands or on SPU loads and stores from LS, both of which use LSAs to access the LS.

#### 5.3.1.2 *Coherence*

All SPE memory areas are architecturally defined as noncoherent. Privileged software on the PPE should access aliased pages of LS through the main-storage domain. The pages must either have the caching-inhibited attribute (*Section 4 Virtual Storage Environment* on page 79), or software must explicitly manage the coherency of LS with other system caches.

#### 5.3.1.3 *Data-Transfer Performance*

Because aliased LS must be treated as noncacheable, transferring a large amount of data using the PPE load and store instructions can result in poor performance. Data transfers between the LS domain and the main-storage domain should use the MFC DMA commands to avoid stalls.

#### 5.3.1.4 *LS-Access Errors*

MFC commands that access an EA range that maps to its own LS can produce an error or unpredictable results. If the two address ranges of a DMA transfer (translated EA and LSA) overlap and the source is a lower address than the destination, the DMA transfer results in the corruption of the source data. Address overlap is not detectable and does not generate an exception. Therefore, it is software's responsibility to avoid an unintended overlap.

### 5.3.2 SPE Memory-Mapped Registers

*Figure 5-1* on page 122 and *Table 5-1* on page 123 show the SPE memory areas and their offsets from the value in the `SPE BE_MMIO_Base` registers that are filled by the configuration ring at POR. All locations within this area that are not populated as described in this section are reserved.

*Table 5-5* lists the groups of SPE problem-state registers, using SPE0 as an example for address offsets (see *Table 5-2* on page 124 for details). *Table 5-6* lists the groups of SPE privilege state 2 registers, using SPE0 as an example. *Table 5-6* lists the groups of SPE privilege state 1 registers, using SPE0 as an example. For details about the registers in these groups, see the *Cell Broadband Engine Registers* document.

*Table 5-5. SPE Problem-State Memory-Mapped Register Groups*

| Offset From Base Register + x‘4000’ | | Register Group |
|---|---|---|
| Start | End | |
| x‘0000’ | x‘2FFF’ | SPE Multisource Sync Register |
| x‘3004’ | x‘3017’ | MFC Command Parameter Registers |
| x‘3104’[1] | x‘322F’ | MFC Command Queue Control Registers |
| x‘4004’ | x‘4038’ | SPU Control Registers |
| x‘1 400C’ | x‘1 FFFF’ | SPU Signal Notification Registers |

1. This address is shared by multiple registers. See *Table 17-2* on page 450 for details.

*Table 5-6. SPE Privilege-State-2 Memory-Mapped Register Groups* (Sheet 1 of 2)

| Offset From Base Register + x‘6000’ | | Register Group |
|---|---|---|
| Start | End | |
| x‘1100’ | x‘1FFF’ | Segment Lookaside Buffer (SLB) Management Registers |
| x‘2000’ | x‘2FFF’ | Context Save and Restore Register |
| x‘3000’ | x‘3FFF’ | MFC Control Register |

*Table 5-6. SPE Privilege-State-2 Memory-Mapped Register Groups* (Sheet 2 of 2)

| Offset From Base Register + x'6000' | | Register Group |
|---|---|---|
| Start | End | |
| x'4000' | x'4008' | Interrupt Mailbox Register |
| x'4040' | x'4FFF' | SPU Control Registers |
| x'5000' | x'1 FFFF' | Context Save and Restore Registers |

*Table 5-7. SPE Privilege-State-1 Memory-Mapped Register Groups*

| Offset From Base Register + x'40 0000' | | Register Group |
|---|---|---|
| Start | End | |
| x'0000' | x'00FF' | MFC Registers |
| x'0100' | x'01FF' | Interrupt Registers |
| x'0200' | x'028F' | Atomic Unit Control Registers |
| x'0290' | x'03B7' | Fault Isolation Registers |
| x'03B8' | x'03C7' | Miscellaneous Registers |
| x'03C8' | x'03FF' | Reserved |
| x'0400' | x'0547' | MFC Translation Lookaside Buffer (TLB) Management Registers |
| x'0580' | x'0587' | Memory Management Register |
| x'0600' | x'06FF' | MFC Status and Control Registers |
| x'0710' | x'07FF' | Replacement Management Table (RMT) Registers |
| x'0800' | x'087F' | MFC Command Data-Storage Interrupt Registers |
| x'0880' | x'08FF' | DMAC Unit Performance Monitor Control Register |
| x'0900' | x'0BFF' | Real-Mode Support Registers |
| x'0C00' | x'0FFF' | MFC Command Error Register |
| x'1000' | x'101F' | SPU ECC and Error Mask Registers |
| x'1028' | x'13FF' | SPU Performance Monitor Control Registers |
| x'1400' | x'1FFF' | Performance Monitor Register |

## 5.4 BEI Memory-Mapped Registers

The Cell Broadband Engine interface (BEI) unit memory space contains registers for the following functional units:

- Element Interconnect Bus (EIB)
- Internal Interrupt Controller (IIC)
- I/O Controller (IOC) I/O-Address Translation
- I/O Controller (IOC) I/O Commands
- Bus Interface Controller (BIC)

For details about these registers, see the *Cell Broadband Engine Registers* document. For details about I/O functions performed by these units, see *Section 7 I/O Architecture* on page 161.

### 5.4.1 I/O

The IOC command (IOCmd) unit handles commands that originate from the EIB and go to IOIF0 or IOIF1. The IOCmd unit also acts as a proxy for commands that originate from IOIF0 or IOIF1 and go to the EIB. Commands that originate from IOIF0 or IOIF1 are sent with an I/O address. If I/O-address translation is enabled by setting `IOC_IOCmd_Cfg[16]` = '1', the IOCmd unit provides buffering while the I/O addresses are being translated into RAs.

The IOCmd unit accepts commands from the EIB for which the addresses match the RA ranges specified for each IOIF, using the following registers:

- IOC Base Address Register 0 (`IOC_BaseAddr0`)
- IOC Base Address Register 1 (`IOC_BaseAddr1`)
- IOC Base Address Mask Register 0 (`IOC_BaseAddrMask0`)
- IOC Base Address Mask Register 1 (`IOC_BaseAddrMask1`)

These IOC base-address registers are filled, in part, by configuration ring fields at POR. Specifically, hardware copies the following Configuration-Ring fields into the IOC base-address registers:

- Configuration-Ring IOIF0 Base Address and Replacement field is copied to `IOC_BaseAddr0[22:32,53:63]`.
- Configuration-Ring IOIF1 Base Address and Replacement field is copied to `IOC_BaseAddr1[22:32,53:63]`.
- Configuration-Ring IOIF0 Base Address Mask field is copied to `IOC_BaseAddrMask0[0,22:32]`.
- Configuration-Ring IOIF1 Base Address Mask field is copied to `IOC_BaseAddrMask1[0,22:32]`.

The IOC base-address registers are writable by privileged software, so that the I/O portion of the memory map can be relocated after POR.

# 6. Cache Management

The Cell Broadband Engine Architecture (CBEA) processors[1] have several types of caches. Primary among them are the PowerPC Processor Element (PPE) level-1 (L1) instruction and data caches, and its unified level-2 (L2) cache. The contents of these caches are maintained coherently with the contents of main storage on a 128-byte cache-line basis, and the *PowerPC Architecture* cache-control instructions support user-level operations on cache lines. In addition to the L1 and L2 caches, the PPE and the SPEs have other caches, queues, and arrays that enhance performance and can be controlled by software.

The basic PPE cache architecture is described in the *PowerPC Virtual Environment Architecture, Book II*, *PowerPC Operating Environment Architecture, Book III*, and *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors*. Features of the Synergistic Processor Element (SPE) caches are described later in this section. For detailed descriptions of the caches used in memory management, see *Section 4 Virtual Storage Environment* on page 79.

## 6.1 PPE Caches

In addition to the L1 and L2 caches, the PPE also has other caches, queues, and arrays that support memory management and act as predecessors and extensions to the L1 and L2 caches. For example, the PPE contains store queues for holding data that is in flight between the load and store units and the L2 cache, and castout queues for holding modified data that has been pushed out of the L2 cache. All of the L1 and L2 caches and some of the other storage structures are shared between PPE threads, but some storage structures are duplicated for each thread.

Both PPE threads share the execution units, microcode engine, instruction-fetch control, PowerPC processor storage subsystem (PPSS), plus the following caches, arrays, queues, and other storage structures:

- Effective-address-to-real-address translation arrays (ERATs)—I-ERAT (instructions) and D-ERAT (data) and their associated ERAT miss queues
- Translation lookaside buffer (TLB)
- L1 caches (ICache and DCache)
- L2 cache
- Instruction prefetch request queue (IPFQ)
- Load-miss queue (LMQ)
- Core interface unit (CIU) store queue
- Noncacheable unit (NCU) store queue
- L2 store queue

Each PPE thread has its own copy of the following arrays and queues, so that each thread can use that resource independently from the other thread:

- Segment lookaside buffer (SLB)
- Branch history table (BHT), with global branch history

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

- Instruction buffer (IBuf)
- Link stack

### 6.1.1 Configuration

Of the PPE's storage structures, most can be configured by software to some extent using the following registers:

- Hardware Implementation Registers 1 and 4 (`HID1`, `HID4`)
- CIU Mode Setup Register (`CIU_ModeSetup`)
- NCU Mode Setup Register (`NCU_ModeSetup`)

However, software management of the following structures can be more extensive, as described later in this section:

- L1 caches
- L2 cache, and its read-and-claim (RC) queues and replacement management table (RMT)
- ERATs
- TLB and its RMT
- Instruction prefetch queue
- Load-miss queue

All of the instructions for managing the L1 and L2 caches are accessible by problem-state (user) software, but RMTs (*Section 6.3* on page 154) are only accessible by privilege-2-state (supervisor) software, and ERATs (*Section 6.1.7* on page 150) are only accessible by privilege-1-state (hypervisor) software.

### 6.1.2 Overview of PPE Cache

The PPE's several storage structures are distributed throughout its PowerPC processor unit (PPU) and its PowerPC processor storage subsystem (PPSS), as shown in *Figure 6-1* on page 135. This figure differentiates the structures according to their speed with respect to the full core clock speed.

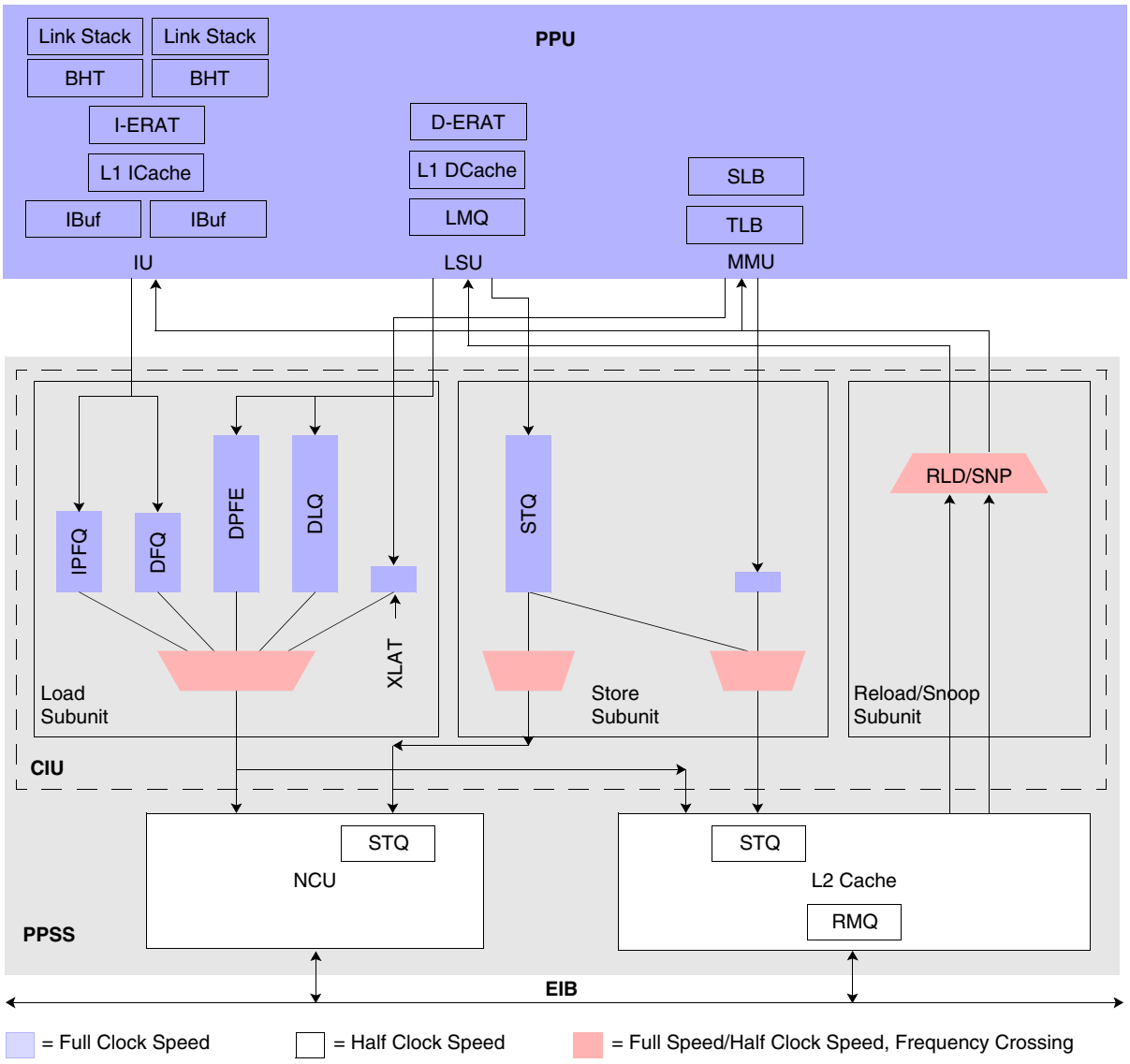
#### 6.1.2.1 *PowerPC Processor Unit*

The major blocks of the PPU that support caches, arrays, queues, and other storage structures are the instruction unit (IU), load and store unit (LSU), and memory management unit (MMU):

- The IU contains the L1 ICache (32 KB), I-ERAT (64 entries), two BHTs (4096 × 2-bit entries and 6 bits of global branch history, per thread), two link stacks (4 entries, per thread), and two instruction buffers (IBufs, 4 instructions × 5 entries, per thread).
- The LSU contains the L1 DCache (32 KB), D-ERAT (64 entries), and load-miss queue (LMQ, 8 entries).
- The MMU contains two SLBs (64 entries, per thread), and the TLB (1024 entries).

*Section 4 Virtual Storage Environment* on page 79 describes the workings of MMU caches. The PPU and all its subunits operate at the full frequency of the core clock.

*Figure 6-1. PPE Cache*




| | |
|---|---|
| BHT | Branch History Table (with global history) |
| DFQ | Demand Instruction Request Queue |
| DLQ | Demand Data load Request Queue |
| DPFE | Data Prefetch Engine |
| EIB | Element Interconnect Bus |
| ERAT | Effective-to-Real Address Translation |
| IPFQ | Instruction Prefetch Request Queue |
| IU | Instruction Unit |
| L1 DCache | Level 1 Data Cache |
| L1 ICache | Level 1 Instruction Cache |
| L2 | Level 2 Cache |
| LMQ | Load Miss Queue |
| LSU | Load and Store Unit |
| MMU | Memory Management Unit |
| NCU | Noncacheable Unit |
| PPSS | PowerPC Processor Storage Subsystem |
| PPU | PowerPC Processor Unit |
| RLD | Reload |
| RMQ | Reload Miss Queue |
| SLB | Segment Lookaside Buffer |
| SNP | Snoop |
| STQ | Store Queue |
| TLB | Translation Lookaside Buffer |
| XLAT | Translate Request Queue |

#### 6.1.2.2 *PowerPC Processor Storage Subsystem*

The major blocks of the PPSS, shown in *Figure 6-1* on page 135, are the core interface unit (CIU), the noncacheable unit (NCU), and L2 cache.

The CIU contains load, store, and reload subunits, which—with the exception of a load-hit-store comparison—work independently of the each other. In the CIU:

- The load subunit contains the instruction prefetch request queue (IPFQ, 2 entries), demand instruction request queue (DFQ, 2 entries), data prefetch engine (DPFE, 8 **dcbt** instructions), demand data load request queue (DLQ, 8 entries), and MMU translate request queue (XLAT, 1 entry).
- The store subunit contains a store queue (STQ, 8 × 16-byte entries).
- The reload/snoop subunit contains a reload/snoop block (RLD/SNP) that performs 2:1 clock-frequency conversion and multiplexing of data from the L2 to the 32-byte reload bus that is shared by the LSU, IU, and MMU. A cache line is sent over four subsequent full-speed cycles totaling 128 bytes. Noncacheable data is sent in one cycle totaling 32 bytes. The reload/snoop subunit also performs 2:1 frequency conversion for the snoop bus.

The NCU has a store queue (STQ, 4 × 64-byte entries), and the L2 has a store queue (STQ, 8 × 64-byte entries). The L2 also has a reload miss queue (RMQ, 6 × 128-byte cache-line entries). The CIU accesses the NCU or L2 based on the Caching-Inhibited (I) storage-attribute bit in the page table entry (PTE).

### 6.1.3 L1 Caches

The PPE has level-1 (L1) instruction and data caches that support the *PowerPC Architecture* cache-control instructions. Cacheability of main-storage accesses (loads, stores, and instruction fetches) can be controlled by the caching-inhibited (I) bit in the page-table entry (PTE) for the accessed page. To be cacheable, the I bit must be cleared to ‘0’.

#### 6.1.3.1 *Features*

- *L1 Instruction Cache (ICache):*
  - 32 KB, 2-way set-associative
  - Read-only
  - 128-byte cache-line size
  - 5-entry, 4-instruction-wide instruction buffer, duplicated for each thread
  - Out-of-order completion of cache-miss load instructions (otherwise in-order execution)
  - Effective-addressed index, real-addressed tags
  - One read port, one write port (one read or write per cycle)
  - Does not support snoop operations
  - Contents not guaranteed to be included in the L2 cache
  - Cache bypass (inhibit) configuration options

- *L1 Data (DCache):*
  - 32 KB, 4-way set-associative
  - Write-through
  - Not write-allocate
  - 128-byte cache-line size
  - 8-entry load-miss queue, shared between threads
  - 16-entry store queue, shared between threads
  - Nonblocking (hit under miss)
  - Supports back-invalidation for snooping
  - Effective-addressed index, real-addressed tags
  - One read port, one write port (one read or write per cycle)
  - Contents fully included in the L2 cache
  - Cache bypass (inhibit) configuration options

Both L1 caches are dynamically shared by the two PPE threads; a cache block can be loaded by one thread and used by the other thread. The coherence block, like the cache-line size, is 128 bytes for all caches. The L1 ICache is an integral part of the PPE's instruction unit (IU), and the L1 DCache is an integral part of the PPE's load/store unit (LSU), as shown in *Figure 2-2 PPE Functional Units* on page 52. Accesses by the processing units to the L1 caches occur at the full frequency of the core clock.

The L1 caches support the *PowerPC Architecture* cache-management instructions. These perform operations such as write-back, invalidate, flush (write-back and invalidate), clear the contents to zeros, touch (which can be used by the compiler to speculatively load the cache line), and hint (that the cache line will probably be accessed soon). See *Section 6.1.6 Instructions for Managing the L1 and L2 Caches* on page 146 for a description. All are problem-state (user) instructions.

#### 6.1.3.2 *ICache Instruction Prefetch*

The PPE uses speculative instruction prefetching (including branch targets) for the L1 ICache and instruction prefetching for the L2 cache. Because the PPE can fetch up to four instructions per cycle and can execute up to two instructions per cycle, the IU will often have several instructions queued in the instruction buffers. Letting the IU speculatively fetch ahead of execution means that L1 ICache misses can be detected early and fetched while the processor remains busy with instructions in the instruction buffers.

In the event of an L1 ICache miss, a request for the required line is sent to the L2. In addition, the L1 ICache is also checked to see if it contains the next sequential cache line. If it does not, a prefetch request is made to the L2 to bring this next line into the L2 (but not into the L1, to avoid L1 ICache pollution). This prefetch occurs only if the original cache miss is committed (that is, all older instructions must have passed the execution-pipeline point—called the *flush point*—at which point their results can be written back to architectural registers). This is especially beneficial when a program jumps to a new or infrequently used section of code, or following a task switch, because prefetching the next sequential line into the L2 hides a portion of the main-storage access latency.

Prefetch requests are not made when a page boundary may potentially be crossed, meaning that the next sequential cache line is on a different aligned 4 KB boundary. Prefetching is also not performed on caching-inhibited instructions (I = '1' in the PTE).

#### 6.1.3.3 *ICache Fetch Misses*

When an L1 ICache miss occurs, a request is made to the L2 for the cache line. This is called a *demand fetch*. To improve performance, the first beat[2] of data returned by the L2 contains the fetch group with the address that was requested, and so returns the data *critical-sector first*. To reduce the latency of an L1 miss, this critical sector of the cache line is sent directly to the instruction pipeline, instead of first writing it into the L1 cache and then rereading it (this is termed *bypassing the cache*).

Each PPE thread can have up to one instruction demand-fetch and one instruction prefetch outstanding to the L2 at a time. This means that in multithread mode, up to four total instruction requests can be pending simultaneously. In multithreading mode, an L1 ICache miss for one thread does not disturb the other thread.

#### 6.1.3.4 *ICache Replacement Algorithm*

The L1 ICache uses a true binary least recently used (LRU) replacement policy, without replacement management table (*Section 6.3* on page 154) locking.

#### 6.1.3.5 *ICache Aliasing*

The ICache is 2-way set associative, so a line that is loaded can be put into one of two ways of a congruence class (set). The smallest page size for the 32 KB L1 ICache is 4 KB, so the same real address can map up to four different locations in the L1 ICache (four congruence classes). This condition is referred to as *aliasing*. Because the ICache is read-only, no particular hazards exist as a result of this. However, the effectiveness of the L1 ICache might be lessened if a great deal of localized aliasing occurs, which is not common in most cases.

Aliasing in the L1 DCache is treated differently, as described in *Section 6.1.3.9* on page 139.

#### 6.1.3.6 *ICache Invalidation*

To help invalidate the contents of the ICache efficiently, the PPE implements a special mode of operation for the **icbi** instruction. This mode can be selected with the Enable Forced **icbi** Match (`en_icbi`) bit in the Hardware Implementation Register 1 (`HID1`). In this mode, all directory lookups on behalf of an **icbi** act as though there were a real-address match, so that all cache lines looked at by the **icbi** are invalidated. As a result, the entire L1 ICache can be invalidated by issuing a series of **icbi** instructions that step through each congruence class of the ICache.

Another way to clear the ICache is to fill it with a set of known values by executing code that touches each cache line. One way to write this code is to have a series of 256 branches to branches whose EAs are sequentially separated by 128 bytes (the line size of the ICache).

The entire ICache can also be invalidated by first clearing the `HID1[flush_ic]` bit to '0' and then setting the same bit to '1'.

2. A data beat is the data sent in one L2 clock cycle, which is two core-clock cycles.

#### 6.1.3.7 *DCache Load Misses*

Loads that miss the L1 DCache enter a load-miss queue for processing by the L2 cache and main storage. Data is returned from the L2 in 32-byte beats on four consecutive cycles. The first cycle contains the critical section of data, which is sent directly to the register file. The DCache is occupied for two consecutive cycles while the reload is written to the DCache, half a line at a time. The DCache tag array is then updated on the next cycle, and all instruction issue is stalled for these three cycles. In addition, no instructions can be recycled during this time. The load-miss queue entry can then be used seven cycles after the last beat of data returns from the L2 to handle another request.

Instructions that are dependent on a load are issued speculatively, assuming a load hit. If it is later determined that the load missed the L1 DCache, any instructions that are dependent on the load are flushed, refetched, and held at dispatch until the load data has been returned. This behavior allows the PPE to send multiple overlapping loads to the L2 without stalling if they are independent. In multithreading mode, this behavior allows load misses from one thread to occur without stalling the other thread.

In general, write-after-write (WAW) hazards do not cause penalties in the PPE. However, if the target register for an instruction is the same as the target register for an outstanding load miss, the new instruction will be flushed, refetched, and held at dispatch until the older load writes its result to the register file, to avoid having stale values written back to the register file.

Some old software performs prefetching by executing load instructions to bring cache lines into the DCache, and the results of the loads are discarded into a “garbage” general purpose register (GPR). This is not effective on the PPE, because this WAW condition effectively serializes the prefetch. This serialization is not a problem if **dcbt** instructions are used instead.

#### 6.1.3.8 *DCache Replacement Algorithm*

The L1 DCache uses a pseudo least recently used (p-LRU) replacement policy. It is a 4-way policy, comparable to the L2 pseudo-LRU 8-way policy described in *Section 6.1.5.3* on page 143, but without replacement management table (*Section 6.3* on page 154) locking. The policy is updated whenever a line is read in the cache. The L1 DCache never needs to cast out data because it uses a write-through scheme (stores are sent to the L2 as well as the L1 DCache).

#### 6.1.3.9 *DCache Indexing*

The DCache is indexed with effective addresses. The data array and cache directory both have a single read and write port (one read or write per cycle). The data and tag arrays operate independently, giving the cache some of the features of having dual ports. Specifically, the cache allows a store to check the directory tags while a previous store writes into the data array.

The DCache tags and the D-ERAT are checked in parallel, similar to the way the ICache is accessed.

#### 6.1.3.10 *DCache Aliasing*

DCache aliasing occurs when separate translation entries have the same real address (RA) mapped from different effective addresses (EAs). Data can be in any of four ways at each of two congruence classes. If the data is not at the congruence class determined by EA[51], but there is

a hit in the cache, the data is found at the other congruence class, addressed by the NOT EA[51]. That line is invalidated and the data is brought in from the L2 and placed at the congruence class determined by EA[51].

If shared memory is used for communication between simultaneously running threads or two processes on the same PPE thread, mapping the shared memory area to the same effective address (EA) in both threads will prevent unnecessary cache invalidations.

#### 6.1.3.11 *DCache Write-Back and Invalidation*

The L1 DCache is a write-through design, so that it never holds data in a modified state. As a result, to perform a write-back of the L1 DCache, the only instruction required is a **sync**, which will force any pending stores in the store queue above the L1 cache to become globally coherent before the **sync** is allowed to complete.

A sticky mode bit is provided in the Hardware Implementation Register 4 (`HID4`) to invalidate the entire contents of the L1 DCache. To invalidate the cache, software must clear the `HID4[l1dc_flsh]` bit to '0', then set the same bit to '1', then execute a **sync** instruction.

#### 6.1.3.12 *DCache Boundary-Crossing*

The L1 DCache is physically implemented with 32-byte sectors. Thus, conditionally microcoded load or store instructions (*Table A-1 PowerPC Instructions by Execution Unit* on page 723) that attempt to perform a L1 DCache access that crosses a 32-byte boundary must be split into several instructions. When one of these misaligned loads or stores first attempts to access the L1 DCache, the misalignment is detected and the pipeline is flushed. The flushed load or store is then refetched, converted to microcode at the decode stage, and split into the appropriate loads or stores, as well as any instructions needed to merge the values together into a single register.

Doubleword integer loads that cross a 32-byte alignment boundary are first attempted as two word-sized loads or stores. If these still cross the 32-byte boundary, they are flushed and attempted again at byte granularity. The word and halfword integer loads behave similarly.

Doubleword floating-point loads and stores that are aligned to a word boundary, but not to a doubleword boundary, are handled in microcode. If these loads or stores are not word aligned, or if they cross a virtual page boundary, an alignment interrupt is taken.

Integer loads and stores that are misaligned but do not cross a 32-byte boundary are not converted into microcode and have the same performance characteristics as aligned loads and stores.

The load string instruction first attempts to use load word instructions to move the data. If the access would cross a 32-byte boundary when it accesses the L1 DCache, the load is flushed and refetched, and it proceeds byte-by-byte. The store string instruction behaves similarly.

#### 6.1.3.13 *L1 Bypass (Inhibit) Configuration*

To bypass the L1 and L2 data caches (that is, to inhibit their use for load/store instructions), set the `HID4[enb_force_ci]` bit to '1' and the `HID4[dis_force_ci]` bit to '0'.

### 6.1.4 Branch History Table and Link Stack

The PPE's instruction unit (IU) is shown in *Figure 6-1* on page 135. For each thread, the IU has a branch history table (BHT) with 4096 × 2-bit entries and 6 bits of global branch history, and a 4-entry link stack. The link stack is used to predict the target address of branch-to-link instructions. The following mode bits enable these functions:

- Branch History Table: Hardware Implementation Register 1, `HID1[bht_pm]` bit
- Global Branch History: Hardware Implementation Register 1, `HID1[dis_gshare]` bit
- Link Stack: Hardware Implementation Register 1, `HID1[en_ls]` bit

Static branch prediction can reduce pollution in the BHT, improving accuracy on branches elsewhere in the code. The BHT is accessed in parallel with the L1 ICache. The global branch history is constantly updated with the latest speculatively-predicted history value. In the event of a pipeline flush including branch misprediction, the correct history state is restored. Using two bits for each BHT entry allows the processor to store weakly and strongly taken or not taken states.

Only instructions that relied on the BHT to predict their direction will ever cause an update to the BHT, meaning that unconditional branches (including conditional branches with BO = 1z1zz), and statically predicted conditional branches do not update the BHT.

Some conditional branches are known ahead of time to be almost always unidirectional. Software can communicate this to the processor by setting the "a" bit in the BO field of the branch instruction. If the "a" bit is '0', then the branch is predicted dynamically, using the BHT as described in the preceding paragraphs. If the "a" bit is '1' in the branch, the BHT is not used, and the BHT will not be updated if this branch is mispredicted. The direction of the static prediction is in the "t" bit of the BO field, which is set to '1' for taken and '0' for not-taken. Software is expected to use this feature for conditional branches for which it believes that the static prediction will be at least as good as the hardware branch prediction. See *PowerPC User Instruction Set Architecture, Book I* for additional details.

### 6.1.5 L2 Cache

The PPE contains a 512 KB unified level-2 (L2) cache. The L2 supports the *PowerPC Architecture* cache-control instructions and it controls cacheability with the caching-inhibited (I) bit in PTEs.

#### 6.1.5.1 *Features*

- 512 KB, 8-way set-associative or direct-mapped[3]
- Unified instruction and data caching
- Inclusive of L1 DCache. Not inclusive of L1 ICache
- Write-back (copy-back)
- 128-byte cache-line size
- 128-byte coherence granularity
- Allocation on store miss

3. See *Section 6.1.5.2* on page 142 for configuring the L2 cache to operate in direct-mapped mode.

- Choice of replacement algorithms: True binary LRU, pseudo-LRU with replacement management table (RMT) locking[4], or direct-mapped
- Hardware coherency (modified, exclusive, recent, shared, invalid [MERSI] plus unsolicited modified [Mu] and tagged [T]), with separate directory for snoops
- Critical quadword-forwarding on data loads and instruction fetches
- 6-entry, 128-byte reload/store-miss queue (read-and-claim queues)
- 6-entry, 128-byte castout queue
- 8-entry, 64-byte fully associative store queue
- 4-entry, 128-byte snoop intervention/push queue
- Store-gathering
- Nonblocking L1 DCache invalidations
- Real-address index and tags
- One read port, one write port (one read or write per cycle)
- Separate snoop directory for all system bus snoops
- Bypass (inhibit) configuration option
- Error correction code (ECC) on data
- Parity on directory tags (recoverable using redundant directories)
- Global and dynamic power management

The L2 maintains full cache-line coherency within the system and can supply data to other processor elements. Logically, the L2 is an inline cache. It is a write-back cache that includes all of the L1 DCache contents but is not guaranteed to include the L1 ICache contents. The L2 is dynamically shared by the two PPE execution threads; a cache block can be loaded by one thread and used by the other thread.

The L2 cache handles all cacheable loads and stores (including **lwarx**, **ldarx**, **stwcx.**, and **stdcx.** instructions), data prefetches, instruction fetches, instruction prefetches, cache operations, and barrier operations. The L2 cache supports the same cache-management instructions as are supported by the L1 caches, as described in *Section 6.1.6 Instructions for Managing the L1 and L2 Caches* on page 146.

#### 6.1.5.2 *Cache Configuration*

The basic L2 cache operation can be configured in the following registers:

- As an alternative to the 8-way set-associative configuration, the L2 cache can be configured to operate in direct-mapped mode by setting the `L2_ModeSetup1[RMT_Mode]` register field to '11'.
- L2 cache replacement algorithm can be configured in the `L2_ModeSetup1[RMT_Mode]` register field. See *Section 6.1.5.3 Replacement Algorithm* on page 143.

4. See *Section 6.3.2* on page 157 for a description of RMT.

- L2 cache operation during errors can be configured in the `L2_ModeSetup1[L2_Stop_Data_Derr_En]`, `L2_ModeSetup1[L2_Stop_Ins_Derr_Dis]`, and `L2_ModeSetup1[L2_Stop_TO_Dis]` register bits.
- The L1 and L2 caches can be disabled using the `HID4[dis_force_ci]` and `HID4[enb_force_ci]` register bits.

#### 6.1.5.3 *Replacement Algorithm*

One of three L2 cache-replacement algorithms can be selected by the `L2_ModeSetup1[RMT_Mode]` field:

- *Binary Least-Recently Used (LRU) Mode*—This algorithm is based on the binary tree scheme.
- *Pseudo LRU (p-LRU) Mode*—This algorithm is described in the next paragraph.
- *Direct-Mapped Mode*—This algorithm uses three tag address bits to map an address to one of eight congruence-class members.

The pseudo LRU (p-LRU) mode uses software-configurable address range registers and a replacement management table (*Section 6.3.1.3 Address Range Registers* on page 155 and *Section 6.3 Replacement Management Tables* on page 154) to lock cache ways to a particular Replacement Class Identifier (`RclassID`). The other two modes cannot use these facilities.

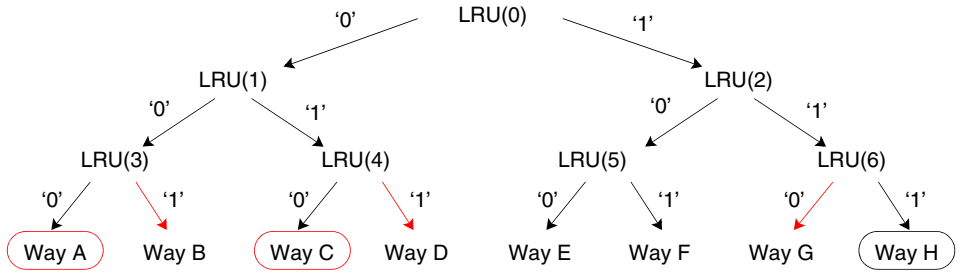
The p-LRU mode supports the following features:

- *Multiple Cache-Line Locking*—The L2 LRU replacement requirement is used to perform cache-line replacement with the normal LRU method and with cache ways locked by the RMT.
- *RclassID from PPU to LRU Controller*—The LRU controller uses the `RclassID` from the PPU to determine which way or ways are locked by the RMT.
- *Cache-Line Replacement Involves RMT in the LRU Controller*—The LRU controller must look up the LRU value for the corresponding cache line and then update the LRU value based on the look-up value from the RMT, as described in *Section 6.3.2 PPE L2 Replacement Management Table* on page 157.
- *Updating LRU Data (Pointer) for Unlocked Cache Lines*—The LRU controller also must update the least-recently-used pointer to the next cache line that is unlocked by the RMT.

*Figure 6-2* on page 144 shows an example of Pseudo LRU replacement. In this figure, if the way replacement value in the RMT is ‘0101 1110’, then:

- Ways A, C, and H are locked from replacement by the value in the RMT.
- The RMT value overrides the following LRU data: LRU(3) = ‘1’, LRU(4) = ‘1’, LRU(6) = ‘0’ (The left arrows correspond to ‘0’ and the right arrows correspond to ‘1’).
- The LRU control must apply pseudo-LRU replacement only on ways B, D, E, F, and G in a miss, or if there is a hit on any unlocked way.
- For a hit on a locked way, the LRU controller does not perform an LRU update, and only the way that is hit is replaced.

*Figure 6-2. Pseudo-LRU Binary Tree for L2-Cache LRU Binary Tree with Locked Ways A, C, and H*




### 6.1.5.4 *Cache Writeback*

The following procedure writes back the entire L2 cache to main storage by means of software. The other thread should be quiesced or guaranteed not to be performing cacheable operations.

1. Set the `L2_ModeSetup1` register mode bits to direct-mapped mode (`L2_ModeSetup1[RMT_Mode]` = '11'). Switching between LRU mode and direct-mapped mode can be done at any time. Data integrity is maintained, although performance might be affected.
2. Pick a 4 MB region of main storage to hit all eight ways in direct-mapped mode.
3. Execute one load, store, or **dcbz** instruction to bytes in each 128-byte cache line in the chosen main-storage region. This fills the L2 with the contents of the chosen region and forces out all other data contained in the cache. However, if the address range of the current contents of the L2 cache is unknown, and the L2 cache needs to be written back to memory as is, then use only a load instruction (not a store or **dcbz**).
4. Execute one **dcbf** instruction for each 128-byte cache line in the specified main-storage region. This writes back and invalidates the specified region. If the contents of the region were resident and modified in the L2, the **dcbf** operation ensures that the data is copied back to main storage. Instruction fetches can cause the L2 cache to retain some valid data even after the **dcbf** operation is performed.

### 6.1.5.5 *Cache Queues*

The 512 KB L2 cache is physically implemented as four quadrants, each quadrant containing 32 bytes of the 128-byte cache line. The cache contains the following queues:

- CIU Load Queue
  - I-demand queue (2 entries)
  - I-prefetch queue (2 entries)
  - Translate queue (1 entry)
  - D-demand queue (8 entries)
  - D-prefetch queue (8 entries)
- CIU store queue: 8 × 16 bytes
- L2 store queue: 8 × 64 bytes, gather enable

- L2 RC reload queue: 6 × 128 bytes
- L2 snoop queue: 4 × 128 bytes
- L2 castout queue: 6 × 128 bytes

A cache-line read from the L2 requires two data beats (two L2 clock cycles, which is four core-clock cycles) and returns 16 bytes of data per beat from each quadrant. The L2 then sends two 64-byte beats to the CIU data-reload bus. A cache-line read request that misses the L2 is sent to the element interconnect bus (EIB). Reload data coming from the EIB is received in eight back-to-back 16-byte data beats. Because all reload data is sent back-to-back, the L2 must wait for all missed data from the EIB before sending data to the PPU. The L2 allows the PPU to request which quadword it wants to receive first from the L2 (called the *critical quadword*). A cache-line write requires two write accesses to the L2 cache, performing a 64-byte write for each access.

#### 6.1.5.6 *Read-and-Claim State Machines*

The L2 cache has six read-and-claim (RC) state machines, with corresponding address and data queues. They manipulate data in and out of the L2 cache in response to PPU or snoop requests. These queues have the following features:

- *Configuration*—The number of RC state machines used for prefetches can be configured in the Prefetch Outstanding Request Limit (`PR`) field of the CIU Mode Setup Register (`CIU_ModeSetup`). When the RC state-machine resource that is servicing data or instruction prefetches reaches a threshold, the data and instruction prefetches are temporarily suspended until the RC state-machine resources become available again. This threshold and other parameters of prefetching can be set in the `CIU_ModeSetup` register. See *Section 6.1.9* on page 150 for more information.
- *Congruence Classes*—Two RC state machines never simultaneously work on the same L2 congruence class. This prevents LRU victim conflicts where both RC state machines have chosen the same victim.
- *Castouts*—If an RC state machine starts on a miss and the victim is valid, the castout (CO) state machine must be available for the RC state machine to process the access. If the CO state machine is busy working on a prior castout, the RC state machine does not acknowledge the access if it is a miss (it will be reissued by the CIU). If it is a miss and the CO state machine is available to process the castout, the RC state machine accepts the access (acknowledge the CIU request). After the RC state machine has finished, it can process a new access that results in a hit, even though the CO state machine is still processing a castout from a prior access.

#### 6.1.5.7 *Store Queue*

The L2 store queue consists of eight fully associative 64-byte sectors. Each sector is byte writable, so gathering occurs at the point the buffer is written. This is effectively a store cache with the following features:

- *Store Gathering Logic*—Gathering is supported for nonatomic, cacheable stores. Gathering of nonstore opcodes is not supported (for example, a request that involves a **dcbz**, PTE update, or atomic stores is not gathered). Gathering is done into any of the store queue entries, to any 64-byte combination within an entry, for any store operation, assuming no barrier operations prevent gathering to a previously stored entry.

- *Store Queue*—A request from the CIU store queue is sent in two consecutive cycles (address, then data) to the L2 and is stored in one of the eight 64-byte store queue entries. Unloading of the eight entries in the store queue is done in a first-in-first-out (FIFO) manner.
- *Store Queue Request by Read-and-Claim (RC) State Machine*—No two RC state machines (*Section 6.1.5.6* on page 145) are simultaneously active to the same L2 congruence class. Therefore, the store queue sometimes does not dispatch until a previous store is completed by an RC state machine, even if another RC state machine is available.

If the store queue detects and sends a full-line store indication with its request, the RC state machines handle back-to-back writes to the same cache line. This completes a cache-line write in 13 cycles (12 cycles for the normal 64-byte write, plus one cycle for the adjacent 64-byte data). Only one directory update is necessary because this is a single cache line.

Store-queue data parity and other errors can be read in the L2 Cache Fault Isolation Register (`L2_FIR`).

#### 6.1.5.8 *Reservations*

The L2 cache supports two reservation bits (one per thread) to handle a load and reserve request by each thread. The snoop logic compares against both reservation addresses, but this is not a thread-based comparison.

#### 6.1.5.9 *L2 Bypass (Inhibit) Configuration*

To bypass the L1 and L2 caches (that is, to inhibit their use), set the `HID4[enb_force_ci]` bit to '1' and clear the `HID1[en_if_cach]` bit to '0'.

### 6.1.6 Instructions for Managing the L1 and L2 Caches

The PPE supports six *PowerPC Architecture* instructions for management of the L1 and L2 caches. All are problem-state (user) instructions. They are described in the sections that follow.

The *PowerPC Architecture* provides cache management instructions that allow user programs to:

- Invalidate an ICache line (**icbi**)
- Hint that a program will probably access a specified DCache line soon (**dcbt**, **dcbtst**)
- Set the contents of a DCache line to zeros (**dcbz**)
- Copy a modified DCache line to main storage (**dcbst**)
- Copy a modified DCache line to main storage and invalidate the copy in the DCache (**dcbf**)

The cache-management instructions are useful in optimizing the use of main-storage bandwidth in such applications as graphics and numerically intensive computing. The functions performed by these instructions depend on the storage-control attributes associated with the specified storage location (see *Section 4.2.6.7 WIMG-Bit Storage Control Attributes* on page 91). Cache management is also needed when a program generates or modifies code that will be executed (that is, when the program modifies data in storage and then attempts to execute the modified data as instructions).

#### 6.1.6.1 *Instruction Cache Block Invalidate*

The L1 ICache line size for the **icbi** instruction is 128 bytes. The L1 ICache does not support a snoop operation. Because of this, the L1 ICache is not necessarily consistent with modifications to those storage locations.

The effective address of the **icbi** instruction is translated as a data address by the load store unit (LSU). Because of this, data-address translation rules are used, and any exception that results is handled as a data storage exception. The address translation is done using the D-ERAT, not the I-ERAT (see *Section 4.2.4 Effective-to-Real-Address Translation* on page 83). Because this instruction invalidates cache lines based on real addresses, it must search and invalidate up to eight cache lines in the L1 ICache to account for the potential aliasing conditions described in *Section 6.1.3.4* on page 138. Two passes are done to invalidate the entries, one to generate a vector containing the lines that will be invalidated, and another pass to do the actual invalidation.

In addition to checking and invalidating lines L1 ICache lines, the **icbi** instruction and the real address of the line to be invalidated are broadcast to all PPEs in the system. This PPE must also accept **icbi** instructions from other PPE processors.

#### 6.1.6.2 *Data Cache Block Touch*

The DCache line size for the **dcbt** and **dcbtst** instructions is 128 bytes. The behavior of **dcbtst** is the same as the behavior of **dcbt** with the Touch Hint (TH) field equal to ‘0000’, as described in the rest of this section.

These instructions act like a load instruction from the viewpoint of the cache hierarchy and the TLB, with the following exceptions:

- If data translation is enabled (`MSR[DR]` = ‘1’), and a segment or page fault results, the instruction is cancelled and no exception occurs.
- If a TLB miss results and `LPCR[TL]` = ‘0’ (hardware table-lookup mode), the TLB is reloaded and the corresponding reference bit is set if a matching PTE is found.
- If the page-protection attributes prohibit access, the page is marked caching-inhibited (I = ‘1’), `LPCR[TL]` = ‘1’ (software table-lookup mode), or the page is marked guarded (G = ‘1’), the instruction is cancelled and does not reload the cache.

The instructions check the state of the L1 data cache. If the data requested is not present, a reload request is sent to the L2 cache. If the data is not present in the L2 cache, then the L2 cache requests the data from memory. Only the L2 cache is reloaded (the data is not forwarded to the L1 cache) unless TH = ‘0000’ in which case the L1 cache is reloaded as well. Similarly, if the data is in the L2 cache already, data is not forwarded to the L1 cache unless TH = ‘0000’ in which case the L1 cache is reloaded. If the **dcbt** or **dcbtst** instruction reloads a cache line, the replacement attributes for the cache (for example, the least-recently used information) are updated.

The PPE implements the optional form of the **dcbt** instruction defined in the *PowerPC Virtual Environment Architecture, Book II*. This form of the instruction can set up a *data stream*. Each data stream is managed independently. There are eight streams available in total. These are split into four streams for each thread. The Touch Hint (TH) field of the instruction has the following

implementation-specific characteristics:

- The PPE implements the TH = ‘0000’, ‘1000’, and ‘1010’ forms.
- All other TH(0) = ‘0’ forms are not supported and are implemented the same as TH = ‘0000’.
- All other TH(0) = ‘1’ forms are not supported and are ignored.
- For the TH = ‘1010’ form, the Transient field (bit 57 of the effective address) is ignored.

The PPE’s data-prefetch engine (DPFE) implements the **dcbt** and **dcbst** instructions. The prefetch engine services all active streams simultaneously in a round-robin fashion. Each time a stream is serviced, a prefetch request is generated for the stream and issued to the L2 cache. One cache line is requested at a time. By default, data prefetch has the lowest priority of any queue in the PPE’s load subunit. This default behavior can be changed. See *Section 6.1.10 Load Subunit Management* on page 150.

The basic function of the **dcbt** and **dcbtst** instructions can be changed in the following ways:

- Force all load instructions to be treated as loading to guarded storage (thus causing data cache touch instructions, such as **dcbt**, to be treated as nop instructions) using the `HID4[force_geq1]` register bit.
- Force the x-form of the **dcbt** and **dcbtst** instructions to be treated as nop instructions using the `HID4[tch_nop]` register bit.

#### 6.1.6.3 *Data Streams*

As described in *Section 6.1.6.2* on page 147, the PPE implements the optional *data stream* form of the **dcbt** and **dcbtst** instructions. To improve prefetch efficiency, the L2 notifies the CIU, *Section 6.1.2.2* on page 136, when it detects that a demand load to the same prefetch address is being serviced by the L2. In this case, the demand load is ahead of the data prefetch and the prefetch is terminated. The CIU does not terminate the stream, but it can then send another load request.

Stopping a stream terminates the stream and relinquishes its data prefetch engine (DPFE) queue entry. The following conditions can terminate a stream:

- When software issues an all stop command (**dcbt** or **dcbtst** instruction TH field = ‘1010’, GO = ‘0’, Stop = ‘11’), all active streams for the issuing thread are terminated.
- When software issues a single stop command (TH = ‘1010’, GO = ‘0’, Stop = ‘10’) to a specific stream, the specified stream is terminated.
- When software issues a **dcbt** instruction with either TH = ‘1000’ or TH = ‘1010’ that has the same stream ID as an already active stream, the active stream is terminated and restarted with the state indicated by the new instruction.
- When a stream set up with TH = ‘1010’ reaches the UNITCNT (number of units in data stream), it is terminated.
- When a stream reaches its page boundary, the stream is terminated. For pages larger than 64 KB, the stream is nonetheless terminated at each 64 KB boundary.
- When a thread enters real-mode operation (`MSR[DR]` = ‘0’), all streams for the thread are terminated.

- When a thread is no longer enabled, all streams for the thread are terminated.
- When a **tlbie** instruction or a snooped TLBIE bus command is issued, all streams for both threads are terminated.

#### 6.1.6.4 *Data Cache Block Zero*

In the CBEA processors, the DCache line size for the **dcbz** instruction is 128 bytes. The function of **dcbz** is performed in the L2 cache. If the line addressed by the **dcbz** is present in the L1 DCache, the line is invalidated before the operation is sent to the L2 cache. The L2 cache gains exclusive access to the line, without actually reading the old data, and performs the zeroing function in a broadside manner. That is, all bits in the cache line are cleared to '0', which means the data cannot be discovered by another process. The **dcbz** instruction is handled like a store and is performed in the appropriate storage-model order. It does not bypass any of the store queues. In general, the performance of **dcbz** is better than a sequence of 8-byte stores, because it zeros 128-bytes at a time.

Regardless of whether the cache line specified by the **dcbz** instruction contains an error (even one that is not correctable with ECC), the contents of the line are cleared to '0' in the L2 cache. If the line addressed by the **dcbz** instruction is in a memory region marked caching-inhibited (I = '1'), or if the L1 DCache is disabled (`HID4[en_dcway]` = '0000'), the instruction causes an alignment interrupt.

#### 6.1.6.5 *Data Cache Block Flush and Data Cache Block Store*

The L1 DCache line size for the **dcbf** and **dcbst** instructions is 128 bytes. A **dcbf** instruction causes all caches for all processor elements in the system to write modified contents back to main storage and invalidate the data in the cache for the line containing the byte addressed by the instruction. A **dcbst** behaves the same, except that the data in the cache remains valid (if present).

The **dcbf** and **dcbst** instructions will probably complete before the operation they cause has completed. A context-synchronizing instruction (see *Section 20 Shared-Storage Synchronization* on page 561) ensures that the effects of these instructions are complete for the processor element issuing the synchronizing instruction.

#### 6.1.6.6 *Vector/SIMD Multimedia Extension Intrinsics*

The following instructions defined in the *AltiVec Technology Programming Interface Manual* are deprecated and are implemented as nops on the CBEA processors:

- vec_dss(a)—Vector Data Stream Stop
- vec_dssall()—Vector Stream Stop All
- vec_dst(a,b,c)—Vector Data Stream Touch
- vec_dstst(a,b,c)—Vector Data Stream Touch for Store
- vec_dststt(a,b,c)—Vector Data Stream Touch for Store Transient
- vec_dstt(a,b,c)—Vector Data Stream Touch Transient

Software should use the PPE the **dcbt** and **dcbtst** instructions for stream control.

### 6.1.7 Effective-to-Real-Address Translation Arrays

The PPE contains two 64-entry, 2-way set-associative, effective-to-real-address translation (ERAT) arrays: an instruction ERAT (I-ERAT) and a data ERAT (D-ERAT). Both are shared by both threads. The ERATs store recently used EA-to-RA translations, and ERAT operation is maintained by hardware, as described in *Section 4.2.4 Effective-to-Real-Address Translation* on page 83.

### 6.1.8 Translation Lookaside Buffer

The PPE's memory management unit (MMU) contains a1024-entry, 4-way set associative TLB cache that holds most-recently used page table entries (PTEs). There are 256 congruence classes (sets), each requiring eight bits for the TLB index. Hardware or software can be used to replace the TLB entry. Privileged PPE software can maintain the TLB based on searches of the PPE's page table, using access to memory-mapped I/O (MMIO) registers in the MMU. For details, see *Section 4.2.7 Translation Lookaside Buffer* on page 93.

### 6.1.9 Instruction-Prefetch Queue Management

The instruction unit (IU) hardware examines each committed instruction fetch and checks to see if the next instruction is in the L1 ICache. If not, it issues a prefetch request.

When the L2 cache is available for a request, and no higher-priority request exists (see *Section 6.1.10*), the instruction prefetch queue forwards the prefetch request to the L2. The L2 loads the line from main storage, if necessary. There is no subsequent reload for this transaction, and data is not loaded into the L1.

When the L2 read-and-claim (RC) state machine (*Section 6.1.5.6* on page 145) resource that is servicing data prefetches or instruction prefetches reaches a threshold, the data and instruction prefetch are temporarily suspended until the RC state machine resources become available again. This threshold and other parameters of prefetching can be set in the `CIU_ModeSetup` register.

### 6.1.10 Load Subunit Management

The load subunit in the CIU (*Figure 6-1* on page 135) employs a two-level arbitration scheme. The first level of arbitration is within each separate subqueue of the load subunit. This arbitration is always from oldest entry to youngest entry. The second level of arbitration enforces a priority among queues. The default priority order is:

- Translate request queue (XLAT)
- Demand data load request queue (DLQ)
- Demand instruction request queue (DFQ)
- Instruction prefetch request queue (IPFQ)
- Data prefetch engine (DPFE)

This second-level priority scheme can be programmed in the `CIU_ModeSetup` register as follows (both settings are independent of the other):

- The priority of DLQ and DFQ can be swapped to change the order to: XLAT, DFQ, DLQ, IPFQ, DPFE.
- To prevent starvation, the DPFE priority can be raised to the second priority: XLAT, DPFE, DLQ, DFQ, IPFQ. Alternatively, the DPFE priority can be periodically toggled from second to fifth priority after 32 load requests are issued to the L2.

## 6.2 SPE Caches

Each SPE's memory flow controller (MFC) has the following caches, both of which can be managed by privileged PPE software:

- Translation lookaside buffer (TLB)
- Atomic (ATO) unit cache, called the *atomic cache*

### 6.2.1 Translation Lookaside Buffer

Each SPE's synergistic memory management (SMM) unit, in the MFC, contains an 256-entry, 4-way, set-associative TLB cache that holds most-recently used PTEs. There are 64 congruence classes (sets), each requiring six bits for the TLB index. The TLB array is parity protected. Parity generation and checking can be disabled in the `SMM_HID` register. Hardware or software can be used to replace the TLB entry. For details, see *Section 4.3.5 Translation Lookaside Buffer* on page 108.

### 6.2.2 Atomic Unit and Cache

Each SPE's MFC contains an atomic unit that handles semaphore operations for the synergistic processor unit (SPU) and provides PTE data to the SMM. More specifically, the atomic unit:

- Provides atomic operations for the **getllar, putllc, putlluc**, and **putqlluc** MFC atomic commands.
- Provides PTEs to the SMM for hardware table lookups and updates to the Reference (R) or Change (C) bits in PTEs (called an *RC update*).
- Maintains cache coherency by supporting snoop operations.

The atomic cache stores six 128-byte cache lines of data. Four of those cache lines support semaphore operations, one cache line supports PTE accesses, and one cache line supports reloading data from cache-miss loads like **getllar**.

MFC atomic commands are issued from the DMA controller (DMAC) to the atomic unit. These requests are executed one at a time. The atomic unit can have up to two outstanding requests: one immediate-form command (**getllar**, **putllc**, or **putlluc**) and one queued-form command (**putqlluc**). The SMM's PTE request can occur at the same time as MFC atomic requests. However, the atomic unit handles this situation by alternating SMM requests and MFC atomic requests.

#### 6.2.2.1 *Atomic Unit Operations*

The atomic unit accepts MFC atomic requests, SMM requests, and snoop requests. The MFC atomic requests are **getllar**, **putllc**, and **putlluc**. The SMM request is for a PPE page-table lookup with RC update. These requests are executed in the atomic unit RC machine. Each of these requests can come at any time. Two MFC atomic requests can be queued in the atomic unit while the first request executes. Only one SMM request can be pending at a time. Snoop requests can come every two cycles, at most.

The get lock line and reserve command, **getllar**, is similar to the *PowerPC Architecture* **lwarx** instruction. A 128-byte cache line of atomic data is loaded into the local storage (LS) and sets a reservation in the atomic unit. Only one reservation is allowed at a time in the atomic unit. In the case of a **getllar** miss with the atomic cache full, room is created for this load by casting out LRU data. Looping on the **getllar** MFC atomic command is not necessary, because software can look for the reservation lost event in the `MFC_RdAtomicStat` channel's lock-line command status.

The put lock line conditional command, **putllc**, is similar to the *PowerPC Architecture* **stwcx** instruction. This is a store command with a condition on reservation active. If the **putllc** data is a cache hit, the atomic unit executes the store by sending store data to the atomic cache from the LS. In addition, it resets the reservation and sends success status to the DMAC, which, in turn, is routed to the SPU.

The put lock-line unconditional command, **putlluc**, is a store command without any condition on the reservation. The atomic unit executes this request as it does **putllc**, except that it does not send any status to the DMAC. Also, the atomic unit executes the store, regardless of the reservation flag.

The put queued lock-line unconditional command, **putqlluc**, is issued to the atomic unit as a **putlluc** MFC atomic request. The atomic unit can have up to two outstanding requests: one immediate form command (***getllar***, ***putllc***, or ***putlluc***) and one queued form command (**putqlluc**).

#### 6.2.2.2 *Lock Acquisition Sequence and its Cache Effects*

Significant performance improvement is realized by using the **putlluc** command to release locks. Using the **putlluc** command rather than a DMA **put** command results in the modification of the data in the atomic cache. There is no need to invalidate the cache and update main storage if the lock-line block is still in the cache.

Because the SPE does not implement an SPU Level 1 (SL1) cache, the **putlluc** command results in a direct store to the referenced real memory and the invalidation of all other caches that contain the affected line. The **getllar**, **putllc**, and **putlluc** commands support high-performance lock acquisition and release between SPEs and the PPE by performing direct cache-to-cache transfers that never leave the CBEA processors.

The possible atomic-cache state transitions based on a typical lock code sequence are:

1. Issue **getllar** command.
2. Read the `MFC_RdAtomicStat` channel to see if **getllar** done.
3. Compare (check for active lock).
4. Branch if not equal (If the lock is still set, go to the step 1).

5. Issue **putllc** (attempt to set lock).
6. Check to see if the ***putllc*** is done.
7. Branch if not equal (If the **put** conditional is not successful, go to step 1).
8. Critical code.
9. Issue ***putlluc*** (release lock).
10. Check if ***putlluc*** is done (for future atomic command issue).

#### 6.2.2.3 *Reservation Lost Events and their Cache Effects*

A reservation can be lost (cleared) in any of the following situations; some of these situations affect the state of the atomic cache:

- One ***getllar*** request (***getllar*** A) of atomic unit 1 sets the reservation. Before the first ***putllc*** request (***putllc*** A) is issued by atomic unit 1, another ***getllar*** request (***getllar*** B) is received in atomic unit 1 to clear the reservation of ***getllar*** A and set a new reservation for a second ***getllar*** request (***getllar*** B).
- A ***putllc*** request is received for a different storage location for which the previous ***getllar*** established a reservation. This scenario occurs only if ***getllar*** A is executed, a reservation is set, and software is interrupted inside the interrupt routine. A ***putllc*** request to a different storage location resets the reservation, but the store is not performed. If there is no interrupt, a scenario in which a ***getllar*** and ***putllc*** pair does not address the same storage location is a software error.
- SPU1 issued ***getllar*** command A and got the reservation. Processor Element 2 wants to perform a store that hits a Shared line in its Processor Element 2's cache. The SPU1 copy is invalidated before the ***putllc*** instruction is issued and performed.

See the *PowerPC Virtual Environment Architecture, Book II* for more details on reservations.

#### 6.2.2.4 *Atomic-Cache Commands*

The *Cell Broadband Engine Architecture* defines an SL1 cache for DMA transfers between LS and main storage. The architecture also defines five MFC commands for the SL1 cache: **sdcrt**, **sdcrtst**, **sdcrz**, **sdcrst**, and **sdcrf**. However, the CBEA processors do not implement an SL1 cache. The **sdcrt** and **sdcrtst** commands are implemented in the SPE as nops, but three of the commands—**sdcrz**, **sdcrst**, and **sdcrf**—affect the SPE atomic cache, as shown in *Table 6-1*:

*Table 6-1. Atomic-Unit Commands*

| CBEA SL1 Command | SPE Atomic-Unit Command | Definition |
|---|---|---|
| **sdcrf** | **sdcrf** | Atomic-cache range write-back and invalidate. |
| **sdcrt** | nop | Undefined. |
| **sdcrtst** | nop | Undefined. |
| **sdcrst** | **sdcrst** | Atomic-cache range store. |
| **sdcrz** | **sdcrz** | Atomic-cache range cleared to '0'. |

#### 6.2.2.5 *Atomic-Cache Flush*

Setting the Flush (F) bit of the MFC_Atomic_Flush register activates an atomic-cache flush that copies cache lines in the Modified state back to main storage. The RC machine executes an atomic-cache flush by dispatching the castout machine six times to copy all six entries of the atomic cache back to main storage.

While the atomic unit performs a cache flush, the SPU must suspend the DMAC so that no DMAC atomic or SMM requests can enter the atomic unit. Otherwise, the atomic unit flags a hardware error called *ATO flush collision*, reported in the MFC Fault Isolation (MFC_FIR) register. See the *Cell Broadband Engine Registers* for details.

## 6.3 Replacement Management Tables

The PPE and SPEs provide a method for supervisor-state software to control L2 and TLB cache replacements, based on a replacement class ID (RclassID). This class ID is used as an index into a replacement management table (RMT), which is used to lock entries in the L2 and TLB caches. The locking function of the RMTs modifies the way in which the L2 and TLB pseudo-LRU algorithms operate.

The PPE has an RMT for managing its TLB, described in *Section 6.3.1*, and for its L2 cache, described in *Section 6.3.2* on page 157. Each SPE also has an RMT for managing its TLB, described in *Section 6.3.3* on page 158.

The RMTs are useful when a small set of pages is frequently accessed by application software and need to be locked in the L2 or TLB or both to prevent misses—for example, in real-time applications (see *Section 11.2.3.2* on page 342 for more detail). They are also useful for preventing streaming data from casting out other data. However, overlocking resources can negatively impact performance.

### 6.3.1 PPE TLB Replacement Management Table

The PPE's TLB RMT is stored in the PPE Translation Lookaside Buffer RMT Register (PPE_TLB_RMT). This RMT can lock translation entries into the TLB by reserving particular ways of the TLB for specific EA ranges in a program. TLB replacement management applies only to the hardware-managed TLB mode, as described in *Section 4.2.7.1 Enabling Hardware or Software TLB Management* on page 93 and *Section 4.2.7.3 TLB Replacement Policy* on page 95. In the software-managed TLB mode, software controls TLB-entry replacement.

#### 6.3.1.1 *RClassID Generation*

In the PPE, a 3-bit replacement class ID (RclassID) is generated from the effective address specified by instruction fetches or by data loads or stores. Instruction fetches include fetches for sequential execution, speculative and nonspeculative branch targets, prefetches, and interrupts.

The Address Range Registers, described in *Section 6.3.1.3* on page 155, are used together with the RclassIDs defined in the following registers (two pairs, one for each PPE thread) to identify addresses covered by the RMT:

- Instruction Class ID Register 0 (ICIDR0)
- Instruction Class ID Register 1 (ICIDR1)
- Data Class ID Register 0 (DCIDR0)
- Data Class ID Register 1 (DCIDR1)

Cache management instructions (*Section 6.1.6* on page 146) are treated like load or store instructions with respect to RClassID.

#### 6.3.1.2 *Index and Table Structure*

The PPE's TLB RMT is implemented by the PPE_TLB_RMT register, and the 3-bit RClassID is used as an index into this 8-entry RMT. For each RclassID, the PPE_TLB_RMT defines which ways of the TLB are eligible to be replaced when a TLB miss occurs. The high-order 32 bits of this register are not implemented and are reserved for future use. The low-order 32 bits are divided into eight 4-bit RMT entries, corresponding to eight RclassIDs.

Each of the eight 4-bit RMT entries has a one-to-one correspondence to the four ways in the TLB. For each RMT-entry bit with a value of '1', the table-lookup algorithm can replace an entry in that way if a TLB miss occurs in the EA range that maps to the RclassID for that RMT entry. If multiple ways are indicated, they are replaced in priority order beginning with invalid entries first, and then valid entries according to the pseudo-LRU policy described in *Section 4.3.5.2 Hardware TLB Replacement* on page 110.

For example, if a translation with an RclassID of 5 is requested, that class corresponds to the RMT5 field of the PPE_TLB_RMT register, which has a value of '1001'. If a TLB miss occurs, TLB ways 0 and 3 are valid candidates for replacement (pseudo-LRU initially chooses way 0 and then points to way 3 when another miss occurs). If all bits in an RMT field are written by software with zero ('0000'), then the hardware treats this RMT field as if it had been set to '1111', therefore allowing replacement to any way. (This corresponds to the method used for SPE TLB replacement described in *Section 6.3.3* on page 158.)

The RMT0 field of the PPE_TLB_RMT register is the default entry and should be used by software to specify the replacement policy for any EA range that is not mapped by the Address Range Registers.

#### 6.3.1.3 *Address Range Registers*

The PPE has a set of Address Range Registers for supporting the pseudo-LRU TLB replacement algorithm (the registers are also used for L2 replacement). An address range is a naturally-aligned range that is a power-of-2 size between 4 KB and 4 GB, inclusive. An address range is defined by two types of registers; a Range Start Register (RSR) and a Range Mask Register (RMR). For each address range, there is an associated ClassID Register (CIDR) that specifies the RclassID. The names RSR, RMR, and CIDR are generic labels for groups of registers that are differentiated according to instructions or data (I or D prefix).

As shown in *Table 6-2* on page 156, for each PPE thread there are two sets of RSRs, RMRs, and CIDRs for load/store data accesses and two sets for instruction fetches. The Address Range Registers are accessible by the **mtspr** and **mfspr** privileged instructions.

*Table 6-2. Address Range Registers (One per PPE Thread)*

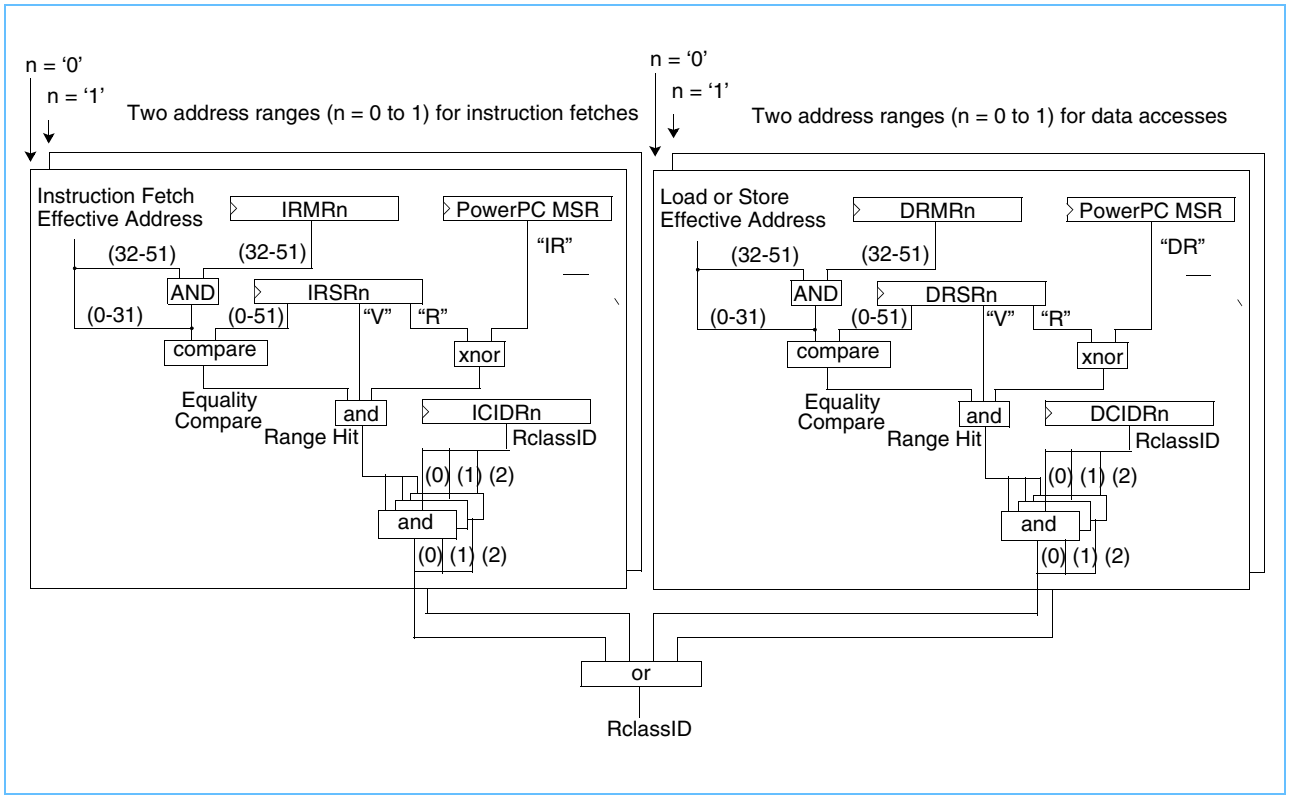
| Type of Operation | Address Range | | Replacement ClassID (RclassID) |
|---|---|---|---|
| | Range Start Registers | Range Mask Registers | |
| Data Access | Data Address Range Start Register 0 (DRSR0) | Data Address Mask Register 0 (DRMR0) | Data Class ID Register 0 (DCIDR0) |
| | Data Address Range Start Register 1 (DRSR1) | Data Address Mask Register 1 (DRMR1) | Data Class ID Register 1 (DCIDR1) |
| Instruction Fetch | Instruction Range Start Register 0 (IRSR0) | Instruction Range Mask Register 0 (IRMR0) | Instruction Class ID Register 0 (ICIDR0) |
| | Instruction Range Start Register 1 (IRSR1) | Instruction Range Mask Register 1 (IRMR1) | Instruction Class ID Register 1 (ICIDR1) |

If all the following conditions are met, the particular address range defined by an RSR and RMR pair applies to a given EA and a *range hit* is said to have occurred:

- RSR[63] = '1'
- RSR[0:51] = EA[0:31] || (EA[32:51] & RMR[32:51])
  - If the operation is a load or store, then
    - RSR[62] = MSR[DR] and
    - The RSR and RMR pair used in these conditions must be DRSR0 and DRMR0 or DRSR1 and DRMR1
  - else (the operation is an instruction fetch)
    - RSR[62] = MSR[IR] and
    - The RSR and RMR pair used in these conditions must be IRSR0 and IRMR0 or IRSR1 and IRMR1

If there is no range hit for a given EA, the RclassID has a value of '0'. In effect, RMR defines the size of the range by selecting the bits of an EA used to compare with the RSR. The upper bits of an RSR contain the starting address of the range and the lower bits contain a relocation mode (real or virtual) and a valid bit. The size of the range must be a power-of-2. The starting address of the range must be a range-size boundary.

*Figure 6-3. Generation of RclassID from the Address Range Registers for Each PPE Thread*




To avoid confusion about the `RclassID` value, it is recommended that software ensure that data-address ranges 0 and 1 do not overlap such that both have a simultaneous range hit. Likewise, it is recommended that software ensure that instruction-address ranges 0 and 1 do not overlap such that both have a simultaneous range hit.

### 6.3.2 PPE L2 Replacement Management Table

The PPE supports an 8-entry RMT for the 8-way set associative L2 cache. It is implemented by the L2 RMT Setup Register (`L2_RMT_Data`). The L2 RMT uses the same Address Range Registers that are used by the TLB RMT (see *Section 6.3.1.3* on page 155). The L2 pseudo-LRU replacement algorithm, for which these registers are used, is described in *Section 6.1.5.3 Replacement Algorithm* on page 143. For each `RclassID`, the L2 RMT defines which ways of the L2 cache are eligible to be replaced when an L2-cache miss occurs.

The L2-cache replacement policy is controlled by a replacement management table analogous to that used for the PPE TLB (*Section 6.3.1 PPE TLB Replacement Management Table* on page 154) when the L2 cache is not in direct-mapped mode (see *Section 6.1.5.3 Replacement Algorithm* on page 143). The RMT and LRU functions of the L2 cache can be configured in the L2 Mode Setup Register 1 (`L2_ModeSetup1`).

The workings of the L2 RMT are analogous to those of the TLB RMT, described in *Section 6.3.1.2 Index and Table Structure* on page 155, except that the L2 cache has eight ways instead of four ways. The 3-bit `RclassID` is used as an index into the `L2_RMT_Data` register to select one of the RMT entries. Each entry of the RMT contains a replacement-enable bit for each of the eight ways in the L2 cache. The `RclassID` selects one of the eight fields in this register.

Each of the eight bits of an RMT entry has a one-to-one correspondence to the eight ways in the L2 cache. For each RMT-entry bit with a value of '1', the L2 replacement algorithm can replace an entry in that way if a miss occurs in the EA range that maps to the `RclassID` for that RMT entry. A bit value of '0' indicates that the entry is locked and cannot be replaced. If multiple ways are indicated, they are replaced in priority order beginning with invalid entries first, and then valid entries. If all bits in an RMT field are zero ('0000 0000'), then the hardware treats this RMT field as if it had been set to '1111 1111', therefore allowing replacement to any way. The replacement method is analogous to that used by the pseudo-LRU policy for TLB replacement, described in *Section 4.3.5.2 Hardware TLB Replacement* on page 110.

The L2 RMT can be used by software to achieve various results, including:

- Lock an address range into the cache so that an access to a cache line always gets an L2 cache hit. The address range is *locked* into the L2 cache by first making it valid in the L2 cache and then configuring the L2 RMT to prevent the cache line from being displaced from the L2 cache.
- Limit an address ranges to one or more ways of the L2 cache without locking these locations and without prohibiting accesses for other data from replacing those ways. This is useful for some kinds of accesses over large data structures to prevent such accesses from flushing a large portion of the L2 cache.
- Allow an application to only use a limited number of ways of the L2 cache while reserving other ways for other applications. This can be useful to ensure a certain performance level for applications, and that performance is dependent on the L2-miss rate. Limiting an application to a subset of the L2 cache prevents the application from displacing all data cached by another application. Such data is not necessarily locked into the L2 cache, but can simply be the latest set of data used by the other application.

### 6.3.3 SPE TLB Replacement Management Table

The synergistic memory management (SMM) unit of each SPE supports a 4-entry RMT that can be used to lock TLB entries and prevent their replacement. The SPE's TLB RMT is analogous to that of the PPE's, described in *Section 6.3.1 PPE TLB Replacement Management Table* on page 154.

The SPE RMT is enabled with the `SMM_HID[RMT_Dsbl]` bit. The RMT entries are stored in the MFC TLB Replacement Management Table Data Register (`MFC_TLB_RMT_Data`)[5]. A 2-bit `RClassID` is provided as a parameter in MFC commands. RMT-entry selection is specified by the `RClassID` given to the SMM by the DMAC during a translation request. The `RClassID` is defined in the MFC Class ID or Command Opcode (`MFC_ClassID_Cmd`) register; unlike the PPE RMTs, the SPE RMT does not use the system of Address Range Registers described in *Section 6.3.1.3 Address Range Registers* on page 155.

5. The Cell/B.E. and PowerXCell processors do not implement the RMT Data Register (RMT_Data) defined in the *Cell Broadband Engine Architecture*. Instead, the SMM uses the implementation-dependent MFC TLB Replacement Management Table Data Register (MFC_TLB_RMT_Data).

Each 4-bit RMT entry represents one of the four TLB ways in a congruence class (see *Figure 4-6 SMM Virtual- to Real-Address Mapping* on page 109 for an illustration of the TLB structure). A bit value of ‘1’ indicates that the corresponding TLB way can be replaced. A bit value of ‘0’ indicates that the way is locked and cannot be replaced. If an entire RMT entry is enabled (‘1111’), then all TLB ways in the congruence class represented by the RMT entry are available for replacement, and the LRU algorithm chooses the replacement set, as described in *Section  TLB Replacement Policy* on page 111. If all bits in an RMT field are written by software with zero (‘0000’), then the hardware treats this RMT field as if it had been set to ‘1111’, therefore allowing replacement to any way. This corresponds to the method used for PPE TLB replacement described in *Section 6.3.1.2* on page 155.

## 6.4 I/O Address-Translation Caches

For a description of the caches used in I/O address-translation, and their management, see *Section 7.4 I/O Address Translation* on page 176 and *Section 7.6 I/O Address-Translation Caches* on page 181.

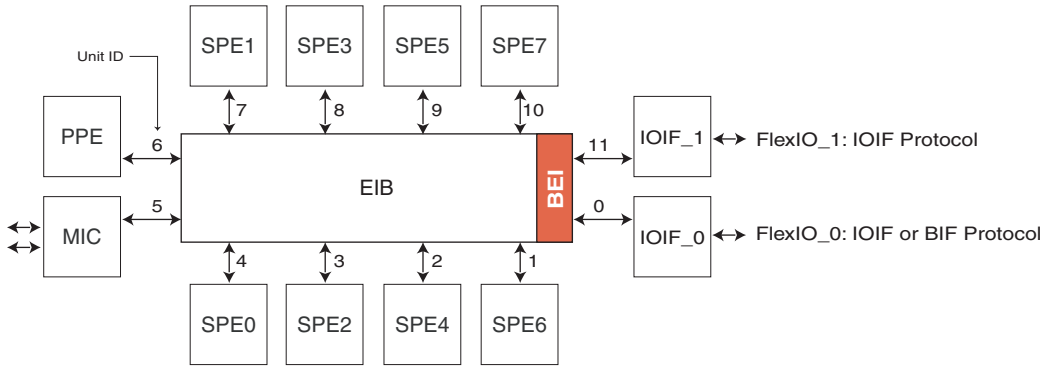
# 7. I/O Architecture

## 7.1 Overview

### 7.1.1 I/O Interfaces

The Cell Broadband Engine interface (BEI) unit, shown in *Figure 7-1*, manages data transfers between the processor elements on the element interconnect bus (EIB) and I/O devices. The BEI supports two Rambus FlexIO I/O bus interfaces, FlexIO_0 and FlexIO_1. FlexIO_0 is also called IOIF0; FlexIO_1 is also called IOIF1.

*Figure 7-1. BEI and its Two I/O Interfaces*




| | | | |
|---|---|---|---|
| BEI | Cell Broadband Engine Interface Unit | IOIF | I/O Interface, or I/O Interface Protocol |
| BIF | Cell Broadband Engine Interface Protocol | MIC | Memory Interface Controller |
| EIB | Element Interconnect Bus | PPE | PowerPC Processor Element |
| FlexIO | Rambus FlexIO Bus | SPE | Synergistic Processor Element |

The input and output bandwidth (number of transmitters and receivers) of each I/O interface (IOIF) is software-configurable during the Cell Broadband Engine Architecture (CBEA) processor[1] power-on reset (POR) sequence, (For a description of the sequence, see the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.) The FlexIO_0 interface can be configured at higher input and output bandwidths than the FlexIO_1 interface. The FlexIO_0 interface can also be configured with either the noncoherent IOIF protocol or the fully coherent Cell Broadband Engine interface (BIF) protocol, which is the EIB's internal protocol. The FlexIO_1 interface supports only the noncoherent IOIF protocol. Each IOIF-protocol interface has four virtual channels.

Both the IOIF and BIF protocols provide fully pipelined, packet-transaction interconnection using credit-based flow control. The noncoherent IOIF protocol can interface to I/O devices, noncoherent memory subsystems, switches, or bridges. The coherent BIF protocol, which supports memory coherency and data synchronization, can be connected to another CBEA processor or coherent memory subsystem, as well as a compatible switch or bridge.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

The POR configuration bits are not visible to software after POR—except in system designs in which the external system controller chip makes them visible—but several related parameters of the configuration can be modified by software after initialization[2].

The BEI has an I/O interface controller (IOC) that supports the two IOIFs with I/O-command processing, I/O-address translation, I/O-interrupt control, programmable memory-mapped I/O (MMIO) registers, and, in general, implements the I/O facilities. Both IOIFs use the same bus address space—a 42-bit real-address space. The BEI also adds a paging mechanism for addresses from I/O devices that is independent of the standard *PowerPC Architecture* memory-paging mechanism. The PowerPC Processor Element (PPE) controls the system-management resources used for both memory paging and I/O paging. It is possible to extend a current kernel's I/O architecture to support this addition of I/O paging without having to re-architect the kernel.

Use of the IOIFs can be controlled using tokens that are allocated by the resource allocation manager (RAM), as described in *Section 8 Resource Allocation Management* on page 203.

### 7.1.2 System Configurations

*Figure 7-2* on page 163 shows two possible system configurations. At the top is a system with a single CBEA processor that connects to coherent I/O devices and noncoherent I/O devices. At the bottom is a system with two CBEA processors that connect together, forming a coherent symmetric multiprocessor (SMP) system, with each CBEA processor also connecting to noncoherent I/O devices.

*Figure 7-3* on page 164 shows a 4-way configuration, which requires a BIF-protocol switch to connect the four CBEA processors. This figure does not necessarily represent future system-product plans.

2. The coherency protocol for the FlexIO_0 interface is software-selectable only at power-on reset (POR). However, if an I/O interface is known to use the IOIF protocol, an I/O device on that interface can access memory pages coherently by setting the M bit in an I/O page table (IOPT) entry.

*Figure 7-2. Single and Dual CBEA Processor System Configurations*

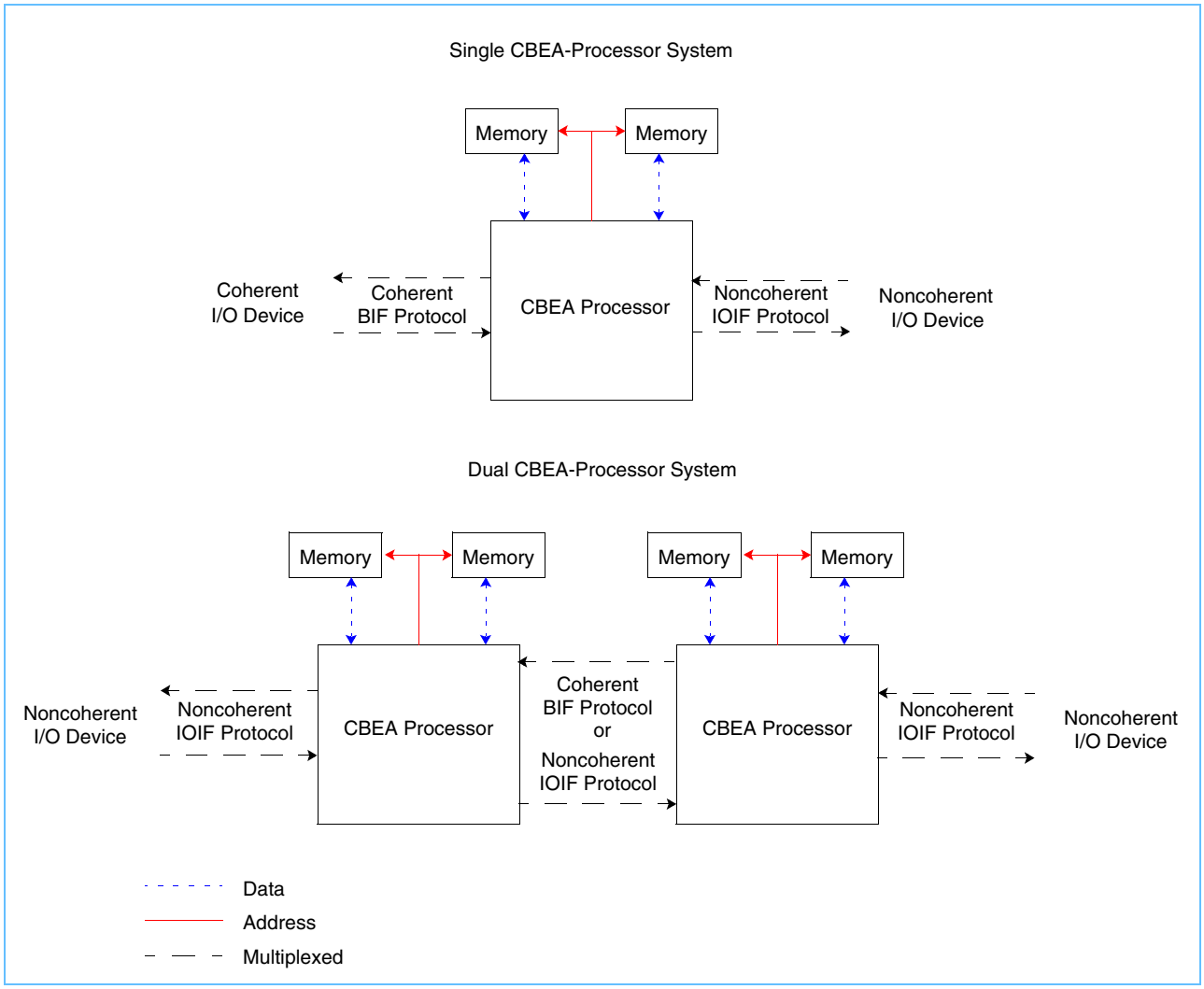
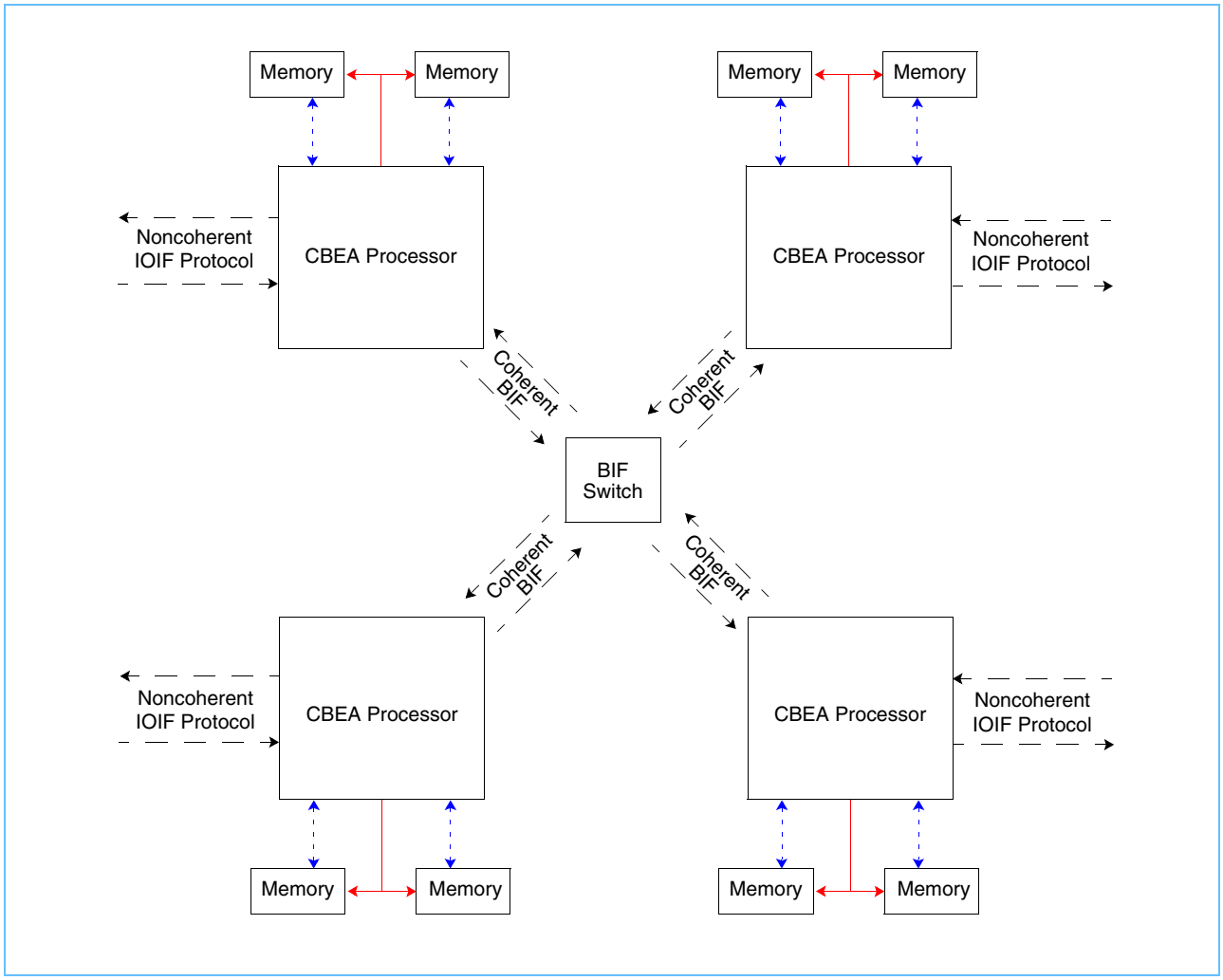


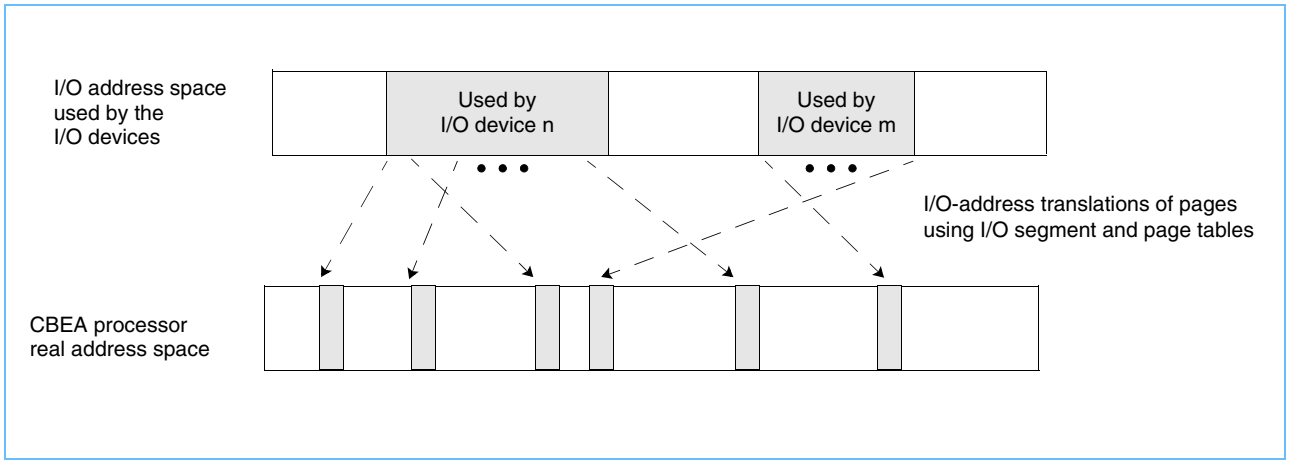
*Figure 7-3. Quad CBEA Processor System Configuration*




### 7.1.3 I/O Addressing

The CBEA processors support I/O-address translation on interfaces configured for the IOIF protocol (but not for the BIF protocol). An I/O address is an address passed from an I/O device to a CBEA processor when the I/O device attempts to perform an access to the CBEA processor real-address space—for example, to an MMIO register, to the local storage (LS) of a Synergistic Processor Element (SPE), or to the memory attached to a CBEA processor. For I/O accesses coming into the CBEA processors on an I/O interface, the CBEA processors provide two modes for handling the incoming address. If I/O-address translation is disabled, the I/O address is used as a real address. If I/O-address translation is enabled, the address is converted from an I/O address into a real address, as shown in *Figure 7-4* on page 165.

*Figure 7-4. Use of I/O Addresses by an I/O Device*




I/O operations often involve multiple, noncontiguous real pages. The I/O address-translation mechanism supports the use of a contiguous I/O address range assigned to noncontiguous real pages. Each I/O device can be assigned a separate range of I/O addresses, either statically or dynamically.

I/O devices can only access pages that have been set up in an I/O page table with their I/O identifier (IOID). This identifier must be passed on the IOIF bus to the CBEA processor, along with the I/O address, during the I/O access. If the identifier from the I/O interface does not match the value in the CBEA processor's I/O page table entry corresponding to the I/O address, an I/O exception occurs. Thus, a form of storage protection is enforced. Pages can also be protected according to their read/write attributes, and pages can be assigned memory-coherence and storage-order attributes.

In some systems, such rigid protection between all I/O devices might not be required. For example, if all I/O devices on a certain I/O bus in the I/O subsystem are only used by one logical partition, it might be satisfactory for the I/O subsystem to group these I/O devices into one identifier.

## 7.2 Data and Access Types

### 7.2.1 Data Lengths and Alignments

The following data lengths and alignments are supported for read and write transfers across an IOIF:

- Read and write transfers of 1, 2, 4, 8, and 16 bytes that are naturally aligned. (An I/O device, but not the CBEA processors, can also perform read and write transfers of 3, 5, 7, 9, 10, 11, 12, 13, 14, and 15 bytes, if they do not cross a 16-byte boundary.)
- Read and write transfers of $q$-times-16 bytes, where $q$ is an integer from 1 to 8, that are quadword-aligned and do not cross a 128-byte boundary.

SPEs and I/O devices can perform DMA transfers of data lengths and alignments that differ from the above by splitting the accesses into multiple transfers across the IOIF.

### 7.2.2 Atomic Accesses

Atomic IOIF accesses are supported for naturally aligned 1, 2, 4, and 16-bytes transfers. These read and write transfers are fully serialized and complete before the next transfer is initiated. For example, suppose a naturally aligned 2-byte memory location initially contains x'0000'. If an I/O device writes x'1111' to this location, a PPE or SPE doing a halfword read of this location will load either x'0000' or x'1111'. It will never load x'0011' or x'1100'.

## 7.3 Registers and Data Structures

The IOC supports several registers and data structures used for I/O-address translation (*Section 7.4 I/O Address Translation* on page 176) and related purposes. The bit fields and values for these registers are described in the *Cell Broadband Engine Registers* specification, although in some cases the descriptions in this section have additional detail.

### 7.3.1 IOCmd Configuration Register

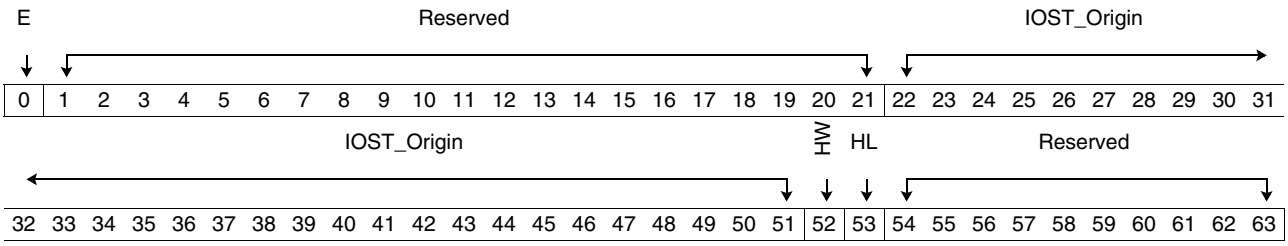
The IOCmd Configuration Register (`IOC_IOCmd_Cfg`) configures basic settings for the IOC. These settings include the IOID for each I/O interface, I/O address-translation enable, token parameters, timeout, read intervention, and others.

### 7.3.2 I/O Segment Table Origin Register

The 64-bit I/O Segment Table Origin Register (`IOC_IOST_Origin`) enables I/O-address translation, specifies the origin and size of the I/O segment table (IOST), specifies the hint-or-lock function of the Hint bit in IOST and I/O page table (IOPT) entries, and enables hardware cache-miss handling for the IOST and IOPT caches.

The `IOC_IOST_Origin` register should be initially configured at boot time, before there are any I/O accesses, and it can be modified after boot. If a hypervisor is implemented (see *Section 11* on page 331), software typically maps the `IOC_IOST_Origin` register using a page that is only accessible to the hypervisor or firmware.

#### 7.3.2.1 *Register Layout and Fields*



| Bits | Field Name | Description |
|---|---|---|
| 0 | E | I/O-Address Translation Enable:<br>0 Disabled (default). I/O addresses are treated as real addresses, and the I/O-address translation unit is disabled to save power. The only IOC address translation MMIO register that can be accessed is this register.<br>1 Enabled. I/O addresses are translated using the IOST and IOPTs.<br>The IOC_IOCmd_Cfg[TE] bit must be set to the same value as this E bit. |
| 1:21 | Reserved | Bits are not implemented; all bits read back as '0'. |
| 22:51 | IOST Origin | The real page number (RPN) of the IOST. The IOST must be integrally aligned and must be aligned on a 4 KB page. These are the 30 high-order bits of the 42-bit real address of the IOST. |
| 52 | HW | Hardware Cache-Miss Handling Enable for IOST and IOPT caches:<br>0 Hardware miss handling is disabled.<br>1 Hardware miss handling is enabled. |
| 53 | HL | Hint or Lock for IOST and IOPT caches:<br>0 The Hint bit for an IOST cache entry or IOPT cache entry is treated as just a hint for valid entries. If there are no invalid entries in the applicable congruence class, hardware miss handling should replace an entry whose Hint bit is '0' instead of an entry whose Hint bit is '1'. If there are no invalid entries in the congruence class and all entries in the congruence class have Hint bits equal to '1', hardware miss handling will replace a cache entry whose Hint bit is '1'.<br>1 The Hint bit is treated as a lock. If hardware miss handling is enabled and an IOST or IOPT cache miss occurs, hardware must not replace a cache entry whose Hint bit is '1', even if the valid (V) bit in the entry is '0'. |
| 54:57 | Reserved | |
| 58:63 | IOST Size | This is the base-2 logarithm of the number of entries in the IOST.<br>Both the first implementation of the CBEA, the Cell/B.E. processor, and the PowerXCell 8i processor treat this field as a fixed value of 7, so that the IOST has 128 entries.<br>Any value set in the IOST Size field is ignored. When this field is read, it returns zero. |

#### 7.3.2.2 ***Register Functions***

*I/O-Address Translation*

The `IOC_IOST_Origin` register's E bit enables I/O-address translation on interfaces configured for the IOIF (but not BIF) protocol. It determines whether addresses generated by I/O units will be used as real addresses (untranslated) or whether the translation mechanism will produce the real address based on the IOST and IOPT. The primary advantages of enabling I/O-address translation include:

- Identity-matching of the IOID and the real address in the IOPT
- Read/write protection of the main-storage space during accesses by I/O units
- Memory coherence and storage ordering of I/O accesses
- Hardware control of misses in the IOST and IOPT caches
- Sizing of the IOST and its pages

Unlike PowerPC virtual-memory address translation, the use of I/O-address translation does not include access to a larger address space; in fact, the translated address space is smaller ($2^{35}$ bytes) than the untranslated real-address space ($2^{42}$ bytes).

*IOST Origin*

The IOST_Origin field is architected to contain up to 50 high-order bits of the 62-bit real address of the IOST. However, the CBEA processors support 42-bit real addresses, and only the right-most (low-order) 30 bits (bits 22:51) of the IOST_Origin field are used; the upper 20 bits of the field are reserved and treated as zeros.

*Hardware Cache-Miss Handling*

The hardware cache-miss (HW) bit enables or disables hardware handling of misses in the IOST cache (*Section 7.6.1* on page 181) and IOPT cache (*Section 7.6.2* on page 183). When HW = '1', the Hint or Lock (HL) bit indicates to hardware whether IOPT entries with their Hint (H) bit = '1' (*Section 7.3.4* on page 171) can be replaced.

If HL = '1', hardware miss handling avoids loading an IOPT-cache entry whose H = '1' in the IOPT. In this case, hardware will select an entry using the following criteria order if there is a cache miss:

1. An entry in the applicable congruence class (cache index) whose valid (V) bit in the IOPT and Hint (H) bit are '0'.
2. An entry in the applicable congruence class whose Hint (H) bit is '0'. A pseudo-LRU algorithm is used in conjunction with this criterion to select an entry among those that meet the criteria.

If no entry is found based on the preceding criteria, and if a cache miss occurs in that congruence class, an I/O segment fault (*Section 9.6.4.4* on page 274) occurs, instead of an I/O page fault, even if the entry exists in the IOST. Likewise, if all the H bits in the applicable congruence class of an IOPT cache are '1', and if a cache miss occurs in that congruence class, an I/O page fault (*Section 9.6.4.4* on page 274) occurs even if the entry exists in the IOPT.

If HL = '0', hardware miss handling will select an entry using the following criteria order:

1. An entry in the applicable congruence class whose V bit is '0'.
2. An entry in the applicable congruence class whose H bit is '0'.
3. An entry in the applicable congruence class whose H bit is '1'.

A pseudo-LRU algorithm is used in conjunction with criteria 2 and 3 to select an entry among those that meet the criteria.

See *Section 7.6.1.3* on page 183 for further details about hardware and software cache-miss handling.

*IOST Size*

The IOST Size field in the CBEA processor is treated by the hardware as if the value is always 7. Thus, the IOST has 128 entries. The corresponding bits in the `IOC_IOST_Origin` are not implemented.

#### 7.3.2.3 *Register Modifications*

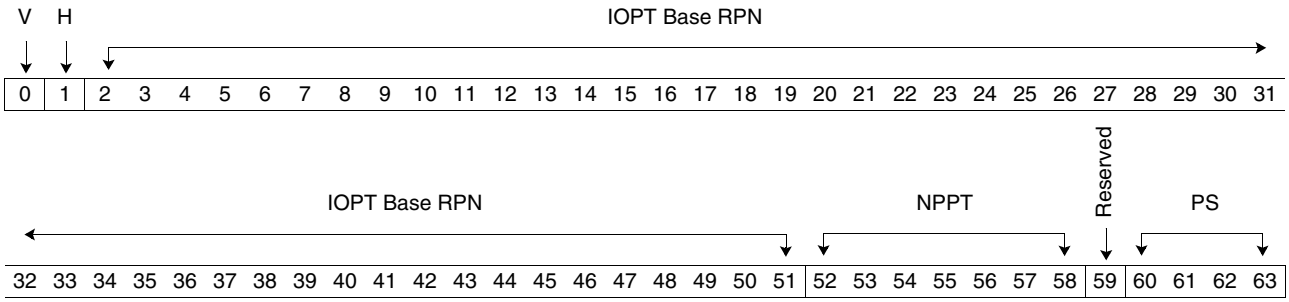
Software must ensure that there are no I/O accesses occurring when it modifies the `IOC_IOST_Origin` register. Also, the `IOC_IOCmd_Cfg[TE]` and `IOC_IOST_Origin[E]` bits must both be set to the same value.

To guarantee that the new value of a store to `IOC_IOST_Origin` register is used on future I/O-address translations, software should load from the `IOC_IOST_Origin` register and ensure that the load has completed before a subsequent store that might cause the new `IOC_IOST_Origin` value to be used. A synchronization instruction or command (*Section 20* on page 561) that orders a load from storage that is both caching-inhibited and guarded versus the subsequent store can be used for this purpose.

### 7.3.3 I/O Segment Table

The IOST is a variable-sized data structure in memory that defines the first phase of the translation of an I/O address to a real address. For each I/O segment, the table contains an entry that points to a page table that is used to complete the translation. There are $2^7$ segments, each one containing $2^{28}$ bytes, for a total translated address space of $2^{35}$ bytes.

#### 7.3.3.1 *Table-Entry Layout and Fields*




| Bits | Field Name | Description |
|---|---|---|
| 0 | V | Valid<br>0 The IOST entry is not valid.<br>1 The IOST entry is valid. |
| 1 | H | Hint<br>1 *Not valid in the CBEA processors.* (Architecturally, if IOC_IOST_Origin[HL] = '0', this is a hint to the IOST cache to retain this entry if possible. If IOC_IOST_Origin[HL] = '1', this entry is never displaced from an IOST cache by hardware when handling an IOST cache miss.)<br>0 *The IOST Hint bit is treated as a '0' in the CBEA processors.* No cache hint. |
| 2:51 | IOPT Base RPN | IOPT-base real page number<br>This is the RPN of the first entry in the IOPT for this I/O segment. Only the right-most 30 bits of the IOPT Base RPN are implemented. The upper 20 bits of the IOPT Base RPN are treated as zeros. |
| 52:58 | NPPT | Number of 4 KB pages (minus one) in the page table for this I/O segment.<br>For IOPT format 1 (*Section 7.3.4* on page 171), the number of IOPT entries for the segment is equal to 512 times the sum of the NPPT plus 1. The IOPT entries are for the lowest-numbered pages in the segment. (IOPT format 2 is not supported in the CBEA processors). |
| 59 | Reserved | |

| Bits | Field Name | Description |
|---|---|---|
| 60:63 | PS | Page size of the pages in this I/O segment. The page size is $4^{PS}$ KB. PS values are:<br>• 0001 = 4 KB page size<br>• 0011 = 64 KB page size<br>• 0101 = 1 MB page size<br>• 0111 = 16 MB page size<br>All four page sizes can be used, regardless of the page sizes used by the PPE and SPEs. |

#### 7.3.3.2 *Table-Entry Functions*

IOST entries specify the properties of a segment and its page tables, including the validity of the IOST entry, the RPN of the first page table in this I/O segment, the number of pages in a page table, and the page size.

The IOST can contain up to 128 64-bit entries. The IOST must start on a 4 KB boundary, at a minimum, and it must be naturally aligned.

*Validity*

Software can control access to pages in a segment by writing the IOST entry's valid (V) bit. If an I/O translation attempts to use an IOST entry that has V = '0', an I/O segment-fault type of I/O exception occurs.

*Cache Hint*

In the CBEA processors, the IOST Hint (H) bit is always treated as a '0'. Thus, hints or locks for the IOST cache (*Section 7.6.1* on page 181) are not supported, irrespective of the value of the `IOC_IOST_Origin[HL]` bit.

*Page Size*

The IOST page-size (PS) field defines the size of the pages in the I/O segment. Four page sizes are supported—4 KB, 64 KB, 1 MB, and 16 MB. However, because it is possible to disable I/O-address translation and use an external I/O address-translation mechanism, the page sizes supported by an external I/O-address translation mechanism might differ from those supported by the PPE and SPEs.

#### 7.3.3.3 *Table-Entry Modifications*

An IOST entry should only be modified when there is no simultaneous I/O access that uses the IOST entry; if this condition is not met, it is unpredictable whether the I/O access will use the old or new value of the IOST entry. If software chooses to stop simultaneous I/O accesses, such accesses must be stopped at the I/O unit. To prevent an I/O-address-translation access of the IOST from getting part of an old value and part of a new value when there is a simultaneous I/O access using the entry, software should take either of the following actions:

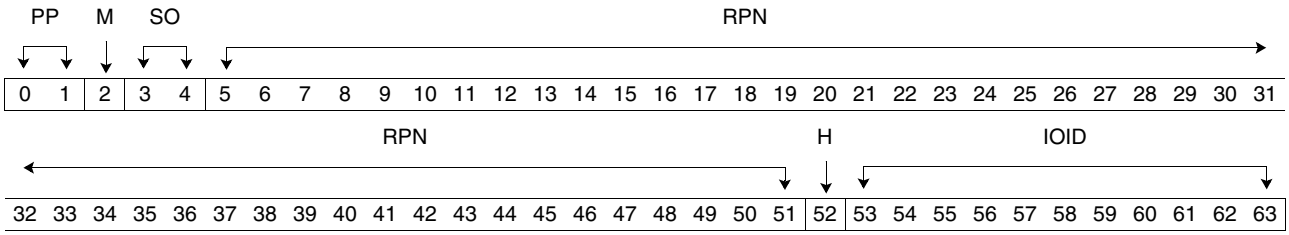
- Use a doubleword atomic store to the IOST entry. According to *PowerPC Architecture*, this is a single-register doubleword store that is naturally aligned.
- Use an atomic store to set the IOST valid (V) bit to '0', execute an **eieio**, **lwsync**, or **sync** instruction, and then store to other bytes in the IOST entry. After making the updates, exe-

cute an **eieio**, **lwsync**, or **sync** instruction, and then execute an atomic store to set the IOST V bit to '1'. The **lwsync** instruction can only be used for this purpose if the IOST is in storage that is not caching-inhibited. The **eieio** and **sync** instructions are unnecessary if the IOST is in storage that is both caching-inhibited and guarded.

### 7.3.4 I/O Page Table

The IOPT is a variable-sized data structure in memory. This table defines the second phase of the translation of an I/O address to a real address. For each I/O page, the table can contain an entry that specifies the corresponding real page number. The IOPT must start on a 4 KB boundary.

#### 7.3.4.1 *Table-Entry Layout and Fields (Format 1)*




| Bits | Field Name | Description |
|---|---|---|
| 0:1 | PP | Page Protection<br>00 No access. (This can also be used to designate an invalid entry.)<br>01 Read access only.<br>10 Write access only.<br>11 Read and write access. |
| 2 | M | Coherence Required<br>0 Memory coherence not required.<br>1 Memory coherence required. |
| 3:4 | SO | Storage Ordering for accesses performed on one IOIF with the same IOID and virtual channel.<br>00 Previous reads and writes are not necessarily performed before a write.<br>01 Reserved.<br>10 Previous writes are performed before a write.<br>11 Previous reads and writes are performed before a read or write. |
| 5:51 | RPN | Real Page Number of the translated I/O address:<br>• For 64 KB page sizes, the 4 least-significant bits of the RPN must be cleared to '0'.<br>• For 1 MB page sizes, the 8 least-significant bits of the RPN must be cleared to '0'.<br>• For 16 MB page sizes, the 12 least-significant bits of the RPN must be cleared to '0'.<br>The real address supported has only 42 bits; therefore, bits[0:16] of the RPN field must be zeros in IOPT entry format 1. |
| 52 | H | Hint<br>1 If IOC_IOST_Origin[HL] = '0', this is a hint to the IOPT cache to retain this entry if possible. If IOC_IOST_Origin[HL] = '1', this entry is never displaced from an IOPT cache by hardware when handling an IOPT cache miss.<br>0 No cache hint. |
| 53:63 | IOID | I/O Device Identifier. Only the I/O device specified can have its I/O address translated using this IOPT entry. |

#### 7.3.4.2 *Table-Entry Functions*

The entries in this table specify the properties of an I/O page, including its read-write protection, memory coherence, storage-ordering, RPN of the page's translated I/O address, hint for entry-caching, and the IOID of the devices to which this page applies. This format is called format 1 for IOPTs; it is the only format supported in the CBEA processors.

Generally, software should not create IOPT entries that give I/O devices access to the IOST or IOPT, because this allows the I/O device to circumvent I/O-address translation storage protection.

*Page Protection*

The page protection (PP) bits in the IOPT entry define the type of access allowed for the I/O unit. If PP = '00', no access is allowed. If PP = '01' or '11', an I/O read of the page is permitted. If PP = '10' or '11', an I/O write of the page is permitted. If the access is attempted but not permitted, an I/O page-fault type of I/O exception occurs. Software can set PP = '00' to mark invalid entries.

*Memory Coherence*

The IOPT M bit is similar to the PowerPC memory-coherence-required storage attribute. If M = '1' and the IOIF C bit = '1' (*Section 7.7.1* on page 188), memory coherence is maintained by hardware. If M = '0' or if the IOIF C bit = '0', memory coherence is not maintained by hardware. For storage that does not have the memory-coherence-required attribute, software must explicitly manage memory coherence to the extent required for the program to run correctly. See *PowerPC Virtual Environment Architecture, Book II* for a detailed definition of the memory-coherence-required storage attribute.

*Storage Order*

The Storage Ordering (SO) bits define the storage order rules that apply to reads and writes from one IOIF with the same virtual channel and same IOID, unless the IOIF S bit (*Section 7.7.2.2* on page 190) allows a more relaxed ordering. Previous writes on the same IOIF with the same virtual channel and the same IOID will be performed before a write with a corresponding SO = '10' or '11' and with the IOIF S bit = '1'. Previous reads and writes on the same IOIF with the same virtual channel and same IOID will be performed before a read or write with a corresponding SO = '11' and with the IOIF S bit = '1'.

If SO = '00' or the IOIF S bit = '0', the CBEA processors do not order the access relative to previous accesses. See *Section 7.7.2.2* on page 190 for additional details, including complexities associated with locations that do not have the memory-coherence-required storage attribute.

*Real Page Number*

An I/O access will read or write the RPN of the translated I/O address if the page-protection (PP) bits allow the access for the specified I/O device. If $2^p$ bytes is the page size (as specified in the IOST PS field, *Section 7.3.3* on page 169), and $p$ is greater than 12 (page size greater than 4 KB), then the least-significant ($p$ - 12) bits of the RPN must be '0'. The CBEA processors

support 42 real address bits; bits 0 through 16 of the RPN field must be zeros in IOPT entry. The I/O address is checked to ensure that it does not access a page beyond the number of pages in the segment, as defined by the NPPT field of the IOST.

Assume a page size of $2^p$ bytes is specified in the IOST PS field. If the value of the most-significant (28 - *p*) bits of the 28 least-significant I/O address bits is greater than 512 times the NPPT, an I/O page fault (*Section 9.6.4.4* on page 274) occurs.

*Cache Hint*

Depending on the `IOC_IOST_Origin[HL]` bit (*Section 7.3.2* on page 166), the IOPT Hint (H) bit is treated as either a hint or a lock. If H = '1', it is a hint to the hardware to retain a valid IOPT-cache entry, if possible; if H = '1' and `IOC_IOST_Origin[HL]` = '1', the IOPT entry is never displaced by hardware during an IOPT cache miss.

The IOPT cache (*Section 7.6.2* on page 183) is 4-way associative with 64 congruence classes. For a given congruence class, if no valid IOPT entries have H = '1' and `IOC_IOST_Origin[HW]` = '1', then an IOPT-cache miss in this congruence class causes hardware to select one of the four ways in the congruence class, based on a pseudo-least-recently-used (LRU) algorithm. For an IOPT-cache miss in a second congruence class with two valid entries having H = '1' and two entries having H = '0', the IOPT entry will be placed in one of the cache entries having H = '0'. For an IOPT-cache miss in a third class with three valid entries having H = '1' and one entry having H = '0', the IOPT entry will be placed in the cache way with the entry having H = '0'.

*I/O Device Identifier*

The IOID bits specify the I/O device identifier. Inbound accesses from an I/O unit to the EIB contain an IOID as part of their access command, and hardware compares this access-command IOID to the IOID in the IOPT. Only the I/O device with this IOID can have its I/O address translated using the IOPT entry. If other I/O devices attempt translation using this IOPT entry, it will result in an I/O exception.

*Looped-Back Operations*

The CBEA processors do not support looped-back operations from an I/O device.

Architecturally, however, support for looped-back operations is optional for each implementation of the CBEA. When supported, commands that are received by the CBEA processors from an IOIF device that are routed back out the same IOIF are called *looped-back operations*. To prevent IOIF looped-back operations when I/O-address translation is enabled, one option is to ensure that the RPNs in the IOPT are not in the range of real addresses mapped to an IOIF. If direct accesses from one IOIF to another IOIF are required, the RPNs corresponding to the second IOIF need to be in the IOPT. In this case, looped-back operations from this second IOIF can be prevented by using IOIDs for the IOIF-to-IOIF accesses that are different from IOIDs used on the second IOIF.

#### 7.3.4.3 *Table-Entry Modifications*

An IOPT entry should only be modified when there is no simultaneous I/O access that uses the IOPT entry; if this condition is not met, it is unpredictable whether the I/O access will use the old or new value of the IOPT entry. If software chooses to stop simultaneous I/O accesses, such accesses must be stopped at the I/O unit. To prevent an I/O-address translation access of the IOPT from getting part of an old value and part of a new value when there is a simultaneous I/O access using the entry, software should take one of the following actions:

- Use a doubleword atomic store to the IOPT entry. According to *PowerPC Architecture*, this is a single-register doubleword store that is naturally aligned.
- Use an atomic store to set the IOPT valid (V) bit to '0', execute an **eieio**, **lwsync**, or **sync** instruction, and then store to other bytes in the IOPT entry. After making the updates, execute an **eieio**, **lwsync**, or **sync** instruction, and then execute an atomic store to set the IOPT V bit to '1'. The **lwsync** instruction can only be used for this purpose if the IOPT is in storage that is not caching-inhibited (but see *Section 7.7.1* on page 188). The **eieio** and **sync** instructions are unnecessary if the IOPT is in storage that is both caching-inhibited and guarded.
- Invalidate the IOPT cache to ensure that the new IOPT entry is read into the cache.

### 7.3.5 IOC Base Address Registers

The IOC accepts commands from the EIB in which the address of the command matches the real-address ranges specified for each IOIF in the following MMIO registers:

- IOC Base Address Register 0 (`IOC_BaseAddr0`)
- IOC Base Address Register 1 (`IOC_BaseAddr1`)
- IOC Base Address Mask Register 0 (`IOC_BaseAddrMask0`)
- IOC Base Address Mask Register 1 (`IOC_BaseAddrMask1`)

The `IOC_BaseAddr`*n* and `IOC_BaseAddrMask`*n* together identify the range of real addresses to be routed to IOIF*n*, where the value of *n* is '0' or '1' and is equal to the number of the I/O Interface. The IOC supports two IOIF buses, two `IOC_BaseAddr` registers, and two `IOC_BaseAddrMask` registers.

Because the CBEA processors support only 42-bits of real address, only the least-significant 11 bits of the `IOC_BaseAddr`*n* registers' Base Real Address and Base Replacement Address fields are implemented. Likewise, only the least-significant 11 bits of the `IOC_BaseAddrMask`*n* registers' mask fields are implemented.

Valid values in the mask field and the corresponding address range routed to an IOIF are shown in the *Table 7-1*.

*Table 7-1. IOIF Mask and Size of Address Range (Sheet 1 of 2)*

| Mask in `IOC_BaseAddrMask`*n* Registers | Size of Address Range |
|---|---|
| x'000007FF' | 2 GB |
| x'000007FE' | 4 GB |
| x'000007FC' | 8 GB |
| x'000007F8' | 16 GB |

*Table 7-1. IOIF Mask and Size of Address Range (Sheet 2 of 2)*

| Mask in IOC_BaseAddrMask*n* Registers | Size of Address Range |
|---|---|
| x'000007F0' | 32 GB |
| x'000007E0' | 64 GB |
| x'000007C0' | 128 GB |
| x'00000780' | 256 GB |
| x'00000700' | 512 GB |
| x'00000600' | 1 TB |
| x'00000400' | 2 TB |
| x'00000000' | 4 TB |

When modifying IOC_BaseAddr*n* or IOC_BaseAddrMask*n*, software should ensure that no accesses to the previous or new address range occur by using the following procedure:

1. Ensure that pages for the old and new address ranges have the guarded storage attribute.
2. Ensure that SPE accesses do not occur while these registers are modified, for example, by suspending memory flow controller (MFC) DMA transfers using the MFC_CNTL register.
3. Ensure that previous PPE accesses are performed before storing to IOC_BaseAddr*n* or IOC_BaseAddrMask*n* by using one of the synchronization facilities described in *Section 20* on page 561 (for example, the **sync** instruction).
4. Ensure that read and write commands from IOIF*n* are performed on the EIB before storing to IOC_BaseAddr*n* or IOC_BaseAddrMask*n*. One potential option is to stop such accesses at the IOIF device or I/O device, if such a facility is provided (an I/O device, here, means a device that is connected to an IOIF indirectly by an IOIF device). If the IOIF device or I/O device supports a facility to send an interrupt after its last access, then reception of the external interrupt by the PPE is a guarantee that the previous command from the IOIF has been performed on the EIB. Although this does not guarantee the data transfer has completed, it is sufficient to ensure that the command phase has completed. Alternatively, if reads and writes from the IOIF are translated by the IOC (IOC_IOCmd_Cfg[TE] = '1' and IOC_IOST_Origin[0] = '1'), then the pages corresponding to the range specified by IOC_BaseAddr*n* and IOC_BaseAddrMask*n* can be invalidated in the IOST (*Section 7.3.3* on page 169) or IOPT (*Section 7.3.4* on page 171) and IOST cache (*Section 7.6.1* on page 181) or IOPT cache (*Section 7.6.2* on page 183), and, if the IOIF device or I/O device supports a facility to send an interrupt after this invalidation, reception of this external interrupt by the PPE is a guarantee that the previous command from the IOIF has been performed on the EIB.
5. Disable accesses to the IOIF by setting IOC_BaseAddrMask*n*[E] = '0'.
6. Store to IOC_BaseAddr*n* with the new value.
7. Store to IOC_BaseAddrMask*n* with the new value, setting IOC_BaseAddrMask*n*[E] = '1'.
8. Load from IOC_BaseAddrMask*n* to ensure the effect of the previous store to IOC_BaseAddrMask*n*.
9. Execute a **sync** instruction to ensure the completion of the previous load before subsequent accesses.

### 7.3.6 I/O Exception Status Register

The I/O Exception Status Register (`IOC_IO_ExcpStat`) captures error information for an I/O exception. This register is described in *Section 7.5 I/O Exceptions* on page 180.

## 7.4 I/O Address Translation

### 7.4.1 Translation Overview

I/O address translation is available as an option on interfaces configured for the IOIF (but not BIF) protocol. If I/O address translation is enabled, the IOC supports an I/O address space of $2^{35}$ bytes, which is translated into a 42-bit real addresses. In this case, the IOC uses the least-significant 35 bits of the IOIF address field and ignores the remaining most-significant bits[3].

The I/O address space is divided into segments. Each segment has $2^{28}$ bytes, and there are $2^{7}$ segments. The supported I/O page sizes are 4 KB, 64 KB, 1 MB, and 16 MB. One of the four page sizes can be selected for each I/O segment. One, two, three or four of these pages sizes can be used concurrently in the system.

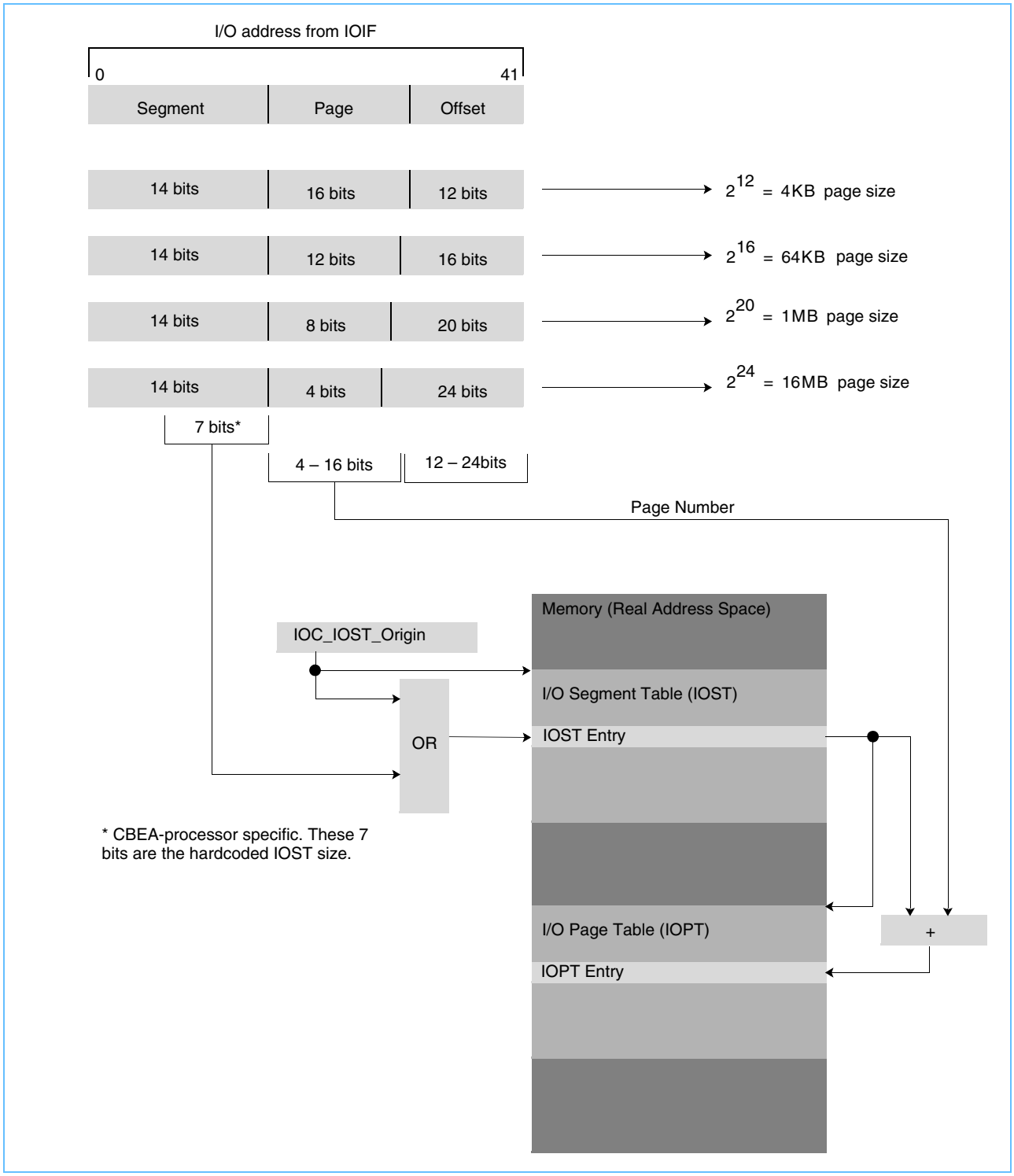
The IOC supports an IOST cache (*Section 7.6.1* on page 181) and an IOPT cache (*Section 7.6.2* on page 183), which can be loaded by hardware or software. Multiple coherency and ordering modes support both optimized performance for I/O accesses to memory as well as coherent accesses with strong ordering when needed.

All I/O accesses by the CBEA processors are big-endian accesses. If an I/O access is to a little-endian device, such as a little-endian MMIO register, software must reorder the bytes into big-endian order.

*Figure 7-5* on page 177 shows an overview of the translation process in accessing the IOST and IOPT. For each of the four possible page sizes, this figure shows which bits of the I/O address are used in the CBEA processors to select the IOST entry and IOPT entry. The OR operation on the `IOC_IOST_Origin` and the 7-bit segment values is equivalent to an add operation because `IOC_IOST_Origin` is naturally aligned to 4 KB boundaries.

3. Bit numbering of I/O addresses in this section is based on a 42-bit I/O address with bit 0 being the most-significant and bit 41 being the least-significant.

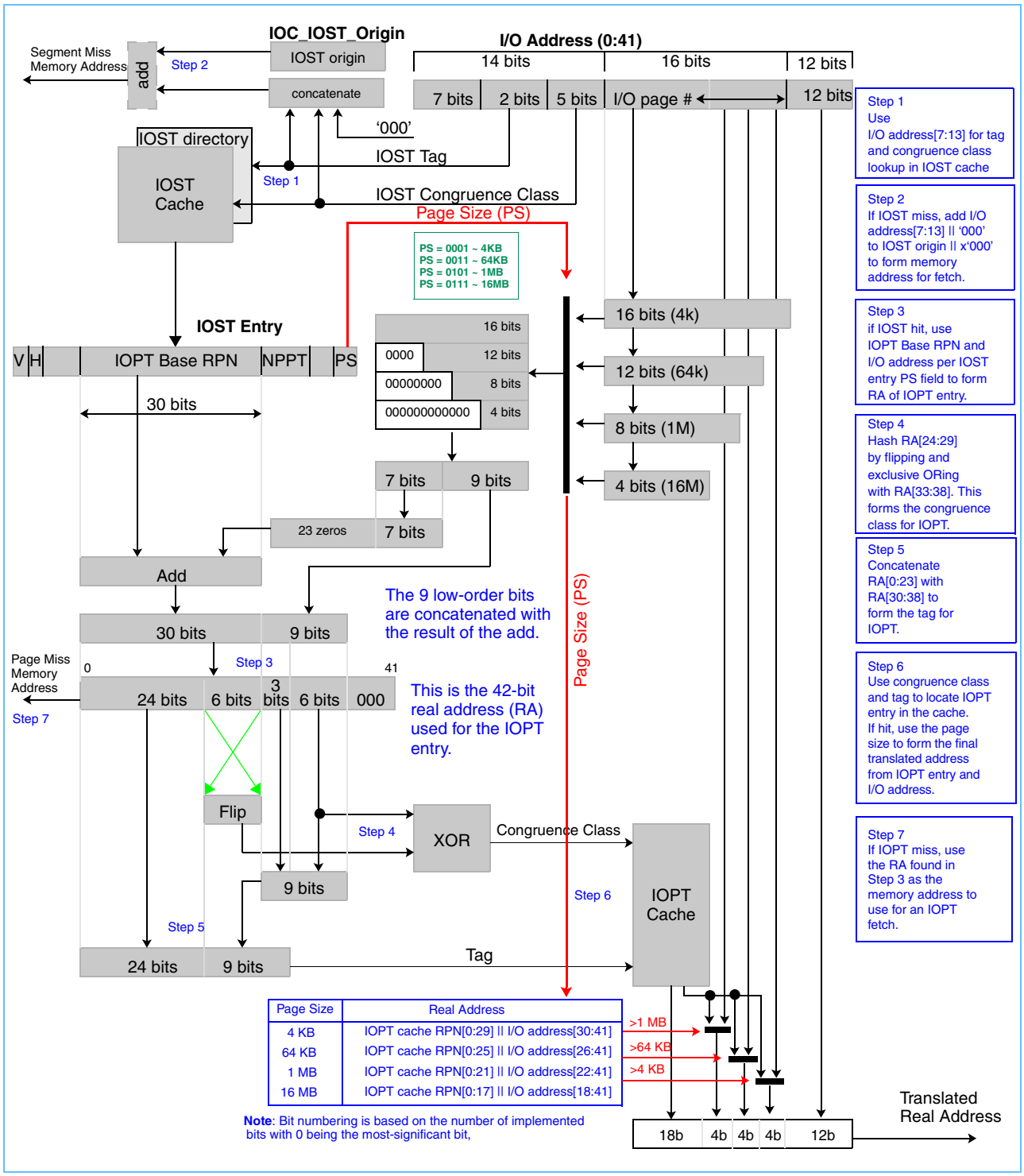
*Figure 7-5. I/O-Address Translation Overview*



### 7.4.2 Translation Steps

*Figure 7-6* shows a detailed view of the I/O translation process, including details about the IOST cache and IOPT cache.

*Figure 7-6. I/O-Address Translation Using Cache*



I/O-address translation is enabled if IOC_IOST_Origin[E] = '1' and IOC_IOCmd_Cfg[TE] = '1'. I/O-address translation is disabled if IOC_IOST_Origin[E] = '0' and IOC_IOCmd_Cfg[TE] = '0'. In either case, the IOC_IOCmd_Cfg[TE] bit must be set to a value consistent with IOC_IOST_Origin[E].

If I/O-address translation is disabled, the I/O address from the IOIF is treated as a real address. In this case, the IOC uses the least-significant 42 bits of the IOIF address and ignores the most-significant 22 bits.

If I/O-address translation is enabled, the following steps are taken by the I/O address-translation mechanism if there are misses in the IOST or IOPT cache:

1. If an IOST-cache miss occurs and IOC_IOST_Origin[HW] = '0', an I/O segment-fault type of I/O exception occurs (*Section 7.5* on page 180), and I/O-address translation stops. If IOC_IOST_Origin[HW] = '1', the IOST entry (*Section 7.3.3* on page 169) is loaded using the I/O address, the IOST origin field in the IOC_IOST_Origin register, and an index based on the IOST size field in the IOC_IOST_Origin register. Concatenating forms the real address of the IOST entry, as follows:
   - The most-significant (50 - $s$) bits of the IOST origin, where $s$ is equal to the maximum of 0 and (IOST size - 9)
   - $n$ '0' bits, where $n$ is equal to the maximum of 0 and (9 - IOST size)
   - $t$ least-significant bits of the I/O address-segment bits, where $t$ equals the value of the IOST size field
   - Three '0' bits
2. If the IOST-entry valid bit is '0', an I/O segment-fault type of I/O exception occurs and I/O-address translation stops. If the IOST-entry-valid bit is '1', the I/O translation process continues with the following step.
3. The I/O address is checked to ensure that it does not access a page beyond the number of pages in the segment as defined by the NPPT field of the IOST. A page size of $2^p$ is specified in the IOST page-size field. For format 1 (the only format supported in the CBEA processors), if the value of the most-significant (28 - $p$) bits of the 28 least-significant I/O address bits is greater than 512 times the NPPT, an I/O page-fault type of I/O exception occurs and I/O-address translation stops. If an IOPT-cache miss occurs and IOC_IOST_Origin[HW] = '0', an I/O page-fault type of I/O exception occurs and the I/O-address translation stops.
4. A real address is formed and used to access the IOPT entry (*Section 7.3.4* on page 171). The real address is formed by adding the following two values:
   - The IOPT-base RPN concatenated with x'000'.
   - The most-significant (28 - $p$) bits of the 28 least-significant I/O address bits concatenated with '000', where $2^p$ is the page size specified in the IOST page-size field.
5. If any of the following apply, an I/O page-fault type of I/O exception occurs and I/O-address translation stops:
   - The IOPT PP bits are '0'.
   - The IOPT PP bits are '01' (read only), and the I/O access was a write operation.
   - The IOPT PP bits are '10' (write only), and the I/O access was a read operation.
   - The IOID does not match the I/O device.
6. If the access is permitted by the IOPT PP bits, a real address is formed to use for the I/O access. For IOPT format 1, the real address is formed by concatenating bits 0 through (58 -

*p*) of the IOPT RPN with the least-significant *p* bits of the I/O address. The page size specified in the IOST is $2^p$.

## 7.5 I/O Exceptions

### 7.5.1 I/O Exception Causes

If I/O-address translation is enabled (`IOC_IOST_Origin[E]` = '1' and `IOC_IOCmd_Cfg[TE]` = '1'), an I/O exception (*Section 9.6.4.4 I/O Address Translation* on page 274) can occur due to one of the following events:

- An I/O segment fault occurs in the following cases when a translation is attempted for an I/O address:
  - `IOC_IOST_Origin[HW]` = '0', and an IOST cache miss occurs.
  - `IOC_IOST_Origin[HW]` = '1', an IOST cache miss occurs, and there is no corresponding entry in the IOST.
- An I/O page fault occurs in the following cases when a translation is attempted for an I/O address and there is no I/O segment fault:
  - `IOC_IOST_Origin[HW]` = '0' and an IOPT cache miss occurs.
  - `IOC_IOST_Origin[HW]` = '1', an IOPT cache miss occurs, and there is no corresponding entry in the IOPT.
  - `IOC_IOST_Origin[HW]` = '1', `IOC_IOST_Origin[HL]` = '1', an IOPT cache miss occurs and all 4 entries for the IOPT cache congruence class have the H bit equal to '1'.
- An I/O access that is either translated or not translated and that accesses an invalid real address.
- A CBEA processor access to an invalid location in the range of addresses mapped to an IOIF bus.

An I/O exception results in an external interrupt (*Section 9.6 Direct External Interrupts* on page 265). I/O exceptions caused by I/O accesses are reported to the I/O subsystem that attempted the I/O-address translation. If any of the events in the preceding list are reported as an I/O exception, system error interrupt, or machine-check interrupt, and if the access is an SPE DMA get or a PPE load, the data returned contains all ones. If reported as an I/O exception, system error interrupt, or machine-check interrupt, and if the access is an SPE DMA put or a PPE store, the data is discarded. When an I/O exception occurs on an I/O access, additional information is captured in the I/O Exception Status Register (`IOC_IO_ExcpStat`); see *Section 7.5.2* on page 181.

***Note:*** *Software should avoid I/O exceptions. If I/O exceptions occur, they should only be a result of a software bug or a result of hardware associated with an IOIF that is not working. Whether software can or should continue operation after an I/O operation causes an I/O exception is up to the programmer or the system designer.*

### 7.5.2 I/O Exception Status Register

The I/O Exception Status Register (`IOC_IO_ExcpStat`) captures error information for an I/O exception. The bit fields and values for this register are described in the *Cell Broadband Engine Registers* specification.

An *IOIF device* is a device that is directly connected to a CBEA processor's IOIF port. When an I/O exception occurs on an I/O access, an Error response is sent to the IOIF device, and the following information related to the fault is captured in the `IOC_IO_ExcpStat` register:

- I/O address[7:29].
- Access type (read or write).
- I/O device ID. (An *I/O device*, in this context, is a device that is indirectly connected to a CBEA processor's IOIF by an IOIF device.)

### 7.5.3 I/O Exception Mask Register

The `IOC_IO_ExcpMask`[1:2] bits are ANDed with associated I/O segment fault and I/O page fault conditions to set `IIC_ISR`[61], but these mask bits do not gate the setting of the `IOC_IO_ExcpStat` bits. See *Section 9 PPE Interrupts* on page 239 for information about when an external interrupt signal to the PPE is asserted based on these bits.

### 7.5.4 I/O-Exception Response

Unlike PowerPC page tables, software should normally not expect to set up an IOST or IOPT entry after an I/O exception. There is no hardware support to allow an IOIF command to be suspended at the point of an I/O exception and later restarted.

For many types of I/O operations, overruns and underruns occur if there is significant delay associated with an I/O exception. Also, the I/O reads and writes to one page might be delayed by the I/O exception while subsequent I/O reads, writes, and interrupt operations might complete, thereby violating any applicable ordering requirements and potentially giving software an erroneous indication that some I/O operation has successfully completed.

## 7.6 I/O Address-Translation Caches

To improve performance, the IOC supports an IOST cache and an IOPT cache, shown in *Figure 7-6* on page 178. These caches can be loaded by hardware or software. Both software and hardware can control the contents of these caches as described in the following sections.

### 7.6.1 IOST Cache

The IOST cache is direct-mapped and has 32 entries, indexed by I/O address bits [9:13]. These address bits correspond to the least-significant five bits of the I/O segment number. If `IOC_IOST_Origin[HW]` = '1', a valid IOST cache entry for the I/O address does not exist, and a valid IOST entry for the I/O address exists, the IOST cache is loaded automatically by hardware. If `IOC_IOST_Origin[HW]` = '0' and a valid IOST cache entry for the I/O address does not exist, an I/O segment fault occurs.

#### 7.6.1.1 *IOST Cache Access by Software*

A doubleword load or store can be used to access each IOST cache entry in the MMIO memory map. For information about this map, see the *Cell Broadband Engine Registers* specification, which includes the offset range at which the IOST cache is accessible in the IOC address-translation MMIO memory map. For a specific 42-bit I/O address, the corresponding IOST cache-entry offset in this range is formed by using I/O address[9:13] concatenated with '000'.

The fields in the IOST cache entry have the same meaning as the corresponding fields in the IOST, except for the following differences:

- The H bit is not implemented and is treated as a '0'.
- Only the least-significant 30 bits of the IOPT Base RPN are implemented.
- Only the least-significant 3 bits of the PS field are implemented.
- There is an additional valid bit and tag field.

The IOST cache-entry valid bit and tag field correspond to values that are typically stored in a cache directory. The IOST cache-entry valid bit indicates that the IOST cache entry is valid, whereas an IOST valid bit indicates that the IOST entry is valid. There is no tag field in an IOST entry. For a given 42-bit I/O address with a corresponding valid entry in the IOST cache, the IOST cache-entry tag field has the same value as I/O address[7:8]. The valid bit should be set to '1' to allow I/O-address translation to use the entry for translating I/O accesses.

After a store to the IOST cache, to guarantee that the new value of the IOST cache entry is used on future I/O-address translations, software should load from any IOST cache entry and ensure the load has completed before a subsequent store that might cause the new value to be used. See *Section 7.6.1.3 Simultaneous Hardware and Software Reload of IOST Cache* on page 183 for other considerations when `IOC_IOST_Origin[HW]` = '1'.

Software should ensure that, for IOST cache entries that software reloads, it does not attempt to have two or more valid entries in the IOST with the same value in the least-significant 5 bits of the I/O segment number, because the IOST cache is direct-mapped and has only 32 entries indexed by I/O address[9:13].

If software fails to pre-load an IOST cache entry, and either the hardware IOST cache-miss handling is disabled (`IOC_IOST_Origin[HW]` = '0') or the IOST table does not contain the corresponding entry, an I/O segment fault occurs when an I/O device attempts to read or write the location with the corresponding the missing entry.

#### 7.6.1.2 *IOST Cache Invalidation by Software*

An MMIO store to the IOST Cache Invalidate Register (`IOC_IOST_CacheInvd`) invalidates one entry in the IOST cache. The I/O segment to be invalidated is defined by `IOC_IOST_CacheInvd`[29:35]. Because the IOST cache is direct-mapped, the least-significant 5 bits of `IOC_IOST_CacheInvd`[29:35] specify the index of the IOST cache entry. Software must set `IOC_IOST_CacheInvd`[B] to '1' when storing to the IOST Cache Invalidate Register; otherwise, the result of the store is undefined.

When the specified I/O segment has been invalidated in the IOST cache, hardware sets `IOC_IOST_CacheInvd`[B] to '0'. Software must only store to the IOST Cache Invalidate Register when `IOC_IOST_CacheInvd`[B] is '0'; otherwise, the results of the previous store and current store

to the IOST Cache Invalidate Register are undefined. When IOC_IOST_CacheInvd[B] is set to '0' by the IOC, this guarantees that I/O accesses no longer use the old cached IOST value. However, it does not guarantee that all previous I/O accesses that used the old cached IOST value have been performed. Typically, software can depend on other facilities to know that these I/O accesses have been performed—for example, notification by the I/O device by means of an external interrupt.

Software must ensure that the IOST cache invalidation has completed (IOC_IOST_CacheInvd[B] = '0') before writing to the IOC_IOST_CacheInvd register, IOPT cache, IOC_IOPT_CacheDir register, IOST cache, or IOST cache directory. Software must also ensure that a store to the IOPT cache, IOC_IOPT_CacheDir, IOST cache, or IOST cache directory completes before storing to IOC_IOST_CacheInvd by executing a **sync** instruction between the last store to one of these arrays and the store to IOC_IOST_CacheInvd.

When IOC_IOST_Origin[HW] = '1', software typically modifies IOST entries before invalidating the IOST cache; otherwise, the I/O-address translation facility would reload the old values if needed by an I/O access. To ensure the last store to the IOST is performed before the store to IOC_IOST_CacheInvd, software should use PowerPC memory-barrier instructions, such as **sync**, or should write the IOST using the storage attributes of both caching-inhibited and guarded.

***Note:*** *Due to the asynchronous nature of DMA transfers, it is important to ensure that DMA transfers have completed before removing an IOST-cache or IOPT-cache invalidation. Otherwise, an I/O exception might occur.*

#### 7.6.1.3 *Simultaneous Hardware and Software Reload of IOST Cache*

If IOC_IOST_Origin[HW] and IOC_IOST_Origin[HL] are both '1', both software and the I/O address-translation hardware can load the IOST cache when a miss occurs. This might result in hardware and software writing the same IOST cache entries.

To prevent hardware from overwriting an entry written by software, software can ensure that, for any IOST cache entry that software reloads, there are no valid entries in the IOST with the same cache congruence class. If there ever were any such entries since boot, and if software wants to start using the IOST entry with the same congruence class, software must first make these entries in the IOST invalid, execute an enforce in-order execution of I/O (**eieio**) instruction or a synchronize (**sync**) instruction, invalidate the IOST cache, and wait for invalidation to complete before writing the IOST entry.

Software can modify an IOST cache entry when there are no outstanding I/O accesses. How software ensures that there are no outstanding I/O accesses is application-dependent.

### 7.6.2 IOPT Cache

The IOPT cache is 4-way set-associative and has 256 entries, indexed by a 6-bit hash formed from the I/O address, as shown in *Figure 7-7* on page 185. For each index (congruence class), a pseudo-LRU algorithm is used to determine which entry is to be replaced when hardware cache-miss handling is enabled (that is, when IOC_IOST_Origin[HW] = '1'). If IOC_IOST_Origin[HL] = '0' and any entry in the congruence class has its Hint (H) bit or Valid (V) bit set to '0', this pseudo-LRU algorithm avoids replacing a valid entry that has its H bit set to '1'. If IOC_IOST_Origin[HL] = '1', this pseudo LRU algorithm does not replace an entry that has its H bit set to '1', even if the valid bit for the IOPT cache entry is '0'.

#### 7.6.2.1 *IOPT Cache Access by Software*

Doubleword loads or stores can be used to access the IOPT cache directory (IOC_IOPT_CacheDir) register in the MMIO memory map. A load of IOC_IOPT_CacheDir simultaneously causes the corresponding IOPT cache entry to be read into the IOPT Cache Register (IOC_IOPT_Cache). Similarly, a store to IOC_IOPT_CacheDir simultaneously causes the corresponding IOPT cache entry to be written from IOC_IOPT_Cache.

As shown in the *Cell Broadband Engine Registers* specification, IOC_IOPT_CacheDir register is accessible at a set of offset ranges in the IOC address-translation MMIO memory map. Each way of IOC_IOPT_CacheDir is accessed using a different 512-byte range. For a specific 42-bit I/O address, the corresponding IOC_IOPT_CacheDir entry offset in this range is formed by adding of two quantities:

- x‘0200’ times 1 minus the number of the way of IOC_IOPT_CacheDir
- 6 bits of the IOPT cache hash concatenated with ‘000’

Software can read an entry in the IOPT cache by:

1. Performing a doubleword load from the IOC_IOPT_CacheDir entry. This simultaneously causes the corresponding IOPT cache entry to be read into IOC_IOPT_Cache.
2. Performing a doubleword load from IOC_IOPT_Cache.

Software can write an entry in the IOPT cache by:

1. Storing a doubleword to IOC_IOPT_Cache. This store does not cause an IOPT cache entry to be updated.
2. Storing a doubleword to the IOC_IOPT_CacheDir entry. This store causes IOC_IOPT_CacheDir to be loaded with the store data and causes the IOPT cache to be simultaneously loaded with the value in IOC_IOPT_Cache. Hardware-generated parity bits are also automatically loaded. The tag, as determined in *Section 7.6.2.2 IOPT Cache Hash and IOPT Cache Directory Tag* on page 185, should be written into the tag field of IOC_IOPT_CacheDir entry and the valid bit should be set to ‘1’ to allow I/O-address translation to use this new entry for translating I/O accesses.

***Note:*** *Software must not store to the IOPT Cache Invalidate (IOC_IOPT_CacheInvd) register between the preceding steps 1 and 2 of the read or write sequence; otherwise, the result of step 2 is undefined.*

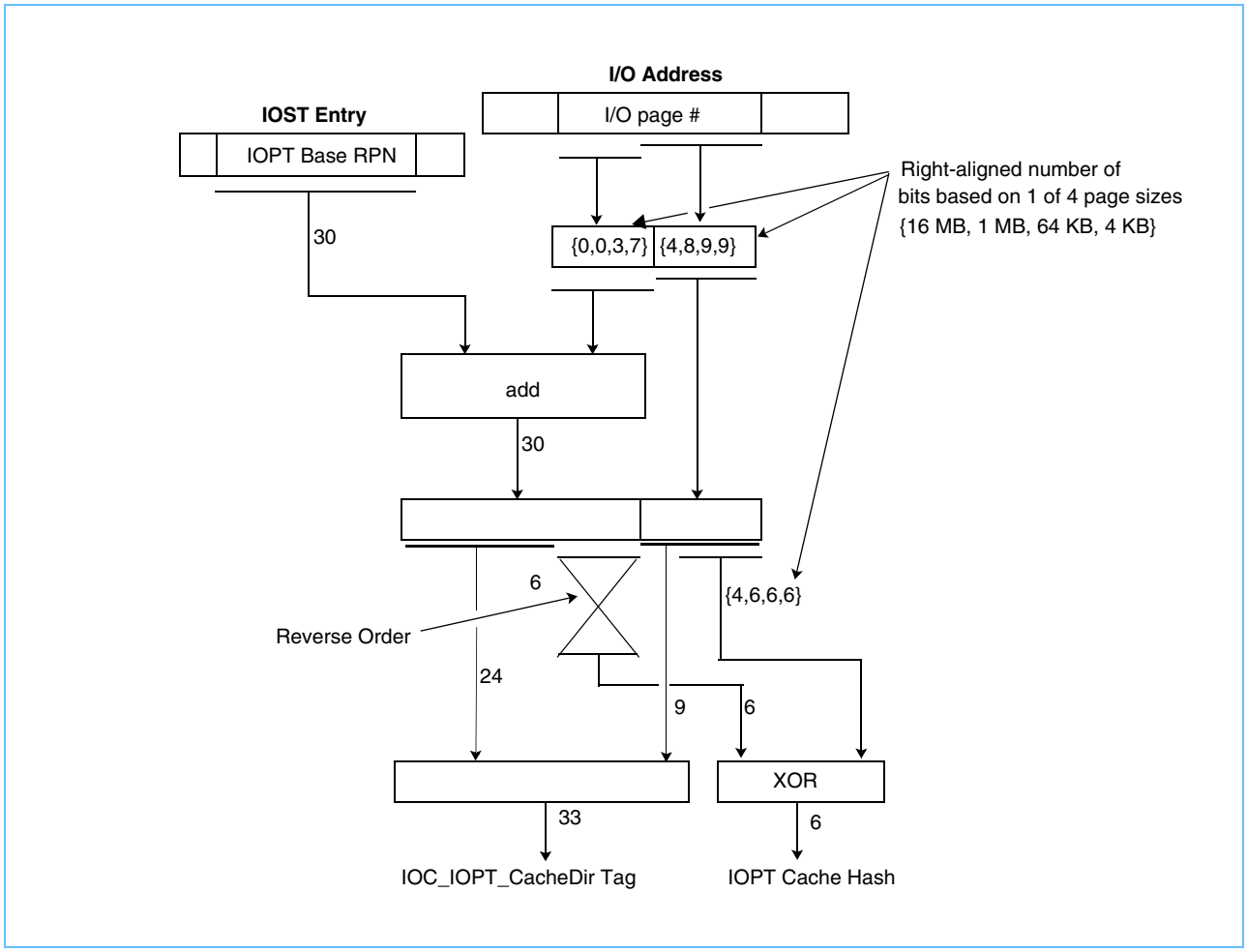
After a store to IOC_IOPT_CacheDir, to guarantee that the new value of the IOPT cache entry is used on future I/O-address translations, software should load from any IOC_IOPT_CacheDir entry and ensure the load has completed before a subsequent store that might cause the new value to be used. See *Section 7.6.2.4 Simultaneous Hardware and Software Reload of IOPT Cache* on page 187 for other considerations when IOC_IOST_Origin[HW] = ‘1’.

Software must not write two or more ways of the IOC_IOPT_CacheDir for the same congruence class with valid entries having the same tag value; otherwise, the results of subsequent I/O-address translations are undefined.

### 7.6.2.2 *IOPT Cache Hash and IOPT Cache Directory Tag*

If hardware I/O-address translation is enabled, software must know how the IOPT cache hash is created to avoid setting too many bits in the same congruence class. The 6-bit IOPT cache hash is used by the hardware to select the IOPT cache congruence class for hardware reloads. The generation of this hash value is shown in *Figure 7-7*.

*Figure 7-7. IOPT Cache Hash and IOPT Cache Directory Tag*




The generation of the hash value in *Figure 7-7* is based on the following steps:

1. Select 4, 8, 9, or 9 bits of the page number from the I/O address, based on the page size of 16 MB, 1 MB, 64 KB, or 4 KB, respectively. The page size is defined by the IOST entry PS field.
2. Select the most-significant 0, 0, 3, or 7 bits of the page number from the I/O address, based on the page size of 16 MB, 1 MB, 64 KB, or 4 KB, respectively.
3. Add the value selected in step 2 to the 30-bit IOPT Base RPN from the IOST entry.
4. Select the least-significant 6 bits of the result from the previous step.
5. Put these bits in reverse order; that is, swap bits 0 and 5, swap bits 1 and 4, and swap bits 2 and 3.

6. From the bits selected in step 1, select the least-significant 4, 6, 6, or 6 bits, based on the page size of 16 MB, 1 MB, 64 KB, or 4 KB, respectively.
7. If the page size is 16 MB, create a 6-bit value from two ‘0’ bits concatenated with the 4 bits from the previous step; otherwise, create a 6-bit value equivalent to the value created in the previous step.
8. Exclusive-OR the 6 bit value from step 5 with the 6-bit value from the previous step.

The result of the exclusive-OR in step 8 is the IOPT cache hash.

The IOC_IOPT_CacheDir tag is determined by the following additional steps:

a. Create a 9-bit value by padding as many ‘0’ bits as necessary to the left of the bits from step 1.
b. Form a 33-bit IOC_IOPT_CacheDir tag by concatenating the most-significant 24 bits from the addition result in step 3 with the 9-bit value from the previous step.

#### 7.6.2.3 *IOPT Cache Invalidation by Software*

An MMIO store to IOC_IOPT_CacheInvd invalidates IOPT cache entries corresponding to a number of IOPT entries. The value in the Number of Entries (NE) field in IOC_IOPT_CacheInvd is equal to one less than the number of IOPT entries to be invalidated in the IOPT cache. The MMIO store invalidates all IOPT cache entries corresponding to IOPT entries located at the doubleword addresses between (IOPTE_DRA || ‘000’) and ((IOPTE_DRA + NE) || ‘000’), inclusive. The IOPTE_DRA field in IOC_IOPT_CacheInvd specifies the I/O Page Table Entry Doubleword Real Address, where (IOPTE_DRA || ‘000’) is the real address of the first IOPT entry to be invalidated. Software must set IOC_IOPT_CacheInvd[B] to ‘1’ when storing to IOC_IOPT_CacheInvd; otherwise, the result of the store is undefined.

When the IOPT cache entries have been invalidated for the specified IOPT entries, hardware sets the busy bit, IOC_IOPT_CacheInvd[B], to ‘0’. Software must only store to the IOPT Cache Invalidate Register when IOC_IOPT_CacheInvd[B] is ‘0’; otherwise, the results of the previous store and current store to IOC_IOPT_CacheInvd are undefined. When IOC_IOPT_CacheInvd[B] is set to ‘0’ by the IOC, this guarantees that I/O accesses no longer use the old cached IOPT value. However, it does not guarantee that all previous I/O accesses that used the old cached IOPT value have been performed. Typically, software can depend on other facilities to know that these I/O accesses have been performed—for example, notification by the I/O device by means of an external interrupt.

An invalidation of an IOPT cache entry does not change the value of the H bit in the cache.

If the I/O-address translation facility already has an outstanding load of an IOPT entry when software does an IOPT cache invalidation that includes that IOPT entry, then hardware will not set IOC_IOPT_CacheInvd[B] to ‘0’ until the outstanding load completes and either the reloaded IOPT cache entry is invalidated or the reload is somehow aborted.

Software must ensure that the IOPT cache invalidation has completed (IOC_IOPT_CacheInvd[B] = ‘0’) before writing IOC_IOST_CacheInvd, IOPT cache, IOC_IOPT_CacheDir, IOST cache, or IOST cache directory. Software must ensure that a store to the IOPT cache, IOC_IOPT_CacheDir, IOST cache, or IOST cache directory completes before storing to IOC_IOPT_CacheInvd by executing a **sync** instruction between the last store to one of these arrays and the store to IOC_IOPT_CacheInvd.

When IOC_IOST_Origin[HW] = ‘1’, software typically modifies IOPT entries before invalidating the IOPT cache; otherwise, the I/O-address translation facility would reload the old values if needed by an I/O access. Software should use PowerPC memory-barrier instructions, such as **sync**, to ensure that the last store to the IOPT is performed before the store to IOC_IOPT_CacheInvd or software should write the IOPT using the storage attributes of both caching-inhibited and guarded.

#### 7.6.2.4 *Simultaneous Hardware and Software Reload of IOPT Cache*

If IOC_IOST_Origin[HW] and IOC_IOST_Origin[HL] are both '1', both software and the I/O address-translation hardware can load the IOPT cache when a miss occurs. This might result in hardware and software writing the same IOPT cache entries.

To prevent hardware from overwriting an entry written by software, software must set the H bits in the IOPT cache to ‘1’. To prevent hardware from writing an entry that software is changing at the time, software must only modify IOPT cache entries that already have the H bit equal to ‘1’. At boot time, if both software and hardware reloads are to be used, software should reserve a number of entries in each congruence class by setting their H bits to ‘1’ before setting IOC_IOST_Origin[HW] to ‘1’.

Software should only modify an IOPT cache entry for which the H bit is a ‘0’ when there are no outstanding I/O accesses. How software ensures that there are no outstanding I/O accesses is application-dependent. If there is an outstanding IOPT cache miss at the time software writes the H = ‘0’ entry, hardware cache-miss handling might overwrite the entry that software modified even though software set the H bit to ‘1’.

Alternatively, software can use the techniques described in *Section 7.6.1.3* on page 183 for IOST cache loading.

Invalidation of an IOPT cache entry will not change the value of the H bit in the cache. This allows software to use the hardware-invalidation mechanisms and still retain the ability to later load the entries that were reserved for software use (that is, entries with H = ‘1’).

Software can also preload the IOPT cache even when IOC_IOST_Origin[HL] = ‘0’. However, if IOC_IOST_Origin[HW] = ‘1’ and there is a simultaneous I/O access that has an IOPT cache miss in the same congruence class, then hardware miss handling might cause the software entry to be overwritten. If an IOPT cache miss occurred just before software wrote the entry, hardware miss handling might have already selected the entry to be written when the IOPT data arrives from memory. At the time hardware detects an IOPT cache miss, hardware selects an entry in the congruence class, using the first entry to meet certain criteria according to the following order and criteria:

1. An entry with the valid bit equal to ‘0’.
2. An entry with the H bit equal to ‘0’, using the pseudo-LRU algorithm among such entries.
3. An entry with the H bit equal to ‘1’, using the pseudo-LRU algorithm among such entries.

If hardware selects an entry that software subsequently writes before hardware updates the entry from the IOPT, the entry that software writes is overwritten when the IOPT data arrives from memory.

## 7.7 I/O Storage Model

The *PowerPC Architecture* and the *Cell Broadband Engine Architecture* (CBEA) define the order of storage operations by the PPE and SPEs within the processor coherence domain (all CBEA processor elements and all interfaces to main storage). These documents also define the order of storage operations by the PPE and SPEs relative to I/O mechanisms. This section clarifies the order of storage operations by the PPE and SPEs relative to an IOIF, because ordering operations defined by the *PowerPC Architecture* (such as, **sync**) have no corresponding equivalent on the IOIFs.

This section also describes the order of operations on an IOIF and its relationship to the order in which the operations are performed within the processor coherence domain or on another IOIF. Accesses to locations outside the processor coherence domain are considered by the coherence domain to be performed when they are considered performed on the IOIF. From the perspective of the CBEA, these accesses are performed coherently in the coherence domain of the I/O subsystem in the same order as they are performed on the IOIF. However, an I/O subsystem might have a different perspective. It is the responsibility of the I/O subsystem that attaches to the IOIF to define ordering rules for the completion of IOIF operations in the coherence domain of the I/O subsystem.

The processor systems map I/O units to main storage. Storage operations that cross the processor coherence-domain boundary are referred to as *I/O operations*. These accesses can be initiated either internally in the CBEA processor (called *outbound* accesses) or externally (called *inbound* accesses). An access initiated externally is called an *I/O access*, and the units that initiate them are called *I/O units*. An I/O unit can be one or more physical I/O devices, I/O bridges, or other functional units attached to an IOIF, in which one value of the IOID described in the IOPT is used for all accesses by these physical entities.

A physical IOIF can be divided into several virtual IOIFs called *virtual channels*. An access on one virtual channel must be allowed to pass accesses for other virtual channels so that a slower virtual channel does not block a faster virtual channel. (See the *Hardware Initialization Guide* for a description of I/O virtual channels.)

### 7.7.1 Memory Coherence

I/O accesses—inbound accesses from an I/O unit to the CBEA processor's element interconnect bus (EIB)—can be treated as either memory-coherence-required or not, depending on the following control bits:

- *IOC TE Bit*—This is the enable bit for I/O-address translation. It is located in the IOCmd Configuration Register (`IOC_IOCmd_Cfg`). Translation is enabled when TE = '1'. See the *Cell Broadband Engine Registers* specification for details.
- *IOIF C Bit*—This is the memory-coherence bit sent with each command on the IOIF. Memory coherence is required on any read or write when C = '1', and memory coherence is not required when C = '0'.
- *IOPT M Bit*—This is the memory-coherence bit from the IOPT entry or IOPT-cache entry corresponding to the accessed location, as described in *Section 7.3.4* on page 171. Memory coherence is required when M = '1', and memory coherence is not required when M = '0'.

*Table 7-2* summarizes the effects on memory coherence of I/O accesses for the various combinations of control bits. If either C = '0' or M = '0', the access is performed on the EIB as not coherent. For incoming I/O addresses that are not translated, the IOIF C bit is used to determine coherence requirements. If this bit is '0', the access is performed on the EIB as not coherent. The access is also performed on the EIB as not coherent if the address is in the range defined by the EIB Local Base Address Register 0 (`EIB_LBAR0`) and the EIB Local Base Address Mask Register 0 (`EIB_LBAMR0`)—or by the `EIB_LBAR1` and `EIB_LBAMR1` registers.

*Table 7-2. Memory Coherence Requirements for I/O Accesses*

| Control Bit Combination[1] | | | Memory-Coherence-Required Attribute Used for the I/O Access |
|---|---|---|---|
| TE | C | M | |
| x | 0 | x | Not coherent |
| 0 | 1 | x | Memory-coherence-required |
| 1 | 1 | 0 | Not coherent |
| 1 | 1 | 1 | Memory-coherence-required |

1. An "x" indicates don't care or not applicable.

There is no hardware support for maintaining I/O-subsystem memory coherence relative to caches in the processors. Any data in the I/O subsystem must be accessed with the storage attribute of caching-inhibited, or software must manage its coherency. For storage that does not have the memory-coherence-required attribute, software must manage memory coherence to the extent required for the program to run correctly. The operations required to do this can be system-dependent. Because the memory-coherence-required attribute for a given storage location is of little use unless the PPE, all SPEs, and all I/O units that access the location do so coherently, statements about memory-coherence-required storage elsewhere in this section generally assume that the storage has the memory-coherence-required attribute for all the PPE, SPEs, and I/O units that access it.

### 7.7.2 Storage-Access Ordering

The following sections describe both outbound and inbound ordering of storage accesses between the EIB and the IOIFs. For background about storage ordering, see *Section 20.1 Shared-Storage Ordering* on page 561.

#### 7.7.2.1 *Outbound (EIB to IOIF) Accesses*

All outbound read and write accesses to I/O locations mapped to main storage locations are performed in program order. The IOIF C bit (*Section 7.7.1* on page 188) and the IOIF S bit (*Section 7.7.2.2* on page 190) are driven to indicate memory-coherence-required on all accesses by the PPE or SPEs. Outbound write data associated with write accesses is sent on the IOIF in the order in which acknowledgment responses are received from the I/O unit on the IOIF.

To be more specific, processor outbound accesses are performed on the IOIF in a weakly consistent storage order, except as ordered by one of the following rules:

- Accesses performed by a PPE or SPE to the same location beyond the IOIF are performed in a strongly ordered manner on the IOIF. (The "same location" is defined as any byte locations that match.) Thus, the same locations might be accessed even if the initial addresses used for the accesses do not match.

- If a PPE performs two stores to different locations across the same IOIF, and those storage locations are both caching-inhibited and guarded, the stores are performed on the IOIF in program order.
- If an SPE performs two writes to different locations across the same IOIF, if the locations are caching-inhibited, and if the writes are ordered by a **fence** or **barrier**, the writes are performed on the IOIF in program order. The SPE writes can be DMA transfers or send-signal instructions with fence or barrier options (**sndsig*<f,b>***). See the *Cell Broadband Engine Architecture* for details on which writes are ordered by a **fence** or **barrier**, details on tag-specific ordering, and details on SPE read ordering. Because reads are necessarily completed on the IOIF before they can be completed in the CBEA processors, *Cell Broadband Engine Architecture* ordering rules are sufficient to establish read order on the IOIF.
- When a PPE executes a **sync** or **eieio** instruction, a memory barrier is created which orders applicable storage accesses pairwise as defined by the *PowerPC Architecture*. The ordering defined by the *PowerPC Architecture* applies to the order in which the applicable storage accesses are performed on the IOIF.
- When an SPE executes an MFC synchronize (**mfcsync**) or MFC enforce in-order execution of I/O (**mfceieio**) command, a memory barrier is created for DMA transfers within the same DMA tag group[4]. This orders applicable storage accesses pairwise as defined by the *Cell Broadband Engine Architecture* and *PowerPC Architecture*. The ordering defined by these architectures applies to the order in which the applicable storage accesses are performed on the IOIF.
- If an interrupt is transmitted by a PPE or an SPE to a unit on the IOIF, previous loads, stores, or DMA transfers by that PPE or SPE to locations beyond the same IOIF are performed on the IOIF before the interrupt is transmitted on the IOIF.

In general, nonoverlapping accesses caused by different PPE load instructions that specify locations in caching-inhibited storage can be combined into one access. Nonoverlapping accesses caused by separate store instructions that specify locations in caching-inhibited storage can also be combined into one access. Such combining does not occur if the load or store instructions are separated by a **sync** instruction, or by an **eieio** instruction if the storage is also guarded. Such combining does not occur for SPE accesses to locations beyond the IOIF that are ordered by a **barrier**, **mfceieio**, or **mfcsync** command, or by another MFC command that has a barrier or fence option.

For ordering accesses, the IOC makes no distinction between a specific PPE or SPE. All EIB read and write commands to the IOIF are transmitted on the IOIF in the order in which they are performed on the EIB, regardless of the unit initiating the access. If **eieio** instructions are configured to be propagated to an IOIF, all **eieio**, read, and write commands from the EIB are transmitted on the IOIF in the order in which they are performed on the EIB. An I/O device can use **eieio** to determine which read and write commands need to be ordered instead of performing all read and write commands in the order in which they are received.

#### 7.7.2.2 *Inbound (IOIF to EIB) Accesses*

The ordering of inbound storage accesses is determined by the following control bits:

- *IOC TE Bit*—See *Section 7.7.1* on page 188.

4. The *Cell Broadband Engine Architecture* specifies that the **mfcsync** and **mfceieio** commands are tag-specific, but the CBEA processors treat all three barrier command identically, having no tag-specific effects.

- *IOIF S Bit*—This is the strict-ordering bit sent with each command on the IOIF. Strict ordering is required on any read or write when S = '1', unless the SO bits (see the next list item) override this requirement. Weak ordering is permitted when S = '0'.
- *IOPT SO Bits*—These are the storage-ordering bits from the IOPT entry or IOPT-cache entry corresponding to the accessed location, described in *Section 7.3.4* on page 171.

*Table 7-3* summarizes the effects on inbound read and write storage order for the various combinations of control bit. The table assumes that the accesses are from one I/O unit on one virtual channel on one IOIF. Inbound data transfers from the IOIF to the EIB that are associated with these commands are kept in the order in which they are received by the EIB.

*Table 7-3. Ordering of Inbound I/O Accesses from Same or Different I/O Addresses*

| Control Bit Combination[1] | | | Same IOID[2] and VC[3], Same I/O Address | Same IOID and VC, Different I/O Address, Previous Command Was a Write | Same IOID and VC, Different I/O Address, Previous Command Was Not a Write |
|---|---|---|---|---|---|
| TE | S | SO | | | |
| 0 | 0 | x | Ordered | Not Ordered | Not Ordered |
| 0 | 1 | x | Ordered | Ordered | Ordered |
| 1 | 0 | 00 | Ordered | Not Ordered | Not Ordered |
| 1 | 0 | 01 | Undefined (SO = '01' is reserved) | | |
| 1 | 0 | 10 | Ordered | Not Ordered | Not Ordered |
| 1 | 0 | 11 | Ordered | Not Ordered | Not Ordered |
| 1 | 1 | 00 | Ordered | Not Ordered | Not Ordered |
| 1 | 1 | 01 | Undefined (SO = '01' is reserved) | | |
| 1 | 1 | 10 | Ordered | Ordered | Not Ordered |
| 1 | 1 | 11 | Ordered | Ordered | Ordered |

1. An "x" indicates don't care or not applicable.
2. IOID = I/O identifier.
3. VC = I/O virtual channel. See the *Hardware Initialization Guide* for a description of I/O virtual channels.

*Table 7-4* on page 192 differs from *Table 7-3* in that the accesses are (a) only from different I/O addresses, (b) to locations that are either not cached by the PPE or have the memory-coherence-required attribute, and (c) made with no assumption about whether the previous access was or was not a write.

*Table 7-4. Ordering of Inbound I/O Accesses from Different I/O Addresses*

| Control Bit Combination[1] | | | Ordering of I/O Accesses[2] |
|---|---|---|---|
| TE | S | SO | |
| x | 0 | x | Previous reads and writes NOT necessarily performed before this access. |
| 0 | 1 | x | Previous reads and writes performed before this access. |
| 1 | 1 | 00 | Previous reads and writes NOT necessarily performed before this access. |
| 1 | 1 | 01 | Undefined (SO = '01' is reserved). |
| 1 | 1 | 10 | Previous writes performed before this access. |
| 1 | 1 | 11 | Previous reads and writes performed before this access. |

1. An "x" indicates don't care or not applicable.
2. Assumes accesses are reads or writes from different I/O addresses, accesses are from one IOIF bus with the same IOID and virtual channel, and addressed locations either have the storage attribute of memory-coherence-required or are not in a PPE cache.

If either S = '0' or SO = '00', the access on the EIB may be performed out of order relative to a preceding access, unless the address, IOID, and I/O virtual channel (VC) of both accesses match. For incoming I/O addresses that are not translated, the IOIF S bit is used to determine ordering requirements.

Aside from interrupt commands, only read or write commands that have the same IOID and VC are ordered. Interrupts are ordered behind all previous writes from the same IOIF interface. Interrupt-reissue commands are intercepted before they reach the EIB, and therefore have no ordering properties.

I/O accesses to the EIB from an IOIF are ordered based on the implicit storage ordering for real addresses if the I/O address is not translated, or are ordered based on the SO bits if the I/O address is translated. This ordering can be further relaxed if S = '0'.

### *Ordering Rules*

Inbound I/O accesses are unordered relative to the PPE and SPEs, except in the following cases; the SO value shown in these exceptions is either the implicit SO for real addresses, if the I/O address is not translated, or the IOPT SO bits if the I/O address is translated:

- Accesses by a specific I/O unit on one IOIF virtual channel, to the same location in the CBEA processor, are performed (with respect to the PPE and SPEs) in the order in which they are initiated on the IOIF if the location has the storage attribute of memory-coherence-required or if the location is not in a CBEA processor cache. The address is used to determine whether the "same location" that is accessed is the I/O address passed on the IOIF. (The "same location" is defined as any byte locations that match.)

  ***Note:*** *If an I/O device reads from a system memory location using one I/O address, and an I/O device writes to the same system memory location using a different address, these operations need not be kept in order. To have these operations performed in the expected I/O order, the system architect must use the same virtual channel and the same IOID on the IOIF for all I/O devices on an I/O bus with a shared-memory address space, such as PCI Express.*

- An I/O unit write, with corresponding SO = '10' and S = '1', will be performed with respect to the PPE and SPEs after all previous writes by the same I/O unit on the same IOIF virtual

channel have been performed with respect to the PPE and SPEs. This is true only for previous writes to locations that have the storage attribute of memory-coherence-required, or that are not in a PPE cache. This is true regardless of the SO values of the previous writes.

- An I/O unit read or write, with a corresponding SO = '11' and S = '1', will be performed with respect to the PPE and SPEs after all previous reads and writes on the same IOIF virtual channel by the same I/O unit. This is true only for previous accesses to locations that have the storage attribute of memory-coherence-required or that are not in a PPE cache. This is true regardless of the SO values of the previous accesses.
- Data from an I/O read must be allowed to pass on the IOIF before a CBEA-processor access that has been previously queued in the IOC, to prevent a possible deadlock. For example, read data from an SPE's LS must be allowed to pass to the IOIF before a write command that is queued in the IOC waiting for a command credit.
- Data from an I/O write must be allowed to pass on the IOIF to system memory before another I/O access to the same coherency block is performed on the internal system interconnect bus, to prevent a possible deadlock. For example, an I/O device might send two write commands to the SPU Signal Notification Register 1. Until data for the first write is received, the second write might never get a response on the IOIF. After the SPE receives the data for the first write, the second write command gets acknowledged on the IOIF. Thus data for the first write must be sent without waiting for a response for the second write command. A *coherency block* is a collection of memory bytes corresponding to a cache line.
- Typically, software is notified of the completion of an I/O operation by means of an interrupt transmitted on the IOIF after the last I/O read or write required by the I/O operation.
  - Software must ensure that, after the interrupt, software sees the effect of the previous I/O writes performed by the I/O unit assuming that all writes and the interrupt are transmitted on the same IOIF virtual channel. If the I/O subsystem does not ensure that previous I/O writes have occurred on the IOIF before a subsequent interrupt is transmitted on the IOIF, additional system-specific or I/O-subsystem-specific methods might be required to ensure that software sees the effect of previous I/O writes.
  - The I/O unit must not transmit an interrupt on the IOIF until the last I/O read associated with the interrupt has been performed on the IOIF.

The preceding ordering rules do not apply to an access that causes an I/O exception. Other preceding and subsequent accesses are still ordered according to these rules.

Read or write accesses from an IOIF to a CBEA processor are never combined into one access.

*Cache Writebacks*

Before an I/O write is performed to a storage location that does not have the memory-coherence-required attribute, software must ensure that modified versions of the memory location in any PPE cache or SPE atomic cache[5] are written to memory by using PPE cache-management instructions or SPE synchronization instructions or MFC commands. Otherwise, if the cached copy is modified, the modified line might subsequently be written to memory, overlaying the memory copy written by the I/O device. After the modified cache copies are written to memory,

5. Four cache lines are available in each SPE for atomic operations.

software must ensure that a cached copy does not become modified until after the I/O write is performed. If deemed necessary, privileged software can use the following methods to ensure this:

- Software can assign the location the storage attribute of caching-inhibited.
- Software can remove the corresponding valid entry from the page table and translation lookaside buffers (TLBs).
- Software can use the page protection bits to prevent writes to the location.

See the *PowerPC Architecture* for details about caching-inhibited, page table, TLB, and page protection bits.

*Interrupts*

Regardless of the virtual channel and IOID, interrupt commands from an IOIF are not performed on the EIB until all previous writes from the IOIF are successfully performed on the EIB. “Successfully performed” means the writes have obtained a combined response of acknowledged and not retry. This does not guarantee the data transfer for a previous write has completed. However, if the PPE accesses such a location after reading the corresponding interrupt information by means of the MMIO register `IIC_IPP0` or `IIC_IPP1`, this PPE access is performed after the previous IOIF write. Thus, if this PPE access to such a location is a load, the load returns the data written by the previous IOIF write, assuming there is no intervening write to the location.

Because the data transfer for a write from the IOIF is not guaranteed to be completed when the subsequent interrupt is presented to the PPE, and because the SPEs do not operate in the EIB coherence domain, and because the interrupt does not have a cumulative ordering property, if a subsequent SPE access to its LS is a result of a subsequent PPE store, the SPE access can still occur before the data transfer for the write from the IOIF. For example, if the PPE stores to a synergistic processor unit (SPU) signal-notification register after receiving the interrupt, and the SPE accesses LS after the signal notification event, the SPE access might occur before the data transfer for the write from the IOIF. As described in the *Section 20.1.5* on page 577, the MFC multisource synchronization facility can be used to ensure that such previous data transfers have completed.

### 7.7.3 I/O Accesses to Other I/O Units through an IOIF

***Note:*** *The following description conforms to the I/O architecture. However, unlike the I/O architecture, the CBEA processors do not support operations from an IOIF that access the same IOIF.*

An I/O access from an IOIF may access the same or a different IOIF. The IOIF on which the I/O command occurred is called the incoming IOIF, or the IOIF in bus. The IOIF that is the target of the I/O command is called the outgoing IOIF, or the IOIF out bus.

I/O accesses from an IOIF to a different IOIF are ordered based on the implicit SO for real addresses, if the I/O address is not translated, or based on the IOPT SO, if the I/O address is translated. This ordering can be further relaxed if the S bit = ‘0’. I/O accesses from an IOIF to a different IOIF are generally unordered on the outgoing IOIF, except in the following cases:

- Accesses by a specific I/O unit on one IOIF virtual channel that access the same location beyond the outgoing IOIF are performed on the outgoing IOIF in the order they are initiated

on the incoming IOIF. For I/O accesses, the address used to determine whether the “same location” is accessed is the I/O address passed on the IOIF. (The “same location” is defined as any byte locations that match.)

- An I/O-unit write with a corresponding SO = ‘10’ and S = ‘1’ will be performed on the outgoing IOIF after all previous writes by the same I/O unit on the same IOIF virtual channel have been performed on the outgoing IOIF, regardless of the storage attribute of memory-coherence-required or the SO values of the previous writes.
- An I/O-unit read or write with a corresponding SO = ‘11’ and an S = ‘1’ will be performed on the outgoing IOIF after all preceding reads and writes on the same IOIF virtual channel by the same I/O unit, regardless of the storage attribute of memory-coherence-required or the SO values of the previous accesses.
- An I/O-unit interrupt will be performed on the outgoing IOIF after all previous writes by the same I/O unit on the same IOIF virtual channel have been performed on the outgoing IOIF, regardless of the storage attribute of memory-coherence-required or the SO values of the previous writes.
  - If the I/O subsystem does not ensure that previous I/O writes have occurred on the IOIF before a subsequent interrupt packet is transmitted on the IOIF, additional system-specific or I/O subsystem-specific methods might be required to ensure that software in the destination IOIF sees the effect of previous I/O writes.
  - The I/O unit must not transmit an interrupt on the IOIF until the last I/O read associated with the interrupt has been performed on the IOIF.

Read or write accesses from an IOIF to a different IOIF are never combined into one access.

### 7.7.4 Examples

The examples in this section reference combinations of the following control bits:

- *IOIF C Bit*—This is the memory-coherence bit sent with a command on the IOIF. Memory coherence is required on any read or write when C = ‘1’, and memory coherence is not required when C = ‘0’.
- *TLB I Bit*—This is the caching-inhibited bit from the TLB Real Page Number Register (`TLB_RPN`), as described in the *Cell Broadband Engine Registers* specification. Caching is inhibited when I = ‘1’, and caching is permitted when I = ‘0’.
- *IOPT M Bit*—This is the memory-coherence bit from the IOPT entry or IOPT-cache entry corresponding to the accessed location, as described in *Section 7.3.4* on page 171. Memory coherence is required when M = ‘1’, and memory coherence is not required when M = ‘0’.
- *IOIF S Bit*—This is the strict-ordering bit sent with a command on the IOIF. Strict ordering is required on any read or write when S = ‘1’, unless the SO bits (see the next list item) override this requirement. Weak ordering is permitted when S = ‘0’.
- *IOPT SO Bits*—These are the storage-ordering bits from the IOPT entry or IOPT-cache entry corresponding to the accessed location, described in *Section 7.3.4* on page 171.

In the first example that follows, the ordering of an outgoing DMA transfer with respect to a subsequent send-signal is described. The subsequent examples demonstrate the order of commands originating externally on an IOIF relative to commands originating inside the CBEA processor. Unless otherwise stated, locations A and B in these examples refer to memory that is in the processor coherence domain.

#### 7.7.4.1 *Example 1*

Consider the following example of an SPE initiating an outbound DMA transfer to an I/O device, followed by the SPE sending a signal notification to a different location on that I/O device. On all outgoing I/O accesses, the C and S bits are set to '1' by hardware; in this example, the I bit in the TLB is assumed to be '1':

1. DMA transfer to location A in the I/O device.
2. **sndsig <f,b>** to location B in the I/O device.

The write command to location A is performed on the IOIF before the write command to location B is performed on the IOIF. There is no guarantee that the data is transferred on the IOIF by the time that the write command to location B is performed. However, the data is transmitted on the IOIF in the order the commands are acknowledged on the IOIF. In some cases the write to location B might cause some operation to be initiated in an I/O bridge or I/O device, and this operation might have a dependency on the data for location A. In this case, the I/O bridge or I/O device must ensure the data for location A has arrived at the necessary destination point before initiating the operation triggered by the write to location B.

#### 7.7.4.2 *Example 2*

Software must ensure that PPE thread 0 does not have a modified copy of location A in its cache when the write to location A is performed. If PPE thread 0 might have an unmodified copy of location A, then software must execute a **dcbf** instruction to the location to ensure that PPE thread 0 loads data from memory. For IOIF1 virtual channel 3, C and S are the IOIF address-modifier bits; SO and M are bits from the IOPT entries corresponding to the locations A and B. For PPE thread 0, the loads use translated addresses and use M bits from the *PowerPC Architecture* page table entry (PTE) corresponding to locations A and B. M = '1' means the access has a storage attribute of memory-coherence-required.

| IOIF1 Virtual Channel 3 | PPE Thread 0 |
|---|---|
| 1. Write location A (IOPT[M] = '0') | 1. Load location B (PTE[M] = '1') |
| | 2. Any type of serialization facility that orders these loads. |
| 2. Write location B(IOPT[M] = '1', C = '1', SO = '10', S = '1') | 3. Load location A (PTE[M] = '1') |

If PPE thread 0 loads the new value of location B, it must also load the new value of A. On behalf of IOIF1, the IOC performs the write to memory location A before performing the write to location B.

#### 7.7.4.3 *Example 3*

Location A is a memory location that has the caching-inhibited attribute and is therefore not in any cache. For PPE thread 0, the load uses a translated address and uses an I bit from the *PowerPC Architecture* page table entry corresponding to location A. Location A is in the noncoherent range defined by `EIB_LBAR1` and `EIB_LBAMR1` so the M bit used for both the IOIF1 and PPE thread 0 is forced to indicate the location is not coherent.

| IOIF1 Virtual Channel 3 | PPE Thread 0 |
|---|---|
| 1. Write location A (IOPT[M] = ‘0’) | 1. External interrupt occurs |
| 2. Send interrupt | 2. Load location A (IOPT[M] = ‘0’, IOPT[I] = ‘1’) |

After the interrupt, PPE thread 0 loads the new value of A from memory.

#### 7.7.4.4 *Example 4*

For IOIF1 virtual channel 3, C is an IOIF address-modifier bit, M is the bit from the IOPT entry corresponding to locations A. For PPE thread 0, the load uses a translated address and uses an M bit from the *PowerPC Architecture* page table entry corresponding to location A.

| IOIF1 Virtual Channel 3 | PPE Thread 0 |
|---|---|
| 1. Write location A (IOPT[M] = ‘1’, IOPT[C] = ‘1’) | 1. External interrupt occurs |
| 2. Send interrupt | 2. **lwsync** or **sync** |
| | 3. Load location A (IOPT[M] = ‘1’) |

After the interrupt, the **lwsync** or **sync** ensures that if location A was initially in the cache, the cache line invalidation due to the I/O write occurred before the load by PPE thread 0. Thus, PPE thread 0 loads the new value of A from memory.

#### 7.7.4.5 *Example 5*

When IOPT SO equals ‘10’ or '11' and the IOIF S bit is '1' for an I/O write to memory in the CBEA processor, a previous I/O write to memory in the CBEA processor must be performed in order, with respect to loads performed by a PPE.

Typically, I/O writes will not be separated by an **eieio** or **sync** instruction. This requires special consideration in the system-hardware design if cache-invalidate operations due to snoops are delayed. For example, if an I/O write hits a valid cache line in the PPE—which requires the cache entry to be invalidated—and the PPE delays invalidating the line, the PPE must be able to handle any future operations that hit the same line while the invalidation is still pending.

| IOIF1 Virtual Channel 3 IOID 9 | PPE Thread 0 |
|---|---|
| 1. Write location A<br>(M = ‘1’, C = ‘1’, SO = ‘00’, S = ‘0’). | 1. Load location B (M = ‘1’). |
| | 2. Use any synchronization mechanism that causes these loads to be performed in order (see *Section 20 Shared-Storage Synchronization* on page 561). |
| 2. Write location B<br>(M = ‘1’, C = ‘1’, SO = ‘10’, S = ‘1’). | 3. Load location A (M = ‘1’). |

If the PPE loads the new value of location B, it must also load the new value of location A. Suppose location A is in the PPE’s L1 or L2 cache, but location B is not. The PPE loads the new value of B from memory. The PPE must not load the old value of A from its cache. This implies a guarantee that the cache line containing B must be invalidated before the load of A.

Several methods follow that describe using PowerPC assembler code for serializing loads to cacheable, coherent memory. In these methods, r1, r2, r3, and r4 are GPR locations, r2 contains the address of location B, and r1 contains the address of location A.

Method 1 (**lwsync**):

```
lwz r3,0(r2)
lwsync
lwz r4,0(r1)
```

Method 2 (**sync**):

```
lwz r3,0(r2)
sync
lwz r4,0(r1)
```

Method 3 (dependent branch and **isync**):

```
lwz r3,0(r2)
xor. r4,r3,r3
bne $-4       #branch never taken
isync
lwz r4,0(r1)
```

Method 4 (dependent operand address)

```
lwz r3,0(r2)
xor r4,r3,r3
lwz r4,r4,r1
```

#### 7.7.4.6 *Example 6*

| IOIF1 Virtual Channel 3 IOID 9 | PPE Thread 0 |
|---|---|
| 1. Write location A<br>(M = ‘1’, C = ‘1’, SO = ‘11’, S = ‘1’). | 1. Load location B (M = ‘1’). |
| | 2. Use any serialization mechanism that orders these loads. |
| 2. Write location B<br>(M = ‘1’, C = ‘1’, SO = ‘00’).<br>The value of the S bit is irrelevant in this instance. | 3. Load location A (M = ‘1’). |

Even if the PPE loads the new value of location B, it may load either the new or old value of A. The IOC can perform the write to memory location A after performing the write to location B.

#### 7.7.4.7 *Example 7*

| IOIF1 Virtual Channel 3 IOID 9 | PPE Thread 0 |
|---|---|
| 1. Read location A<br>(M = ‘1’, C = ‘1’, SO = ‘00’, S = ‘0’). | 1. Store to location B (M = ‘1’). |
| | 2. Use any serialization mechanism that causes these stores to be performed in order. |
| 2. Read location B<br>(M = ‘1’, C = ‘1’, SO = ‘11’, S = ‘1’). | 3. Store to location A (M = ‘1’). |

If the read of A returns the new value of location A, the read of B must return the new location of B.

Any of the following PowerPC instructions can be used to serialize stores to cacheable, coherent memory: **lwsync**, **eieio**, **sync**.

#### 7.7.4.8 *Example 8*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Write location A<br>(M = ‘1’, C = ‘1’). | 1. DMA from location B (M = ‘1’). |
| | 2. Use any serialization mechanism that causes these DMA transfers to be performed in order (see the *Cell Broadband Engine Architecture*). |
| 2. Write location B<br>(M = ‘1’, C = ‘1’, SO = ‘10’, S = ‘1’). | 3. DMA from location A (M = ‘1’). |

If the SPE 0 DMA gets the new value of location B, its other DMA must get the new value of A.

#### 7.7.4.9 *Example 9*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Write location A in SPE 0 LS<br>(SO = ‘00’, S = ‘0’). | 1. Read channel for SPE 0’s SPU Signal-Notification Register 1. |
| 2. Write SPE 0’s SPU Signal-Notification Register 1<br>(SO = ‘10’, S = ‘1’). | 2. Load from location A in SPE 0 LS. |

If SPE 0 loads the new value of Signal-Notification Register 1, it must also load the new value of A.

#### 7.7.4.10 *Example 10*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Write location A<br>(M = ‘0’, C = ‘0’, SO = ‘00’, S = ‘0’). | 1. Read channel for SPE 0’s SPU Signal-Notification Register 1. |
| 2. Write SPE 0’s SPU Signal-Notification Register 1<br>(M = ‘0’, C = ‘0’, SO = ‘10’, S = ‘1’). | 2. DMA from location A. |

Location A is assumed to be in a caching-inhibited page in memory. Thus, there is no copy of memory location A in any cache when the write to location A occurs. If SPE 0 loads the new value of Signal-Notification Register 1, SPE 0 must also load the new value of A.

#### 7.7.4.11 *Example 11*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Write location A<br>(M = ‘1’, C = ‘1’, SO = ‘11’, S = ‘1’). | 1. DMA from location B (M = ‘1’). |
| | 2. Use any serialization mechanism that causes these DMA transfers to be performed in order. |
| 2. Write location B<br>(M = ‘1’, C = ‘1’, SO = ‘00’). | 3. DMA from location A (M = ‘1’). |

Even if SPE 0 DMA-transfers the new value of location B, it may transfer either the new or old value of A. The IOC may perform the write to memory location A after performing the write to location B.

#### 7.7.4.12 *Example 12*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Read location A<br>(M = ‘1’, C = ‘1’, SO = ‘00’, S = ‘0’). | 1. DMA write to location B (M = ‘1’). |
| | 2. Use any serialization mechanism that causes these DMA transfers to be performed in order. |
| 2. Read location B<br>(M = ‘1’, C = ‘1’, SO = ‘11’, S = ‘1’). | 3. DMA write to location A (M = ‘1’). |

If the read of A returns the new value of location A, the read of B must return the new value of B.

#### 7.7.4.13 *Example 13*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Write location A in SPE 0 LS<br>(M = ‘0’, C = ‘0’, SO = ‘00’, S = ‘0’). | 1. Load from location B in SPE 0 LS. |
| 2. Write location B in SPE 0 LS<br>(M = ‘0’, C = ‘0’, SO = ‘10’, S = ‘1’). | 2. Load from location A in SPE 0 LS. |

If SPE 0 loads the new value of location B, it must also load the new value of A.

#### 7.7.4.14 *Example 14*

| IOIF1 Virtual Channel 3 IOID 9 | SPE 0 |
|---|---|
| 1. Read location A in SPE 0 LS<br>(M = '0', C = '0', SO = '00', S = '0'). | 1. Store to location B in SPE 0 LS. |
| 2. Read location B in SPE 0 LS<br>(M = '0', C = '0', SO = '11', S = '1'). | 2. **sync** |
| | 3. Store to location A in SPE 0 LS. |

If the read of A returns the new value of location A, the read of B must return the new value of B.

#### 7.7.4.15 *Example 15*

A value of x'1' in location B indicates to an I/O device when to perform an I/O operation. The I/O device polls this location. Initially, location B has a '0' value.

| PPE | I/O Device |
|---|---|
| 1. MMIO store to location A in I/O device. | 1. Read from location B.<br>If the I/O device received a new data value from location B, it performs step 2. |
| 2. **sync** | |
| 3. Store x'1' to location B. | 2. Perform I/O operation, assuming it has new data for A. |

There is no guarantee that the I/O device will receive the most recent data for location A. There is no guaranteed ordering of reads from one end of the IOIF versus writes from the other end of the IOIF.

# 8. Resource Allocation Management

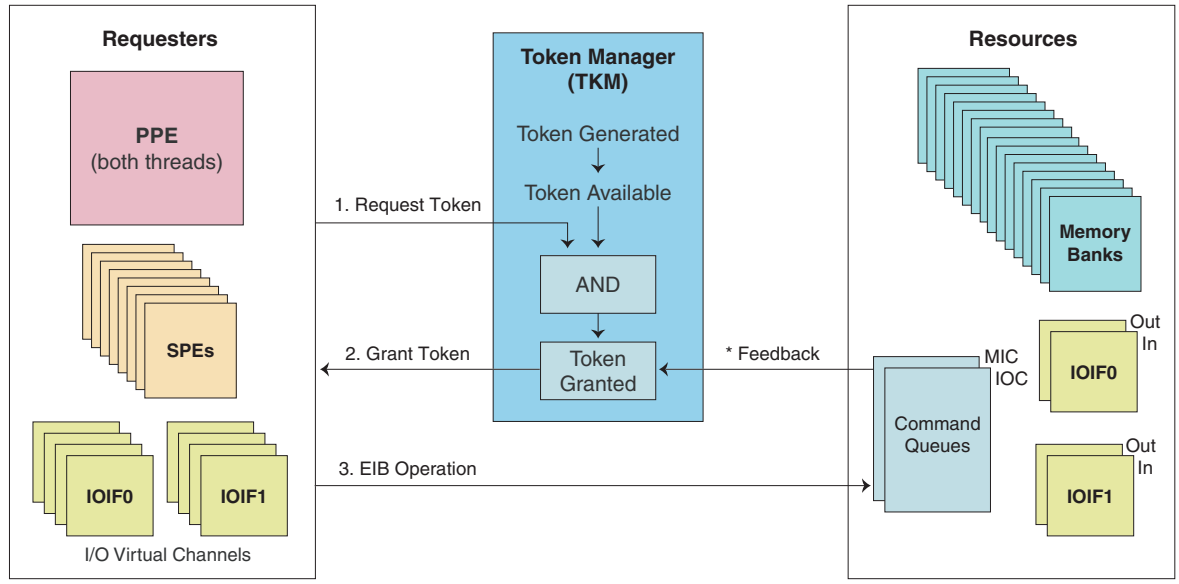
## 8.1 Introduction

The Cell Broadband Engine Architecture (CBEA) processors[1] resource allocation management (RAM) facility can allocate time portions of managed *resources* (memory banks and I/O interfaces) to *requesters* (PowerPC Processor Element [PPE], Synergistic Processor Elements [SPEs], or I/O devices). The requesters can be assigned to one of four resource allocation groups (RAGs). A portion of a resource's time is represented by a *token*. The RAM facility's token manager (TKM) manages the amount of each resource's time that is allocated to each RAG by controlling the generation of tokens for a RAG at a programmable rate and the granting of generated tokens to a requester. To access a resource, a requester must first acquire a corresponding token from the TKM.

*Figure 8-1* on page 204 shows an overview of the RAM facility. Here, I/O resources are represented by the I/O interface controller (IOC) command queue that supports the two I/O interfaces (IOIFs), and memory resources are represented by the memory interface controller (MIC) command queue. Before performing a read or write operation on the element interconnect bus (EIB) that accesses a managed resource, requesters request a token for that particular resource from the TKM. The TKM generates tokens and can retain a limited number of available tokens for brief amount of time. If a token request arrives when the corresponding token is available, the token is granted, unless (for SPE and I/O requesters only) there is feedback from the MIC (for a memory bank token) or IOC (for an IOIF token) command queue to block the grant. After a token is granted, the requester can access the resource on the EIB.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

*Figure 8-1. Resource Allocation Overview*

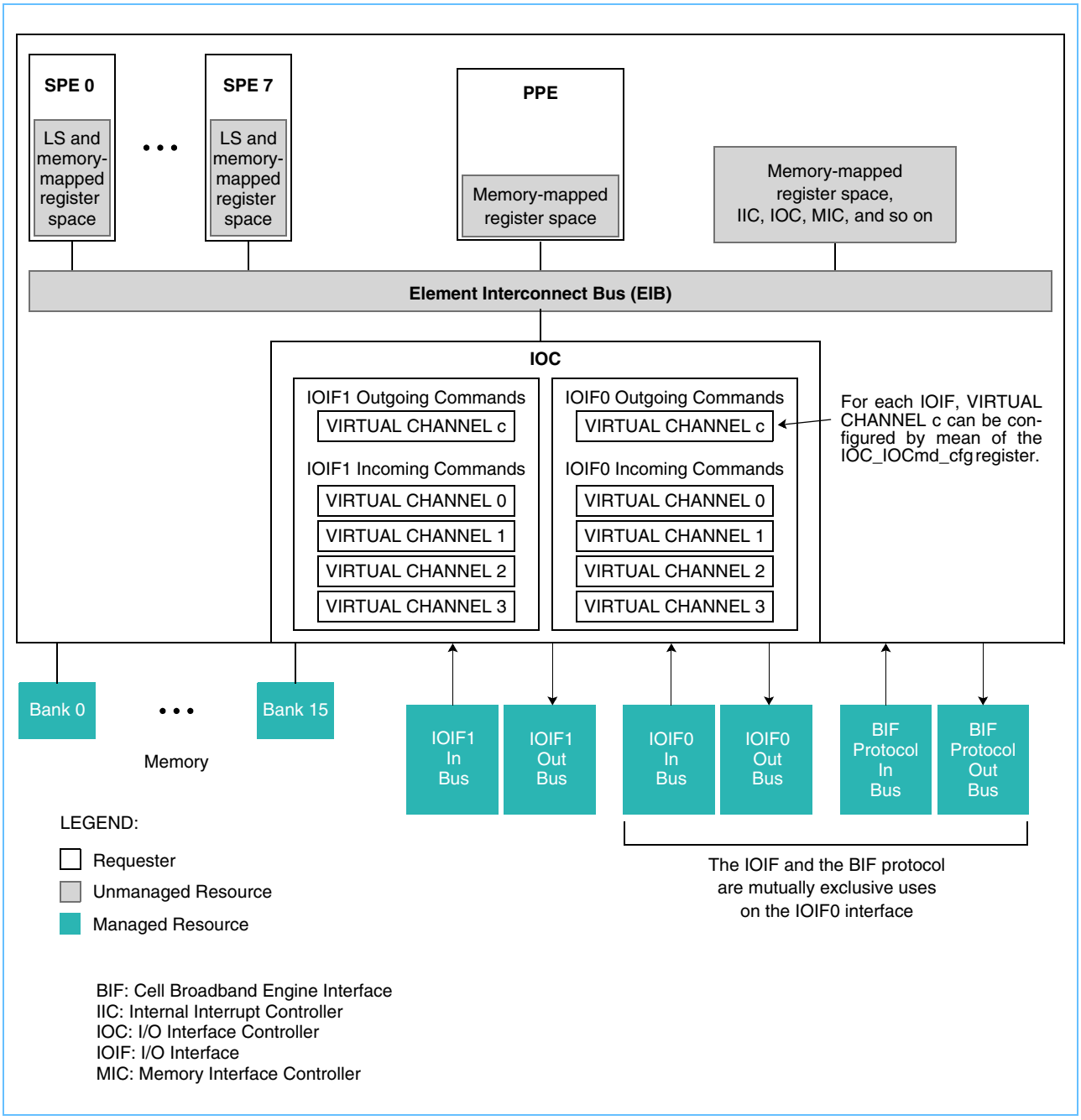



* Feedback from the IOC command queue only occurs when the IOIF (not the BIF) protocol is used.

| | | | |
|---|---|---|---|
| IOC | I/O Interface Controller | PPE | PowerPC Processor Element |
| IOIF | I/O Interface | SPE | Synergistic Processor Element |
| MIC | Memory Interface Controller | TKM | Token Manager |

PPE token requests are somewhat different than shown in *Figure 8-1*, as described in *Section 8.4* on page 208. *Figure 8-2* on page 205 shows a more detailed view of the requesters and resources. The SPE local storages (LSs), EIB, and the processor internal arrays and registers in memory-mapped address spaces are unmanaged resources.

*Figure 8-2. Managed Resources and Requesters*



## 8.2 Requesters

The resource allocation groups (RAGs) are groups of zero or more *requesters*. These requesters are physical or virtual units that can initiate load or store requests or DMA read or write accesses. There are 17 requesters, as shown in *Figure 8-2* on page 205:

- One PPE (both threads)
- Eight SPEs.
- Four virtual I/O channels for physical interface IOIF0
- Four virtual I/O channels for physical interface IOIF1

*Table 8-5* on page 231 lists the configuration registers used to assign requesters to RAGs.

### 8.2.1 PPE and SPEs

When resource allocation is enabled, each requester belongs to one and only one RAG at any given time. However, the RAG assigned to a requester can be changed by software over the course of time. Each RAG is identified by its resource allocation group ID (RAID). The PPE and each SPE are assigned a RAID by software, as described in *Section 8.6* on page 229 and *Section 8.7* on page 233.

### 8.2.2 I/O

Each IOIF interface has four virtual channels, one for each of the four RAGs. For each IOIF, RAM virtual channel `r` is associated with RAID `r`, for `r` = 0 to 3. An external bridge or I/O device must be able to identify the virtual channel associated with each physical transfer on the IOIF. A virtual channel can be used for both reads and writes. This information is passed to the CBEA processor with the address and command. The IOC acts as the agent on behalf of the virtual channels associated with incoming commands. The static or dynamic assignment of an I/O device to a RAG is the responsibility of the external bridge or I/O device.

#### 8.2.2.1 *Virtual Channel Configuration*

As described in *Cell Broadband Engine Registers* specification, the memory-mapped I/O (MMIO) IOCmd Configuration Register(`IOC_IOCmd_Cfg`) defines which outgoing virtual channel is used for outgoing accesses—that is, read or write commands to the IOIF from the PPE, an SPE, or the other IOIF interface. The `IOC_IOCmd_Cfg[CVCID0]` field specifies the virtual channel for outgoing commands on IOIF0 and the `IOC_IOCmd_Cfg[CVCID1]` field specifies the virtual channel for outgoing commands on IOIF1. The outgoing virtual channel has no significance with respect to the RAM unit in the CBEA processor where the command originated. However, if two CBEA processors are connected by means of an IOIF, the outgoing virtual channel on one CBEA processor becomes an incoming virtual channel on the other CBEA processor. The incoming virtual channel is associated with a RAG in the CBEA processor that is the target of the incoming command.

For an access from one IOIF to another IOIF, the virtual channel used for the incoming access on one IOIF is independent of the virtual channel used for the outgoing command on the other IOIF.

#### 8.2.2.2 *Interchannel Delay*

Delays occur when outbound and inbound requests are competing for the same resource. For example, accesses from two virtual channels might contend for the same memory bank token; because there is only one physical link, accesses on the link are serialized even if they are to different virtual channels.

Accesses on different virtual channels might delay each other due to:

- EIB-address or data contention
- SPE LS contention
- A miss in the I/O-address translation cache for one virtual channel
- IOIF contention

The last type of contention (IOIF) should be minimal because the IOIF resource is managed by the RAM facility.

## 8.3 Managed Resources

As shown in *Figure 8-2* on page 205, the RAM facility can manage up to the following 20 resources:

- *Memory banks 0 through 15*—For the PPE, both L2-cache and noncacheable unit (NCU) loads and stores.
- *IOIF0 In Bus*—IOIF or Cell Broadband Engine interface (BIF) protocol. For the IOIF protocol on the PPE, only noncacheable loads; for BIF protocol on the PPE, both cacheable and non-cacheable loads.
- *IOIF0 Out Bus*—IOIF or BIF protocol. For the IOIF protocol on the PPE, only noncacheable stores; for BIF protocol on the PPE, both cacheable and noncacheable stores.
- *IOIF1 In Bus*—IOIF protocol. For the PPE, only noncacheable loads.
- *IOIF1 Out Bus*—IOIF protocol. For the PPE, only noncacheable stores.

The RAM facility allocates resources' time over long time periods. As the time period decreases, there is less certainty about allocation effects. RAM is not a solution for queueing effects when resources are highly used—in which case, latency increases. Although RAM is expected to be used for bandwidth management, RAM does not guarantee that the RAGs will get the allocated bandwidths.

RAM allocates the time that a resource is available for a RAG, without regard to whether the allocated time is actually used for requested data transfers. RAM allocates the time of resources, but does not guarantee the use by a RAG of the allocated time if this allocation exceeds the capabilities of units in the RAG or exceeds the use of the resource by the software. It is up to software to use the time of a resource in an optimal manner to achieve a required bandwidth.

The following types of actions might help achieve optimal bandwidth:

- Writing 128 bytes in a naturally aligned 128-byte block.
- Sequential accesses that read or write equal amounts of data to all memory banks.
- Limited use of the caching-inhibited storage attribute for memory.

- Infrequent translation lookaside buffer (TLB) misses; that is, maintain locality in loads, stores, and DMA transfers, or pin the TLB entries using the replacement management table (RMT).
- Updating a cache line before it is displaced from the PPE L2 cache.
- Limiting the sharing of data among the PPE and SPEs or I/O devices.

## 8.4 Tokens

Managed resources are allocated on the basis of portions of time, represented by tokens. Before using a managed resource, a requester obtains a token for the resource, as shown in *Figure 8-1* on page 204. One token is obtained for every read or write access, of up to 128 bytes, that is wholly contained within a naturally aligned 128 bytes. Tokens are consumed when the command that required the token has been successfully transferred on the EIB.

A requester may have only one outstanding token requests per resource, with the exception that an SPE and IOC may have two outstanding token requests for the IOIF0 In and Out buses. The token manager (TKM) grants tokens at a programmable rate for each RAG.

### 8.4.1 Tokens Required for Single-CBEA-Processor Systems

A requester might need multiple tokens for a single data transfer, one for each resource that its access uses. For example, an I/O device on the IOIF0 interface needs two tokens to transfer data to memory, one for the specific memory bank stored to and one for the IOIF0 In data bus. *Table 8-1* shows combinations of tokens that must be obtained for operations on the EIB and IOIF interfaces in a single-CBEA-processor system.

*Table 8-1. Tokens Needed for EIB or IOIF Operations in Single-CBEA-Processor Systems*
(Sheet 1 of 3)

| Requester and Operation | Tokens Required | | | | |
|---|---|---|---|---|---|
| | Target Memory Bank (1 of 16) | IOIF0 In | IOIF0 Out | IOIF1 In | IOIF1 Out |
| **PPE** | | | | | |
| Memory write of 128 bytes | 1 | | | | |
| Memory write of 1 to 16 bytes | 1 or 2[2] | | | | |
| **PPE or SPE** | | | | | |
| Read from memory of 1 to 128 bytes | 1 | | | | |

1. The IOC requests memory-write tokens according to the following modes, configured by means of `IOC_IOCmd_Cfg`:
   – One token regardless of the write length (1 to 128 bytes)
   – Two tokens for lengths less than 16 bytes and one token for 16 to 128 bytes (mode for memory error-correcting code [ECC] and Extreme Data Rate [XDR] DRAMs with write-masked)
   – Two tokens for lengths less than 128 bytes and one token for 128 bytes (mode for XDR DRAMs that does not support write-masked)
2. For caching-inhibited writes to memory, the PPE requests 1 or 2 tokens, as specified by the bus interface unit (BIU) configuration ring bit “2 Token Decode for NCU Store” at power-on reset (POR). If this bit is ‘0’, one token is obtained; if this bit is ‘1’, two tokens are obtained. Also see the `IOC_IOCmd_Cfg[SXT]` bit.

*Table 8-1. Tokens Needed for EIB or IOIF Operations in Single-CBEA-Processor Systems*
(Sheet 2 of 3)

| Requester and Operation | Tokens Required | | | | |
|---|---|---|---|---|---|
| | Target Memory Bank (1 of 16) | IOIF0 In | IOIF0 Out | IOIF1 In | IOIF1 Out |
| 1 to 128 byte read from I/O bridge or I/O device attached to IOIF0 | | 1 | | | |
| 1 to 128 byte write to I/O bridge or I/O device attached to IOIF0 | | | 1 | | |
| 1 to 128 byte read from I/O bridge or I/O device attached to IOIF1 | | | | 1 | |
| 1 to 128 byte write to I/O bridge or I/O device attached to IOIF1 | | | | | 1 |
| **SPE** | | | | | |
| Memory write of 16 to 128 bytes | 1 | | | | |
| Memory write of 1 to 8 bytes | 2 | | | | |
| **IOC as Agent for I/O Device on IOIF0** | | | | | |
| Write to LS or CBEA processor memory-mapped registers | | 1 | | | |
| Read from LS or CBEA processor memory-mapped registers | | | 1 | | |
| Write to memory of 1 to 16, 32, 48, 64, 80, 96, or 112 bytes | 1 or 2[1] | 1 | | | |
| Write to memory of 128 bytes | 1 | 1 | | | |
| Read of 1 to 128 bytes that does not cross a 128-byte boundary | 1 | | 1 | | |
| Write to I/O bridge or I/O device attached to IOIF1 | | 1 | | | 1 |
| Read from I/O bridge or I/O device attached to IOIF1 | | | 1 | 1 | |
| Read or write to I/O bridge or I/O device attached to IOIF0 | | 1 | 1 | | |
| **IOC as Agent for I/O Device on IOIF1** | | | | | |
| Write to LS or CBEA processor memory-mapped registers | | | | 1 | |
| Read from LS or CBEA processor memory-mapped registers | | | | | 1 |
| Write to memory of 1 to 16, 32, 48, 64, 80, 96, or 112 bytes | 1 or 2[1] | | | 1 | |
| Write to memory of 128 bytes | 1 | | | 1 | |
| Read from memory of 1 to 128 bytes that does not cross a 128-byte boundary | 1 | | | | 1 |

1. The IOC requests memory-write tokens according to the following modes, configured by means of `IOC_IOCmd_Cfg`:
   – One token regardless of the write length (1 to 128 bytes)
   – Two tokens for lengths less than 16 bytes and one token for 16 to 128 bytes (mode for memory error-correcting code [ECC] and Extreme Data Rate [XDR] DRAMs with write-masked)
   – Two tokens for lengths less than 128 bytes and one token for 128 bytes (mode for XDR DRAMs that does not support write-masked)
2. For caching-inhibited writes to memory, the PPE requests 1 or 2 tokens, as specified by the bus interface unit (BIU) configuration ring bit “2 Token Decode for NCU Store” at power-on reset (POR). If this bit is ‘0’, one token is obtained; if this bit is ‘1’, two tokens are obtained. Also see the `IOC_IOCmd_Cfg[SXT]` bit.

*Table 8-1. Tokens Needed for EIB or IOIF Operations in Single-CBEA-Processor Systems*
(Sheet 3 of 3)

| Requester and Operation | Tokens Required | | | | |
|---|---|---|---|---|---|
| | Target Memory Bank (1 of 16) | IOIF0 In | IOIF0 Out | IOIF1 In | IOIF1 Out |
| Write to I/O bridge or I/O device attached to IOIF0 | | | 1 | 1 | |
| Read from I/O bridge or I/O device attached to IOIF0 | | 1 | | | 1 |
| Read or write to I/O bridge or I/O device attached to IOIF1 | | | | 1 | 1 |

1. The IOC requests memory-write tokens according to the following modes, configured by means of `IOC_IOCmd_Cfg`:
   – One token regardless of the write length (1 to 128 bytes)
   – Two tokens for lengths less than 16 bytes and one token for 16 to 128 bytes (mode for memory error-correcting code [ECC] and Extreme Data Rate [XDR] DRAMs with write-masked)
   – Two tokens for lengths less than 128 bytes and one token for 128 bytes (mode for XDR DRAMs that does not support write-masked)
2. For caching-inhibited writes to memory, the PPE requests 1 or 2 tokens, as specified by the bus interface unit (BIU) configuration ring bit “2 Token Decode for NCU Store” at power-on reset (POR). If this bit is ‘0’, one token is obtained; if this bit is ‘1’, two tokens are obtained. Also see the `IOC_IOCmd_Cfg[SXT]` bit.

#### 8.4.1.1 *PPE Access Considerations*

The PPE has a cache with 128-byte cache lines. When software executes a load, store, **dcbt**, or **dcbtst** that is not caching-inhibited, the PPE only performs a memory read operation on the EIB if an L2 miss occurs. Any L2 miss can also cause a castout of a modified cache line, which results in a 128-byte memory write. These L2-miss memory read and write operations require a 128-byte EIB operation.

An instruction fetch can also cause an L2 miss. A load, store, or instruction fetch for which the address is translated might have a TLB miss. An L2 miss can also occur if hardware sets the page-table Change bit for a store or searches the page table for the TLB miss and the page-table entry is not in the L2 cache. A castout can also occur due to a **dcbf** instruction.

When software executes a load or a store that is caching-inhibited, the PPE does an EIB operation to memory, to an IOIF, or to the MMIO register space, depending on the destination address. The length of the load or store operation is defined by the instruction.

When software executes a store that is caching-inhibited and not guarded, the PPE can combine multiple stores to the same quadword into one store operation on the EIB. The length of this store operation can be 1, 2, 4, 8, or 16 bytes, aligned to the operand size, depending on the number of stores combined.

#### 8.4.1.2 *SPE Access Considerations*

An SPE maintains a 6-cache-line buffer (the memory flow controller’s [MFC’s] atomic cache). Some buffer entries are used for atomic commands and some for page-table accesses. Page-table accesses are needed when a page-table Change bit must be set or a hardware page-table search is needed for a TLB miss. When the target of an atomic command or page-table access is

in this buffer, there is no read operation on the EIB. But when the target is not in this buffer, a 128-byte memory read operation occurs on the EIB. When this buffer-miss occurs, a castout of another buffer entry might occur, resulting in a 128-byte write operation on the EIB.

MFC DMA operations might need to be split into multiple EIB operations because the EIB only supports operations of 1 to 15 bytes without crossing a quadword boundary, and it only supports quadword integer multiples without crossing a 128-byte boundary.

#### 8.4.1.3 *I/O Access Considerations*

An IOIF device can initiate read and write operations on the IOIF interface. All real-mode IOIF reads and writes and successfully translated IOIF reads and writes result in EIB operations. The IOIF device operations might need to be split into multiple IOIF operations because the IOIF interface and the EIB only support operations of 1 to 15 bytes without crossing a quadword boundary and integer multiples of quadwords without crossing a 128-byte boundary.

#### 8.4.1.4 *Multiple Token Requests*

Some operations at the software level or I/O-device level might result in multiple token requests. For example:

- Writes to memory that are less than 16 bytes in size require a read-modify-write operation in the MIC, and this results in two token requests from the requester—one for read and one for write.
- SPE 128-byte DMA writes to aligned caching-inhibited storage memory locations using a translated address that gets a TLB miss, which is handled by a search of the hardware page table. In this case, one or two tokens are needed for the page-table access (one for the primary hash and, potentially, one for the secondary hash) and one token is needed for the memory write.
- PPE writes to memory that experience an instruction TLB miss (up to 2 tokens), an ICache or L2 miss of the instruction (1 token), a write-address data TLB miss (up to 2 tokens), and a read (1 token). In this case, up to 6 tokens are needed. Read latency becomes long when all memory accesses are to the same memory bank.

Because the TLB miss is not known when the access is initially attempted, and the final target resource of a DMA transfer is not known until the translation is successful, the tokens are obtained one at a time, as needed, until the real address is known. When the real address is known, at least one token is obtained, per resource used by an EIB operation, before initiating the access on the EIB.

#### 8.4.1.5 *PPE Replacement Tokens*

PPE token requests are treated as high-priority requests by the TKM. With the exception of the PPE, tokens are not requested before the requester has a specific access to make. When the CBEA processor is powered up, the PPE resets to a state where it has one token for each of the 20 possible resources at boot time. The PPE requests a replacement token in parallel with the EIB access that uses the same resource. This allows the PPE to bank one token for each resource before there is a specific access that needs the token. This reduces latency for TLB misses and L2 cache misses. If the access requires two memory tokens, the PPE might need to obtain the second token before performing the operation.

In the case of a TLB miss that requires a search of the hardware page table using the secondary hash, the accesses to the primary and secondary page-table entry group (PTEG) might occur immediately if the PPE already has tokens for these or if one of the PTEGs is in the L2 cache. However, if the page-table accesses for both hashes are to the same memory bank, the second access can only be made after the PPE has replaced the token used on the first access.

If the target is a managed resource, tokens are required for the following EIB operations: read, read with intent to modify, read with no intent to cache, write, write with kill, write with clean, and write with flush (see the *Glossary* on page 835). A memory token is consumed even though the EIB operation was satisfied by intervention and the memory bank might not actually be accessed. If speculative-read mode is enabled (`MIC_Ctl_Cnfg_0[1]` = '0'), a memory read might occur even though intervention occurs or the EIB command receives a retry response.

### 8.4.2 Operations Requiring No Token

In general, a token must be obtained by a requester for every use of memory. However, there are some exceptions to this rule. The amount of total resource time allocated should be less than the theoretical maximum resource time to allow for accesses that do not request tokens.

Operations that do not require tokens include:

- L2 cache writes back a modified cache line to memory in response to another requester doing one of the following operations:
  - EIB Write with Flush operation that hits a modified cache line (partial cache-line write due to small SPE DMA or a small or unaligned I/O device access to memory)
  - EIB reads of fewer than 128 bytes that hit a modified cache line (partial cache-line reads due to SPE DMA or I/O device access to memory; intervention on the EIB is not supported for these)
  - EIB Flush (**dcbf**) operation that hits a modified cache line
  - EIB Clean (**dcbst**) operation that hits a modified cache line

  For the preceding cases, the L2 cache writes back the modified cache line to memory without obtaining a token because there is no facility in the EIB for the PPE to obtain the RAID of the requester causing the cache line to be written back to memory. If software and I/O devices infrequently cause these situations, the effect on the accuracy of the RAM facility is small.
- The IOC performs memory accesses for I/O-address translation. If software carefully manages the I/O segment table (IOST) cache and I/O page table (IOPT) cache either by using the hint bits or by preloading these caches, the number of misses might be small and the effect introduced by this exception can be minimized.
- DMA transfer between the LSs of two SPEs (the LS is an unmanaged resource).
- Accesses to memory-mapped registers within the CBEA processor.
- All EIB requests that do not require a data transfer, including **sync**, **tlbie**, **tlbsync**, **eieio**, and **larx**.

### 8.4.3 Tokens Required for Multi-CBEA-Processor Systems

The RAM facility is supported in a limited manner for multiple CBEA processors connected through the IOIF0 interface using the BIF protocol. In a multi-CBEA-processor system, the use of resources in or attached to the CBEA processor (also referred to as the *local CBEA processor*) are managed by the TKM within that CBEA processor for accesses by requesters within that CBEA processor. The use of memory and IOIF1 resources in or attached to the CBEA processor by requesters in the other CBEA processor are not directly managed. Such resources use can only be indirectly managed by managing the use of the BIF protocol. Resource allocations for IOIF0 effectively become allocations for the BIF protocol, although configuration bits in the PPE, SPEs, and IOC can be set so these units do not depend on obtaining tokens.

The PPE, SPEs, and IOC have configuration registers that recognize the real addresses corresponding to:

- Memory in the local CBEA processor
- IOIF1 attached to the local CBEA processor
- 8 MB MMIO register space within the local CBEA processor

Addresses associated with IOIF0 or the BIF protocol are addresses that do not correspond to one of these three local spaces. In a multi-CBEA-processor environment, this allows all addresses associated with resources in or attached to the other the CBEA processor to be identified as using the BIF protocol. Because the BIF-protocol IOIF resource is used by these accesses, by the command and response phase of any coherent read or write, and by command-only transactions, the availability of the IOIF resource for read and write data cannot be definitively stated. However, an IOIF0 token is required for BIF-protocol reads and writes that might cause data to be transferred.

## 8.5 Token Manager

The token manager (TKM) is shown in *Figure 8-1* on page 204 It receives token requests, generates tokens at a software-programmable rate, and grants tokens to requesters when the tokens are available. In each clock cycle, each physical requester might request a token. Token requests are latched by the TKM and potentially granted at a later time.

### 8.5.1 Request Tracking

The TKM tracks 17 virtual requesters (*Section 8.2* on page 206), each with up to 20 resource targets. The TKM also tracks the priority of requests for the SPEs. For PPE and SPE requests, the TKM records the RAID for the RAG. Internally, the TKM keeps track of requests by RAG. For each RAG, the TKM records nine high-priority requests (PPE and SPE 0:7) for each resource and 10 low-priority requests (SPE 0:7 and IOC 0:1) for each resource. The TKM allows both IOIF and memory token requests on both IOC0 and IOC1.The TKM uses round-robin pointers for selecting among the high-priority and low-priority requests within each RAG. For more information, see *Section 8.5.5.4 Memory-Token Grant Algorithm* on page 219 and *Section 8.5.6 I/O Tokens* on page 220.

### 8.5.2 Token Granting

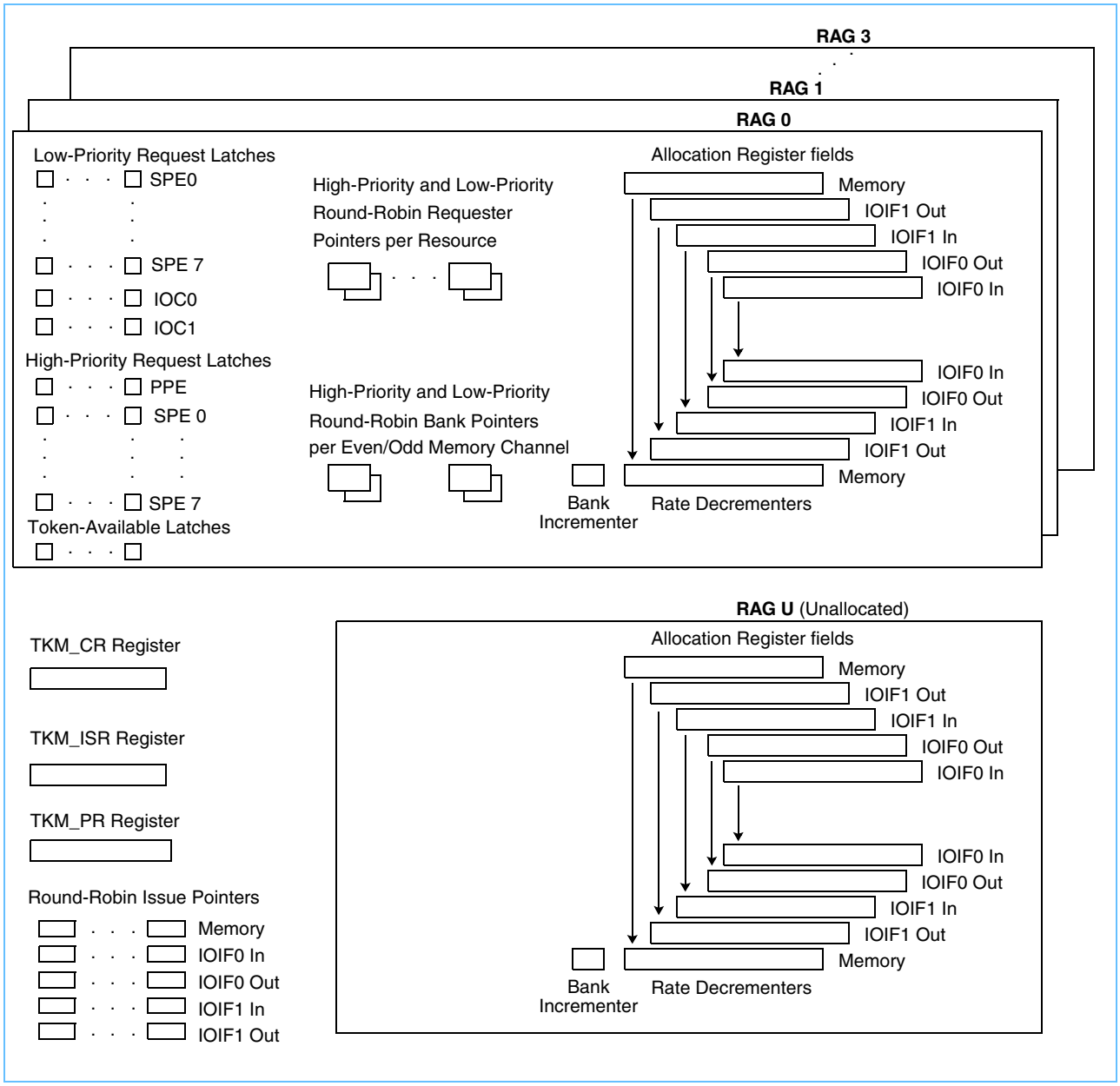
*Figure 8-3* on page 215 shows the organization of the TKM with respect to token granting. For each RAG, the TKM uses Allocation Registers and rate decrementers to determine how often and exactly when to grant tokens, and it uses Token-Available Latches to record when tokens are available for the RAG. It has round-robin pointers per priority, per resource, and per RAG to track which of multiple requesters in a RAG should be granted a token. For tokens that are not used by a RAG and that may be granted to other requesters in other RAGs—the Unallocated RAG (RAG U) described in *Section 8.5.3* on page 215—the TKM Priority Register (TKM_PR) defines the priority of other RAGs in obtaining these unused available tokens.

TKM_CR[44:47] enables the granting of tokens to RAG 0:3, respectively. When the bit for a RAG is disabled, no tokens are granted for requesters assigned to the RAG. If a particular RAG is not used, software can disable the RAG to save power by clearing the RAG-enable in TKM_CR[44:47]. However, disabling the RAG does not prevent token-generation within the RAG; token generation is disabled by disabling the rate decrementer, which is done by clearing the appropriate bit in TKM_CR[55:58] for the required RAG. TKM_CR[54] controls token generation for RAG U (see *Section 8.5.3* on page 215). Any token generated by a RAG but not granted to a requester can eventually be classified as unused, as described in *Section 8.5.7 Unused Tokens* on page 220. Unused tokens can be shared and granted to another RAG if enabled by TKM_CR[51].

To grant tokens at a programmable rate, the Allocation Registers define the period of time between token-generation for a RAG. The TKM uses the TKM Memory Bank Allocation (TKM_MBAR), TKM IOIF0 Allocation (TKM_IOIF0_AR), and TKM IOIF1 Allocation (TKM_IOIF1_AR) MMIO registers for this purpose. The rate decrementers track the periods of time. Rate decrementers cannot be directly loaded by software. These decrementers are decremented at a rate specific to the resource and at a rate established by a programmable prescaler specified by an Allocation Register. When a rate decrementer is ‘0’ and is to be decremented again, the rate decrementer expires. Not only must the rate decrementer reach ‘0’, it also must be decremented again, at which time it is loaded again, based on the allocation.

When a rate decrementer for memory expires, a token becomes available for the associated RAG and the 3-bit interval value from the appropriate Allocation Register field is loaded into the rate decrementer. The most-significant bit of the interval value loaded into a rate decrementer is an implied '1', with the following exception for memory: when TKM_MBAR[33] = ‘0’ and a RAG's prescaler value for memory is '0', a '0' is loaded into the RAG's rate decrementer for memory.

*Figure 8-3. Token Manager*




### 8.5.3 Unallocated RAG

In addition to RAGs 0:3, there is a RAG called Unallocated, or RAG U. There are no requesters in this RAG. Tokens are generated for this RAG using a rate decrementer and an Allocation Register similar to other RAGs. Because there are no RAG U requesters, all tokens generated for RAG U become unused tokens that can be granted to other requesters in other RAGs. The `TKM_PR` register defines the other RAGs' priority in obtaining Unallocated tokens. See *Section 8.5.7 Unused Tokens* on page 220 for an explanation of the unused tokens.

RAG U can take care of occasional spikes in token request rates—for example, if RAG 0 typically needs only 10% of memory bandwidth but occasionally needs 20% for a brief time. In such a case, it might be best to allocate only 10% to RAG 0, but allow RAG U tokens be given to RAG 0 to cover the occasional difference. RAG U is also a buffer that is taken from circulation because the MIC queues might fill up due to random-access patterns from RAGs. Unallocated tokens are sensitive to feedback that represents the use of a resource. See *Section 8.5.12 Feedback from Resources to Token Manager* on page 228. RAG U can also be used to allocate resource time that is more likely to be used by multiple RAGs.

An Unallocated token is only given to a RAG that has an outstanding request at the instant the token is generated, whereas a token generated for RAG 0:3 is made available to the RAG for which it is generated if there is no such token already available for the same resource.

### 8.5.4 High-Priority Token Requests

SPEs indicate priority along with their token requests. High-priority requests by the SPE are only used for memory accesses. All PPE token requests have an implicit high priority, even though the PPE does not drive the high-priority signal. PPE token requests have high priority for IOIF0 In and Out, IOIF1 In and Out, and memory tokens. SPE high-priority token requests are only enabled if the Priority Enable bit, `TKM_CR[PE]`, is '1'. However, if `TKM_CR[PE]` is '0', all SPE token requests are treated as low-priority.

Except for `TKM_CR[PE]`, software has no control on high-priority requests. Instead, hardware controls which token requests are high-priority. If there are multiple high-priority requests for a resource from requesters in a RAG, their requests are granted tokens on a round-robin basis. Low-priority requests from requesters in a RAG are only granted tokens when there are no high-priority requests from requesters in the RAG for the same resource. If there are no high-priority requests from requesters in a RAG for a resource and there are multiple low-priority requests from requesters in the RAG for the resource, their requests are granted tokens on a round-robin basis.

An SPE drives the priority signal high when it requests a token for any of the following reasons:

- The **getllar**, **putllc**, or **putlluc** MFC atomic update commands
- A TLB miss
- A memory write of a modified cache line

If a requester already has a low-priority token request for a resource outstanding, and it has another access to the same resource that warrants high priority, the requester cannot issue another token request for the same resource.

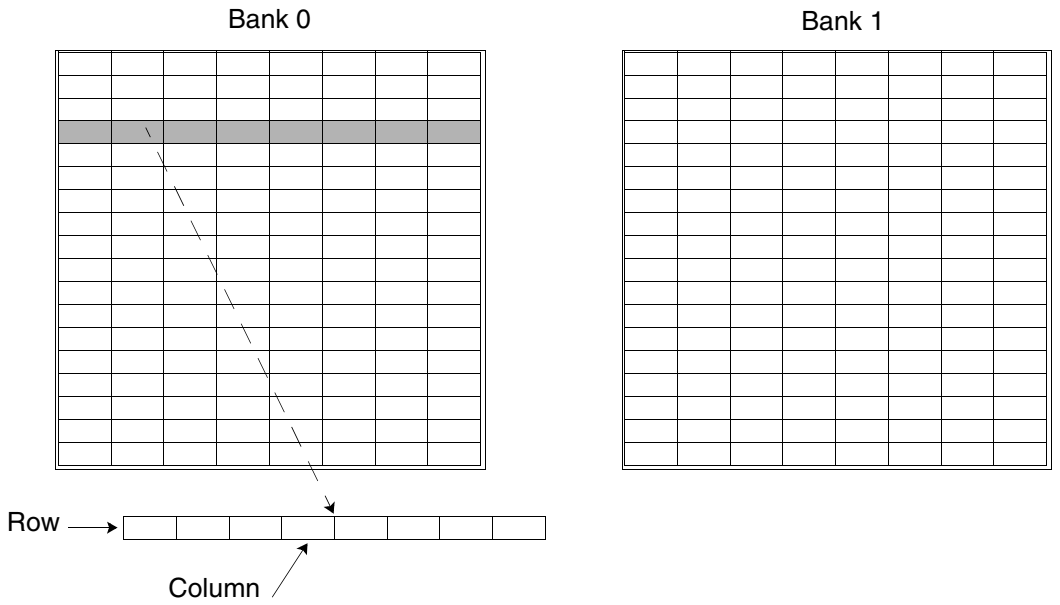
### 8.5.5 Memory Tokens

#### 8.5.5.1 *Memory Banks*

Inside each memory device, there are independent memory banks. Each memory bank can handle only one operation at a time. The operations they handle are reads, writes, and refreshes.

The MIC supports 4, 8, and 16 memory banks on each channel. Because the memory interfaces are independent, the number of effective memory banks supported is double the number each memory device contains when both channels are used. *Figure 8-4* illustrates the differences between rows, columns, and banks. Two banks are shown in this logical representation of the internal organization of a DRAM. One row of Bank 0 is highlighted.

*Figure 8-4. Banks, Rows, and Columns*




If there are 16 memory banks, real addresses are interleaved across the 16 memory banks on a naturally aligned 128-byte basis. Bank 0 holds the first 128-byte block and every sixteenth 128-byte block after the first. Bank 1 holds the second 128-byte block and every sixteenth 128-byte block after the second, and so forth. If there are eight memory banks, real addresses are interleaved across the eight memory banks on a naturally aligned 128-byte basis. Thus, regardless of the actual number of memory banks, memory accesses are managed as if there are 16 logical banks. If there are only eight banks, each physical bank is treated as two logical banks from a RAM-facility perspective.

When a memory device is busy doing an access to a bank, other accesses to the same bank are delayed until that access is complete. The MIC starts accesses to other banks if there are commands for those banks queued inside the MIC and those banks are free. Therefore, to maximize performance, addresses should be spread across all banks. The RAM facility understands memory banks and allocates tokens based on this same definition. The more banks available, the more parallel memory operations can occur at the same time.

#### 8.5.5.2 *Memory-Bank Tokens*

A pair of memory-bank tokens, one even and one odd, are generated when the rate decrementer would decrement below the value of '0'. These generated bank tokens become available to be granted to requesters in the RAG for which they are generated. The choice of banks for which a token becomes available is based on the value of the bank incrementer of a RAG. In the same clock cycle in which a bank token is generated, the bank incrementer is incremented by one.

#### 8.5.5.3 *Memory-Channel Tokens*

When there have been no requests for tokens for a long period of time, the TKM can accumulate up to one token for each managed resource, thus up to 16 tokens for all memory banks. To reduce the rate at which memory tokens are granted if there is a sudden burst of requests for the available bank tokens, the TKM also maintains internal *memory-channel 0 tokens* and *memory-channel 1 tokens*. These memory-channel tokens represent XDR DRAM bus time and are used to pace the rate at which bank tokens are granted.

Memory-channel tokens also allow accesses to be concentrated in a subset of the banks while preventing the XDR DRAM bus from being overused. For example, for certain XDR DRAM $t_{RC}$ values, memory bandwidth might be fully used even though all the accesses went to half the banks. The memory-channel tokens are maintained internal to the TKM and are not tokens that are requested or granted to requesters. The TKM only grants a bank token if both a bank token and a corresponding memory-channel token are available. For even-numbered banks, a memory channel 0 token is needed. For odd-numbered banks, a memory channel 1 token is needed.

See the memory interface controller (MIC) MMIO registers section in the *Cell Broadband Engine Registers* document for configuration information about XDR DRAM $t_{RC}$. To allow for different XDR DRAM $t_{RC}$ values, the Mode for Token-Generation bit, `TKM_CR[MT]`, defines the relative rate at which memory-channel tokens are generated versus memory-bank tokens. Three rates are possible:

- When `TKM_CR[MT]` = '00', memory-channel tokens are made available at the same rate that bank tokens are made available. In this case maximum memory bandwidth can be achieved only if all the banks are equally accessed.
- When `TKM_CR[MT]` = '01', memory-channel tokens are made available at 75% of the rate that bank tokens are made available. Thus maximum memory bandwidth can be achieved even if only three quarters of the banks are accessed.
- When `TKM_CR[MT]` = '10', memory-channel tokens are made available at half the rate that bank tokens are made available. When `TKM_CR[MT]` = '11', operation of the TKM is undefined.

*Table 8-2* summarizes when memory bank and memory-channel tokens become available for the modes defined by `TKM_CR[MT]`.

*Table 8-2. Memory Tokens Available Based on TKM_CR[MT]*

| Bank Incrementer | Memory Tokens That Become Available When Rate Decrementer Decrements Below 0 | | | | |
|---|---|---|---|---|---|
| | Bank Tokens | Memory Channel 0 and Memory Channel 1 Tokens | | | |
| | | MT = '00' | MT = '01' | MT = '10' | MT = '11' |
| 0 | 0 and 1 | memory channel 0, and memory channel 1 | | | Undefined |
| 1 | 2 and 3 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | Undefined |
| 2 | 4 and 5 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | | Undefined |
| 3 | 6 and 7 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | Undefined |
| 4 | 8 and 9 | memory channel 0, and memory channel 1 | | | Undefined |
| 5 | 10 and 11 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | Undefined |
| 6 | 12 and 13 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | | Undefined |
| 7 | 14 and 15 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | memory channel 0, and memory channel 1 | Undefined |

Although the memory-channel tokens prevent bank tokens from being granted in big bursts, it also implicitly requires that tokens be requested at a more steady rate. For example, if the PPE was alone in one RAG with 100% allocation of memory, a pair of memory-channel tokens are made available every interval of time as defined by the `TKM_MBAR` register. An even-bank and odd-bank token must be requested in every such interval, or some of these tokens become unused.

Memory-bank tokens are made available for banks in increasing order. If software fetches data structures from memory in this order, the result might be better use of available tokens.

#### 8.5.5.4 *Memory-Token Grant Algorithm*

The following description of memory-token granting applies to one RAG, the requests in the RAG, and the round-robin pointers in that RAG:

For each memory bank and for each RAG, there is a round-robin pointer to select the next high-priority request in the RAG for the bank, if there is more than one high-priority request active. Likewise, there is a round-robin pointer to select the next low-priority request in the RAG for the bank, if there is more than one low-priority request active. A request is only eligible to be granted a token if the round-robin pointer points to the request.

Tokens for odd banks are granted in a manner similar to even banks, except that a memory-channel 1 token must be available and odd-bank pointers are used to select the bank and requester.

### 8.5.6 I/O Tokens

For each RAG 0:3 and RAG U, there are separate rate decrementers for each IOIF bus direction, in and out. This is necessary because each IOIF and each direction can have different bus widths. When an IOIF0 in rate decrementer, IOIF0 out rate decrementer, IOIF1 in rate decrementer, or an IOIF1 out rate decrementer of a RAG is '0' and is about to be decremented again, a token becomes available for the corresponding bus.

The following description applies to one RAG, the requests in the RAG, and the round-robin pointers in that RAG:

The token-grant facilities for all IOIF buses (IOIF0 In, IOIF0 Out, IOIF1 In, and IOIF1 Out) are alike. The PPE is the only high-priority requester for IOIF buses. For each IOIF bus and for each RAG, there is a round-robin pointer to select the next low-priority request in the RAG for the IOIF bus, if there is more than one low-priority request active. A request is only eligible to be granted a token if the round-robin pointer points to the request.

If an IOIF bus token is available and there is either a high or low-priority eligible request for the IOIF bus, the IOIF bus token is granted in the following manner.

- If there is a high-priority PPE request for the token, the IOIF bus token is granted to the PPE.
- If there is no high-priority eligible request, the IOIF bus token is granted to the requester whose request is pointed to by the low-priority request round-robin pointer for the IOIF bus.
- The Token-Available Latch for the IOIF bus is reset. The request latch corresponding to the requester granted the token is reset.

### 8.5.7 Unused Tokens

If there is no token for a resource already available for a RAG 0:3 whose rate decrementer generates a token, the new token is made available only to that RAG and remains available until the token is granted to the RAG. If there is a token for the resource already available to that RAG, the new token is called an *unused token*.

This unused token is available to other RAGs that have outstanding requests for the resource, if this function is enabled by the Unused Enable bit, `TKM_CR[UE]`. Only RAGs with outstanding requests are eligible for an unused token. If there are multiple RAGs with outstanding requests, the token is given to the RAG with the highest priority, per the `TKM_PR` register, for the specific RAG with the unused token.

Tokens issued from the staging buffer are ready to be used. The tokens can be granted if an outstanding token request exists for that resource, retained in the Token-Available Latch for that resource if there is no existing token, or shared with another RAG as an unused token under the guidelines given in this section.

### 8.5.8 Memory Banks, IOIF Allocation Rates, and Unused Tokens

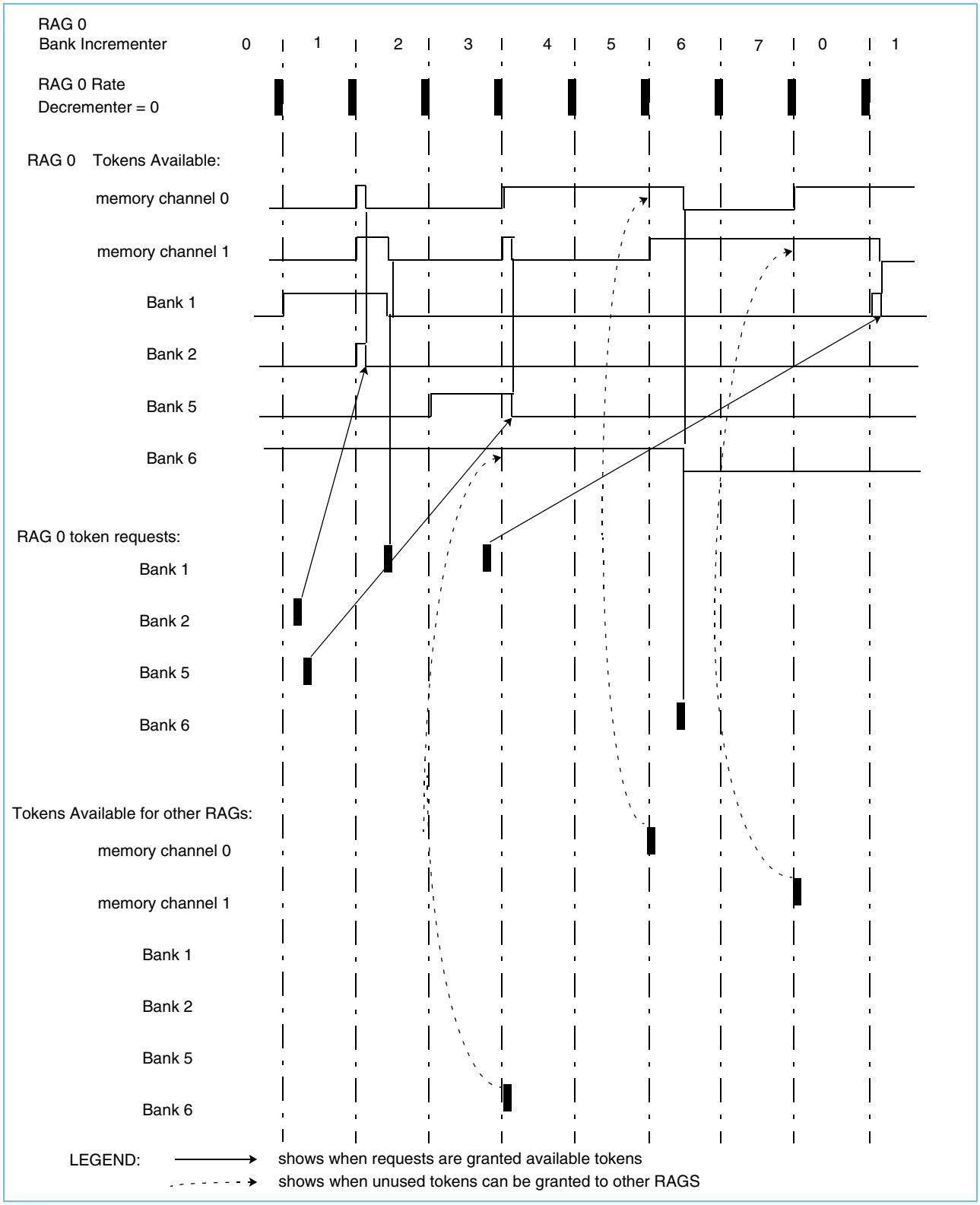
Tokens for memory banks are described in *Section 8.5.5* on page 216, tokens for I/O tokens are described in *Section 8.5.6* on page 220, and unused tokens are described in *Section 8.5.7* on page 220. *Section 8.5.10.1 Memory Allocation* on page 225 and *Section 8.5.10.2 IOIF Allocation* on page 225 describe the calculation of allocation rates for memory and IOIF.

The sum of all RAGs' allocation rates for any given resource should not exceed 100% of the resource. Otherwise, tokens might be lost in the token generator and the resource might be overcommitted. If the resource is overcommitted, the exact amount of the resource that can be used by any given RAG is unpredictable.

### 8.5.9 Token Request and Grant Example

*Figure 8-5* on page 222 shows an example of token requests and grants when `TKM_CR[MT]` = '10'. Bold vertical lines indicate the TKM clock cycles when the RAG 0 rate decrementer has a value of '0'.

*Figure 8-5. Memory-Token Example for TKM_CR[MT] = '10'*



For the example in *Figure 8-5* on page 222, the prescaler for RAG 0 (PO) field of the TKM Memory Bank Allocation Register (TKM_MBAR) is '1', so the RAG 0 rate decrementer contains a '0' value for two NClk/2 cycles. After two cycles, the RAG 0 rate decrementer, instead of being decremented, is reloaded from the Interval value for RAG 0 (IO) field of the TKM_MBAR register, and the RAG 0 bank incrementer is incremented.

With TKM_CR[MT] = '10', tokens for memory channel 0 and 1 become available when the RAG 0 bank incrementer changes from an odd value to an even value. In the cycle when the RAG 0 rate decrementer is reloaded from the TKM_MBAR register, a token becomes available for two banks based on the RAG 0 bank incrementer value. When the bank incrementer value changes from i to i+1, tokens for banks 2i and 2i + 1 become available. For this example, TKM_CR[UE] = '1' to allow sharing of unused tokens.

The following explanation is divided into time periods that correspond to the RAG 0 bank incrementer value:

- RAG 0 Bank Incrementer = 1:
  - When RAG 0 bank incrementer changes from 0 to 1, a token for Bank 1 is generated and the Token-Available Latch is set for this bank. Although not shown in *Figure 8-5*, a token for Bank 0 also becomes available.
  - During this period, there is a RAG 0 request for Bank 2. The tokens for Bank 2 and memory channel 0 are not available. Therefore, a token for Bank 2 cannot be immediately granted.
  - During this period, there is also a RAG 0 request for Bank 5. The tokens for Bank 5 and memory channel 1 are not available. Therefore, a token for Bank 5 cannot be immediately granted.
- RAG 0 Bank Incrementer = 2:
  - When RAG 0 bank incrementer changes from 1 to 2, tokens for Bank 2, memory channel 0 and memory channel 1 are generated, and the Token-Available Latches are set for Bank 2, memory channel 0, and memory channel 1. Although not shown in *Figure 8-5*, a token for Bank 3 also becomes available.
  - Because tokens for both Bank 2 and memory channel 0 are now available, a token for Bank 2 is granted to satisfy the outstanding RAG 0 request from the previous period, and the RAG 0 tokens for Bank 2 and memory channel 0 are no longer available for further grants until the TKM generates new tokens for these resources, based on the RAG 0 bank incrementer and rate decrementer. The token for memory channel 0 is not explicitly granted, but the available token is internally consumed by the TKM as a prerequisite for granting a token for Bank 2.
  - Tokens for both Bank 1 and memory channel 1 are available, so a token for Bank 1 is granted as soon as a RAG 0 requester requests this token in the period of time when the RAG 0 bank incrementer is 2. The Token-Available Latches for Bank 1 and memory channel 1 are reset.
- RAG 0 Bank Incrementer = 3:
  - When the RAG 0 bank incrementer changes to 3, a token for Bank 5 becomes available. Although not shown in *Figure 8-5*, a token for Bank 4 also becomes available.

  - Although the token for Bank 5 is available and there is an outstanding request from the previous period, a token for Bank 5 cannot be granted because there is no available token for memory channel 1.
  - During this period, there is an RAG 0 request for Bank 1. Because the tokens for Bank 1 and memory channel 1 are not available, the request cannot be immediately granted.

- RAG 0 Bank Incrementer = 4:
  - When the RAG 0 bank incrementer changes to 4, tokens for Bank 6, memory channel 0 and memory channel 1 become available. Although not shown in *Figure 8-5*, a token for Bank 7 also becomes available.
  - Because the tokens for both memory channel 1 and Bank 5 are now available, the outstanding RAG 0 request for Bank 5 (from period 1) is granted a token for Bank 5.
  - Because the token for Bank 6 was already available, this newer token (unused token) becomes available for another RAG with an outstanding request for Bank 6.
- RAG 0 Bank Incrementer = 5:
  - When the RAG 0 bank incrementer changes to 5, although not shown in *Figure 8-5*, tokens for banks 8 and 9 become available at this time.
- RAG 0 Bank Incrementer = 6:
  - When the RAG 0 bank incrementer changes to 6, tokens for memory channel 0 and 1 become available. Although not shown in *Figure 8-5*, tokens for Banks 10 and 11 become available at this time.
  - Because the token for memory channel 0 was already available, this newer token becomes available for another RAG with an outstanding request needing a token for memory channel 0.
  - When a RAG 0 request for a token for Bank 6 occurs, the token is immediately granted because tokens for both Bank 6 and memory channel 0 are available.
- RAG 0 Bank Incrementer = 7:
  - When the RAG 0 bank incrementer changes to 7, although not shown in *Figure 8-5*, tokens for Banks 12 and 13 become available at this time.
- RAG 0 Bank Incrementer = 0:
  - When the RAG 0 bank incrementer changes to 0, tokens for memory channel 0 and 1 become available. Although not shown in *Figure 8-5*, tokens for Banks 14 and 15 also become available.
  - Because the token for memory channel1 was already available, this newer token becomes available for another RAG with an outstanding request needing a token for memory channel 1.
- RAG 0 Bank Incrementer = 1:
  - When the RAG 0 bank incrementer changes to 1, a token for Bank 1 becomes available. Although not shown in *Figure 8-5*, a token for Bank 0 also becomes available.
  - Because there are available tokens for both Bank 1 and memory channel 1, the outstanding request for Bank 1 (from period 3) is granted a token for Bank 1.

## 8.5.10 Allocation Percentages

### 8.5.10.1 *Memory Allocation*

The total memory-allocation percentage equals the sum of percentages of all RAGs. The percentage of memory-resource time allocated to a single RAG can be calculated using the following equations, in which the I and P variables are the Interval and Prescaler values from the appropriate fields of the TKM_MBAR register, and MT is the Mode-for-Token-generation field in the TKM_CR register.

These equations assume that TKM_CR[MT] is configured consistently with XDR DRAM $t_{RC}$. The maximum value of the $t_{RC}$ for reads and writes must be used in this calculation. CFM and CFMN are differential input clock signals to the CBEA processor. See the *Rambus XDR Architecture (DL-0161)* document for more information about the clock from master (CFM) signals. The number of memory channels in the equations is the number of memory channels in the Cell/B.E. processor that are used to attach XDR DRAM.

NClk is the core clock. The EIB, TKM, and portions of MIC and Cell Broadband Engine interface (BEI) unit run at half this frequency (NClk/2). See *Section 13 Time Base and Decrementers* on page 381.

The equations are:

Allocated bandwidth capability =

(128 bytes/memory channel token) ×
(((4-MT) / 2) memory channel tokens/decrementer expiration) ×
(1 decrementer expiration / (NClk/2 period × (I + 1 + 8 ×
((P>0) OR TKM_MBAR[33]) × $2^P$)))

= 64 bytes × (4-MT) × NClk/2 frequency / (I + 1 + 8 × ((P>0) OR TKM_MBAR[33]) × $2^P$)

### 8.5.10.2 *IOIF Allocation*

The total IOIF-allocation percentage equals the sum of percentages of all RAGs. The percentage of IOIF-bus resource time allocated to a single RAG can be calculated using the following equations, in which the I and P variables are the Interval and Prescaler values from the appropriate fields of the TKM_IOIF0_AR or TKM_IOIF1_AR registers. The IOIF0 and IOIF1 allocation percentages are different for the same values of I and P.

The equations are:

Allocated% = 100 × allocated bandwidth capability / available bandwidth

IOIF0 allocated bandwidth capability

= (128 bytes / token) × (1 token per decrementer expiration) × (1 decrementer expiration / (NClk/2 period × (9 + I) × $2^P$))

= 128 bytes × NClk/2 frequency / ((9 + I) × $2^P$)

$$\text{IOIF1 allocated bandwidth capability} = 128 \text{ bytes} \times \text{NClk/2 frequency} / ((9 + I) \times 2^{p+2})$$

### 8.5.11 Efficient Determination of TKM Priority Register Values

A method is described here for calculating the TKM Priority Register **(**TKM_PR) fields using the *PowerPC Architecture* instructions.

There are enables in the TKM_PR register for each RAG to obtain tokens from every other RAG. For example, the TKM_PR register has four enables for obtaining unused tokens from RAG U—one enable for each of RAG 0:3. There are three enables for obtaining unused tokens from RAG 0—one each for RAG 1:3.

The values for the TKM_PR fields can be calculated without a large number of instructions and without calculating each individual bit as a separate calculation. An example is used to explain the algorithm for determining priorities for RAGs 0:3 in obtaining unused tokens from RAG U. For this example, the priority for getting RAG U tokens is given as the following order, from highest-priority to lowest-priority:

```
RAG1 > RAG0 > RAG3 > RAG2
```

RAG1 has highest priority and RAG2 has lowest priority.

*Table 8-3* shows the universal set of binary OR and AND values for all RAG priorities. These values are used in the algorithm steps that follow. The values are also used every time priority is determined, not only for this example but for all RAG priorities. However, the values are used by the algorithm steps in a different order; the order is based on the relative priority of the various RAGs.

*Table 8-3. OR and AND Values for All RAG Priorities*

| Binary Value | Set Priority OR and AND value |
|---|---|
| 111000 | RAG 0 set priority OR_value |
| 111111 | RAG 0 set priority AND_value |
| 000110 | RAG 1 set priority OR_value |
| 011111 | RAG 1 set priority AND_value |
| 000001 | RAG 2 set priority OR_value |
| 101011 | RAG 2 set priority AND_value |
| 000000 | RAG 3 set priority OR_value |
| 110100 | RAG 3 set priority AND_value |

This 3-step algorithm is used for determining priority for RAGs 0:3 in obtaining RAG U unused tokens (with all 6 digits necessary), and for determining the priority for one RAG obtaining unused tokens from another RAG (with only the last three digits used). When the steps are used for determining priority for unused tokens from RAG 0:3, the number of RAGs used in step 2 is one fewer, and only a 3-bit value is produced in the least-significant bits of general-purpose register (GPR) *r* by steps 1 and 2. Similar to step 3, this 3-bit value should be rotated and

inserted into the 3 bit positions in TKM_PR for the appropriate RAG's Unused tokens. Thus, for RAG 0 unused tokens, the 3-bit value should be rotated to bits 43:45 of the data value to be stored to TKM_PR.

The steps for determining priority are:

1. Clear a GPR *r* in the PPE to hold the results for the RAG U priority fields, where the least-significant 6 bits of the GPR will eventually hold the 6-bit value corresponding to TKM_PR[34:39].
2. For each of three RAGs, starting with the next-to-last RAG in the priority order (that is, the second from lowest priority) and working backward to the RAG with the highest priority order:
   a. Logically OR GPR *r* with the OR_value for the RAG.
   b. Logically AND GPR *r* with the AND_value for the RAG.
3. GPR *r* now has the value of the priority fields for RAG U, so rotate and insert this value into a GPR that will hold the data value to be stored to TKM_PR register.

Applying these three steps to the example:

1. Clear the GPR r bits. The least-significant 6 bits of this GPR, corresponding to the priority fields, have the following values:

   0 0 0 0 0 0

2. Apply this step for 3 RAGs:
   a. GPR r least-significant 6 bits after RAG 3 values are ORed and ANDed:

      0 0 0 0 0 0

   b. GPR r least-significant 6 bits after RAG 0 values are ORed and ANDed:

      1 1 1 0 0 0

   c. GPR r least-significant 6 bits after RAG 1 values are ORed and ANDed:

      0 1 1 1 1 0

3. GPR r now has the value to be rotated and inserted into a GPR that holds the data value to be stored to TKM_PR register.

The 6-bit value being determined for RAG U unused-token priorities encompasses three distinct fields of the TKM_PR register—RU0Vr, RU1Vr, and RU2V3. When the algorithm is used to find the 3-bit values for determining priorities from unused tokens from RAG 0:3, those three bits encompass two distinct fields of the TKM_PR register.

The following *PowerPC Architecture* instructions are used to determine the values to be loaded into the RU0Vr, RU1Vr, and RU2V3 fields of the TKM_PR register for this example. GPR 1 is used for calculating the value to be rotated/inserted into GPR 2, which can be used to collect values calculated for other fields before GPR 2 is finally written to the TKM_PR register.

```
xor 1,1,1                   * clear GPR 1
ori 1,1,b'000000'           * logically OR GPR 1 with RAG 3 OR_value
andi. 1,1,b'110100'         * logically AND GPR 1 with RAG 3 AND_value
ori 1,1,b'111000'           * logically OR GPR 1 with RAG 0 OR_value
andi. 1,1,b'111111'         * logically AND GPR 1 with RAG 0 AND_value
```

```
ori 1,1,b'000110'          * logically OR GPR 1 with RAG 1 OR_value
andi. 1,1,b'011111'        * logically AND GPR 1 with RAG 1 AND_value
rldimi 2,1,24,34           * rotate GPR 1 value and insert into GPR 2
                           * to align it with fields in TKM_PR register
```

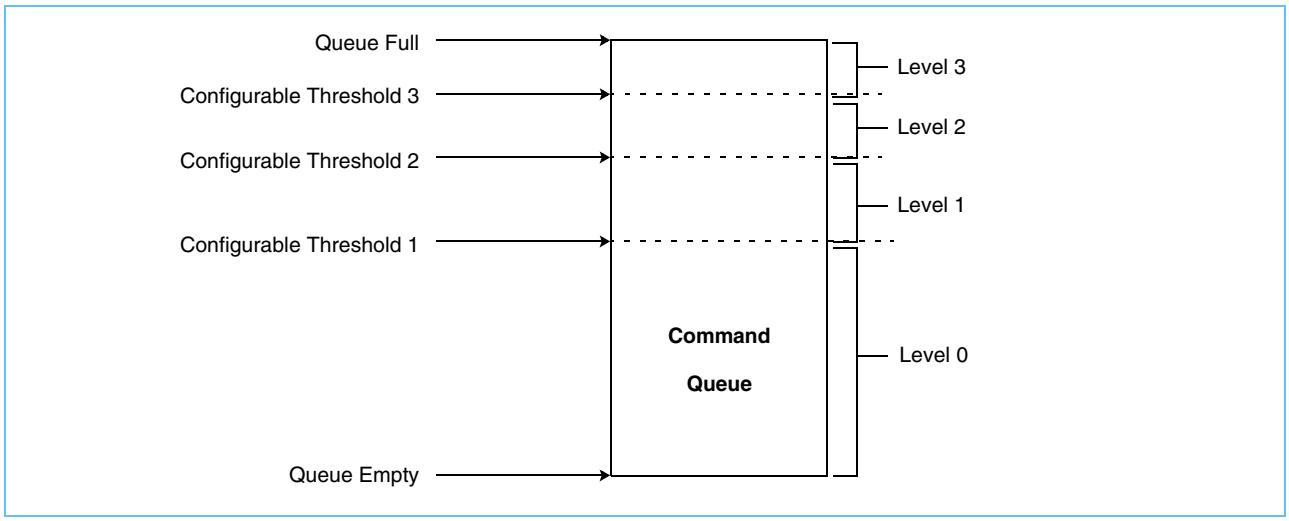
### 8.5.12 Feedback from Resources to Token Manager

The MIC and IOC (when configured for the IOIF protocol) provide feedback to the TKM about the state of their command queues, as shown in *Figure 8-1* on page 204. The MIC and IOC have registers to hold programmable thresholds. The feedback describes command-queue levels relative to these programmable thresholds, except that when an IOIF interface is configured for the BIF protocol, no such feedback is provided.

The command-queue levels, shown in *Figure 8-6* on page 228, are defined as follows, relative to the number of entries in the command queue:

- Level 0: number entries < Queue Threshold 1 field.
- Level 1: Queue Threshold 1 field ≤ number entries < Queue Threshold 2 field.
- Level 2: Queue Threshold 2 field ≤ number entries < Queue Threshold 3 field.
- Level 3: Queue Threshold 3 field ≤ number entries.

*Figure 8-6. Threshold and Queue Levels*




If TKM_CR[FE] enables the use of this queue-level feedback, the TKM grants tokens based on the queue-level feedback shown in *Table 8-4*. Ideally, feedback primarily limits the number of unallocated tokens granted and, secondarily, limits the sharing of unused tokens. If the feedback reaches Level 3, a TKM_ISR bit is set for the resource. If a TKM_ISR bit is '1' and the corresponding interrupt-enable bit in TKM_CR enables the interrupt, an external interrupt condition is reported to the internal interrupt controller (IIC). See *Section 9 PPE Interrupts* on page 239 for more information about IIC external interrupt exceptions and how an external interrupt to the PPE can occur based on TKM_ISR.

*Table 8-4. Token Allocation Actions Per Resource Queue Level*

| Token Allocation Actions | Queue Level for the Resource | | | |
|---|---|---|---|---|
| | 0 | 1 | 2 | 3 |
| Unallocated tokens are given out when available. | Yes | No | No | No |
| Unused tokens from one RAG are given to another RAG when the Allocation Decrementer is 0. | Yes | Yes | No | No |
| Tokens are generated for a RAG. | Yes | Yes | Yes | No |

#### 8.5.12.1 *MIC Feedback to Token Manager*

For each memory channel, the MIC has a command queue for reads and a command queue for writes. The MIC maintains counts for valid entries in these queues. These counts are compared to programmable threshold values in the MIC token manager Threshold Levels Register (`MIC_TM_Threshold_n` [`n = '0','1'`]). The results of these comparisons are reported by the MIC to the TKM and used by the TKM as indicated in *Table 8-4*.

The highest-level information for any of these counts is reported by the MIC to the TKM. For example, if the command queue for reads has more entries than indicated by the value of the read-threshold-level 2 field in `MIC_TM_Threshold_n`, but fewer than the value of the read-threshold-level 3 field, and the command queue for writes has fewer entries than the value of the write-threshold-level 1 field, then the level reported to the TKM is level 2.

#### 8.5.12.2 *IOC Feedback to Token Manager*

The IOC has logically separate command queues for each of the two IOIF buses, and for each bus there are separate queues for the commands received from the EIB and the IOIF. The commands remain in these queues until data is transferred on the IOIF bus. Thus, these queue entries represent the amount of outstanding data traffic for which a token has already been granted.

There is separate feedback to the TKM for each IOIF bus. For each IOIF bus, combined counts of the commands from the EIB and IOIF are used for the feedback, and the commands in the IOIF queue are only counted when a token has been obtained for their use of the IOIF bus. For commands in these queues, the IOC determines the number of commands that transfer data on the In bus and commands that transfer data on the Out bus. The IOIF0 counts are compared to programmable threshold values in the IOIF0 Queue Threshold Register (`IOC_IOIF0_QueThshld`) and the IOIF1 counts are compared to programmable threshold values in the IOIF1 Queue Threshold Register (`IOC_IOIF1_QueThshld`). The results of these comparisons are reported by the IOC to the TKM and used by the TKM as indicated in *Table 8-4* on page 229.

## 8.6 Configuration of PPE, SPEs, MIC, and IOC

### 8.6.1 Configuration Register Summary

*Table 8-5* on page 231 shows the configuration registers or bits in the PPE, SPEs, MIC, and IOC units for controlling functions related to resource allocation. This table does not include all configuration registers that affect resource bandwidth. It only lists those related to resource allocation.

The Configuration Purpose column in *Table 8-5* shows the following functions that are controlled by these configuration registers:

- *Token-Request Enable or Override*—Configuration values shown in this row are reset to a state in which they do not require tokens. A configuration bit allows software to enable their dependency on obtaining tokens. Memory-mapped configuration registers holding these bits are in 4 KB pages with other registers that only need to be accessed by the hypervisor.
- *RAID Assignment*—The PPE and SPEs each have a configuration register that identifies the resource allocation group ID (RAID) of the RAG associated with a resource access. The IOC has a configuration register that specifies the virtual-channel ID for outbound commands—that is, commands initiated by the PPE or an SPE. This virtual-channel ID can be associated with a RAID by the IOIF device, if that device implements Resource Allocation Management.
- *Address Map for Resources*—These registers provide a simplified view of the memory-map sufficient to recognize the addresses that access memory and IOIF1. The address ranges are defined either by MMIO registers accessible to software, or by nonmapped registers accessible only by the configuration ring bits at power-on reset (POR), as described in the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

  If MMIO configuration registers are implemented, these registers are in 4 KB pages with other registers that only need to be accessed by the hypervisor (see *Section 11* on page 331). Addresses associated with IOIF0 or the BIF protocol are assumed to be addresses that do not correspond to one of the three local address spaces (memory, IOIF1, or MMIO registers). In a multi-CBEA-processor environment, this allows all addresses associated with resources in or attached to the other CBEA processor to be identified as using the BIF protocol.

  See *Section 8.6.2* on page 231 for details about the SPE mapping of addresses to resources. Like the SPE, the PPE has a similar coupling of IOIF1 address space to the CBEA processor MMIO address space.
- *Multiple Tokens for Short Writes*—For memory-write operations shorter than 128 bytes, these configuration settings define whether multiple tokens are required for the PPE and IOC. The SPE always obtains two tokens for memory writes of 8 bytes or less.
- *Feedback Thresholds*—Configuration registers for the MIC and IOC that define the thresholds for reporting feedback to the TKM.

*Table 8-5. Configuration Bits in the PPE, SPEs, IOC, and MIC*

| Purpose | Configuration Bits | | | |
|---|---|---|---|---|
| | PPE | SPEs | IOC | MIC |
| Token Request Enable or Override | BIU Configuration Ring:[1]<br>• Resource Allocation Override<br>MMIO Register:<br>• BIU Mode Setup 1 (BIU_ModeSetup1[1]) | SPE Configuration Ring:[1]<br>• Resource Allocation Override<br>MMIO Register:<br>• Resource Allocation Enable (RA_Enable[63]) | MMIO Register:<br>• IOCmd Configuration (IOC_IOCmd_Cfg[17]) | Not applicable (the MIC is not a requester). |
| RAID Assignment | MMIO Register:<br>• BIU Mode Setup 1 (BIU_ModeSetup1[2:3]) | MMIO Register:<br>• RAID (RA_Group_ID[62:63]) | MMIO Register:<br>• Virtual Channel For Outbound Accesses: IOCmd Configuration (IOC_IOCmd_Cfg[32:34] and IOC_IOCmd_Cfg[49:51]) | |
| Address Map for Resources | BIU Configuration Ring:[1]<br>• MMIO Base Address<br>MMIO Registers:<br>• for memory: BIU Mode Setup Register 1 (BIU_ModeSetup1[46:60] and BIU_ModeSetup1[14:28])<br>• for IOIF1: BIU Mode Setup Register 2 (BIU_ModeSetup2[22:31]) | SPE Configuration Ring:[1,2]<br>• MC_BASE<br>• MC_COMP_EN<br>• IOIF1_COMP_EN<br>• BE_MMIO_Base | Configuration Ring:[1]<br>• BEI Base Address<br>• IOIF1 Base Address<br>• IOIF1 Base Address Mask<br>MMIO Registers:<br>• Memory Base Address (IOC_MemBaseAddr)<br>• IOC Base Address 1 (IOC_BaseAddr1)<br>• IOC Base Address Mask 1 (IOC_BaseAddrMask1) | |
| Multiple Tokens for Short Writes | BIU Configuration Ring:[1]<br>• Two-token decode for noncacheable unit (NCU) store | Not applicable | MMIO Register:<br>• IOCmd Configuration (IOC_IOCmd_Cfg[18] and IOC_IOCmd_Cfg[19]) | |
| Feedback Thresholds | Not applicable | | MMIO Register:<br>• IOIF0 Queue Threshold (IOC_IOIF0_QueThshld)<br>• IOIF1 Queue Threshold (IOC_IOIF1_QueThshld) | MMIO Register:<br>• Token Manager Threshold Levels n [n = 0,1] (MIC_TM_Threshold_n [n = 0,1]) |

1. For details about the configuration ring bits and the power-on reset (POR) sequence, see the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.
2. See *Section 8.6.2* on page 231.

### 8.6.2 SPE Address-Range Checking

Each SPE has a defined real-address range for resources managed by the RAM facility. Before an SPE requests a token from the TKM, it checks the address range of the accessed resource. If the address falls into the SPE's defined range for a resource managed by the RAM facility, the SPE requests a token.

The range is defined in four registers that are written by the configuration ring bits. These registers, listed in *Table 8-6*, are not memory-mapped, but are only accessible by means of the configuration ring. The registers and the configuration ring bits are described in the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

*Table 8-6. SPE Address-Range Registers*

| Configuration-Ring Register | Description |
|---|---|
| BE_MMIO_Base | 19-bit register that defines the BE_MMIO_Base starting address. |
| MC_BASE | 15-bit register that defines the starting address of the MIC external memory. |
| MC_COMP_EN | 15-bit register used to specify the memory controller size. If ((RA(0:14) AND MC_COMP_EN) equal (MC_BASE AND MC_COMP_EN), then the RA is in the memory controller (MC) range. See *Table 8-7*. |
| IOIF1_COMP_EN | 10-bit register used to specify the IOIF1 size. See *Table 8-8* on page 233. |

The `MC_COMP_EN` register is assigned for system memory that belongs to the on-chip MIC. *Table 8-7* specifies the size of the `MC_COMP_EN`.

*Table 8-7. MC_COMP_EN Setting*

| MC_COMP_EN | Memory Size |
|---|---|
| x‘4000’ | 2 TB |
| x‘6000’ | 1 TB |
| x‘7000’ | 512 GB |
| x‘7800’ | 256 GB |
| x‘7C00’ | 128 GB |
| x‘7E00’ | 64 GB |
| x‘7f00’ | 32 GB |
| x‘7f80’ | 16 GB |
| x‘7FC0’ | 8 GB |
| x‘7FE0’ | 4 GB |
| x‘7FF0’ | 2 GB |
| x‘7FF8’ | 1 GB |
| x‘7FFC’ | 512 MB |
| x‘7FFE’ | 256 MB |
| x‘7FFF’ | 128 MB |

The IOIF1 range must start from the aligned address set by the `IOIF1_COMP_EN` register, and it is adjacent to `BE_MMIO_Base`. To guarantee that `BE_MMIO_Base` and the IOIF1 range do not touch each other, use the settings shown in *Table 8-8* on page 233. The IOIF0 range and IOIF1 range cannot be set independently. The IOIF0 range is the remaining address space (from the 42-bit real address [RA]) not defined by BE_MMIO_Base, IOIF1 range, and `MC_COMP_EN`.

*Table 8-8. IOIF1_COMP_EN and BE_MMIO_Base Settings*

| IOIF1_COMP_EN | Size of IOIF1 | BE_MMIO_Base | IOIF1 Start |
|---|---|---|---|
| 1111111111 | 2 GB | be_mmio_base(10)= 0 | be_mmio_base(0:9)\|1 \| 31'0 |
| 1111111110 | 4 GB | be_mmio_base(9) = 0 | be_mmio_base(0:8)\|1 \| 32'0 |
| 1111111100 | 8 GB | be_mmio_base(8) = 0 | be_mmio_base(0:7)\|1 \| 33'0 |
| 1111111000 | 16 GB | be_mmio_base(7) = 0 | be_mmio_base(0:6)\|1 \| 34'0 |
| 1111110000 | 32 GB | be_mmio_base(6) = 0 | be_mmio_base(0:5)\|1 \| 35'0 |
| 1111100000 | 64 GB | be_mmio_base(5) = 0 | be_mmio_base(0:4)\|1 \| 36'0 |
| 1111000000 | 128 GB | be_mmio_base(4) = 0 | be_mmio_base(0:3)\|1 \| 37'0 |
| 1110000000 | 256 GB | be_mmio_base(3) = 0 | be_mmio_base(0:2)\|1 \| 38'0 |
| 1100000000 | 512 GB | be_mmio_base(2) = 0 | be_mmio_base(0:1)\|1 \| 39'0 |
| 1000000000 | 1 TB | be_mmio_base(1) = 0 | be_mmio_base(0)\|1 \| 40'0 |

## 8.7 Changing Resource-Management Registers with MMIO Stores

In some situations, it is not necessary to ensure the exact time at which a change to a resource-allocation facility has occurred and the side effects of changing such a facility occur. If there is a need to synchronize such changes or limit the amount of variability of the side effects, PPE software can use the procedures described in this section.

### 8.7.1 Changes to the RAID

When a PPE context switch is to be made from one application to another application, or from one operating system to another operating system, it might be necessary to change the RAID used by the PPE or an SPE. (When an SPE context switch is to be made, use the procedures described in *Section 12 SPE Context Switching* on page 357.)

When a PPE context switch is made, the following sequence can be used to change the RAID:

1. Read the PPE's `BIU_ModeSetup1` register, and save this as part of the state.
2. After saving state, execute a **sync** instruction to ensure previous storage accesses are performed before the RAID change in the following step. (This **sync** is only useful if more precise use of the RAID is required.)
3. Store the new RAID in the PPE's `BIU_ModeSetup1` register.
4. Execute a **sync** instruction to ensure the store to change the RAID is performed before subsequent storage accesses. (This **sync** is only useful if more precise use of the RAID is required.)
5. Do memory accesses to restore the state of the next task.

This sequence assumes that the resource used up to the step in which the RAID is modified is attributed to the previous RAID, and the resource used after this step is attributed to the new RAID. This sequence does not guarantee that resource use is attributed to the correct RAG with 100% accuracy, because the PPE uses an implied token for each resource. Complete accuracy might be unimportant, because resource usage by the privileged software doing the context

switch and handling interrupts cannot be accurately attributed either. If complete accuracy is required, software can use the more lengthy sequence defined in *Section 8.7.6.1 Changes Involving Zero Allocation* on page 236.

## 8.7.2 Changing a Requester's Token-Request Enable

*Table 8-5* on page 231 shows the configuration registers that control whether a requester requires a token. To synchronize the side effects of a store to such a configuration register, software should subsequently load from the register and synchronize that load using a synchronization instruction such as a **sync**.

Before enabling token requests of a requester, software should configure the TKM_CR, TKM_MBAR, TKM_IOIF0_AR, and TKM_IOIF1_AR registers so that there are tokens granted to the RAG of the requester. Failure to do this might cause the requester to request a token that will not be granted, potentially resulting in a timeout and checkstop.

### 8.7.2.1 *Changing PPE Token-Request Enable*

Software cannot turn off the PPE's token request after it is enabled. Software must only change BIU_ModeSetup1[ResAllocEn] when BIU_ModeSetup1[ResAllocEn] = '0'. If BIU_ModeSetup1[ResAllocEn] = '1', software must not change BIU_ModeSetup1[ResAllocEn]; otherwise, an outstanding storage access might timeout and a checkstop might occur. To ensure that the PPE uses the new value for accesses, software should load from BIU_ModeSetup1 after storing to it and software should synchronize that load using a synchronization instruction such as a **sync**.

### 8.7.2.2 *Changing SPE Token-Request Enable*

Software can change the RA_Enable register to enable SPE token requests at any time. Token-issue can be enabled by setting the Resource Allocation Enable bit, RA_Enable[M], to '1' and the Resource Allocation Override bit, RA_Enable[C], to '1'. The RA_Enable[C] field is read-only; its state can only be set by means of the Configuration-Ring load at POR, as described in the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

Resource-allocation override is set by the configuration ring. RA_Enable[M] is cleared to '0' after POR. Software must set this bit to '1' to enable token-request. After RA_Enable[M] is set to '1', software should not reset this register unless all direct memory access controller (DMAC) and atomic commands are complete. If there are any outstanding DMAC or atomic commands and RA_Enable[M] is '0', the behavior of the SPE is unpredictable.

Software must only change the RA_Enable register to disable SPE token requests when SPE accesses that require tokens are not occurring. To disable SPE token requests, PPE software can use the following sequence:

1. Perform an MMIO write of the MFC_CNTL register with the suspend DMA-queue (Sc) bit set to '1'.
2. Poll the MFC_CNTL register using an MMIO read until the MFC suspension requested in the previous step is complete.
3. Flush the contents of the MFC atomic unit to prevent a push from occurring while changing RA_Enable. The contents of the atomic unit can be flushed by writing a '1' to

MFC_Atomic_Flush[F] and then polling MFC_Atomic_Flush using an MMIO read until ‘0’ is returned.

4. Perform an MMIO write of the RA_Enable register to disable SPE token requests.
5. Perform an MMIO read of the RA_Enable register to ensure that the SPE uses the new value of RA_Enable for accesses.
6. To resume normal MFC operations, perform a MMIO write of the MFC_CNTL register with the restart MFC command (R) bit set to '1'.

#### 8.7.2.3 *Changing I/O Token-Request Enable*

Software can change IOC_IOCmd_Cfg[TK] to ‘1’ at any time. However, this change only affects the token requests for new I/O accesses. I/O accesses that are queued in the IOC waiting for tokens are not affected by the change. To ensure that the IOC uses the new value for new reads and writes from the IOIF, software should load from IOC_IOCmd_Cfg after storing to it and synchronize that load using a synchronization instruction such as a **sync**.

Software must only change IOC_IOCmd_Cfg[TK] to ‘0’ when there are no concurrent or outstanding accesses from IOIF0 or IOIF1. Changing IOC_IOCmd_Cfg[TK] to ‘0’ when there is a read or write from an IOIF might cause the access to timeout because the read or write does not get the token that the IOC previously detected was needed.

### 8.7.3 Changing a Requester’s Address Map

*Table 8-5* on page 231 shows the configuration registers that controls the address map for resources. The CBEA processor MMIO address map and the SPEs can only be configured by means of the configuration ring, so its configuration must be established at power-on reset.

To ensure that any side-effect of changing an IOC MMIO configuration register that controls the address map for resources has occurred, software should load from the last register modified and synchronize that load using a synchronization instruction such as a **sync** (the **sync** is needed after the read because these are pipelined registers). This change only affects the memory map for new I/O accesses. I/O accesses that are queued in the IOC waiting for tokens are not affected by the changing address map. Because the address map for IOIF1 is controlled by two registers, it is not possible to change both at the same instant. As a result, software should only change the IOC_BaseAddr1 register or the IOC_BaseAddrMask1 register when either IOC_IOCmd_Cfg[TK] is ‘0’ or there are no concurrent or outstanding accesses to IOIF1 from IOIF0 while the IOC_BaseAddr1 register or IOC_BaseAddrMask1 register is being modified.

Changing the PPE address map for resources might require changing both the BIU_ModeSetup1 register and the BIU_ModeSetup2 register. After changing one of these registers, software must avoid having a temporary memory map that has a single address mapped to two resources. To ensure that any side-effect of changing either the BIU_ModeSetup1 register or the BIU_ModeSetup2 register has occurred, software should load from the modified register after storing to it and synchronize that load using a synchronization instruction such as a **sync**. If modifying both of these registers, software only needs to load from the register modified last and then synchronize that load.

### 8.7.4 Changing a Requester's Use of Multiple Tokens per Access

For some period of time after performing a store to change IOC_IOCmd_Cfg[RMW:SXT], the IOC token requests can be based on either the new or old value of IOC_IOCmd_Cfg[RMW:SXT]. If software needs to ensure the new value is used, software can load from the modified register and synchronize that load using a synchronization instruction such as a **sync**. This change only affects the token requests for new I/O accesses. I/O accesses that are queued in the IOC waiting for tokens are not affected by the change.

### 8.7.5 Changing Feedback to the TKM

After a store to MIC_TM_Threshold_n [n = '0','1'], IOC_IOIF0_QueThshld, or IOC_IOIF1_QueThshld, to ensure that the new values for these registers are used for feedback to the TKM, software can load from a modified register and synchronize that load using a synchronization instruction such as a **sync**.

### 8.7.6 Changing TKM Registers

Before changing TKM_CR, TKM_MBAR, TKM_IOIF0_AR, TKM_IOIF1_AR, or TKM_PR, software must ensure that there are no outstanding requests in the token manager. Furthermore, the allocation registers (TKM_IOIF0_AR and TKM_IOIF1_AR) must be written with zeros before a new allocation rate is set. To ensure that the new register values are used after a store to the TKM_CR, TKM_PR, or TKM_PMCR registers, software can load from a modified register and synchronize that load using a synchronization instruction such as a **sync**.

After a store to the TKM_MBAR, TKM_IOIF0_AR, or TKM_IOIF1_AR registers, the TKM uses the new allocation rate after the corresponding RAG rate decrementer expires. To ensure that the new values for these registers are used, software can load from a modified register, synchronize that load using a synchronization instruction such as a **sync**.

#### 8.7.6.1 *Changes Involving Zero Allocation*

Before changing the TKM_CR, TKM_MBAR, TKM_IOIF0_AR, or TKM_IOIF1_AR registers so that a RAG is no longer allocated some resource, software must ensure that there are no outstanding token requests for that resource from that RAG and that there will be no such future requests. Requesters in such a RAG for such a resource are not granted tokens and might time-out, potentially resulting in a checkstop.

Software can ensure that this is the case for SPEs by using the sequences in the *Section 12 SPE Context Switching* on page 357. For some value r, if RAG r is no longer going to be allocated some resource, then during the restore sequence software loads a value other than r into the RAID field of the SPE RA_Group_ID register.

Ensuring that the PPE has no outstanding request for RAG r is more complicated, because the PPE uses an implied token for each resource. Ensuring that previous storage accesses are performed does not guarantee that replacement tokens have been obtained. The following sequence should be executed from the PPE before changing the TKM_CR, TKM_MBAR, TKM_IOIF0_AR, or TKM_IOIF1_AR registers to eliminate a resource allocation for RAG r:

1. Store the new RAID, where this RAID is not r, in the PPE's BIU_ModeSetup1 register.
2. Read the BIU_ModeSetup1 register.

3. Execute a **sync** instruction to ensure the read has been performed before subsequent storage accesses.
4. If RAG `r` is no longer allocated a memory resource, then, for each memory bank, execute a **dcbf** instruction and a load instruction where both instructions use the same operand address corresponding to the memory bank. Software controls whether a RAG is allocated tokens for a memory resource. For example, setting any of the `TKM_CR[Ei]` bits[44:47] to '0' is one way to stop granting tokens to the corresponding RAG.
5. If RAG `r` is no longer allocated an IOIF Out data bus resource, then store to a location attached to the IOIF.
6. If RAG `r` is no longer allocated an IOIF In data bus resource, then load from a location attached to the IOIF.
7. Execute a **sync** instruction to ensure that accesses in the preceding steps have been performed.

## 8.8 Latency Between Token Requests and Token Grants

Enabling Resource Allocation Management adds latency because of the time required to request and grant tokens. If a token is already available and there are no other requests in the same RAG for the same resource, there is a 6 `NClk`/2 cycle latency in the TKM from request to grant. If the token is not already available, there might be additional latency due to the wait time in the TKM request queue. There is also latency in each requester, delay in transferring the request from the requester to the TKM, and delay in transferring the grant from the TKM to the requester.

## 8.9 Hypervisor Interfaces

The CBEA processor hypervisor facility is described in *Section 11 Logical Partitions and a Hypervisor* on page 331. Hypervisor software might find it useful to define a few hypervisor function calls that manage the relationship between the RAM facility's RAG allocations and the hypervisor's logical-partitioning facility. This might be useful, for example, when quality of service (QoS) must be guaranteed for specific resources used by a specific logical partition.

For example, hypervisor function calls might perform the following functions:

- *Acquire a RAG*—This function acquires a RAG for a logical partition. Function arguments might include the logical-partition ID, minimum acceptable memory-access allocation, and minimum acceptable I/O-access allocation. Return values might include the allocated RAID (or more detailed RAG descriptor) and various success or failure indicators. All resource requesters belonging to that logical partition become a member of the new RAG. When associating a RAG with a logical partition that is currently able to run, the partition should be immediately removed from an unallocated RAG and placed into the new RAG. The acquisition of a new RAG in this manner requires some initial allocation for memory and I/O resources, or the partition does not make forward progress.
- *Allocate a RAG*—This function assigns a portion of resource's time that a RAG can use the resource. Function arguments might include the RAID or descriptor of the RAG that was acquired by calling the *Acquire a RAG* function, the resource to be allocated, and the portion of the resource's time to allocate. Return values might include the remaining portion of the resource's time, after successful allocation, and various success or failure indicators.

- *Release a RAG*—This function releases a RAG associated with a logical partition. Function arguments include the RAID or descriptor of the RAG to release. Return values might include various success or failure indicators. Any requester owned by the logical partition associated with this RAG has its resources allocated by the Unallocated RAG.

See *Section 11.5.1 Combining Logical Partitions with Resource Allocation* on page 352 for a short overview of how these two facilities can be used together.

# 9. PPE Interrupts

This section describes all interrupts that are handled by the PowerPC Processor Element (PPE), or by an external interrupt controller. Additional interrupts caused by Synergistic Processor Element (SPE) events, such as signal-notification events, are handled locally by each SPE as described in *Section 18.3 SPU Interrupt Facility* on page 476.

## 9.1 Introduction

Exceptions and interrupts are closely related. The *PowerPC Architecture* and the *PowerPC Compiler Writer's Guide* define an *exception* as an error, external event, or other condition requiring attention, such as an instruction that cannot be executed, the assertion of an external signal, or a timer event. An exception can set status information in a register, indicating the nature of the exception, and it can cause a corresponding *interrupt,* if a corresponding interrupt is enabled and (if applicable) unmasked. It is possible for a single instruction to cause multiple exceptions and for multiple exceptions to cause a single interrupt. By masking interrupts, software can also support polling as an alternative to interrupts.

An interrupt is handled by an *interrupt handler* (sometimes called an *interrupt service routine*, *error handler*, or an *exception handler*)—a privileged program that determines the nature of the exception and decides how to respond. Interrupts are thus the mechanisms by which a processor identifies the exception and passes control to the interrupt handler.

Because they are closely related, the terms *exception* and *interrupt* are sometimes used interchangeably—either one meaning what the other was originally intended to mean in the *PowerPC Architecture*. For example, the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors* document uses the term *exception* to mean what the *PowerPC Architecture* distinguishes as *exception* and *interrupt*. That PEM document also uses the term *interrupt* to mean an *asynchronous exception*. The original meanings, which are used in this section, are that one term (*exception*) names the cause of an event, and the other term (*interrupt*) names the response to that event.

In the Cell Broadband Engine Architecture (CBEA) processors[1], exceptions can arise due to the execution of PPE instructions or from the occurrence of external events, including the execution of synergistic processor unit (SPU) instructions or memory flow controller (MFC) commands or from an I/O or memory device. All interrupts, except those resulting from SPE events (*Section 18 SPE Events* on page 471), are directed to and handled by the PPE or an external interrupt controller.

The CBEA processors provide facilities for:

- Routing interrupts and interrupt-status information to the PPE or to an external interrupt controller attached to one of the two I/O interfaces
- Prioritizing interrupts presented to the PPE or to an external interrupt controller
- Generating interprocessor interrupts between one PPE thread and another

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

The PPE is supported by two internal interrupt controllers (IICs), one for each PPE thread. The IICs generate, route, and maintain interrupt information for the PPE threads. Each SPE has a set of interrupt registers for masking interrupt conditions, holding the status of interrupt conditions, and routing interrupts.

The *PowerPC Operating Environment Architecture, Book III* and the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors* documents describe PowerPC exception ordering, interrupt definitions, interrupt priorities, interrupt synchronization, and interrupt processing. Similar details regarding CBEA-processor-specific exceptions and interrupts can be found in the *Cell Broadband Engine Architecture* document.

## 9.2 Summary of Interrupt Architecture

The PowerPC Architecture defines two classes of interrupts:

- *Instruction-Caused*—Interrupts caused by the execution of a PowerPC instruction, and thus synchronous with the PowerPC instruction pipeline. In the CBEA processors, these include the execution of all PPE instructions.
- *System-Caused*—Interrupts caused by an event external to the PowerPC instruction pipeline, or otherwise not synchronous with the execution of PowerPC instructions. From the viewpoint of the PPE, interrupts caused by the execution of an SPU instruction or an MFC command are system-caused.

In addition, the PowerPC Architecture defines two kinds of interrupts caused by instructions:

- *Precise*—These are interrupts in which the architecturally visible processor state (in the CBEA processors, this is the PPE state) is known at the time of the exception that caused the interrupt. After a precise interrupt is handled, execution can be resumed where it left off. A precise interrupt is caused by an exception that was generated when the instruction was fetched or executed.
- *Imprecise*—These are interrupts in which the architecturally visible processor state (PPE state) is not guaranteed to be known at the time of the exception that caused the interrupt.

The PPE exception and interrupt architecture is consistent with these PowerPC Architecture definitions. In the CBEA processors, instruction-caused (synchronous) interrupts are precise with respect to the PPE[2], and system-caused (asynchronous) interrupts are imprecise with respect to the PPE[3].

*Table 9-1* on page 242 summarizes the characteristics of interrupts supported by the PPE:

- *Precise*—As defined. For consistency across all interrupts, this parameter is identified in *Table 9-1* for both instruction-caused and system-caused interrupts.
- *Maskable*—With an interrupt-mask bit and a corresponding exception-status bit that can be read by polling, as an alternative to taking an interrupt when an exception occurs. Software can delay the generation of these interrupts by masking them.
- *Can Be Disabled*—Prevented from generating an exception or a corresponding interrupt.

2. In the *PowerPC Architecture*, certain floating-point exceptions can be imprecise, based on mode settings. However, in the PPE, floating-point exceptions are always precise.
3. But see *Section 9.5.2* on page 249 for details about the precise form of the machine check interrupt.

- *Context-Synchronizing*—Previously issued instructions are completed before continuing with program execution. Context-synchronizing instructions empty (flush) the instruction-prefetch queue and start instruction fetching in the context established after all preceding instructions have completed execution. See *Section 20.1.2.4* on page 566 for details about context-synchronization of PPE interrupts.

In *Table 9-1* on page 242, the ability to mask and disable interrupts assumes that the underlying function that can cause such an interrupt is enabled. For example, floating-point unavailable interrupts can be prevented by disabling the underlying function of executing any floating-point instructions (`MSR[FP]` = '0'); however, *Table 9-1* assumes that these functions are enabled.

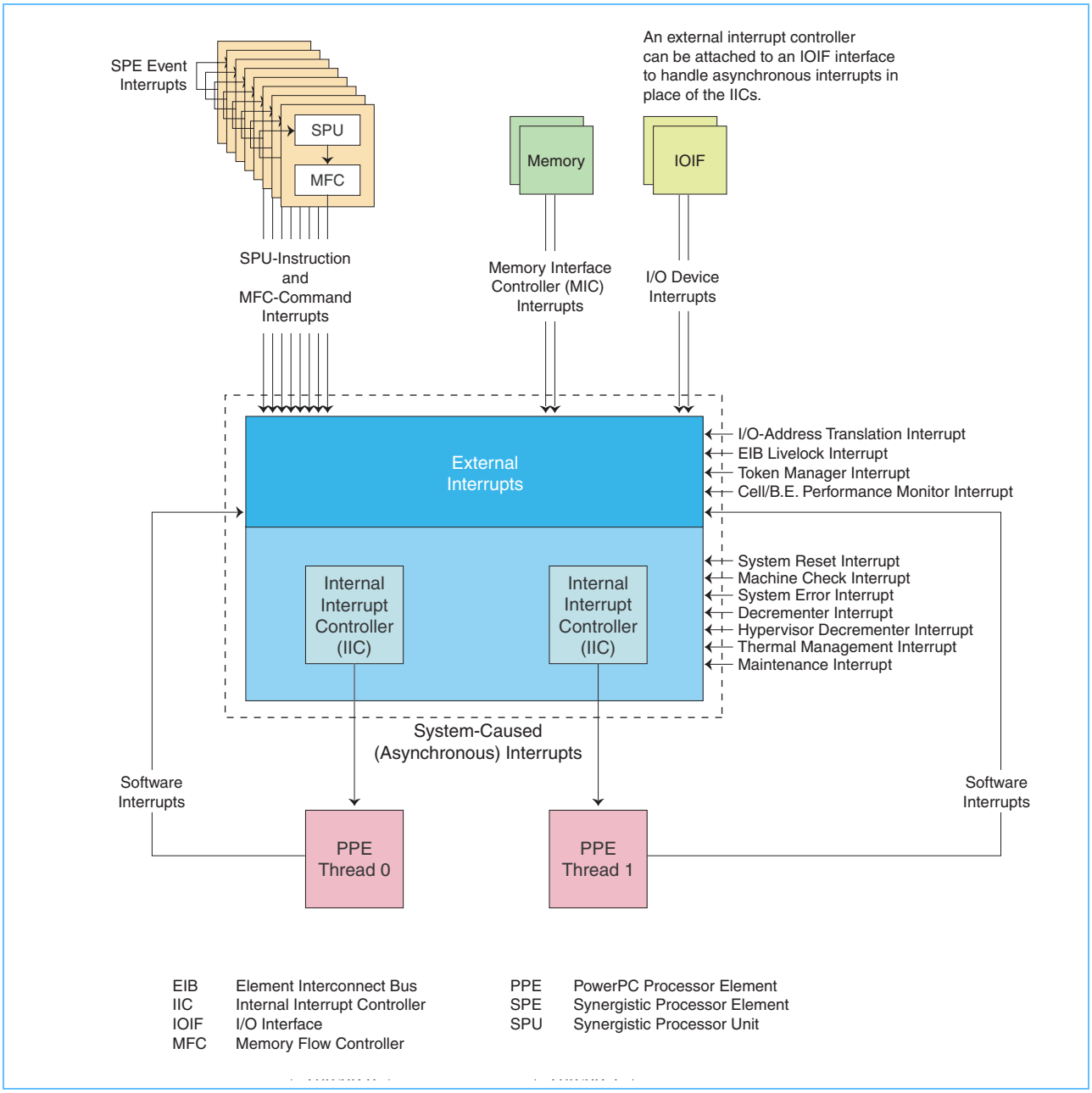
*Table 9-1. PPE Interrupts*

| Interrupt Class | Interrupt | | Precise | Can Be Masked | Can Be Disabled | Context Synchronizing | Reference |
|---|---|---|---|---|---|---|---|
| Instruction-Caused (Synchronous) | Data Storage | | Yes | No | No | Yes | *Section 9.5.3* on page 251 |
| | Data Segment | | | | | | *Section 9.5.4* on page 252 |
| | Instruction Storage | | | | | | *Section 9.5.5* on page 253 |
| | Instruction Segment | | | | | | *Section 9.5.6* on page 254 |
| | Alignment | | | | | | *Section 9.5.8* on page 255 |
| | Program | | | | | | *Section 9.5.9* on page 256 |
| | Floating-Point Unavailable | | | | | | *Section 9.5.10* on page 257 |
| | System Call | | | | | | *Section 9.5.13* on page 258 |
| | Trace | | | | Yes | | *Section 9.5.14* on page 259 |
| | VXU Unavailable | | | | | | *Section 9.5.15* on page 260 |
| | Maintenance (instruction-caused) | | | Yes | | | *Section 9.5.17* on page 261 |
| System-Caused (Asynchronous) | System Reset | | No[1] | No | No[2] | Yes | *Section 9.5.1* on page 248 |
| | Machine Check | | Yes[3] | | Yes | Yes[3] | *Section 9.5.2* on page 249 |
| | System Error | | No[1] | Yes | | Yes | *Section 9.5.16* on page 260 |
| | Decrementer | | | | Yes[4] | | *Section 9.5.11* on page 257 |
| | Hypervisor Decrementer | | | No | Yes | | *Section 9.5.12* on page 258 |
| | Thermal Management | | | Yes | | | *Section 9.5.18* on page 263 |
| | Maintenance (system-caused)[5] | | | No | | | *Section 9.5.17* on page 261 |
| | External (direct and mediated)[6] | SPU Instructions | | Yes[7,8] | Yes[4] | | *Section 9.6.3* on page 271 |
| | | MFC Commands | | | | | |
| | | Memory Interface Controller (MIC) Auxiliary Trace Buffer Full | | Yes | | | *Section 9.6.4.2* on page 273 |
| | | I/O devices (IOIF0, IOIF1) | | Yes[8,9] | Yes[4,8] | | *Section 9.6.4.3* on page 273 |
| | | I/O-Address Translation | | Yes[8] | Yes[4] | | *Section 9.6.4.4* on page 274 |
| | | Element Interconnect Bus (EIB) Possible Livelock Detection | | | | | *Section 9.6.4.5* on page 275 |
| | | Token Manager | | | | | *Section 9.6.4.6* on page 276 |
| | | Performance Monitor | | Yes[7,8] | | | *Section 9.6.4.7* on page 276 |
| | | Software Interrupt[10] | | Yes[8] | | | *Section 9.6.2.4* on page 270 |

1. These interrupts are recoverable.
2. But one of the several sources for this interrupt can be disabled.
3. See *Section 9.5.2* on page 249.
4. The interrupt can be disabled with `MSR[EE]`, but it cannot be *independently* disabled without also disabling all External interrupts.
5. A system-caused maintenance interrupt cannot occur without an appropriate boundary-scan debug tool.
6. See *Section 9.5.7* on page 254 for the difference between direct and mediated external interrupts.
7. By means of specific mask and status registers.
8. By the minimum-priority mechanism of the `IIC_CPL0` and `IIC_CPL1` registers (*Section 9.6.2* on page 266).
9. A Southbridge (I/O interface chip) can support enabling and masking.
10. An interrupt to one of the PPE threads using an interrupt generation port (IGP), `IIC_IGP0` or `IIC_IGP1`. Also called an Interprocessor Interrupt (IPI). There is no corresponding exception, and software can disable such interrupts by simply not issuing them.

*Figure 9-1* illustrates the organization of interrupt handling, assuming that the PPE (rather than an external interrupt controller attached to an I/O interface) handles interrupts. The CBEA processors have two internal interrupt controllers (IICs), one for each PPE thread. These controllers service all system-caused (asynchronous) interrupts, including external interrupts.

*Figure 9-1. Organization of Interrupt Handling*




The only interrupts shown in *Figure 9-1* that are not handled by the PPE are those arising from local SPE events. These are shown as “SPE Event Interrupts” in *Figure 9-1*, and they are handled locally by each SPE as described in *Section 18.3 SPU Interrupt Facility* on page 476.

## 9.3 Interrupt Registers

*Table 9-2* summarizes the interrupt-related registers in the PPE. The interrupt registers allow privileged software on the PPE to select which PPE exceptions are allowed to generate an interrupt to the PPE. For details, see the *PowerPC Operating Environment Architecture, Book III,* the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors* document, and the *Cell Broadband Engine Registers* document.

*Table 9-2. PPE Interrupt Register Summary (Sheet 1 of 2)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| **Thread-Dependent Registers[1] (*PowerPC Architected Registers)*** | | | |
| *Machine State Register* | MSR | Defines the state of the processor. When an exception occurs, the contents of MSR are saved in SRR1, and a new set of bits are loaded into MSR as determined by the exception. | R/W |
| *Machine Status Save/Restore Register 0* | SRR0 | For nonhypervisor interrupts, records the effective address (EA) at which an interrupt occurs, and provides the return address to the interrupted program when an **rfid** (or **rfi**) instruction is executed. It also holds the EA for the instruction that follows a System Call (**sc**) instruction. | R/W |
| *Machine Status Save/Restore Register 1* | SRR1 | For nonhypervisor interrupts, records the exception status and the contents of the MSR when an **rfid** (or **rfi**) instruction is executed. | R/W |
| *Hypervisor Machine Status Save/Restore Register 0* | HSRR0 | For hypervisor interrupts, records the EA at which an interrupt occurs, and provides the return address to the interrupted program when an **hrfid** instruction is executed. | R/W |
| *Hypervisor Machine Status Save/Restore Register 1* | HSRR1 | For hypervisor interrupts, records the exception status and the contents of the MSR when an **hrfid** instruction is executed. | R/W |
| *Floating-Point Status and Control Register* | FPSCR | Enables floating-point interrupts. Records exceptions generated by floating-point operations and the type of result produced by a floating-point operation. | R/W |
| *Data Storage Interrupt Status Register* | DSISR | Records the cause of data-storage interrupt (DSI) and alignment exceptions. | R/W |
| *Decrementer* | DEC | Causes a decrementer exception after a programmable delay. | R/W |
| *Data Address Register* | DAR | Records the effective address generated by a memory-access instruction if the access causes an exception (for example, an alignment exception). | R/W |
| *Data Address Breakpoint Register* | DABR<br>DABRX | Specifies an EA for load and store accesses to generate a data-address breakpoint interrupt. | R/W |
| *Address Compare Control Register* | ACCR | Specifies a Data Address Compare match for instruction storage interrupts. | R/W |

1. These registers are duplicated, one for each PPE thread.
2. These registers serve interrupt-handling for both PPE threads.

*Table 9-2. PPE Interrupt Register Summary (Sheet 2 of 2)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| **Thread-Dependent Registers**[1] *(Implementation-specific Registers)* | | | |
| *Logical Partition Control Register* | LPCR | Controls logical (hypervisor) partitioning. Enables mediated external exceptions. Controls whether external interrupts set MSR[HV] to '1' or leave it unchanged. The LPCR is a partially shared between the two PPE threads. | R/W |
| *Thread Status Register Local* | TSRL | Specifies thread priority, and reports forward-progress timer. When a thread reads its own Thread Status Register (TSR), this register is called the Thread Status Register Local (TSRL). When a thread reads the TSR for the other thread, this register is called the Thread Status Register Remote (TSRR). | R/W |
| *Thread Status Register Remote* | TSRR | | |
| **Thread-Independent Registers**[2] *(**PowerPC Architected Registers**)* | | | |
| *Hypervisor Decrementer* | HDEC | Provides a means for the hypervisor to manage timing functions independently of the decrementer, which is managed by virtual partitions. | R/W |
| *Control Register* | CTRL | Control and status for threads. | R/W |
| **Thread-Independent Registers**[2] *(Implementation-specific Registers)* | | | |
| *Hardware Implementation Dependent Register 0* | HID0 | Enables thermal management interrupt, system error interrupt, extended external interrupt, precise or imprecise form of the machine check interrupt, and attention instruction. | R/W |
| *Hardware Implementation Dependent Register 1* | HID1 | Enables configuration ring system reset interrupt address register. Forces instruction-cache parity error. Controls trace bus. | R/W |
| *Thread Switch Control Register* | TSCR | Enables decrementer wakeup, priority boost for system-caused interrupts. | R/W |
| *Thread Switch Time-Out Register* | TTR | Specifies thread time-out flush values. | R/W |

1. These registers are duplicated, one for each PPE thread.
2. These registers serve interrupt-handling for both PPE threads.

## 9.4 Interrupt Handling

PPE interrupts are handled according to the rules defined in *PowerPC Operating Environment Architecture, Book III*. The interrupts are handled in either privileged (supervisor) or hypervisor mode, depending on the interrupt. In this section, the phrase *the interrupt is taken* refers to the *PowerPC Architecture* interrupt facility.

Upon entering the privileged state, a small portion of the processor's current state is saved in the Machine Status Save/Restore Registers (SRR0, SRR1, HSRR0, and HSRR1). The Machine State Register (MSR) is updated, and instruction fetch and execution resumes at the real address associated with the interrupt (called the *interrupt vector*). Because the Machine Status Save/Restore

Registers are serially reused by the PPE, all interrupts except System Reset and Machine Check are ordered as defined in the *PowerPC Architecture*. The save/restore registers are described in the *Cell Broadband Engine Registers* specification.

Either the SRR0 or the HSRR0 register is set to the effective address of the instruction where processing should resume when returning from the interrupt handler. Depending on the interrupt type, the effective address might be the address of the instruction that caused the exception or the next instruction the processor would have executed if the exception did not exist. All interrupts, except the imprecise form of the machine check interrupt, are context-synchronizing as defined in *PowerPC Operating Environment Architecture, Book III*. Essentially, all instructions in the program flow preceding the instruction pointed to by SRR0 or HSRR0 have completed execution and no subsequent instruction has begun execution when the interrupt is taken. The program restarts from the address of SRR0 or HSRR0 when returning from an interrupt (the execution of an **rfid** or **hrfid** instruction).

Because the MSR setting is modified by hardware when an interrupt is taken, the SRR1 or the HSRR1 register is used to save the current MSR state and information pertaining to the interrupt. Bits [33:36] and [42:47] of the SRR1 or HSRR1 registers are loaded with information specific to the interrupt type. The remaining bits in the SRR1 or the HSRR1 register (bits [0:32], [37:41], and [48:63]) are loaded with a copy of the corresponding bits in the MSR. The MSR bits saved in the SRR1 or HSRR1 registers are restored when returning from an interrupt (the execution of a **hrfid** or **rfid** instruction).

Each PPE thread is viewed as an independent processor complete with separate exceptions and interrupt handling. Exceptions can occur simultaneously for each thread. The PPE supports concurrent handling of interrupts on both threads by duplicating some registers defined by the PowerPC Architecture. The registers associated with Interrupt handling are summarized in *Table 9-2* on page 244. See *Table 9-38* on page 291 for the tread targets of various interrupt types.

See *Section 9.6 Direct External Interrupts* on page 265 for further details about handling external interrupts.

For details about PowerPC interrupt handling, see the *PowerPC Operating Environment Architecture, Book III* and the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors* documents. For details about CBEA-specific interrupt handling, see the *Cell Broadband Engine Architecture* document.

## 9.5 Interrupt Vectors and Definitions

*Table 9-1* on page 242 lists the types of interrupts caused by PPE instructions and system conditions. *Table 9-3* on page 247 provides references to interrupt vectors and the detailed subsections that describe the PPE interrupt causes and behaviors.

The PPE does not implement the optional performance monitor interrupt defined in the PowerPC Architecture (although the CBEA processors do implement an unrelated performance monitor interrupt). In addition, the PPE does not implement the optional example extensions to the trace facility as outlined in an appendix to the *PowerPC Operating Environment Architecture, Book III*.

Further descriptions of the PowerPC interrupts and their causes are given in the PowerPC Architecture*, Book I and Book III* and the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors* document. For a list identifying PowerPC and CBEA processor interrupt differences, see *Section 9.17 Relationship Between CBEA Processor and PowerPC Interrupts* on page 296.

*Table 9-3. Interrupt Vector and Exception Conditions*

| Interrupt Type | Interrupt Vector Effective Address | Causing Conditions (Exceptions) |
|---|---|---|
| System Reset | Selectable, based on `HID1[dis_sysrst_reg]` | See *Section 9.5.1* on page 248. |
| Machine Check | x‘00..00000200’ | See *Section 9.5.2* on page 249. |
| Data Storage | x‘00..00000300’ | See *Section 9.5.3* on page 251. |
| Data Segment | x‘00..00000380’ | See *Section 9.5.4* on page 252. |
| Instruction Storage | x‘00..00000400’ | See *Section 9.5.5* on page 253. |
| Instruction Segment | x‘00..00000480’ | See *Section 9.5.6* on page 254. |
| External | x‘00..00000500’ | See *Section 9.5.7* on page 254. |
| Alignment | x‘00..00000600’ | See *Section 9.5.8* on page 255. |
| Program | x‘00..00000700’ | See *Section 9.5.9* on page 256. |
| Floating-Point Unavailable | x‘00..00000800’ | See *Section 9.5.10* on page 257. |
| Decrementer | x‘00..00000900’ | See *Section 9.5.11* on page 257. |
| Hypervisor Decrementer | x‘00..00000980’ | See *Section 9.5.12* on page 258. |
| System Call | x‘00..00000C00’ | See *Section 9.5.13* on page 258. |
| Trace | x‘00..00000D00’ | See *Section 9.5.14* on page 259. |
| Performance Monitor | x‘00..00000F00’ | This exception is optional in the *PowerPC Architecture*. It is not implemented by the PPE. |
| VXU Unavailable | x‘00..00000F20’ | See *Section 9.5.15* on page 260. |
| System Error | x‘00..00001200’ | See *Section 9.5.16* on page 260.<br>This exception is not defined in the *PowerPC Architecture*; it is specific to the processor implementation. |
| Maintenance | x‘00..00001600’ | See *Section 9.5.17* on page 261.<br>This exception is not defined in the *PowerPC Architecture*; it is specific to the processor implementation. |
| Thermal Management | x‘00..00001800’ | See *Section 9.5.18* on page 263.<br>This exception is not defined in the *PowerPC Architecture*; it is specific to the processor implementation. |

**Note:** Interrupt vectors not listed in this table are reserved.

### 9.5.1 System Reset Interrupt (Selectable or x‘00..00000100’)

The system reset interrupt can be used for power-on reset (POR), to restart a running processor, and to resume a suspended thread. A system reset interrupt is caused by the assertion of one of the system reset interrupt signals from the internal logic or if an exception occurs for a suspended thread.

On a system reset interrupt, the PPE starts thread 0 from the address specified in the *PPE SReset Vector* field of the configuration ring (see the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*). As a result of the system reset interrupt, the Hypervisor and 64-Bit-Mode bits, `MSR[HV]` and `MSR[SF]`, are both set to ‘1’, so that the PPE comes up in hypervisor mode.

A system reset interrupt is generated for a suspended thread if a thermal management, system error, external, or decrementer interrupt exception exists, or if a thread is enabled by a write to the Control Register (`CTRL`). The system reset interrupt vector’s real address for thread 0 is selected based on the setting of `HID1[dis_sysrst_reg]`. If `HID1[dis_sysrst_reg]` is set to ‘1’, the interrupt vector is x‘000..0100’. If `HID1[dis_sysrst_reg]` = ‘0’, bits [22:61] of the interrupt vector are defined by the configuration ring and bits [0:21] and bits [62:63] are set to ‘0’. The interrupt vector for thread 1 is fixed at x‘000..0100’. For details on configuration-ring settings, see *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

When a thread is resumed, the PPE sets the corresponding Thread Enable bit in the `CTRL` register to ‘1’. Resuming a thread due to external and decrementer interrupts is maskable by `TSCR[WEXT]` plus `TSCR[WDEC0]` for thread 0 or `TSCR[WDEC1]` for thread 1. Resuming a thread due to a system error is maskable by `HID0[syserr_wakeup]`; resuming a thread due to a thermal management interrupt is maskable by `HID0[therm_wakeup]`.

The PPE treats the system reset interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a system reset interrupt, the thread’s priority is set to High, the register state is altered as defined in *Table 9-4* on page 249, and instruction fetch and execution resume at the effective address specified in *Table 9-3* on page 247.

*Table 9-4. Registers Altered by a System Reset Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| Thread-Dependent Registers | | |
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. |
| SRR1 | 0:41<br>45:63 | These bits of SRR1 are set as defined in the PowerPC Architecture.<br>SRR1[62] is always a copy of the MSR[RI] bit at the time of the interrupt. |
| | 42:44 | System Reset Reason Code<br>000 Reserved<br>001 Reserved<br>010 System Reset due to a thermal management interrupt to a suspended thread (HID0[therm_wakeup] must be set to '1')<br>011 System Reset due to a decrementer interrupt to a suspended thread (TSCR[WDEC0] for thread0 or TSCR[WDEC1] for thread1 must be set to '1')<br>100 System Reset due to an external interrupt to a suspended thread (TSCR[WEXT] must be set to '1')<br>101 Thread resumed due to 1) a System Reset due to a power-on reset (POR) or 2) a write to the Control Register (CTRL) with the thread enable bit set. Software can check HID0 and HID1 to see if they are set to all zeros to determine if condition 1 occurred. If HID0 and HID1 are not all zeros, then condition 2 occurred.<br>110 Thread resumed due to a System Error (HID0[syserr_wakeup] must be set to '1')<br>111 Reserved |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |
| TSRL | 11:12 | 11 High thread priority |
| Thread-Independent Register | | |
| CTRL | 8:9 | Bit 8 is the thread enable for thread0, and bit 9 is the thread enable for thread1. Depending on which thread took the interrupt, the corresponding thread enable bit will be set. |

### 9.5.2 Machine Check Interrupt (x'00..00000200')

A machine check interrupt occurs when no higher-priority interrupts exist and a machine check condition occurs. In the PPE, there is one possible source for a machine check interrupt—a data error on a caching-inhibited load (such as a nonexistent memory address). If two Machine Checks occur too close together, a checkstop can occur (if there is not enough time to bring down a partition). A checkstop is a full-stop of the processor that requires a System Reset to recover. If either thread has Machine Checks disabled (MSR[ME] = '0') and a Machine Check occurs, the PPE also enters a checkstop condition. The PPE does not modify the Data Storage Interrupt Status Register (DSISR) or Data Address Register (DAR) when a machine check interrupt is taken.

The *PowerPC Architecture* defines the machine check interrupt as a system-caused interrupt. In the CBEA processors:

- If HID0[en_prec_mchk] = '1' ("precise"), only the thread that causes a machine check interrupt—that is, thread issuing the offending caching-inhibited load instruction—will take the interrupt, and the interrupt will appear to be instruction-caused (synchronous) from the viewpoint of that thread (at the expense of PPE performance), and the interrupt will be context-synchronizing, and the interrupt will be recoverable.

- If HID0[en_prec_mchk] = ‘0’ (“imprecise”), both threads will take a machine check interrupt, and the interrupt will appear to be system-caused (asynchronous) to both threads, and the interrupt will not be context-synchronizing, and the interrupt will not be recoverable.

However, the interrupt—in either its precise or imprecise form—is system-dependent in the sense that it reflects the configuration of the system; thus, the same code sequence on a different system might not cause a machine check interrupt.

The advantage of a precise machine check is that a software-defined interrupt handler can potentially recover from the interrupt with predictable results. An imprecise machine check has the advantage of yielding better hardware performance (for caching-inhibited loads) with the risk of nonrecoverability if a machine check occurs. Precise machine check interrupts are enabled depending on the setting of HID0[en_prec_mchk]. If HID0[en_prec_mchk] = ‘0’, then the processor is in normal operation and machine checks are imprecise and have the potential to cause a checkstop if multiple loads outstanding can cause a machine check. Otherwise, if HID0[en_prec_mchk] = ‘1’, then the processor will flush all instructions after any caching-inhibited load is issued and block them at dispatch until the load completes so that machine checks can be precise.

Software running in imprecise mode that depends on the return of a data error (DERR) on a caching-inhibited load—and thus the imprecise Machine Check—to indicate the existence of a device should perform the load with only a single thread active. If the load is performed while both threads are active, it is possible to receive the interrupt when the other thread is in an unrecoverable state resulting in a checkstop condition. In precise mode, software can run with both threads enabled to poll devices that return a DERR.

The PPE treats the precise machine check interrupt condition as a context-synchronizing operation as defined in the PowerPC Architecture. Imprecise machine checks are not context synchronizing. To guarantee the exception when machine checks are imprecise, a **sync** (L=0) instruction is needed after the caching-inhibited load.

When a condition exists for an imprecise machine check, both threads of the PPE, if active, take the imprecise machine check interrupt. For the thread that issues the imprecise machine-check-causing load, the interrupt appears to be instruction-caused; on the other thread, the interrupt is system-caused.

When the PPE takes a machine check interrupt, the register state is altered as defined in *Table 9-5* on page 251, and instruction fetch and execution resume at the effective address specified in *Table 9-3* on page 247.

*Table 9-5. Registers Altered by a Machine Check Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. |
| SRR1 | 0:63 | All bits cleared to '0' except those set by the MSR.<br>SRR1 Specifics:<br>[33:36] = 0<br>[42:47] = 0<br>[0:32], [37:41], [48:61], [63] = MSR values.<br>SRR1[62] is set to '0' for an imprecise Machine Check. For a precise Machine Checks, SRR1[62] is set to the MSR value. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.3 Data Storage Interrupt (x'00..00000300')

A data storage interrupt occurs when no higher-priority interrupts exist and an exception condition related to a data access occurs. Data storage interrupts are precise and are not maskable. The data storage interrupt is implemented as defined in the PowerPC Architecture. The following conditions cause a data storage exception:

- When any byte of the storage location specified by a load, store, **icbi**, **dcbz**, **dcbst**, or **dcbf** instruction cannot be translated from a virtual address to a real address. Data translation must be enabled for this exception to occur (MSR[DR] set to '1').
- When the data access violates the storage protection as defined in the PowerPC Architecture. The storage protection is based on the [Ks] and [Kp] bits in the segment lookaside buffer (SLB), and the [PP] bits in the page table entry (PTE) when MSR[DR] is set to '1'. When MSR[DR] is set to '0', a data storage protection violation occurs if LPCR[LPES] bit 1 = '0' and MSR[HV] = '0' or if LPCR[LPES] bit 1 = '1' and the data effective address is greater than or equal to the LPCR[RMLS].
- When a data address compare or data address breakpoint match occurs.

The PowerPC Architecture describes the processor behavior for the following situation as implementation-specific: a data storage interrupt is generated because an **stwcx.** or **stdcx.** instruction is executed, storage is not updated, and there is a data storage exception. For this case, the PPE produces a data storage interrupt, regardless of the condition for update of storage.

A data storage interrupt in the PPE can occur at two points during the translation of a virtual address to a real address, depending on the translation lookaside buffer (TLB) management mode. For software TLB management, the data storage interrupt occurs if the translation is not found in the TLB. Software TLB management is controlled by the LPCR[TL] bit. For hardware TLB management, the interrupt occurs after an unsuccessful search of the page table.

The PPE treats the data storage interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a data storage interrupt, the register state is altered as defined in *Table 9-6* on page 252, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-6. Registers Altered by a Data Storage Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that caused the exception condition. |
| SRR1 | 0:63 | Set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |
| DSISR | 0:31 | Set as defined in the PowerPC Architecture. |
| DAR | 0:63 | Set as defined in the PowerPC Architecture. |

#### 9.5.3.1 *Data Address Breakpoint Exception*

The PPE supports the optional data address breakpoint facility as defined in *PowerPC Operating Environment Architecture, Book III*. This facility is controlled by the Data Address Breakpoint Register (DABR) and the Data Address Breakpoint Register Extension (DABRX). A data address breakpoint match occurs if any byte accessed matches the doubleword address specified in the upper 61 bits of the DABR. The lower three bits of the DABR and the DABRX control the conditions under which a data address breakpoint exception occurs. Each thread has an independent DABR. See *PowerPC Operating Environment Architecture, Book III* for more information about the data address breakpoint facility.

For nonmicrocoded instructions, the PPE does not alter or access any bytes of storage when a DABR match occurs.

In this processor implementation, a DABR match occurs for an **stwcx.** or **stdcx.** instruction regardless of the reservation.

### 9.5.4 Data Segment Interrupt (x‘00..00000380’)

A data segment interrupt occurs when no higher-priority interrupts exist and the effective address of the data access cannot be translated to a virtual address. Data segment interrupts are precise and are not maskable. The data segment interrupt is implemented as defined in the PowerPC Architecture with one exception. The data segment interrupt modifies the DSISR.

The PowerPC Architecture describes the processor behavior for the following situation as implementation-specific: an **stwcx.** or **stdcx.** instruction is executed and a data segment exception is generated. For this case, storage is not updated and the processor will take the data segment interrupt.

The PPE treats the data segment interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a data segment interrupt, the register state is altered as defined in *Table 9-7* on page 253, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-7. Registers Altered by a Data Segment Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that caused the exception condition. |
| SRR1 | 0:63 | Set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |
| DSISR | 0:5 | Set to '0'. |
| | 6 | Set to '1' on any store or **dcbz** instruction. |
| | 7:9 | Set to '0'. |
| | 10 | Set to '1'. |
| | 11:31 | Set to '0'. |
| DAR | 0:63 | Set as defined in the PowerPC Architecture. |

### 9.5.5 Instruction Storage Interrupt (x'00..00000400')

An instruction storage interrupt (ISI) occurs when no higher-priority exception exists and the next instruction cannot be fetched. ISIs are precise and are not maskable. The ISI is implemented as it is defined by the PowerPC Architecture. In general, an instruction storage exception is caused by the following conditions:

- When the virtual address of the next instruction cannot be translated to a real address
- When an instruction fetch violates the storage protection as defined in the *PowerPC Operating Environment Architecture, Book III*.

  The storage protection is based on the [Ks] and [Kp] bits in the SLB, and the [PP] bits in the PTE when MSR[IR] is set to '1'. When MSR[IR] is set to '0', an instruction storage protection violation occurs if either:

  – LPCR[LPES] bit 1 = '0' and MSR[HV] = '0'

  – LPCR[LPES] bit 1 = '1', MSR[HV] = '0', and the instruction effective address is greater than or equal to the LPCR[RMLS]
- When an instruction fetch is attempted from a no-execute page (PTE[N] = '1')
- When an instruction fetch is attempted with translation on (MSR[IR] = '1') from a guarded page (PTE[G] = '1')
- When an instruction is fetched from a no-execute segment (SLBE[N] = '1')

The PPE treats the instruction storage interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes an instruction storage interrupt, the register state is altered as defined in *Table 9-8* on page 254, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-8. Registers Altered by an Instruction Storage Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. If the interrupt occurs due to a branch target fetch, SRR0 is set to the branch target. |
| SRR1 | 0:63 | Set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |
| DSISR | 0:31 | Set as defined in the PowerPC Architecture. |

### 9.5.6 Instruction Segment Interrupt (x‘00..00000480’)

An instruction segment interrupt (ISEG) occurs when no higher-priority exception exists and the fetch the next instruction cannot be performed. Instruction segment interrupts are precise and are not maskable. The ISEG is implemented as defined by the PowerPC Architecture. In general, an ISEG is caused when the effective address of the next instruction cannot be translated to a virtual address.

The PPE treats the ISEG conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes an ISEG, the register state is altered as defined in *Table 9-8*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-9. Registers Altered by an Instruction Segment Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. If the interrupt occurs due to a branch target fetch, SRR0 is set to the branch target. |
| SRR1 | 0:63 | Set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.7 External Interrupt (x‘00..00000500’)

An external interrupt occurs when higher-priority interrupts do not exist, an external interrupt signal is asserted for one of the two execution threads, and external interrupts are enabled. External interrupts are system-caused and can be masked.

The PPE supports two kinds of external interrupts—direct and mediated. Details about these interrupts are given in the following sections:

- *Section 9.6 Direct External Interrupts* on page 265
- *Section 9.7 Mediated External Interrupts* on page 276
- *Section 9.8 SPU and MFC Interrupts Routed to the PPE* on page 280

### 9.5.8 Alignment Interrupt (x‘00..00000600’)

An alignment interrupt occurs when no higher-priority interrupts exist and a data access cannot be performed for one of the following reasons:

- The operand of a floating-point load or store is not aligned, or crosses a virtual page boundary.
- The instruction is **lfd**, **lfdx**, **lfdu**, **lfdux**, **stfd**, **stfdx**, **stfdu**, or **stfdux**, and the operand is in storage that is caching-inhibited, and the operand is word-aligned but not doubleword-aligned.
- The operand of **lmw**, **stmw**, **lwarx**, **ldarx**, **stwcx.**, or **stdcx.** is not aligned.
- The instruction is **lmw**, **stmw**, **lswi**, **lswx**, **stswi**, or **stswx**, and the operand is in storage that is caching-inhibited.
- The operand of a load or store is not aligned and is in storage that is caching-inhibited.
- The operand of **lwarx**, **ldarx**, **stwcx.**, or **stdcx.** is in storage that is caching-inhibited.
- The instruction is **lha**, **lhau**, **lhaux**, **lhax**, **lwa**, **lwaux**, or **lwax**, and the operand is in storage that is caching-inhibited.
- The instruction is **dcbz** and the operand is in storage that is caching-inhibited or the L1 data cache is disabled by means of HID4[en_dcway] = ‘0000’.

Any load or store instruction that does not cause an alignment interrupt, but crosses a virtual page boundary is flushed and sent to microcode. This lowers the performance of such instructions. However, misalignments that are handled by microcode can be forced to cause an alignment interrupt by setting HID4[force_ai] to ‘1’. This rule applies to the following cases:

- Cacheable (I = ‘0’) load/store instruction crossing a 32-byte boundary
- Cacheable (I = ‘0’) floating-point doubleword that is load/store word aligned, but not doubleword aligned
- Cacheable (I = ‘0’) load/store instruction that crosses a virtual page boundary

The PPE treats the alignment interrupt condition as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes an alignment interrupt, the register state is altered as defined in *Table 9-10*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247. The PPE does not modify the DSISR when an alignment interrupt is taken.

*Table 9-10. Registers Altered by an Alignment Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that caused the exception condition. |
| SRR1 | 0:63 | SRR1 is set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294 |
| DAR | 0:63 | Set as defined in the PowerPC Architecture. |

### 9.5.9 Program Interrupt (x‘00..00000700’)

A program interrupt occurs when no higher-priority interrupts exist and one of the following conditions exists (as defined in the *PowerPC Architecture*):

- Floating-point enabled exception
- Illegal instruction
- Privileged instruction
- Trap

The PPE treats all floating-point exception modes, if enabled, as precise.

The PPE invokes the Illegal Instruction type of program interrupt when it detects any instruction from the illegal instruction class, defined in *PowerPC User Instruction Set Architecture, Book I*. An **mtspr** or **mfspr** instruction with a reserved special purpose register (SPR) selected causes an Illegal Instruction type of program interrupt. In addition to these conditions, the PPE can be configured using the hardware implementation dependent (HID) registers to cause an Illegal Instruction type of program interrupt for the following condition:

- If HID0[en_attn] = ‘0’, the **attn** instruction causes an illegal instruction type of program interrupt.

The following cases list other implementation-specific causes of the illegal instruction type of program interrupt (in addition to the cases listed in the architecture):

- If instructions are stored into the instruction stream (for example, self-modifying code), it is possible for an application error to cause a nonmicrocoded instruction to be sent to the microcode engine.
- Load or Store with Update invalid forms (when RA = ‘0’ or RA = RT).
- Load Multiple invalid forms (when RA is in the range of registers to be loaded or when RA = ‘0’).
- Load String invalid forms (when RA or RB is in the range of registers to be loaded, including the case in which RA = ‘0’, or when RT = RA or RT = RB).

  ***Note:*** *For the invalid forms where RA is in the range of registers to be loaded, the PPE completes the load up unto the point at which the range collision occurs. The interrupt is then taken at the point of collision.*
- Load or Store Floating-Point with Update invalid forms (when RA = ‘0’).
- **bcctr[l]** invalid form (BO[2] = ‘0’).

The PPE invokes the privileged instruction type of program interrupt for the following cases (in addition to any cases listed in the architecture):

- A read or write is performed in problem state to a supervisor or hypervisor special purpose register.
- A read or write is performed in problem or supervisor state to a hypervisor special purpose register.
- A privileged instruction such as **rfid** or **hrfid** is executed without proper privileges.

The PPE treats the program interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a program interrupt, the register state is altered as defined in *Table 9-11*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-11. Registers Altered by a Program Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | SRR0 is set as defined in the PowerPC Architecture. |
| SRR1 | 0:63 | SRR1 is set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.10 Floating-Point Unavailable Interrupt (x‘00..00000800’)

The floating-point unavailable interrupt is implemented as defined in the PowerPC Architecture.

The PPE treats the floating-point interrupt as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a floating-point unavailable interrupt, the register state is altered as defined in *Table 9-12*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-12. Registers by a Floating-Point Unavailable Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that caused the exception condition. |
| SRR1 | 0:63 | SRR1 is set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.11 Decrementer Interrupt (x‘00..00000900’)

The decrementer interrupt is implemented as defined by the PowerPC Architecture. Decrementer interrupts are system-caused. A decrementer exception condition exists if the most-significant bit of the decrementer is set to ‘1’. A decrementer interrupt is taken if the decrementer interrupt is enabled (`MSR[EE]` = ‘1’).

The PPE treats the decrementer interrupt as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a decrementer interrupt, the register state is altered as defined in *Table 9-13*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-13. Registers Altered by a Decrementer Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception conditions were present. |
| SRR1 | 0:63 | Set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.12 Hypervisor Decrementer Interrupt (x‘00..00000980’)

The PPE implements a hypervisor decrementer that is shared for both threads and a corresponding hypervisor decrementer interrupt. These facilities are part of the basic logical partitioning (LPAR) extension to the PowerPC Architecture. A hypervisor decrementer interrupt occurs when no higher-priority exception exists, a hypervisor decrementer exception condition exists, and hypervisor decrementer interrupts are enabled. Hypervisor interrupts are enabled when the following expression is equal to ‘1’:

```
(MSR[EE] | NOT(MSR[HV])) & LPCR[HDICE]
```

A hypervisor decrementer exception condition exists if the most-significant bit of the hypervisor decrementer is set to ‘1’. Hypervisor decrementer interrupts are system-caused and are maskable. If a thread is suspended, it will not be resumed by a hypervisor decrementer exception. If both threads are active and have interrupts enabled, then both threads will take the hypervisor decrementer interrupt.

The PPE treats the hypervisor decrementer interrupt as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a hypervisor decrementer interrupt, the register state is altered as defined in *Table 9-14*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-14. Registers Altered by a Hypervisor Decrementer Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| HSRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception conditions were present. |
| HSRR1 | 0:63 | HSRR1[33] is set to ‘1’.<br>**Note:** Setting HSRR1[33] to ‘1’ for this interrupt is a deviation from the PowerPC Architecture for this implementation.<br>HSRR1[34:36] and [42:47] are set to ‘0’.<br>All other bits are set to the corresponding value in the MSR. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.13 System Call Interrupt (x‘00..00000C00’)

The system call interrupt is implemented as defined in the PowerPC Architecture. System call interrupts are caused by the **sc** instruction; they are precise and cannot be masked.

The PPE treats the execution of the **sc** instruction, which is the condition for the interrupt, as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a system call interrupt, the register state is altered as defined in *Table 9-15* on page 259, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

The **sc** instruction has a LEV bit which, when set to ‘1’, causes the Hypervisor bit, `MSR[HV]`, to be set to ‘1’ so that the CBEA processor runs in hypervisor mode.

*Table 9-15. Registers Altered by a System Call Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception conditions were present. |
| SRR1 | 0:63 | Set as defined in the PowerPC Architecture. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.14 Trace Interrupt (x'00..00000D00')

A trace interrupt occurs when no higher-priority interrupts exist, and any instruction except **rfid**, **hrfid**, **sc**, or **attn** completes successfully, and MSR[SE] is set to '1'; or a branch instruction completes successfully and MSR[BE] is set to '1'. The trace interrupt is implemented as defined in the PowerPC Architecture. Trace interrupts are instruction-caused, precise, and cannot be masked. The PPE does not implement the example extensions (optional) to the trace facility as outlined in the Example Trace Extensions (Optional) section of *PowerPC Operating Environment Architecture, Book III*.

The PPE generates a trace interrupt when an **mtmsrd** instruction sets MSR[SE] to '1'. In this case, SRR0 is set to the effective address of the **mtmsrd** instruction instead of the next instruction that the processor would have attempted to execute if no exception condition were present. This occurs only if the prior state of MSR[SE] was set to '0'. If an **mtmsrd** instruction changes the MSR[SE] bit from '1' to '0', the PPE does not generate a trace interrupt on the **mtmsrd** instruction.

The architecture specifies that the value of the MSR[SE] bit before an **mtmsrd** is executed is used to determine if a trace interrupt is taken after the **mtmsrd** instruction is executed. This implementation does the opposite; the value after the **mtmsrd** instruction executes is used to determine if a trace interrupt is taken after the **mtmsrd** instruction.

If trace interrupts are enabled and an **mtctrl** instruction that disables a thread is issued, then the thread will be disabled and the value of SRR0 is undefined.

The PPE treats the trace interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a trace interrupt, the register state is altered as defined in *Table 9-16*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-16. Registers Altered by a Trace Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception conditions were present. In the case of an **mtmsrd** instruction that causes a trace interrupt by changing the state of MSR[SE] from '0' to '1', SRR0 is set to the effective address of the **mtmsrd** instruction. |
| SRR1 | 0:63 | SRR1 bits [33:36], [42], and [44:47] are set to '0'.<br>SRR1 bit [43] is set to '1'.<br>All other bits are set to the corresponding value in the MSR. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.15 VXU Unavailable Interrupt (x‘00..00000F20’)

A vector/SIMD multimedia extension unit (VXU) unavailable interrupt occurs when a higher-priority interrupt does not exist, an attempt is made to execute a VXU instruction, and MSR[VXU] is set to ‘0’. The VXU unavailable interrupt is a PPE-specific interrupt. The interrupt is not defined in the *PowerPC Architecture*, but it is defined in the vector/SIMD multimedia extension to the *PowerPC Architecture*. VXU unavailable interrupts are instruction-caused, precise, and cannot be masked.

The PPE treats the VXU unavailable interrupt as a context-synchronizing operation as defined in the *PowerPC Architecture*. When the PPE takes a VXU unavailable interrupt, the register state is altered as defined in *Table 9-17*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-17. Registers Altered by a VXU Unavailable Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. |
| SRR1 | 0:63 | SRR1[33:36] and [42:47] are set to ‘0’. All other bits are set to the corresponding value in the MSR. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

For additional information about this interrupt, see the *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual*.

### 9.5.16 System Error Interrupt (x‘00..00001200’)

A system error interrupt occurs when no higher-priority interrupts exist, a system error interrupt signal is asserted for one of the execution threads, HID0[en_syserr] = ‘1’, and `MSR[EE]` = ‘1’ or `MSR[HV]` = ‘0’. The system error interrupt is a PPE-specific interrupt. The interrupt is not defined in the PowerPC Architecture. System error interrupts are system-caused and are enabled if the following expression is equal to ‘1’:

```
(MSR[EE] | NOT(MSR[HV])) & HID0[en_syserr]
```

The PPE treats the system error interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a system error interrupt, the register state is altered as defined in *Table 9-18*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-18. Registers Altered by a System Error*

| Register | Bits | Setting Description |
|---|---|---|
| HSRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. |
| HSRR1 | 0:63 | HSRR1 bit 33 is set to ‘1’ and bits [34:36] and [42:47] are set to ‘0’. All other bits are set to the corresponding value in the MSR. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294 |

The IOC Fault Isolation Register (IOC_FIR) acts as a collection point for error conditions from several sources within the CBEA processors when these errors are external to the PPE. See the *Cell Broadband Engine Registers* document for more information about the IOC_FIR.

Bits in the IOC_FIR are set as a result of errors detected by the EIB, bus interface controller (BIC), and I/O interface controller (IOC). A subset of the bits can cause a system error interrupt in the PPE. The following IOC-detected errors can result in a System Error signal being asserted if the corresponding IOC Fault Isolation Error Mask Register (IOC_FIR_ErrMask) bit is '0', the corresponding IOC Fault Isolation Checkstop Enable Register (IOC_FIR_ChkStop_en) bit is '0', and the corresponding IOC Fault Isolation System Error Enable Register (IOC_SysErrEn) bit is '1':

- The time that a CBEA-processor-originated command to IOIF0 (if IOC_FIR_ErrMask[yi5] = '0', IOC_FIR_ChkStop_en[yi5] = '0', and IOC_SysErrEn[Q0] = '1') or to IOIF1 (if IOC_FIR_ErrMask[yi11] = '0', IOC_FIR_ChkStop_en[yi11] = '0', and IOC_SysErrEn[Q1] = '1') has been in the IOC outbound command queue exceeds the time specified in the memory-mapped I/O (MMIO) register, IOCmd Configuration (IOC_IOCmd_Cfg). See the *Cell Broadband Engine Registers* document for more information about IOC_IOCmd_Cfg.
- The time that a command from IOIF0 (if IOC_FIR_ErrMask[yi4] = '0', IOC_FIR_ChkStop_en[yi4] = '0', and IOC_SysErrEn[A0] = '1') or from IOIF1 (if IOC_FIR_ErrMask[yi10] = '0', IOC_FIR_ChkStop_en[yi10] = '0', and IOC_SysErrEn[A1] = '1') has been in the IOC inbound command queue exceeds the time specified in IOC_IOCmd_Cfg.
- An ERR response was received on a command from the CBEA processor to IOIF0 (if IOC_FIR_ErrMask[yi3] = '0', IOC_FIR_ChkStop_en[yi3] = '0', and IOC_SysErrEn[E0] = '1') or from the CBEA processor to IOIF1 (if IOC_FIR_ErrMask[yi9] = '0', IOC_FIR_ChkStop_en[yi9] = '0', and IOC_SysErrEn[E1] = '1').

If a system error occurs, both threads are signaled. If the system error is signaled to a PPE thread, whether a system reset interrupt occurs first or if a system error interrupt is taken depends on the settings of various SPR bits in the PPE. See *Section 9.5.1 System Reset Interrupt (Selectable or x'00..00000100')* on page 248 for more information about interrupts due to an external interrupt signal's being asserted. If a system error interrupt occurs, another system error interrupt is not taken until the System Error signal is deasserted and is then asserted again.

The System Error signal goes through the power management logic to synchronize any change in this signal relative to power management state changes. The Pause(0) state is exited when a system error interrupt occurs. See *Section 15 Power and Thermal Management* on page 429 for more information about power management.

### 9.5.17 Maintenance Interrupt (x'00..00001600')

The maintenance interrupt is intended for hardware debugging purposes only. This interrupt is not intended for normal programming use. The maintenance interrupt is a PPE-specific interrupt. The interrupt is not defined in the PowerPC Architecture. System-caused maintenance interrupts are enabled if the following expression is equal to '1':

```
(MSR[EE] | NOT(MSR[HV])
```

Instruction-caused maintenance interrupts cannot be masked.

A system-caused maintenance interrupt cannot occur without an appropriate boundary-scan debug tool. A system-caused maintenance interrupt occurs when a higher-priority interrupt does not exist, a trace signal is asserted, and system-caused maintenance interrupts are enabled. The PPE has two trace signals. Each trace signal can be configured through the hardware-debug scan chain to trigger a system-caused maintenance interrupt on one or both threads.

An instruction-caused maintenance interrupt occurs when a higher-priority interrupt does not exist, and the interrupt is triggered by one of the following cases:

1. A completing instruction matches one of two opcode compare and mask registers (opcode compare)
2. An internal instruction address breakpoint (IABR) that is accessed by the hardware-debug scan chain
3. An internal data address breakpoint (DABR) also accessed by the scan chain.

The internal hardware data address breakpoints (set using TDABR and TDABRX) are not the same as the software data address breakpoints (set using DABR and DABRX) defined in the PowerPC Architecture. The facilities are similar in both name and function; however, one is meant for software debug and the other is meant for hardware debug. The internal hardware instruction address breakpoint facility (set using TIABR and TIABRX) works identically to the hardware data address breakpoint except that it is triggered by instruction addresses rather than data addresses.

The PPE treats the maintenance interrupt as a context synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a maintenance interrupt, the register state is altered as defined in *Table 9-19* on page 263, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247. When a maintenance interrupt occurs, HSRR1 reports all conditions that are true for the instruction taking the interrupt, not just the condition that caused it.

*Table 9-19. Registers Altered by a Maintenance Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| HSRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. |
| HSRR1 | 0:32 | Set to the corresponding value in the MSR. |
| | 33 | Source of maintenance interrupt<br>0 The interrupt is caused by an instruction (also called an instruction-caused interrupt).<br>1 The interrupt is caused by the system (also called a system-caused interrupt). |
| | 34 | Set to '0'. |
| | 35:36 | If the interrupt is an instruction-caused interrupt (HSRR1[33] = '0'):<br>10 Opcode 0 Compare<br>01 Opcode 1 Compare<br>11 Opcode 0 and Opcode 1 Compare<br>**Note:** If a trace match is set up to cause a maintenance interrupt, and also **iabr**, **dabr**, or opcode compare is set up, then when a maintenance interrupt caused by an **iabr**, **dabr**, or opcode compare occurs, the opcode compare bits can be falsely set by a previous maintenance interrupt caused by a trace match.<br>If the interrupt is a system-caused interrupt (HSRR1[33] = '1'):<br>00 Reserved<br>01 Trace 0 input asserted<br>10 Trace 1 input asserted<br>11 Trace 0 and Trace 1 input asserted |
| | 37 | Set to '0'. |
| | 38 | Set to the corresponding value in the MSR. |
| | 39:44 | Set to '0'. |
| | 45 | Set to '1' for a hardware data address breakpoint (HSRR1[33] is set to '0')<br>The breakpoint address is loadable only through scan.<br>**Note:** The hardware data address breakpoint reported by this bit is not the same as the DABR defined in *Section 9.5.3.1 Data Address Breakpoint Exception* on page 252. |
| | 46 | Set to '1' for a hardware instruction address breakpoint. (HSRR1[33] is set to '0')<br>The breakpoint address is loadable only through scan. |
| | 47 | Set to '0'. |
| | 48:63 | Set to the corresponding value in the MSR. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

### 9.5.18 Thermal Management Interrupt (x'00..00001800')

A thermal management interrupt occurs when no higher-priority interrupts exist, a thermal management interrupt signal is asserted for one of the two execution threads, and the thermal management interrupts are enabled. Thermal management interrupts are enabled if the following expression is equal to '1':

```
((MSR[EE] | NOT(MSR[HV])) & (HID0[therm_intr_en]))
```

The thermal management interrupt is a PPE-specific interrupt. This interrupt is not defined in the PowerPC Architecture. Thermal management interrupts are system-caused and can be masked.

The PPE treats the thermal management interrupt as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes a thermal management interrupt, the register state is altered as defined in *Table 9-20*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-20. Registers Altered by a Thermal Management Interrupt*

| Register | Bits | Setting Description |
|---|---|---|
| HSRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if no exception condition were present. |
| HSRR1 | 0:32 | Set to the corresponding value in the MSR. |
| | 33 | Set to '1'. |
| | 34:47 | Set to '0'. |
| | 48:63 | Set to the corresponding values in the MSR. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

Ten thermal sensors in the CBEA processors are used to detect the temperature of various parts of the chip. The thermal management control unit (TMCU) provides software with an interface to these sensors. The TMCU can be configured to cause a thermal management interrupt to the PPE. There is a set of registers that provides status of the thermal sensors, a set of interrupt status registers indicating an interrupt condition, and a set of mask registers for the interrupts.

The Thermal Sensor Current Temperature Status Register One (TS_CTSR1) and Thermal Sensor Current Temperature Status Register Two (TS_CTSR2) provide the current temperature information using encoded values to represent specific temperature ranges for the ten thermal sensors. Larger values indicate higher temperatures. The Thermal Sensor Interrupt Temperature Register One (TS_ITR1) and Thermal Sensor Interrupt Temperature Register Two (TS_ITR2) contain the corresponding encoding for the temperature that causes a thermal management interrupt to the PPE. The Thermal Sensor Global Interrupt Temperature Register (TS_GITR) contains a second interrupt temperature level. This register specifies one interrupt temperature level that applies to all sensors in the CBEA processors. When the encoded temperature value in the TS_CTSR1 or TS_CTSR2 for a sensor is greater than or equal to the corresponding sensor's interrupt temperature encoding in the TS_ITR1 and TS_ITR2, the corresponding bit in TS_ISR is set. When the temperature encoding in the TS_CTSR1 and TS_CTSR2 register for any sensor is greater than or equal to the global interrupt temperature encoding in the TS_GITR, the corresponding status bit in TS_ISR is set. If any TS_ISR bit is '1' and the corresponding TS_IMR bit is '1', a thermal management interrupt signal is asserted to the PPE.

To clear the interrupt condition, privileged software should set the corresponding mask bits in the TS_IMR to '0'. To enable a thermal management interrupt, privileged software should ensure the temperature is below the interrupt temperature for the corresponding sensors and then perform the following steps.

1. Write a '1' to the corresponding status bit in the TS_ISR.
2. Write a '1' to the corresponding mask bit in the TS_IMR.

When the temperature is not below the interrupt temperature, enabling the interrupt can result in the generation of an immediate management interrupt.

See *Section 15 Power and Thermal Management* on page 429 for additional details on thermal management interrupts.

## 9.6 Direct External Interrupts

This section describes the direct external interrupts supported by the PPE.

As shown in *Table 9-1* on page 242, the sources or causes of direct external interrupts include:

- SPU Instructions—*Section 9.6.3* on page 271
- MFC Commands—*Section 9.6.3* on page 271
- Memory Interface Controller (MIC) Auxiliary Trace Buffer Full—*Section 9.6.4.2* on page 273
- I/O Devices (IOIF0, IOIF1)—*Section 9.6.4.3* on page 273
- I/O-Address Translation—*Section 9.6.4.4* on page 274
- Element Interconnect Bus (EIB) Possible Livelock Detection—*Section 9.6.4.5* on page 275
- Token Manager—*Section 9.6.4.6* on page 276
- Performance Monitor—*Section 9.6.4.7* on page 276
- Software Interrupt, also called interprocessor interrupt (IPI)—*Section 9.6.2.4* on page 270

The PPE receives these interrupts from:

- An IIC, by means of the EIB, from:
  - An SPU or MFC (all such packets are actually sent by the MFC)
  - An I/O device by means of the IOIF0 or IOIF1 interface
  - The IIC, due to an event recorded by the MIC, EIB, token manager, I/O-address translation logic, or performance monitor
- Software interrupts from one PPE thread to another PPE thread. They are generated by an MMIO write to an Interrupt Generation Port Register (IGP), `IIC_IGP0` or `IIC_IGP1`, in the IIC.

As shown in *Figure 9-1* on page 243, the IICs receive and route these interrupts and interrupt status information to the PPE. The sections that follow describe the organization and handling of the Interrupts.

### 9.6.1 Interrupt Presentation

Interrupts generated by an SPU or MFC are sent to the IICs (or an external interrupt controller) as interrupt packets or dedicated signals on the EIB. External I/O devices send interrupts to the CBEA processors using either interrupt packets or dedicated signals on an I/O interface (IOIF). When an IIC receives an interrupt packet it signals an external interrupt to a PPE thread. If an external interrupt controller is used, the SPEs can be individually configured to send specific classes of interrupts to this external controller.

An IIC receives interrupt packets and signals an external interrupt to the appropriate PPE thread. These IIC signals go through the power management logic to synchronize any changes in these signals relative to power-management state changes. The power management logic also detects that the PPE Pause(0) power management state should be exited when an external interrupt

occurs. When an external interrupt is signaled to a PPE thread, the type of interrupt taken depends on the settings of the special-purpose register (SPR) bits in the PPE. See *Section 15 Power and Thermal Management* on page 429 for more information about power management.

An incoming interrupt packet—generated within the CBEA processor or external to it—destined for a PPE thread will have a unit ID of x'E' or x'F', as shown in *Table 9-21*.

*Table 9-21. Interrupt-Destination Unit ID*

| Destination Unit | Unit ID |
|---|---|
| PPE Thread 0 | x'E' |
| PPE Thread 1 | x'F' |
| External interrupt controller attached to IOIF1, if any | x'B' |
| External interrupt controller attached to IOIF0, if any | x'0' |

The registers related to external interrupt generation, routing, and presentation are summarized in *Section 9.6.2.1 IIC Register Summary* on page 266 and *Section 9.6.4 Other External Interrupts* on page 272. The Current Priority (`IIC_CPL0` and `IIC_CPL1`), Interrupt Pending Port (`IIC_IPP0` and `IIC_IPP1`), Interrupt Generation Port (`IIC_IGP0` and `IIC_IGP1`), IIC Interrupt Routing (`IIC_IRR`), and IIC Interrupt Status Register (`IIC_ISR`) are the most important. For details, see the *Cell Broadband Engine Registers* specification and the *Cell Broadband Engine Architecture* document.

## 9.6.2 IIC Interrupt Registers

### 9.6.2.1 *IIC Register Summary*

*Table 9-22* summarizes the interrupt-related registers—one for each PPE thread—in the IICs, which themselves are part of the I/O interface controller (IOC). For details, see the *Cell Broadband Engine Registers* document and the *Cell Broadband Engine Architecture* document.

*Table 9-22. IOC Interrupt Register Summary (Sheet 1 of 2)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| **Interrupt Registers** *(Implementation-specific Registers)* | | | |
| *IIC Thread 0 Interrupt Pending Port* (nondestructive) | `IIC_IPP0` | Allows software to read the interrupt source and other information about pending interrupts. This register is sometimes referred to as the interrupt pending port (IPP).<br>After software reads this register destructively or nondestructively, the next highest interrupt is loaded into the register. When reading `IIC_IPP0` nondestructively, the value of the associated `IIC_CPL0` register is not updated with the interrupt in `IIC_IPP0`. When reading `IIC_IPP0` destructively, the `IIC_CPL0` register takes on the priority of the interrupt in `IIC_IPP0`.<br>If a nondestructive read of the interrupt pending port that returns a valid interrupt is followed by another read of the interrupt pending port (destructive or nondestructive), the IIC returns the same data value unless another interrupt of a higher priority has been received by the IIC. | R |
| *IIC Thread 0 Interrupt Pending Port* (destructive) | | | R |

*Table 9-22. IOC Interrupt Register Summary (Sheet 2 of 2)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| *IIC Thread 0 Interrupt Generation Port* | IIC_IGP0 | Allows privileged software to generate interrupts to PPE thread 0. | W |
| *IIC Thread 0 Current Priority Level* | IIC_CPL0 | Allows software to mask interrupts of a specified priority for PPE thread 0.<br>Software directly loads the current priority level (CPL) by means of an MMIO write of the register or indirectly by means of destructive read of IIC_IPP0. When executing a destructive read of IIC_IPP0, the priority of the interrupt in IIC_IPP0 is written to the CPL if that interrupt is valid. Otherwise, the content is unchanged. | R/W |
| *IIC Thread 1 Interrupt Pending Port* (nondestructive) | IIC_IPP1 | Same as IIC_IPP0, except for PPE thread 1. | R |
| *IIC Thread 1 Interrupt Pending Port* (destructive) | | | R |
| *IIC Thread 1 Interrupt Generation Port* | IIC_IGP1 | Same as IIC_IGP0, except for PPE thread 1. | W |
| *IIC Thread 1 Current Priority Level* | IIC_CPL1 | Same as IIC_CPL0, except for PPE thread 1. | R/W |
| *IIC Interrupt Routing* | IIC_IRR | Configures the priority and destination of certain interrupts reported in the IIC_ISR. | R/W |
| *IIC_Interrupt Status* | IIC_ISR | Records interrupt conditions from the memory interface, I/O interface, element interconnect bus (EIB), performance monitor, and token manager.<br>If the IIC_ISR is nonzero, the IIC creates a class 1 interrupt to either the EIB, IIC_IPP0, or IIC_IPP1, depending on how the IIC_IRR is configured.<br>After resetting the interrupt condition in the source unit, software can reset the interrupt by writing '1' to the corresponding bit in the IIC_ISR. Writing this register must occur to confirm handling of interrupts in the IIC_ISR. | R/W |
| **Fault Isolation Registers** *(Implementation-specific Registers)* | | | |
| *IOC Fault Isolation* | IOC_FIR | Records, sets, resets, and masks IOC faults. Enables checkstop interrupt for faults.<br>IOC_FIR_SysErrEnbl is not part of other fault isolation registers (FIRs); it is used to generate an enable for the system error interrupt. | R/W |
| *IOC Fault Isolation Register Set* | IOC_FIR_Set | | R/W |
| *IOC Fault Isolation Register Reset* | IOC_FIR_Reset | | R/W |
| *IOC Fault Isolation Error Mask* | IOC_FIR_ErrMask | | R/W |
| *IOC Checkstop Enable* | IOC_FIR_ChkStpEnbl | | R/W |
| *IOC System Error Enable* | IOC_FIR_SysErrEnbl | | R/W |

### 9.6.2.2 *Current Priority Level*

There is one current priority level (CPL) register for each PPE thread in a CBEA processor. IIC_CPL0 holds the current priority level for thread 0, and IIC_CPL1 holds the current priority level for thread 1. The registers are intended to hold the priority level at which software is currently operating. The lower the numeric value of the priority field, the higher the priority level (highest priority corresponds to a numeric value of '0'). If the priority of an interrupt is numerically less than the priority in the CPL register, the interrupt is signaled to the PPE.

The priority-level value can be written explicitly by software to the CPL registers or indirectly by software doing a destructive read of a valid interrupt in the Interrupt Pending Port Register (IPP). Only the most-significant 5 bits of the CPL register priority field are implemented. Unimplemented bits of the interrupt pending port are read as '0'. This implies that `IIC_CPL0[Priority[0:3]]` is set to `IIC_IPP0[Priority[0:3]]`, and `IIC_CPL0[Priority[4]]` is set to '0' when software performs a destructive read of a valid interrupt from `IIC_IPP0`. Likewise, `IIC_CPL1[Priority[0:3]]` is set to `IIC_IPP1[Priority[0:3]]`, and `IIC_CPL1[Priority[4]]` is set to '0' when software performs a destructive read of a valid interrupt from `IIC_IPP1`.

Although the CPL registers have five bits, the IPP registers have only four bits. This extra bit in the CPL registers allows software to set the current priority level lower than any pending interrupt in the IPP registers. This allows for 16 priority levels to be used for these interrupts, whereas only 15 priority levels could be used if the CPL registers had only four bits. If a CPL register had four bits, the lowest priority (highest numeric value) in this register is all ones, and this masks an interrupt in the corresponding IPP register with the same priority value.

#### 9.6.2.3 *Interrupt Pending Port*

The interrupt pending port (IPP) registers allow software to read the interrupt packet data and other information about the highest priority interrupt pending for each PPE thread. The Thread 0 Interrupt Pending Port Register is `IIC_IPP0` and the Thread 1 Interrupt Pending Port Register is `IIC_IPP1`. *Table 9-23* shows the bit fields the registers (reserved fields are not shown).

*Table 9-23. IIC_IPP0 and IIC_IPP1 Interrupt Pending Port Bit Fields*

| Bits | Name | Description |
|---|---|---|
| 32 | V | Interrupt Valid:<br>0 No interrupt pending.<br>1 Interrupt pending. |
| 33 | T | Interrupt Type:<br>0 SPE, external device, or external interrupt controller.<br>1 Thread 0 interrupt generation port. |
| 46:47 | Class | Interrupt Class. Returns zeros when T = '1'. |
| 48:51 | Src_Node_ID | Interrupt Source Node ID. Returns zeros when T = '1'. |
| 52:55 | Src_Unit_ID | Interrupt Source Unit ID. Returns zeros when T = '1'. |
| 56:59 | Priority | Interrupt Priority. |

The interrupt source (ISRC) of an external interrupt is designated as the two 4-bit fields, bits 48:55, shown in *Table 9-24*.

*Table 9-24. Values for Interrupt Source (Sheet 1 of 2)*

| Interrupt Source | Most-Significant Four Bits of Interrupt Source (ISRC) (BIF Node ID) | Least-Significant Four Bits of Interrupt Source (ISRC) (BIF Unit ID) |
|---|---|---|
| SPU(0) | x'4' | x'4' |
| SPU(1) | x'7' | x'7' |
| SPU(2) | x'3' | x'3' |
| SPU(3) | x'8' | x'8' |

*Table 9-24. Values for Interrupt Source (Sheet 2 of 2)*

| Interrupt Source | Most-Significant Four Bits of Interrupt Source (ISRC) (BIF Node ID) | Least-Significant Four Bits of Interrupt Source (ISRC) (BIF Unit ID) |
|---|---|---|
| SPU(4) | x‘2’ | x‘2’ |
| SPU(5) | x‘9’ | x‘9’ |
| SPU(6) | x‘1’ | x‘1’ |
| SPU(7) | x‘A’ | x‘A’ |
| IOC-IOIF1 | x‘B’ | x‘B’ |
| IOC-IOIF0 | x‘0’ | x‘0’ |
| IIC | not applicable | x‘E’ |

There are two types of read, one is nondestructive and the other is destructive. The destructive-read register is located at a doubleword offset to the nondestructive read address. Because of the IICs’ priority system, lower-priority interrupts will not be realized until the status is cleared (that is, reading the status by means of destructive reads).

The following description applies to interrupts for one PPE thread and that thread’s IPP register, CPL register, and IGP register.

The interrupt priority is 8 bits in the *Cell Broadband Engine Architecture*. The CBEA processors implement the 4 most-significant bits of interrupt priority and ignores the 4 least-significant bits of interrupt priority in the IPP register. There are 16 interrupt priorities decoded from these four bits. The IIC implements one interrupt-queue entry for each of the 16 priorities for each thread. In addition, for interrupts created by the MMIO writes to the IGP register, the IIC implements one queue entry for each of the 16 priorities.

When the priority of the highest-priority, valid interrupt in the interrupt-pending queue is higher (lower numeric value) than the priority in the current priority level, the External Interrupt signal to the PPE thread is activated. When the priority of the highest-priority, valid interrupt in the interrupt-pending queue is the same or lower (equal or higher numeric value) than the priority in the CPL register, the External Interrupt signal to the PPE thread is deactivated. If the PPE does not take an interrupt immediately because the external interrupt is disabled or the interrupt is blocked because its priority is not higher than the priority in the CPL register, the interrupt is not lost. In this case, the interrupt is simply deferred until the external interrupt is enabled and the priority in the CPL register is lowered.

After a destructive read of the IPP register, the IIC performs an interrupt-reissue transaction on the EIB if the IIC responded to an EIB interrupt command with a retry snoop response since the last interrupt-reissue transaction that the IIC performed. This means that the number of interrupt-reissue transactions might be less than the number of interrupt transactions that have been retried.

The IIC ignores priority in determining whether an interrupt-reissue transaction should be performed. Thus, the IIC can perform an interrupt-reissue transaction even when the priority of the interrupt that was given a retry response is different from the priority read from the interrupt pending port. In determining whether an interrupt-reissue transaction should be performed, the IIC also ignores the thread destination of the retried interrupt versus the thread whose interrupt pending port was read.

The IIC makes its determination of when an interrupt-reissue transaction occurred at a different point in an EIB transaction than when it determines its retry response for an interrupt command. The result is that more interrupt-reissue transactions might occur than necessary; however, there is not an interrupt transaction that received a retry response without a subsequent interrupt-reissue transaction on the EIB if there is no valid entry in the interrupt pending queue for the same priority as the retried interrupt.

The CBEA processors do not support the *Cell Broadband Engine Architecture* optional interrupt-packet data field in the IPP.

The four most-significant bits of the interrupt source (ISRC) have the same value as the Cell Broadband Engine interface (BIF) node ID of the interrupt source. The four least-significant bits have the same value as the BIF unit ID of the interrupt source in most cases; however, for interrupts that the IIC generates by means of its Interrupt Routing Register (IRR), the least-significant four bits of the ISRC are not related to a BIF unit ID, as shown in *Table 9-24* on page 268.

For interrupt packets from an IOIF, the IOC adjusts the source information when forwarding the interrupt packet to the EIB so that the IOC unit ID for that IOIF appears to be the source. Thus, an IPP register with an ISRC corresponding to IOC-IOIF0 in a specific BIF node represents an interrupt from a device on the IOIF0 connected to that IOC. An IPP value with an ISRC corresponding to IOC-IOIF1 in a specific BIF node represents an interrupt from a device on the IOIF1 connected to that IOC. When the MFC sends the interrupt packet on the internal EIB, the MFC inserts the ISRC that corresponds to itself.

#### 9.6.2.4 *Interrupt Generation Port*

The interrupt generation port (IGP) registers allows privileged software to generate an interrupt packet to a PPE thread. There is one IGP register for each PPE thread—Thread 0 Interrupt Generation Port Register (`IIC_IGP0`) and Thread 1 Interrupt Generation Port Register (`IIC_IGP1`). Software can generate an interrupt packet to a PPE thread by storing to the PPE thread's IGP register. When the interrupt packet is read by means of the appropriate Interrupt Pending Port Register (`IIC_IPP0` or `IIC_IPP1`), the interrupt packet data, class information, and ISRC are read as zeros. This interrupt packet does not need to be transmitted on the internal EIB because the IGP register and the destination of the interrupt packet are both in the same IIC. The least-significant 8 bits written to this register contain the interrupt priority; however, only the most-significant 4 bits of priority are implemented. For each priority level, a latch is used to represent a pending interrupt. A second store to the IGP with the same priority as a pending IGP interrupt results in no state change and, effectively, the interrupts are merged.

An IGP interrupt is treated like any other PPE external interrupt, including the interrupt priority. Unlike other external interrupts, however, this interrupt cannot be routed outside the CBEA processor. For a multi-CBEA-processor system, the registers can be memory mapped, and accessing the appropriate register will interrupt the required PPE thread.

The `IIC_IGP0` and `IIC_IGP1` are write-only MMIO registers. The only field is the priority field, as shown in *Table 9-25* on page 271. Because this interrupt is meant to be a quick way to communicate with the other thread, it is recommended that the external interrupt handler check this status before others.

*Table 9-25. IIC_IGP0 and IIC_IGP1 Interrupt Generation Port Bit Field*

| Bits | Names | Descriptions |
|---|---|---|
| 56:59 | Priority | Interrupt Priority. |

### 9.6.3 SPU and MFC Interrupts

Interrupt packets from the SPUs and MFCs are routed to the PPE by each MFC's Interrupt Routing Register (`INT_Route`). There are no implementation-dependent MFC interrupts; only the interrupts defined in the *Cell Broadband Engine Architecture* are supported. This section gives an overview of the status, masking, and routing of SPU and MFC interrupts. For details about these interrupts, see *Section 9.8 SPU and MFC Interrupts Routed to the PPE* on page 280.

The SPU and MFC interrupts routed to the PPE are independent from and unrelated to the event interrupts local to each SPE, which are described in *Section 18.3 SPU Interrupt Facility* on page 476.

#### 9.6.3.1 *Status and Masking*

There are three interrupt status registers (`INT_Stat_class0`, `INT_Stat_class1`, and `INT_Stat_class2`) and three interrupt mask registers (`INT_Mask_class0`, `INT_Mask_class1`, and `INT_Mask_class2`) in each MFC. Thus, there is an interrupt status register and an interrupt mask register for each class of interrupt (error, translation, and application).

When using these registers, software must adhere to these requirements:

1. When multiple CBEA processors are connected through the BIF and an external interrupt occurs due to a nonzero MFC Interrupt Status Register (INT_Stat_class0, INT_Stat_class1, or INT_Stat_class2), then software must issue a **sync** or **eieio** instruction before a write to the MFC Interrupt Status Register.
2. In both a single-CBEA-processor or multiple-CBEA-processor configuration, if software performs writes to the MFC Interrupt Status Register at any time other than during the external interrupt handling after an external interrupt, then an additional test and another operation are needed. After all writes to the MFC Interrupt Status Register are completed, software must read the MFC Interrupt Status Register and perform a bit-wise AND masking of the status with the contents of the MFC Interrupt Mask Register. If the result is nonzero, software must issue a **sync** or **eieio** instruction and write the MFC Interrupt Status Register again. In such a situation, it is unpredictable whether the correct number of external interrupts required by the Cell Broadband Engine Architecture were triggered. Software should take appropriate actions for the interrupt conditions indicated by the status read and should not assume that an external interrupt occurred for all the status conditions.

   If software writes the MFC Interrupt Status Register in more than one place in the external interrupt handler and cannot guarantee that one of these writes is the first since the external interrupt, then the second requirement applies.

#### 9.6.3.2 *Routing*

The MFC Interrupt Routing Register (INT_Route), described in *Section 9.8.2.3 Interrupt Routing Register* on page 285, allows privileged software to select the PPE thread or external interrupt controller that is to service an interrupt. For each class of interrupt, the register contains a priority and interrupt destination.

Only the most-significant four bits of interrupt priority are implemented. The interrupt destination indicates which logical PPE thread or external interrupt controller is to receive interrupt packets for MFC interrupts of the corresponding interrupt class. The four most-significant bits of the interrupt destination have the same value as the BIF node ID of the interrupt destination. See *Table 9-21* on page 266 for the least-significant four bit values of the required destination. This table shows both the MFC setting of the destination IDs and the IIC interpretation of the destination IDs. When the MFC sends the interrupt packet on the internal EIB, the MFC inserts the interrupt source (ISRC) that corresponds to itself.

### 9.6.4 Other External Interrupts

The IIC Status Register (IIC_ISR) acts as a collection point for interrupt conditions from the MIC, EIB, token manager, I/O address translation, and performance monitor. There are internal signals from the MIC, EIB, token manager, I/O address translation logic, and performance monitor, and these signals indicate the corresponding interrupt conditions to the IIC. When one of these interrupt signals is asserted, the corresponding bit in the IIC_ISR register is set to a ‘1’. When IIC_ISR becomes nonzero, the IIC creates a class 1 interrupt packet.

After resetting the interrupt condition in the MIC, EIB_Int, TKM_ISR, IOC_IO_ExcpStat, or the Performance Monitor Status/Interrupt Mask Register (pm_status), software can reset IIC_ISR by writing a ‘1’ to the corresponding bit position in IIC_ISR.

The MIC Auxiliary Trace Buffer Full interrupt condition can be reset by either:

- Setting MIC_Ctl_Cnfg2[8] = '0'
- If MIC_CTL_CNFG2[7] = '1', waiting for auxiliary trace current address to wrap so that it is no longer equal to the auxiliary trace max address
- If MIC_CTL_CNFG2[7] = '0', writing the auxiliary trace current address so that it is no longer equal to the Auxiliary trace max address

Writing a ‘0’ to the corresponding bit has no effect on that bit.

#### 9.6.4.1 *Routing*

The IIC_IRR register provides the priority and destination of the interrupt. The 4-bit priority field in the IIC_IRR register corresponds to the four bits of the priority field implemented in the IIC_IPP0 and IIC_IPP1 registers. If the interrupt is routed to the PPE thread in the same chip as the IIC, then the interrupt is placed in the appropriate interrupt pending queue for the thread indicated by the IIC_IRR register. If the interrupt is routed to some other destination, an interrupt is sent on the internal EIB. When the IIC sends the interrupt packet on the internal EIB, the IIC inserts the ISRC that corresponds to itself. The IIC’s unit ID in this ISRC is x‘E’, as shown in *Table 9-24* on page 268. If the IIC_IRR register destination node ID specifies some other node, then the inter-

rupt packet is sent to that node across the BIF protocol. The unit ID specifies a destination within a node or to an IOIF attached to the node, according to *Table 9-21* on page 266. If the destination is an external interrupt controller, the interrupt packet is sent across the appropriate IOIF.

The programmable fields in the `IIC_IRR` are shown in *Table 9-26*.

*Table 9-26. IIC_IRR Routing Fields*

| Bits | Names | Descriptions |
|---|---|---|
| 48:51 | Priority | Interrupt Priority. |
| 56:59 | Dst_Node_ID | Destination node ID for interrupts in the IIC_ISR. |
| 60:63 | Dst_Unit_ID | Destination unit ID for interrupts in the IIC_ISR. Interrupt packets routed to IOIF0 have the same destination unit ID as IOC0. Those that are routed to IOIF1 have the same destination unit ID as IOC1. The valid values are:<br>x‘0’ IOC 0 (IOIF0).<br>x‘B’ IOC 1 (IOIF1).<br>x‘E’ PPE thread 0.<br>x‘F’ PPE thread 1. |

### 9.6.4.2 *MIC Auxiliary Trace Buffer Full*

The performance monitor (PM) facility provides a means of tracing various events and counts. The MIC provides a facility to store this trace in a region of memory. The memory trace region is considered full when the Auxiliary Trace Current Address Register is equal to the Auxiliary Trace Max Address. If the memory trace region is full and the MMIO register `MIC_Ctl_Cnfg2[8]` = ‘1’, an MIC Auxiliary Trace Buffer Full signal to the IIC is asserted. When the MIC Auxiliary Trace Buffer Full signal is asserted, `IIC_ISR[60]` is set to ‘1’.

### 9.6.4.3 *I/O Devices*

I/O devices attached to the IOIF0 or IOIF1 interfaces can initiate and receive interrupts, and those interrupts can also be reissued. There are two types of interrupt transactions defined for the IOIF interfaces: Interrupt and Interrupt-reissue.

Interrupt transactions are initiated by sending an interrupt packet on the EIB. *Table 9-37* on page 290 shows the format of the I/O interrupt packet. After receiving such a packet, the IIC asserts the External Interrupt signal to the appropriate PPE thread. The first-level interrupt handler (FLIH) in that PPE thread then reads the Interrupt Pending Port Register (`IIC_IPP0` or `IIC_IPP1`) to determine the interrupting device.

If the target device cannot accept the interrupt, the interrupt command gets a retry response. The interrupting device must hold this interrupt pending until it sees an interrupt-reissue transaction. When an interrupt-reissue occurs, pending interrupts that received a retry since the last interrupt-reissue must be retransmitted if the BIF node ID field in the I/O identifier of the interrupt-reissue command sent by the CBEA processor matches the destination BIF node ID of the pending interrupts.

For additional information about the I/O architecture in general, see *Section 7 I/O Architecture* on page 161.

#### 9.6.4.4 *I/O Address Translation*

An I/O address translation exception can occur due to either an I/O segment fault or I/O page fault. An I/O segment fault type of I/O exception occurs for the following reason:

- There is an I/O translation that attempts to use an I/O segment table entry (IOSTE) that is not valid.

An I/O page fault type of I/O exception occurs for the following reasons when an external device attempts a read or write on the IOIF with translation enabled:

- Assume a page size of $2^p$ is specified in IOSTE[PS]. The CBEA processors support I/O page-table (IOPT) entry format 1. See*Section 7.4 I/O Address Translation* on page 176 for details of the I/O segment table (IOST) and IOPT. If the value of the most-significant bits (28 - p) of the 28 least-significant I/O address bits is greater than 512 times the IOSTE[NPPT], an I/O page-fault type of I/O exception occurs.
- The I/O address is translated using an I/O page-table entry whose I/O Identifier (IOID) field does not match the IOID of the I/O device.
- The I/O address is translated using an I/O page-table entry whose page protection (PP) bits equal ‘00’, or whose PP bits equal ‘01’ and the IOIF transaction is a write, or whose PP bits equal ‘10’ and the IOIF transaction is a read.

When an I/O exception occurs on an I/O access, the following information related to the fault is captured in the I/O Exception Status Register (IOC_IO_ExcpStat):

- I/O Segment and I/O Page Number Address bits
- Access type (read or write)
- I/O device ID

Once an I/O exception occurs, IOC_IO_ExcpStat[V] is set to ‘1’. When IOC_IO_ExcpStat[V]=‘1’, errors that occur are not captured. After handling an I/O exception, software should set the valid bit to ‘0’ to enable capturing of error information about a subsequent I/O exception.

If an I/O segment fault occurs and IOC_IO_ExcpMask[1]=‘1’, `IIC_ISR[61]` is set to ‘1’. If an I/O page fault occurs and IOC_IO_ExcpMask[2]=‘1’, `IIC_ISR[61]` is set to ‘1’. After handling an I/O address translation exception, software should write IOC_IO_ExcpStat[0] to '0' and write a ‘1’ to `IIC_ISR[61]` to reset this bit, It is recommended that software store to IOC_IO_ExcpStat to reset it first, store to `IIC_ISR` next, execute an **eieio** instruction, and then read IOC_IO_ExcpStat again to verify that another I/O Address Translation did not occur between the two stores. Assuming stores to these registers are Caching Inhibited and Guarded, the register bits are reset in program order.

IOC_IO_ExcpMask[1:2] is ANDed with associated fault conditions to set `IIC_ISR[61]`, but these mask bits do not gate the setting of the I/O Exception Status bits. As a result, if software were to set one mask bit to 1 and the other mask bit to 0, the MMIO register IOC_IO_ExcpStat might be set by the masked fault. In this case, as expected, software is not alerted to the occurrence. IOC_IO_ExcpStat captures and retains the exception status information for this masked fault. Typically, software uses the same value for both mask bits. When only one mask bit in IOC_IO_ExcpMask is set to '1', it is recommended that software frequently poll and reset IOC_IO_ExcpStat to avoid losing the exception status information for the enabled fault.

For additional details about I/O and I/O-address translation, see *Section 7 I/O Architecture* on page 161.

#### 9.6.4.5 *EIB Possible Livelock Detection*

The EIB possible livelock detection interrupt detects a condition that might cause a livelock. It is a useful tool for programmers to use to detect whether they have created an unexpected bottleneck in their code during development. During normal running, this interrupt can hinder performance and should be disabled or masked. The EIB possible livelock detection interrupt signals that a possible livelock condition has been detected on the EIB due to over-utilization of the previous adjacent address match (PAAM) wait buffer. This can occur due to bad or malicious software or to frequent accesses by multiple units to common locations. Global commands that have a PAAM collision with an outstanding command are placed in the PAAM wait buffer. When this outstanding command gets its combined response, one command from the wait buffer that had a PAAM collision gets reflected on the EIB.

After the PAAM wait buffer completely fills, various things can happen, depending on how the MMIO register bits `EIB_AC0_CTL[13:15]`, the Livelock Avoidance Operating Mode field, are configured. When the wait buffer is full, the default mode ('110') is to allow the EIB command pipe to continue issuing EIB commands, but to force all commands that get a PAAM hit to receive a combined snoop response or retry.

When a wait buffer entry becomes available, a command that gets a PAAM hit is again placed in the wait buffer. There is logic that counts commands that exit the wait buffer and re-enter it due to yet another PAAM hit. This happens when multiple commands get a PAAM hit against the same cache line address within the same PAAM window. This is called a loopback. `EIB_AC0_CTL[4:5]`, the AC0 Wait Buffer Flush Delay field, defines an encoded threshold for these loopbacks. For example, when the default threshold of 64 loopbacks is hit, the command pipe is frozen until all the wait buffer commands are reflected and the wait buffer is empty. Then, normal operation resumes.

`EIB_AC0_CTL[10:11]` specifies threshold criteria for the wait buffer's becoming full again since it was last full. If the wait buffer completely fills again and the number of commands with the address modifier M bit = '1' that have occurred is less than the threshold, `EIB_Int[0]` is set to '1', provided that `EIB_AC0_CTL[13:16]` specifies that this interrupt is not disabled. `EIB_Int[0]` is not set to '1' if `EIB_AC0_CTL[13:15]` = '011', '100', '101', or '111'. If `EIB_Int[0]` changes from '0' to '1' and `EIB_AC0_CTL[16]` is set to '0', then `IIC_ISR[59]` is set to '1'.

When the EIB possible livelock detection interrupt occurs, privileged software should determine the cause of the interrupt and take appropriate action. Also, if `EIB_AC0_CTL[13:15]` is set to '010', privileged software must reset the MMIO register `EIB_Int[0]` for hardware to resume its use of the PAAM wait buffer.

Using the PAAM wait buffer will probably result in better performance for most applications. Overuse of the PAAM wait buffer is typically caused by poorly written or malicious software that creates a situation in which multiple devices simultaneously try to access the same 128-byte cache line.

#### 9.6.4.6 *Token Manager*

The MIC and IOC provide feedback to the token manager on their command queues. The MIC and IOC have registers to hold programmable thresholds. The MIC and IOC provide the token manager information about the command queue levels relative to these programmable thresholds. The IOC provides this information for each IOIF. These queue levels are defined as follows, relative to the number of entries in the command queue: If feedback for a managed resource reaches Level 3, the corresponding Token Management (TKM) Interrupt Status bit is set. Whenever the MMIO register TKM Interrupt Status (TKM_ISR) has a nonzero bit and the corresponding enable bit in the MMIO register TKM_CR[59:63] enables the interrupt, an exception signal from the TKM to the IIC is asserted. When this signal is asserted, `IIC_ISR[TMI]` is set to '1'.

#### 9.6.4.7 *Performance Monitor*

The performance monitor (PM) facility (an entirely different facility from the performance monitor facility described in the *PowerPC Architecture*) allows acquisition of performance information from multiple logic islands. It can be used to acquire a number of programmable counts. The PM facility provides four 32-bit event counters, which can alternatively be configured as eight 16-bit event counters. Up to 8 events from up to 8 different logic islands can be counted simultaneously. A 32-bit interval timer is also provided.

A performance monitor exception can be caused by the trace buffer being full, the interval timer overflow, or an overflow of any of the eight event counters. The Performance Monitor Status/Interrupt Mask Register (pm_status) is a dual function register. Writes to this address store data in the PM Interrupt Mask Register. Reads from this address return PM Status Register data and reset all PM Status Register bits. Software can read the PM Status Register to determine which exception occurred since the last read. The occurrence of a subsequent exception causes the corresponding bit of the PM Status Register to be set to '1'. The PM Interrupt Mask Register provides the ability to individually mask any one or more of these 10 conditions. If the logical AND of the PM Status Register and the PM Interrupt Mask Register is a nonzero value, and PM Control bit [0] is '1', the PM interrupt signal to the IIC is asserted; otherwise, the PM interrupt signal to the IIC is deasserted. When the PM interrupt signal is asserted, `IIC_ISR[PMI]` is set to '1'.

## 9.7 Mediated External Interrupts

### 9.7.1 Mediated External Interrupt Architecture

The PPE implements a mediated external interrupt extension to the *PowerPC Architecture* for external interrupts. On a shared processor (that is, a processor on which virtual partitions are dispatched), these external interrupt enhancements can reduce interrupt latency. The new kind of external exception is called a *mediated external exception*. An external interrupt that is caused by a mediated external exception is called a *mediated external interrupt*. This section defines the extension; page 279 provides implementation details.

Bit 52 of the Logical Partition Control Register (LPCR) is defined as the "Mediated External Exception Request" [`MER`] bit. The value '1' means that a mediated external exception is requested. The value '0' means that a mediated external exception is not requested.

On a shared processor, the hypervisor sets bit zero of the logical partitioning environment selector to '0' (LPCR[LPES] bit 0 = '0') to cause external interrupts to go to the hypervisor. In the current architecture, external interrupts are disabled if MSR[EE] = '0', and MSR[EE] can be altered by the operating system. Thus, by running with MSR[EE] = '0', an operating system can delay the presentation of external interrupts to the hypervisor for nontrivial periods of time.

On a shared processor with LPCR[LPES] bit 0 set to '0', when redispatching a partition if an external interrupt has occurred for the partition and has not yet been presented to the partition (by passing control to the operating system's external interrupt handler as described later), the hypervisor proceeds as follows:

- If the external interrupt was direct and the partition has MSR[EE] = '1',
  - Set registers (MSR and SRR0/1, and the external interrupt hardware registers as appropriate) to emulate the external interrupt.
  - Restore the partition's LPCR[MER].
  - Return to the operating system's external interrupt handler.
- If the external interrupt was direct but the partition has MSR[EE] = '0',
  - Set LPCR[MER] to '1'.
  - Return to the partition at the instruction at which it was interrupted.
- If the external interrupt was mediated and the partition now has MSR[EE] = '1'),
  - Set registers (MSR and SRR0/1, and external interrupt hardware registers as appropriate) to emulate the original direct external interrupt.
  - Restore the partition's LPCR[MER].
  - Return to the operating system's external interrupt handler.
  - When the operating system interrupt handler calls the hypervisor to obtain the interrupting condition/status, set LPCR[MER] to '0' if all external interrupts (for the partition) now have been presented to the partition. Otherwise, restore the partition's LPCR[MER]. Return from the hypervisor call with the interrupting condition/status.

In all three cases, the partition is redispatched with MSR[EE] = '0'. When the partition is to be redispatched at the operating system's external interrupt handler (as in the first and third cases), the hypervisor sets the MSR and SRR0/1 as if the original direct external interrupt occurred when LPCR[LPES] bit 0 was set to '1' and the partition was executing. In particular, no indication is provided to the operating system (for example, in an SRR1 bit) regarding whether the external interrupt that is now being presented to the partition was direct (first case) or mediated (third case).

The MER bit has the same relationship to the existence of a mediated external exception as bit zero of the decrementer (DEC[0]) or hypervisor decrementer (HDEC[0]) has to the existence of a decrementer or hypervisor decrementer exception—a value of '1' indicates an exception. (See *PowerPC Architecture, Book III* for more information.) The exception effects of LPCR[MER] are considered consistent with the contents of LPCR[MER] if one of the following statements is true.

- LPCR[MER] = '1' and a mediated external exception exists.
- LPCR[MER] = '0' and a mediated external exception does not exist.

A context synchronizing instruction or event that is executed or occurs when LPCR[MER] equals '0' ensures that the exception effects of LPCR[MER] are consistent with the contents of LPCR[MER]. Otherwise, when an instruction changes the contents of LPCR[MER], the exception effects of LPCR[MER] become consistent with the new contents of LPCR[MER] reasonably soon after the change.

The following changes are made in the description of the external interrupt.

- Direct external interrupts are enabled if the value of this expression is '1':

  ```
  MSR[EE] | (^(LPCR[LPES] bit 0) & (^(MSR[HV]) | MSR[PR]))
  ```

  In particular, if LPCR[LPES] bit 0 = '0' (directing external interrupts to the hypervisor), direct external interrupts are enabled if the processor is not in hypervisor state.

  ***Note:*** *Because the value of MSR[EE] is always '1' when the processor is in problem state, the simpler expression: MSR[EE] | ^(LPCR[LPES] bit 0 | MSR[HV]) is equivalent to the preceding expression.*

  Mediated external interrupts are enabled if the value of this expression is '1':

  ```
  MSR[EE] & (^(MSR[HV]) | MSR[PR])
  ```

  In particular, mediated external interrupts are disabled if the processor is in hypervisor state.

- There is no relative priority between direct and mediated external exceptions. If an external interrupt occurs when both kinds of external exceptions exist and are enabled, the exception that actually caused the interrupt can be either.
- When an external interrupt occurs, the state is saved in HSRR[0 and 1] if LPCR[LPES] bit 0 = '0', and in SRR[0 or 1] if LPCR[LPES] bit 0 = '1'. The state that is saved is independent of the contents of LPCR[LPES] bit 0 except as described in the next bullet (that is, the setting described for SRR0/1 in the current architecture continues to apply to SRR0/1 if LPCR[LPES] bit 0 = '1', and applies instead to HSRR[0 or 1]if LPCR[LPES] bit 0 = '0' except as modified as described in the next bullet).

  Similarly, the hypervisor external interrupt handler uses the Hypervisor Software-Use Special-Purpose Registers (HSPRGs) as scratch registers if LPCR[LPES] bit 0 = '0', and uses the Software-Use Special-Purpose Registers (SPRGs) if LPCR[LPES] bit 0 = '1'.
- If LPCR[LPES] bit 0 = '0', when an external interrupt occurs HSRR1[42] is set to '1' for a mediated external interrupt; otherwise it set to '0'. (If LPCR[LPES] bit 0 = '1', when an external interrupt occurs SRR1[42] is set to '0' as in the current architecture.)
- The hypervisor must ensure that mediated external interrupts do not occur when LPCR[LPES] bit 0 = '1'. The hypervisor can accomplish this by setting LPCR[MER] to '0' whenever it sets LPCR[LPES] bit 0 to '1'. (Mediated external interrupts are disabled when the processor is in hypervisor state, and leaving hypervisor state is necessarily accomplished by means of a context synchronizing instruction, which, if LPCR[MER] = '0', ensures that no mediated external exception exists.) If the hypervisor violates this requirement the results are undefined.

### 9.7.2 Mediated External Interrupt Implementation

The mediated external interrupt extension allows external interrupts to be serviced by hypervisor software even when external interrupts are disabled (MSR[EE] = '0'). This feature reduces the external interrupt latency for applications running in a partitioned system.

This extension defines two conditions that cause an external interrupt—the assertion of an External Interrupt signal or the hypervisor setting a mediated external interrupt request (by setting LPCR[MER] = '1'). An external interrupt caused by the assertion of an External Interrupt signal is called a direct external interrupt. An external interrupt caused by setting LPCR[MER] = '1' is called a mediated external interrupt. Each of these conditions can be masked by the following equations:

- A mediated external interrupt is taken if the following expression is equal to '1':

  ```
  MSR[EE] & (NOT(MSR[HV]) | MSR[PR]) & HID0[extr_hsrr]
  ```

- A direct external interrupt is taken if the following expression is equal to '1':

  ```
  MSR[EE] | ((NOT(LPCR[LPES] bit 0) & (NOT(MSR[HV]) | MSR[PR])) & HID0[extr_hsrr])
  ```

For direct external interrupts, the External Interrupt signals must remain asserted until the interrupting condition has been reset by the interrupt handler. Deasserting an External Interrupt signal before the PPE interrupt handler resets the interrupt condition can lead to undefined results. After the PPE accepts the interrupt by reading and resetting the interrupting condition or status, the External Interrupt signal can be deasserted.

This "level" type of direct external interrupt behavior, where the External Interrupt signal must remain asserted until the software has reset the interrupt condition, should also be applied to Mediated External interrupts. Deasserting a Mediated External Interrupt signal (by setting LPCR[MER] = '0') before the PPE interrupt handler resets the interrupting condition can lead to undefined results. After the PPE accepts the interrupt by reading and resetting the interrupting condition or status, the Mediated External Interrupt signal can be deasserted by software setting LPCR[MER] = '0'.

The PPE treats the external interrupt conditions as a context-synchronizing operation as defined in the PowerPC Architecture. When the PPE takes an external interrupt, the register state is altered as defined in *Table 9-27*, and instruction fetch and execution resume at the effective address specified in *Table 9-3 Interrupt Vector and Exception Conditions* on page 247.

*Table 9-27. Registers Altered by an External Interrupt* (Sheet 1 of 2)

| Register | Bits | Setting Description |
|---|---|---|
| SRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if exception conditions were not present.<br>SRR0 is updated if LPCR[LPES] bit 0 = '1' or HID0[extr_hsrr] = '0'. |
| HSRR0 | 0:63 | Set to the effective address of the instruction that the processor would have attempted to execute next if exception conditions were not present.<br>HSRR0 is updated if LPCR[LPES] bit 0 = '0' and HID0[extr_hsrr] = '1'. |

*Table 9-27. Registers Altered by an External Interrupt* (Sheet 2 of 2)

| Register | Bits | Setting Description |
|---|---|---|
| SRR1 | 0:63 | If LPCR[LPES] bit 0 = '1' or HID0[extr_hsrr] = '0', SRR1 is set as defined in the PowerPC Architecture.<br>Otherwise, SRR1 is not altered. |
| HSRR1 | 0:63 | If LPCR[LPES] bit 0 = '0' and HID0[extr_hsrr] = '1', HSRR1[42] is set to '1' if the external interrupt is caused by LPCR[MER].<br>Otherwise, HSRR1[42] is set to '0'.<br>HSRR1[33] is set to '1'.<br>**Note:** Setting HSRR1[33] to '1' for this interrupt is a deviation from the PowerPC Architecture for this implementation.<br>All other HSRR1 bits are set as defined in the PowerPC Architecture for SRR1. |
| MSR | 0:63 | See *Table 9-40 Machine State Register Bit Settings Due to Interrupts* on page 294. |

## 9.8 SPU and MFC Interrupts Routed to the PPE

This section describes the External interrupts (*Section 9.6.3 SPU and MFC Interrupts* on page 271) caused by SPU instructions and MFC commands and routed by each SPE, to either the PPE or an external interrupt controller. Interrupts caused by SPU events are handled locally by each SPU and are described in *Section 18.3 SPU Interrupt Facility* on page 476.

### 9.8.1 Interrupt Types and Classes

*Table 9-28* summarizes the SPU and MFC interrupt types and classes. All SPU and MFC interrupts share the same PowerPC External interrupt vector, defined in *Section 9.5.7 External Interrupt (x'00..00000500')* on page 254.

*Table 9-28. SPU and MFC External Interrupt Definitions*

| Class[1] | Interrupt Type | Description | Enabling and Status Bit[2] |
|---|---|---|---|
| 0 | SPU Error (SPU_Error) | The SPU has encountered an error condition. These conditions include:<br>• Invalid SPU instruction<br>• Illegal channel instruction<br>• Uncorrectable error-correcting code (ECC) error | SE |
| | Invalid DMA Command (Invalid_DMACmd) | Software attempted to execute an MFC command with an invalid opcode, an MFC command not supported for the specific queue (for example, MFC atomic commands in the MFC proxy command queue), or an invalid form of an MFC command. | C |
| | DMA Alignment Error (DMA_Alignment) | Software attempted to execute a DMA command that is not aligned to a 128-byte boundary | A |
| | MFC Fault Isolation Register Interrupt (MFC_FIR) | MFC Fault Isolation Register (FIR) interrupt (see *Section 9.8.3.2* on page 288). | MF[3] |
| 1 | MFC Data Segment Error (Data_Segment_Interrupt) | A DMA effective address cannot be translated to a virtual address. (segment fault) | SF |
| | MFC Data Storage Error (Data_Storage_Interrupt) | A DMA effective address cannot be translated to a real address. (mapping fault) | MF[4] |
| | MFC Local Storage Address Compare Suspend on **get** (Local-Storage Addr Comp Susp on get) | A DMA local storage (LS) address-compare stop from an LS write has occurred. | LG |
| | MFC Local Storage Address Compare Suspend on **put** (Local-Storage Addr Comp Susp on put) | A DMA LS address-compare stop from an LS read has occurred. | LP |
| 2 | Mailbox | An SPU writes to an SPU Write Outbound Interrupt Mailbox channel (SPU_WrOutIntrMbox). | M |
| | SPU Stop-and-Signal Instruction Trap (SPU_Pgm_Stop) | An SPU executes a stop-and-signal (**stop**) instruction. | S |
| | SPU Halt Instruction Trap (SPU_Trapped) | An SPU executes a halt conditional instruction, and the condition is met. | H |
| | Tag-Group Completion (Tag_group_completion) | A DMA command for a tag group has completed.<br>The interrupt generation is dependent on the Proxy Tag-Group Query Mask Register (Prxy_QueryMask) and theProxy Tag-Group Query Type Register (Prxy_QueryType). | T |
| | SPU Inbound Mailbox Threshold (SPU_mailbox_threshold) | The number of valid entries in the SPU inbound mailbox queue has dropped from 1 to 0. See *Section 17* on page 447 for details. | B |

1. The *Cell Broadband Engine Architecture* also names the classes as 0 = Error, 1 = Translation, 2 = Application. In the CBEA processors, however, the MFC_FIR interrupt is moved from CBEA class 1 to CBEA processor class 0. There are also differences in interrupt causes between the CBEA and CBEA processor versions.
2. The bit field in the INT_Mask_class and INT_Stat_class registers (*Table 9-29* on page 282).
3. This is the MF bit in the INT_Stat_class0 and INT_Mask_class0 registers.
4. This is the MF bit in the INT_Stat_class1 and INT_Mask_class1 registers.

### 9.8.2 Interrupt Registers

*Table 9-29* on page 282 summarizes the interrupt-related registers in each SPE. The interrupt registers allow privileged software on the PPE to select which MFC and SPU exceptions are allowed to generate an external interrupt to the PPE. There are three interrupt mask and interrupt status registers, one for each class of interrupt (details of class membership are listed in *Table 9-28* on page 281). There are also fault-isolation registers and other registers that report details on specific types of interrupts.

For details about these registers, see the *Cell Broadband Engine Registers* document and the *Cell Broadband Engine Architecture* document.

*Table 9-29. SPE Interrupt Register Summary (Sheet 1 of 3)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| **Interrupt Registers (*Cell Broadband Engine Architecture Registers)*** | | | |
| *Class 0 Interrupt Mask Register* | `INT_Mask_class0` | Enables the following interrupts:<br>• MFC_FIR (MF)<br>• SPU error (SE)<br>• Invalid DMA command (C)<br>• MFC DMA alignment (A) | R/W |
| *Class 1 Interrupt Mask Register* | `INT_Mask_class1` | Enables the following interrupts:<br>• MFC local-storage address compare suspend on **put** (LP)<br>• MFC local storage address compare suspend on **get** (LG)<br>• MFC data-storage interrupt—mapping fault (MF)<br>• MFC data segment interrupt—segment fault (SF) | R/W |
| *Class 2 Interrupt Mask Register* | `INT_Mask_class2` | Enables the following interrupts:<br>• SPU mailbox threshold (B)<br>• Tag-group completion (T)<br>• SPU Halt instruction trap (H)<br>• SPU Stop-and-Signal instruction trap (S)<br>• Mailbox (M) | R/W |
| *Class 0 Interrupt Status Register* | `INT_Stat_class0` | Records status of the following interrupts:<br>• MFC_FIR (MF)<br>• SPU error (SE)<br>• Invalid DMA command (C)<br>• MFC DMA alignment (A) | R/W |
| *Class 1 Interrupt Status Register* | `INT_Stat_class1` | Records status of the following interrupts:<br>• MFC local-storage address compare suspend on **put** (LP)<br>• MFC local storage address compare suspend on **get** (LG)<br>• MFC data-storage interrupt—mapping fault (MF)<br>• MFC data segment interrupt—segment fault (SF) | R/W |

*Table 9-29. SPE Interrupt Register Summary (Sheet 2 of 3)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| *Class 2 Interrupt Status Register* | INT_Stat_class2 | Records status of the following interrupts:<br>• SPU Mailbox threshold (B)<br>• Tag-group completion (T)<br>• SPU Halt instruction trap (H)<br>• SPU Stop-and-signal instruction trap (S)<br>• Mailbox (M) | R/W |
| *Interrupt Routing Register* | INT_Route | Configures the priority and destination of class 0, 1, and 2 interrupts. Interrupts can be routed to either PPE thread or either IOIF interface. | R/W |
| **Fault Isolation Registers** *(Implementation-specific Registers)* | | | |
| *MFC Fault Isolation* | MFC_FIR | Records, sets, resets, and masks MFC faults. Enables checkstop interrupt for faults. | R/W |
| *MFC Fault Isolation Set* | MFC_FIR_Set | | R/W |
| *MFC Fault Isolation Reset* | MFC_FIR_Reset | | R/W |
| *MFC Fault Isolation Error Mask* | MFC_FIR_Err | | R/W |
| *MFC Fault Isolation Error Set Mask* | MFC_FIR_Err_Set | | R/W |
| *MFC Fault Isolation Error Reset Mask* | MFC_FIR_Err_Reset | | R/W |
| *MFC Checkstop Enable* | MFC_FIR_ChkStpEnbl | | R/W |
| **Miscellaneous Registers** *(Implementation-specific Registers)* | | | |
| *MFC SBI Data Error Address* | MFC_SBI_Derr_Addr | Records the destination address (as a subordinate device) when the SPE receives data that has an error and the error is recorded in MFC_FIR register | R |
| *MFC Command Queue Error ID* | MFC_CMDQ_Err_ID | Records the MFC command queue entry index number that caused a DMA bus transaction request error | R |
| **MFC Status and Control Registers** *(**Cell Broadband Engine Architecture Registers**)* | | | |
| MFC Address Compare Control | MFC_ACCR | Allows the detection of DMA access to a virtual page with the address compare (AC) bit set in the page-table entry (PTE) and a range within the LS. | R/W |
| *MFC Data-Storage Interrupt Status* | MFC_DSISR | Records status relating to data-storage interrupts (DSIs) generated by the synergistic memory management unit (SMM). | R/W |
| *MFC Data Address* | MFC_DAR | Records the 64-bit EA from a DMA command. | R/W |
| **MFC Command Data-Storage Interrupt Registers** *(Implementation-specific Registers)* | | | |
| *MFC Data-Storage Interrupt Pointer* | MFC_DSIPR | Records the index (pointer) to the command that has an MFC data-storage interrupt (DSI) or an MFC data-segment interrupt. The cause of an MFC data-storage interrupt is identified in the MFC_DSISR register. | R |
| *MFC Local Storage Address Compare* | MFC_LSACR | Records the LS address and LS-address mask to be used in the LS address compare operation selected by the MFC_ACCR register | R/W |
| *MFC Local Storage Compare Results* | MFC_LSCRR | Records the LS address that triggered the compare, as well as the MFC command queue index of the DMA command that triggered the compare stop. | R |

*Table 9-29. SPE Interrupt Register Summary (Sheet 3 of 3)*

| Register Name | Mnemonic | Description | Read/ Write |
|---|---|---|---|
| *MFC Transfer Class ID* | MFC_TClassID | Enables transfer class ID and specifies issue quota and slot alternation. | R/W |
| **MFC Command Error Register** ***(Cell Broadband Engine Architecture Registers)*** | | | |
| *MFC Command Error* | MFC_CER | Records the MFC command-queue-entry index of the command that generated the invalid DMA-command interrupt or the DMA-alignment interrupt. | R |
| **SPU Error Mask Registers** *(Implementation-specific Registers)* | | | |
| *SPU Error Mask* | SPU_ERR_Mask | Enables SPU invalid-instruction interrupt. (Illegal-operation detection and SPU halting are always enabled.) | R/W |
| **MFC Control Register** ***(Cell Broadband Engine Architecture Registers)*** | | | |
| *MFC Control* | MFC_CNTL | Controls and reports decrementer status, and MFC command and command-queue status, | R/W |

### 9.8.2.1 *Interrupt Mask Registers*

*Table 9-29* lists three Interrupt Mask Registers (INT_Mask_class0, INT_Mask_class1, and INT_Mask_class2), one for each class of interrupts. Class 0 interrupts can be masked by means of the INT_Mask_class0 register, as shown in *Table 9-30*.

*Table 9-30. INT_Mask_class0 Bit Fields*

| Bits | Names | Descriptions |
|---|---|---|
| 31 | MF[1] | Enable for MFC_FIR interrupt. |
| 61 | SE | Enable for SPU error Interrupt. |
| 62 | C | Enable for invalid DMA command interrupt. |
| 63 | A | Enable for MFC DMA alignment interrupt. |

1. This is the MF bit in the INT_Stat_class0 and INT_Mask_class0 registers, which is different than the MF bit in the INT_Stat_class1 and INT_Mask_class1 registers.

The INT_Mask_class0[SE] bit, shown in *Table 9-30*, can be individually masked by means of SPU_ERR_Mask[I] bit, as shown in *Table 9-31*.

*Table 9-31. SPU_ERR_Mask Bit Fields*

| Bits | Name | Description |
|---|---|---|
| 63 | I | Invalid instruction interrupt enable:<br>0 Invalid instruction interrupt generation disabled.<br>1 Invalid instruction interrupt generation enabled. Generates a class 0 SPU error interrupt in INT_Stat_class0[61].<br>Illegal operation detection and SPU halting are always enabled. |

Class 1 interrupts can be masked by means of the INT_Mask_class1 register, as shown in *Table 9-32* on page 285.

*Table 9-32. INT_Mask_class1 Bit Fields*

| Bits | Names | Descriptions |
|---|---|---|
| 60 | LP | Enable for MFC local-storage compare suspend on put interrupt. |
| 61 | LG | Enable for MFC local-storage compare suspend on get interrupt. |
| 62 | MF[1] | Enable for MFC data-storage interrupt (mapping fault). |
| 63 | SF | Enable for MFC data-segment interrupt (segment fault). |

1. This is the MF bit in the `INT_Stat_class1` and `INT_Mask_class1` registers, which is different than the MF bit in the `INT_Stat_class0` and `INT_Mask_class0` registers.

Class 2 interrupts can be masked by means of the `INT_Mask_class2` register, as shown in *Table 9-33*.

*Table 9-33. INT_Mask_class2 Bit Fields*

| Bits | Names | Descriptions |
|---|---|---|
| 59 | B | Enable for SPU mailbox threshold interrupt. |
| 60 | T | Enable for DMA tag group completion interrupt. |
| 61 | H | Enable for SPU halt instruction trap or single instruction step complete. |
| 62 | S | Enable for SPU stop-and-signal instruction trap. |
| 63 | M | Enable for mailbox interrupt. |

#### 9.8.2.2 *Interrupt Status Registers*

Each SPE has three Interrupt Status Registers (`INT_Stat_class0`, `INT_Stat_class1`, and `INT_Stat_class2`), which represent the status of the interrupts in the same bit position as the corresponding mask registers. A status-register bit that is set to '1' represents an interrupt that is pending.

#### 9.8.2.3 *Interrupt Routing Register*

SPE interrupts can be prioritized and routed using the `INT_Route` register, as shown in *Table 9-34*.

*Table 9-34. INT_Route Bit Fields*

| Bits | Names | Descriptions |
|---|---|---|
| 0:3 | Class 0 priority | Priority for class 0 interrupts. |
| 8:15 | Class 0 destination | Destination ID for class 0 interrupts. |
| 16:19 | Class 1 priority | Priority for class 1 interrupts. |
| 24:31 | Class 1 destination | Destination ID for class 1 interrupts. |
| 32:35 | Class 2 priority | Priority for class 2 interrupts. |
| 40:47 | Class 2 destination | Destination ID for class 2 interrupts. |

Each class priority is four bits (the highest priority is 0, and the lowest priority is 15). The destination consists of two values: the 4 most-significant bits are the BIF node ID of the interrupt destination, and the 4 least-significant bits specify the destination ID of the target unit, as shown in *Table 9-35*.

*Table 9-35. CBEA Processor Unit Values*

| Value | Destination ID |
|---|---|
| x'0' | External interrupt controller attached to IOIF0, if any. |
| x'B' | External interrupt controller attached to IOIF1, if any. |
| x'E' | PPE thread 0. |
| x'F' | PPE thread 1. |

### 9.8.3 Interrupt Definitions

The following sections give details on Interrupts related to storage protection and MFC Fault-Isolation Register Interrupts. For details about other SPU and MFC Interrupts and the interrupt-generation process, see the *Cell Broadband Engine Architecture*.

#### 9.8.3.1 *Storage-Protection Errors*

The SPE's synergistic memory management unit (SMM) checks for storage protection errors during both virtual-mode and real-mode address translation, although real-mode checking is very limited. The DMA controller (DMAC) also checks for LS-address compares. *Table 9-36* summarizes how the contents of the `INT_Stat_class1` and `MFC_FIR` registers are affected by the results of these checks.

*Table 9-36. Exceptions Generated by Translation, MMIO, and Address Matches* (Sheet 1 of 2)

| Exception Name | INT_Stat_class1 Register | | | MFC_DSISR Register | DMAC Suspend | Data Error (DERR) | MFC_FIR Register |
|---|---|---|---|---|---|---|---|
| | LP or LG Bits[60:61] | MF (DSI) Bit[62] | SF Bit[63] | | | | |
| SLB Segment Fault | | | X | | | | |
| Atomic access to cache-inhibited page when translation is turned on | | X | | A bit[37] | X | | |
| Atomic access to cache-inhibited page when translation is turned off (real mode) | | X | | A bit[37] | X | | |
| Page protection violation (PP) | | X | | P bit[36] | | | |
| LS address compare (reported by DMAC and not reported as SMM exception) | X | | | | X | | |
| Data address compare (DAC) | | X | | C bit[41] | X | | |

*Table 9-36. Exceptions Generated by Translation, MMIO, and Address Matches* (Sheet 2 of 2)

| Exception Name | INT_Stat_class1 Register | | | MFC_DSISR Register | DMAC Suspend | Data Error (DERR) | MFC_FIR Register |
|---|---|---|---|---|---|---|---|
| | LP or LG Bits[60:61] | MF (DSI) Bit[62] | SF Bit[63] | | | | |
| TLB page fault (hardware or software) | | X | | M bit[33] | | | |
| Parity error detected during SLB/TLB translation operation | | | | | X | | SLB or TLB bits [43:44} |
| Parity error detected during MMIO read to SLB or TLB array (only reported through DERR bit attached to the MMIO data returned to the requester) | | | | | | X | |

For all translation exceptions reported by the SMM, except for parity errors detected on segment lookaside buffer (SLB) or translation lookaside buffer (TLB) MMIO reads, the 64-bit EA is saved to the MFC_DAR register. If it is a DSI exception, information about which kind of DSI error and other related information is written into the MFC_DSISR register, and the MFC command-queue index is saved to the MFC Data Storage Interrupt Pointer Register (MFC_DSIPR). The MFC_DAR, MFC_DSIPR, and MFC_DSISR are locked to the value of the current outstanding exception or miss.

The SMM only supports one outstanding translation miss or exception, excluding MMIO read parity errors. The SMM has the ability to serve hits under one outstanding translation miss or exception. After the first outstanding miss or exception is detected by the SMM, any subsequent exception or miss of a DMA translation request is reported back to the DMAC as a miss. The DMAC sets the corresponding memory management unit (MMU) dependency bit associated to the MFC command and stalls the command until the MMU dependency bit is cleared. When the outstanding miss or exception condition is cleared, all 24 (16-entry MFC SPU command queue and 8-entry MFC proxy command queue) MMU dependency bits are cleared, and the DMAC re-sends the oldest MFC command to the SMM for translation. Exceptions are cleared by setting the Restart (R) bit in the MFC_CNTL register. Setting MFC_CNTL[R] also unlocks the MFC_DAR and MFC_DSISR registers and allows the SMM to report a new exception condition. The MFC_CNTL[SC] bit should be cleared by privileged software (after clearing the MFC_CNTL[SM] bit) to unsuspend the DMAC only after all exception conditions have been cleared. The MFC_DSIPR register is unlocked by reading the register or purging the command queue.

When a 16 KB DMA transfer crosses page boundary and causes a page fault, it is not restarted from the beginning; instead, the transfer is resumed where it left off. The DMA controller unrolls DMA transfers into cache-line-sized transfers. When a fault occurs, the DMA controller only knows which transfers are still outstanding. It does not retain the original transfer request, so is can only resume the transfer where it left off.

LS address compare exceptions occur when the LS is accessed within the address range specified in the MFC_LSACR register. The DMAC writes the MFC_LSCRR register with the LS address that met the compare conditions set in the MFC_LSACR register along with the MFC command queue index and the DMA command type. This exception causes the DMAC to suspend. To clear the error, the LP or LG bit in the INT_Stat_class1 MMIO register should be cleared and the DMAC

unsuspended. See the MFC Local Storage Compare Results Register (MFC_LSCRR) and the MFC Local Storage Address Compare Register (MFC_LSACR) in the *Cell Broadband Engine Registers* document for more information.

SLB segment faults occur when there is no matching effective segment ID (ESID) in the array for the given EA from the translation request, or the valid bit is not set in the array for the matching ESID entry. When this fault occurs, the MFC_DAR register contains the 64-bit EA, as described previously. The MFC_DAR register and the SLB array can be read by means of MMIO to determine which entry to replace in the SLB. Then, the entry should be invalidated with an MMIO SLB_Invalidate_Entry or SLB_Invalidate_All write, followed by SLB_Index, SLB_ESID, and SLB_VSID writes to put the proper entry into the SLB. The INT_Stat_class1[SF] should then be cleared and the MFC_CNTL[R] bit set to resume normal SMM and DMAC operation.

All DSI exceptions are defined in the *PowerPC Operating Environment Architecture, Book III*. When these exceptions occur, the SMM writes the 64-bit EA to the MFC_DAR register and records the DSI exception type in the MFC_DSISR register, as shown in *Table 9-36* on page 286. This table also shows which DSI exceptions cause the DMAC to suspend. MFC_DSISR[S] is also set if the DMAC translation request was a **put[rlfs]**, **putll[u]c**, or **sdcrz** operation. The SMM reports the DSI to the synergistic bus interface (SBI) unit to set the INT_Stat_class1[MF] bit. The MFC_DAR and MFC_DSISR registers and the SLB and TLB arrays can be read by means of MMIO to determine the type of DSI fault and the method for fixing it. After the error condition has been fixed, the INT_Stat_class1[MF] bit should be cleared and the MFC_CNTL[R] bit set.

SLB and TLB parity errors might be recoverable by replacing the faulty entry by means of MMIO. Parity errors detected during address translation set bits in the MFC_FIR register, as shown in *Table 9-36* on page 286. See the *Cell Broadband Engine Registers* document for more information about FIR registers. Only parity errors detected on an address translation cause the MFC_DAR to be locked with the EA value captured in it. After the error has been fixed, the MFC_FIR bits should be cleared and the MFC_DSIPR register written to restart the SMM and unsuspend the DMAC. Parity errors detected during MMIO reads on the SLB or TLB cause the DERR bit to be set on the EIB when the MMIO data is returned to the requester. The SMM can support an unlimited number of MMIO read parity errors because they do not lock the SMM or DMAC interrupt registers or affect address translation.

#### 9.8.3.2 *MFC Fault-Isolation Register Interrupts*

The MFC Fault Isolation Register (MFC_FIR) Interrupt is controlled by the MF bit in the INT_Stat_class0 and INT_Mask_class0 registers. It is generated in the following cases:

- *DMA Command Alignment Error:*
  - Transfer size which is:
    - Greater than 16 KB.
    - Size neither partial nor quadword.
    - **sndsig** size not 4 bytes.
    - Partial check for effective address failed.
  - List transfer size which is:
    - Greater than 2 K list elements.
    - Bits [13:15] nonzero.

  - LS address (LSA) in which:
    - Bits [28:31] are not equal to the effective address.
    - Bits [60:63] for **put**, **get**, or **sndsig** are not all zeros for quadword transfers.
  - List address in which the lower three bits are nonzero.
- *DMA Command Error:*
  - Atomic command received while one is still pending in the queue.
  - List or atomic commands present in the MFC proxy command queue.
  - Start modifier present in the MFC SPU command queue.
  - Upper 8 bits of opcode are not all '0'.
  - Invalid opcode.
  - Transfer tag (0:10) is nonzero.

### 9.8.4 Handling SPU and MFC Interrupts

When an SPE receives an interrupt request from its SPU, DMA controller (DMAC), or synergistic memory management unit (SMM), it sends an interrupt packet to an IIC, as shown in *Figure 9-1* on page 243. The SPE sends an address-only transaction. There is no data packet associated with the interrupt packet. After receiving such a packet, the IIC asserts the external interrupt signal to the appropriate PPE thread. The first-level interrupt handler (FLIH) in that PPE thread then reads the Interrupt Pending Port Register (`IIC_IPP0` or `IIC_IPP1`) to determine the interrupt class and source. Then, the PPE reads the SPU Interrupt Status Registers for each class of SPU interrupt (`INT_Stat_class0`, `INT_Stat_class1`, and `INT_Stat_class2`) to determine the cause of the interrupt.

#### 9.8.4.1 *Interrupt Packets*

*Table 9-37* on page 290 shows the format for an address-only interrupt packet. Bits 24 through 31 of the interrupt packet's address field identify the destination of the interrupt. Interrupts should not be processed by a device unless bits 24:31 of the interrupt packet's address field match its device ID (that is, its BIF node ID concatenated with its unit ID). The address (bits 0:63) is not used for an interrupt-reissue command and is undefined in that case.

Only the target device can reject an interrupt packet. In that case, the BIF device that issued the interrupt packet must hold the interrupt pending until the target device issues an interrupt-reissue command. All BIF devices must re-send the pending interrupts if the BIF node field in the transfer tag of the interrupt-reissue command matches the BIF node of the pending interrupts. Such pending interrupts must be resent after the combined snoop response of the interrupt-reissue command. Although the resulting performance might be less than optimal, BIF devices can optionally resend the pending interrupts even if the BIF node field in the transfer tag of the interrupt-reissue command does not match the BIF node of the pending interrupts. Devices other than the target must provide a null snoop response for the interrupt transaction. All devices must provide a null or acknowledgment snoop response for the interrupt-reissue transaction. However, the master must give an acknowledgment snoop response.

*Table 9-37. Format for Interrupt Packet*

| Bits | Width | Definition |
|---|---|---|
| **Address (0:63)** | | |
| 0:21 | 22 | Reserved. |
| 22:23 | 2 | Interrupt class. |
| 24:27 | 4 | Destination BIF node ID. |
| 28:31 | 4 | Destination unit ID. |
| 32:39 | 8 | Interrupt priority. |
| 40:63 | 24 | Data. |
| **Transfer Size (0:5)** | | |
| 0 | 1 | Reserved. |
| 1:4 | 4 | Source unit ID (typically, the same as bits 4 through 7 of the transfer tag). |

As with all commands on the BIF, the device that issues an interrupt command uses its device ID as part of the transfer tag. The device is also required to put the unit ID of the device that requested the interrupt in the transfer size field of the interrupt command. In most cases, the unit ID in the size field matches the unit ID in the transfer tag. When the two values are different, the device that issues the interrupt command does so on behalf of the device identified by the unit ID in the transfer size field. This is referred to as an interrupt proxy. A device can only be an interrupt proxy for other devices within its BIF node.

The interrupt class is identified by bits 22:23 of the address. The meaning of the interrupt class is defined by the source of the interrupt and is typically provided to software as an indication of the type of interrupt. See the *Cell Broadband Engine Architecture* for the use of the interrupt class.

The interrupt priority is identified by bits 32:39 of the address. The interrupt priority is typically used by an interrupt controller to determine if this interrupt should be presented to the processor, or to the device, or rejected. See *Cell Broadband Engine Architecture* for the use of the interrupt priority.

For additional information about the I/O architecture in general, see *Section 7 I/O Architecture* on page 161.

#### 9.8.4.2 *Clearing Interrupts*

An interrupt status bit that has been set to ‘1’ in the Interrupt Status Registers (`INT_Stat_class0`, `INT_Stat_class1`, and `INT_Stat_class2`) must be cleared after the corresponding interrupt is handled. The only case in which this is not true is for class 1 interrupts in which the MFC Data-Storage Interrupt Status Register (`MFC_DSISR`) needs to be cleared to ‘0’ before any of the status register bits for class 1 are cleared. Also, the `MFC_CNTL[R]` bit needs to be set to ‘1’ to resume DMA operation after status-register bits have been cleared.

Software should not clear the source of an interrupt until the Interrupt Status Register (`INT_Stat_class0`, `INT_Stat_class1`, and `INT_Stat_class2`) bits are read and cleared. If an Interrupt Status Register is cleared on the same cycle as the interrupt, the source sends a signal to the SBI, and this interrupt source can be lost. In the case of a MFC Fault Isolation Register

(MFC_FIR) class 0 interrupt, software should clear the FIR bits first and then clear the Interrupt Status Register. If there is any FIR bit set when the Interrupt Status Register is cleared, the SBI sends a new interrupt packet.

## 9.9 Thread Targets for Interrupts

*Table 9-38* summarizes the PPE thread targets of the various interrupt types. The nonrecoverable checkstop error (which is not an interrupt) is described in *Section 9.15* on page 295.

*Table 9-38. PPE-Thread Target for Interrupts*

| Exception Type | Interrupt | Target PPE Thread |
|---|---|---|
| Instruction-Caused (Synchronous) | All instruction-caused interrupts | Thread causing interrupt |
| System-Caused (Asynchronous) | System Reset, when a thread is started or restarted | Selectable by software |
| | System Reset, at power-on reset (POR) | Thread 0 |
| | Machine Check[1] | Both threads |
| | System Error | One or both threads, depending on the error |
| | Decrementer | Thread causing interrupt |
| | Hypervisor Decrementer | Both threads |
| | Thermal Management | Selectable by software |
| | Maintenance (instruction-caused) | One or both threads, depending on the error |
| | External Interrupt (direct or mediated)[2] | Selectable by software |

1. If HID0[en_prec_mchk] = '1' ("precise"), only the thread that causes a machine check interrupt will take the interrupt, and the interrupt will appear to be instruction-caused (synchronous) from the viewpoint of that thread (see *Section 9.5.2* on page 249). If HID0[en_prec_mchk] = '0' ("imprecise"), both threads will take a machine check interrupt, and the interrupt will appear to be system-caused (asynchronous) to both threads. A suspended thread will not resume.
2. See *Section 9.5.7* on page 254 for the difference between direct and mediated external interrupts.

## 9.10 Interrupt Priorities

The following exception conditions are ordered from highest to lowest priority.

*Table 9-39. Priority Order of Interrupt Conditions (Sheet 1 of 3)*

I. System reset interrupt (highest priority exception)

II. Machine check interrupt

III. Instruction-dependent

*Table 9-39. Priority Order of Interrupt Conditions (Sheet 2 of 3)*

| | |
|---|---|
| A. Fixed-Point Loads and Stores | 1. Data storage interrupt (DSI) caused by the DABR<br>2. Alignment interrupts caused by fixed-point load and store instructions (**lwarx**, **ldarx**, **stwcx.**, **stdcx.**, **lmw**, or **stmw**).<br>3. DSI caused by a logical partitioning (LPAR) error in real mode (MSR[HV] = '0' and LPCR[HDICE] = '0')<br>4. Data segment interrupt<br>5. DSI caused by a page fault<br>6. Alignment interrupts (referencing caching-inhibited storage)<br>a. Alignment interrupts caused by fixed-point load instructions that reference caching-inhibited storage.<br>b. Alignment interrupts caused by fixed-point store instructions that reference caching-inhibited storage, if the C-bit in the page table entry is set to '1'.<br>7. DSI caused by a page-protection violation<br>8. Alignment interrupt caused by a fixed-point store instruction referencing caching-inhibited storage, if the C-bit in the page table entry is set to '0'.<br>9. DSI caused by the ACCR<br>10. Trace interrupt |
| B. Floating-Point Loads and Stores | 1. Floating-point unavailable interrupt<br>2. DSI caused by the DABR<br>3. Alignment interrupt caused by a floating-point load and store instruction that is not word aligned.<br>4. DSI caused by an LPAR error in real mode (MSR[HV] = '0' and LPCR[HDICE] = '0')<br>5. Data segment interrupt<br>6. DSI caused by a page fault<br>7. Alignment interrupts (referencing caching-inhibited storage)<br>a. Alignment interrupt caused by a double-precision floating-point load instruction that is word-aligned, but not doubleword-aligned, referencing caching-inhibited storage.<br>b. Alignment interrupt caused by a double-precision floating-point store instruction that is word-aligned, but not doubleword-aligned, referencing caching-inhibited storage and the C bit in the page table entry is set to '1'.<br>8. DSI caused by a page-protection violation<br>9. Alignment interrupt caused by a double-precision floating-point store instruction that is word-aligned, but not doubleword-aligned, referencing caching-inhibited storage, and the C bit in the page table entry is set to '0'.<br>10. DSI caused by the ACCR<br>11. Trace interrupt |
| C. Other Floating-Point Instructions | 1. Floating-point unavailable interrupt<br>2. Precise-mode floating-point enabled exceptions type of program interrupt<br>3. Trace interrupt |
| D. VXU Loads and Stores | 1. VXU unavailable interrupt<br>2. DSI caused by the DABR<br>3. DSI caused by an LPAR error in real mode (MSR[HV] = '0' and LPCR[HDICE] = '0')<br>4. Data segment interrupt<br>5. DSI caused by a page fault<br>6. Alignment interrupts (referencing caching-inhibited storage)<br>a. Alignment interrupts caused by vector scalar unit (VSU) load and store (left or right) instructions that reference caching-inhibited storage.<br>b. Alignment interrupts caused by VSU load and store (left or right) instructions that reference caching-inhibited storage, if the C-bit in the page table entry is set to '1'. |

*Table 9-39. Priority Order of Interrupt Conditions (Sheet 3 of 3)*

7. DSI caused by a page-protection violation
8. Alignment interrupt caused by a VSU load and store (left or right) instruction referencing caching-inhibited storage, if the C-bit in the page table entry is set to '0'.
9. DSI caused by the ACCR
10. Trace interrupt

E. Other VXU Instructions
1. VXU unavailable interrupt
2. Trace interrupt

F. **rfid**, **hrfid**, and **mtmsr**[**d**]
1. Precise-mode floating-point enabled exceptions type of program interrupt
2. Trace interrupt (for **mtmsr**[**d**] only)

G. Other Instructions
1. Exceptions that are mutually exclusive and the same priority:
   a. Trap type of program interrupt
   b. System call
   c. Privileged Instruction type of program interrupt
   d. Illegal Instruction type of program interrupt
2. Trace interrupt

H. Instruction segment interrupt

I. Instruction storage interrupt

IV. Thermal management interrupt

V. System error interrupt

VI. Maintenance interrupt

VII. External interrupt

A. Direct

B. Mediated

VIII. Hypervisor decrementer Interrupt

XI. Decrementer Interrupt

## 9.11 Interrupt Latencies

Latencies for taking various interrupts are variable, based on the state of the machine when conditions exist for an interrupt to be taken. Instruction dispatch is blocked for both threads when the conditions exist for a system-caused exception to be taken. Because microcoded instructions must complete, the interrupt latency depends on the current instructions being processed. The interrupt is not taken until both threads are flushed, no higher-priority exceptions exist, and the exception condition, if enabled, exists.

## 9.12 Machine State Register Settings Due to Interrupts

*Table 9-40* on page 294 shows the settings of Machine State Register (MSR) bits for Hypervisor (HV), Machine-Check Enable (ME), and Recoverable Interrupt (RI), relative to the interrupt type. *Table 9-3* on page 247 list the interrupts supported by the PPE, the effective address of the interrupt handler (interrupt vector), and summarizes the conditions that cause the interrupt.

*Table 9-40. Machine State Register Bit Settings Due to Interrupts*

| Interrupt Type | MSR Bit[1] | | |
|---|---|---|---|
| | HV | ME | RI[2] |
| System Reset | 1 | — | 0 |
| Machine Check (imprecise form) | 1 | 0 | 0[3] |
| Precise Machine Check | 1 | 0 | 0[3] |
| Data Storage | m | — | 0 |
| Data Segment | m | — | 0 |
| Instruction Storage | m | — | 0 |
| Instruction Segment | m | — | 0 |
| External | e | — | 0[4] |
| Alignment | m | — | 0 |
| Program | m | — | 0 |
| Floating-Point Unavailable | m | — | 0 |
| Decrementer | m | — | 0 |
| Hypervisor Decrementer | 1 | — | — |
| System Call | s | — | 0 |
| Trace | m | — | 0 |
| System Error | 1 | — | — |
| Maintenance | 1 | — | — |
| Thermal Management | 1 | — | — |
| VXU Unavailable | m | — | 0 |

Legend:

- 0 Bit is set to ‘0’.
- 1 Bit is set to 1.
- — Bit is not altered
- m MSR[HV] is set to ‘1’ if LPCR[LPES] bit 0 = ‘0’ and LPCR[LPES] bit 1 = ‘0’; otherwise, the state of MSR[HV] is not altered.
- e MSR[HV] is set to ‘1’ if LPCR[LPES] bit 0 = ‘0’; otherwise, the state of MSR[HV] is not altered.
- s MSR[HV] is set to ‘1’ if LEV = ‘1’ or LPCR[LPES] bit 0 = ‘0’ and LPCR[LPES] bit 1= ‘0’; otherwise, the state of MSR[HV] is not altered.

**Note:**

1. MSR[BE], MSR[EE], MSR[FE0], MSR[FE1], MSR[FP], MSR[PR], MSR[SE], MSR[DR], and MSR[IR] are cleared to ‘0’, and MSR[SF] is set to ‘1’.
2. MSR[RI] is not changed for interrupts that use HSRR0 or HSRR1.
3. MSR[RI] is cleared to ‘0’. SRR1[62] is cleared to ‘0’ for an imprecise Machine Check and is not altered for a precise Machine Check.
4. MSR[RI] is not altered if LPCR[LPES] bit 0 is set to ‘0’ and HID0[extr_hsrr] is set to ‘1’; otherwise, MSR[RI] is cleared to ‘0’.

## 9.13 Interrupts and Hypervisor

The hypervisor feature supported by the PPE is described in *Section 11* on page 331. This section and *Table 9-40* on page 294 contain details relating to the requirements for and effects of interrupts in hypervisor mode.

The main setup for interrupt handling is to place code in the interrupt vector offset. Each logical partition (LPAR) can have its own interrupt vector using the Real Mode Offset Register (RMOR). This register points to the start of the interrupt vector (defaults to 0), so that each LPAR—by means of a hypervisor (*Section 11* on page 331)—can have its own dedicated vectors. There is also a hypervisor-only interrupt vector offset register called HRMOR which points to the interrupt vector when MSR[HV] = '**1**'.

The following areas need to be initialized:

- HRMOR for setting hypervisor interrupt vector area.
- RMOR for pointing to the current supervisor's interrupt vector area (which should change for each LPAR).
- Each enable bit (and any mask bits) listed in *Table 10-9* on page 329.
- External Interrupt settings.
- Interrupt priority settings by means of IIC_CPL0 and IIC_CPL1 (internal interrupt controller current priority level for thread 0 and thread 1). Default is 0, which is the highest priority.

## 9.14 Interrupts and Multithreading

The multithreading features supported by the PPE are described in *Section 10* on page 299. For details on the entry condition, activated thread, mask condition, and suspended-thread behavior for each interrupt type, see *Section 10.8.4 Thread Targets and Behavior for Interrupts* on page 328.

## 9.15 Checkstop

A checkstop is a special nonrecoverable condition (rather than an interrupt) in which the PPE and the rest of the CBEA processor shut down. There are multiple causes for a checkstop. It can be caused, for example, by two successive machine check interrupts occurring close together, with the second occurring before the first one completes. It can also be caused by a watchdog timeout or a hardware signal from an external device.

Software cannot accept nor mask a checkstop event. After the CBEA processor hardware initiates the checkstop sequence, software has no role until System Reset occurs. See the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide* for more information about the checkstop.

## 9.16 Use of an External Interrupt Controller

The primary purpose of the IIC is to allow interrupts from SPEs and other components in the CBEA processors to be handled internally, without using a external interrupt controller. In a system with many I/O devices, however, an external interrupt controller might be desirable. In such a configuration, the IICs can function as second-level interrupt controllers, handling all interrupts internal to the CBEA processor or within a multi-CBEA-processor system, while an external system interrupt controller handles interrupts external to the CBEA processor or multi-CBEA-processor system.

When handling an interrupt in a system with an external main system interrupt controller, PPE software must first check the IICs to determine if the interrupt was sourced internally or from an external system interrupt controller.

## 9.17 Relationship Between CBEA Processor and PowerPC Interrupts

The following differences exist between CBEA processor and PowerPC Architecture interrupts:

- PowerPC Architecture interrupts not supported by the CBEA processors:
  - Imprecise-mode floating-point enabled
  - Performance monitor (However, see the unrelated CBEA processor performance monitor interrupt later in this section.)
  - Optional example extensions to the trace facility
- CBEA processor interrupts supported differently than in the PowerPC Architecture:
  - In the PowerPC Architecture, the machine check interrupt is implementation-dependent. In the CBEA processors, the interrupt is supported in two forms, *precise*[4] and *imprecise*, neither of which are defined in the PowerPC Architecture.
  - In the PowerPC Architecture, external interrupts have a single form. In the CBEA processors, the interrupts are supported in two forms, *direct* and *mediated*, neither of which are defined in the PowerPC Architecture.
- CBEA processor interrupts that are not defined in the PowerPC Architecture or the vector/SIMD multimedia extension to the PowerPC Architecture:
  - System error interrupt
  - Maintenance interrupt (instruction-caused)
  - Maintenance interrupt (system-caused)
  - Thermal management interrupt
  - External interrupts:
    - SPU instruction interrupt.
    - MFC command interrupt.
    - MIC auxiliary trace buffer full interrupt.
    - EIB possible livelock detection interrupt.

4. The precise machine check interrupt is a system-caused interrupt but appears to be an instruction-caused interrupt from the viewpoint of the PPE thread that causes it (see *Section 9.5.2* on page 249).

- Token manager interrupt.
- CBEA processor performance monitor interrupt (unrelated to PowerPC performance monitor interrupt).
- Software interrupt.

# 10. PPE Multithreading

The Cell Broadband Engine Architecture (CBEA) processor[1] elements use several implementation-specific techniques to speed program execution. In the PowerPC Processor Element (PPE), simultaneous multithreading is one of the most important techniques. To software, the PPE appears to provide two independent instruction-processing units. The threads appear to be independent because the PPE provides each thread with a copy of architectural state, such as general-purpose registers, but the threads are not completely independent because many execution resources are shared by the threads to reduce the hardware cost of multithreading.

To software, the PPE implementation of multithreading looks similar to a multiprocessor implementation, but there are several important differences. *Table 10-1* compares the PPE multithreading implementation to a conventional dual-core microprocessor.

*Table 10-1. PPE Multithreading versus Multi-Core Implementations*

| Characteristic | PPE Multithreading | Multiple Cores on a Single Chip |
|---|---|---|
| Hardware cost | Low to moderate cost; the processor replicates only architectural state. | High cost; the entire processor core and caches are replicated. |
| Performance/Cost | A typical 10% to 30% performance gain (application-dependent) for 5% hardware cost. | Up to 100% performance gain for 100% hardware cost. |
| Hardware efficiency | Execution-unit utilization increased compared to single-thread of execution. | Execution-unit utilization unchanged. |
| Software restrictions | Threads must execute in same logical partition (must share memory-management tables, and so forth). | Depending on implementation, threads and processes might be able to execute in different logical partitions. |
| Thread competition | Threads can compete for execution resources, such as cache residency and execution units. | Threads do not compete; they have private, replicated execution resources. |
| System throughput | System throughput improved at possible expense of slightly slower single-thread latency. | System throughput improved without reducing single-thread latency. |

Because most of the PPE hardware is shared by the two threads of execution, the hardware cost of the PPE multithreading implementation is dramatically lower than the cost of replicating the entire processor core. The PPE's dual-threading typically yields approximately one-fifth the performance increase of a dual-core implementation, but the PPE achieves this 10%-to-30% performance boost at only one-twentieth of the cost of a dual-core implementation.

## 10.1 Multithreading Guidelines

In the PPE multithreading implementation, the two hardware threads share execution resources. Consequently, concurrent threads that tend to use different subsets of the shared resources will best exploit the PPE's multithreading implementation. Pointer chasing and scattered memory accesses are an excellent examples of applications in which multithreading really shines, even if both threads are running the same type of memory-latency-sensitive application. For example, a speedup of 2x is entirely possible when running memory-pointer chases on both threads.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

When one thread experiences high miss rates in the L1 or L2 caches, the other execution resources in the PPE, such as the fixed-point and floating-point units, will be completely free to execute instructions from the other thread. Two examples of typical software functions that are likely to incur cache misses on memory references are the process of gathering data from sparse or scattered vectors in preparation for single instruction, multiple data (SIMD) processing and the pointer chasing performed when searching a linked list. For example, searching a linked list for an element of a particular type might be done with a loop like the following example:

```
ptr = list_head;
while (ptr != NULL)
{
        if (ptr->type == desired_type) {
            break;
        }
        ptr = ptr->next;
}
```

This code accesses list elements in link order, but the physical-memory addresses of adjacent list elements might not be adjacent. Consequently, on each iteration of the loop shown in the code sample, every execution of the statement `ptr = ptr->next` is likely to cause a cache miss the first time through the loop. If the list is large enough to overwhelm the data caches, it might cause cache misses every time a link is accessed.

Floating-point operations also often leave most of the other execution resources sitting idle. When one thread makes heavy use of the floating-point execution unit, a second, fixed-point-intensive thread will be able to execute freely and make significant forward progress when the pipeline would otherwise be idle.

Certain types of source code result in instruction sequences that are dominated by serial dependencies (each successive instruction depends on the previous one) and so have little instruction-level parallelism. Similarly, when compute-intensive applications have loops that cannot be effectively unrolled or that prevent successful software pipelining, the resulting thread of execution will not be able to exploit fully the available parallelism in the PPE execution pipeline. In these situations, when a single thread lacks enough instruction-level parallelism, a second simultaneous thread can help increase the use of execution resources and therefore system throughput.

Another specific situation that leaves execution resources idle is a thread that has a lot of branches that cannot be successfully predicted by the PPE branch-prediction hardware. A mispredicted branch leaves the execution resources idle for many cycles. A second, more-predictable thread can make good forward progress and increase pipeline use and system throughput.

Whenever the overriding goal for an application is to maximize system throughput, multithreading should be used. When the overriding goal is to minimize the latency of a particular task, the system should be programmed to disable the second thread or to set its priority low enough to reduce competition for execution resources. However, setting the second thread's priority low will not eliminate latency-inducing effects on the other thread. If the target is minimum *guaranteed* latency, the processor should ideally be in single-thread mode or should not have a second thread running at all (that is, the thread context is sleeping), because some operations on the low-priority thread—for example, a **sync** operation—can still cause the entire pipeline to stop.

PPE multithreading will not typically help code that is tuned precisely to the characteristics of the pipeline, because such a thread can probably use all or almost all of the execution resources. Even in this case, using multithreading will not necessarily compromise system throughput unless performance of the tuned thread also depends on the sizes of the L1 and L2 caches; a thread that is tuned to the cache sizes can experience thrashing if a second thread is allowed to run and pollute the cache. Cache misses are so expensive that any increase in misses will probably more than negate the benefit gained from multithreading.[2]

## 10.2 Thread Resources

For each thread, the PPE has duplicate copies of all architected state, such as general-purpose registers and thread-dependent special-purpose registers (SPRs). Exceptions to this rule are registers that deal with system-level resources, such as logical partitions and memory management. To software, the PPE appears to offer two separate virtual processors for instruction execution.

Resources that are not described in the *PowerPC Architecture*, such as the hardware execution units, are typically shared by the two threads, but nonexposed resources are duplicated in cases where the resource is small or offers a performance improvement that is critical to multithreaded applications.

The two PPE threads are independent in most aspects, but they share logical-partitioning resources. Thus, the threads always execute in the same logical-partition context and share structures such as virtual-memory-mapping tables. Threads also share most of the large arrays and queues, such as caches, that consume significant amounts of chip area.

### 10.2.1 Registers

The following architected registers are *duplicated* for multithreading and used by software running in any privilege state (including problem-state):

- General-Purpose Registers (GPRs) (32 entries per thread)
- Floating-Point Unit Registers (FPRs) (32 entries per thread)
- Vector Registers (VRs) (32 entries per thread)
- Condition Register (CR)
- Count Register (CTR)
- Link Register (LR)
- Fixed-Point Exception Register (XER)
- Floating-Point Status and Control Register (FPSCR)
- Vector Status and Control Register (VSCR)
- Decrementer (DEC)

2. Multithreading might result in fewer cache misses if both threads use the same code, as for example when using shared library routines.

In addition, see *Section 10.4 Thread Control and Status Registers* on page 306 for additional architected registers that are duplicated for multithreading but used specifically for thread control and status monitoring. All memory mapped I/O (MMIO) registers in the PPE deal with aspects of the power processor storage system (PPSS) that are shared for both threads, so all MMIO registers are shared for both threads as well.

### 10.2.2 Arrays, Queues, and Other Structures

The following arrays, queues, and structures are fully *shared* between threads running in any privilege state:

- L1 instruction cache (ICache), L1 data cache (DCache), and L2 cache
- Instruction and data effective-to-real-address translation tables (I-ERAT and D-ERAT)
- I-ERAT and D-ERAT miss queues
- Translation lookaside buffer (TLB)
- Load miss queue and store queue
- Microcode engine
- Instruction fetch control
- All execution units:
  - Branch (BRU)
  - Fixed-point integer unit (FXU)
  - Load and store unit (LSU)
  - Floating-point unit (FPU)
  - Vector media extension unit (VXU)

The following arrays and queues are *duplicated* for each thread.

- Segment lookaside buffer (SLB)
- Branch history table (BHT), with global branch history
- Instruction buffer (IBuf) queue
- Link stack queue

Duplicating the instruction buffer allows each thread to dispatch regardless of any dispatch stall in the other thread. Duplicating the SLB is convenient for the implementation because of the nature of the *PowerPC Architecture* instructions that access it and because it is a relatively small array.

The instruction-fetch control is shared by both threads because the instruction cache has only one read port and so fetching must alternate between threads every cycle. Each thread maintains its own BHT and global branch history (GBH) to allow independent and simultaneous branch prediction.

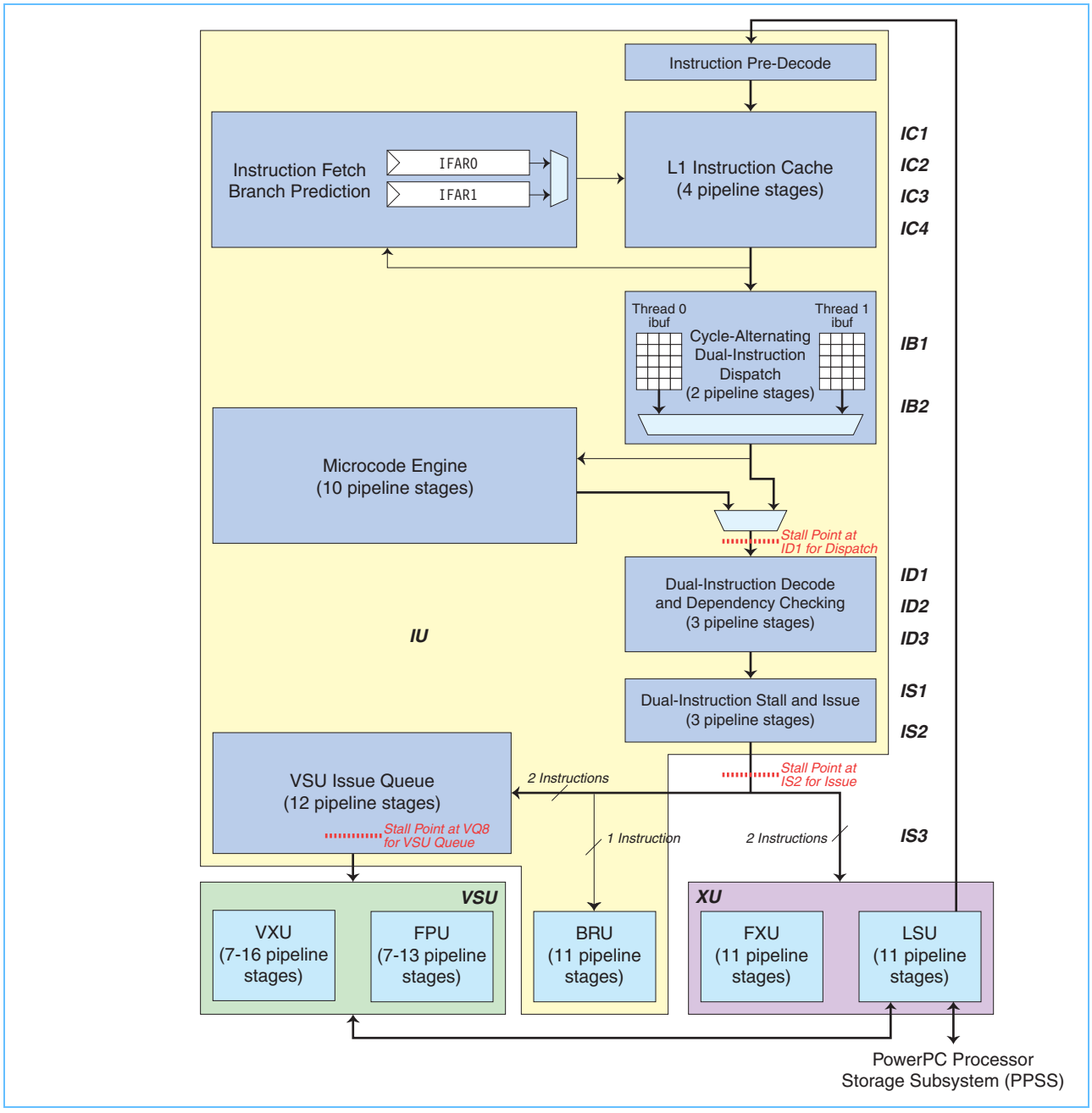
### 10.2.3 Pipeline Sharing and Support for Multithreading

*Figure 10-1* on page 304 shows a high-level, multithreading view of instruction flow and resources in the PPE pipeline. The units shown in this figure include the instruction unit (IU), execution unit (XU), fixed-point unit (FXU), load-store unit (LSU), branch unit (BRU), vector scalar unit (VSU), floating-point unit (FPU), and vector/SIMD multimedia extension unit (VXU). The pipeline stages shown for the execution units at the bottom of this figure are the total number of stages through completion; bypass paths are reached in fewer stages. For details on actual latency and throughput for each instruction, see *Table A-1 PowerPC Instructions by Execution Unit* on page 723.

Instruction-fetch hardware autonomously operates and maintains an Instruction-Fetch Address Register for each thread (`IFAR0` and `IFAR1`), which are not visible to software. In contrast to instruction dispatch, which can repeatedly favor one thread over another, instruction fetching strictly alternates on successive cycles. The purpose of instruction fetching is to keep each thread's instruction buffer full with useful instructions that are likely to be needed by the thread. A thread's Instruction Fetch Address Register (`IFAR`) is distinct from its program counter: the program counter tracks *actual* instruction-execution flow while an IFAR tracks *predicted* instruction-execution flow. Instruction fetching is pipelined; starting the fetch of a new sequential stream after a predicted-taken branch requires eight cycles.

The dispatch hardware sends instructions from the instruction buffers to the shared decode, dependency-checking, and issue pipeline. The dispatch hardware selects instructions from the buffer of a particular thread based on thread priority and other factors such as stall conditions. Each pipeline stage beyond the dispatch point contains instructions from one thread only.

*Figure 10-1. PPE Multithreading Pipeline Flow and Resources*




### 10.2.3.1 *Pipeline Stall Points*

There are three stall points in the pipeline: dispatch (ID1), issue (IS2), and VSU issue (VQ8). The dispatch stall point at ID1 (*Figure 10-1*) is separate for each thread so that one thread can dispatch even if the other is stalled.

The issue and VSU-issue stall points (see *Appendix A.5 Issue Rules* on page 760 and *Appendix A.6.2 Vector/Scalar Unit Issue Queue* on page 764) are not separate for each thread. A stall at IS2 requires both threads to stall, and a stall at VQ8 will stall both threads if and only if the stall also results in a stall at IS2. To reduce the number of VQ8 stalls that also stall at IS2, the VSU issue queue has a 2-stage buffer that can expand the pipeline length temporarily by two stages. When VQ8 stalls and the 2-stage buffer is free, the VSU issue queue can accept up to four more instructions (2 entries × 2 instructions) before it must signal IS2 to stall.

Most stalls do not cause flushes, refetches, or dispatch blocking. Almost all stalls just cause the stall to occur such that issue is stopped for a few cycles (see *Section 10.7* on page 322). However, for very long-latency conditions, such as when an instruction is issued that is dependent on an L1 cache miss or on a divide that is being computed, stall conditions might cause the following sequence of actions in the pipeline:

- Instructions younger than the stalled instruction are flushed from the pipeline.
- Instructions starting with the stalled instruction are refetched.
- The thread is stalled at the dispatch (ID1) stage until the stall condition is removed.

When an instruction is dependent on an older caching-inhibited load instruction that has not completed, stall conditions cause the following sequence of actions in the pipeline:

- The dependent instruction and all younger instructions for the same thread are flushed. The flush takes place 10 cycles after the dependent instruction is issued from the IS2 point.
- The flushed instructions are refetched.
- The thread is stalled at the ID1 stage until the caching-inhibited load is complete.

## 10.3 Thread States

The operating state of each thread has several properties that affect instruction dispatch, including:

- The thread's privilege state
- Whether the thread is suspended or enabled
- Whether or not the thread is blocked by a stall condition not related to its priority
- The thread's priority

### 10.3.1 Privilege States

A thread can run in any of three privilege states: hypervisor state, privileged state (supervisor), and problem (user) state. Hypervisor state is the most privileged, and some system operations require the initiating thread to be in hypervisor state. Hypervisor state exists to run a meta-operating system that manages logical partitions in which multiple operating system instances can run. Supervisor state is the state in which an operating system instance is intended to run when multiple instances are supported. User state (problem state) is for running application programs. *Table 10-2* on page 306 summarizes the Machine State Register (MSR) bit settings needed for each state.

*Table 10-2. Privilege States*

| MSR[HV] | MSR[PR] | Privilege State |
|---|---|---|
| 1 | 0 | Hypervisor |
| 0 | 0 | Privileged State (also called Supervisor) |
| Don't care | 1 | Problem State (also called User) |

As shown in the table, when MSR[PR] is set to '1', the user state (problem state) is active, regardless of the state of the hypervisor bit in the MSR register.

After power-on reset (POR), the Problem State Change Thread Priority Enable bit, (TSCR[UCP]) and the Privileged But Not Hypervisor State Change Thread Priority Enable bit (TSCR[PSCTP]) are cleared to '0', so thread priority can be changed only in hypervisor state. See *Section 10.4.6* on page 312.

### 10.3.2 Suspended or Enabled State

A thread can be in the suspended state if its Thread Enable bit is cleared to '0' in the Control Register (CTRL[TE0] or CTRL[TE1]). A thread can suspend only itself (a thread cannot suspend the other thread), and a thread can only write to the Control Register (CTRL) to change a thread-enable bit when the thread is in hypervisor state (MSR[PR,HV] = 0,1).

### 10.3.3 Blocked or Stalled State

In general, *blocking* occurs at the instruction-dispatch stage and stops only one of the two threads, whereas *stalling* occur at the instruction-issue stage and stop both threads. A thread can be blocked or stalled for any one of several conditions, including the following conditions:

- The thread executes one of a set of special **nop** instructions that block instruction dispatch for a specific number of cycles (see *Table 10-4* on page 314).
- The pipeline blocks the thread at instruction dispatch because the combination of thread priorities puts in force an instruction-dispatch policy that selects the other thread for dispatch (see *Section 10.5* on page 313).
- The pipeline is forced to stall instruction progress due to dependencies.
- A context-synchronizing event or instruction occurs.
- The HID0[issue_serialize] bit is set to '1'.

## 10.4 Thread Control and Status Registers

*Section 10.2.1* on page 301 lists the architected registers that are duplicated, per thread, and used by software running in any privilege state (including problem state). This section describes the primary control and status registers that affect and monitor the behavior of the multithreading facilities in the PPE. These registers allow software to enable threads, set relative thread priorities, set the duty cycle for threads of unequal priority, determine thread response to interrupts, and set timer values so that system software will be informed when a thread is starved of execution resources.

Each thread is viewed as an independent processor, complete with separate exceptions and interrupt handling. The threads can generate exceptions simultaneously, and the PPE supports concurrent handling of interrupts on both threads by duplicating some registers defined by the PowerPC Architecture.

The following registers associated with exceptions and interrupt handling are duplicated or are thread-dependent:

- Machine State Register (MSR)
- Machine Status Save/Restore Registers (SRR0 and SRR1)
- Hypervisor Machine Status Save/Restore Registers (HSRR0 and HSRR1)
- Floating-Point Status and Control Register (FPSCR)
- Data Storage Interrupt Status Register (DSISR)
- Decrementer (DEC)
- Logical Partition Control Register (LPCR)
- Data Address Register (DAR)
- Data Address Breakpoint Register (DABR and DABRX)
- Address Compare Control Register (ACCR)
- Thread Status Register Local (TSRL)
- Thread Status Register Remote (TSRR)

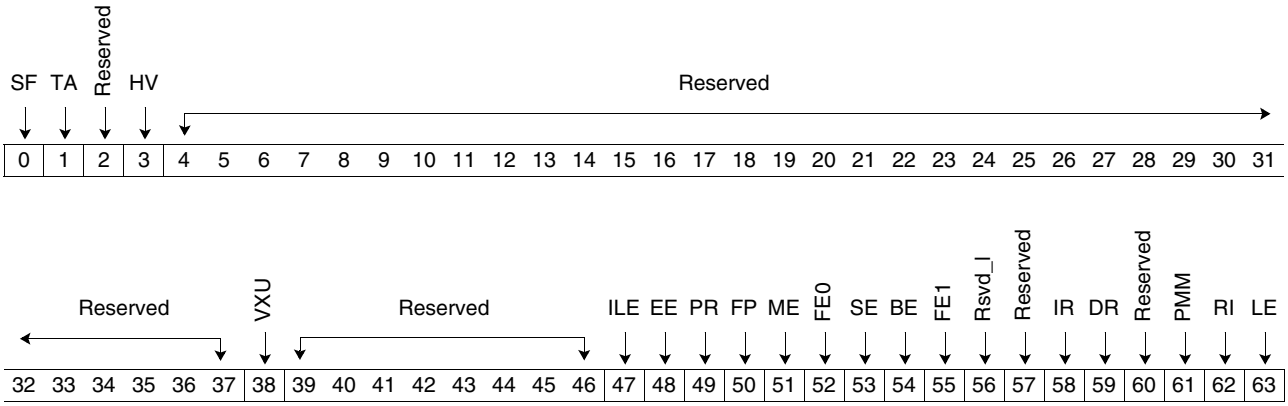
In addition, the following thread-independent registers also are associated with exceptions and interrupt handling on both threads:

- Hypervisor Decrementer (HDEC)
- Control Register (CTRL)
- Hardware Implementation Dependent Registers 0 and 1 (HID0 and HID1)
- Thread Switch Control Register (TSCR)
- Thread Switch Time-Out Register (TTR)

The following sections describe processor registers that play a central role in controlling and monitoring multithreading activity. The tables that describe register fields have been edited for brevity and clarity (for example, the reserved fields and fields not relevant to multithreading are not described). For complete register descriptions, see the *Cell Broadband Engine Registers* specification.

### 10.4.1 Machine State Register (MSR)

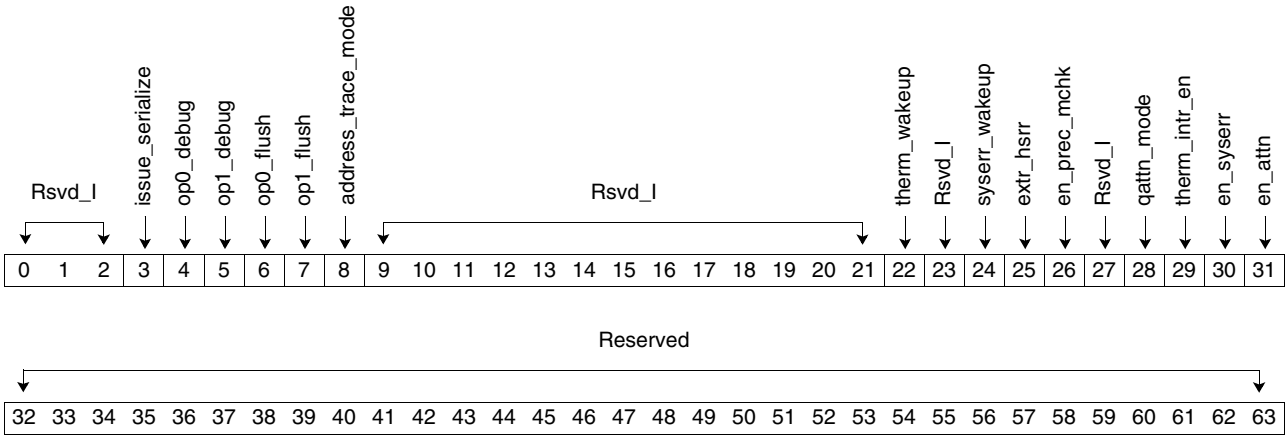
The Machine State Register (MSR) contains the hypervisor-state and problem-state (user-state) bits for setting the privilege state of a thread. See *Section 9.5.7* on page 254 for how external interrupts are masked when MSR[EE] is cleared to ‘0’. When MSR[ME] is set to ‘1’, the thread can take a machine-check interrupt. See *Section 11.2.1* on page 336 for how the MSR[HV] bit is set to enable hypervisor state.




| Bits | Field Name | Description |
|---|---|---|
| 3 | HV | Hypervisor State (see *Table 10-2* on page 306) |
| 48 | EE | External interrupt enable |
| 49 | PR | Problem state (see *Table 10-2* on page 306) |
| 51 | ME | Machine check enable<br>**Note:** This bit can only be modified by the **rfid** and **hrfid** instructions while in hypervisor mode (see *PowerPC Operating Environment Architecture, Book III*). |

### 10.4.2 Hardware Implementation Register 0 (HID0)

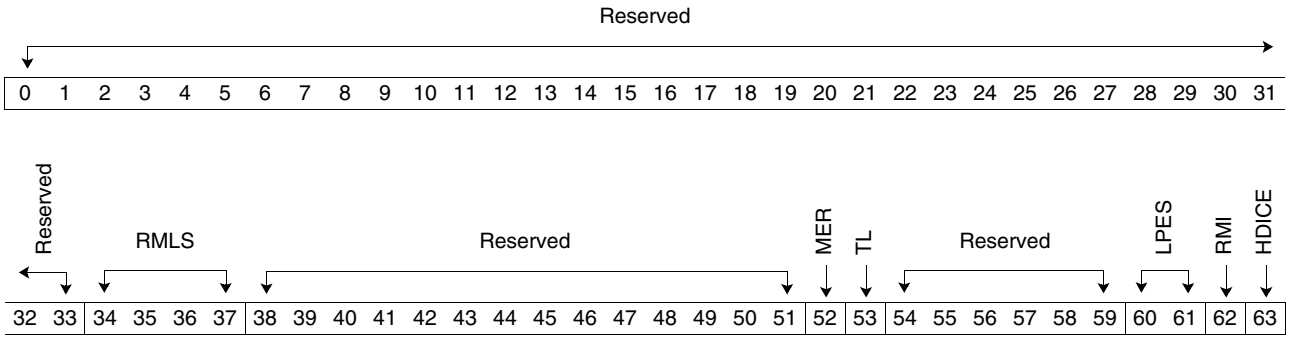
The Hardware Implementation Register 0 (HID0) holds several bits that determine whether a thread will take an interrupt in response to some exceptions and a bit to set issue-serialize mode. Issue-serialize mode, which is enabled with the issue_serialize bit is set to '1', causes the processor to wait until all previous instructions have completed before issuing the next instruction; this mode also prevents dual instruction issue.



| Bits | Field Name | Description |
|---|---|---|
| 3 | issue_serialize | Issue Serialize mode<br>0 Normal operation<br>1 Next instruction is not issued until all previous instructions have completed (no dual-issue either). |
| 22 | therm_wakeup | Enable thermal management interrupt to wake up suspended thread<br>**Note:** Wakeup occurs even if HID0[therm_intr_en] = '0'. |
| 24 | syserr_wakeup | Enable system error interrupt to wakeup suspended thread<br>Allows the system error interrupt to wake up either thread if it is dormant. When a system error interrupt is received, if this bit is enabled the interrupt wakes up the thread that the interrupt was for (can be both threads). For example. if thread 0 is dormant, thread 1 is active, syserr_wakeup is set, and the interrupt is for thread 0, then thread 0 is awakened (the active thread [thread1] is unaffected because it is already awake). If both threads are dormant and if syserr_wakeup is set, the interrupt awakens both threads.<br>**Note:** Wakeup occurs even if HID0[en_syserr] = '0'. |
| 26 | en_prec_mchk | Enable precise Machine Check<br>**Note:** All loads to caching-inhibited (I = '1') space cause the instructions after the load to be flushed, and the thread to be blocked, at dispatch until data is returned for the load. |
| 29 | therm_intr_en | Master thermal management interrupt enable<br>Clearing this bit disables all thermal management interrupts regardless of MSR state.<br>0 Disables all thermal management interrupts regardless of MSR state<br>1 Enabled |
| 30 | en_syserr | Enable system errors generated from outside the PPE.<br>0 Disabled<br>1 Enabled |

### 10.4.3 Logical Partition Control Register (LPCR)

The Logical Partition Control Register (LPCR) holds several partition-related processor control bits. LPCR is a partially shared register: the LPCR[RMLS], LPCR[TL], and LPCR[LPES] fields are shared between threads; the LPCR[MER], LPCR[RMI], and LPCR[HDICE] fields are duplicated per thread.



| Bits | Field Name | Shared or Duplicated | Description |
|---|---|---|---|
| 34:37 | RMLS | Shared by all threads | Real mode limit selector |
| 52 | MER | Duplicated for each thread | Mediate external exception request (interrupt enable) |
| 53 | TL | Shared by all threads | TLB load<br>0 TLB loaded by processor<br>1 TLB loaded by software |
| 60:61 | LPES | Shared by all threads | Logical partitioning (environment selector) |
| 62 | RMI | Duplicated for each thread | Real-mode caching (caching inhibited) |
| 63 | HDICE | Duplicated for each thread | Hypervisor decrementer interrupt control enable |

### 10.4.4 Control Register (CTRL)

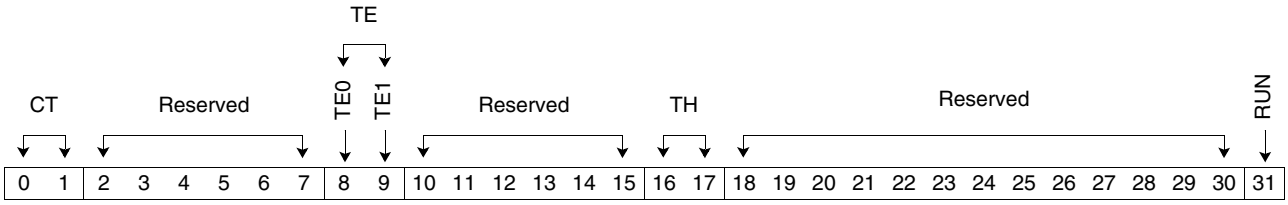
The Control Register (CTRL) holds basic thread status and control bits. Software can use the TE0 and TE1 bits to suspend and resume threads. Hypervisor software executing in one thread can:

- Suspend the present thread
- Resume the other thread, but cannot suspend the other thread

*Table 10-3* shows the effect on thread state when thread 1 or thread 0 sets the thread enable (TE) field. In this table, "—" means the setting has no effect.

*Table 10-3. Effect of Setting the CTRL[TE0] and CTRL[TE1] Bits*

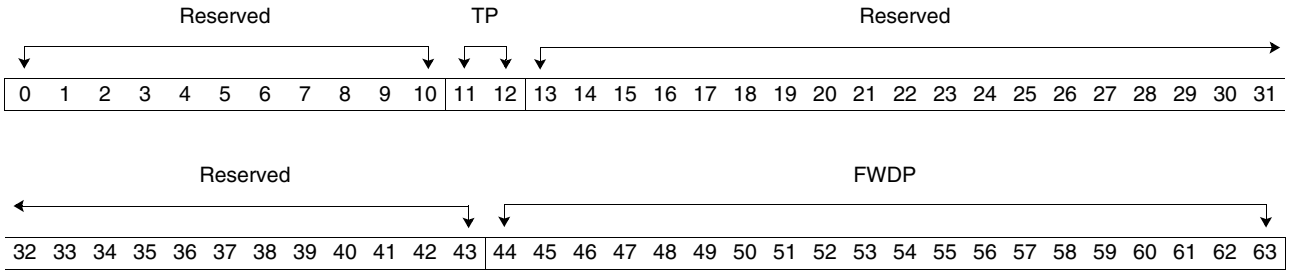
| CTRL[TE0] | CTRL[TE1] | If thread 0 sets TE bits | | If thread 1 sets TE bits | |
|---|---|---|---|---|---|
| | | Thread 0 | Thread 1 | Thread 0 | Thread 1 |
| 0 | 0 | Suspend | — | — | Suspend |
| 0 | 1 | Suspend | Resume | — | — |
| 1 | 0 | — | — | Resume | Suspend |
| 1 | 1 | — | Resume | Resume | — |



| Bits | Field Name | Description |
|---|---|---|
| 0:1 | CT | Current thread active (Read Only)<br>These read-only bits contain the current thread bits for threads 0 and 1. Software can read these bits to determine which thread they are operating on. Only one current thread bit is set at a time.<br>00 Reserved<br>01 Thread 1 is reading CTRL<br>10 Thread 0 is reading CTRL<br>11 Reserved |
| 8:9 | TE | Thread enable bits (Read/Write); see *Table 10-3* on page 310.<br>**Note:** Software should not disable a thread when in trace mode (MSR[SE] or MSR[BE] set to '1'). Doing so causes SRR0 to be undefined and can cause a system livelock hang condition. |
| 16:17 | TH | Thread history (Read only)<br>If thread A writes CTRL[RUN] then CTRL[16] is set; otherwise if thread B writes CTRL[31] then CTRL[17] is set.<br>These bits cannot be set directly by writing bits 16 or 17 with a **mtctrl** instruction. They are only set when a thread writes CTRL[RUN]. |
| 31 | RUN | Run state bit. |

### 10.4.5 Thread Status Register Local and Remote (TSRL and TSRR)

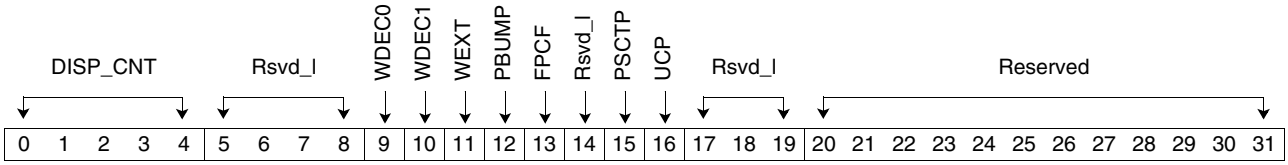
Each thread has a Thread Status Register Local (TSRL) and a Thread Status Register Remote (TSRR). For each thread, the TSRR is a read-only copy of the other thread's TSRL. A thread uses its TSRL to set its priority; a thread uses the TSRR to read the other thread's priority.




| Bits | Field Name | Description |
|---|---|---|
| 11:12 | TP | Thread priority (read/write); see *Table 10-6* on page 320.<br>00 Disabled<br>01 Low<br>10 Medium<br>11 High |
| 44:63 | FWDP | Forward progress timer (Read Only); see *Section 10.6.3* on page 321. |

### 10.4.6 Thread Switch Control Register (TSCR)

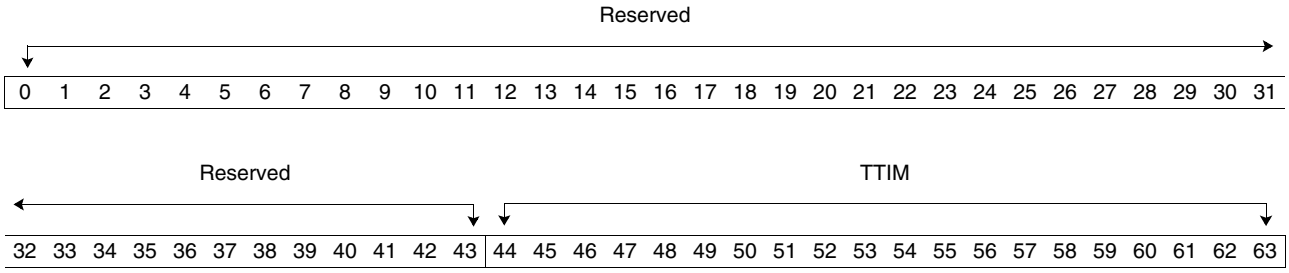
The Thread Switch Control Register (TSCR) has bits to control the thread's response to interrupts, determine its duty cycle in multithreading mode, and so on.




| Bits | Field Name | Description |
|---|---|---|
| 0:4 | DISP_CNT | Thread dispatch count; see *Section 10.5* on page 313. |
| 9 | WDEC0 | Decrementer wakeup enable for thread 0<br>0 Disabled<br>1 If a decrementer exception exists, and the corresponding thread is suspended, then the thread is activated. |
| 10 | WDEC1 | Decrementer wakeup enable for thread 1<br>0 Disabled<br>1 If a decrementer exception exists, and the corresponding thread is suspended, then the thread is activated. |
| 11 | WEXT | External interrupt wakeup enable<br>0 Disabled<br>1 If an external exception exists, and the corresponding thread is suspended, then the thread is activated. |
| 12 | PBUMP | Thread priority boost enable<br>0 Disabled<br>1 If a system-caused exception is presented, the corresponding interrupt is not masked, and the priority of the corresponding thread is less than medium, sets the priority of the thread to medium.<br>The hardware internally boosts the priority level to medium when the interrupt is pending. This does not change the value in the TSRL[TP] bits for the affected thread. The internal priority remains boosted to medium until an **mttsrl** or priority changing **nop** instruction occurs. |
| 13 | FPCF | Forward progress count flush enable<br>This bit only enables or disables the flush from occurring. The forward progress timer does not stop decrementing when set to zero. During normal operation, this bit should be set to '1'.<br>If set to '1', the nonstarved thread will be blocked at dispatch until one instruction completes from the starved thread. See TTR[TTIM] and TSRL[FWDP] and *Section 10.6.3* on page 321 for more information. |
| 15 | PSCTP | Privileged but not hypervisor state change thread priority enable; See *Section 10.6.2.1* on page 319.<br>Enables the privileged state (MSR[HV, PR] = '00') to change priority with "**or** Rx, Rx, Rx" **nop** instructions or writes to TSRL[TP].<br>0 The ability of the privileged state to change thread priority is determined by TSCR[UCP].<br>1 The privileged state can change thread priority to low, medium and high. |
| 16 | UCP | Problem state change thread priority enable; see *Section 10.6.2.1* on page 319.<br>Enables the problem state to change priority with "**or** Rx, Rx, Rx" **nop** instructions or writes to TSRL[TP].<br>0 The problem state cannot change thread priority.<br>1 The problem state can change thread priority to low and medium only. |

### 10.4.7 Thread Switch Time-Out Register (TTR)

The Thread Switch Time-Out Register (TTR) has a single field, TTIM. The value in TTIM is copied into the TSRL[FWDP] field every time an instruction from the corresponding thread completes. Thus, with each instruction completion, the forward-progress timer is reset.




| Bits | Field Name | Description |
|---|---|---|
| 44:63 | TTIM | Thread time-out flush value; see *Section 10.6.3* on page 321 for more information. |

## 10.5 Thread Priority

Each thread has an instruction-dispatch priority determined by the value of the TSRL[TP] field. This 2-bit field can encode four possibilities:

- Thread disabled
- Low priority
- Medium priority
- High priority

Software, in particular operating-system software, sets thread priorities to cause hardware to allocate instruction-dispatch bandwidth according to the throughput requirements of the programs running in the threads. For example, software can set thread priorities to cause hardware to favor a foreground thread over a background thread when allocating instruction-dispatch slots (see *Table 10-5* on page 315).

### 10.5.1 Thread-Priority Combinations

*Table 10-4* on page 314 lists all possible combinations of thread priorities and values of TSCR[DISP_CNT] and describes the instruction-dispatch policy that results from each combination. For some combinations of thread priorities, the value in TSCR[DISP_CNT] allows software to set the duty cycles of instruction dispatching for the two threads.

Though the effects of some of the priority combinations are subtle, there are some basic concepts governing the response of hardware to thread priority settings:

- When the threads have equal priorities, hardware uses fair, round-robin scheduling.
- Dispatch is attempted every cycle except when both threads are set at low priority and the TSCR[DISP_CNT] field is not equal to one.

- When the threads have unequal priorities, the lower-priority thread can be prevented from dispatching any instructions for long periods of time when TSCR[DISP_CNT] is equal to one or when one thread is set to high priority and the other is set to low priority. See *Section 10.6.3 Preventing Starvation: Forward-Progress Monitoring* on page 321 for details.

*Table 10-4. Relative Thread Priorities and Instruction-Dispatch Policies*

| TSCR[DISP_CNT] | Thread 0 Priority | Thread 1 Priority | Attempt to Dispatch | Dispatch Policy Description |
|---|---|---|---|---|
| == 1 | Same | | Every Cycle | Even priority; a round-robin algorithm is used to choose a thread for dispatch. |
| == 1 | Different | | Every Cycle | Uneven priority; the higher-priority thread always gets the dispatch slot if it can use it; otherwise, the lower-priority thread can use the dispatch slot. |
| != 1 | Low | Low | Once every DISP_CNT cycles | Dispatch attempted once in TSCR[DISP_CNT] cycles; no dispatch is attempted on the other TSCR[DISP_CNT]–1 cycles. Even priority; a round-robin algorithm is used to choose a thread for dispatch; if the chosen thread in a given cycle cannot dispatch for other reasons (for example, instruction buffer empty), the other thread is allowed to dispatch (priority can toggle between threads). |
| != 1 | Low | Medium | Every Cycle | Out of every TSCR[DISP_CNT] cycles, one dispatch cycle is given to the lower-priority thread. The other TSCR[DISP_CNT]–1 dispatch cycles are given to the higher priority thread. |
| != 1 | Medium | Low | | |
| != 1 | Medium | High | | |
| != 1 | High | Medium | | |
| != 1 | Low | High | Every Cycle | Out of every TSCR[DISP_CNT] cycles, one dispatch cycle is given to the lower priority thread only if the higher-priority thread cannot dispatch during that cycle. The other TSCR[DISP_CNT]–1 dispatch cycles are given to the higher priority thread. |
| != 1 | High | Low | | |
| != 1 | Medium | Medium | Every Cycle | Even priority; a round-robin algorithm is used to choose a thread for dispatch. If dispatch is stalled for the chosen thread in a cycle, no dispatch occurs in that cycle (priority does not toggle between threads). |
| != 1 | High | High | | |

There are two special cases for the value of TSCR[DISP_CNT]. When the actual value is cleared to '0', hardware uses a dispatch count of 32, the maximum value. When TSCR[DISP_CNT] is set to '1', the higher-priority thread always gets the dispatch slot if it can use it, otherwise the lower-priority thread gets the slot; there is no blocking.

### 10.5.2 Choosing Useful Thread Priorities

This section describes thread priorities from a high-level perspective. A full description of priorities, including low-level details, is given in *Section 10.5* beginning on page 313 and *Section 10.6.2* beginning on page 319.

Each thread can be given one of three priorities: low, medium, or high. Thread priority is normally determined by the value in the Thread Priority field of the Thread Status Register Local (TSRL[TP]). For certain combinations of thread priorities, the Dispatch Count field in the Thread Switch Control Register (TSCR[DISP_CNT]) sets the duty cycle of instruction dispatch for the threads. See *Section 10.5.3* on page 316 for examples of setting TSCR[DISP_CNT].

Together with the value of the TSCR[DISP_CNT] field, there are many possible combinations of thread priorities (see *Table 10-4* on page 314). In practice, thread priorities will most often be set to one of three useful combinations, as shown in *Table 10-5*.

*Table 10-5. Three Useful Thread Priority Combinations*

| Purpose | Thread 0 Priority | Thread 1 Priority | TSCR[DISP_CNT] | Comment |
|---|---|---|---|---|
| Two Foreground Threads (normal operation) | Medium | Medium | Don't care | PPE attempts to dispatch on every cycle; threads alternate dispatch on each attempt. Software can boost a thread to high when minimum latency is important. |
| One Foreground Thread One Background Thread | Medium | Low | Higher value puts Low thread farther in the background | PPE attempts to dispatch on every cycle; medium-priority thread gets most dispatch attempts; low-priority thread gets one dispatch attempt out of every TSCR[DISP_CNT] cycles. |
| | Low | Medium | | |
| Two Background Threads (save power) | Low | Low | Higher value decreases power consumption | PPE attempts to dispatch only once every TSCR[DISP_CNT] cycles; threads alternate dispatch on each attempt. |

The first combination sets both threads to medium priority. This combination causes the PPE to attempt to use all available dispatch opportunities and use a fair, round-robin scheduling algorithm. When a thread needs a temporary boost in priority to, for example, execute a time-critical interrupt handler, it can assume high priority to cause the PPE to choose it for most dispatch opportunities.

The second combination sets one thread at medium priority and one to low. This combination causes the PPE to attempt to use almost all available dispatch opportunities for the medium-priority thread; thus, this combination is appropriate to run a low-priority program in the background during otherwise normal operation. The background program will get some execution time, but it will minimally hinder the foreground thread's progress.

The third combination sets both threads at low priority. This combination causes the PPE to attempt dispatch only once every TSCR[DISP_CNT] cycles with a fair, round-robin scheduling algorithm. With a high value in TSCR[DISP_CNT], the PPE will be mostly idle, which will reduce power consumption and heat production while keeping the two threads alive and responsive to changes in system conditions.

One combination that might cause problems—and therefore is not explicitly supported—is a high-priority thread plus a low-priority thread. In such a case, the PPE would never dispatch instructions from the low-priority thread unless the cycle was the TSCR[DISP_CNT] cycle and the high-priority thread was unable to dispatch. The low priority thread would thus starve for dispatch opportunities. Because of this, the combination of high-priority plus low-priority is explicitly not supported.

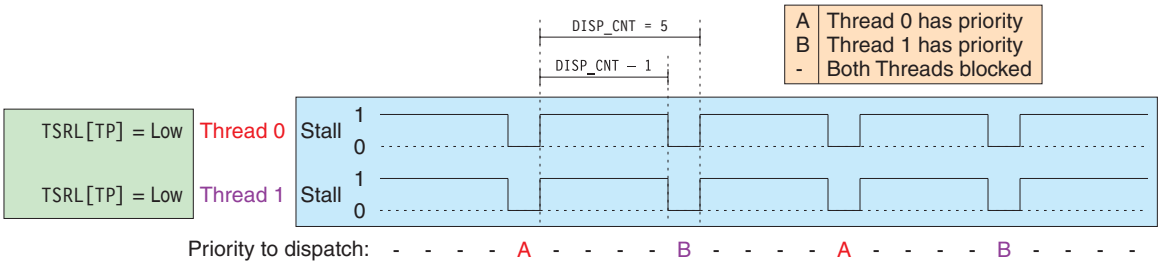
The PPE has a completion-count time-out feature to catch the extreme cases of thread starvation that can result from using this combination. See *Figure 10-4* on page 317 for an illustration of how this combination can cause starvation. See *Section 10.6.1* on page 319 for a description of how to prevent starvation.

### 10.5.3 Examples of Priority Combinations on Instruction Scheduling

#### 10.5.3.1 *Low Priority Combined with Low Priority*

*Figure 10-2* illustrates nominal instruction-dispatch activity when both threads are set at low priority. Because the threads have equal priority, the scheduling is fair, with each thread getting an equal chance to dispatch. The result of this combination is that instructions from thread 0 are dispatched only once every ten cycles; similarly, instructions from thread 1 are dispatched only once every ten cycles.

*Figure 10-2. Thread Priorities and Instruction Dispatch: Low Priority with Low Priority*



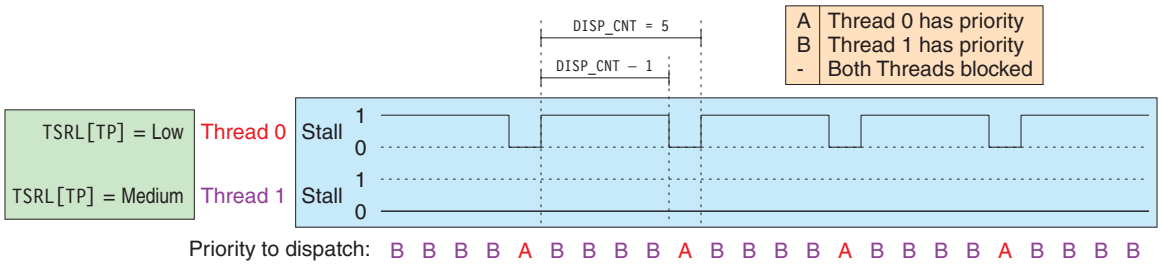

The duty cycle, which is the value in `TSCR[DISP_CNT]`, is set at five. With both threads set at low priority, the processor only dispatches up to two instructions once every five cycles; for this combination, the duty-cycle value describes the ratio of cycles spent attempting to dispatch to cycles spent idling. A given attempt to dispatch an instruction might not succeed if the threads are blocked for other reasons.

The dispatch stall signal from the priority logic releases the dispatch logic only every fifth cycle, and it releases the stall for both threads. If one thread cannot dispatch for other reasons, the other thread will be allowed to dispatch even if it was selected for dispatch on the previous attempt.

#### 10.5.3.2 *Low Priority Combined with Medium Priority*

*Figure 10-3* on page 317 illustrates nominal instruction-dispatch activity when one thread is set at low priority and the other one at medium. Because the threads have unequal priority, the scheduling favors the higher-priority thread. The result of this combination is that instructions from thread 1 are dispatched on four out of five cycles while instructions from thread 0 are dispatched on only one out of five cycles.

*Figure 10-3. Example 2 of Thread Priorities and Instruction Dispatch*



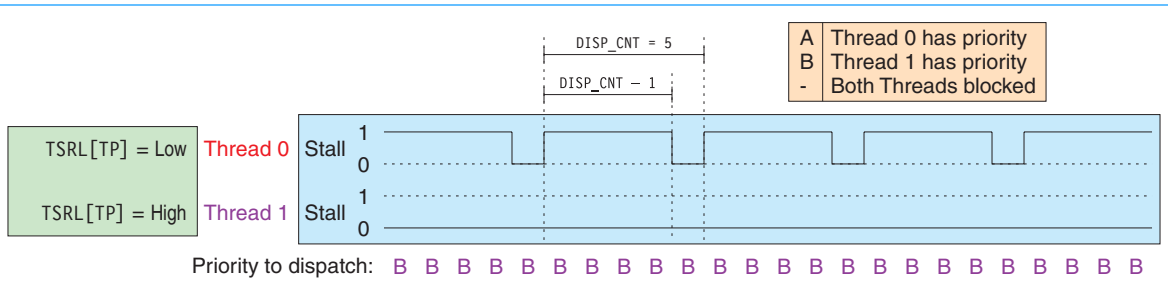

The duty cycle, TSCR[DISP_CNT], is again set at five. With this combination of priorities, the processor dispatches up to two instructions every cycle, and the duty-cycle value describes the ratio of cycles spent on attempts to dispatch exclusively from thread 1 (the higher-priority thread) to cycles spent giving instructions from thread 0 a chance to dispatch.

The dispatch stall signal from the priority logic for thread 1 is always false; with this combination of priorities, instruction dispatch from thread 1 is always allowable. The dispatch stall signal for thread 0 is false (dispatch allowed) only one cycle in five. On this cycle, priority is given to thread 0 if dispatch is not otherwise stalled. If thread 0 is stalled for other reasons, the dispatch logic will attempt to dispatch from thread 1.

#### 10.5.3.3 *Low Priority Combined with High Priority*

*Figure 10-4* illustrates nominal instruction-dispatch activity when one thread is set at low priority and the other at high. The scheduling favors the higher-priority thread to the exclusion of the low-priority thread; in the segment of activity shown, instructions from thread 0 are never given a opportunity to dispatch.

*Figure 10-4. Example 3 of Thread Priorities and Instruction Dispatch*



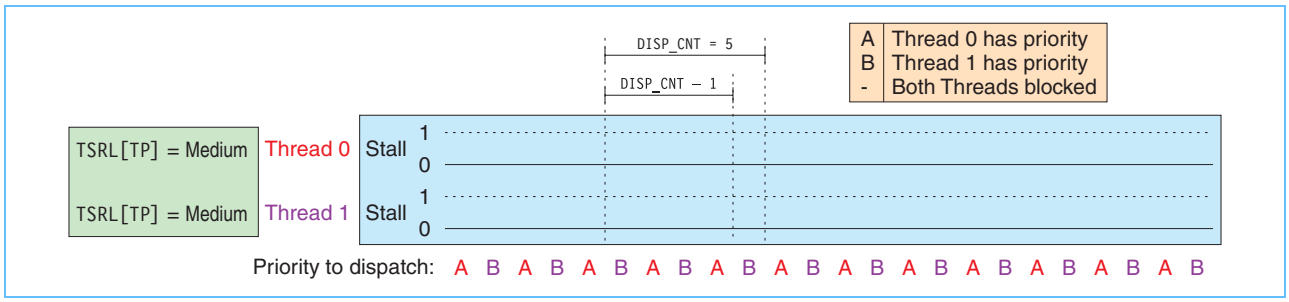

The duty cycle is again set at five. With this combination of priorities, the processor dispatches up to two instructions every cycle, and the duty-cycle value describes the ratio of cycles spent on attempts to dispatch exclusively from thread 1 to cycles spent giving instructions from thread 0 a chance to dispatch. For this combination, however, dispatch from thread 0 is attempted only if thread 1 is unable to dispatch. Consequently, it is possible that no instructions from thread 0 would ever be dispatched because it is possible that thread 1 will always be ready for dispatch. To keep thread 0 from being completely starved in a scenario like this, the PPE can be programmed to monitor forward-progress and force a dispatch opportunity for thread 0 when starvation exceeds a programmable limit (see *Section 10.6.3* on page 321).

Although the stall signal from the priority logic for thread 0 is false every fifth cycle, no instructions from thread 0 are dispatched because, in this example, thread 1 is never stalled for other reasons on the cycles when thread 0 is considered for dispatch.

#### 10.5.3.4 *Medium Priority Combined with Medium Priority*

*Figure 10-5* illustrates nominal instruction-dispatch activity when both threads are set at medium priority; the same dispatch activity applies when both threads are set to high priority. Because the priorities are equal, the scheduling is fair with each thread getting an equal chance for dispatch. The result of this combination of priorities is that the dispatch alternates threads so that a given thread is allowed to dispatch once every other cycle.

*Figure 10-5. Example 4 of Thread Priorities and Instruction Dispatch*

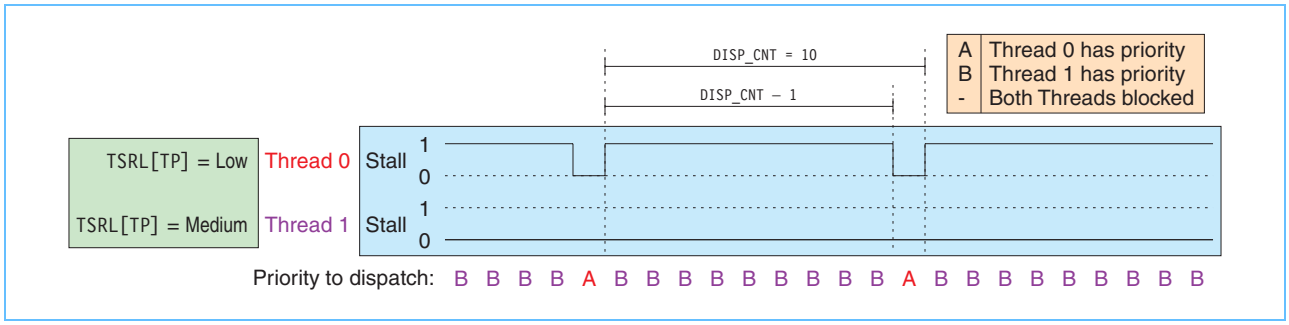



The duty cycle is again five, but with this combination of priorities, the duty cycle value, `TSCR[DISP_CNT]`, plays no part in choosing a thread for dispatch consideration. Dispatch is attempted on every cycle, and the dispatch logic alternates the thread it chooses for consideration. If the thread under consideration in a given cycle is unable to dispatch any instructions, the other thread will be allowed to dispatch even if it was selected for dispatch on the previous attempt.

#### 10.5.3.5 *Low Priority Combined with Medium Priority and Larger `TSCR[DISP_CNT]`*

*Figure 10-6* illustrates nominal instruction-dispatch activity when one thread is set at low priority and the other at medium. The situation and resulting dispatch behavior are identical to that explained in the second example in *Section 10.5.3.2* on page 316, but the duty-cycle value is set at ten, which causes the dispatch logic to give an even greater preference to instructions from thread 1. The ratio is nine-to-one instead of four-to-one as in the previous example.

*Figure 10-6. Example 5 of Thread Priorities and Instruction Dispatch*



## 10.6 Thread Control and Configuration

Software can configure and control thread behavior in several ways, including:

- Enabling and suspending a thread
- Setting the instruction-dispatch policy by setting thread priority
- Executing a special **nop** instruction to cause temporary dispatch stalls
- Setting an appropriate forward-progress timer value to allow hardware to guarantee a minimum amount of execution time
- Switching the state of the processor between single-threaded and multithreaded mode

### 10.6.1 Resuming and Suspending Threads

The only way software can suspend a thread is to execute an **mtspr** in hypervisor state that clears its corresponding Thread Enable bit to zero in the CTRL register (CTRL[TE0] or CTRL[TE1]).

Software executing in hypervisor state can resume execution for a suspended thread by setting the appropriate Thread Enable bit in the CTRL register (CTRL[TE0] or CTRL[TE1]). Hardware can resume a suspended thread with certain enabled interrupts. In either case, when the thread resumes, it begins execution at the system reset interrupt vector location. See *Section 10.7.2.9* on page 324 for detailed information about resuming and suspending threads.

### 10.6.2 Setting the Instruction-Dispatch Policy: Thread Priority and Temporary Stalling

Software running in a thread can influence how hardware dispatches the thread's instructions into the pipeline in two ways:

- The thread can set its priority level to low, medium, or high (depending on privilege state).
- The thread can execute a special **nop** instruction that suspends instruction dispatch for a specific number of processor cycles.

The priority level of one thread relative to the priority level of the other thread determines how the processor will allocate instruction dispatch bandwidth to the thread. If the thread has a higher priority than then the other thread, the hardware will, in general, dispatch the thread's instructions more frequently than it will dispatch the other thread's instructions. See *Section 10.5* on page 313 for a complete description of thread priorities and instruction-dispatch policy. The next section explains how a thread can change its priority and what changes are allowed.

A thread can also deliberately stall the dispatch of its instructions by executing a special **nop** instruction, as explained later.

#### 10.6.2.1 *Allowed Thread Priority Changes*

*Table 10-6* on page 320 shows the changes to thread priority that a thread is allowed to make. The priority changes that are allowed depend on the privilege state of the thread and the priority-change bits in the TSCR. A thread must be in hypervisor or supervisor state to set its priority to high.

*Table 10-6. Privilege States and Allowed Thread-Priority Changes*

| MSR[HV] | MSR[PR] | Privilege State | TSCR[PSCTP] | TSCR[UCP] | Can Suspend? | Thread Priority Change Allowed | | |
|---|---|---|---|---|---|---|---|---|
| | | | | | | High | Medium | Low |
| 1 | 0 | Hypervisor | Don't care | Don't care | Yes | Yes | Yes | Yes |
| 0 | 0 | Supervisor | 0 | 0 | No | — | — | — |
| | | | 0 | 1 | No | — | Yes | Yes |
| | | | 1 | Don't care | No | Yes | Yes | Yes |
| Don't care | 1 | User (Problem State) | Don't care | 0 | No | — | — | — |
| | | | | 1 | No | — | Yes | Yes |

A thread can attempt to change its priority level in two ways:

- The thread can write to its TSRL[TP] field with the required priority encoding. If the change is not allowed, no action is taken. Writing '0' to this field will not disable the thread regardless of operating mode; writing '0' is always ignored.
- The thread can execute one of the special **nop** instructions listed in *Table 10-7* on page 320 that change thread priority as a side-effect.

By default, the hardware initializes the priority-change bits to allow priority changes only in hypervisor state. Hypervisor software can set the priority-change bits to allow software in other privilege states to change thread priority.

One thread cannot change the other thread's priority level. If software attempts this or some other priority change that is not allowed, the attempt is silently ignored by the processor.

If the TSCR[PBUMP] bit is set to '1' when an enabled interrupt is pending and not masked as a result of a system-caused exception, hardware will temporarily boost the thread's priority to medium if it was less than medium.

#### 10.6.2.2 *nop Instructions that Change Thread Priority and Dispatch Policy*

In addition to the instruction-dispatch policy set by the combination of thread priorities, a thread can execute a special **nop** instruction to cause the processor hardware to suspend the dispatch of the thread's instructions for a specific number of cycles. Also, a thread can attempt to change its priority by executing one of a set of special **nop** instruction. See *Table 10-7*.

*Table 10-7. **nop** Instructions to Set Thread Priority and Stall Instruction Dispatch (Sheet 1 of 2)*

| OR Form | Extended Mnemonic | Description/Side-Effects |
|---|---|---|
| ori 0, 0, 0 | **nop** | Preferred form for **nop** instruction without side-effects. |
| or 1, 1, 1 | cctpl | Change current thread priority to low. |
| or 2, 2, 2 | cctpm | Change current thread priority to medium. |
| or 3, 3, 3 | cctph | Change current thread priority to high. |
| or 28, 28, 28 | db8cyc | Cause the current thread to block at dispatch for eight cycles. Used for multithread compiler optimizations of long-latency vector/scalar unit (VSU) dependencies. |

*Table 10-7. **nop** Instructions to Set Thread Priority and Stall Instruction Dispatch (Sheet 2 of 2)*

| OR Form | Extended Mnemonic | Description/Side-Effects |
|---|---|---|
| `or 29, 29, 29` | `db10cyc` | Cause the current thread to block at dispatch for 10 cycles. Used for multithread compiler optimizations of long-latency VSU dependencies. |
| `or 30, 30, 30` | `db12cyc` | Cause the current thread to block at dispatch for 12 cycles. Used for multithread compiler optimizations of long-latency VSU dependencies. |
| `or 31, 31, 31` | `db16cyc` | Cause the current thread to block at dispatch for 16 cycles. Used for multithread compiler optimizations of long-latency VSU dependencies. |
| `ori 1-31, 1-31, 0`<br>`or 0, 0, 0`<br>`or 4-27, 4-27, 4-27` | N/A | Nonspecial forms for **nop** instruction. |

### 10.6.3 Preventing Starvation: Forward-Progress Monitoring

When one thread is assigned high priority and the other low priority, the low-priority thread will be given at most one chance out of `TSCR[DISP_CNT]` cycles to dispatch an instruction, and on that cycle, the low-priority thread will be denied the dispatch opportunity if the high-priority thread can use the dispatch slot. This situation can lead to complete starvation of the low-priority thread; it might make no forward progress for an excessively long time.

To prevent undesirable starvation, the PPE can be programmed to monitor the forward progress of each thread and take action when one thread has made no forward progress and the other has made more than a pre-set limit. The monitoring hardware uses the `TTR[TTIM]` field, the `TSRL[FWDP]` field, and the `TSCR[FPCF]` bit. Each thread has a dedicated set of these resources.

To start the hardware forward-progress monitor for a thread, an instruction sequence like the following can be used:

```
li      r10,0x1000    // other thread allowed 4096 instruction completions
mtspr   TTR,r10       // Set TTR[TTIM]
li      r10,(1<<19)   // FPCF bit position
mfspr   r11,TSCR      // get current TSCR
or      r11,r11,r10   // set the FPCF bit position to 1
mtspf   TSCR,r11      // set FPCF bit to 1 in TSCR
```

The forward-progress monitoring hardware will copy the `TTR[TTIM]` field into the `TSRL[FWDP]` field each time the thread successfully completes an instruction; thus, each time a thread completes an instruction, its corresponding forward-progress timer is reset to the initial value. In the preceding instruction sequence, when the second **li** instruction completes, the value in the `TTR[TTIM]` field—4096—will be copied into the `TSRL[FWDP]` field.

On cycles when the thread does not complete an instruction, the hardware will decrement the `TSRL[FWDP]` field by one if the other thread does complete an instruction; if neither thread completes an instruction, the `TSRL[FWDP]` value does not change. Thus, a thread's `TSRL[FWDP]` will be exhausted only if that thread completes no instructions while the other thread completes enough instructions to reduce the value in the `TTR[TTIM]` field to one.

The action taken when the TSRL[FWDP] field reaches one is determined by the TSCR[FPCF] bit: if this bit is set to '1', the other thread is flushed after the next instruction completes; if this bit is cleared to '0', no action is taken regardless of the value in TSRL[FWDP]. After the other thread is flushed, its dispatch is stalled until the starved thread completes an instruction.

The TSCR[FPCF] is only an enable for the hardware flush of the opposite thread; even when TSCR[FPCF] is cleared to '0' (flush of opposite thread disabled), the thread's TSRL[FWDP] field will decrement and be reloaded as previously described (assuming the thread is enabled).

### 10.6.4 Multithreading Operating-State Switch

When software wants to change the operating mode of the PPE from single-threaded to multi-threaded, or vice versa, a set of events must occur in the proper sequence to guarantee correct hardware and software behavior.

#### 10.6.4.1 *Changing from Single-Threaded to Multithreaded Operation*

The steps, in hypervisor mode, to change the PPE operating mode from single-threaded to multi-threaded are as follows:

1. Perform a normal context-serializing operation such as **isync** that waits for all pipelines and queues to drain.
2. Issue the **mtspr** instruction for the CTRL register to enable the second thread.

#### 10.6.4.2 *Changing from Multithreaded to Single-Threaded Operation*

The steps, in hypervisor mode, to change the PPE operating mode from multithreaded to single-threaded are as follows:

1. Perform a normal context-serializing operation such as **isync** that waits for all pipelines and queues to drain.
2. Disable interrupts and trace mode (MSR[EE], MSR[SE], and MSR[BE] cleared to '0').
3. Issue an **mtspr** CTRL instruction to disable your thread.

## 10.7 Pipeline Events and Instruction Dispatch

### 10.7.1 Instruction-Dispatch Rules

In a single-threaded processor implementation, instruction dispatch is governed by rules that consider many conditions in the processor such as:

- The availability of instructions to dispatch
- The availability of execution units to process instructions ready for dispatch
- Dependencies on previously dispatched instructions
- Branch misprediction

In the multithreaded PPE, instruction dispatch is governed by additional rules that consider thread-related conditions such as:

- Relative thread priority.
- A thread has made insufficient forward progress (forward-progress monitor decrements to a count of one).
- A thread experiences a cache miss.

The following sections describe internal conditions that stall instruction dispatch.

## 10.7.2 Pipeline Events that Stall Instruction Dispatch

### 10.7.2.1 *Dependency on a Load that Misses the DCache*

When an instruction is detected that has a dependency on an older load instruction that missed the DCache, the dependent instruction and all younger instructions for the same thread are flushed (10 cycles after the dependent instruction is issued from the IS2 stage). The flushed instructions are refetched, and dispatch for the flushed thread is then stalled until the load data is available in the DCache.

### 10.7.2.2 *Dependency on a Store-Conditional Instruction*

If an instruction is issued with a dependency on an **stdcx.** or **stwcx.** instruction that has not yet updated the CR, the dependent instruction and all younger instructions for the same thread are flushed (10 cycles after the dependent instruction is issued from the IS2 stage). The flushed instructions are refetched, and dispatch for the flushed thread is then stalled until the CR is updated.

### 10.7.2.3 *Load or Store Instruction D-ERAT Miss*

When a load or store instruction causes a D-ERAT miss, all younger instructions for the same thread are flushed once the instruction causing the miss is 10 cycles past the IS2 stage. The flushed instructions are refetched, and dispatch for the flushing thread is then stalled until the load data is available in the D-ERAT.

Only one outstanding D-ERAT miss from either thread is allowed. If a subsequent D-ERAT miss occurs for the other thread, both the instruction causing the miss and all younger instructions for the same thread are flushed, and dispatch is then stalled for both threads until the first D-ERAT miss is resolved.

### 10.7.2.4 *Microcoded Instruction*

The microcode engine is shared between both threads. When a microcode instruction is dispatched, both threads are stalled at dispatch until the microcode sequence is complete. The microcode sequence is complete when the last microcode operation is in the ID1 stage. In the following cycle, instructions from the other (nonmicrocode) thread can be dispatched and placed into ID1. Therefore, invoking microcode will cause a performance penalty on both threads.

#### 10.7.2.5 *System-Caused Interrupt Pending*

When a system-caused interrupt is pending and not masked, dispatch for both threads is stalled until the interrupt is masked or serviced. The processor waits for the current instructions that are already committed to finish. Once it has been determined whether or not the committed instructions take an interrupt that masks the system-caused interrupt, or the instructions finish cleanly and the system-caused interrupt is taken, the dispatch stall is removed.

#### 10.7.2.6 *Forward-Progress Timer Timeout*

The forward-progress monitoring facility is described in *Section 10.6.3* on page 321.

When the TSRL[FWDP] count becomes one, the monitoring logic records that a forward-progress timeout condition has occurred. The TSRL[FWDP] counter continues to decrement while committed instructions of the opposite thread complete; the TSRL[FWDP] counter stops decrementing when it reaches x‘00001’. When the already committed instructions complete, the opposite thread will be flushed and dispatch will be blocked until the starved thread completes an instruction.

#### 10.7.2.7 *Caching-Inhibited Load Issued with Precise Machine Check Interrupts Enabled*

If precise machine check interrupts are enabled (HID0[en_prec_mchk] = 1) and a caching-inhibited (I = 1) load instruction is issued, then all younger instructions for the thread that issued the load are flushed, refetched, and then stalled at dispatch. The stall is released when the caching-inhibited load instruction completes.

#### 10.7.2.8 *Thermal Throttling*

When the thermal sensors indicate the chip is overheating, or to save power, the processor can detect and signal a thermal throttling condition. When signalled, dispatch is blocked for one or both threads.

#### 10.7.2.9 *FPU Busy*

When the FPU is executing a long-latency floating point instruction (**fdiv** or **fsqrt**) or any floating point instruction whose operand is denormalized number, the PPE recycles the operation in the FPU issue queue, which causes a stall that might propagate back to dispatch.

#### 10.7.2.10 *Sync*

A **sync** (L = 0 or 1) instruction blocks both threads until the **sync** instruction completes. The **eieio** and **tlbsync** instructions do not cause threads to stall. The **isync** instruction stalls both threads until all previous instructions complete, then the **isync** causes a flush and refetch on its own thread. The instructions from the other thread that are already fetched or dispatched are not refetched after the **isync** completes.

## 10.8 Suspending and Resuming Threads

By adjusting thread priorities and other instruction-dispatch controls, software can reduce the rate of forward progress of a thread so that it is active but nearly disabled. To disable a thread completely, software can clear to '0' the thread's CTRL[TE0] or CTRL[TE1]. When a thread is disabled, the fetch hardware will not fetch instructions into its instruction buffer, and the instruction-dispatch logic will not consider it for priority when making decisions for instruction dispatch. This means that the other, enabled thread (if not in low priority) will be granted all available instruction-fetch and instruction-dispatch bandwidth.

### 10.8.1 Suspending a Thread

The only way software can suspend a thread is to execute the **mtctrl** instruction with an operand that leaves the thread's Thread Enable bit cleared to '0'.

A write to CTRL[TE0] or CTRL[TE1] will succeed only when the thread is executing in hypervisor state (MSR[PR,HV] == 0,1). If the write is attempted in supervisor or problem state, the write is ignored silently. Attempting to clear the Thread Enable bit for the thread that is not executing the **mtctrl** instruction will be ignored, regardless of the thread's privilege state.

If a thread suspends itself and the other thread remains active, the thread suspension will be treated as a context-synchronizing event for the active thread.

### 10.8.2 Resuming a Thread

The actions taken to resume a suspended thread are essentially the same as the actions taken to start the initial thread after the POR sequence completes. In both cases, the thread resumes execution at the system reset interrupt vector location. The system reset interrupt handler can determine whether the System Reset handler is running because of a POR event or a normal thread-resuming event and handle the event appropriately. The handler makes the determination based on the values in the Save and Restore Registers.

#### 10.8.2.1 *Thread-Resuming Events*

A previously suspended thread can be resumed by any of the following events:

- A thread executing in hypervisor state (MSR[PR,HV] == 0,1) can resume the other thread by setting to one the corresponding Thread-Enable bit (CTRL[TE0] or CTRL[TE1]) with the mtctrl instruction.
- The occurrence of a thermal management interrupt can resume a thread if the HID0[therm_wakeup] bit is set to '1'.
- The occurrence of a system error interrupt can resume a thread if the HID0[syserr_wakeup] bit is set to '1'.
- The occurrence of an external interrupt can resume a thread if the TSCR[WEXT] bit is set to '1'.
- The occurrence of a decrementer interrupt for the suspended thread can resume the thread if the corresponding Decrementer Wakeup Enable bit (TSCR[WDEC0] or TSCR[WDEC1])is set to '1'.

A machine-check interrupt cannot occur for a suspended thread, and while a hypervisor decrementer interrupt can occur for a thread, it does not resume a suspended thread.

When a system-caused exception occurs, the thread resumes even if the interrupt that corresponds to the exception is masked.

After POR, the following interrupt mask bits are set to mask interrupts:

- HID0[therm_wakeup,syserr_wakeup] = '00'
- TSCR[WDEC0,WDEC1] = '00'
- TSCR[WEXT] = '0'

With these mask bits cleared to '0', a suspended thread will not resume on any of the corresponding exceptions.

#### 10.8.2.2 *Hardware Resume Process*

A resumed thread starts execution at the system reset interrupt vector location. Hardware sets the TSRL[TP] field to '11', which sets the thread priority to high. The thread priority is set to high regardless of the condition that caused the thread to resume.

When the resumed thread begins executing at the system reset interrupt vector location, the System Reset Reason Code field of the Machine Status Save Restore Register 1 (SRR1[42:44]) is set to one of the thread resume reason codes shown in *Table 10-8*.

*Table 10-8. Thread Resume Reason Code*

| SRR1[42:44] | Resume Reason | Condition |
|---|---|---|
| 101 | System Reset | Write to CTRL[TE] field |
| 010 | Thermal Management | HID0[therm_wakeup] |
| 110 | System Error | HID0[syserr_wakeup] |
| 100 | External | TSCR[WEXT] |
| 011 | Decrementer | TSCR[WDEC] |

If multiple reasons exist for the thread to resume, the reason code reported in SRR1[42:44] will respond to one of the reasons that caused the wakeup (there is no predetermined priority).

The fields in SRR1 other than the System Reset Reason Code field contain values from the MSR as if the thread had never been suspended. The Machine Status Save/Restore Register 0 (SRR0) similarly contains the address of the instruction that would have been executed next if the thread had not been suspended.

The system reset interrupt handler can use the information in SRR0 and SRR1 to resume the suspended thread properly. Because the same reason code is used for both System Reset resume and a software-initiated mtctrl resume, the system reset interrupt handler must examine the value of HID0 or HID1 to determine the cause of the thread resume action.

Because a system reset event can happen only at POR, if software finds HID0 or HID1 set to its default POR values, software knows the thread resume was the result of a system reset event. If software finds that the values in the Hardware-Implementation Dependent (HID) registers differ from their POR default values, software knows the thread resume was the result of a **mtctrl** instruction.

#### 10.8.2.3 *Accepting Interrupts after Resuming a Thread*

When a thread is resumed by a thermal management, system error, external, or decrementer interrupt, the resuming interrupt is then masked because `MSR[EE,HV]` will be set to '01' when the resumed thread is started at the system reset interrupt vector location. If software later enables the interrupt by changing the `MSR` settings and the exception is still pending, the interrupt that caused the thread to be resumed will be taken. In other words, a resuming interrupt will cause a suspended thread to resume but will not cause the resumed thread to take the resuming interrupt until the resumed thread enabled interrupts.

When two threads are resumed in the same clock cycle (a rare event), it is indeterminate which thread is resumed first, but both threads will resume properly. When one thread is active and the other is inactive and the inactive thread is resumed, the active thread will view the thread resumption as a context-synchronizing event.

Between the time the thread is resumed and the time the thread enables interrupts, additional exceptions can occur. The interrupt that will be taken after the thread enables interrupts is not necessarily the interrupt that caused the thread to resume; the interrupt taken by the thread will be determined according to the interrupts pending and the interrupt priorities at the time the interrupt is taken.

### 10.8.3 Exception and Interrupt Interactions With a Suspended Thread

The following thread-dependent interrupts are handled by the thread that causes the exception. These interrupts are enabled for both threads as shown below:

- External — enabled by `TSCR[WEXT]`
- Precise Machine Check — enabled by `HID0[en_prec_mchk]`
- System Reset — enabled by `HID0[en_syserr]`
- Thermal Management — enabled by `HID0[therm_wakeup]`
- System Error — enabled by `HID0[syserr_wakeup]`

#### 10.8.3.1 *Decrementer Interrupt*

Each thread has a separate Decrementer Register that creates a separate decrementer interrupt for the corresponding thread. `TSCR[WDEC0]` or `TSCR[WDEC1]` enables waking up a suspended thread due to a decrementer interrupt.

#### 10.8.3.2 *Hypervisor Decrementer Interrupt*

The single Hypervisor Decrementer Register is shared by both threads. A hypervisor decrementer interrupt does not resume a suspended thread.

#### 10.8.3.3 *Imprecise Machine Check Interrupt*

Imprecise machine check interrupts do not have thread information associated with them. All enabled threads take an imprecise machine check interrupt regardless of which thread caused the exception. The MSR[ME] bit must be set to '1' for both threads (even suspended threads) to avoid causing a checkstop. An imprecise machine check interrupt does not resume a suspended thread. If the thread is later resumed, it will not take an imprecise machine check interrupt.

#### 10.8.3.4 *Maintenance Interrupt*

An instruction-caused maintenance interrupt is thread-dependent and is handled by the thread that caused the exception. A system-caused maintenance interrupt is shared by both threads and occurs when one or both of the two trace input signals, which are shared by both threads, is asserted. The system-caused maintenance interrupt is handled by the thread that enables the interrupt.

#### 10.8.3.5 *Floating-Point Enabled Exceptions*

When either floating-point exception mode bit (MSR[FE0] and MSR[FE1]) is set for either thread (it does not matter which thread is executing), the PowerPC processor unit (PPU) switches into a single-instruction-issue mode. In this mode, only one instruction is active in execute at a time, and all other instruction from either thread are held at issue until the previous instruction completes.

#### 10.8.3.6 *Other Interrupts*

All other interrupt are thread-dependent and are handled by the thread that caused the exception. For details, see *Section 9 PPE Interrupts* on page 239.

### 10.8.4 Thread Targets and Behavior for Interrupts

When an exception is signalled or an interrupt is taken, the PPE interrupt logic selects a thread to handle the event. The chosen thread depends on the exception type and the interrupt type if the exception results in an interrupt. *Table 9-38* on page 291 shows the target threads for each exception and interrupt type.

The thread that the interrupt logic selects as the target can be in the suspended state or enabled state. If the thread is enabled, it will take the interrupt normally. If the thread is suspended, it will either be resumed or allowed to remain suspended.

*Table 10-9* on page 329 shows the entry condition, activated thread, mask condition, and suspended-thread behavior for each interrupt type. For example, a machine-check interrupt will target both threads, but it will be recognized only by the threads that are enabled; suspended threads are not resumed by machine-check interrupts. In this table, a "—" symbol means "not applicable".

*Table 10-9. Interrupt Entry, Masking, and Multithreading Behavior (Sheet 1 of 2)*

| Interrupt Type | Entry Condition | Activated Thread | Mask Condition | Suspended Thread Behavior |
|---|---|---|---|---|
| System Reset | POR | thread 0 | — | If SRR1[42:44] = '101', thread resumed due to system reset. |
| | wakeup | specific thread | — | This exception is where suspended thread resumes by means of a write to CTRL[TE01] and SRR1[42:44] = '101' |
| Machine Check | MC interrupt | all active threads | MSR[ME] | Remain suspended. |
| Data Storage Interrupt | memory access | specific thread | — | — |
| Data Segment Interrupt | memory access | specific thread | — | — |
| Instruction Storage Interrupt | memory access | specific thread | — | — |
| Instruction Segment Interrupt | memory access | specific thread | — | — |
| External Direct | external interrupt if HID0[extr_hsrr] = '0' | programmable | MSR[EE] | If TSCR[WEXT] = '1', wake up at System Reset and set SRR1[42:44] = '100'. |
| External Direct (Mediated mode) | external interrupt if HID0[extr_hsrr] = '1' | programmable | MSR[EE] \| (LPCR[LPES] bit 0 = '0' & MSR[HV] = '0') \| MSR[SV] | If TSCR[WEXT] = '1', wake up at System Reset and set SRR1[42:44] = '100'. |
| External Mediated | HID0[extr_hsrr] = '1' & LPCR[MER] = '1' | programmable | MSR[EE] & (!MSR[HV] \| MSR[PR]) | If TSCR[WEXT] = '1', wake up at System Reset and set SRR1[42:44] = '100'. |
| Alignment | misaligned access | specific thread | — | — |
| Program | Invalid operation | specific thread | — | — |
| Floating-Point Unavailable | any floating operation if MSR[FP] = '0' | specific thread | — | — |
| Decrementer | decrementer timeout | specific thread | MSR[EE] | If TSCR[WDEC] = '1', wake up at System Reset and set SRR1[42:44] = '011'. |
| Hypervisor Decrementer | hypervisor decrementer timeout | all active threads | MSR[EE] \| !MSR[HV] & LPCR[HDICE] | remain suspended |
| System Call | **sc** instruction | specific thread | — | — |
| Trace | any operation if MSR[SE] = '1' or MSR[BE] = '1' | specific thread | — | — |
| VXU Unavailable | VXU operation if MSR[VXU] = '0' | specific thread | — | — |
| System Error | Machine Check outside PPE if HID0 [en_mchk] = '1' | all active threads | MSR[EE] & !MSR[HV] | If HID0[syserr_wakeup] = '1', wake up at System Reset and set SRR1[42:44] = '110'. |

*Table 10-9. Interrupt Entry, Masking, and Multithreading Behavior (Sheet 2 of 2)*

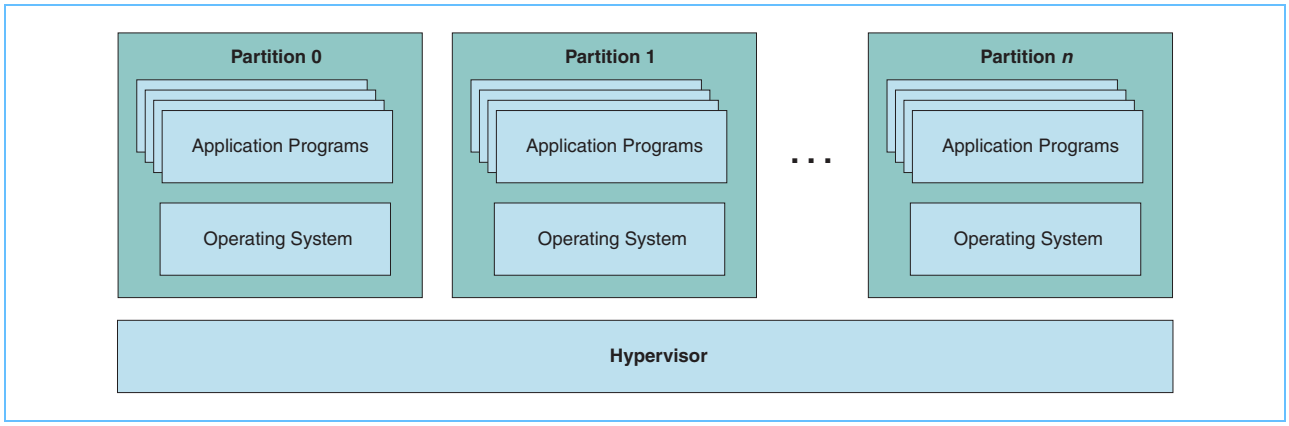
| Interrupt Type | Entry Condition | Activated Thread | Mask Condition | Suspended Thread Behavior |
|---|---|---|---|---|
| Maintenance | instruction caused (HSRR1[33] = '0') | specific thread | — | — |
| Maintenance | system (trace sig) (HSRR1[33] = '1') | all threads | MSR[EE] & !MSR[HV] | — |
| Thermal Management | thermal condition | all active threads | HID0[therm_intr_en] | If `HID0[therm_wakeup]` = '1', wake up at System Reset and set SRR1[42:44] = '010'. |

# 11. Logical Partitions and a Hypervisor

## 11.1 Introduction

Logical partitioning refers to the capability to logically subdivide a single system into multiple partitions, each capable of running its own operating system and user environment such that a program executing in one partition cannot interfere with any program executing in a different partition. *Figure 11-1* illustrates the concept of logical partitioning and the hypervisor partition-management software that controls it.

*Figure 11-1. Overview of Logical Partitions*




Logical partitioning can be useful for several purposes, including:

- In software-production systems, to temporarily dedicate system resources to a testing effort for new software or applications before bringing them fully online in mission-critical environments.
- In business environments to reduce cost by means of system consolidation; the consolidation involves moving workloads from many systems to a single system that is logically partitioned to mimic the environment of the many systems.
- In consumer-electronic and embedded systems, to support environments such as a real-time operating system executing in one partition, a conventional operating system executing in another partition, and a security subsystem executing in a third partition.
- In aggregated computer centers, Internet service providers, and application-service providers, where fractional provisioning of computing and other system resources can support many different and perhaps competing workloads.
- In mission-critical applications, to support a complete hot-plug architecture, including hot-plug replacement of DRAM modules and I/O devices without the operating system being involved or even knowing what happened.

In these environments, separation of applications and data must be maintained while the required performance and system resources are provided for each partition. Reduced cost through logical partitioning can be realized, because hardware resources can be supplied to

each partition, and dynamic provisioning can be used to efficiently move resources from lightly-used partitions to those that are experiencing high or peak demand. This kind of resource allocation reduces a system's requirement to over-provision to meet all worst-case loads.

### 11.1.1 The Hypervisor and the Operating Systems

In a logically-partitioned system, partition-management software—called a *hypervisor*—creates the partitions by assigning portions of the total system resources (processor-element execution time, real memory, and I/O devices) to each partition. The hypervisor executes at a more-privileged level than an operating system in a partition.

Separate instances of an operating system, or different operating systems, can execute in each partition. The hypervisor, with hardware assistance, ensures that the operating system and applications in one partition have no knowledge of, or access to, the resources of another partition, unless specifically granted by the hypervisor. Exceptions to this rule can be made in which the hypervisor brokers a communication between partitions, or one partition publishes services that can be used by other partitions.

### 11.1.2 Partitioning Resources

System resources are typically statically assigned by a primary partition to other partitions or by an administrator's use of a hardware management console (typically a serial port). In these environments, the policy of how many partitions to create, and with what resources, is separated from the mechanism for partitioning. For consumer-electronic or embedded-system applications, for example, a primary partition can first be created by the hypervisor. This partition can then use hypervisor interfaces to create and allocate resources for other partitions. Dynamic-provisioning techniques can move resources from one partition to another, based on demand and the need to meet user-level requirements, such as interactive or real-time response guarantees.

The hypervisor is typically implemented in a small executive that is packaged as firmware. It provides a newly-created partition with access to resources and it harvests those resources when that partition is deleted.

If the hypervisor is unable to dedicate an entire physical resource to a partition, it can virtualize the physical resource into a set of logical resources that are time-shared. For example, if a fraction of the PowerPC Processor Element (PPE) is provided to a partition, the hypervisor can manifest the physical PPE as multiple virtual PPEs that are assigned to the logical partitions in time-multiplexed fashion—for example, by using hypervisor decrementer (HDEC) interrupts to drive the transitions. If an I/O device cannot be logically partitioned, such that each operating system has direct access to its assigned I/O, the hypervisor can virtualize the I/O device so that the device can be shared, or so that an I/O-hosting partition can own the physical device and host other logical partitions' I/O requirements.

The Cell Broadband Engine Architecture (CBEA) processors[1] provide hardware facilities for logical partitioning in the PPE, the SPEs, the I/O controller, and the element interconnect bus (EIB). The standard PowerPC-based hypervisor logical-partitioning facilities are supported, as well as extended facilities for the SPEs, I/O resources, and EIB. The CBEA processors also support a resource allocation management (RAM) facility that can be used in conjunction with the

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

logical-partitioning facilities to control the amount of bandwidth allocated to units in partitions accessing shared resources. The RAM facility is described in *Section 8 Resource Allocation Management* on page 203.

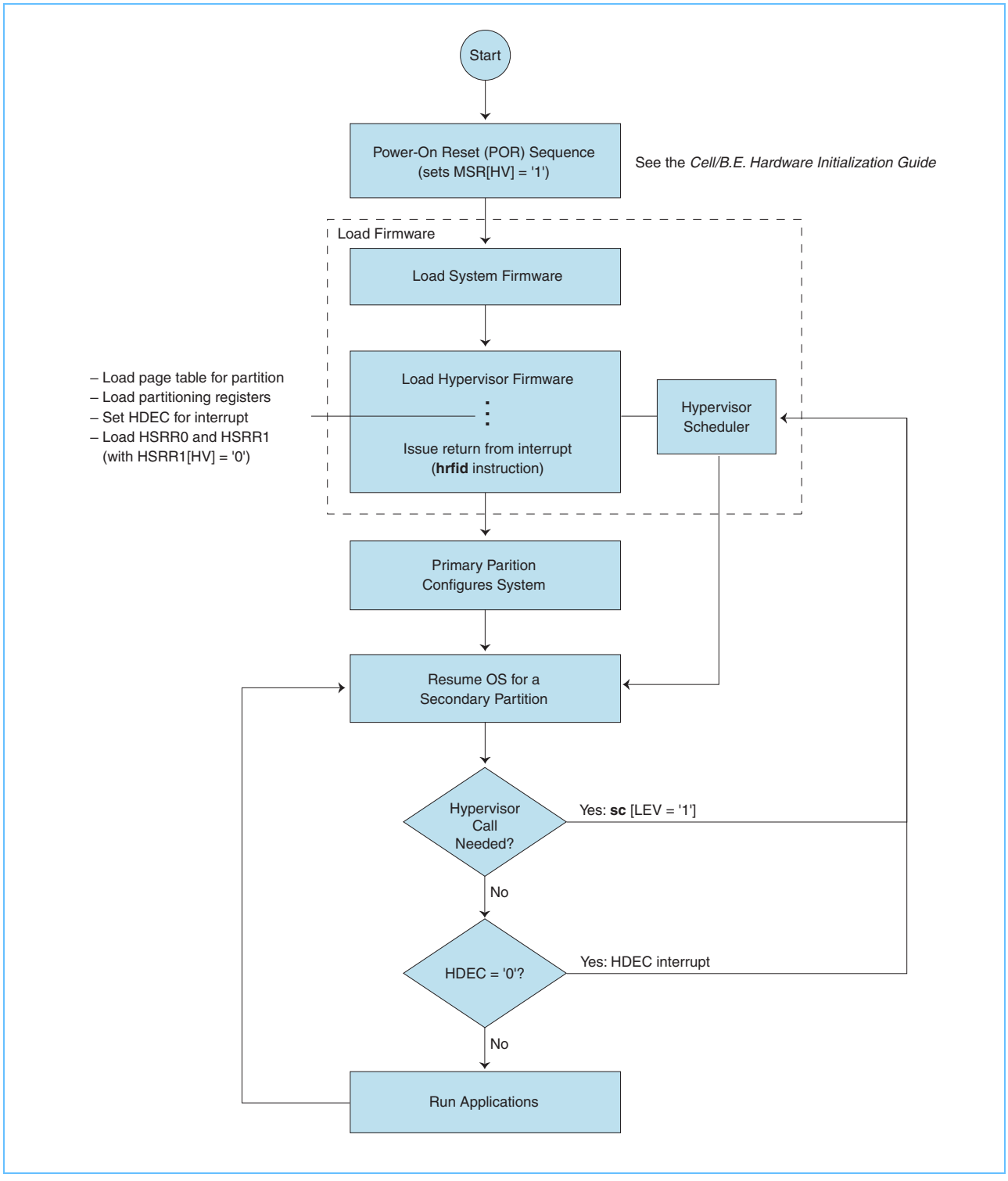
### 11.1.3 An Example Flowchart

*Figure 11-2* on page 335 shows how a hypervisor might boot and manage logical partitions. The events shown here cover only the most basic workings of a hypervisor, the primary partition (also called an alpha or I/O-hosting partition), and a secondary partition and its applications. Other methods of implementing a hypervisor are possible.

The system reset interrupt can be used for power-on reset (POR). This interrupt sets both the Hypervisor bit, `MSR[HV]`, and the 64-bit mode bit, `MSR[SF]`, to '1', so that the CBEA processor comes out of POR in hypervisor mode. After loading system firmware, the CBEA processor loads the hypervisor firmware. This in turn loads the hardware-managed page table for the first partition (the hypervisor maintains such a page table for each partition), the registers used for hypervisor control of the partition, the HDEC decrementer, and the `HSSR0` and `HSSR1` registers that contain the program counter and Machine State Register (`MSR`) values for the first partition. The hypervisor then executes a return from interrupt (**hrfid**) to begin execution of the first partition.

The primary partition is the first partition created by the hypervisor. This partition has a small operating system (OS) and its mission is to discover the resources in the system (memory, I/O, and so forth), and manage the creation of secondary partitions for the hypervisor. After the primary partition has configured the system, it passes control to the OS of a secondary partition. This partition's OS runs its applications until it either needs help from the hypervisor (such as an address translation) or until its time-share slice is finished (signaled by the HDEC count of '0'). It obtains help from the hypervisor by executing a *hypervisor call* (hcall), which is the **sc** instruction with the LEV bit = '1'. When either of these events occurs, the hypervisor loads the necessary information (such as address-translation information) and returns control to the partition, again by executing a return from interrupt.

*Figure 11-2. Flowchart of Hypervisor Management of Logical Partitions*



## 11.2 PPE Logical-Partitioning Facilities

The PPE's standard PowerPC-based logical-partitioning architecture is defined in the *PowerPC Operating Environment Architecture, Book III*. The PPE can be assigned to one partition at a given time. If active, both PPE threads execute in the same logical partition.

### 11.2.1 Enabling Hypervisor State

The PowerPC Architecture supports three privilege states, which are controlled by MSR[HV] and MSR[PR] bits. Memory-mapped I/O (MMIO) registers are classified according to these states. The states are designed for the following uses:

- *Hypervisor State*—MSR[HV] = '1' and MSR[PR] = '0'. The most-trusted state, used by the hypervisor. Access to all system facilities is provided at this level of privilege. In the *PowerPC Architecture*, this state is occasionally referred to as the "privileged and hypervisor state"; for convenience, this document always refers to the state as "hypervisor state".
- *Privileged State*—MSR[HV] = '0' and MSR[PR] = '0'. Used by the operating system within a logical partition. In the *PowerPC Architecture*, this state is occasionally referred to as the "privileged and nonhypervisor state" or "supervisor state"; for convenience, this document always refers to the state as "privileged state".
- *Problem State*—MSR[HV] = '0' and MSR[PR] = '1', or MSR[HV] = '1' and MSR[PR] = '1'. The least-trusted state, used by application software in a logical partition.

The MSR[PR] bit can be set using the **mtmsr** instruction. Privileged or hypervisor state is entered by executing the **sc** (system call) instruction or as a result of some interrupts. The **sc** instruction has a LEV bit which, when set to '1', causes MSR[HV] to be set to '1'. Exiting the privileged or hypervisor states is accomplished by executing the **rfid** or **hrfid** instruction. For details, see *Table 9-40* on page 294 and the *PowerPC User Instruction Set Architecture, Book I* and *PowerPC Operating Environment Architecture, Book III*.

If no hypervisor exists on the CBEA processor, the entire system is normally run at MSR[HV] = '1', and only two privilege states exist, PR = '0' for firmware and operating-system privileged state, and PR = '1' for application-software problem state.

If address-translation is enabled, privileged software can control, by means of page-table entries, whether application software is given access to particular problem-state MMIO registers; access to the MMIO registers in this mode is not directly enforced by hardware.

### 11.2.2 Hypervisor-State Registers

The following registers are used by the hypervisor to configure operations in logical partitions. Most of these registers can only be accessed in hypervisor state:

- *Logical-Partition Identification Register (LPIDR)*—Specifies the logical partition to which the PPE (both threads) is currently assigned. Up to 32 partitions can be used.
- *Processor Identification Register (PIR)*—Specifies the system-wide processor ID for each PPE thread.
- *Hypervisor Status Save and Restore Registers (HSRR0, HSRR1)*—Specifies the state of the PPE when an interrupt occurs or the state from which the PPE should resume after a hypervisor return from interrupt (**hrfid**).

- *Storage Description Register 1 (SDR1)*—Specifies the real address and size of the hardware-accessed page table for the current partition. The hypervisor provides each partition with its own hardware page table.
- *Real-Mode Offset Register (RMOR)*—Specifies the offset from real address x'0' at which real-mode memory begins for nonhypervisor real-mode accesses. See *Figure 11-3* on page 340 for how an OS and the hypervisor view this memory space.
- *Logical-Partitioning Control Register (LPCR)*—Specifies the largest effective address that can be used by a logical partition in real mode, how storage is accessed and cached in real mode, and certain effects of interrupts. See *Figure 11-3* on page 340 for how an OS and the hypervisor view the real-mode address limit.
- *Hypervisor Real-Mode Offset Register (HRMOR)*—Specifies the offset from real address x'0' at which real-mode memory begins for hypervisor real-mode accesses.
- *Hypervisor Decrementer (HDEC)*—Provides a means for the hypervisor to manage timing functions independently of the Decrementer (DEC) registers, which are managed by user or supervisor software running in a partition. Like the DECs, the HDEC provides a means of signaling an interrupt after a specified amount of time has elapsed. This is useful, for example, when time-sharing resources.
- *Machine-Check Enable Bit (MSR[ME])*—Enables machine check interrupts.
- *Time Base (TB) Register*—Specifies the value of the time base.
- *Data-Address Breakpoint Register and Data-Address Breakpoint Extension Register (DABR, DABRX)*—Specifies the effective address and other conditions of a load or store access at which a data storage exception should occur.
- *Hypervisor Software-Use Special-Purpose Registers (HSPRG0, HSPRG1)*—Available for any use by the hypervisor.
- *Hardware-Implementation-Dependent (HID) Registers*—Specify a variety of PPE-state and debug functions.

See the *Cell Broadband Engine Registers* specification for details about these registers. See the *PowerPC Operating Environment Architecture, Book III* and *Section 20 Shared-Storage Synchronization* on page 561 for the required synchronization instructions and sequences that must be used by the hypervisor with these facilities. Failure to properly synchronize these changes can lead to inconsistent state and failure of the application or hypervisor.

### 11.2.3 Controlling Real Memory

When supporting logical partitioning, the hypervisor must itself use some real memory and also provide some real memory to each logical partition. The hypervisor must control access to real memory so that no logical partition can access another partition's real memory or the hypervisor's real memory.

In the Instruction Relocate mode (MSR[IR]) and Data Relocate mode (MSR[DR])—described in the *Cell Broadband Engine Registers* specification and the *PowerPC Operating Environment Architecture, Book III*—this access to real memory is controlled by the hypervisor, which is responsible for the hardware-accessed page table and translation lookaside buffer (TLB) for each partition. This hardware-accessed page table is described by the contents of the Storage Description Register 1 (SDR1). The SDR1, the TLB facilities, and the real memory reserved by the hypervisor for the hardware page table must be accessible only to the hypervisor.

The operating system in a logical partition normally manages the virtual-address (VA) space for the logical partition and the mapping between effective-address (EA) space and VA space through direct manipulation of its segment lookaside buffer (SLB). However, the relationship between the VA space and the real-address (RA) space is managed by the operating system indirectly through hypervisor calls.

The hypervisor typically uses its own real-memory allocation table to manage which blocks of real memory have been allocated to a logical partition. As shown in *Figure 11-3* on page 340, the hypervisor allocates to a partition the real memory that the partition software uses as address 0, and this memory appears (to the partition software) contiguous up to the total amount of memory allocated to the partition.

When managing its own real memory, the operating system in a logical partition typically makes a hypervisor call to request a page-table update for its VA-to-RA mapping. The hypervisor then checks the requested real-page mapping, and if valid for that logical partition, it converts the RA supplied by the operating system to the true RA and manages the required changes to the hardware page table and TLBs.

The PPE and the Synergistic Processor Element (SPE) MMUs support three simultaneous pages sizes—4 KB, and two sizes selectable from 64 KB, 1 MB, or 16 MB. Large page sizes are very useful in reducing thrashing of TLBs when applications access large amounts of main storage. The PPE provides 1,024 TLB entries and each SPE provides 256 TLB entries.

#### 11.2.3.1 *Memory-Management Facilities*

The hypervisor typically executes in real mode (`MSR[IR]` = ‘0’ for instructions, `MSR[DR]` = ‘0’ for data). In this mode, the real-mode (nonrelocate) translation facility is employed. However, some hypervisor implementations might choose to have components that operate with instruction relocation (`MSR[IR]` = ‘1’) or data relocation (`MSR[DR]` = ‘1’) or both enabled. When relocation is enabled, the hypervisor can choose to have VA-to-RA translations supported by the hardware page table by setting `LPCR[TL]` = ‘0’, or it can specify the software-managed TLB mode by setting `LPCR[TL]` = ‘1’. This can be varied on a logical-partition basis.

*Software-Managed TLB Mode*

In software-managed TLB mode, software directly manages the TLBs and the hardware memory management unit (MMU) never directly accesses a page table in main storage to reload a TLB entry. An instruction storage or data storage interrupt is generated immediately following any TLB miss by the VA-to-RA translation hardware (see *Section 11.2.4 Controlling Interrupts and Environment* on page 343 and *Section 11.3.3 Controlling Interrupts* on page 349).

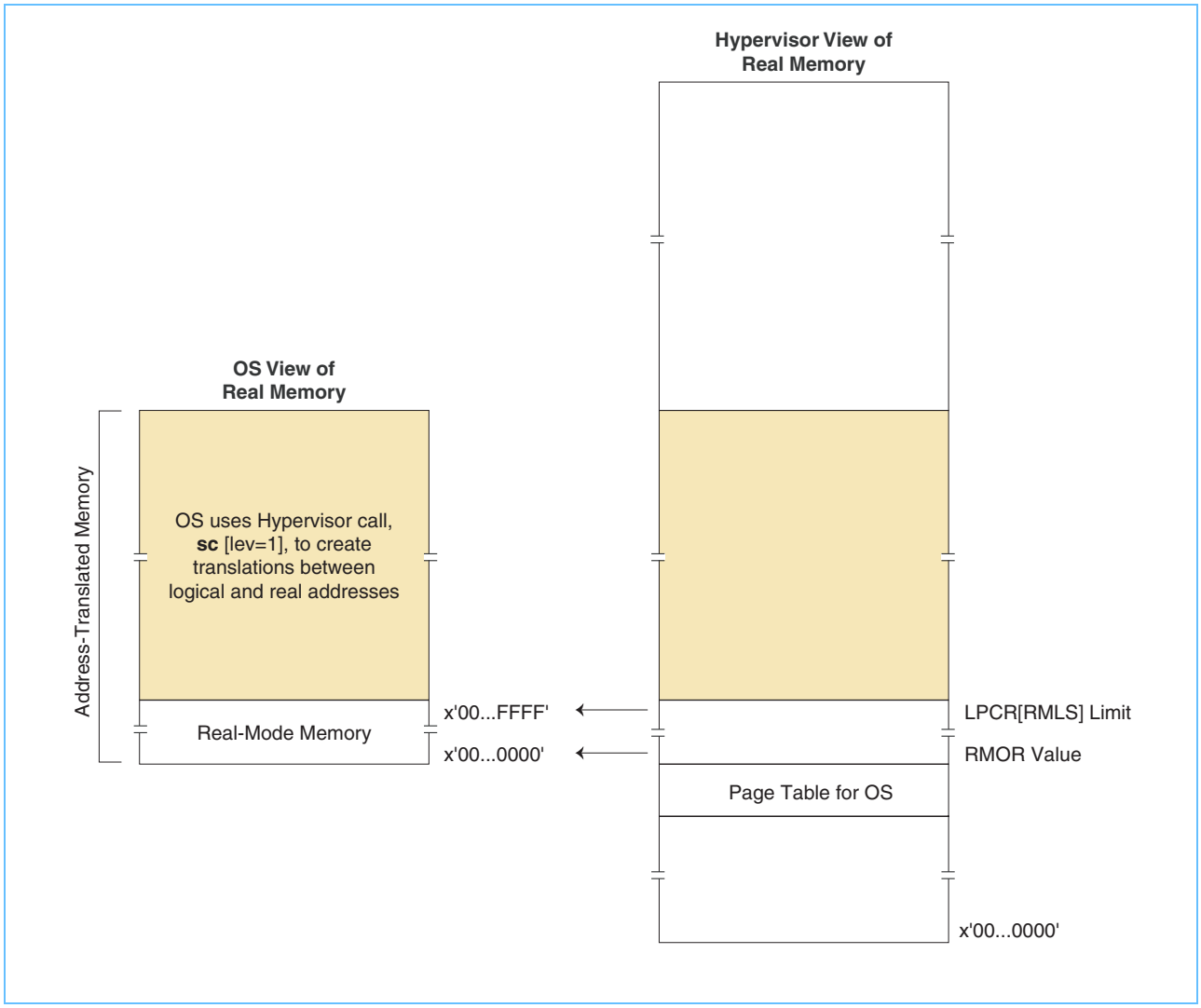
Hypervisor software uses several registers to directly manage the TLBs, including the PPE TLB Index Hint Register (`PPE_TLB_Index_Hint`), PPE TLB Index Register (`PPE_TLB_Index`), PPE TLB Virtual-Page Number Register (`PPE_TLB_VPN`), and PPE TLB Real Page Number Register (`PPE_TLB_RPN`). In addition to using these registers to reload a TLB entry after a miss, software is also able to preload translations directly into the TLB with these registers for improved real-time response.

The **tlbie(l)** instruction must be used to invalidate the contents of a TLB entry in both hardware-managed and software-managed TLB modes. The hypervisor also uses Hardware Implementation-Dependent Register 6 (HID6) to select the sizes of up to two concurrent large page sizes, if large-page translations are to be supported by the hypervisor or provided to one or more partitions.

*Access by Operating Systems to Physical Memory in Real Mode*

Many existing operating systems require some access to physical memory in real mode, typically an area in low memory for interrupt handling and typically starting at real address x‘0’. Unrestricted access to physical memory by a logical partition operating in real mode cannot be allowed. This is handled through the use of the real addressing mode (real mode) facility, which includes (among other resources) the Real Mode Offset Register (RMOR) and the Real Mode Limit Selector field in the Logical-Partitioning Control Register (LPCR[RMLS]) when MSR[HV] = ‘0’ and LPCR[LPES] bit 1 = ‘1’. Using these facilities, the hypervisor can provide a logical partition with access to a limited amount of physical memory while accessing storage in real mode, as shown in *Figure 11-3* on page 340.

*Figure 11-3. Operating-System and Hypervisor Views of Physical Memory*




*Access by Hypervisor to Physical Memory in Real Mode*

A similar facility exists, using the Hypervisor Real Mode Offset Register (HRMOR), for the hypervisor itself to provide a real-mode offset when accessing storage in real mode with MSR[HV] = '1'. This can be useful in certain system configurations—such as multiprocessor nonuniform memory access (NUMA) environments—in which a copy of hypervisor code or data should be replicated for performance or reliability considerations. The HRMOR value is only effective when bit 0 (the most-significant bit) of the EA is '0'. If bit 0 of the effective address is a '1', bits 2 through 63 of the EA are used as the RA for the access.

*Real-Mode Limit Select*

The Real Mode Limit Selector field in the Logical-Partitioning Control Register, LPCR[RMLS], controls the size of the partition's real-mode offset. The Logical Partition Environment Selector field, LPCR[LPES] (see *Section 11.2.4.1* on page 343), controls how external interrupts and a subset of internal interrupts are handled.

*Table 11-1* summarizes the supported Real Mode Limit Select (RMLS) values and the corresponding limits chosen. The CBEA processors support an additional high-order bit of the RMLS field not listed in *PowerPC Operating Environment Architecture, Book III*, for a total of four bits rather than three. When in real offset mode and the requested effective address is beyond the chosen RMLS limit, an instruction or data storage interrupt occurs, and SRR1[36] is set to '1'.

*Table 11-1. Summary of Real Mode Limit Select (RMLS) Values*

| LPCR[RMLS] Field | Effective Address Limit |
|---|---|
| 0000 | 256 GB |
| 0001 | 16 GB |
| 0010 | 1 GB |
| 0011 | 64 MB |
| 0100 | 256 MB |
| 0101 | 16 MB |
| 0110 | 8 MB |
| 0111 | 128 MB |
| 1000 | Reserved[1] |
| 1001 | 4 MB |
| 1010 | 2 MB |
| 1011 | 1 MB |
| 11xx | Reserved[1] |

1. All reserved values are treated as 1 MB.

*Real-Mode Caching*

The Real Mode Caching-Inhibited bit in the LPCR register (LPCR[RMI]) controls whether a real-mode data access is treated as cacheable or caching-inhibited. When the real mode storage control facility (HID6[RMSC]) is used, the LPCR[RMI] bit is only used for real-mode accesses with addresses above the address specified in the HID6[RMSC] field. This control is used to avoid a cache paradox (the inadvertent caching of noncacheable data) when accessing data storage in real mode.

*Logical Partition Identification Register*

The PPE has a Logical Partition Identification Register (LPIDR) that the hypervisor writes to assign the PPE to a partition. In a multi-CBEA-processor system, or when the PPE is time-shared or virtualized so that it can serially support multiple partitions (fractional assignment of the PPE), the hypervisor updates the LPIDR register with the identifier of the logical partition to which the PPE has been assigned. This value is used by the hardware to avoid unnecessary removal of TLB entries.

The LPIDR register is used by the hardware under three circumstances:

- To tag TLBs with the current logical partition ID (LPID) value when the TLB entry for a page translation is loaded by the hardware.
- To invalidate only TLB entries that match the criteria of the **tlbie(l)** instruction when the current LPIDR contents match the LPID tag value of the matching entries in the TLB. This prevents a previous partition's TLB entries from being discarded from the TLB unnecessarily by the partition currently using the PPE. This also prevents a broadcast **tlbie** instruction issued on a processor element from unnecessarily invalidating TLB entries on another processor element that is operating in a different partition.
- To prevent the hardware from using a TLB entry for translation when the current contents of the LPIDR does not match the LPID tag of the TLB entry. This prevents a partition from using a different partition's address translation. It also removes the requirement for the hypervisor to invalidate all TLB entries when switching the PPE between partitions. However, the LPID tag in the TLB entry does not affect the hardware's TLB replacement algorithm.

#### 11.2.3.2 *Real-Time Memory Management*

In addition to using standard cache-management and TLB-management instructions, the hypervisor can also control the replacement policy for cache and TLB sets on behalf of itself or at the request of the operating system in a logical partition. Hardware uses a least-recently used (LRU) algorithm to determine which cache line or TLB entry to replace. In many real-time applications, such replacement can result in nondeterministic behavior or poor performance, due to application-specific access patterns. One solution is to use the software-managed TLB mode, described in *Section 11.2.3.1 Memory-Management Facilities* on page 338. Another is the software locking of TLB entries, as described in *Section 4.2.7.7 TLB Replacement Management Table* on page 99.

*Data-Address and Instruction-Address Range Registers*

The hypervisor has access to the PPE's Address Range Registers, which support the pseudo-LRU replacement algorithm for L2 and TLB replacement. An address range is a naturally-aligned range that is a power-of-2 size between 4 KB and 4 GB, inclusive.

Each hardware thread has two sets of Data-Address Range Start Registers (DRSR0 and DRSR1), Data-Address Range Mask Registers (DRMR0 and DRMR1), and associated Class ID registers. Each hardware thread also has two sets of Instruction Address Range Start Registers (IRSR0 and IRSR1), Instruction Address Range Mask Registers (IRMR0 and IRMR1), and associated Class ID Registers.

These registers define four windows of EA space (two for data accesses and two for instruction fetches) in which the accesses can be associated with a replacement-class identifier. This identifier is then used to index into one of two replacement management tables (RMTs)—an L2 replacement management table or a TLB replacement management table—to control which sets within the L2 cache or TLB can be used by the access. This provides software with the ability to define an EA region in which data is continuously reused and should be retained or locked in the L2 cache and TLBs. Software can also define an EA region in which data is streamed with little to no reuse, such that only a small area of the L2 cache and TLB is used and the rest of the L2 cache and TLB entries are not replaced by reloads due to this streaming data.

The Address Range Start Registers also provide a real-mode or relocation-mode bit to specify whether the address match is a real address or a translated address. Storage accesses that do not match any range are assigned to replacement class 0. Eight replacement classes are supported (class 0 through 7).

For further details about the Address Range Registers, see *Section 6.3.1.3 Address Range Registers* on page 155.

*Replacement Management Tables*

The hypervisor assigns L2-cache sets or TLB sets or both to each class ID using the replacement management tables, described in *Section 6.3 Replacement Management Tables* on page 154. The PPE MMU supports an 8-entry RMT for the L2 cache and an 8-entry RMT for the TLB. A 3-bit class ID is used to index into the table to select one of the RMT entries. Each entry of the RMT contains a replacement-enable bit for each way of the cache or TLB. For the L2 cache, the RMT is set up and controlled through the `L2_RMT_Data` register and the `L2_ModeSetup1[RMT_mode]` bit. For the TLB, the RMT is setup and controlled through the `PPE_TLB_RMT` register. The `PPE_TLB_RMT` register is only effective when the hardware-managed TLB mode is enabled (`LPCR[TL]` = '0').

## 11.2.4 Controlling Interrupts and Environment

### 11.2.4.1 *Logical-Partition Environment Selector*

One of the bits in the Logical-Partition Environment Selector field, `LPCR[LPES]`, controls whether external interrupts force `MSR[HV]` to '1' or leave it unmodified. This control is used by the hypervisor to determine whether external interrupts are always delivered to the hypervisor or directly to the software being interrupted (the hypervisor or the operating system in the logical partition). Another `LPCR[LPES]` control bit is used to determine whether the real mode offset facility is to be used by the logical partition, and to direct some of the interrupts to the hypervisor instead of the software being interrupted—one bit for both purposes[2]. The system reset, machine check, system error, and hypervisor decrementer (`HDEC`) interrupts are always presented to the hypervisor.

2. See *PowerPC Operating Environment Architecture, Book III* for details on the LPCR register.

#### 11.2.4.2 ***External Interrupts***

*External-Interrupt Handling*

The hypervisor is typically responsible for managing external interrupts (see *Table 9-1 PPE Interrupts* on page 242). This can be done in one of two ways:

- The logical partition supplies the first-level interrupt handler (FLIH) and the interrupt vector in the real-mode offset (RMO) region.
- The hypervisor supplies the FLIH and interrupt vector in the low real-address range or the hypervisor real-mode offset (HRMO) region.

In the first case, although the logical partition provides the FLIH, it typically will make hypervisor calls to determine the interrupting condition, mask interrupts, generate interrupts, and acknowledge interrupts. In either case, the hypervisor must retain direct control over the interrupt masks, interrupt status, interrupt routes, interrupt-generation registers, interrupt priorities, and interrupt-pending registers of the two internal interrupt controllers (IICs), one for each PPE hardware thread, or of any external interrupt controllers attached to the I/O interfaces (IOIFs). All external interrupts ultimately must be presented to one of the IICs. The IICs, in turn, provide the interrupt to the PPE thread.

*Real-Time Processing of External Interrupts*

In real-time applications, tasks must be completed within a specific period of time. In particular, interrupt processing for high-priority interrupts must be deterministic and worst-case latency of interrupt handling must be managed accordingly. In a logically partitioned system, the PPE might be currently assigned to one partition when a high-priority interrupt occurs that must be serviced by an operating system in another partition. Hypervisors traditionally defer external-interrupt processing until the PPE is dispatched on the partition required to handle the external interrupt. This can result in extremely long interrupt-processing delays, which are unacceptable to real time applications.

Setting the `LPCR[LPES]` bit 0 = ‘0’ allows external interrupts to be directly handled by the hypervisor, instead of by the operating system in the partition currently executing. However, the executing partition retains control over the `MSR[EE]` which enables or disables external interrupts, with the potential to significantly delay the handling of a high-priority external interrupt.

*Mediated External Interrupts*

To solve the potential problem described in the preceding paragraph, the PPE supports a mediated external exception mode (see *Section 9.5.7* on page 254 for details). This mode is enabled by a hardware implementation register (`HID0[extr_hsrr]` = ‘1’). In this mode, with `LPCR[LPES]` bit 0 = ‘0’ to direct external interrupts to the hypervisor, the `MSR[EE]` state does not inhibit external interrupt delivery while the PPE is executing in nonhypervisor state. Instead of saving the interrupting PPE state in the `SRR0` and `SRR1` registers, the interrupting state is saved in `HSRR0` and `HSRR1` to avoid loss of context. In addition, the hypervisor uses the `HSPRG` registers for scratch registers instead of the `SPRG` registers. This provides the capability for the hypervisor to respond to the reception of a high-priority external interrupt.

The hypervisor can preemptively swap the targeted PPE context to the partition that must service the interrupt. The hypervisor then generates a mediated external-exception request by setting the Mediated External-Exception Request (LPCR[MER]) bit to '1'. This mediated external exception condition honors the setting of the MSR[EE] state for the logical partition, and it will not provide a mediated external interrupt to the partition until MSR[EE] = '1'. The occurrence of a mediated external interrupt saves the processor-element state in the SRR0 and SRR1 registers. A mediated external-exception request is outstanding when LPCR[MER] = '1', however the exception will never interrupt the processor element while in hypervisor mode (MER[HV] = '1' and MER[PR] = '0').

The mediated external interrupt facility is also available when enabled by HID0[extr_hsrr] = '1' and LPCR[LPES] bit 0 = '1'. In this mode, direct external exceptions are delivered as external interrupts directly to the currently executing environment (hypervisor or the logical-partition software). However the interrupted processor-element state is saved in SRR0 and SRR1, and the interrupt handler uses the SPRG registers as scratch registers. The setting of MSR[EE] enables or disables the external interrupt delivery in both hypervisor and the logical partition. The hypervisor can generate a mediated external exception and associated external interrupt by using the LPCR[MER] bit. This is particularly useful, for example, when the hypervisor wants to request the addition or removal of a resource for dynamic provisioning.

The LPCR[LPES] setting applies to both PPE threads, whereas the MSR[EE] and the LPCR[MER] are per thread.

For further details about interrupt handling, see *Section 9 PPE Interrupts* on page 239.

#### 11.2.4.3 *Hypervisor Decrementer Interrupts*

To provide periodic maintenance tasks or timeshare resources among partitions, the hypervisor can use the hypervisor decrementer (HDEC). This decrementer will interrupt the processor element regardless of the MSR[EE] setting when MSR[HV] = '0'. An HDEC interrupt will not be delivered when MSR[EE]= '0' in hypervisor state. In addition, the LPCR register provides a Hypervisor Decrementer Interrupt Conditionally Enable bit (LPCR[HDICE]) to disable the hypervisor decrementer in all states.

#### 11.2.4.4 *Machine Check Interrupts*

Only the hypervisor has the capability to enable or disable machine check interrupts by altering the MSR[ME] bit in the Machine State Register. This bit is used to control whether a subsequent machine check event causes a machine check interrupt or a system wide checkstop.

#### 11.2.4.5 *Time Base*

The Time Base Register (TB) is readable in all privilege states. However, because the TB is used to manage system-wide time, it is only writable in the hypervisor state.

#### 11.2.4.6 *Processor Identification*

The Processor Identification Register (PIR) is used by software to distinguish one PPE from another in a multi-CBEA-processor system. It is typically used to identify communicators during communication between processor-elements and I/O devices. The register can be written only by the power-on reset configuration bits, after which it is read-only. (For details about configuration bits, see the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.)

#### 11.2.4.7 *Data Address Breakpoints*

The Data Address Breakpoint Register (DABR) and Data Address Breakpoint Register Extended (DABRX) are used for system-level debugging and can only be accessed by the hypervisor. This prevents unauthorized breakpoints being set in hypervisor data.

## 11.3 SPE Logical-Partitioning Facilities

The hypervisor can allocate SPEs among logical partitions. The hypervisor can also dynamically reassign an SPE from one partition to another. The SPEs can operate in a partition which does not—due to PPE-sharing—currently have a physical PPE assigned to it, until a PPE interrupt is generated by an SPE.

SPE software can only fetch instructions and load and store data directly from its local storage (LS) using an LS address (LSA). SPE software has access to a limited set of communications and DMA facilities directly through the use of synergistic processor unit (SPU) channel instructions. Software in the SPU itself has no direct access to main storage. It must use the memory flow controller (MFC) DMA facilities to move data between the LS and the effective-address (EA) space. Each MFC contains an MMU, very similar to the MMU in the PPE, for EA-to-RA translation.

### 11.3.1 Access Privilege

Accesses to an SPE from all other system units (PPE, other SPEs, and I/O devices) are provided through MMIO registers. Therefore, with respect to logical partitions executing on the PPE, an SPE can be allocated to the logical partition by the hypervisor granting access to the SPE MMIO registers. These registers are divided into three groups, based on privilege, as specified in the *Cell Broadband Engine Architecture* document:

- *Privilege 1 Registers* (most privileged)—Used by the hypervisor to manage the SPE on behalf of a logical partition.
- *Privilege 2 Registers*—Used by the operating system in a logical partition to manage the SPE within the partition.
- *Privilege 3 Registers* (least privileged)—Used by problem-state (application) software, if direct access to the SPE from user space is supported by the operating system.

A hypervisor should not grant logical-partition access to the privileged 1 group of SPE MMIO registers. These registers provide control over SPE memory management, interrupt management, versioning, SPE cache management, and SPE status and control.

### 11.3.2 Memory-Management Facilities

As with storage access by the PPE, the hypervisor must control storage access by an SPE allocated to a logical partition. The following hypervisor (privilege 1) facilities are replicated for each SPE to provide this control:

- *MFC Logical Partition ID Register (*`MFC_LPID`*)*—Identifies the partition to which an SPE is assigned.
- *MFC State Register 1 (*`MFC_SR1`*)*—Controls MFC access to storage.
- *MFC TLB Management Registers (*`MFC_SDR`*,* `MFC_TLB_Index_Hint`*,* `MFC_TLB_Index`*,* `MFC_TLB_VPN`*,* `MFC_TLB_RPN`*, and* `MFC_TLB_Invalidate_Entry`*)*—Specify page-table properties and state.
- *MFC TLB Replacement Management Table Data Register (*`MFC_TLB_RMT_Data`*)*—Enables up to four replacement management tables (RMTs).

The contents of the `MFC_SR1` register controls the following functions:

- Software-managed or hardware-managed MFC TLB reload on translation miss
- SPU master run control
- Relocation (translation) of MFC EAs to RAs
- Problem-state or privileged-state access by MFC
- Ignore or honor PPE-issued **tlbie** instructions in the same logical partition
- Aliasing of LS space into the main-storage EA space

As in logical partitioning for the PPE, the hypervisor must control an MFC DMA controller's access to physical storage through the use of the hardware page table or TLBs, as described in *Section 11.2.3* on page 337. Unlike the PPE, the MFCs do not support a real mode offset facility, because there is no operating system executing on an SPE and there is no requirement to access a low-memory area in real mode.

The hypervisor should not disable MFC Address Relocation for an SPE assigned to a logical partition, because this provides uncontrolled access to all of the EA space. Relocation should only be disabled on SPEs that are allocated for hypervisor use, requiring real-mode access.

#### 11.3.2.1 *Page Tables*

The hypervisor can use the `MFC_SR1` to specify that TLB (VA-to-RA) reloads should be handled either by a hardware lookup in a hashed page table in main storage or by direct software load of a TLB entry. When the hardware TLB reload is used, the MFC Storage Description Register (`MFC_SDR`) must be set by the hypervisor to describe the hardware-accessible page table. This page table has the same format as the PPE page table. Hardware page tables can be common for all PPE and SPE resources in a logical partition, or they can be separate. When a common hardware page table is used, the `MFC_SR1` register should be set so that PPE-issued **tlbie** instructions are honored by all SPEs in the logical partition. In addition, the `MFC_LPID` register must be set by the hypervisor to the LPID for the logical partition that the SPE is assigned to. The SPE will ignore PPE-issued **tlbie** instructions if the PPE's LPID does not match the MFC LPID.

If per-SPE page tables are used, the `MFC_SR1` register should be set so that PPE-issued **tlbie** instructions are ignored by the SPEs in the logical partition.

#### 11.3.2.2 *MFC TLB Management Registers*

The hypervisor can use the software TLB-reload mode on one or more of the SPEs by setting `MFC_SR1[TL]` = '1'. This can be useful in small-memory configurations to avoid the need for a large hardware-accessible page table in main storage. In this mode, when a page-translation miss occurs, the PPE is interrupted to load the TLB with the necessary translation and resume the operation. The TLB-management registers in the privilege 1 area are used for this purpose.

When a miss occurs, the MFC TLB Index Hint Register (`MFC_TLB_Index_Hint`) specifies which TLB congruency class needs to be loaded to resolve the miss. An LRU hint is also available in this register, indicating which TLB entry in the class is the least-recently used. The MFC TLB Invalidate Entry Register (`MFC_TLB_Invalidate_Entry`) must be used to invalidate the TLB in the SPE before reloading a new one[3]. The MFC TLB Index Register (`MFC_TLB_Index`) is used to specify which TLB entry is read or written by accesses to the MFC TLB VPN (`MFC_TLB_VPN`) and MFC TLB RPN (`MFC_TLB_RPN`) Registers.

#### 11.3.2.3 *TLB Entries*

Unlike the PPE TLB entries, the SPE TLB entries are not tagged with a logical partition ID[4], because the frequency of switching an SPE between logical partitions should be much lower than switching the PPE between logical partitions. Switching an SPE between logical partitions requires that all entries in the SPE TLB be invalidated when the SPE is switched from one partition to another. This is required because each partition manages its own VA space, therefore page translations are unique to the partition.

Like the PPE, in the hardware TLB-reload mode (`MFC_SR1[TL]` = '0') the MMU hardware uses a least-recently used (LRU algorithm) to determine which entry of the 4-way set-associate TLB will be replaced during a TLB reload. For finer control of this facility, the MFC MMU also supports a TLB replacement management table (RMT), similar to that in the PPE MMU. This table is set up by the hypervisor for each SPE, using the MFC TLB Replacement Management Table Data Register (`MFC_TLB_RMT_Data`). This register has four fields, one for each replacement-class ID (RClassID), with each field containing four set-enable bits to indicate whether the set is to be used for the corresponding class. The RClassID for DMA commands is provided by the application program when queuing a DMA command.

#### 11.3.2.4 *Access to Privileged Storage*

When assigning the SPE to a logical partition, the hypervisor can also specify whether the SPE is to have access to privileged storage as well as problem-state storage. This is specified by a the `MFC_SR1[PR]` bit. Storage-access privilege is specified by the storage key in the segment look aside buffers (SLBs) that the operating system controls in the logical partition, in conjunction with the storage-protection bits in the hardware page table or TLB entries that the operating system specifies and the hypervisor controls.

3. See *Section 4.3.5.3* on page 112 for details.
4. SPE TLB entries are invalidated by writing the MFC_TLB_Invalidate_Entry register, as opposed to the PPE method of issuing a **tlbie**, **tlbiel**, or **tlbia** instruction to invalidate PPE TLB entries.

#### 11.3.2.5 *Aliasing LS to Main Storage*

To facilitate LS-to-LS DMA transfers, or the capability of an I/O device to access LS, the operating system can request the hypervisor to alias the allocated SPE's LS into the real-address space available to the logical partition or to the hypervisor. The enabling and disabling of LS aliasing is controlled by the `MFC_SR1[D]` bit.

### 11.3.3 Controlling Interrupts

The hypervisor has exclusive access to the SPE interrupt routing, interrupt mask and interrupt status registers, described in *Section 9.6.3 SPU and MFC Interrupts* on page 271. These registers are replicated for each SPE, and within each SPE there is one for each class of interrupt (error, translation, and application):

- Interrupt Routing Register (`INT_Route`)
- Interrupt Mask Registers (`INT_Mask_class0`, `INT_Mask_class1`, and `INT_Mask_class2`)
- Interrupt Status Registers (`INT_Stat_class0`, `INT_Stat_class1`, and `INT_Stat_class2`)
- SPU Identification Register (`SPU_IDR`)

By setting the unit ID in the `INT_Route` register, SPE-initiated interrupts can be routed to a specific PPE thread or to an external interrupt controller attached to the IOIF0 or IOIF1 interface. This register allows each of the three classes of SPE-initiated interrupts to be routed to the same or different units. This is useful if one unit is to handle error interrupts while a different unit is to handle translation or application interrupts.

As described in *Section 11.2.4* on page 343, when SPE interrupts (which are seen as external interrupts) are routed to a PPE thread, the hypervisor can handle the interrupts directly or it can allow the operating system in the currently executing partition to handle the interrupt. In either case, the hypervisor should provide a set of SPE interrupt services to the logical partition. These services should include masking an allocated SPE interrupt, querying an SPE's interrupt status, and resetting an SPE interrupting condition.

### 11.3.4 Other SPE Management Facilities

The following additional facilities provide the hypervisor with per-SPE information and control:

- MFC and SPU Version Registers
- MFC Cache Management Registers
- MFC and SPU Error Management Registers
- MFC and SPU Debug and Performance Monitor Registers

#### 11.3.4.1 *SPU Run Control*

The hypervisor can inhibit the operating system or an application in the logical partition from starting an SPU that has been allocated to it. This control is provided by the SPU Master Run Control bit, `MFC_SR1[S]`. It can also be used to stop a currently executing SPU that has been allocated. This can be useful in certain high-thermal conditions, or can be used to keep the SPU resources allocated to a logical partition for the purposes of retrieving the SPU context, but preventing additional tasks being started on the SPU. If an SPU was stopped with the `MFC_SR1[S]`

bit, the SPU will not resume execution until the bit has been enabled and a run request to the SPU Run Control Register (SPU_RunCntl) has been initiated by the hypervisor, operating system, or application program.

#### 11.3.4.2 *SPE Identification*

The SPU Identification Register (SPU_IDR) is set by hardware to a unique value for each physical SPE and can be used by the hypervisor to index per-SPE structures. The MFC and SPU Version Registers (MFC_VR and SPU_VR) are also set by hardware and are used by the hypervisor to manage the features and functions available for a given version. It is suggested that the hypervisor propagate this information to the software in the logical partitions as configuration information, as required.

#### 11.3.4.3 *Cache Management*

Each SPE contains a small cache to support TLB loads from the hardware page table and the caching of atomic-update lock lines. The atomic cache can be flushed by the hypervisor using the MFC Atomic Flush Register (MFC_Atomic_Flush). This should always be done before putting an unallocated SPE into one of the low-power states (see *Section 15.1 Power Management* on page 429) because the contents of the cache can become stale or lost in these states. The cache uses real-address tags, so it is not necessary to flush the cache when switching the SPE between logical partitions.

#### 11.3.4.4 *Error Management*

To support error detection, the hypervisor can use the SMM Hardware Implementation Dependent Register (SMM_HID) to enable SLB and TLB parity generation and checking, and to control the MFC performance, trace, and trigger settings. In addition, SPU illegal instruction and LS error-correcting code (ECC) error detection and correction functions can be enabled and controlled through the SPU ECC and Error Mask Registers (SPU_ECC_Cntl, SPU_ECC_Stat, SPU_ECC_Addr, SPU_ERR_Mask). In addition, a set of seven MFC Fault Isolation, Error Mask, Checkstop Enable Registers are provided.

If a DMA transfer is halted due to an illegal command, an SPU-initiated interrupt is generated and the hypervisor can access the MFC Command Error Register (MFC_CER) to determine which DMA command in the MFC queue caused the error.

#### 11.3.4.5 *Debug and Performance Monitoring*

Several trace, trigger, and event control registers are provided to support debug and performance monitoring of all major logic blocks in the SPEs.

The hypervisor can enable DMA debug facilities through the use of the MFC address-compare facilities. The facilities are provided by the MFC Address Compare Control Register (MFC_ACCR), the MFC Local Storage Address Compare Register (MFC_LSACR), and the MFC Local Storage Compare Result Register (MFC_LSCRR). These registers are replicated for each SPE and allow halting of MFC DMA transfers when a specified LS address or main-storage page is read or written by a DMA transfer. The MFC Data Storage Interrupt Pointer Register (MFC_DSIPR) provides an index to the DMA command in the queue that triggered the address compare.

## 11.4 I/O-Address Translation

When a hypervisor creates logical partitions, it is important to restrict main-storage access by bus-mastering I/O devices (that is, I/O devices that can initiate DMA transfers to and from main storage), if the I/O devices are to be controlled directly by an operating system within the logical partition. Although a hypervisor can support I/O devices directly and it can fully virtualize them to software in the logical partition, this adds unacceptable latency and reduced throughput in many real-time or high-performance systems.

### 11.4.1 IOC Memory Management Units

The I/O interface controller (IOC) has an MMU that provides I/O-address translation and access protection for main-storage accesses initiated by I/O devices. The hypervisor must exercise full and exclusive control over the IOC's MMU to ensure that I/O devices supporting a logical partition cannot access resources in another logical partition. I/O devices access main storage through an I/O-address space that is separate from the real-address (RA) space.

The hypervisor can set up the IOC's MMU to support translation of I/O addresses to RAs. To do this, the hypervisor must enable I/O translation by setting the Enable bit of the I/O Segment Table Origin Register (`IOC_IOST_Origin[E]`) to '1' and the Translate Enable bit of the IOCmd Configuration Register (`IOC_IOCmd_Cfg[TE]`) to '1'. See *Section 7.4 I/O Address Translation* on page 176 for details.

### 11.4.2 I/O Segment and Page Tables

When an I/O device initiates a DMA transfer, the IOC MMU looks the supplied I/O address up in an I/O segment table (IOST), then uses the segment table to look up the RA in an I/O page table (IOPT). Unlike PPE MMU and SPE MMU translations, I/O translations must be present in the IOST or IOPT cache entries before an I/O device attempts access. The CBEA processors do not support the suspension and resumption of an I/O device's bus-mastering operation in order for the hypervisor to dynamically load IOST or IOPT entries for translations.

The hypervisor can enable the use of an IOST and IOPT in main storage—which the IOC MMU will access when an address is not found in the IOST or IOPT cache—by setting the `IOC_IOST_Origin[HW]` bit = '1'. If this option is not enabled, the hypervisor must directly manage the IOST and IOPT caches and set the `IOC_IOST_Origin[SW]` bit = '1'. Each IOST entry contains a real page number (RPN) of the start of the IOPT for the segment, the number of 4 KB pages used for the segment's IOPT, the I/O page size used for all accesses in the segment, a Valid (V) bit and a Hint (H) bit. The Hint bit can be used by the hypervisor to keep frequently-used I/O-segment translations in the IOST cache.

When a valid IOST entry is found for an I/O address, bits in the I/O address are then used to index into the IOPT to look for the corresponding page-table entry (PTE). The PTE contains page-protection bits (restricting read and write access), a bit controlling cache coherency for the transfer, bits specifying storage-access ordering for the transfer, a Hint bit, an I/O Identifier (IOID), and the RPN corresponding to the RA to be used for the transfer. The IOID in each PTE prevents one I/O device from using another I/O device's address translation. This provides the necessary protection when I/O devices are allocated to different partitions. These IOIDs are normally assigned by the system supplier and are generated by the host bridges attached to the I/O Interfaces.

If an I/O-device access does not find a valid IOST or IOPT entry, or if the access is prohibited by the IOPT protection bits, or if the IOID supplied by the device does not match the IOID in the IOPT, an I/O-device exception occurs and an external interrupt is presented to the PPE.

## 11.5 Resource Allocation Management

The resource allocation management (RAM) facility is described in *Section 8 Resource Allocation Management* on page 203. The hypervisor can apply the RAM facility to logical partitions. This facility can allocate specific amounts of memory or I/O bandwidth to the requesters (PPE, SPEs, or I/O devices) in each logical partition. The combined capabilities of logical partitioning and resource allocation management provide a level of resource control that is not typically found in conventional logical-partitioning processors.

### 11.5.1 Combining Logical Partitions with Resource Allocation

Conventional logical-partitioning processors divide up real memory, processing units, and I/O devices among logical partitions. This works well in most environments, but when one or more operating environments in the logical partitions must meet real-time deadlines, this type of resource allocation is insufficient.

For example, if partition 1 is allocated 70% of the PPE and five of the eight SPEs, and partition 2 is allocated 30% of the PPE and three of the eight SPEs, it is possible for applications in partition 2 to use so much main-storage bandwidth or I/O bandwidth that logical partition 1, which has more processing units, cannot meet its real-time schedule because the remaining memory or I/O bandwidth is insufficient. To alleviate this problem, the hypervisor can use the EIB resource allocation facility to manage the amount of real memory and I/O capacity the units allocated to a logical partition receive.

### 11.5.2 Resource Allocation Groups and the Token Manager

In the CBEA processors, the requesters assigned to each logical partition can be assigned to one of four managed resource allocation groups (RAGs). The hardware supports a token manager that generates tokens for the managed resources at hypervisor-programmed rates for each RAG. *Section 8.5* on page 213 describes the details.

Typically, the system supplier provides configuration information to the hypervisor in order for the hypervisor to determine maximum usable capacity of the managed resources. The hypervisor uses this information to ensure that the token-generation rate for the resource in all concurrently active RAGs does not exceed the maximum capacity of the resource.

Each managed resource also provides a feedback mechanism to the token manager, indicating that it is becoming overcommitted. The hypervisor can enable an external interrupt to the PPE when a resource becomes overcommitted. The hypervisor programs queue-threshold levels in the IOC units (for IOIF0 and IOIF1) and in the memory interface controller (MIC) for the DRAM memory banks.

The hypervisor uses these methods to assign resources to managed logical partitions:

- *PPE*—When the hypervisor assigns the PPE to a managed logical partition, the hypervisor enables the RAG requester and sets a resource allocation identifier (RAID) for that logical partition using the BIU Mode Setup Register 1 (`BIU_ModeSetup1`).

- *SPE*—When the hypervisor assigns an SPE to a managed logical partition, the hypervisor enables the RAG requester and sets an RA Identifier for that logical partition using the `RA_Enable` register and the `RA_Group_ID` register.
- *I/O*—When the hypervisor assigns an IOIF unit to a managed logical partition, the hypervisor enables the RAG requester and sets a virtual channel for outbound transfer from that logical partition using the `IOC_IOCmd_Cfg` register.

For complete details, see *Section 8 Resource Allocation Management* on page 203.

## 11.6 Power Management

The CBEA processors can support multiple power management states, as described in *Section 15 Power and Thermal Management* on page 429. Many of the units in the CBEA processors support dynamic power management when the units are not being used, and this requires no explicit action on the part of software. However, the CBEA processors do support explicit power management for the low-power states—the MFC Pause, SPU Pause, SPE State Retained and Isolated (SRI), PPE Pause (0), and Slow states.

### 11.6.1 Entering Low-Power States

When multiple logical partitions are present, only the hypervisor can determine if conditions are right to enter or exit these low-power states. Under certain system-specified conditions, it might be advisable to slow the processor core clock (`NClk`) frequency, as described in *Section 13 Time Base and Decrementers* on page 381.

When a logical partition yields both of its PPE threads, and no other logical partitions are dispatchable (all waiting on external events), the hypervisor can elect to enable the PPE Pause (0) state and suspend each thread. Before entering this state, the hypervisor must write a PPE control register to designate the conditions under which the PPE Pause (0) state is to be exited. Machine check, system error, and thermal interrupts must always exit the PPE Pause (0) state. The hypervisor can conditionally enable resumption on a decrementer event or on an external interrupt.

A low-power, system-specific hibernation state might or might not be supported by the system (it is not specifically supported by the Cell Broadband Engine Architecture [CBEA]). If hibernation state is supported by the system, it requires support from the system controller and system designer, as well as the hypervisor and in many cases the I/O devices. Hibernation mode consists of saving all volatile states into system DRAM, switching the DRAM to self-timed refresh mode, and disabling the power to the CBEA processor. It also requires that an 8 KB block of contiguous physical DRAM be left unused by system firmware or software, and dedicated to the memory interface controller (MIC) recalibration-pattern buffer.

### 11.6.2 Thread State Suspension and Resumption

Unlike previous PowerPC processing units, the PPE retains the full hardware thread state when the thread is suspended. When the thread is resumed, it starts executing from the reset vector, and the reset interrupt handler can choose to resume the execution at the instruction following the instruction that suspended the thread.

When the hypervisor has determined that an SPE is no longer being used by a logical partition (it has freed the resource), the SPE can be held in a low-power SPU Pause state, with the SPU stopped and the MFC suspended. If the hypervisor determines that the SPE will not be used for a relatively long period (several seconds or minutes), the hypervisor can place the SPU in the SPE State Retained and Isolated (SRI) state.

## 11.7 Fault Isolation

Many faults detected by the CBEA processors, including programming errors related to hypervisor resources, are nonrecoverable and indicate a significant failure that might compromise data integrity. These faults result in a checkstop (see *Section 9.5.2* on page 249), which immediately halts all CBEA processor elements and must be handled or recorded by a system controller external to the CBEA processor.

However, some faults are recovered by the CBEA processor hardware, are software recoverable, or might be caused by software errors in the hypervisor or logical partition. Faults detected while the PPE is accessing a resource generate a machine check interrupt (also *Section 9.5.2* on page 249) that must be handled by the hypervisor. Some faults detected by the SPE generate a system error interrupt (also *Section 9.5.16* on page 260) that also must be handled by the hypervisor or logical partition.

## 11.8 Code Sample

The following code sample illustrates one method of making an ABI-compliant hypervisor call. In particular, it demonstrates making a example (H_EX) request that accepts two input parameters (`arg1` and `arg2`) and returns two resulting words. The hypervisor call is invoked by executing the system call instruction (**sc**) with the LEV field of '1'. The hypervisor request is specified in register 3.

### 11.8.1 Error Codes and Hypervisor-Call (hcall) Tokens

An example set of error codes and hypervisor call (hcall) tokens:

```
#define H_Success;        0
#define H_Hardware;      -1 /* Hardware error */
#define H_Function;      -2 /* Not supported */
#define H_Privilege;     -3 /* Caller not in privileged mode */
#define H_Parameter;     -4 /* Partition violation/conflict. */
```

### 11.8.2 C Functions for PowerPC 64-bit ELF Hypervisor Call

Examples of creating C-callable functions that perform a hypervisor call for PowerPC 64-bit ELF:

```
filename: hcall_ex.s
        .set     H_EX, 0x104     # example hcall token
        .section ".opd","aw"
        .align 3
```

```
        .globl hcall_ex
hcall_ex:
        .quad .hcall_ex,.TOC.@tocbase, 0
        .previous
        .size hcall_ex,24
        .section ".text"
        .align 2
        .globl .hcall_ex
.hcall_ex:
        std     r3,-8(1)   # r3 (array of values) stored in stack
        li      r3,H_EX    # load r3 with hypervisor code
        sc 1               # Hypervisor Trap
        ld      r12,-8(1)  # reload array into r12
        cmpi    0,12,0     # only store return regs if array is non-NULL
        bne     ret2       # this &hcall; only returns contents of r4,r5
        blr                # return no values

ret8:   std r11,(7 * 8)(r12)
ret7:   std r10,(6 * 8)(r12)
ret6:   std r9,(5 * 8)(r12)
ret5:   std r8,(4 * 8)(r12)
ret4:   std r7,(3 * 8)(r12)
ret3:   std r6,(2 * 8)(r12)
ret2:   std r5,(1 * 8)(r12)
ret1:   std r4,(0 * 8)(r12)
        blr
```

Example of how to call the preceding assembler code:

```
filename: ex.c
void example(unsigned long long arg1, unsigned long arg2)
{
        long  rc;
        unsigned long results[2];

        rc = hcall_ex(results, arg1, arg2);
        if (rc != H_Success) {
                ... Failure Case ...
        }
        ...
}
```

# 12. SPE Context Switching

## 12.1 Introduction

The ability to save and restore the context of a Synergistic Processor Element (SPE) is important for two primary reasons: preemptive or cooperative context-switching SPE tasks (programs) for virtualization of the SPEs, and debugger access and interaction with an SPE's execution state.

Saving and restoring an SPE context can be very expensive in terms of both time and system resources. The complete SPE context consists of:

- 256 KB of local storage (LS)
- 128 128-bit general purpose registers (GPRs)
- Special purpose registers (SPRs)
- Interrupt masks
- Memory flow controller (MFC) command queues
- SPE channel counts and data
- MFC synergistic memory management (SMM) state
- Other privileged state

In all, the SPE context occupies more than 258 KB of storage space.

A programming model that allocates SPEs in a serially reusable manner—in which an SPE is assigned a task until it completes, then it is assigned another task—generally results in more efficient use of SPEs than does a programming model that implements frequent context switches. The serially reusable model requires fewer context switches, and less context needs to be saved and restored during a context switch.

A *preemptive context switch* facilitated by privileged software on the PPE is the most costly form of context save and restore, because the context is not known and a worst-case save and restore of all context state must be performed. An *application-yielding* context switch, in which the application using the SPE determines the timing and amount of context to be saved and restored, can minimize the overhead in terms of cycles and space used to save and restore context.

***Note:*** *If an I/O device interfaces directly to an SPE, the SPE's physically mapped LS interacts with that device. Because the LS is part of the context that is switched, preemptive context switching of an SPE that interfaces directly with an I/O device should be avoided, unless the I/O device can first be reliably quiesced.*

Several machine resources must be saved and restored when switching an SPE's context. This section provides a basic overview of resources and save-restore sequences for a preemptive context switch. The sample sequences assume the SPE is already stopped at context-save time. Some operating environments can forgo restoring the context of an SPE that has been stopped, especially when stopped on error conditions.

## 12.2 Data Structures

The data that must be saved to suspend and later resume execution of an SPE task is stored in two related structures. One is in the LS of the SPE being saved or restored; the other is in main storage:

- *Local Storage Context Save Area (LSCSA)*—A data structure in the SPE's Local Storage (LS). It contains the local problem state that the SPE itself saves to and restores from during a context switch. The LSCSA is a temporary data structure that is copied to/from the permanent copy in the CSA.
- *Context Save Area (CSA)*—A data structure in main storage that contains all of the SPE state, both problem state and privileged state. The CSA includes a copy of the LSCSA.

### 12.2.1 Local Storage Context Save Area

The LSCSA is the area of LS in which the synergistic processor unit (SPU) temporarily stores local problem state—primarily the GPRs, SPRs, and channel state. This is by far the smaller of the two data structures, requiring only a few KB of space.

The SPU saves to the LSCSA as part of a code sequence during a context-save, and it restores from the LSCSA as part of another code sequence during a context-restore. During a context-save, after saving to its LSCSA in the LS, the SPU then saves the LSCSA (along with the high 240 KB of the LS) to the CSA in main storage. Likewise, during a context-restore, the SPU first loads the LSCSA (along with the high 240 KB of the LS) into the LS from the CSA in main storage, and then the SPU restores its own local state from the LSCSA in LS.

Due to DMA-alignment restrictions, the LSCSA memory region should, at a minimum, be aligned to a 128-byte boundary. For many environments, it might be more practical to align the region to a page boundary.

### 12.2.2 Context Save Area

The CSA is a main-storage structure that contains the entire SPE context state, including the 256 KB LS, the 128 GPRs, the SPRs, the state of all status registers, interrupt masks, MFC command queues, SPE channel counts and data, MFC SMM state, and other privileged state. During save and restore sequences, small parts of the CSA are written and read by the PPE, and large parts are written and read by the SPE.

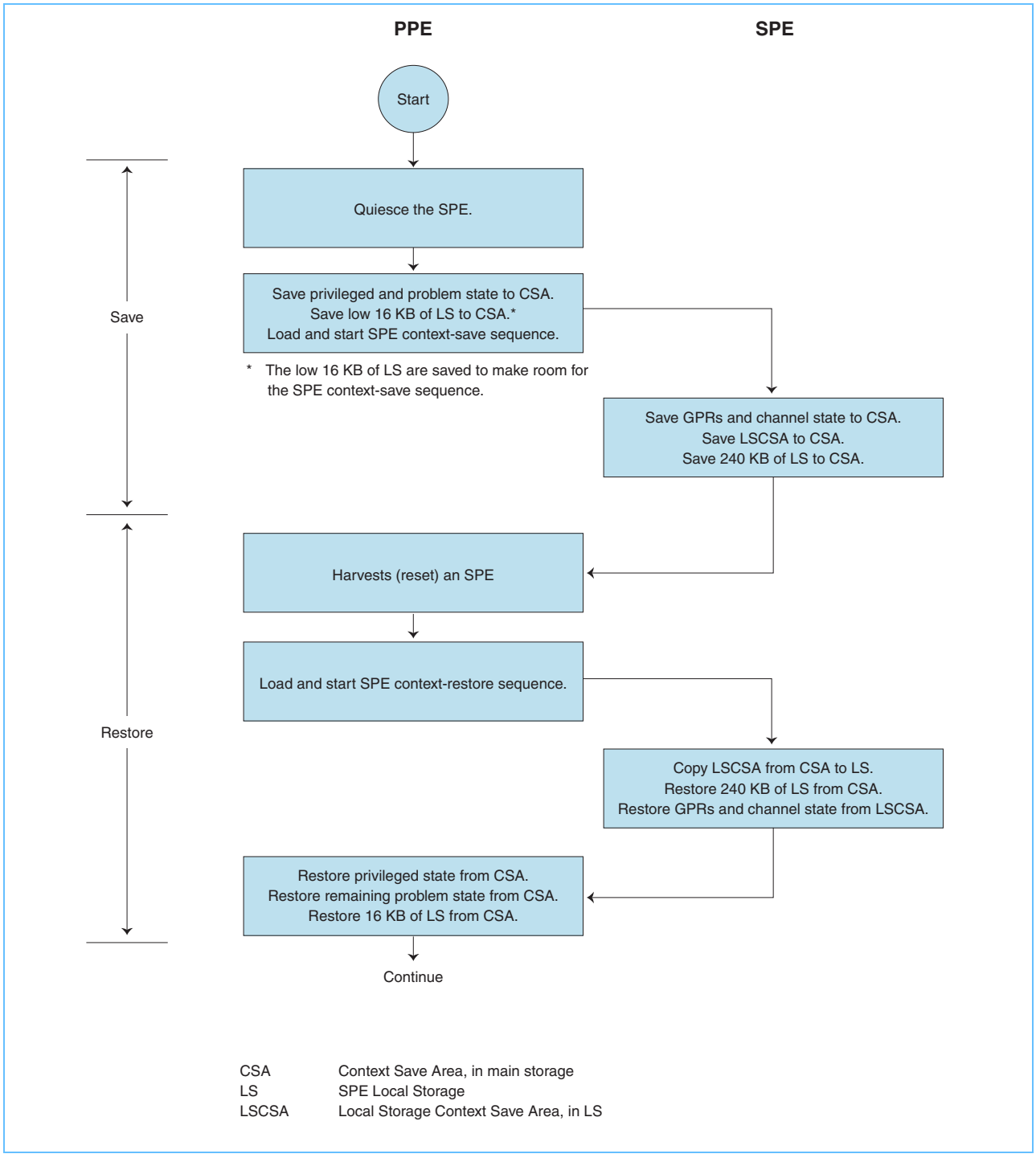
It might be best to order the CSA data so that the save areas for the LSCSA copy come first, followed by the save area for the LS. Otherwise, the SPE context-switch code might need to load immediate values for offsets that exceed the size of the LS, and thus require multiple assembly instructions to code.

## 12.3 Overview of SPE Context-Switch Sequence

This section gives an overview of the steps necessary for a preemptive context switch of an SPE task. A preemptive context switch consists of the combination of a context-save (*Section 12.3.1* on page 360) followed by a context-restore (*Section 12.3.2* on page 360). A full, detailed sample sequence is provided in *Section 12.5* on page 365.

The entire SPE context is not accessible from the PPE. As such a SPE context save or restore requires the assistance of specialized SPE code. A pictorial overview of the SPE context switch sequence is provided in *Figure 12-1*.

*Figure 12-1. SPE Context-Switch Sequence*



### 12.3.1 Save SPE Context

1. *Quiesce the SPE* (performed by the PPE):
   a. Disable SPE interrupts.
   b. Disable problem-space access to this SPE (unmap pages).
   c. Disable MFC command queues.
   d. Disable SPU execution.
   e. Disable all access to this SPE from other SPEs (unmap pages).
   f. Handle pending interrupts from this SPE.
2. *Save Privileged and Critical Problem State in CSA, Load and initiate SPU Context-Save Sequence on the SPU* (performed by the PPE):
   a. Save MFC command state, including command queues.
   b. Save privileged state, including segment lookaside buffer (SLB) entries, resource allocation masks, and so forth.
   c. Save problem state, including mailbox data and SPE channel counts.
   d. Purge the MFC command queues.
   e. Resume MFC command processing.
   f. Invalidate MFC SLB entries, and reload with SLB entry for loading the SPU context-save sequence code (step 3).
   g. Invalidate MFC translation lookaside buffer (TLB) entries.
   h. Reset interrupts.
   i. Save the low 16 KB of LS in the CSA.
   j. Place CSA effective-address pointer in SPU Signal Notification Registers.
   k. Copy the SPU context-save sequence code into the low 16 KB of LS, and start its execution.
3. *SPU Context-Save Sequence* (performed by the SPE):
   a. Save the low 16 GPRs in the LSCSA.
   b. Read the SPU Signal Notification Registers to obtain effective address of the CSA.
   c. Save the SPU accessible channel state in the LSCSA.
   d. Save the high 240 KB of LS in the CSA.
   e. Save the high 112 GPRs in the LSCSA.
   f. Save the LSCSA in the CSA.

### 12.3.2 Restore SPE Context

1. *Harvest (reset the state of) an SPE* (performed by the PPE):
   a. Disable SPE interrupts.
   b. Disable problem-space access to this SPE (unmap pages).
   c. Disable all access to this SPE from other SPEs (unmap pages).

d. Disable MFC command queues.
e. Disable SPE execution.
f. Invalidate MFC TLB entries.
g. Handle pending interrupts from this SPE.
h. Purge MFC command queues.
i. Invalidate MFC SLB entries.
j. Reset SPE channel state.
k. Reset interrupts.

2. *Load and initiate SPU Context-Restore Sequence on the SPU* (performed by the PPE):
   a. Invalidate MFC SLB entries, and reload with entry for SPU context-restore sequence code access.
   b. Copy mailbox state from the CSA to the CSA's copy of the LSCSA.
   c. Place CSA effective-address pointer in SPU Signal Notification Registers.
   d. Copy the SPU context-restore sequence code into the low 16 KB of LS, and start execution.
3. *SPU Context-Restore Sequence* (performed by the SPE):
   e. Read SPU Signal Notification Registers to obtain effective address of the CSA.
   a. Restore the high 240 KB of the LS from the CSA.
   b. Restore the high 112 GPRs from the LSCSA.
   c. Restore channel and other problem state from the LSCSA.
   d. Restore low 16 GPRs from the LSCSA.
4. *Restore Privileged and Remaining Problem State* (performed by the PPE):
   a. Restore SPU run control state.
   b. Restore low 16 KB of LS from the CSA.
   c. Restore interrupt state.
   d. Restore MFC command queues from CSA.
   e. Restore privileged state from the CSA.
   f. Restore additional problem state from the CSA.
   g. Restore MFC SLB entries from the CSA.
   h. Enable SPU execution.
   i. Restore MFC control state from the CSA.
   j. Enable interrupts.

## 12.4 Implementation Considerations

### 12.4.1 Locking

Due to the length of the context-switch sequence, it might be appropriate to use semaphores rather than spin locks for access-locking. This allows other operations (for example, context-switch on other SPEs) to proceed when multiple threads are attempting to exclusively access the same SPE during context switch.

### 12.4.2 Watchdog Timers

Non-real-time systems might not be able to set reasonable time-out values for the entire context-switch sequence. This is due to the potential for paging activity, unexpected external events (for example, network), and so forth. It might be more appropriate to set timers on individual operations—for example, waiting for a purge of MFC command queues to finish.

### 12.4.3 Waiting for Events

Non-real-time systems might suffer significantly from excessive busy-polling for SPE context-switch events. It might be more appropriate to yield or sleep on events, rather than busy-poll. Enabling handlers for SPE class 2 interrupts (stop-and-signal, tag-group completion) might be useful. See *Section 9.6.3 SPU and MFC Interrupts* on page 271 for details. For a description of SPE events, see *Section 18 SPE Events* on page 471.

### 12.4.4 PPE's SPU Channel Access Facility

Privileged PPE software has access to the four event-management channels and the hidden Pending Event Register (*Section 18.2 Events and Event-Management Channels* on page 472) through the PPE's SPU channel access facility. The SPU channel access facility can initialize, save, and restore the SPU channels.

The facility consists of three memory-mapped I/O (MMIO) registers:

- SPU Channel Index Register (`SPU_ChnlIndex`)
- SPU Channel Count Register (`SPU_ChnlCnt`)
- SPU Channel Data Register (`SPU_ChnlData`)

The SPU Channel Index Register is a pointer to the channel whose count is accessed by the SPU Channel Count Register and whose data is accessed by the SPU Channel Data Register. See the *Cell Broadband Engine Architecture* document for details about the PPE's SPU channel access facility.

### 12.4.5 SPE Interrupts

It is possible to perform a context switch with SPE interrupts disabled. In this case, memory regions containing the context-switch program or data should be pinned (locked), and the MFC's SMM should be pre-loaded with the SLB and TLB entries necessary for these memory regions. Polling of MMIO registers is necessary to determine the completion status for data transfers, SPE execution, and so forth.

Non-real-time environments might prefer to enable interrupt handlers for address-translation (SPE class 1) interrupts, so that code or data can be accessed as virtual memory rather than pinned real memory. Handlers for error (SPE class 0) interrupts and application (SPE class 2) interrupts can be used to detect potential errors or completion events (for example, tag-group complete, stop-and-signal). See *Section 9.6.3 SPU and MFC Interrupts* on page 271 for details.

### 12.4.6 Suspending the MFC DMA Queue

Depending on the operating system (OS), address-translation (SPE class 1) interrupts might be processed by a deferred interrupt handler. The use of deferred interrupt handling complicates the context-save operation in the way that the MFC command queues are suspended and saved.

In this case, it is not enough for the context-save code to simply disable SPE interrupts. It is also necessary to prevent a deferred interrupt handler from restarting an MFC command after the queue has been suspended. This can be done by setting a per-SPE save-pending flag, indicating to the deferred handler that the MFC command should not be restarted.

### 12.4.7 SPE Context-Save Sequence and Context-Restore Sequence Code

The SPE's context-save sequence and context-restore sequence code should probably be written in SPE assembly language, so that register and memory use can be carefully controlled. The sequence described in *Section 12.5* on page 365 assumes the code is written to execute in a portion of the low 16 KB of LS and uses only the first 16 GPRs.

A straight-forward implementation of either the save or restore sequence will typically use less than 1 KB for the program text. Because the save and restore sequences have much in common, it is possible to write a combined switch program that fits into roughly the same space. A further 2 KB will be used at run time for saving or restoring registers.

Special **stop** instruction codes should be reserved for indicating the successful completion of a save and a restore.

### 12.4.8 SPE Parameter Passing

The only parameter needed by the save and restore code is the location of the CSA region in main storage. This can be communicated either by having privileged software running on the PowerPC Processor Element (PPE) store this address to a known location in LS before starting the SPE or by writing this 64-bit address to the 32-bit SPU Signal Notification 1 and 2 Registers as is done in *Section 12.5* on page 365.

### 12.4.9 Storage for SPE Context-Save Sequence and Context-Restore Sequence Code

The SPE context-switch code is trusted by the OS and must therefore reside in privileged memory. The code must be copied by the PPE from main storage into LS, so it should be, at a minimum, aligned to a 128-byte boundary for efficient transfer by the MFC. It might be possible to dump the SPE context-switch object code into raw hexadecimal form, so that it can be initialized and aligned in memory at compile time, rather than run time.

### 12.4.10 Harvesting an SPE

There might be situations in which an SPE context is not immediately restored to the SPE on which it was originally executing. This might be the case when the context is executing under control of a debugger. There might also be conditions that cause the SPE context-save operation to fail (for example, a bad address translation for a save to the CSA in the user-memory region).

In such cases, this or another SPE needs to be harvested. Harvesting means resetting (clearing) the state of an SPE, and it allow a subsequent restore operation to be performed to that SPE. The steps in the context-restore operation in *Section 12.5.2.1* on page 371 describe how an SPE can be harvested.

### 12.4.11 Scheduling

Sharing SPEs, as opposed to serially reusing SPEs, can be very expensive. A fully preemptive SPE context switch can take many μsecs. It might, however, still be desirable for an OS to share SPEs among multiple user programs. What an individual SPE gives up in terms of context-switch overhead, the Cell Broadband Engine Architecture (CBEA) processors[1] can make up for in the number of available SPEs. Taken together, overall system interactivity might be considered reasonable, even if the latency for an individual SPE switch is somewhat high.

To avoid thrashing, the SPE-switch time should be factored into the time-slice quantum that is assigned to an SPE context. Typically, shorter quanta are used for interactive processing, and longer quanta are used for batch processing. Selection of time-slice quanta is a always compromise between these two processing goals. The best approach is to use a quantum that is short enough to seem interactive, but long enough to allow forward progress. A good rule of thumb for minimum quanta on a more traditional processor might be 10 times the switch-time. However, given the abundance of SPEs and relatively long switch time, it is probably best to err on the side of longer rather than shorter quanta.

### 12.4.12 Light-Weight SPE Context Save

The sequence outlined in *Section 12.5* on page 365 describes a fully preemptive SPE context switch. In certain environments, however, it might be desirable to support light-weight context switching. This can be implemented in a cooperative manner by having the PPE send a privileged attention event to a program executing on an SPE. If the SPE program responds to the event within a fixed time-out—for example, by executing the stop and signal (**stop**) instruction with special return code—then the PPE might perform a simplified context-save sequence. Otherwise, the PPE performs a full context save, and might subsequently penalize that program by lowering its priority within the system.

To implement such a policy, there would be some cooperation between the OS and SPE application code. For example, upon receipt of the privileged attention event, the SPE program might wait for all pending MFC commands to complete before executing the **stop** instruction. This allows the PPE to avoid suspending and saving the MFC command queues.

Another alternative is to support reentrant contexts that are never saved, provided that they respond to the privileged attention event within a fixed time-out period. There are several variations possible between a fully preemptive switch and reentrant contexts.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

## 12.5 Detailed Steps for SPE Context Switch

*Section 12.3* on page 358 gave an overview of the sequence of steps involved in preemptively saving and restoring an SPE context. The sequences in this section expand on that overview by providing more detail.

These sequences depict the saving and restoring of SPE context by an OS that is running at both privilege 1 and privilege 2 levels—that is, an operating environment without a separate hypervisor and without support for logical partitioning (see *Section 12.6 Considerations for Hypervisors* on page 379). The sequences also assume that the OS uses hardware-managed TLBs and a common hardware-managed page table, and that it has allocated an area of main storage for the CSA, as described in *Section 12.2 Data Structures* on page 358.

Although the sequences handle preemptive saving and restoring, they can also handle cases in which the SPE is already stopped at context-save time or has stopped itself after context-save read its status. Some operating environments might forego restoring the context of an SPE that was not running at context-save time (especially when stopped on error conditions). However, support for restoring is included in this example for completeness.

The sequences do not include saving and restoring performance-monitor and trace registers or application-debug facilities. Interaction between the context-switching system and the debug and performance-monitor subsystems dictate how these registers are dealt with during context switch.

### 12.5.1 Context-Save Sequence

#### 12.5.1.1 *PPE Context-Save Sequence for SPE*

This section outlines how to fully save the context of an active SPE:

1. Acquire a software SPE mutual-exclusion lock (if multithreading is being used).
2. If any `SPU_Status[E, L, IS]` field is set to ‘1’, this SPE is in the isolate state and cannot be context-saved at this time:
   - If multithreading is being used, release software SPE mutual-exclusion lock.
   - If the isolate state is an unexpected condition, harvest an SPE by going to step 1 of the PPE Context-Restore Sequence on page 371.
   - If the isolate state is allowed, exit the context-save sequence.
3. Disable all interrupts (class 0, 1, and 2) from this SPE:
   a. Save the `INT_Mask_class0` register in the CSA. Then, write the `INT_Mask_class0` register with ‘0’.
   b. Save the `INT_Mask_class1` register in the CSA. Then, write the `INT_Mask_class1` register with ‘0’.
   c. Save the `INT_Mask_class2` register in the CSA. Then, write the `INT_Mask_class2` register with ‘0’.
4. Set a software watchdog timer, specifying the maximum time for a context-save sequence.
5. Inhibit problem-space access to this SPE, if provided, by unmapping the virtual pages assigned to the SPE MMIO problem space.

6. Read the SPU_Status register again. If any SPU_Status[E, L, IS] fields are active, this SPE is isolated and cannot be context-saved at this time. If the isolate state is an expected condition, the context-save sequence should be exited by skipping to step 74 on page 378. If this is an unexpected condition, skip to step 4 on page 371 to reclaim this SPE and restore a context to it.
7. Set a software Context Switch Pending flag.
8. Read the MFC_CNTL[Ss] register field:
   - If MFC_CNTL[Ss] is set to '01', poll MFC_CNTL[Ss] until it is set to '11'. Then, save the MFC_CNTL register in the CSA with MFC_CNTL[Sc] set to '1' and MFC_CNTL[Sm] set to '1'.
   - If MFC_CNTL[Ss] is set to '00', set MFC_CNTL[Sc] to '1'. Then, poll MFC_CNTL[Ss] until it is set to '11', and save the MFC_CNTL register in the CSA with MFC_CNTL[Sc] set to '0' and MFC_CNTL[Sm] set to '0'.
   - If MFC_CNTL[Ss] is set to '11', save the MFC_CNTL register in the CSA with MFC_CNTL[Sc] set to '1' and MFC_CNTL[Sm] set to '1'.
9. Save the SPU_RunCntl register in the CSA. This value contains the Application Desired State.
10. Save the MFC_SR1 register in the CSA.
11. Read the SPU_Status[R] register field:
    - If SPU_Status[R] is set to '0', save the SPU_Status register in the CSA.
    - If SPU_Status[R] is set to '1':
      a. Set SPU_RunCntl[Run] to '00'.
      b. Issue a PPE **eieio** instruction.
      c. Poll SPU_Status[R] until it is set to '0'. If SPU_Status[I,S,H, and P] are all set to '0', write the SPU_Status register in the CSA with the R bit set to '1'. Otherwise, write the SPU_Status register in the CSA with the R bit reset to '0'.
12. Read the MFC_CNTL[Ds] register field. Then, update the saved copy of this field in the CSA.
13. Write the MFC_CNTL[Dh] register field to '1' to halt the decrementer.
14. Read the special purpose Time Base Register (TB) and save it in the CSA.
15. Prevent other SPE accesses to this SPE by unmapping this SPE's pages from their address spaces.
16. Write the MFC_MSSync register. Then, poll MFC_MSSync[P] for a value of '0'.
17. Write all of the MFC_TLB_Invalidate_Entry[IS,VPN,L,Lp] register fields to '0'. Then, issue a PPE **sync** instruction.
18. Handle any pending interrupts from this SPE. This is OS-specific or hypervisor-specific. One option is to re-enable interrupts to handle any pending interrupts, with the interrupt handlers recognizing the software Context Switch Pending flag, to ensure that SPE execution or the MFC command queue is not resumed.
19. If the MFC_CNTL[Q] register field is set to '0' (MFC command queues not empty), then save 96 doublewords from the MFC_CQ_SR registers in the CSA.
20. Save the Prxy_QueryMask register in the CSA.
21. Save the Prxy_QueryType register in the CSA.
22. Save the MFC_CSR_TSQ register in the CSA.

23. Save the MFC_CSR_Cmd1 and MFC_CSR_Cmd2 registers in the CSA.
24. Save the MFC_CSR_ATO register in the CSA.
25. Save the MFC_TClassID register and in the CSA.
26. Write the MFC_TClassID register with the value x‘10000000’ (T0 Quota = 16).
27. Write the MFC_CNTL[Pc] register field to ‘1’ (purge queue).
28. Poll the MFC_CNTL[Ps] register field until the value ‘11’ is read (purge complete).
29. If the MFC_SR1[R] register field is set to ‘1’:
    a. Save the SLB_Index register in the CSA.
    b. For index = 0 through 7:
        1. Write the index value to SLB_Index. Then, issue a PPE **eieio** instruction.
        2. Save the SLB_ESID register in the CSA.
        3. Save the SLB_VSID register in the CSA.
30. Write the MFC_SR1 register with MFC_SR1[D] set to ‘0’, MFC_SR1[S] set to ‘1’, and MFC_SR1[TL,R,PR, and T] set correctly for the OS-specific environment.
31. Save the SPU_NPC register in the CSA.
32. Save the SPU_PrivCntl register in the CSA.
33. Write the SPU_PrivCntl[S,Le,A] register field to ‘0’.
34. Save the SPU_LSLR register in the CSA.
35. Set the SPU_LSLR register to the size of implemented LS, minus one.
36. Save the SPU_Cfg register in the CSA.
37. Save the PM_Trace_Tag_Wait_Mask register in the CSA.
38. Save the RA_Group_ID and RA_Enable registers in the CSA.
39. For the following channels:
    - SPU_RdEventStat (channel index 0)
    - SPU_WrEventMask (channel index 1)
    - SPU_RdSigNotify1 (channel index 3)
    - SPU_RdSigNotify2 (channel index 4)
    - MFC_RdTagStat (channel index 24)
    - MFC_RdListStallStat (channel index 25)
    - MFC_RdAtomicStat (channel index 27)
        a. Write the channel index to the SPU_ChnlIndex register. Then, issue a PPE **eieio** instruction.
        b. Save the channel data from SPU_ChnlData register in the CSA.
        c. Save the channel count from SPU_ChnlCnt register in the CSA.
        d. Write ‘0’ to the SPU_ChnlData register.
        e. Write ‘0’ to the SPU_ChnlCnt register.
40. Save the SPU_Mbox_Stat register in the CSA.

41. Save the SPU_Out_Mbox register in the CSA.
42. Save the SPU_Out_Intr_Mbox register in the CSA.
43. For the SPU_RdInMbox channel:
    a. Write the channel index value 29 to the SPU_ChnlIndex register. Then, issue a PPE **eieio** instruction.
    b. Save the SPU_ChnlCnt register in the CSA.
    c. Save the SPU_ChnlData register in the CSA's SPU_RdInMbox data0 area.
    d. Save the SPU_ChnlData register in the CSA's SPU_RdInMbox data1 area.
    e. Save the SPU_ChnlData register in the CSA's SPU_RdInMbox data2 area.
    f. Save the SPU_ChnlData register in the CSA's SPU_RdInMbox data3 area.
44. For the MFC_Cmd channel:
    a. Write the channel index value 21 to the SPU_ChnlIndex register. Then, issue a PPE **eieio** instruction.
    b. Save the SPU_ChnlCnt register in the CSA.
45. For the following channels:
    - MFC_Cmd (channel index 21, count = 16)
    - MFC_WrTagUpdate (channel index 23, count = 1)
    - SPU_WrOutMbox (channel index 28, count = 1)
    - SPU_WrOutIntrMbox (channel index 30, count = 1)
        a. Write the channel index to the SPU_ChnlIndex register. Then, issue a PPE **eieio** instruction.
        b. Write the count to the SPU_ChnlCnt register.
46. Write the MFC_CNTL[Sc] register field to '0' and MFC_CNTL[Sm] to '0' (resume queue processing).
47. If the MFC_SR1[R] register field is set to '1', write '0' to SLB_Invalidate_All register. Then:
    a. Write x'0' to the SLB_Index register.
    b. Write the SLB_VSID register with the virtual segment ID (VSID) to the SPU context-save sequence code.
    c. Write the SLB_ESID register with the effective segment ID (ESID) to the SPU context-save sequence code.
    d. Write x'1' to the SLB_Index register.
    e. Write the SLB_VSID register with the VSID to the CSA.
    f. Write the SLB_ESID register with the ESID to the CSA.
48. Change the software Context Switch Pending flag to Context Switch Active.
49. If the OS is using interrupts instead of polling:
    a. Write the INT_Stat_class0 register with -1 to reset all class 0 interrupts.
    b. Write the INT_Stat_class1 register with -1 to reset all class 1 interrupts.
    c. Write the INT_Stat_class2 register with -1 to reset all class 2 interrupts.

d. Write the INT_Mask_class0 register with an OS-specific value to re-enable OS-supported class 0 interrupts.

e. Write the INT_Mask_class1 register with an OS-specific value to re-enable OS-supported class 0 interrupts.

f. Write the INT_Mask_class2 register with an OS-specific value to re-enable OS-supported class 0 interrupts.

50. Issue a DMA command to save the low 16 KB of LS in the CSA:
    a. Write x‘0000_0000’ to MFC_LSA register.
    b. Write the CSA effective-address bits[0:32] to the MFC_EAH register.
    c. Write the CSA effective-address bits[33:63] to the MFC_EAL register.
    d. Write x‘4000’ (transfer size of 16 KB) to MFC_Size register.
    e. Write x‘00’ to MFC_TagID register.
    f. Write '0020' to the MFC_ClassID_Cmd register (**put** command, with TClassID and RClassID of '00').
    g. Read the MFC_CmdStatus[Rc] register field and, if value is not equal to ‘00’, go back to step 50a on page 369.
51. Write ‘0’ to the SPU_NPC[IE] register field, and write the SPU_NPC[LSA] register field to the LS entry-point address of the SPU context-save sequence code[2].
52. Write the SPU_Sig_Notify_1 register with the upper 32 bits of the effective address in the CSA where the copy of the LSCSA is stored. (This register is read in step 6 on page 370.)
53. Write the SPU_Sig_Notify_2 register with the lower 32 bits of the effective address in the CSA where the copy of the LSCSA is stored. (This register is read in step 8 on page 370.)
54. Issue a DMA command to the copy the SPU context-save sequence code from main-storage to the low 16 KB of the LS and start the SPU:
    a. Write the LS beginning address of the SPU context-save sequence code to the MFC_LSA register.
    b. Write the most-significant word of the main-storage address of the SPU context-save sequence code to the MFC_EAH register.
    c. Write the least-significant word of the main-storage address of the SPU context-save sequence code to the MFC_EAL register.
    d. Write the SPU context-save sequence code size to the MFC_Size register.
    e. Write x‘00’ to MFC_TagID register.
    f. Write '004A' to the MFC_ClassID_Cmd register (**getfs** command, with TClassID and RClassID of '00').
    g. Read the MFC_CmdStatus[Rc] register field and, if value is not equal to ‘00’, go back to step 54a on page 369.
55. Write the Prxy_QueryMask register to ‘1’ (enable Tag-Group 0). Then, issue a PPE **eieio** instruction.

2. See *Section 12.5.1.2* on page 370 for this sequence.

56. Poll the Prxy_TagStatus[$g_0$] register field until it reads ‘01’ (tag-group 0 complete), or write Prxy_QueryType[TS] to ‘01’ and wait for a tag-group complete interrupt. Then, write the INT_Stat_class2[T] register field until it reads ‘1’.
57. Poll the SPU_Status[R] register field until it reads ‘0’, or wait for an SPU class 0 or class 2 interrupt. Then, reset interrupt conditions by writing the INT_Stat_class0 register with ‘1’ for each interrupt handled, and by writing the INT_Stat_class2[T] bit with ‘1’.
58. If the SPU_Status[P] register field is set to ‘1’ and SPU_Status[StopCode] reads Success, the context-save has succeeded. Otherwise, the context-save failed.
59. Proceed to step 26 of the PPE Context-Restore Sequence on page 373.

#### 12.5.1.2 *SPU Context-Save Sequence*

The SPU code that performs the following context-save sequence is copied by the PPE into a portion of the low 16 KB of the LS and run in step 54 on page 369. This sequence copies the remainder of the LS and the contents of the local-storage context save area (LSCSA) into the CSA. The LSCSA consists of the GPRs, status registers, and the remaining channel state.

The SPU Context-Save sequence is:

1. Save the low 16 GPRs, using quadword stores, in the LSCSA. They are saved at this time for performance reasons.
2. Read the SPU_RdEventMask channel and save it in the LSCSA.
3. Read the MFC_RdTagMask channel and save it in the LSCSA.
4. Set the SPU_WrEventMask channel to ‘0’ to mask all events.
5. Set the MFC_WrTagMask channel to ‘01’ to unmask only Tag-Group 0.
6. Read the SPU_RdSigNotify1 channel data to obtain the upper 32-bits of the CSA effective address.
7. Write the high address from step 6 to the MFC_EAH channel.
8. Read the SPU_RdSigNotify2 channel data to obtain the low 32-bits of the CSA effective address.
9. Update the low effective addresses in a predefined, 15-element (16 KB each) DMA list that saves the most-significant 240 KB[3] of the LS to the CSA.
10. Enqueue a **putl** (Tag-Group 0) command for the DMA list updated in step 9 on page 370.
11. Save the high 112 GPRs, using quadword stores, in the LSCSA.
12. Issue the floating-point status and control register read instruction (**fscrrsd**) and save the status in the LSCSA.
13. Read and save the SPU_RdDec channel data in the LSCSA.
14. Read and save the SPU_RdSRR0 channel data in the LSCSA.
15. Enqueue a **putllc** command, using an effective address in the CSA, to remove any possible lock-line reservation.
16. Enqueue a **put** (Tag-Group 0) command to save the LSCSA to the CSA.

3. The size can differ from 240 KB, depending on the value of the SPU_LSLR register.

17. Enqueue an **mfcsync** (Tag-Group 0) command to the MFC_Cmd register.
18. Write the MFC_WrTagUpdate channel with '01' (update tag status on any completion).
19. Read the MFC_RdTagStat channel data. This stalls the SPU until DMA Tag-Group 0 is complete.
20. Read the MFC_RdAtomicStat channel data. This stalls until the **putllc** command completes.
21. Issue a stop-and-signal (**stop**) instruction with the status set to a success status. If there is a software-detected error in the SPU save-context code, the stop-and-signal status should be set to a diagnostic status instead of a success status.

## 12.5.2 Context-Restore Sequence

### 12.5.2.1 *PPE Context-Restore Sequence for SPE*

This section outlines how to fully restore the context of an SPE.

If the SPE context was successfully saved, as described in *Section 12.5.1.1* on page 365, proceed to step 26 on page 373. Otherwise, harvest an SPE by performing the following steps:

1. If multithreading is being used, acquire an SPU mutual-exclusion kernel lock.
2. Disable Interrupts from this SPE:
   - Write the INT_Mask_class0 register with '0' (disable all class 0 interrupts).
   - Write the INT_Mask_class1 register with '0' (disable all class 1 interrupts).
   - Write the INT_Mask_class2 register with '0' (disable all class 2 interrupts).
3. Inhibit problem-space access to this SPE, if provided, by unmapping the virtual pages assigned to the SPE MMIO problem space.
4. If required, notify the using application that the SPE task has been terminated.
5. Set software Context Switch Pending flag.
6. Remove other SPEs' access to this SPE by unmapping this SPE's pages from their address spaces.
7. Write the MFC_CNTL[Dh, Sc, Sm] register fields to '1', '1', '0', respectively, to suspend the queue and halt the decrementer.
8. Poll the MFC_CNTL[Ss] register field until '11' is returned (queue suspended).
9. If the SPU_Status[R] register field is set to '1' (running):
   - If SPU_Status[E] is set to '1', poll SPU_Status[R] until it reads '0' (stopped).
   - If SPU_Status[L] is set to '1', or if SPU_Status[IS] is set to '1':
     a. Write the SPU_RunCntl[Run] register field to '00' (stop request).
     b. issue a PPE **eieio** instruction.
     c. Poll SPU_Status[R] until it reads '0' (stopped).
     d. Write SPU_RunCntl[Run] to '10' (isolate exit request).
     e. Issue a PPE **eieio** instruction
     f. Poll SPU_Status[R] until it reads '0' (stopped).

- If SPU_Status[W] is set to ‘1’:
  a. Write SPU_RunCntl[Run] to ‘00’ (stop request).
  b. Issue a PPE **eieio** instruction.
  c. Poll SPU_Status[R] until it reads ‘0’ (stopped).
- Go to step 11 on page 372.

10. If SPU_Status[R] is set to ‘0’ (stopped):
    - If SPU_Status[E] is set to ‘1’:
      a. Write MFC_SR1[S] to ‘1’.
      b. Write SPU_RunCntl[Run] to ‘01’(run request).
      c. Issue a PPE **eieio** instruction.
      d. Poll SPU_Status[R] until it reads ‘0’ (stopped).
    - If SPU_Status[L] is set to ‘1’, or if SPU_Status[IS] is set to ‘1’:
      a. Write MFC_SR1[S] to ‘1’.
      b. Write SPU_RunCntl[Run] to ‘10’ (exit request).
      c. Issue a PPE **eieio** instruction.
      d. Poll SPU_Status[R] until it reads ‘0’ (stopped).
    - If there is any other value, no action is needed; the SPU is stopped in a nonisolated state.
11. Write the MFC_MSSync register. Then, poll MFC_MSSync[P] for a value of ‘0’.
12. Write all of the MFC_TLB_Invalidate_Entry[IS,VPN,L,Lp] register fields to ‘0’. Then, issue a PPE **sync** instruction.
13. Handle any pending interrupts from this SPE. This is OS-specific or hypervisor-specific. One option is to re-enable interrupts to handle any pending interrupts, with the interrupt handlers recognizing the software Context Switch Pending flag, to ensure that the SPE execution or MFC command queue is not resumed.
14. Write the MFC_CNTL[Pc] register field to ‘1’ (purge queue).
15. Poll MFC_CNTL[Ps] until it reads ‘11’ (purge complete).
16. Write the SPU_PrivCntl[S,Le,A] register fields to ‘0’ (reset).
17. Set the SPU_LSLR register to the size of implemented LS, minus one.
18. Write the MFC_SR1[D] bit to ‘0’, the MFC_SR1[S] bit to ‘1’, and the MFC_SR1[TL,R,PR, and T] bits set correctly for the OS environment.
19. If the MFC_SR1[R] bit is set to ‘1’, write ‘0’ to the SLB_Invalidate_All register.
20. For the following channels:
    - SPU_RdEventStat (channel index 0)
    - SPU_WrEventMask (channel index 1)
    - SPU_RdSigNotify1 (channel index 3)
    - SPU_RdSigNotify2 (channel index 4)
    - MFC_RdTagStat (channel index 24)
    - MFC_RdListStallStat (channel index 25)
    - MFC_RdAtomicStat (channel index 27)

a. Write the channel index to the SPU_ChnlIndex register. Then, issue an **eieio** instruction.

b. Write '0' to the SPU_ChnlData register.

c. Write '0' to the SPU_ChnlCnt register[4].

21. For the following channels:
    - MFC_Cmd (channel index 21, count = 16)
    - MFC_WrTagUpdate (channel index 23, count = 1)
    - SPU_WrOutMbox (channel count 28, count = 1)
    - SPU_RdInMbox (channel count 29, count = 0)
    - SPU_WrOutIntrMbox (channel count 30, count = 1)

    a. Write the channel index to the SPU_ChnlIndex register. Then, issue a PPE **eieio** instruction.

    b. Write the count to the SPU_ChnlCnt register.

22. If the OS is using interrupts instead of polling:

    a. Write the INT_Stat_class0 register with -1 (reset all class 0 interrupts).

    b. Write the INT_Stat_class1 register with -1 (reset all class 1 interrupts).

    c. Write the INT_Stat_class2 register with -1 (reset all class 2 interrupts).

    d. Write the INT_Mask_class0 register with OS-specific value (reset OS-supported class 0 interrupts).

    e. Write the INT_Mask_class1 register with OS-specific value (reset OS-supported class 1 interrupts).

    f. Write the INT_Mask_class2 register with OS-specific value (reset OS-supported class 2 interrupts).

23. Set the software Context Switch Active flag.

24. Write the MFC_TClassID register with the value x'10000000' (TClassID0 Quota=16).

25. Write the MFC_CNTL[Sc] register field to '0' and MFC_CNTL[Sm] to '0' (resumes queue processing).

At this point, either the SPE's context has been saved in the CSA or an SPE has been harvested. A context can now be restored from the CSA to the SPE, as described in the following steps:

26. Set a software watchdog timer to the maximum time for the restore sequence.

27. If any of the SPU_Status[I,S,H, P] bits are set to '1' in the CSA, then add the correct instruction sequence to the end of the SPU context-restore sequence code[5], after the context-restore stop-and-signal (step 22 on page 379) to restore the following values to the SPU_Status register:
    - SPU_Status[P] set to '1': Stop-and-signal instruction, followed by a branch instruction to itself[6].
    - SPU_Status[I] set to '1': Illegal instruction, followed by a branch instruction to itself.

4. The channel count need not be restored for the SPU_WrEventMask channel.
5. See *Section 12.5.2.2* on page 378.
6. A branch instruction to itself can be encoded in assembly as "br.".

- SPU_Status[H] set to '1': Unconditional halt instruction[7] followed by a branch instruction to itself.
- SPU_Status[S] set to '1': Two no-op instructions, followed by a branch instruction to itself.
- SPU_Status[S, I] set to '1': Illegal instruction, followed by a branch instruction to itself.
- SPU_Status[S, P] set to '1': Stop-and-Signal instruction with the stop code from the saved SPU_Status, followed by a branch instruction to itself.
- SPU_Status[P, H] set to '1': Unconditional halt instruction, followed by a Stop-and-Signal instruction with the saved stop code from the saved SPU_Status, followed by a branch instruction to itself.
- SPU_Status[P, I] set to '1': Illegal instruction, followed by a Stop-and-Signal instruction with the saved stop code from the SPU_Status, followed by a branch instruction to itself.

28. If all the SPU_Status bits[I, S, H, P, R] bits are set to '0' in the CSA (SPU was not executing), then add a branch instruction to itself to the end of the SPU context-restore sequence code, after the context-restore stop-and-signal (step 22 on page 379).
29. Restore the RA_Group_ID and RA_Enable registers from the CSA.
30. If the MFC_SR1[R] register field is set to '1':
    a. Write '0' to the SLB_Invalidate_All register.
    b. Write x'0' to SLB_Index register.
    c. Write the SLB_VSID register with the VSID to the SPU context-restore code.
    d. Write the SLB_ESID register with the ESID to the SPU context-restore code.
    e. Write '1' to SLB_Index register.
    f. Write the SLB_VSID register to provide access to the CSA.
    g. Write the SLB_ESID register to complete providing access to the CSA.
31. Write the SPU_NPC[IE] register field to '0' and SPU_NPC[LSA] to the LS entry-point address of the SPU context-restore code.
32. Write the SPU_Sig_Notify_1 register with the upper 32 bits of the CSA effective address.
33. Write the SPU_Sig_Notify_2 register with the lower 32 bits of the CSA effective address.
34. If the MFC_CNTL[Ds] is set to '1' in the CSA (decrementer was running), adjust the decrementer value in the LSCSA copy by the amount of time that elapsed between the context-save and the context-restore, as follows:
    a. Read the Time Base (TB) register value.
    b. Subtract the value from the time-base value saved in the CSA during the save sequence.
    c. Subtract that value from the decrementer value in the LSCSA copy.
    d. Save the result in the LSCSA copy.
    e. Write the software Decrementer Status Running flag in the LSCSA copy.
    f. If the updated decrementer value in the LSCSA copy is negative, then set the software Decrementer Wrapped flag in the CSA.
35. Copy the SPU_Out_Mbox data from the CSA into the CSA's copy of the LSCSA.

7. A unconditional halt instruction can be encoded in assembly as "heq r#,r#", where # is any register.

36. Copy the SPU_Out_Intr_Mbox data from the CSA into the CSA's copy of the LSCSA.
37. Issue a DMA command to copy the SPU context-restore sequence code to a portion of the low 16 KB of the LS, and start the SPU:
    a. Write the beginning LS address of the SPU context-restore sequence to the MFC_LSA register.
    b. Write the beginning CSA effective-address bits[0:32] of the SPU context-restore sequence code to the MFC_EAH register.
    c. Write the beginning CSA effective-address bits[33:63] of the SPU context-restore sequence code to the MFC_EAL register.
    d. Write the size of the SPU context-restore sequence code to the MFC_Size register.
    e. Write x'00' to the MFC_TagID register.
    f. Write '0048' to the MFC_ClassID_Cmd register (**gets** command, with TClassID and RClassID of '00').[8]
    g. Read the MFC_CmdStatus[Rc] register field and, if the value is not '00', go back to step 37a on page 375.
38. Write the Prxy_QueryMask register to '1' (enable Tag-Group 0), and issue a PPE **eieio** instruction.
39. Poll the Prxy_TagStatus[$g_0$] register field until it reads '01' (tag-group 0 complete), or write Prxy_QueryType[TS] with '01' and wait for a tag-group complete interrupt. Then, write the INT_Stat_class2[T] register field with '1'.
40. Poll the SPU_Status[R] register field until it reads '0', or wait for an SPU class 0 or class 2 interrupt. Then, reset Interrupt conditions by writing the INT_Stat_class0 register with '1' for each interrupt handled, and by writing the INT_Stat_class2[T] bit with '1'.
41. If the SPU_Status[P] register field is set to '1' and SPU_Status[StopCode] reads Restore Successful, the context-restore succeeded. Otherwise, the context-restore failed, and the sequence should be exited.
42. Restore the SPU_PrivCntl register from the CSA.
43. If any of the SPU_Status[I, S, H, P] bits are set to '1' in the CSA, restore the error or single-step state, as follows:
    a. Write the SPU_RunCntl[Run] register field to '01' (run request).
    b. Issue a PPE **eieio** instruction. Then, poll SPU_Status[R] until it is set to '0'.
    c. Go to step 45 on page 375.
44. If all of the SPU_Status[I, S, H, P, R] bits are set to '0' in the CSA, restore as follows:
    a. Write SPU_RunCntl[Run] to '01' (run request).
    b. Issue a PPE **eieio** instruction. Then, poll SPU_Status[R] until it reads '1' (running).
    c. Write SPU_RunCntl[Run] to '00' (stop request).
    d. Issue a PPE **eieio** instruction. Then, poll SPU_Status[R] until it reads '0' (stopped).
45. Issue a DMA command to restore the low 16 KB of LS from the CSA:
    a. Write x'0000_0000' to MFC_LSA register.

8. See *Section 12.5.2.2* on page 378 for the sequence to be copied and run.

b. Write the effective-address bits[0:32] for the beginning of the CSA's copy of the LS to the MFC_EAH register

c. Write the effective-address bits [33:63] for the beginning of the CSA's copy of the LS to the MFC_EAL register.

d. Write the value 16 KB to MFC_Size register.

e. Write x'00' to MFC_TagID register.

f. Write '0040' to the MFC_ClassID_Cmd register (**get** command, with TClassID and RClassID of '00').

g. Read the MFC_CmdStatus[Rc] register field. If the value is not equal to '00', go back to step 45a on page 375.

46. Poll the Prxy_TagStatus[$g_0$] register field until it reads '01' (tag-group 0 complete), or write the Prxy_QueryType[TS] register field to '01' and wait for tag-group complete interrupt. Then, write the INT_Stat_class2[T] register field to '1'.
47. Write the MFC_CNTL[Sc] register field to '1' and the MFC_CNTL[Sm] to '0'.
48. Poll MFC_CNTL[Ss] until it reads '11' (queue suspended).
49. Write all of the MFC_TLB_Invalidate_Entry[IS,VPN,L, Lp] register fields to zeros. Then, issue a PPE **sync** instruction.
50. If the OS is using interrupts instead of polling:
    a. Write the INT_Mask_class0 register with '0' (disable all class 0 interrupts).
    b. Write the INT_Mask_class1 register with '0' (disable all class 1 interrupts).
    c. Write the INT_Mask_class2 register with '0' (disable all class 2 interrupts).
    d. Write the INT_Stat_class0 register with -1 (reset all class 0 interrupts).
    e. Write the INT_Stat_class1 register with -1 (reset all class 1 interrupts).
    f. Write the INT_Stat_class2 register with -1 (reset all class 2 interrupts).
51. If the MFC_CNTL[Q] bit is set to '0' in the CSA (MFC command queues were not empty), then restore 96 doublewords to the MFC_CQ_SR registers from the CSA.
52. Restore the Prxy_QueryMask register from the CSA.
53. Restore the Prxy_QueryType register from the CSA.
54. Restore the MFC_CSR_TSQ register from the CSA.
55. Restore the MFC_CSR_Cmd1 and MFC_CSR_Cmd2 registers from the CSA.
56. Restore the MFC_CSR_ATO register from the CSA.
57. Restore the MFC_TClassID register from the CSA.
58. Set the lock-line reservation lost event:
    a. If SPU_Channel_0_Count is set to '0' in the CSA, and SPU_WrEventMask[Lr] is set to '1 in the CSA', and SPU_RdEventStat[Lr] is set to '0' in the CSA, then set SPU_RdEventStat_Count to '1' in the CSA.
    b. Set the SPU_RdEventStat[Lr] to '1' in the CSA to set lock-line reservation lost event.

59. If the status of the CSA software Decrementer Wrapped flag is set:
    a. If the SPU_Channel_0_Count is set to '0' in the CSA, and the SPU_WrEventMask[Tm] is set to '1' in the CSA, and the SPU_RdEventStat [Tm] is set to '0' in the CSA, then set the SPU_RdEventStat_Count to '1' in the CSA.
    b. Set the SPU_RdEventStat[Tm] to '1' in the CSA to set a decrementer event.
60. For the following channels:
    - SPU_RdEventStat (channel index 0)
    - SPU_RdSigNotify1 (channel index 3)
    - SPU_RdSigNotify2 (channel index 4)
    - MFC_RdTagStat (channel index 24)
    - MFC_RdListStallStat (channel index 25)
    - MFC_RdAtomicStat (channel index 27)
      a. Write the channel index to the SPU_ChnlIndex register, and issue a PPE **eieio** instruction.
      b. Restore channel data to the SPU_ChnlData register from the CSA.
      c. Restore channel count to the SPU_ChnlCnt register from the CSA.
61. For the following channels:
    - MFC_WrMSSyncReq (channel index 9, count = 1)
    - MFC_Cmd, (channel index 21, count from CSA)
    - MFC_WrTagUpdate, (channel index 23, count = 1)
      a. Write the channel index to the SPU_ChnlIndex register, and issue a PPE **eieio** instruction.
      b. Restore the channel count to the SPU_ChnlCnt register from the count.
62. Restore the SPU_LSLR register from the CSA.
63. Restore the SPU_Cfg register from the CSA.
64. Restore the PM_Trace_Tag_Wait_Mask register from the CSA.
65. Restore the SPU_NPC register from the CSA.
66. For the SPU_RdInMbox channel:
    a. Write the channel index value 29 to the SPU_ChnlIndex register, and issue a PPE **eieio** instruction.
    b. Restore the SPU_ChnlCnt register from the SPU_RdInMbox_Count in the CSA.
    c. Restore the SPU_ChnlData register from the CSA's SPU_RdInMbox data0 area.[9]
    d. Restore the SPU_ChnlData register from the CSA's SPU_RdInMbox data1 area.
    e. Restore the SPU_ChnlData register from the CSA's SPU_RdInMbox data2 area.
    f. Restore the SPU_ChnlData register from the CSA's SPU_RdInMbox data3 area.

9. This was saved in the CSA in step 43c on page 368.

67. If SPU_Mbox_Stat[P] is set to '0' in the CSA (mailbox empty), read from the SPU_Out_Mbox register (discarding the data).
68. If SPU_Mbox_Stat[I] is set to '0' in the CSA (interrupt mailbox empty):
    a. Read from the SPU_Out_Intr_Mbox register (discarding the data).
    b. Write the INT_Stat_Class2[M] register field to '1' (reset mailbox interrupt status).
69. If MFC_SR1[R] is set to '1' in the CSA:
    a. Restore the SLB_Index register from the CSA.
    b. For index=0 through index=7:
        1. Write the index value to the SLB_Index register, and issue a PPE **eieio** instruction.
        2. Restore the SLB_ESID register from SLB_ESID[index] in the CSA.
        3. Restore the SLB_VSID register from SLB_VSID[index] in the CSA.
    c. Restore the SLB_Index register from the CSA.
70. Restore the MFC_SR1 register from the CSA.
71. Restore other SPE mappings to the SPE.
72. If SPU_Status[R] is set to '1' in the CSA (running), write the SPU_RunCntl[Run] register field to '01'.
73. Restore the MFC_CNTL register from the CSA.
74. Restore mapping of the PPE-side problem-state access to the SPE resources.
75. Reset the Context Switch Active flag.
76. Re-enable SPE interrupts:
    - Restore the INT_Mask_class0 register from the CSA.
    - Restore the INT_Mask_class1 register from the CSA.
    - Restore the INT_Mask_class2 register from the CSA.
77. If multithreading is being used, release software SPU mutual-exclusion lock.

#### 12.5.2.2 *SPU Context-Restore Sequence*

The following context-restore sequence is copied by the PPE into a portion of the low 16 KB of the LS and run in step 37f on page 375:

1. Set the SPU_WrEventMask channel to '0' to mask all events.
2. Set the MFC_WrTagMask channel to '01' to unmask only Tag-Group 0.
3. Read the SPU_RdSigNotify1 channel data to obtain the high address of the CSA effective address.
4. Write the high address from step 3 to the MFC_EAH channel.
5. Read the SPU_RdSigNotify2 channel data to obtain the low address of the CSA effective address.
6. Update the low effective addresses in a predefined, 15-element (16 KB each) DMA list that restores the most-significant 240 KB of the LS from the CSA.

7. Enqueue a **get** (Tag-Group 0) command to restore the LSCSA copy from the CSA to the LSCSA (in low LS memory).
8. Enqueue a **getl** (Tag-Group 0) command for the DMA list updated in step 6 on page 378.
9. Write the `MFC_WrTagUpdate` channel with a value of '01' to update the tag status on any completion.
10. Read the `MFC_RdTagStat` channel data. This stalls the SPU until the DMA Tag-Group 0 instruction is complete.
11. Issue quadword loads for the high 112 GPRs from the LSCSA.
12. If the LSCSA software Decrementer Running flag is set, then write the `SPU_WrDec` channel with the decrementer value from the LSCSA.
13. Enqueue a **putllc** command using an effective address in the CSA, to remove any possible lock-line reservation.
14. Read the `MFC_RdAtomicStat` channel data. This stalls the SPU until the **putllc** completes.
15. Write the `SPU_WrOutMbox` channel with the `SPU_WrOutMbox` data from the LSCSA.
16. Write the `SPU_WrOutIntrMbox` channel with the `SPU_WrOutIntrMbox` data from the LSCSA.
17. Restore the Floating-Point Status and Control Register (FPSCR) from the LSCSA using the **fscrwr** instruction.
18. Restore the Save and Restore Register 0 (`SRR0`) data from the LSCSA and write to the `SPU_WrSSR0` channel.
19. Read the `SPU_RdEventMask` data from the LSCSA and write to the `SPU_WrEventMask` channel.
20. Read the `MFC_RdTagMask` data from the LSCSA and write to the `MFC_WrTagMask` channel.
21. Restore the low 16 GPRs from the LSCSA using quadword loads.
22. Issue a stop-and-signal (**stop**) instruction with a success status.
23. There might be additional instructions placed here to restore the correct stopped state (see steps 27 and 28 on page 373).

If there is a software-detected error in the SPU Context-Restore Sequence, the `SPU_RdInMbox` channel can be written with a diagnostic status. A halt conditional instruction should then be issued, followed by a branch instruction to itself.

## 12.6 Considerations for Hypervisors

The preceding sequences assume an OS that is running at both privilege 1 and privilege 2 levels—that is, an operating environment without a separate hypervisor and without support for logical-partition switching. If a hypervisor is used and the SPE to be switched is allocated to a logical partition, an additional state might need to be saved, set, and restored. This state might include some or all of the following:

- MFC Logical Partition ID Register (`MFC_LPID`).
- Interrupt Routing Register (`INT_Route`).
- MFC Storage Description Register (`MFC_SDR`).
- SMM Hardware Implementation Dependent Register (`SMM_HID`).

- MFC Real-Mode Address Boundary Register (`MFC_RMAB`).
- 256 TLB Entries.
- MFC TLB Index Register (`MFC_TLB_Index`).
- MFC TLB Real-Page Number Register (`MFC_TLB_RPN`).
- MFC TLB Virtual-Page Number Register (`MFC_TLB_VPN`).
- MFC TLB Replacement Management Table Data Register (`MFC_TLB_RMT_Data`).
- MFC command request order is lost.
- DMA and register lock state for address translation exceptions is lost.

See *Section 11 Logical Partitions and a Hypervisor* on page 331 for details about these resources.

# 13. Time Base and Decrementers

## 13.1 Introduction

The time-base facility provides the timing functions for the processor core-clock (NClk) domain. The facility includes two software-visible 64-bit time-base registers—one for configuration and one for counting—and eleven software-visible 32-bit decrementers—three in the PowerPC Processor Element (PPE) (two for user software, and one for hypervisor software) and one in each of the eight Synergistic Processor Elements (SPEs) (for user software).

The counting time-base (called the TB register) and decrementer registers count time in *ticks*, at a ticking rate called the *time-base frequency*, regardless of the Slow-State setting[1]. The TB register and all decrementers are updated at the same time-base frequency. The TB register provides a long-period counter; it holds a monotonically increasing value, and it is split into upper and lower halves. the 32-bit decrementers support shorter-period counting and most of them provide interrupt-signaling after a specified amount of time has elapsed.

The TB register and the four PPE decrementers are implemented in the Cell Broadband Engine Architecture (CBEA) processors[2] as PPE Special-Purpose Registers (SPRs). All of these registers are volatile and must be initialized during startup. The eight SPE decrementers are implemented in the CBEA processors as synergistic processor unit (SPU) channel registers that are accessed by SPE software using channel instructions (**rdch** or **wrch**).

## 13.2 Time-Base Facility

### 13.2.1 Clock Domains

The CBEA processors have three clock domains, each running asynchronously to the other two domains:

- *Processor Core Clock (NClk)*—This clock times the PowerPC processor unit (PPU), the SPUs, and parts of the PowerPC processor storage subsystem (PPSS) and the memory flow controllers (MFCs). The processor core clock (NClk) is occasionally referred to as the core clock (CORE_CLK) or core clock frequency (CCF).
- *MIC Clock (MiClk)*—This clock times the memory interface controller (MIC).
- *BIC Core Clock (BClk)*—This clock times the bus interface controller (BIC), which is part of the Cell Broadband Engine interface (BEI) unit to the I/O interface.

The following logic blocks run at half the processor core clock frequency (NClk/2):

- Element interconnect bus (EIB) and interfaces to the EIB (parts of the PPSS and MFCs)
- I/O Interface controller (IOC)
- MIC logic that is part of the CBEA processor core
- BIC logic that is part of the CBEA processor core

1. See *Section 15.1.1 Slow State* on page 430 for information about the Slow State.
2. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

- Pervasive logic, which is the logic that provides power management, thermal management, clock control, software-performance monitoring, trace analysis, and so forth

The time-base and decrementer facilities described in this section fall entirely within the processor core clock (`NClk`) domain.

### 13.2.2 Time-Base Registers

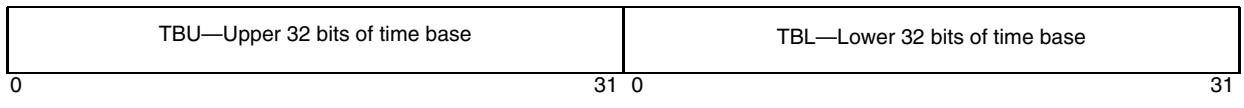
There are two registers with the name “Time Base Register” but they have different acronyms:

- *Time Base Register (TB)*—This SPR register contains the time-base count value, as described in *Section 13.2.3* on page 383. The physical implementation of this register actually includes the following SPR registers:
  - Time Base Register, Read Only (TB), SPR #268
  - Time Base Upper Register, Read Only (TBU), SPR #269
  - Time Base Lower Register, Write Only (TBL), SPR #284
  - Time Base Upper Register, Write Only (TBU), SPR #285
- *Time Base Register (TBR)*—This memory-mapped I/O (MMIO) register specifies the time-base sync mode, as described in *Section 13.2.4* on page 384, and the internal reference-clock divider setting, as described in *Section 13.2.3* on page 383.

For details on the TBR register, see the *Cell Broadband Engine Registers* specification.

The TB register, shown in *Figure 13-1*, is a 64-bit structure that contains a 64-bit unsigned integer. The register is divided into two halves, a 32-bit TB Upper Register (TBU) and a 32-bit TB Lower Register (TBL). Each tick pulse to the TB register adds ‘1’ to the low-order bit, bit[31], of TBL.

*Figure 13-1. Time Base (TB) Register*




The TB register can be read by user or supervisor software but written only by hypervisor software. The register must be software-enabled to allow updates. Set SPR HID6[tb_enable] to 1 to enable the time-base facility. Software reads and writes of the TB register are always allowed, regardless of SPR HID6[tb_enable]. Writes take precedence over updates (increments). Updates are ignored if they collide with a write operation.

The TB register increments until its value becomes x‘FFFFFFFF_FFFFFFFF’ ($2^{64} – 1$). At the next increment its value becomes x‘00000000_00000000’. There is no exception or explicit indication when this occurs. If the update frequency of the TB register is constant, the register can be used as a source of values that increase at a constant rate, such as for time stamps in trace entries. Even if the update frequency is not constant, values read from the TB register are monotonically increasing (except when the value in the TB register wraps from $2^{64} – 1$ to 0). If a trace entry is recorded each time the update frequency changes, the sequence of TB-register values can be postprocessed to become actual time values.

### 13.2.3 Time-Base Frequency

Each tick pulse to the TB register causes it to increment by one, if it is software-enabled to allow updates. There are two modes for updating the TB register:

- *Internal Time-Base Sync Mode:* In this mode, a phase-locked loop (PLL) reference clock (PLL_REFCLK) external input signal and a core clock multiplier (CCM) are used to synchronize and increment the TB register. See *Section 13.2.4.1* on page 384 for details.
- *External Time-Base Sync Mode:* In this mode, the rising edge of the Time-Base Enable (TBEN) external input signal is used to synchronize and increment the TB register. See *Section 13.2.4.3* on page 387 for details.

In the internal time-base sync mode, the processor core clock (NClk) frequency is:

$$NClk = PLL\_REFCLK \times CCM$$

The maximum time-base frequency is constrained by the following factors:

- PLL reference clock (PLL_REFCLK) frequency.
- Core clock multiplier (CCM), also called the PLL multiplier setting (PLLmultiplier).
- Power-management slow state (called *slow mode* in parameters described in this section). For details, see *Section 15.1.1 Slow State* on page 430.

The equations for maximum time-base frequency (Fmax) are:

$$RefDiv \geq \lceil MaxSlowModeNclkDivider \times 11 / CCM \rceil$$

$$Fmax\ (\text{Internal Mode}) = PLL\_REFCLK / RefDiv$$

$$Fmax\ (\text{External Mode}) \leq PLL\_REFCLK / (1 + (11 \times MaxSlowModeNclkDivider / CCM))$$

The preceding equations account for the worst-case Slow State setting, which guarantees that the time-base frequency remains constant during Slow State and during transitions between Slow States.

The term MaxSlowModeNClkDivider denotes the maximum processor core clock (NClk) frequency divider in Slow State configured by the operating system. The Power Management hardware supports eight different Slow-State dividers (also called slow-mode dividers): 0 through 7. The MaxSlowModeNClkDivider values for each of these settings are 1, 2, 3, 4, 5, 6, 8, and 10 respectively. The value is set in the Power Management Control Register (PMCR). See PMSR[BE_Slow_n] for the current Slow State setting, and see *Section 15.1.1 Slow State* on page 430 for information about how to change power management states.

The standard 3.2 GHz processor core clock (NClk) is configured by specifying the Internal Time-Base Sync State, a 400 MHz PLL reference clock (PLL_REFCLK) frequency, and a core clock multiplier (CCM) of x8.

If the operating system uses all the Slow State settings—that is to say, the MaxSlowModeNClkDivider is 10—the maximum time-base frequency is 28.57 MHz in internal time-base sync mode and 27.11 MHz in external time-base sync mode. The time-base frequency of such a system is limited to 200 MHz with Slow State disabled and in internal time-base sync mode.

The maximum time-base frequency limit guarantees a constant time-base frequency during functional operation of the processor (including during all slow states). The time base should be set to a frequency below this limit. The time base should not be changed again in response to changes to the slow mode setting. Changes to the time-base frequency affect the operation and frequency of all time-base facilities including: the time base, PPE DEC 0, PPE DEC 1, HDEC, WDEC, and the SPU decrementers. If it is necessary to change the time-base frequency, all time-base facilities need to be reinitialized.

### 13.2.4 Time-Base Sync Mode Controls

The SPR HID6[tb_enable] bit enables the time-base facility. After enabling, the time-base sync mode (internal or external) can be selected in the Time Base MMIO Register (TBR) using the TBR[Timebase_mode] bit. In the internal time-base sync mode, the TB register should be initialized and TBR[Timebase_setting] is used to specify the required time-base frequency, up to the maximum, Fmax.

#### 13.2.4.1 *Internal Time-Base Sync Mode*

In internal time-base sync mode, TBR[Timebase_setting] is set according to the maximum time-base calculation. The value in TBR[Timebase_setting] can vary from 0 to 255. If it is ‘0’, the internal time-base sync mode is disabled and the TB register is not counting. If it is any other value, the TBR[Timebase_setting] value is used to determine the RefDiv value, as described in *Section 13.2.4.2* on page 387.

Given our previous example of a 400 MHz PLL reference clock (PLL_REFCLK) and x8 core clock multiplier, in which all Slow State settings are supported, then:

$$RefDiv \geq ceil(10 \times 11 / 8) = 14.$$
$$Fmax = 400\ MHz / 14 = 28.57\ MHz.$$

The maximum time-base frequency of 28.57 MHz is configured by setting TBR[Timebase_setting] to x‘94’. This value is determined by finding the Timebase_setting entry in *Table 13-1* on page 385 equal to RefDiv of 14. If a slower time-base is required—for example 5 MHz (a RefDiv value of 80)—then the TBR[Timebase_setting] register field is set to x‘A7’.

To enable internal time-base sync mode operation, follow these steps:

1. Initialize the TB register.
2. Set TBR[timebase_mode] to ‘1’, and TBR[Timebase_setting] to ‘0’.
3. Set TBR[timebase_mode] to ‘0’, and TBR[Timebase_setting] to the value corresponding to the required RefDiv value (see *Table 13-1* on page 385).

A timing diagram of steps 2 and 3 of this sequence is shown in *Figure 13-2* on page 385. In a multiple-CBEA-processor system, TBEN can be used to synchronize the start of updates to TB registers across all CBEA processors. As shown in *Figure 13-2*, the time base of each CBEA processor does not start until after TBEN becomes ‘1’. In a single-CBEA-processor system, TBEN can be tied high and the TB register starts updating after step 2.

*Figure 13-2. Internal Time-Base Sync Mode Initialization Sequence*

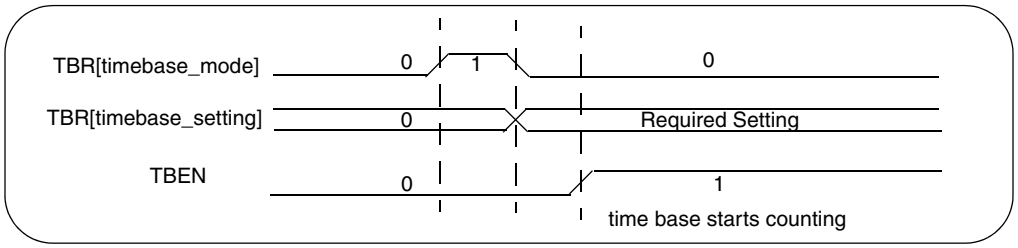



*Table 13-1. TBR[Timebase_setting] for Reference Divider (RefDiv) Setting* (Sheet 1 of 3)

| RefDiv | Timebase_setting | RefDiv | Timebase_setting | RefDiv | Timebase_setting | RefDiv | Timebase_setting |
|---|---|---|---|---|---|---|---|
| 0[1] | 00 | 64 | 7B | 128 | 86 | 192 | 37 |
| 1 | FF | 65 | 3D | 129 | 43 | 193 | 9B |
| 2 | 7F | 66 | 9E | 130 | 21 | 194 | 4D |
| 3 | 3F | 67 | CF | 131 | 10 | 195 | 26 |
| 4 | 9F | 68 | 67 | 132 | 88 | 196 | 93 |
| 5 | 4F | 69 | B3 | 133 | 44 | 197 | 49 |
| 6 | 27 | 70 | D9 | 134 | A2 | 198 | 24 |
| 7 | 13 | 71 | EC | 135 | D1 | 199 | 92 |
| 8 | 09 | 72 | 76 | 136 | E8 | 200 | C9 |
| 9 | 84 | 73 | BB | 137 | 74 | 201 | 64 |
| 10 | 42 | 74 | DD | 138 | BA | 202 | 32 |
| 11 | A1 | 75 | EE | 139 | 5D | 203 | 19 |
| 12 | 50 | 76 | 77 | 140 | AE | 204 | 0C |
| 13 | 28 | 77 | 3B | 141 | D7 | 205 | 06 |
| 14 | 94 | 78 | 9D | 142 | EB | 206 | 03 |
| 15 | CA | 79 | 4E | 143 | F5 | 207 | 81 |
| 16 | E5 | 80 | A7 | 144 | 7A | 208 | C0 |
| 17 | F2 | 81 | 53 | 145 | BD | 209 | E0 |
| 18 | F9 | 82 | A9 | 146 | DE | 210 | 70 |
| 19 | 7C | 83 | 54 | 147 | 6F | 211 | B8 |
| 20 | BE | 84 | 2A | 148 | B7 | 212 | 5C |
| 21 | 5F | 85 | 95 | 149 | DB | 213 | 2E |
| 22 | AF | 86 | 4A | 150 | ED | 214 | 97 |
| 23 | 57 | 87 | A5 | 151 | F6 | 215 | 4B |
| 24 | AB | 88 | 52 | 152 | FB | 216 | 25 |

1. This setting disables internal time-base sync mode.

*Table 13-1. TBR[Timebase_setting] for Reference Divider (RefDiv) Setting* (Sheet 2 of 3)

| RefDiv | Timebase_setting | RefDiv | Timebase_setting | RefDiv | Timebase_setting | RefDiv | Timebase_setting |
|---|---|---|---|---|---|---|---|
| 25 | 55 | 89 | 29 | 153 | 7D | 217 | 12 |
| 26 | AA | 90 | 14 | 154 | 3E | 218 | 89 |
| 27 | D5 | 91 | 8A | 155 | 1F | 219 | C4 |
| 28 | EA | 92 | 45 | 156 | 0F | 220 | E2 |
| 29 | 75 | 93 | 22 | 157 | 87 | 221 | 71 |
| 30 | 3A | 94 | 91 | 158 | C3 | 222 | 38 |
| 31 | 1D | 95 | 48 | 159 | 61 | 223 | 1C |
| 32 | 0E | 96 | A4 | 160 | B0 | 224 | 8E |
| 33 | 07 | 97 | D2 | 161 | 58 | 225 | 47 |
| 34 | 83 | 98 | 69 | 162 | 2C | 226 | 23 |
| 35 | C1 | 99 | B4 | 163 | 96 | 227 | 11 |
| 36 | 60 | 100 | 5A | 164 | CB | 228 | 08 |
| 37 | 30 | 101 | 2D | 165 | 65 | 229 | 04 |
| 38 | 18 | 102 | 16 | 166 | B2 | 230 | 02 |
| 39 | 8C | 103 | 8B | 167 | 59 | 231 | 01 |
| 40 | 46 | 104 | C5 | 168 | AC | 232 | 80 |
| 41 | A3 | 105 | 62 | 169 | D6 | 233 | 40 |
| 42 | 51 | 106 | 31 | 170 | 6B | 234 | A0 |
| 43 | A8 | 107 | 98 | 171 | B5 | 235 | D0 |
| 44 | D4 | 108 | CC | 172 | DA | 236 | 68 |
| 45 | 6A | 109 | E6 | 173 | 6D | 237 | 34 |
| 46 | 35 | 110 | 73 | 174 | B6 | 238 | 1A |
| 47 | 9A | 111 | 39 | 175 | 5B | 239 | 8D |
| 48 | CD | 112 | 9C | 176 | AD | 240 | C6 |
| 49 | 66 | 113 | CE | 177 | 56 | 241 | E3 |
| 50 | 33 | 114 | E7 | 178 | 2B | 242 | F1 |
| 51 | 99 | 115 | F3 | 179 | 15 | 243 | 78 |
| 52 | 4C | 116 | 79 | 180 | 0A | 244 | BC |
| 53 | A6 | 117 | 3C | 181 | 05 | 245 | 5E |
| 54 | D3 | 118 | 1E | 182 | 82 | 246 | 2F |
| 55 | E9 | 119 | 8F | 183 | 41 | 247 | 17 |
| 56 | F4 | 120 | C7 | 184 | 20 | 248 | 0B |
| 57 | FA | 121 | 63 | 185 | 90 | 249 | 85 |
| 58 | FD | 122 | B1 | 186 | C8 | 250 | C2 |
| 59 | 7E | 123 | D8 | 187 | E4 | 251 | E1 |

1. This setting disables internal time-base sync mode.

*Table 13-1. TBR[Timebase_setting] for Reference Divider (RefDiv) Setting* (Sheet 3 of 3)

| RefDiv | Timebase_setting | RefDiv | Timebase_setting | RefDiv | Timebase_setting | RefDiv | Timebase_setting |
|---|---|---|---|---|---|---|---|
| 60 | BF | 124 | 6C | 188 | 72 | 252 | F0 |
| 61 | DF | 125 | 36 | 189 | B9 | 253 | F8 |
| 62 | EF | 126 | 1B | 190 | DC | 254 | FC |
| 63 | F7 | 127 | 0D | 191 | 6E | 255 | FE |

1. This setting disables internal time-base sync mode.

#### 13.2.4.2 *Determining the System Time-Base Frequency*

Assuming the time base is configured as internal time-base sync mode, then the frequency can be determined using the following procedure:

1. Verify that the system is configured in internal time-base sync mode by reading the TBR register and verifying that the TBR[Timebase_mode] field is '0'.
2. Determine the RefDiv value by doing a inverse table lookup of *Table 13-1* using the value stored in TBR[Timebase_setting]. For example, if the TBR[Timebase_setting] value is x‘2A’, then the RefDiv is 84. If TBR[Timebase_setting] is 0, then the time base is disabled.
3. Compute the time-base frequency by multiplying the PLL_REFCLK and RefDiv.

The number of cycles, either PPE or SPE cycles, per time-base tick is further reduced based upon the current Slow State setting, PMSR[BE_Slow_n].

```
cycles per time-base tick = PLL_REFCLK × RefDiv / SlowModeFactor.
```

where the SlowModeFactor is either 1, 2, 3, 4, 5, 6, 8, or 10 corresponding to the PMSR[BE_Slow_n] values of 0 through 7.

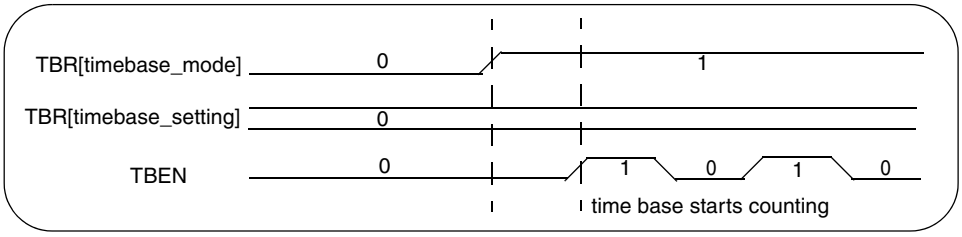
Because the time-base registers are accessible only by privileged software, the operating environment should provide information about the time-base frequency. Consult system documentation for specific details.

#### 13.2.4.3 *External Time-Base Sync Mode*

To select the external time-base sync mode, set TBR[timebase_mode] to ‘1’. The value of TBR[Timebase_setting] has no effect, although the no effect, but the recommended value is x‘00’.

In the external time-base sync mode, the rising edge of the TBEN signal is used to generate the TB-register ticks. Similar to internal time-base sync mode, TBEN can be used to synchronize the start of the time base across different CBEA processors in a multi-CBEA-processor system. A timing diagram is shown in *Figure 13-3* on page 388. In a single-CBEA-processor system, TBEN can be active before TBR[timebase_mode] is set to ‘1’.

*Figure 13-3. External Time-Base Sync Mode Initialization Sequence*




### 13.2.5 Reading and Writing the TB Register

#### 13.2.5.1 *Reading the TB Register*

Either user or supervisor software can read the contents of the TB register into a general-purpose register (GPR). To read the contents of the TB register into register **r**A, execute the following instruction:

```
mftb    rA
```

The preceding example uses the simplified (extended) mnemonic form of the **mftb** instruction (equivalent to **mftb r**A,**268**). Using this instruction copies the entire TB register (TBU || TBL) into **r**A. Reading the TB register has no effect on the value it contains or the periodic incrementing of that value.

If the simplified mnemonic form **mftbu r**A (equivalent to **mftb r**A**,269**) is used, the contents of TBU are copied to the low-order 32 bits of **r**A, and the high-order 32 bits of **r**A are cleared (0 || TBU).

#### 13.2.5.2 *Writing the TB Register*

Only hypervisor software can write the contents of a GPR into the TB register. The simplified (extended) mnemonics, **mttbl** and **mttbu**, write the lower and upper halves of the TB register, respectively. The simplified mnemonics are for the **mtspr** instruction. The **mtspr**, **mttbl**, and **mttbu** instructions treat TBL and TBU as separate 32-bit registers; setting one leaves the other unchanged. It is not possible to write the entire 64-bit TB register in a single instruction.

The instructions for writing the TB register are not dependent on the implementation or mode. Thus, code written to set the TB register on a 32-bit implementation will work correctly on the CBEA processors.

The TB register can be written by a sequence such as:

```
lwz     rx,upper                #load 64-bit value for
lwz     ry,lower                # TB into rx and ry
li      rz,0
mttbl   rz                      #force TBL to 0
mttbu   rx                      #set TBU
mttbl   ry                      #set TBL
```

Provided that no exceptions occur while the last three instructions are being executed, loading 0 into TBL prevents the possibility of a carry from TBL to TBU while the TB register is being initialized.

### 13.2.6 Computing Time-of-Day

Because the update frequency of the TB register is system-dependent, the algorithm for converting the current value in the TB register to time-of-day is also system-dependent.

In a system in which the update frequency of the TB register can change over time, it is not possible to convert an isolated time-base value into time of day. Instead, a time-base value has meaning only with respect to the current update frequency and the time of day at which the update frequency was last changed.

Each time the update frequency changes, either the system software is notified of the change by means of a decrementer interrupt, or the change was instigated by the system software itself. At each such change, the system software must compute the current time of day using the old update frequency, compute a new value of ticks-per-second for the new frequency, and save the time of day, time-base value, and tick rate. Subsequent calls to compute time of day use the current time-base value and the saved data.

A generalized service to compute time-of-day can take the following as input:

- Time-of-day at beginning of current epoch
- Time-base value at beginning of current epoch
- Time-base update frequency
- Time-base value for which time-of-day is required

For a CBEA processor system in which the time-base update frequency does not vary, the first three inputs are constant.

## 13.3 Decrementers

The CBEA processors contain eleven software-visible decrementers, three in the PPE and one in each SPE. All decrementers tick at the same frequency as the same time-base frequency, subject to control commands which turn them on and off.

### 13.3.1 PPE Decrementers

The PPE has three 32-bit software-visible decrementers:

- Two decrementers (DEC), one per PPE thread, accessible to supervisor software
- One hypervisor decrementer (HDEC), accessible to hypervisor software

Each update tick to the PPE decrementers causes their values to decrease by one, at the time-base frequency. SPR `HID6[tb_enable]` must be software-enabled to allow PPE decrementer updates.

The DEC decrementers support short-term counting and provide a means of signaling an interrupt after a specified amount of time has elapsed, unless the decrementer is altered by software in the interim, or the time-base frequency changes.

The HDEC hypervisor decrementer provides a means for hypervisor software to manage timing functions independently of the decrementers, which are managed by user or supervisor software running in virtual partitions. Similar to the decrementer, the HDEC is a counter that is updated at the time-base frequency, and it provides a means of signaling an interrupt after a specified amount of time has elapsed. For more about HDEC, see *Section 11 Logical Partitions and a Hypervisor* on page 331.

#### 13.3.1.1 *Decrementer Operation*

The DEC counts down, causing a decrementer interrupt (unless masked by MSR[EE]) when it passes through '0'. The DEC satisfies the following requirements:

- The operations of the TB and DEC registers are driven by the same time-base frequency.
- Loading a GPR from the DEC has no effect on the DEC.
- Storing the contents of a GPR to the DEC replaces the value in the DEC with the value in the GPR.
- Whenever bit[0] of the DEC is set to '1', a decrementer interrupt is signaled. Multiple decrementer interrupts can be received before the first interrupt occurs; however, any additional requests are canceled when the interrupt occurs for the first request.
- If the DEC is altered by software and the content of bit [0] is set to '1', a decrementer interrupt is signaled.

In systems that change the time-base frequency for purposes such as power management, the decrementer input frequency will also change. Software must be aware of this to set interval timers.

#### 13.3.1.2 *Writing and Reading the Decrementers*

The contents of the DEC can be read or written using the **mfspr** and **mtspr** instructions, both of which are supervisor instructions when they refer to the DEC. Using a simplified (extended) mnemonic for the **mtspr** instruction, the DEC can be written from GPR **r**A with the following:

```
mtdec    rA
```

Using a simplified (extended) mnemonic for the **mfspr** instruction, the DEC can be read into GPR **r**A with the following:

```
mfdec    rA
```

### 13.3.2 SPE Decrementers

#### 13.3.2.1 *Decrementer Operation*

Each SPU contains one 32-bit decrementer implemented as a down-counter. All decrementers in the SPUs and the PPE count down at the same rate. An SPU decrementer is accessed through two channels—the SPU Write Decrementer Channel (SPU_WrDec) and the SPU Read Decrementer Channel (SPU_RdDec)—as described in *Section 17.4 SPU Decrementer* on page 454.

An SPU decrementer can be run, read, and halted by SPU channel instructions, which can also enable or disable decrementer events. A decrementer event becomes pending when the value in the decrementer changes from '0' to negative (the most-significant bit changes from '0' to '1'). See *Section 18.6.9 Procedure for Handling the SPU Decrementer Event* on page 489 for details on event-interrupt handling.

#### 13.3.2.2 *Writing and Reading the Decrementers*

The SPU decrementer is accessed by reading and writing SPU channels using the **rdch** and **wrch** instructions or using the **spu_readch** and **spu_writech** intrinsics. The decrementer value is read from the SPU Read Decrementer Channel, and the SPU Write Decrementer Channel allows the decrementer value to be set. In addition, the SPU Write Event Mask Channel and the SPU Write Event Acknowledgment Channel are used to turn the decrementer on and off, and to handle decrementer events.

For more information about the SPU decrementer channels, see *Section 17 SPE Channel and Related MMIO Interface* on page 447.

### 13.3.3 Using an SPU Decrementer to Monitor SPU Code Performance

The following example demonstrates the use of an SPU decrementer to monitor code performance. In the example, profiling macros are inserted into the source code around the sections of code to be timed. The resulting timing information is output to the console, but it could instead be read and processed by an application program.

#### 13.3.3.1 *Sample Performance-Profiling Code*

The following sample code illustrates a general approach to performance profiling of the SPEs. The sample uses a profile-checkpoint facility which allows timing of code sections. Source code to be timed can be instrumented with macros, such as the `prof_cp(N)` macros shown here (where N is a number) that resolve to special no-op instructions that the profiling program interprets as commands to output instruction-count and cycle-count information. The `prof_clear()`, `prof_start()`, and `prof_stop()` macros have additional meaning: `prof_clear()` clears the performance information, `prof_start()` starts recording of performance data, and `prof_stop()` turns off recording of the data.

The following shows the use of the profiling macros:

```
#include "profile.h"

...
prof_clear();
prof_start();
...
<code to be timed>
...
prof_stop();
```

Additional `prof_cp(N)` macros can be inserted in the timed section for finer-grained timing information.

#### 13.3.3.2 *Hardware Code Performance*

Because the granularity of the decrementer-based profiling implementation is limited by the frequency of the SPU decrementer, the code sections to be timed should typically execute several thousand SPU cycles to obtain reasonable accuracy. The implementation is not suitable for extremely fine-grained timing of SPU code.

#### 13.3.3.3 *SPU Decrementer Profiling*

In the following code sample, the total number of elapsed decrementer ticks is kept in a global variable called `_count` (each SPU has its own copy). Another global variable, `_count_base`, keeps track of the last value read from the decrementer. To allow timing to be turned on and off, a boolean variable called `_counting` keeps track of whether or not profiling is enabled. These variables are defined as:

```
extern volatile unsigned long long _count

extern volatile unsigned int _count_base

extern volatile _Bool _counting;
```

In the body of the `prof_cp()` macro, if profiling is enabled, the decrementer value is read from the SPU Read Decrementer Channel, using the `spu_readch` intrinsic, and subtracted from the `_count_base`. This gives the elapsed number of decrementer ticks since the last profiling-macro invocation (the decrementer counts down with each tick). If the decrementer value is greater than the `_count_base`, an underflow condition has occurred; underflow can also occur if end is less than `_count_base`, but ignore that situation for the moment.

Assuming a single decrementer underflow, the elapsed number of ticks is given by `x'FFFF FFFF' + _count_base - end`. The global `_count` is incremented by the elapsed ticks calculated previously, and a message displaying the count is printed to the console output. Finally, to avoid counting the time spent in the profiling macro itself, the decrementer is read once again and assigned to the `_count_base` variable.

The following code sample implements the routine described in the preceding paragraphs:

```
#define DECR_MAX   0xFFFFFFFF
#define DECR_COUNT DECR_MAX

#define prof_cp(cp_num)                                    \
{                                                          \
   if (_counting) {                                        \
     unsigned int end = spu_readch(SPU_RdDec);             \
     _count += (end > _count_base) ?                       \
       (DECR_MAX + _count_base - end) :                    \
       (_count_base - end);                                \
   }                                                       \
   printf("SPU#: CP"#cp_num", %llu\n", _count);            \
   _count_base = spu_readch(SPU_RdDec);                    \
}
```

The `prof_clear()` macro is implemented similarly, except that it also clears the timing information after displaying it:

```
#define prof_clear()                                          \
{                                                             \
   if (_counting) {                                           \
     unsigned int end = spu_readch(SPU_RdDec);                \
     _count += (end > _count_base) ?                          \
       (DECR_MAX + _count_base - end) :                       \
       (_count_base - end);                                   \
   }                                                          \
   printf("SPU#: clear, %llu\n", _count);                     \
   _count = 0;                                                \
    _count_base = spu_readch(SPU_RdDec);                      \
}
```

The `prof_start()` macro performs the additional work of enabling the decrementer. This is done by using an `spu_writech` intrinsic to write a new decrementer count to the SPU Write Decrementer Channel, and another `spu_writech` to the SPU Write Event Mask Channel to enable decrementer events:

```
#define prof_start()                                          \
{                                                             \
   if (_counting) {                                           \
     unsigned int end = spu_readch(SPU_RdDec);                \
     _count += (end > _count_base) ?                          \
       (DECR_MAX + _count_base - end) :                       \
       (_count_base - end);                                   \
   }                                                          \
   printf("SPU#: start, %llu\n", _count);                     \
   spu_writech(SPU_WrDec, DECR_COUNT);                        \
   spu_writech(SPU_WrEventMask, MFC_DECREMENTER_EVENT);       \
   _counting = 1;                                             \
    _count_base = spu_readch(SPU_RdDec);                      \
}
```

Finally, the `prof_stop()` macro disables counting by writing the SPU Write Event Mask Channel to disable decrementer events, and writing the SPU Write Event Acknowledgment Channel to acknowledge pending events and stop the decrementer, again using the `spu_writech` intrinsic:

```
#define prof_stop()                                           \
{                                                             \
   if (_counting) {                                           \
     unsigned int end = spu_readch(SPU_RdDec);                \
     _count += (end > _count_base) ?                          \
       (DECR_MAX + _count_base - end) :                       \
       (_count_base - end);                                   \
   }                                                          \
   printf("SPU#: stop, %llu\n", _count);                      \
```

```
   _counting = 0;                                          \
   spu_writech(SPU_WrEventMask, 0);                        \
   spu_writech(SPU_WrEventAck, MFC_DECREMENTER_EVENT);     \
   _count_base = spu_readch(SPU_RdDec);                    \
}
```

Commonality can be exploited by defining shared macros for the global `_count` and `_count_base` updates that appear in all of the profiling macros and in the interrupt handler.

#### 13.3.3.4 *Decrementer Underflow Handling*

If the code sections to be profiled are known to execute in less than $2^{32}$ decrementer ticks, the simple implementation shown previously is sufficient for providing timing information. For longer-running code sections, the profiling implementation needs to handle cases in which the decrementer underflows multiple times between profiling-macro invocations. This can be done by installing an interrupt handler that executes whenever the decrementer reaches 0, which triggers an SPU decrementer event.

In the body of the handler, the handler checks for the SPU decrementer event and increments the global `_count` as shown previously. It then handles the event by clearing the SPU decrementer-event bit in the SPU Write Event Mask Channel using the `spu_writech` intrinsic. An `spu_writech` to the SPU Write Event Acknowledgment Channel acknowledges the decrementer event and also causes the decrementer to stop. The decrementer is then restarted by writing a new value to the SPU Write Decrementer Count Channel and setting the SPU decrementer event bit in the SPU Write Event Mask Channel.

The code for the second-level interrupt handler is:

```
static unsigned int spu_decr_profile_slih(unsigned int events)
{
 if (events & MFC_DECREMENTER_EVENT) {
   if (_counting) {
     unsigned int end = spu_readch(SPU_RdDec);
     _count += (end > _count_base) ?
       (DECR_MAX + _count_base - end) :
       (_count_base - end);
   }
    events &= ~MFC_DECREMENTER_EVENT;
   spu_writech(SPU_WrEventMask, events);
   spu_writech(SPU_WrEventAck, MFC_DECREMENTER_EVENT);
   spu_writech(SPU_WrDec, DECR_COUNT);
   spu_writech(SPU_WrEventMask, MFC_DECREMENTER_EVENT);
   _count_base = spu_readch(SPU_RdDec);
 }
 return (events);
}
```

The handler should be installed as an event-specific handler that is dispatched from a common event handler, as described in *Section 18.6.9 Procedure for Handling the SPU Decrementer Event* on page 489. In addition, SPU events must be enabled when profiling is turned on (and can be disabled when profiling is turned off). Enabling and disabling of interrupts is done using the `spu_ienable` and `spu_idisable` intrinsics. The initialization of the handler and enabling of interrupts can be inserted into the body of the `prof_start()` macro (with logic to ensure that the handler only gets initialized once), and disabling of interrupts can be added to the `prof_stop()` macro.

# 14. Objects, Executables, and SPE Loading

This section describes processor-specific aspects of object-file formats that are used for program linking and loading. Topics described include the PowerPC Processor Element (PPE) and Synergistic Processor Element (SPE) extensions to the standard Executable and Linking Format (ELF) object-file formats, runtime loading of programs into an SPE, and the CBEA Embedded SPE Object Format (CESOF).

CESOF allows SPE executable object files to be embedded inside PPE object files. It supports a mechanism for linking SPE and PPE object files that access shared system memory objects through the globally defined symbols. This section also describes basic operating-system protocols for loading and executing programs from a CESOF object file.

## 14.1 Introduction

The Cell Broadband Engine Architecture (CBEA) processors[1] are heterogeneous multiprocessors not only because the SPEs and the PPE have different architectures but also because they have disjointed address spaces and different models of memory and resource protection. The PPE can run a virtual-memory operating system, so it can manage and access all system resources and capabilities. In contrast, the synergistic processor units (SPUs) are not intended to run an operating system, and SPE programs can access the main-storage address space, called the effective-address (EA) space, only indirectly through the DMA controller in their memory flow controller (MFC). The two processor architectures are different enough to require two distinct tool chains for software development.

The tool chains for both the PPE and SPE processor elements produce object files in the ELF format. ELF is a flexible, portable container for relocatable, executable, and shared object (dynamically linkable) output files of assemblers and linkers; this format is described in *Section 14.2* on page 398. The terms PPE-ELF and SPE-ELF are used to differentiate between ELF for the two architectures. CESOF is an application of PPE-ELF that allows PPE executable objects to contain SPE executables; this format is described in *Section 14.5* on page 408.

To ease the development of combined PPE-SPE multiprocessor programs, the operating-system model uses CESOF and provides SPE process-management primitives. Though programmers often keep in mind a heterogeneous model of the CBEA processors when dividing an application program into concurrent threads, the CESOF format and, for example, the Linux operating-system thread application programming interfaces (APIs) allow programmers to focus on application algorithms instead of managing basic tasks such as SPE process creation and global variable sharing between SPE and PPE threads.

The following sections describe the ELF format for PPE and SPE object and executable files, low-level support and requirements for loading SPE programs, the CESOF format for combining PPE and SPE object files, and the operating-system model of loading and running concurrent, cooperative SPE and PPE programs.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

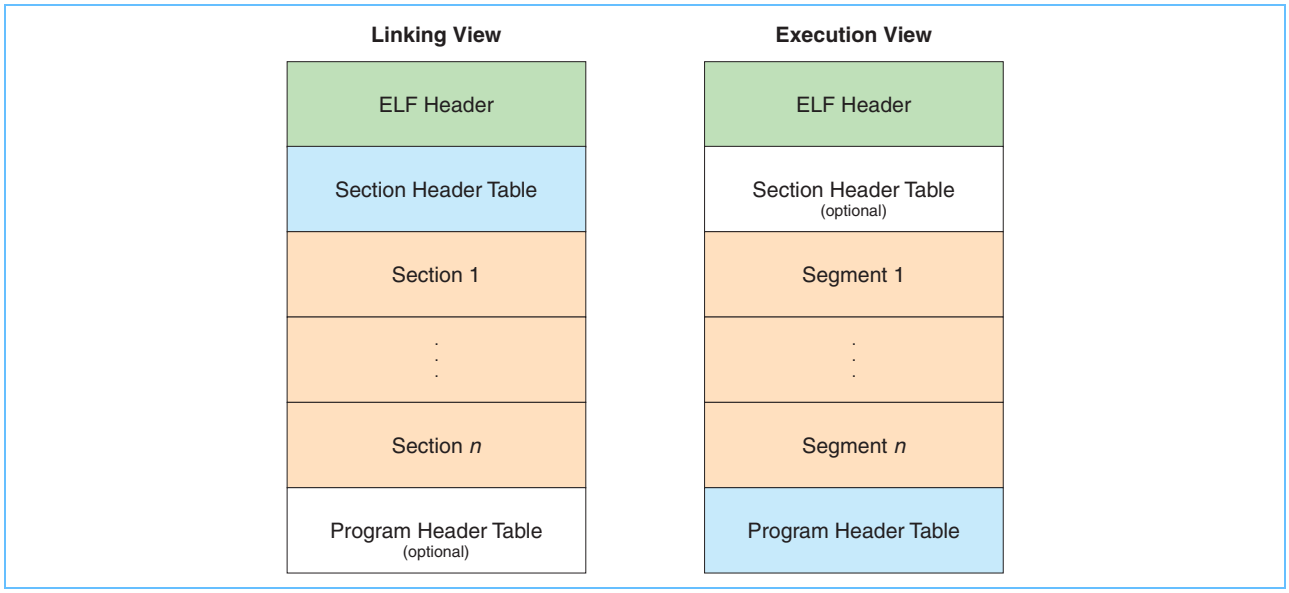
## 14.2 ELF Overview and Extensions

### 14.2.1 Overview

ELF object files are binary representations of programs intended to execute directly on CBEA processor elements. To organize the program information in an object file for efficient processing by linkers and loaders, ELF defines a simple hierarchy of structures. Object files have an ELF header that is used to find a section header or a program header; an object file can have both section and program headers. The section header specifies sections in the file for use in linking; the program header specifies program segments in the file for use in program loading. Much of an object file's information, such as processor instructions and initialized data, is contained in its section and segment data structures, but an object file also contains other data structures such as symbol tables.

The section header and the program header provide two different, parallel views of an object file's contents that reflect the different needs of linking and execution. ELF object files can be used by a linker to build an executable program object (linking view) and by an operating system or loader to build a process image in memory (execution view). *Figure 14-1* shows an object file's organization under the two views.

*Figure 14-1. Object-File Format*




The application binary interface (ABI) specifications for the CBEA processor elements, listed below, extend the basic ELF definition. The extensions result in two ELF definitions specific to the CBEA processors: PPE-ELF and SPE-ELF.

For a complete description of the basic ELF definition and the PPE and SPE APIs, see the following documents:

- *Tool Interface Standard (TIS) Executable and Linking Format (ELF)* specification
- *Linux Standard Base Core Specification for PPC 3.0* (http://www.linuxbase.org/spec)
- *64-bit PowerPC ELF Application Binary Interface Supplement 1.7.1*

- *SYSTEM V APPLICATION BINARY INTERFACE PowerPC Processor Supplement* (This document specifies the 32-bit ABI.)
- *Cell Broadband Engine Linux Reference Implementation Application Binary Interface Specification*
- *SPU Application Binary Interface Specification*

### 14.2.2 SPE-ELF Extensions

#### 14.2.2.1 *ELF Header*

The machine-specific information that applies to SPE-ELF objects is shown in *Table 14-1*.

*Table 14-1. SPE-ELF Header Fields*

| Field | Value | Comments |
|---|---|---|
| `e_ident[EI_CLASS]` | `ELFCLASS32` | Specifies that the objects are 32-bit objects. |
| `e_ident[EI_DATA]` | `ELFDATA2MSB` | Specifies that the objects use big-endian byte addressing. |
| `e_type` | `ET_NONE` | The file has no file type. |
| | `ET_REL` | The file is a relocatable file. The file holds code and data suitable for linking with objects to create an executable or plug-in file. |
| | `ET_EXEC` | The file is an executable file. The file holds a fully linked program suitable for execution. |
| | `ET_DYN` | The file is an SPE plug-in. The plug-in file must contain a SPUNAME note section for each named plug-in. |
| `e_machine` | `EM_SPU` | A value of 23 indicates that the code is to be executed on an SPE processor element. |
| `e_flags` | `0` | No flags have been defined; therefore, this member must contain zero. |

#### 14.2.2.2 *Symbols*

Except for the special `_EAR_` symbol-name mangling protocol used by programmers and tools to build CESOF objects (see *Section 14.5.5* on page 413), the global symbols produced by a compiler should not be *mangled*; that is, a symbol should not have any extra prefix or suffix added by the tool chain.

#### 14.2.2.3 *Sections*

Executable SPE-ELF objects can contain the three typical sections and the special `.toe` section as shown in *Table 14-2* on page 400.

*Table 14-2. SPE-ELF Special Sections*

| Name | Type | Attributes | Description |
|---|---|---|---|
| .bss | SPU_NOBITS | SHF_ALLOC + SHF_WRITE | This zero-length section specifies uninitialized data that contributes to the program's memory image. By definition, the program loader initializes the data with zeros. |
| .data | SPU_PROGBITS | SHF_ALLOC + SHF_WRITE | This section contains initialized data that contributes to the program's memory image. The data is copied from the executable file to EA space and then copied from EA space to an SPE's local storage (LS) using DMA commands. |
| .text | SPU_PROGBITS | SHF_ALLOC + SHF_EXECINSTR | This section contains a program's executable instructions, traditionally called text. The machine code is copied from the executable file to EA space and then copied from EA space to an SPE's LS using DMA commands. |
| .toe | SPU_PROGBITS | SHF_ALLOC | This section contains an array of Effective-Address Reference (EAR) structures. Each EAR contains an eight-byte EA and an eight-byte area reserved for future use. All sections of the toe type are combined into a single segment in the SPE executable, and the SPE loader will overwrite the segment in SPE memory with the toe shadow data segment from the enclosing PPE executable object. |

The .toe section is a key part of the CESOF specification and holds an array of effective address entry (EAR) structures. Each EAR holds the EA of a shared object in main storage. The SPE linker gathers all the .toe sections together into a single loadable toe segment in the linked SPE executable. See *Section 14.5.2* on page 409 and *Figure 14-6* on page 410 for more information about CESOF extensions.

#### 14.2.2.4 *Program Header*

The program header contains the locations of the segment images within the file and other information needed to create the memory image for a program. The SPE ABI specifies two program-header notes, SPE Environment Note and SPE Name Note, which are described next.

*SPE Environment Note*

SPE objects can contain sections of the SHT_NOTE with corresponding program-header elements of type PT_NOTE that define the attributes and runtime environment of an SPE program. *Table 14-3* provides details about the SPE environment note.

*Table 14-3. SPE-ELF Environment Note*

| Field | Size (bytes) | Value |
|---|---|---|
| namesz | 4 | 8 |
| descsz | 4 | sizeof (spu_env) |
| type | 4 | 1 |
| name | 8 | "IBM SPE" |
| desc | sizeof (spu_env) | spu_env structure contents |

The SPE environment note contains an instance of the spu_env structure in the desc field; entries in this structure are shown in *Table 14-4* on page 401.

*Table 14-4. The spu_env Structure*

| Type | Name | Description |
|---|---|---|
| Elf32_Word | revision | Specifies the spu_env structure revision number. The initial revision number is 1. Future additions to this structure are added to the end, and the revision number is incremented. |
| Elf32_Word | ls_size | Specifies the size of SPE LS on which the program is targeted to run; this is the required Local-Storage Limit Register (SPU_LSLR) setting. A size of zero indicates that the SPU_LSLR register must be set to the entire available LS address range. |
| Elf32_Word | stack_size | Specifies the SPE runtime stack size. This value is used for software runtime stack overflow checking. |
| Elf32_Word | flags | One flag is defined: EF_SPU_ENCRYPTED (bit 31). This flag specifies that the SPE ELF program is encrypted and must be decrypted and authenticated before being executed. |

*SPE Name Note*

An SPE object must be identified with a lookup name string, and this name must be contained within a SHT_NOTE and in corresponding program-header elements of type PT_NOTE. The lookup name is used to identify the corresponding SPE plug-in to the operating system.

*Table 14-5. SPE-ELF Name Note*

| Field | Size (bytes) | Value |
|---|---|---|
| namesz | 4 | 8 |
| descrsz | 4 | The number of bytes in the desc field. This value must be a multiple of 4. |
| type | 4 | 1 |
| name | 8 | "SPUNAME" |
| desc | See the descsz field | Contains a null-terminated look-up name string that identifies the object. |

The plug-in name extracted from the PT_NOTE name note is used by a plug-in management mechanism that provides services such as plug-in table construction, plug-in information retrieval, and loading plug-in tables and plug-in images to an SPE's LS.

## 14.3 Runtime Initializations and Requirements

The ABI documents for the CBEA processors specify the initial machine state that the loaders create for program execution on both the PPE and SPE processors. The specifications spell out register initialization, stack frame initialization, and initial arguments passed from the operating system. Programming language systems use the initial machine state requirements to establish a standard environment for their application programs.

### 14.3.1 PPE Initial Machine State

#### 14.3.1.1 *Operating System Interface for PPE Entry Point*

The C-language interface for a PPE program's entry point is conventionally declared as follows:

```
extern int main (int argc, char *argv[]; char *envp[]);
```

*Table 14-6* explains the meaning of the parameters passed by the operating system (OS) on behalf of the program that is invoking the application program.

*Table 14-6. PPE Program Initial Parameters*

| Parameter | Value |
|---|---|
| `argc` | A nonnegative argument count; specifies the number of arguments in the `argv` array. |
| `argv` | An array of argument strings; the OS ensures that `argv[argc] == 0`. |
| `envp` | An array of environment strings; as with `argv`, the `envp` array is terminated by a NULL pointer. |

#### 14.3.1.2 ***Register Initialization***

*Table 14-7* shows the initial register values required by the ABI specifications for the CBEA processors. Registers other than those listed in *Table 14-7* may be cleared when loading PPE programs, but a program that requires registers to have specific values must set them explicitly; a program must not rely on the loader to set any register other than those shown in *Table 14-7*.

*Table 14-7. PPE Initial Register State*

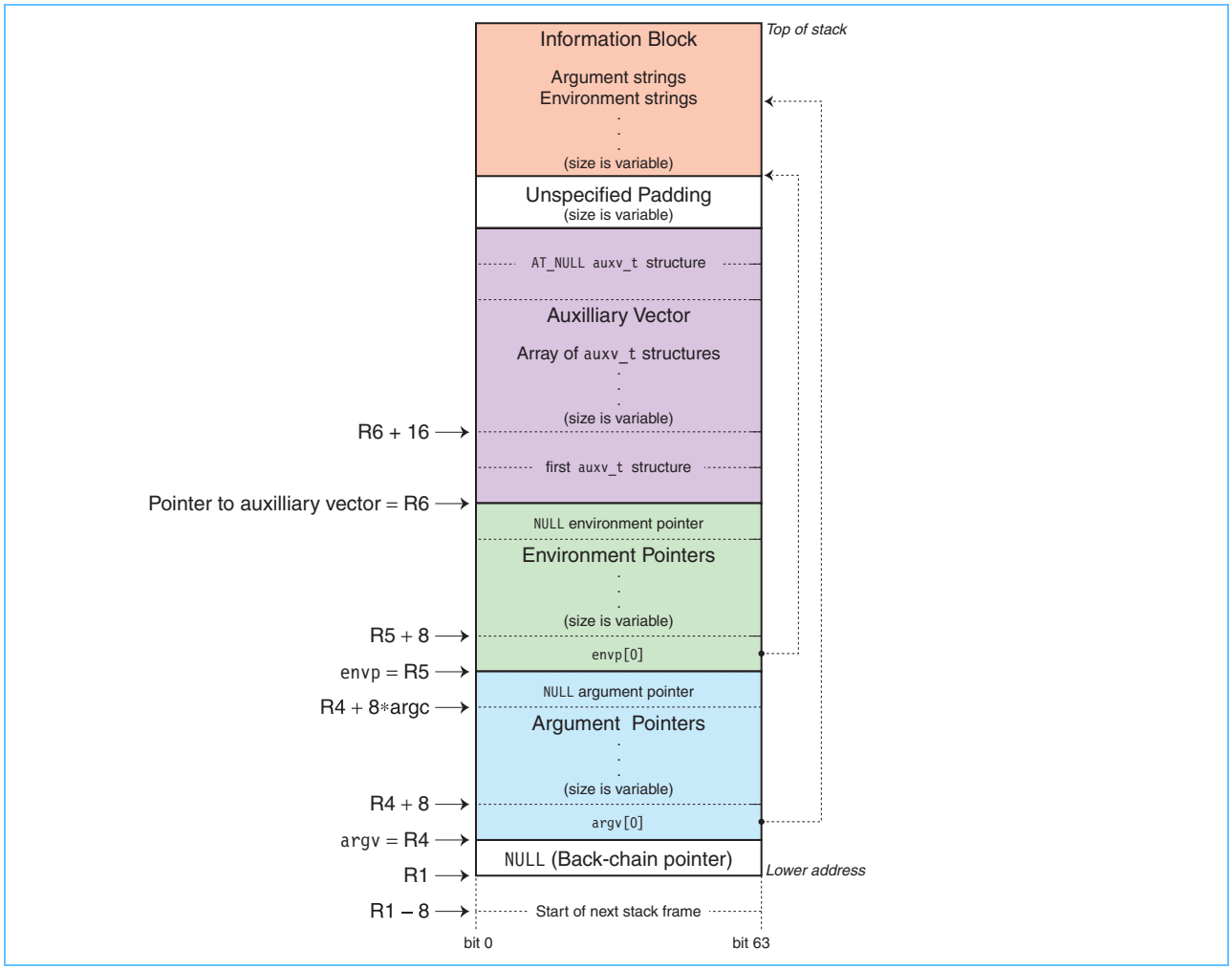
| Register | Contents |
|---|---|
| R1 | Initial stack pointer. The value in R1 is aligned to a quadword boundary and is a pointer to a word containing a `NULL` pointer. |
| R2 | Initial TOC (table of contents) pointer. This value is obtained through the function descriptor pointed at by the `e_entry` field in the ELF header and is used by the dynamic linker. See *64-bit PowerPC ELF Application Binary Interface Supplement 1.7* for more information. |
| R3 | Contains `argc`, the number of arguments passed by the OS. |
| R4 | Contains `argv`, a pointer to the array of pointers to arguments on the stack. The array is immediately followed by a `NULL` pointer; if there are no arguments, R4 points to a `NULL` pointer. |
| R5 | Contains `envp`, a pointer to the array of pointers to environment strings on the stack. The array is immediately followed by a `NULL` pointer; if there is no environment, R5 points to a `NULL` pointer. |
| R6 | Contains a pointer to the auxiliary vector. The auxiliary vector shall have at least one member, a terminating entry with an `a_type` of `AT_NULL`. |
| R7 | Contains a termination-function pointer. If R7 is not zero, its value represents a function pointer that the application should register with `atexit(BA_OS)`. If R7 is zero, no action is required. |
| fpscr | Contains zero, which specifies “round to nearest” rounding mode, IEEE mode, and that floating-point exceptions are disabled. |

#### 14.3.1.3 ***Stack Initialization***

Before the PPE loader jumps to a PPE application’s entry point, the stack is initialized, but the ABI specifications do not define a fixed stack address. A program’s stack can change from one system to another and from one process invocation to another. Thus, process initialization code must use the stack address provided in R1.

The initial stack frame, shown in *Figure 14-2* on page 403, is set up with a `NULL` back-chain pointer (at the address in R1). Above the `NULL` back-chain pointer are the three areas containing explicit and implicit arguments passed by the OS.

*Figure 14-2. PPE 64-Bit Initial Stack Frame*




The first area is an array of pointers to parameter strings. The second area is an array of pointers to environment strings. Both of these areas are terminated by a NULL pointer. These areas contain parameters passed from the invoking application to the invoked application, and each pointer in these areas points to a string in the information block at the top of the stack frame.

The third area is an array of auxiliary information structures; each structure is of type `auxv_t`. The definition of the structure is shown in the following code sample:

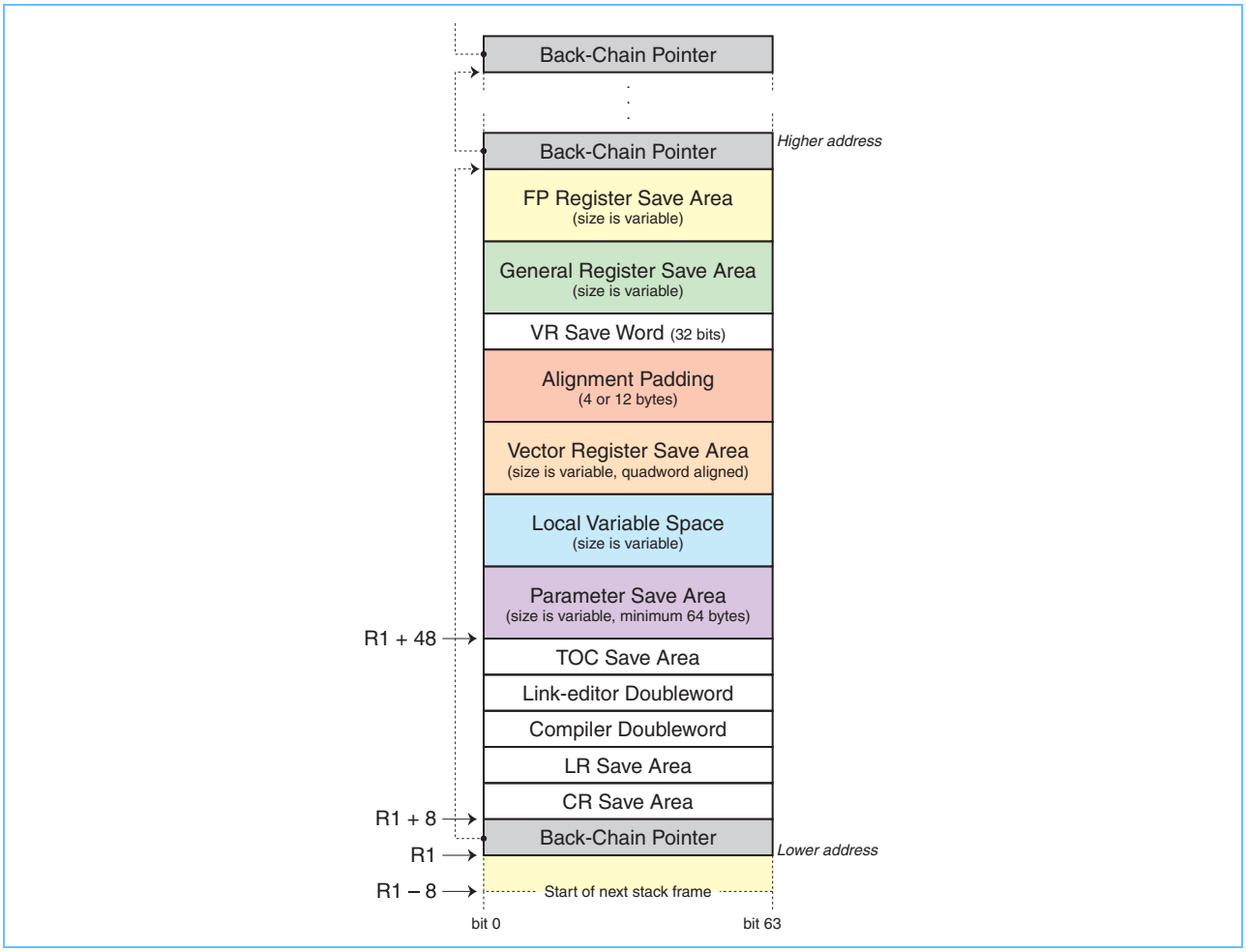
```
typedef struct
{
    int a_type;
    union
    {
        long   a_val;
        void  *a_ptr;
        void (*a_fcn)();
    } a_un;
} auxv_t;
```

Thus, each entry in the auxiliary vector area consists of a type word and a value, pointer, or function pointer. The auxiliary vector contains information passed from the OS to the invoked application program. See the *System V Application Binary Interface PowerPC Processor Supplement* and *64-bit PowerPC ELF Application Binary Interface Supplement* documents for more information about the auxiliary vector area.

*Figure 14-2* on page 403 shows the layout of the PPE initial stack frame in a 64-bit execution environment. The pointers in the `argv` and `envp` arrays point into the information block. Any pointers that occur in the `auxv_t` array point into other areas of main storage.

The first application function call and subsequent calls will set up an instance of the standard SPE stack frame for a 64-bit execution environment, which is shown in *Figure 14-3*. If a calling function needs more than eight doublewords of arguments, it places the additional arguments in its parameter save area. The callee function uses an offset of 48 bytes from the back-chain pointer in its stack frame to find the additional arguments in the caller's stack frame.

*Figure 14-3. PPE 64-Bit Standard Stack Frame*



### 14.3.2 SPE Initial Machine State for Linux

This section specifies the initial SPE machine state for the Linux operating system environment. Other other operating will provide similar operating environments.

#### 14.3.2.1 *Linux-OS Interface for SPE Entry Point*

The C-language interface for an SPE program's entry point is conventionally declared as follows in the Linux runtime environment:

```
extern int
main (unsigned long long spuid, unsigned long long argp, unsigned long long envp);
```

*Table 14-8* explains the meaning of the parameters passed by the OS. This is the interface only for the Linux runtime environment. Other systems might have different interfaces. In addition, in Linux, stand alone SPU programs[2] provide a standard main interface (`int argc, char *argv[]`), where the pointers are LS pointers.

*Table 14-8. SPE Program Initial Parameters*

| Parameter | Value |
|---|---|
| `spuid` | A unique SPU task identifier assigned by the operating system. |
| `argp` | An EA in main storage to application parameters. |
| `env` | An EA in main storage to runtime environment information. |

#### 14.3.2.2 *Register Initialization*

*Table 14-9* on page 406 shows the initial register values required by the *Cell Broadband Engine Linux Reference Implementation Application Binary Interface Specification*. Registers other than those listed in *Table 14-9* may be cleared when loading SPE programs, but a program that requires registers to have specific values must set them explicitly; a program must not rely on the loader to set any register other than those shown in *Table 14-9*.

2. A programming model that allows conventional C programs to be compiled and run on an SPE, unmodified.

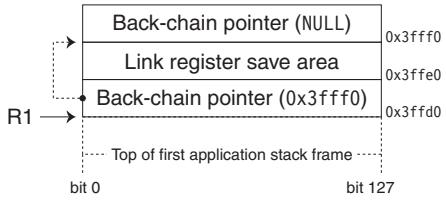
*Table 14-9. SPE Register Initial Values*

| Register | Symbolic Name | Data Type | Comment |
|---|---|---|---|
| R1 | SP | 32-bit unsigned int | Initial stack pointer. crt0 initializes word element 0 of R1 to point to a minimal stack frame starting at the highest LS address. For example, a system with 256 KB LS initializes the stack pointer to x'3FFD0', and the initial, minimal stack frame would be as shown in *Figure 14-4*.<br>The 32-bit unsigned integer in word-element 1 of the SP is used to establish the Available Stack Space. If the runtime stack size is 0, then the value of Available Stack Space is initialized to <top_of_stack> through _end (end of text/data linker symbol). Otherwise, the Available Stack Space is initialized to <stack_size>. |
| R3 | spuid | 64-bit unsigned int | SPU task identifier. The spuid value is a unique identifier assigned by the PPE OS. The spuid allows both the SPE program and the PPE program to identify the SPE thread (program instance) within the system. |
| R4 | argp | 64-bit EA | Pointer into main storage to an array of application-defined program parameters passed from the application that spawned the SPE thread. |
| R5 | envp | 64-bit EA | Pointer into main storage to SPU task environment structure passed from the PPE OS. The content of the structure is not defined by the SPU ABI specification. |

#### 14.3.2.3 ***Stack Initialization***

The stack for an SPU program will be initialized with a single minimal ABI-compliant stack frame that contains three quadwords. In most system implementations, the entry function crt0 or an equivalent will be responsible for setting up this frame. See *Figure 14-4*.

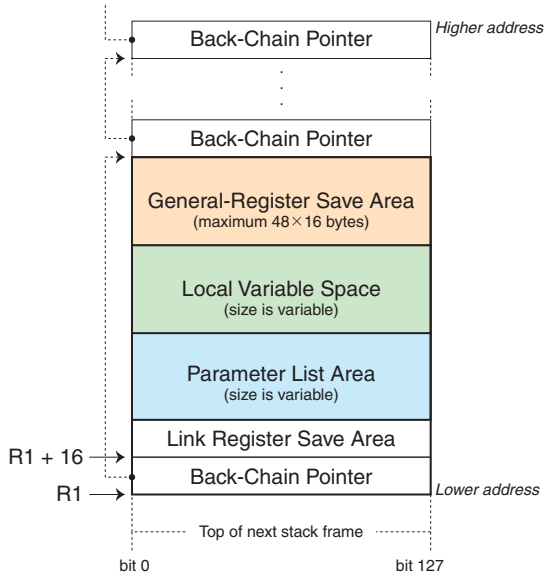
*Figure 14-4. SPE Initial Stack Frame*




Typically, the top of the stack is set to the highest LS quadword memory address; for a 256 KB LS, this address is x'3FFF0'. The quadword at x'3FFF0' contains zero, the NULL pointer. The quadword at x'3FFE0' is the link-register save area for use by the entry function to save R0 when it calls a function. The quadword at x'3FFD0' contains a back chain pointer to the top of the stack. R1 is set to x'3FFD0', which is the start of this initial stack frame.

For compliance with the ABI specification, the stack frame shown in *Figure 14-4* provides a place for the first function call in application code to save the link register. The first application function call and subsequent calls will set up an instance of the standard SPE stack frame shown in *Figure 14-5* on page 407.

*Figure 14-5. SPE Standard Application Function Stack Frame*




Up to 77 quadwords of arguments are passed in SPU registers in R3 through R79. If the calling function needs to pass more than 77 quadwords, it must allocate a parameter list area in its stack frame large enough to hold the additional arguments.

## 14.4 Linker Requirements

The extensions in the CESOF specification require some support from the PPE and SPE tool-chain link editors.

### 14.4.1 SPE Linker Requirements

The linker in an the SPE tool chain has some special requirements that include:

#### 14.4.1.1 *Handling `.toe` Sections*

The SPU linker must recognize and properly handle the `.toe` section type. An SPE-ELF object need not have a `.toe` section or segment. Each `.toe` section holds the array of EAR structures (see *Section 14.5.5* on page 413), and all `.toe` sections, if any are present, must be grouped into a loadable segment in an SPE-ELF executable.

#### 14.4.1.2 *Alignment of Loadable Segments*

The SPE linker must ensure that the loadable SPU segments in an SPU-ELF object are aligned on 16-byte boundaries and padded in length to a multiple of 16 bytes to satisfy the MFC DMA size requirement. The linker may further restrict the alignment of the segments to 128-byte

boundaries for improved transfer performance on some CBEA implementations. These alignment and size requirements allow optimal DMA performance when copying SPE-ELF loadable segments into an SPE's LS storage space.

### 14.4.2 PPE Linker Requirements

The special requirement for the PPE linker is the need to recognize and properly handle the sections (for example, `.spe.elf` or `.rodata.speelf`) wrapping the SPE-ELF executable image; see *Section 14.2.2* on page 399. In a PPE executable, such SPE-ELF image sections must be included by the linker in a per-system loadable segment, such as the `.text` segment.

## 14.5 The CESOF Format

The software-development tool chains for CBEA processor programs conform to three standards:

- Tool Interface Standard (TIS) and Execution specification
- Executable and Linking Format (ELF) specifications
- PPE and SPE Application Binary Interface (ABI) specifications

Together, these standards define an extensible, portable object-file format for relocatable and executable object files that can be statically or dynamically linked.

Because the PPE and SPE processor-element architectures are significantly different, separate tool chains are used for PPE and SPE program development. An important limitation for CBEA processor program development is that these TIS, ELF, and ABI specifications define no way to resolve references between code and data objects for different architectures.

To exploit the performance potential of CBEA processor systems, applications typically contain cooperating PPE and SPE programs that share data variables. Thus, programmers need an object definition beyond the base of TIS-ELF.

The CBEA Embedded SPE Object Format (CESOF) provides three key capabilities for CBEA processor programming environments:

- PPE and SPE programs can co-exist in a single PPE-ELF object file.
- PPE and SPE programs can share global variable references.
- The PPE link editor can accept objects that contain PPE or SPE code.

CESOF uses only the facilities defined in the TIS-ELF standard. Using an embedding approach requires the minimum of changes to existing PPE and SPE development tools. Because CESOF objects are PPE-ELF objects, they can participate in any program linking and execution process, static or dynamic, with other PPE-ELF objects. CESOF files can be archived, linked, and shared just like other PPE-ELF object files.

### 14.5.1 CESOF Overview

With CESOF, programmers can achieve some of the effects of linking PPE and SPE executables. The PPE linker can create a single PPE-ELF executable file that contains code and data for both PPE and SPE processor elements. An OS can load PPE and SPE programs that run concurrently and work cooperatively from an integrated PPE executable image. The CESOF specification includes the following mechanisms:

- The `.spe.elf` section (which can be named differently) to embed SPE-ELF executables in PPE-ELF objects.
- The SPE program handle data structure, which is used by an OS runtime to access the embedded SPE program and data.
- A reserved “shadow” section to keep the EA-space bindings for the shared global data object references of the embedded SPE-ELF executable.
- A runtime requirement that overlays the shadow section onto an SPE executable in LS memory before the SPE executable begins running.

Each of these mechanisms is described in a following section.

CESOF achieves cross-architecture linking by the following method:

- The SPE-ELF executable object defines a loadable segment that contains a copy of the table of effective address references, or TOE (see *Section 14.5.5* on page 413). The TOE segment collects all the `.toe` sections from the linked SPE-ELF linkable objects.
- The CESOF-ELF (PPE-ELF) object, that includes the SPE-ELF executable image, defines another copy of the TOE, called the TOE shadow, in a `.data` section (as per-process data). The TOE shadow allocates the same space as the previous TOE segment of the embedded SPE-ELF executable.
- The PPE-ELF linker binds the references in the TOE shadow to addresses in EA space and fills the TOE shadow with the resolved (or bound) addresses.
- At runtime, the SPE loader loads the SPE-ELF image into an SPE’s LS memory and then replaces the toe segment in LS with the TOE shadow by copying the shadow over the toe segment.
- The running SPE-ELF program uses the bound addresses in its toe segment to reference shared data objects in EA main storage (using DMA commands to access the objects).

Thus, because the SPE program references shared data objects with a level of indirection through the TOE entries and the loader overwrites the toe segment with PPE-linker bound addresses, CESOF achieves the effect of the PPE linker binding PPE addresses into the SPE executable, which is not otherwise supported by the ELF specification. Thanks to the CESOF mechanisms, the linker can do this without knowing anything about the SPE architecture.

### 14.5.2 CESOF Use Convention of ELF

CESOF introduces new sections into PPE-ELF and SPE-ELF object respectively.

- PPE-ELF `.spe.elf` section: a container for an SPE-ELF image in a PPE-ELF object.
- SPE-ELF `.toe` section: a container for an array of EAR structures; each EAR contains an EA pointer to a shared data-object EA space. All `.toe` sections are concatenated into one toe

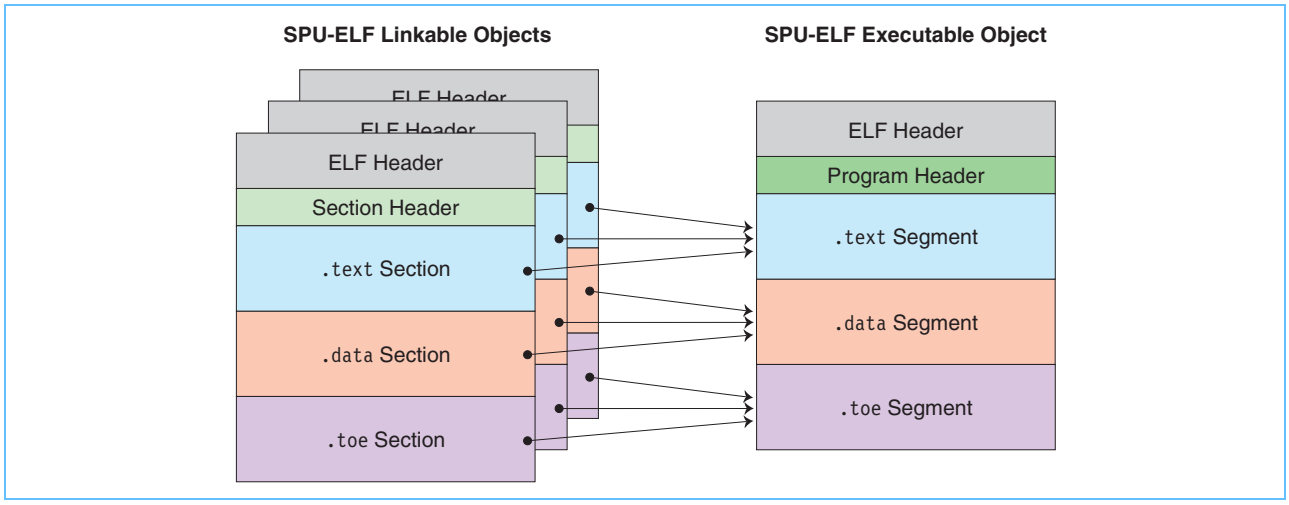
segment in an SPE-ELF executable. `.toe` sections and the segment are filled with empty EAR structures.

- PPE-ELF TOE shadow section: a copy of the SPE-ELF toe segment is allocated in a PPE-ELF data section. This matching TOE shadow data section contains initialized EAR structures containing bound EA pointers instead of the empty EAR structures in the toe segment of the SPE executable.

When the SPE-ELF executable is loaded into an SPE's LS, the TOE shadow is copied directly over the toe segment in the LS. The copying effectively updated the EA addresses to the bound values. At run time, the SPE code executes DMA commands that use the bound addresses in the TOE in LS to access shared data objects in main storage.

*Figure 14-6* shows how the SPE-ELF sections in linkable objects are combined into segments in an executable object.

*Figure 14-6. Linking SPE-ELF Sections into Loadable Segments in an Executable*




### 14.5.3 Embedding an SPE-ELF Executable in a PPE-ELF Object: The `.spu.elf` Section

An SPE program is compiled and linked by the SPE compiler and tool chain into an SPE-ELF executable file. Without the support of the mechanisms like those in CESOF, the only way a PPE program can use an SPE program is to refer to the name of the SPE-ELF executable object file. Because the PPE operating system's linker and runtime loader know nothing about SPE-ELF executables in this case, a PPE program must explicitly copy or memory-map the SPE executable image from a file to main storage and explicitly program an SPE's DMA controller to move the SPE executable to an SPE's LS.

The process of loading an SPE-ELF image from a file to an SPE's LS requires tedious programming that distracts a programmer from coding a solution to an application problem. In addition, if the SPE executable is shared by several PPE processes, storage is wasted by having a separate copy in each PPE process.

A solution to the explicit loading and the storage redundancy is to make the SPE-ELF object part of the PPE-ELF object. CESOF accomplishes this by defining a special section in PPE-ELF to include the image of the SPE-ELF executable file in the section.

The CESOF specification suggests that ".spe.elf" is used as the name for such sections. However, a tool chain may choose to use any name as long as the SPE-ELF image sections are linked and loaded into a per-system segment. The single copy of image can then be shared by more than one processes.

In assembler code, the declaration of the PPE-ELF section and file inclusion can be as shown in the following example:

```
    .global      spu_program

    .align       7
    .section     .spu.elf, "ax", @progbits
spu_program:
    .incbin      "spu_program.elf"
```

In a PPE program, a programmer declares a symbol external to the code module that creates a reference to the SPE image:

```
extern spu_program[];
```

The result is that the SPE-ELF executable object has been embedded (or wrapped) in the `.spe.elf` section. In an executable object, all the `.spe.elf` sections may be collected in the per-system loadable segment, such as the `.text` segment, so that main storage will contain only one copy of an SPE-ELF image shared by multiple processes. The CESOF specification, however, still allows the `.spu.elf` sections to be merged into any segment, such as a data segment, depending on the needs of the hosting OS.

The PPE tool-chain linker can link the PPE program's reference to the global `spu_program` symbol defined in the `.spu.elf` section. The PPE program can use the `spu_program` symbol to access all the information in the SPE-ELF program image using the image's ELF header.

The CESOF specification requires the `.spe.elf` image in the loadable segment in the executable to be aligned to a DMA (128-byte) boundary so that the SPE-ELF image can be copied into an SPE's LS with maximum efficiency.

### 14.5.4 The `spe_program_handle` Data Structure

Using the simple solution shown in *Section 14.5.3* on page 410, a problem arises when the included SPE-ELF object is linked into a shared PPE-ELF library object. If a reference is made directly from other modules in the shared library, the reference will be absolute, which causes the shared library to lose its position independence. Position independence is required to permit a shared library to be dynamically linked at a different address for each process that shares the library.

To maintain position independence in shared libraries that have embedded SPE-ELF objects, CESOF requires a level of indirection through a pointer to the embedded image. The pointer holds the absolute address that is determined by the dynamic linker when it binds the shared library into an in-memory process image. The pointer is held in an `spe_program_handle` data structure with the following definition:

```
typedef struct spe_program_handle {
    int     handle_size;              // sizeof (SPEProgramHandle)
    void*   elf_image;                // pointer to the embedded SPE image
    void*   toe_shadow;               // pointer to the toe overlay image
} SPEProgramHandle;
```

The `spe_program_handle` contains a `handle_size` member so software can determine whether the CESOF object was created for a 32-bit or 64-bit environment. The `elf_image` member holds the address of the embedded SPE-ELF image. The `toe_shadow` member holds the address of the TOE shadow overlay for the toe segment in the SPE-ELF image; this overlay will be copied to the location in LS memory where the toe segment normally resides. The definition of the toe segment and the TOE shadow and their purpose are presented in *Section 14.5.5* on page 413.

Illustrated In assembler language, the 32-bit CESOF wrapping layer now looks as follows:

```
/* per-process section */
    .global     spe_program_handle

    /* new section with pointers for relocatability */
    .section    .data, “a”, @progbits
spe_program_handle:
    .align   2
    .int     12
    .int    _spe_elf_image
    .int    _spe_toe_shadow

    .section    .data, “a”, @progbits
   .align   7
 _spe_toe_shadow:
.extern shared_symbol_name_1
    .int0
    .int    shared_symbol_name_1
    .quad    0
.extern shared_symbol_name_2
    .int0
    .int0   shared_symbol_name_2
   .quad    0

/* per-system section */
    .section    .spe.elf, “ax”, @progbits
    .align      7
_spe_elf_image:
    .incbin     “spe_program.exe”
```

The `spe_program_handle` structure becomes part of a per-process loadable segment so that each in-memory process image has a private copy of the pointer to the SPE-ELF image. The SPE-ELF image is in the `.spe.elf` section, which becomes part of a single shared segment in physical storage through the virtual-memory mapping.

In the case that the SPE-ELF object `spe_program.exe` is wrapped into a PPE-ELF shared library object, the dynamic linker will initialize and map the PPE process image and load initial segments of PPE program text and the shared library into system storage (some or most loading can be delayed until it is demanded by the virtual-memory system). At this point, the dynamic linker can bind the reference to `_spe_elf_image` by computing an absolute (virtual) address value and storing that in the `_spe_elf_image` member of the `spu_program_handle` structure.

Because of the level of indirection through the pointer in the `spe_program_handle` structure, the PPE program code declares a reference to the SPE image as:

```
extern SPEProgramHandle spe_program_handle;
```

The preceding example illustrates the CESOF mechanism with assembler-language code, but the process of embedding SPE-ELF objects into PPE-ELF files can be automated. CBEA processor application programmers can use a few simple utility programs to create PPE-ELF-compatible CESOF objects. See *Section 14.5.5.2* on page 415 and *Section 14.5.5.3* on page 415.

### 14.5.5 The TOE: Accessing Symbol Values Defined in EA Space

Global data objects in CBEA-processor-level programs are frequently declared and represented by PPE level variables and symbols. These data objects are shared by the interacting PPE and SPE programs. The resolved values of the PPE symbols typically represent the addresses of the global data objects. These values are useful to the sharing SPE programs for their DMA target and source.

CESOF allows the PPE symbol values to be resolved into SPE programs statically. CESOF inserts a level of indirection between the SPE program and the address of the shared variables. The level of indirection is an EA variable inside an effective address reference (EAR) structure. All EARs in an SPE program are collected into a table of EARs, or TOE. A TOE is declared in a special SPE-ELF `.toe` section.

An EAR is a simple, 16-byte data structure defined as follows:

```
typedef struct elf_toe_entry
{
        ELF64_Addr    ea_value;
        ELF64_Addr    ea_info;     // reserved for future use
} EAR __attribute__ ((aligned (16)));
```

In assembler language, an EAR structure in a `.toe` section can be declared as follows:

```
.section    .toe
.align      7             # align on 128-byte boundary
.global     _EAR_g_mem_obj_1
_EAR_g_mem_obj_1:
        .quad    0        # ea_value
        .quad    0        # ea_info
```

The EAR structures in the SPE-ELF `.toe` sections (and eventually collected into the loadable toe segment) are empty place-holders. They exist so that the SPE linker can bind references to the symbol labels of the EAR structures to absolute LS addresses and thus build a fully resolved SPE-ELF executable object. At runtime, after the SPE-ELF executable image has been loaded into LS memory but before the SPE program is running, the SPE loader copies the TOE shadow from the PPE-ELF executable over the toe segment in LS. Thereafter, the EAR structures contain valid EAs that have been bound (by the PPE linker) to shared objects in main storage.

Each 16-byte EAR structure is aligned on a 16-byte (quadword) boundary because the entire structure will be loaded into an SPU register at run time. The `ea_value` member is positioned in the structure so that it will be loaded into the register's preferred scalar slot (the most-significant doubleword). Because the EA is already in the preferred scalar slot, an SPU program can transfer the `ea_value` address to the MFC command buffer with minimum overhead.

Using the C-language `spu_mfcdma64` intrinsic, an SPU programmer can use the following code to build a DMA command that copies the shared data object from EA space into LS space:

```
#include <spu_mfcio.h>

extern EAR _EAR_g_mem_obj_1;       // EAR structure to hold the EA of g_mem_obj_1
int g_mem_obj_1;                   // The SPE local (cached) copy of g_mem_obj_1

spu_mfcdma64((volatile void *)&g_mem_obj,              // destination LS
    (unsigned int)(_EAR_g_mem_obj_1.ea_value>>32),     // EA high 32 bits
    (unsigned int)(_EAR_g_mem_obj_1.ea_value),         // EA low 32 bits
    sizeof(g_mem_obj_1),                               // object size in bytes
    0,                                                 // tag id
    MFC_GET);                                          // DMA command
```

#### 14.5.5.1 *Symbol-Name Signatures*

To separate ordinary SPE symbols from those representing the references to the EA memory objects, CESOF defines a special symbol name mangling scheme.

In PPE source code, no special coding guidelines are required for the declaration of a shared global data object. In SPE source code, the symbol representing the reference to an EA memory object must be given a name that contains a special signature. This signature must be used only for this purpose. In this Joint Software Reference Environment (JSRE) implementation, the signature is `_EAR_`. The following is an example of such a declaration:

```
typedef struct elf_toe_entry
{
        ELF64_Addr    ea_value;
        ELF64_Addr    ea_info;         // reserved for future use
} EAR __attribute__ ((aligned (16)));

// the SPE symbol representing the reference to PPE g_mem_obj_1 in EA
const EAR _EAR_g_mem_obj_1;
```

#### 14.5.5.2 *Building the SPE-ELF .toe Section*

When a programmer uses this protocol, the symbol table in an SPE-ELF relocatable object will contain names that have a unique and easily recognizable signature. Assembly language segments can be handcrafted to create this simple array of EAR entries.

Though not provided in the JSRE reference implementation, a simple tool can be built with standard `binutil` commands that can extract these names from the object, create an assembler-language source file for the `.toe` section needed by the relocatable object, and assemble it into a `.toe` object. The relocatable object and the `.toe` object can then be linked into a single relocatable object.

#### 14.5.5.3 *Building a CESOF Object: Embed the SPE Object and Create the TOE Shadow*

When the SPE program has been compiled and linked, with the associated .toe sections, into an SPE-ELF executable object, it is ready to be embedded into a CESOF object, which can then be linked with other PPE-ELF objects to produce an executable suitable for loading and execution by the OS.

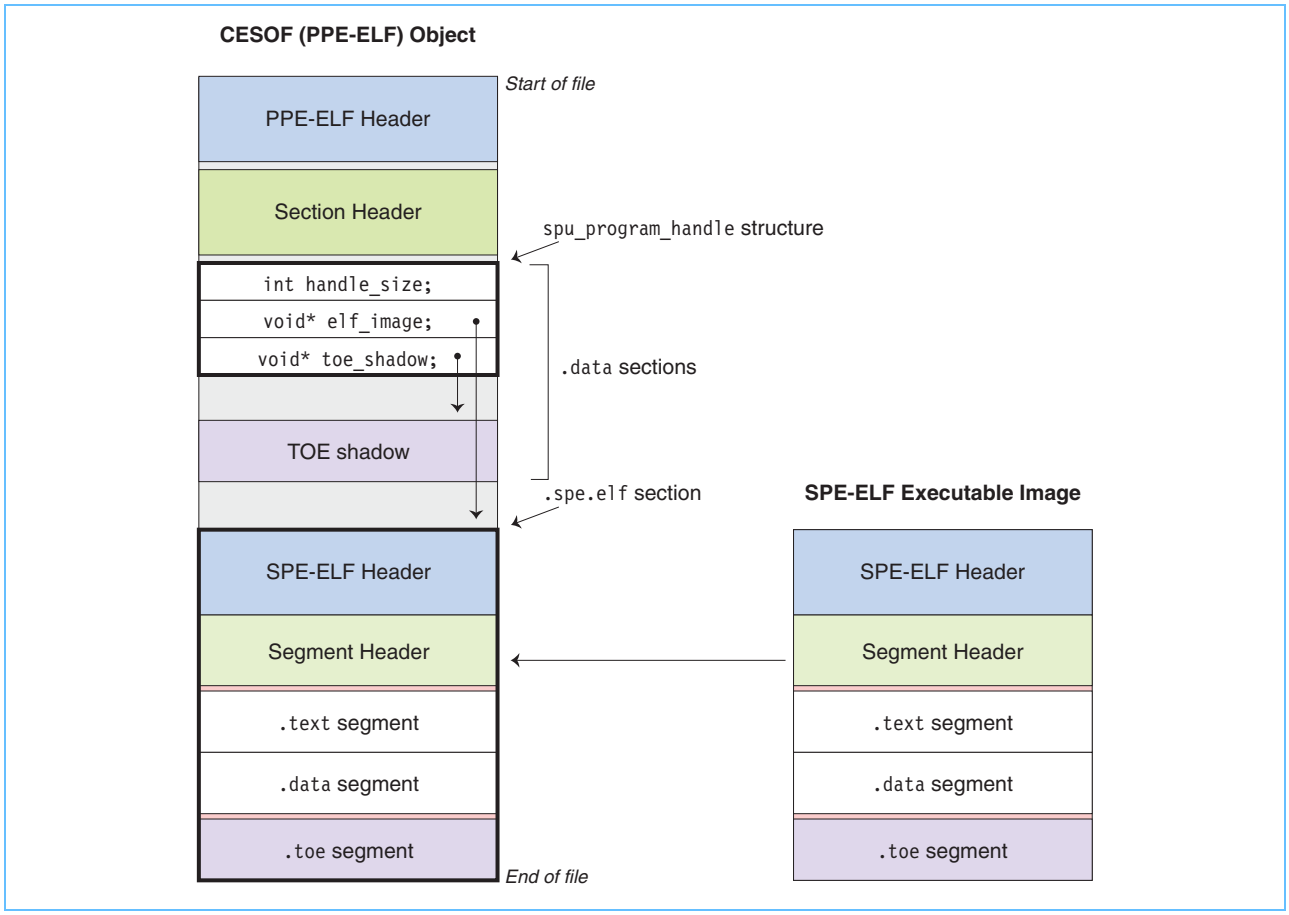
A programmer can build a CESOF relocatable object from the SPE-ELF executable with a tool called “embedspu” in the JSRE reference implementation.

A CESOF object contains three distinct parts:

- The `spe_program_handle` data structure
- The section containing the embedded SPE-ELF executable object
- The section containing the TOE shadow

*Figure 14-7* on page 416 shows the layout of a CESOF relocatable object and how an SPE-ELF object is embedded in the `.spe.elf` section.

*Figure 14-7. CESOF Object Layout*




The following shell-language script illustrates how `embed_spu` can be implemented to construct a CESOF object. This version of the tool builds a CESOF object for a 64-bit PPE OS environment.

```
#!/bin/ksh
#
if [ "x$3" = "x" ]
then
    echo Usage: %0 \<handle_symbol\> \<infile\> \<outfile\>
fi

handle_symbol=$1    # $1 must contain only valid symbol characters

pu-as -o $3 --<<END_xYz
.global ${handle_symbol}
.section .data, "a", @progbits    # definition of the SPE program handle
.align 2
${handle_symbol}:
        .int    20                # 20 bytes for a 64-bit CESOF object
        .quad   _spe_elf_image    # 64-bit pointer
        .quad   _spe_toe_shadow   # 64-bit pointer
```

```
.section .spu.elf, "ax", @progbits
.align 7
_spe_elf_image:
        .incbin "$2"

.section .data, "a", @progbits
.align 7
_spu_toe_shadow:
        `mk_shadow_toe $2`

END_xYz
```

The `embedspu` tool builds assembler source code for an `spe_program_handle` structure, an `.spe.elf` section containing the SPE-ELF executable, and a .data section containing the TOE shadow. The TOE shadow assembler source code is built by the `mk_shadow_toe` tool. The following shows the shell-language source code for a 64-bit version of `mk_shadow_toe`:

```
#!/bin/ksh
# 1. spu-readelf:  dump the symbol table
# 2. egrep:        find the symbols with _EAR_ signature
# 3. sort:         sort the symbols according to their addresses
# 4. awk:          output an EAR entry for each symbol, stripped of signature

spu-readelf -s $1 | egrep '^_EAR_' | sort -k 2 | awk '{ if (index($7, "UND")==0) {\
        print ".quad " substr($8, 6, length($8)); \
        print ".quad  0"; \
}}'
```

After the `mk_shadow_toe` tool creates the TOE shadow code, the `embedspu` tool passes the complete CESOF assembler code to the assembler, which outputs a CESOF relocatable object file. This object can be linked with other PPE-ELF objects to build an executable.

When the PPE-ELF executable is linked together, either statically or dynamically, the TOE shadow will be filled with the bound addresses of shared global data objects, and the SPE loader will write the TOE shadow into LS over the toe segment.

The JSRE reference implementation comes with a more complete implementation of the “embedspu” tool.

### 14.5.6 Future Software Tool Chain Enhancements for CESOF

The software development process described requires a programmer to declare shared data objects in SPE source code with a unique signature and use nonstandard tools to build the CESOF structures. There are two key deficiencies of this process:

- The SPE programmer must code explicit DMA operations in SPE programs to copy shared global objects into LS and copy them back to main storage.

- The source-language compiler knows nothing of the DMA transfers to and from main storage being performed by the SPE program to cache a shared global object in LS.

One way to overcome these deficiencies is to encode, in an object's declaration, the fact that it is a shared global object. The following syntax might be used for such a declaration:

```
remote int    g_mem_obj_1 __attribute__((section (".toe")));
remote double g_mem_obj_2 __attribute__((section (".toe")));
```

In these two declarations, the special `section (".toe")` attribute tells the compiler the following:

- Space in LS must be reserved for the objects—an `int` and a `double`—as usual.
- An entry in the special `.toe` section must be built for each EAR that will contain the EA of a shared object in main storage.
- Whenever the program uses an object declared with the `section (".toe")` attribute, the value of the object must be made coherent with the value in main storage. Thus, if the program uses the object as an r-value (a computation source), the compiler must emit an appropriate MFC DMA **get** command that uses the EA from the object's corresponding EAR. If the program uses the object as an l-value (a computation destination), the compiler must emit an appropriate MFC DMA **put** command.

An unsophisticated compiler will emit DMA **get** or **put** commands whenever the object is referenced in the program, but a sophisticated compiler can perform software caching of objects by performing live analysis of shared objects and using the synchronization facilities of the MFC, such as lock-line reservation, to reduce the number of MFC DMA commands executed at run time.

## 14.6 SPE Runtime Loader

Loading cooperative PPE and SPE programs onto a CBEA processor system, from a combined CESOF object or separately, requires loading both PPE and SPE programs into two distinct memory spaces. First, the entire program is loaded into the EA main storage space by the PPE operating system; second, the SPE program is copied, when needed, from EA space into the LS of the SPE that is selected to run the SPE program.

### 14.6.1 Runtime Loader Overview

A modern operating system, such as Linux, can use a runtime loader to load a binary program into memory for execution. The loader copies or simply maps the program image into the address space of memory that is directly accessible by the processor. In the case of CBEA processor programs, the runtime loader copies or maps code (`.text` in an ELF object), initialized data (`.data` in an ELF object), and uninitialized data (`.bss` in an ELF object) segments into memory and initializes the stack for the process that will contain the running program. There can be more dynamic linking steps in loading programs linked with dynamic or shared libraries.

After the loader has done its work, it jumps to the entry point of the program in the newly created process.

## 14.6.2 SPE Runtime Loader Requirements

### 14.6.2.1 *Static Linking*

An SPE runtime loader can be significantly simpler than a PPE loader because the SPU ABI specifies that programs be statically linked. Static linking results in two major simplifications for an SPE runtime loader: the linker can determine at link time that an SPE program will fit into the SPE's LS, and the SPE loader need not support runtime relocations and dynamic symbol resolution.

### 14.6.2.2 *Resource Conservation*

An SPE loader that conforms to the framework defined here should use as few resources as possible. The loader should use memory in the LS that can be reclaimed when the loaded SPE application starts; no memory should be permanently allocated to the loader. The loader should ensure that the registers it uses—and that are not otherwise defined as initialized in the ABI—are cleared to '0' before it exits. The loader should take a minimum of execution time.

### 14.6.2.3 *Specific Loader Actions*

The SPE runtime loader, either executed by PPE or SPE, must accomplish the following actions:

1. Select an available SPE to execute the program, and perform any needed initialization of its memory and access properties.
2. Transfer the SPE program segments, including overlaying the TOE shadow, from EA space to the selected SPE's LS.
3. Initialize the SPU registers according to the requirements of the *SPU Application Binary Interface Specification*.
4. Enable single-step support if needed.
5. Jump to the SPE application program's entry point (typically `crt0` or an equivalent).

*Select an Available SPE*

This action is handled by the OS running on the PPE or an SPE user-space program. In the latter approach, which has been proven to be a more effective way of loading SPE programs, the PPE first loads a bootstrap loader into the LS, and the bootstrap loader running on the SPE subsequently loads the real SPE program into the LS.

A primary function of the PPE OS is to assign and manage tasks in the system, including tasks for the SPEs. The fundamental responsibilities of the OS for assigning a program to a particular SPE are:

1. Select an SPE to run the program based on the current state of the CBEA processor and other application-dependent criteria.
2. Initialize SPE access properties; this includes setting up the SPE's memory management unit (MMU) and replacement management table (RMT).
3. Initiate SPE program execution by invoking the SPE bootstrap loader.

After the SPE starts executing, the PPE OS is responsible for managing and controlling the progress of the SPE program.

*Transfer Program to LS*

The SPE loader runs either on the PPE or SPE using parameters set up by the PPE loader. The SPE loader uses these parameters to locate the SPE program segments (code, data, and possibly a CESOF toe segment from the TOE shadow area) in EA space and copy them to LS space using MFC DMA commands.

*Initialize SPE Registers*

Certain registers must be initialized before an SPE program is started running. For example, the *Cell Broadband Engine Linux Reference Implementation Application Binary Interface Specification* specifies that R1, the stack pointer (SP), is initialized by the standard `crt0` program entry function to point to the highest quadword address of the SPE's LS. The LS size is 256 KB, so the SP will be initialized to `x'3FFF0'`. In other environments, the initialization address might be less than `x'3FFF0'` to accommodate parameters, isolation open area, and so forth.

The *Cell Broadband Engine Linux Reference Implementation Application Binary Interface Specification* also specifies a convention for passing arguments to an SPE program. These arguments include application arguments and application-environment settings. *Table 14-9* on page 406 shows the registers and their initial values.

*Enable Single-Step*

Single-step execution is required when a debugger is controlling and monitoring an SPU program. Two mechanisms are defined to support SPU single-step operation:

- A privileged-mode register can be set to enable hardware single-step execution. This method is not precise because one or more instructions might be executed before a single-step breakpoint is honored, depending on the fetch and issue rules of the SPU.
- Breakpoints can be set with the SPU **stopd** (stop-and-signal with dependencies) instruction that has the special `x'3FFF'` signal type. This is a method defined by the Cell Broadband Engine Linux Reference ABI and can be as precise as required by replacing each instruction in turn with the **stopd** instruction. In addition, with this mechanism, breakpoints can be set by either the debugger or the programmer.

*Begin Execution*

After the SPE program has been copied to LS and the execution environment has been initialized properly, the SPE loader will set the SPE's Next Program Counter Register (NPC) to the entry point of the program, typically the entry point to `crt0`, and begin execution. The LS address of the entry point is passed to the SPE loader as a parameter (see *Table 14-11* on page 422).

### 14.6.3 Example SPE Runtime Loader Framework Definition

This section describes one possible example of an extensible framework for an SPU runtime loader. This framework uses a small bootstrap loader that runs on the SPU and that copies the SPE application program from main storage into LS and initializes the environment according to the ABI. Many other loader strategies are possible.

With this framework, system programmers can build a compatible replacement for the default SPE loader. The default loader might need to be replaced to accommodate extensions to, or re-definitions of, object-file formats, OS environments, and protocols.

The framework described here defines a set of per-binary, per-loader, and per-thread parameters for the SPU loader; each of these parameter sets is described in the following sections. This description includes the PPE-initiated and SPE-initiated aspects of the loading process and the interactions between the SPE loader and the application's `crt0` entry function.

#### 14.6.3.1 *SPE Runtime Loader Parameters*

The SPE loader framework defines data structures for representing and passing parameters to the SPE loader and application program. All of the EAs for the SPE loader are 64-bit pointers. In 32-bit environments, the high-order 32 bits of an EA are cleared to '0'.

The PPE runtime loader is responsible for copying the SPE loader and its parameters into the target SPE's LS. Thus, when the SPE loader begins execution, all three sets of parameters—per-loader, per-binary, and per-thread—are available in LS. Under the definitions described in the following subsections, the three sets of parameters occupy 112 bytes of LS.

*Per-Loader Parameters*

The per-loader parameters specify the location and size of the SPE bootstrap loader program and the runtime flags the loader should honor. *Table 14-10* explains the per-loader parameters, all of which are 32-bit values.

*Table 14-10. SPE Loader Per-Loader Parameters*

| Parameter | Data Type | Comment |
|---|---|---|
| `ldr_ea_hi` | 32-bit unsigned int | High-order 32 bits of the EA of the SPE loader in main storage. |
| `ldr_ea_low` | 32-bit unsigned int | Low-order 32 bits of the EA of the SPE loader in main storage. |
| `ldr_size` | 32-bit unsigned int | SPE loader program size in bytes. The maximum value for this parameter is currently set at 384 bytes (24 quadwords). |
| `ldr_flags` | 32-bit unsigned int | Runtime flags for the SPE loader.<br>`SPU_LDR_FLAGS_traced (x'1')`: if set, the loader enables single-step operation.<br>`SPU_LDR_FLAGS_multi_seg (x'2')`: if set, the SPE program consists of multiple load segments. |

*Per-Binary Parameters*

The per-binary parameters pass information about the location and size of the SPE program in main storage and in LS. as well as the LS address of the program's entry point. *Table 14-11* on page 422 explains the per-binary parameters; all the parameters are 32-bit values.

*Table 14-11. SPE Loader Per-Binary Parameters*

| Parameter | Data Type | Comment |
|---|---|---|
| `img.size` | 32-bit unsigned int | Image size in bytes. If the `SPU_LDR_FLAGS_multi_seg` flag is set in the per-loader `ldr_flags` parameter (see *Table 14-10* on page 421), `img.size` indicates the number of segments in the list (see `img.ea_hi` and `img.ea_low`) to be loaded into LS. |
| `img.ea_hi` | 32-bit unsigned int | High-order 32 bits of EA of program image in main storage. If the `SPU_LDR_FLAGS_multi_seg` flag is set in the per-loader `ldr_flags` parameter (see *Table 14-10* on page 421), `img.ea_hi` contains the high-order 32-bits of the EA of the list of segments to be loaded into LS. |
| `img.ea_low` | 32-bit unsigned int | Low-order 32 bits of EA of program image in main storage. If the `SPU_LDR_FLAGS_multi_seg` flag is set in the per-loader `ldr_flags` parameter (see *Table 14-10* on page 421), `img.ea_low` contains the low-order 32-bits of the EA of the list of segments to be loaded into LS. |
| `img.dst_ls` | 32-bit unsigned int | LS destination address for the SPE program image segment. |
| `entry` | 32-bit unsigned int | LS address of the SPE program entry point (typically the start of `crt0`). |
| `lslr` | 32-bit unsigned int | LS limit setting; this value is copied from the `spu_env` structure in the SPU Environment Note in the SPE-ELF Program Header; see *Section 14.2.2.4* on page 400. |
| `elf_ea_hi` | 32-bit unsigned int | High-order 32 bits of the EA of the SPE-ELF executable in main storage. |
| `elf_ea_low` | 32-bit unsigned int | Low-order 32 bits of the EA of the SPE-ELF executable in main storage. |

In the simplest case, the SPE program executable contains a single, combined, loadable segment containing the `.text`, `.data`, and `.bss` segments. In this case, the `img.ea_hi`, `img.ea_low`, and `img.size` parameters specify the region of main storage to be copied into LS.

In other cases, particularly for an executable using the capabilities of CESOF, the SPE-ELF executable will have more than one load segment. For multiple load segments, the `img.ea_hi` and `img.ea_low` parameters specify the EA of a list in main storage. The `img.size` parameter specifies the number of elements in the list. The SPE loader uses the list to copy each of the loadable segments from main storage to LS memory.

In the case of a CESOF executable, the list will contain at least two elements: an element for the combined segment that contains the `.text`, `.data`, and `.bss` segments of the SPE program and an element for the toe segment. When the PPE loader builds the list elements, it will replace the main-storage address of the toe segment with the address of the TOE shadow section. Consequently, the SPE loader will overwrite the toe segment in LS with the TOE shadow. This fulfills the CESOF runtime requirement of replacing the `.toe` with the TOE shadow.

Because the framework must support even more sophisticated SPE loaders, the EA in main storage of the complete SPE-ELF executable image is available in the `elf_ea_hi` and `elf_ea_low` parameters.

### *Per-Thread Parameters*

The per-thread parameters are arguments passed to the SPE program because the running instance of the program is an OS thread. These parameters are passed to satisfy the ABI register-initialization requirements described in *Section 14.6.2.3* on page 419, and are explained in *Table 14-9* on page 406.

The `spu_lscsa` in R6 can be used by the PPE OS when it needs to preempt the SPE thread, or the thread itself can use the area to perform a light-weight SPE context save. An SPE thread can perform a context save in response to a privileged-attention event raised by the PPE OS (see *Section 18* on page 471), or the thread can save context as a means to communicate system-call parameters from LS or SPU registers to the PPE OS.

#### 14.6.3.2 *PPE-Initiated Actions for SPE Loading*

The first steps in loading an SPE program into an SPE LS for execution include loading the SPE-ELF executable image from the CESOF object and loading the SPE loader, called `spu_ld.so` in this description, into main storage. These steps are shown graphically in *Figure 14-8*.

*Figure 14-8. PPE OS Loads CESOF Object and spu_ld.so into EA Space*

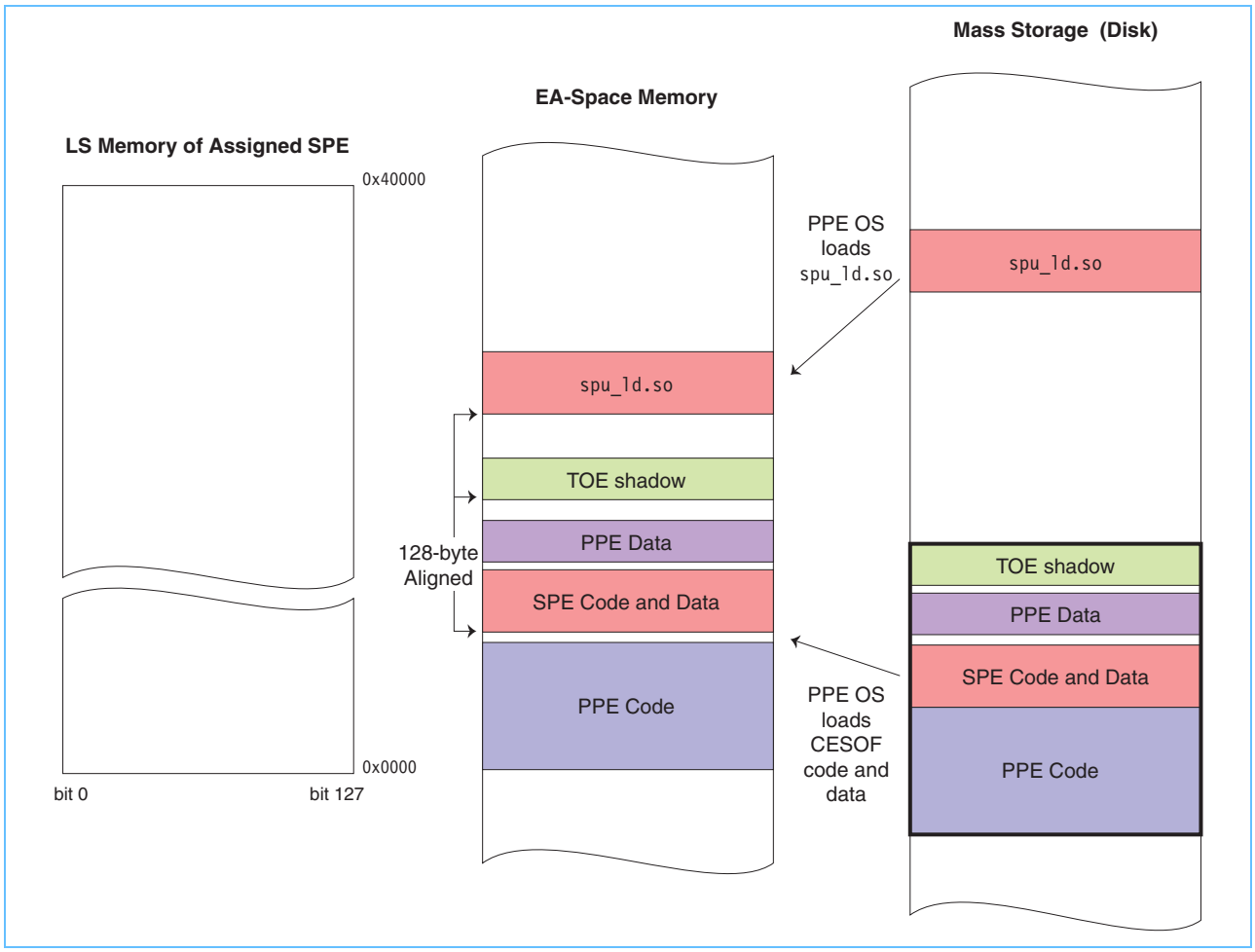
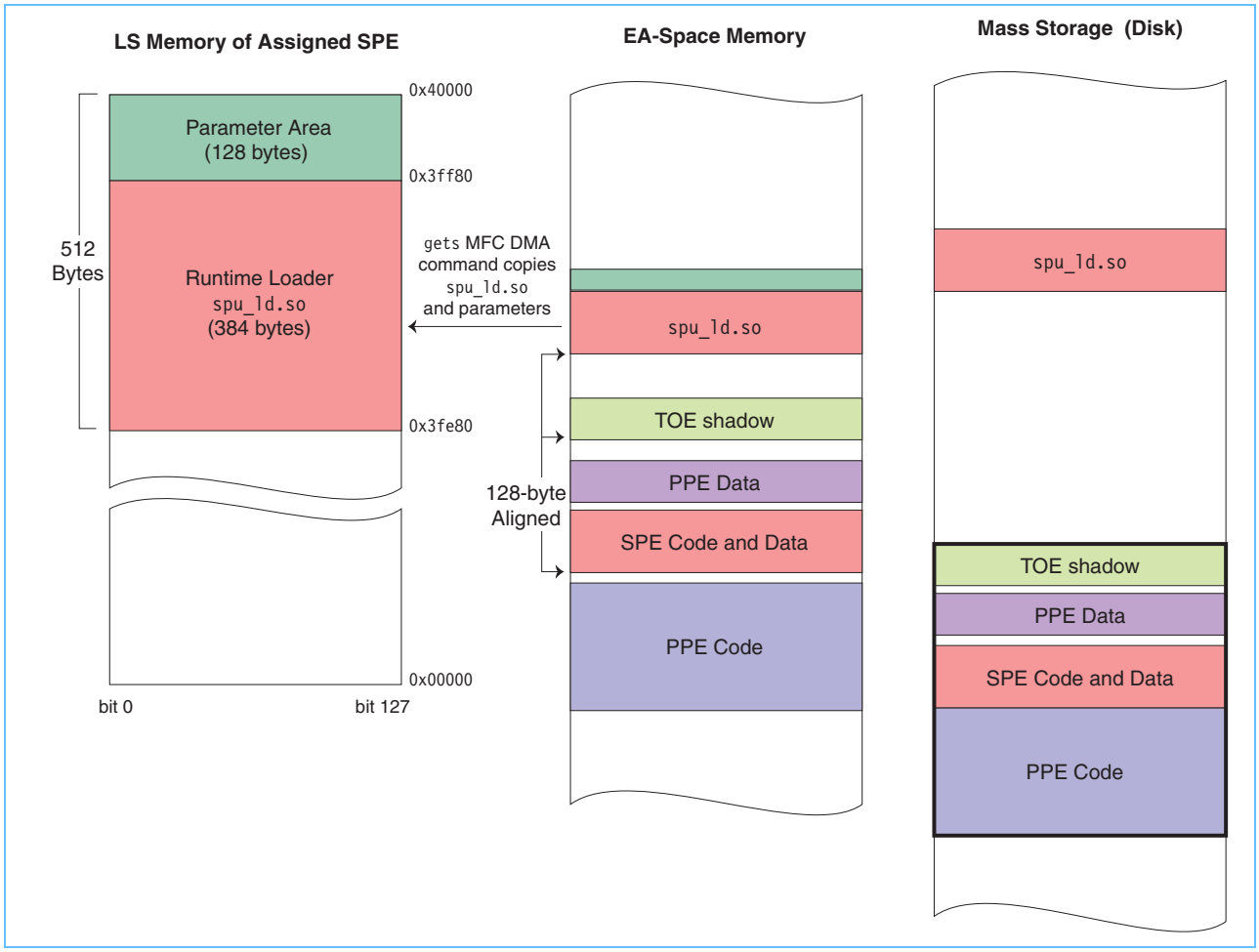



*Figure 14-9* on page 424 shows the location and sizes of the loader and its parameters after they have been copied to LS.

*Figure 14-9. PPE OS Loads `spu_ld.so` and Parameters into SPE LS*




Before the PPE loader can copy the SPE loader and set it running, the PPE OS must assign the SPE program to an available SPE. The assigned SPE's state, including its LS and registers, must be cleaned (for example, initialized to zero), and the SPE's access and memory properties (such as the contents of the MMU) must be properly initialized.

If the PPE loader detects that it is loading an SPE program that has a toe segment, it must create a list in main storage of loadable segments that will cause the SPE loader to copy the TOE shadow from the PPE-ELF wrapper to LS instead of the toe segment from the SPE-ELF executable. Creating the segment list and copying the TOE shadow over the toe segment are the major aspects of the runtime support required by the CESOF specification.

Next, the PPE loader sets the SPE's NPC register to the loader entry point, which, as shown in *Table 14-12* on page 425, is `SPU_LDR_PROG_start` with the value (`LS_SIZE` – `512`). Because `LS_SIZE` is 256 KB, the NPC will be set to `x'3FE00'`.

The next step is invoking the SPE loader with the proper parameters. The PPE loader copies the SPE boot-strap loader and its parameters to LS according to the values shown in *Table 14-12* on page 425.

*Table 14-12. LS Locations and Sizes of SPE Loader and Parameters*

| Attribute | Value | Comment |
|---|---|---|
| SPU_LDR_PROG_start | LS_SIZE – 512 | The SPE loader is copied by the PPE loader into LS at the address of the start of the last 512-byte (32-quadword) block of implemented LS memory. For the CBEA processors, for which LS_SIZE is x'40000', this address is x'3FE00'. |
| SPU_LDR_PROG_size | 384 | The current framework allows the SPE loader to occupy up to 384 bytes (24 quadwords) of LS starting at SPU_LDR_PROG_start. |
| SPU_LDR_PARAMS_start | LS_SIZE – 128 | The SPE loader parameters are copied by the PPE loader into LS at the address of the start of the last 128-byte (8-quadword) block of implemented LS memory.For the CBEA processors, for which LS_SIZE is x'40000', this address is x'3FF80'. |
| SPU_LDR_PARAMS_size | 128 | The current framework allows the SPE loader parameters to occupy up to 128 bytes (8 quadwords) of LS starting at SPU_LDR_PARAMS_start. |

The PPE loader copies the SPE loader parameters to LS by initiating a DMA **get** command in the SPE's associated MFC (see *Section 19* on page 513 for details on DMA commands). Next, the PPE loader copies the SPE loader code to LS by initiating a DMA **gets** command. The **gets** command copies the code and then, after the DMA transfer has completed, starts the SPE's SPU running at the address in the NPC register, which is set to x'3FE00'.

At this point, the state of LS is shown in *Figure 14-9* on page 424, and the work of the PPE loader is finished; the SPE bootstrap loader completes the remaining steps in the loading process. A principle advantage of dividing the work of SPE program loading between the PPE loader and the SPE loader is the fact that after the PPE loader enqueues the **get** and **gets** commands, it need not wait for them to complete; the PPE loader can continue with other aspects of the PPE loading process or terminate and free the PPE for other tasks.

#### 14.6.3.3 *SPE-Initiated Actions for SPE Loading*

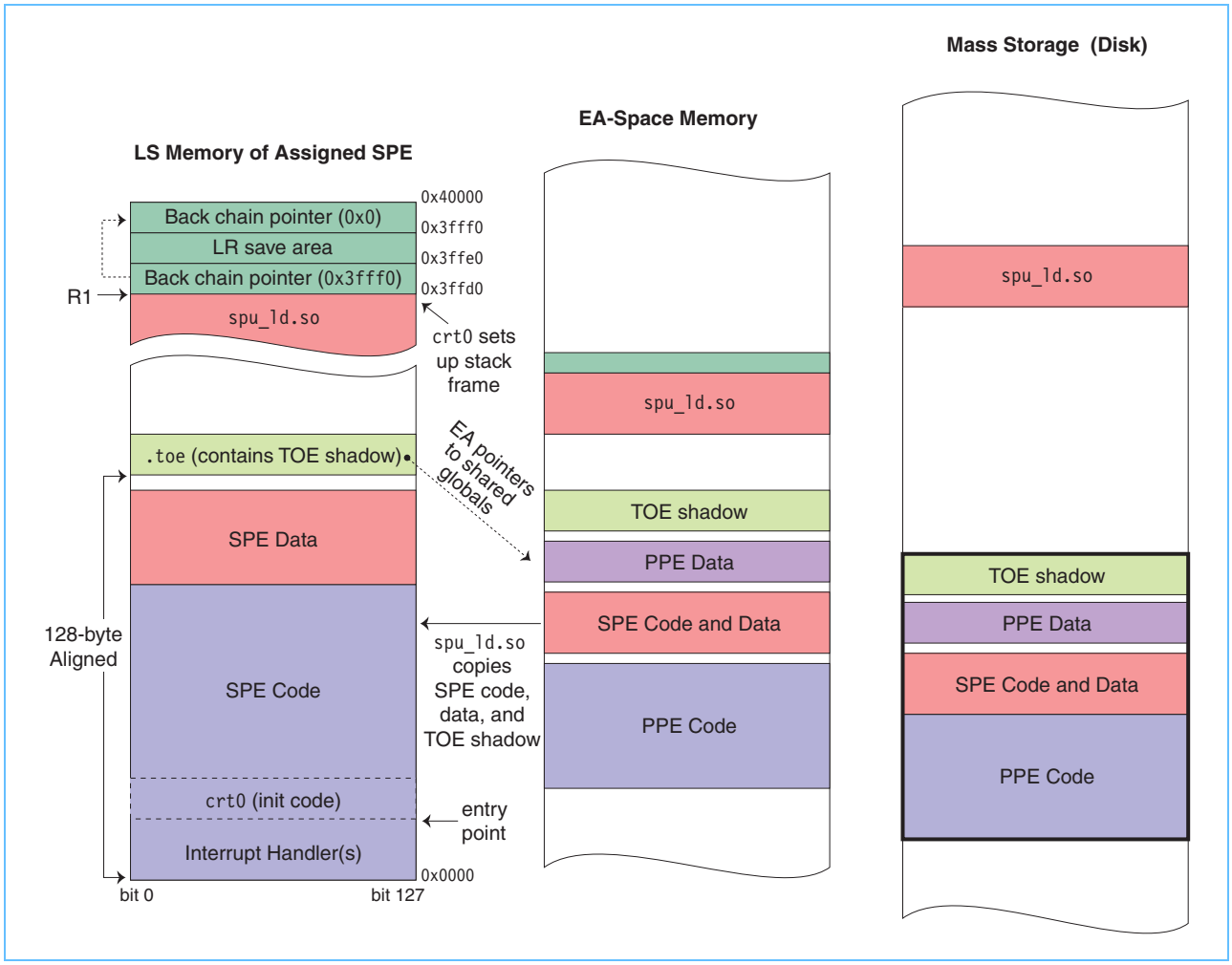
The SPE bootstrap loader must copy the SPE program segments to LS, initialize registers, possibly enable single-step mode (see *Section 14.6.2.3* on page 419), and finally start the SPE executing the loaded program.

*Copy SPE Program Segments to LS*

The sizes and locations of the program segments to be copied are contained in the per-binary section of the parameter area. The parameters can specify either a single loadable segment or a list of segments. If a list of segments is needed to properly load the SPE program, the PPE loader is responsible for setting up the list in main storage. See *Section 14.6.3.1* on page 421 for additional information about per-binary parameters. As summarized in *Section* on page 418, copying the TOE shadow segment is a key part of CESOF runtime support.

In the case of loading a CESOF executable, the final state of LS—with the SPE application loaded and the TOE shadow copied in place of the toe segment—is shown in *Figure 14-10* on page 426.

*Figure 14-10. `spu_ld.so` Loads SPE Code and TOE Shadow*




*Initialize Registers*

See *Section 14-9* on page 406 for details on how R1 (stack pointer) and R2 (stack size) are set. The SPE loader initializes registers R3 through R5 according to the per-thread parameters. See *Section 14.6.2.3* on page 419 for more information.

*Enable Single Step*

This SPE loader framework implements single-stepping and debugger breakpoints with the **stop** instruction. The return code of the **stop** instruction for a breakpoint is `x'3FFF'`.

If the `SPU_LDR_FLAGS_traced` bit is set to '1' in the per-loader `ldr_flags` parameter, the SPE program will be executed in single-step mode. To support single-step mode, the SPE loader must execute the **stop** instruction just before the loader branches to the SPE program's entry point. When controlling software, such as an SPE debugger, takes control after this breakpoint is raised, this software can set additional breakpoints or otherwise modify the SPE program state.

*Begin SPE Program Execution*

The last step is to transfer control to the SPE program's entry point. The address of the entry point is found in the `entry` parameter of the per-binary parameter area. The `entry` parameter typically contains the starting address of the program's `crt0` function. If single-step mode is enabled, the loader executes a breakpoint **stop** instruction before it branches to `crt0`; otherwise it simply branches to `crt0`.

#### 14.6.3.4 *SPE Loader Interactions with `crt0`*

The PPE loader clears to zero the assigned SPE's registers and LS before the SPE loader and its parameters are copied to LS. When the SPE loader has completed its work, the SPE application program is in LS and registers are initialized according to the ABI requirements.

All other LS regions and SPU registers contain zero, except for the last 512 bytes (32 quadwords) of LS, which is where the SPE loader and parameters are located. Although `crt0` need not clear regions of LS that contain uninitialized data segments (the PPE loader already cleared all of LS), some systems might be designed to have `crt0` clear the 512-byte area used by the SPE loader. Explicit clearing of the SPE loader area is probably not necessary, however, because `crt0` is responsible for initializing the SPE stack pointer to the highest LS memory address and setting up the initial stack frame. Consequently, function calls in the SPE application will overwrite the SPE loader and parameters. In any case, it is unlikely that `spu_ld.so` and its parameters will contain any sensitive information.

## 14.7 SPE Execution Environment

This section specifies low-level system information and conventions that are used by the operating system. Some of these conventions are defined by the ABI and are standard for all systems. Others are specific to the Linux operating system environment and are indicated as such. For specific additional details, consult the operating-environments documentation for your specific operating system.

### 14.7.1 Signal Types for the SPE Stop-and-Signal Instruction

The **stop** instruction contains a 14-bit signal-type field. When **stop** is executed, the signal type is written to bits 0 through 13 of the SPE status register. This value is used to indicate why a program stopped. The signal-type values have been partitioned by Linux in the following way:

- Signal-type values with the most-significant bit cleared to '0' are reserved for application use.
- Signal-type values with the most-significant bit set to '1' are reserved for runtime or privileged services.

*Table 14-13* on page 428 describes the reserved signal type values.

*Table 14-13. SPE-ELF Reserved Signal Type Values*

| Signal-Type Value | Description |
|---|---|
| x‘0000’ | Data executed as an instruction. |
| x‘2000’..x‘20FF’ | Linux return from `main()` or `exit()`. Return or `exit()` status is encoded in the least-significant byte of the stop-and-signal value:<br>`exit (EXIT_SUCCESS) == x‘2000’`<br>`exit (EXIT_FAILURE) == x‘2001’` |
| x‘2100’..x‘21FF’ | Linux PPE assisted library calls. See *Cell Broadband Engine Linux Reference Implementation Application Binary Interface Specification* for additional details. |
| x‘3FFE’ | Stack overflow detected. |
| x‘3FFF’ | Debugger breakpoint. |

# 15. Power and Thermal Management

The power-management mechanism provides a means for limiting the amount of power dissipated by a component by limiting its degree of activity. There are five states for power management. Typically, the more aggressive the power-management state, the more time is required to enter and exit the state. These states apply to each of the various processing elements and other logic blocks of the Cell Broadband Engine Architecture (CBEA) processors.[1] Each component can be individually set to a power-management state. The power-management facilities provide privileged software with the ability to manage power while meeting the real-time demands of a system. The CBEA processors implement a subset of the power-management states defined by the *Cell Broadband Engine Architecture*.

The digital thermal-management mechanism provides a means for monitoring and controlling the temperature within the CBEA processor. During the execution of an application, the temperature of areas within the CBEA processor can rise. Left unchecked, the temperature might rise above the maximum specified junction temperature leading to improper operation or physical damage.

Both the power-management and thermal-management facilities are part of the pervasive logic, which provides control, monitoring, and debugging functions.

## 15.1 Power Management

*Table 15-1* summarizes the CBEA processor power-management states. Details on each state are given later in this section.

*Table 15-1. Power-Management States*

| Power Management State | Description |
|---|---|
| Slow State | In Slow state, clock frequency is lowered. A lower clock frequency results in greater potential for performance degradation, and lower power consumption. For more information, see *Section 15.1.1 Slow State* on page 430. |
| PPE Pause (0) State | PowerPC Processor Element (PPE) Pause (0) state is a clock-gating that occurs when both PPE threads are shut down. PPE code execution is suspended during this state. For more information, see *Section 15.1.2 PPE Pause (0) State* on page 431. |
| SPE State Retained and Isolated (SRI) State | In Synergistic Processor Element (SPE) SRI state, all access to the component is inhibited. The state remaining on the component is retained. The component must be prepared by privileged software or hardware to maintain system integrity. The component does not make forward progress. For more information, see the security documentation, available under nondisclosure agreement (NDA). |
| SPU Pause State | In synergistic processor unit (SPU) Pause state, SPU code execution can be suspended to reduce power and stop forward progress until resumed. Memory-mapped I/O (MMIO) registers and the local storage (LS) can still be accessed. For more information, see *Section 15.1.3 SPU Pause State* on page 432. |
| MFC Pause State | In MFC Pause state, memory flow controller (MFC) command queue operations can be suspended to reduce power and to stop forward progress until resumed. For more information, see *Section 15.1.4 MFC Pause State* on page 432. |
| Active State | In Active state, no power-management functions are enabled. This is the state of maximum performance and maximum power-use. |

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

### 15.1.1 Slow State

The main clock (NClk) frequency of the CBEA processors can be changed dynamically in the Slow state. The functions of the CBEA processors are not affected in this state; only the performance is affected. Certain time-base settings are guaranteed to be invariant over all frequency changes. See *Section 13 Time Base and Decrementers* on page 381 for details. Slow state can be run in conjunction with any other power-management state.

#### 15.1.1.1 *Configuration-Ring Settings*

Slow state has one related configuration-ring setting, `slow_mode_delay_setting`. The recommended setting is x'1F'. See the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide* for a description of the configuration ring.

#### 15.1.1.2 *PLL Start*

By default, the CBEA processors start at the maximum frequency. However, if the phase-locked loop (PLL) is configured to not start at the maximum frequency, an extra step is required before scaling the NClk frequency. This required step is to write the Power Management Control Register (`PMCR`) as follows: `PMCR[Override]` = '1' and `PMCR[Curr_Slow_n]` = PLL frequency division. This can be done with one MMIO write. Software must not initiate a frequency scale until a read of the Power Management Status Register (`PMSR[BE_Slow_n]`) reflects the PLL frequency division.

#### 15.1.1.3 *Scaling the Processor Core Clock (NClk) Frequency*

Slow state changes the processor core clock (`NClk`) frequency (see *Section 13.2.3 Time-Base Frequency* on page 383) to a target frequency rate (n). Changing NClk requires a compensating factor (command-pacing rate) in the memory interface controller (MIC) unit to prevent buffer overruns. This is required because the MIC clock (MiClk) runs separately and does not change with this procedure. An initial setting of the `MIC_Slow_Fast_Timer_n` register sets the command-pacing rate during the frequency change. This value is constant for all frequencies and needs to be set up only once. The `MIC_Slow_Next_Timer_n` register sets the command pacing rate after the frequency change; this value is different for each target frequency (n). After all pacing rates are set, a write to the `PMCR[Next_Slow_n]` register field initiates a frequency change. No writes to `MIC_Slow_Fast_Timer_n`, `MIC_Slow_Next_Timer_n`, or `PMCR[Next_Slow_n]` occur until the status of `PMSR[BE_Slow_n]` reflects the target frequency.

To dynamically change to a supported frequency, perform the following steps:

> ***Note:*** *If the PLL is not configured to start at the maximum frequency, perform the step described in Section 15.1.1.2 before performing these steps.*

1. See *Table 15-2* on page 431 for the required target frequency and corresponding recommended register values.
2. Software sets `MIC_Slow_Fast_Timer_0[8:19]` and `MIC_Slow_Fast_Timer_1[8:19]` to x'7FC'. (Software only has to write these registers once.)
3. Software sets the new `MIC_Slow_Next_Timer_0[16:27]`, and `MIC_Slow_Next_Timer_1[16:27]` to the value indicated in *Table 15-2* for the required target frequency.

4. Software sets the new frequency PMCR[Next_Slow_n] to the value indicated in *Table 15-2* for the required target frequency.
5. Software waits for a read of PMSR[BE_Slow_n] to reflect the updated frequency change. Status must reflect the updated frequency before attempting another frequency change.

*Table 15-2* shows the register updates corresponding to required target frequencies for Slow state. Some CBEA processors might not be certified for execution in all of the slow-state frequencies shown in *Table 15-2*. Consult the processor's data sheet for supported frequencies.

*Table 15-2. Register Updates for Slow State*

| PMCR[Next_Slow(n)] [n = 0:7] | MIC_Slow_Fast_Timer_n[8:19] [n = 0,1] | MIC_Slow_Next_Timer_n[16:27] [n = 0,1] |
|---|---|---|
| n = 0 Frequency = Max | x‘7FC’ | x‘240’ |
| n = 1 Frequency = Max/2 | x‘7FC’ | x‘268’ |
| n = 2 Frequency = Max/3 | x‘7FC’ | x‘29C’ |
| n = 3 Frequency = Max/4 | x‘7FC’ | x‘2D0’ |
| n = 4 Frequency = Max/5 | x‘7FC’ | x‘300’ |
| n = 5 Frequency = Max/6 | x‘7FC’ | x‘334’ |
| n = 6 Frequency = Max/8 | x‘7FC’ | x‘39C’ |
| n = 7 Frequency = Max/10 | x‘7FC’ | x‘3FC’ |

**Note:** MIC_Slow_Next_Timer_n table data assumes a NCLK/MiClk ratio of 2:1. For an equation with different ratios see in the *Cell Broadband Engine Registers* document.

For details about the registers referenced in this section, see the *Cell Broadband Engine Registers* document.

### 15.1.2 PPE Pause (0) State

PPE Pause (0) state requires both PPE threads to be shut down and pause state-control is enabled. Cache coherency is maintained during this state. PPE Pause (0) state is exited with interrupts, some of which can be masked individually. PPE Pause (0) state can be run in conjunction with any other power-management state.

#### 15.1.2.1 *Enabling PPE Pause (0) State*

Control of PPE Pause (0) is disabled by default. A one-time write to PMCR[PPE_pause] = ‘1’ is required to enable the state.

#### 15.1.2.2 *Entering PPE Pause (0) State*

Each PPE thread must be shut down for the PPE Pause (0) state to be enabled. After both threads are shut down, the PPE Pause (0) state is entered if no interrupt is pending. See *Section 10 PPE Multithreading* on page 299 for details about resuming and suspending a thread.

#### 15.1.2.3 *Exiting Pause State*

PPE Pause (0) state exits by one of the following types of interrupts: external, decrementer, thermal management, or system error. Pause control masking is possible for external and decrementer interrupts. See *Section 15.1.2.4 Pause Interrupt Resume Masks* for information about masking.

It is possible to exit the PPE Pause (0) state and not wake a thread, because there are separate PPE thread controls. See *Section 10 PPE Multithreading* on page 299 for details about resuming and suspending a thread.

***Note:*** *A pending hypervisor decrementer (HDEC) interrupt prevents entering the PPE Pause (0) state. However, once the state is activated, an HDEC interrupt does not exit the PPE Pause (0) state.*

#### 15.1.2.4 *Pause Interrupt Resume Masks*

The external and decrementer interrupts (per thread) can be selectively masked to determine whether they cause an exit from PPE Pause (0) state. The mask must be set before both threads are shut down by setting a combination of any mask bits located in the Thread Switch Control Register (TSCR) at `TSCR[wdec0]`, `TSCR[wdec1]`, or `TSCR[wext]`. This mask is only for interrupts to exit PPE Pause (0) state; it is possible to exit PPE Pause (0) state and not wake a thread because these masks are different. See *Section 10 PPE Multithreading* on page 299 for details.

### 15.1.3 SPU Pause State

SPU Pause state can be run in conjunction with the Slow, PPE Pause (0), and MFC Pause states. The MMIO registers and the LS can still be accessed in this state. Cache coherency is maintained in this state.

Software pauses (suspends) an SPU by writing `MFC_SR1[S]` = '0'. Software resumes an SPU by writing `MFC_SR1[S]` = '1'.

### 15.1.4 MFC Pause State

MFC Pause state can be run in conjunction with the Slow, PPE Pause (0), and SPU Pause states. The MMIO registers and the LS can still be accessed in this state.

Software pauses (suspends) an MFC's command queues by writing `MFC_CNTL[Sc]` = '1'. Cache coherency is maintained in this state. Software resumes an MFC's command queues by writing `MFC_CNTL[Sc]` = '0'.

## 15.2 Thermal Management

The digital thermal-management unit consists of a thermal-management control unit (TMCU) and ten distributed Digital Thermal Sensors (DTSs). One sensor is located in each of the eight SPEs, one is located in the PPE, and one is adjacent to the linear thermal diode. The linear thermal diode is an on-chip diode that calculates temperature. These sensors are adjacent to areas within the associated unit that typically experience the greatest rise in temperature during the execution of most applications.

The TMCU monitors feedback from each of the sensors. If the temperature of a sensor rises above a programmable point, the TMCU can be configured to cause an interrupt to the PPE and dynamically throttle the execution of the associated PPE or SPE. The throttling is accomplished by stopping and running the PPE or SPE for a programmable number of cycles. The interrupt allows privileged software to take corrective action while dynamic throttling (a hardware device) attempts to keep the temperature below a programmable level without software intervention. Privileged software must set the throttling level equal to or below the recommended settings.

If dynamic throttling or privileged software does not effectively manage the temperature and the temperature continues to rise, the CBEA processor clocks are stopped when the temperature reaches a thermal overload temperature defined by the configuration ring data. The thermal overload feature protects the CBEA processors from physical damage. Recovery from this condition requires a hard reset.

***Note:*** *The temperature of the region monitored by the digital thermal sensors is not necessarily the hottest point within the associated PPE or SPE.*

### 15.2.1 Thermal-Management Operation

In the TMCU, each DTS provides a current temperature-detection signal. This signal indicates that the temperature is equal to or below the current temperature-detection range set by the TMCU. The TMCU uses the state of these signals to continually track the temperature of the PPE's or each SPE's digital thermal sensor. As the temperature is tracked, the TMCU provides the current temperature as a numeric value that represents the temperature within the PPE or SPE. The mapping between this numeric value and the temperature is provided in *Table 15-3*. Internal calibration storage is set in manufacturing to calibrate the individual sensors.

*Table 15-3. Digital Temperature-Sensor Encoding (Sheet 1 of 2)*

| Thermal-Sensor Temperature Encoding | | | |
|---|---|---|---|
| 6-Bit Encode | Temperature Range | 6-Bit Encode | Temperature Range |
| 0 | $temp < 65.0°C$ | 17 | $97.0°C \le temp < 99.0°C$ |
| 1 | $65.0°C \le temp < 67.0°C$ | 18 | $99.0°C \le temp < 101.0°C$ |
| 2 | $67.0°C \le temp < 69.0°C$ | 19 | $101.0°C \le temp < 103.0°C$ |
| 3 | $69.0°C \le temp < 71.0°C$ | 20 | $103.0°C \le temp < 105.0°C$ |
| 4 | $71.0°C \le temp < 73.0°C$ | 21 | $105.0°C \le temp < 107.0°C$ |
| 5 | $73.0°C \le temp < 75.0°C$ | 22 | $107.0°C \le temp < 109.0°C$ |
| 6 | $75.0°C \le temp < 77.0°C$ | 23 | $109.0°C \le temp < 111.0°C$ |
| 7 | $77.0°C \le temp < 79.0°C$ | 24 | $111.0°C \le temp < 113.0°C$ |
| 8 | $79.0°C \le temp < 81.0°C$ | 25 | $113.0°C \le temp < 115.0°C$ |
| 9 | $81.0°C \le temp < 83.0°C$ | 26 | $115.0°C \le temp < 117.0°C$ |
| 10 | $83.0°C \le temp < 85.0°C$ | 27 | $117.0°C \le temp < 119.0°C$ |
| 11 | $85.0°C \le temp < 87.0°C$ | 28 | $119.0°C \le temp < 121.0°C$ |
| 12 | $87.0°C \le temp < 89.0°C$ | 29 | $121.0°C \le temp < 123.0°C$ |
| 13 | $89.0°C \le temp < 91.0°C$ | 30 | $123.0°C \le temp < 125.0°C$ |

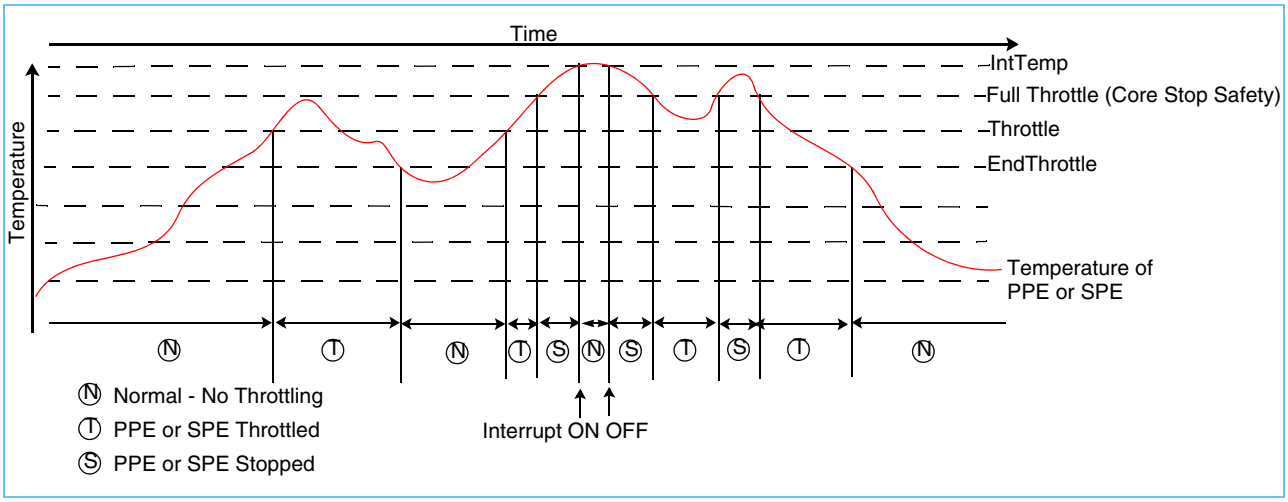
*Table 15-3. Digital Temperature-Sensor Encoding (Sheet 2 of 2)*

| Thermal-Sensor Temperature Encoding | | | |
|---|---|---|---|
| 6-Bit Encode | Temperature Range | 6-Bit Encode | Temperature Range |
| 14 | $91.0°C \le temp < 93.0°C$ | 31-63 | $125.0°C \le temp$ |
| 15 | $93.0°C \le temp < 95.0°C$ | | |
| 16 | $95.0°C \le temp < 97.0°C$ | | |

The TMCU can be configured to cause an interrupt to the PPE and dynamically throttle the execution of the PPE or an SPE. The TMCU compares the numeric value representing the temperature to a programmable interrupt temperature and a programmable throttle point. If the temperature is within the programmed interrupt temperature range, an external interrupt is generated to the PPE, if enabled. In addition, a second programmable interrupt temperature can cause the assertion of an Attention signal to a system controller. See *Section 15.2.5 Thermal Sensor Interrupt Registers* on page 436 for more information about the interrupt capability of the TMCU. If the temperature is equal to or above the throttling point, the TMCU throttles the execution of the PPE or an SPE by starting and stopping the PPE or SPE dynamically. Software can control the ratio and frequency of the throttling using the Dynamic Thermal-Management registers. See *Section 15.2.6 Dynamic Thermal-Management Registers* on page 438 for more information.

*Figure 15-1* illustrates the temperature and the various points at which interrupts and dynamic throttling occur. In this figure, the PPE or SPE is running normally; there is no throttling in the regions marked with an “N”. When the temperature reaches the throttle point, the TMCU starts throttling the execution of the associated PPE or SPE. The regions in which the throttling occurs are marked with a “T”. When the temperature of the PPE or SPE drops below the end throttle point, the execution returns to normal operation. If, for any reason, the temperature continues to rise and reaches a temperature at or above the full throttle point, the PPE or SPE is stopped until the temperature drops below the full throttle point. Regions where the PPE or SPE is stopped are marked with an “S”. Stopping the PPE or SPEs when the temperature is at or above the full throttle point is referred to as the Core Stop Safety. In this illustration, the interrupt temperature is set above the throttle point; therefore, software is notified if the PPE or SPE is ever stopped for this condition, provided that the Thermal Interrupt Mask Register (`TM_ISR`) is set to active, allowing the PPE or SPE to resume during a pending interrupt.

*Figure 15-1. Digital Thermal-Sensor Throttling*



***Note:*** *If dynamic throttling is disabled, privileged software must manage the thermal condition. Not managing the thermal condition can result in improper operation of the associated PPE or SPE or a thermal shutdown by the thermal-overload function.*

### 15.2.2 Configuration-Ring Settings

Thermal management uses the following Configuration-Ring settings:

- `MinStopPPE` specifies the minimum stop time for the PPE. Suggested value: x‘8’.
- `MinStopSPE` specifies the minimum stop time for the SPEs. Suggested value: x‘8’.
- `CFG_TO` specifies the temperature which will cause the Thermal Overload signal to be asserted and the clocks to be stopped. Suggested value: x‘1F’.
- `SenSamp`Time specifies how often a sensor will be sampled. Suggested value: x‘01’.
- `DigFiltDly` specifies the delay time for the sensor-output’s digital filter. Suggested value: x‘0’.

For a description of the configuration ring, see the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

### 15.2.3 Thermal Registers

See the *Pervasive MMIO Registers* section of the *Cell Broadband Engine Registers* document for bit definitions of the thermal-management MMIO registers mentioned in the following sections. See the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide* for a description of the configuration ring.

### 15.2.4 Thermal Sensor Status Registers

The Thermal Sensor Status Registers consist of the Thermal Sensor Current Temperature Status Registers (`TS_CTSR1` and `TS_CTSR2`) and the Thermal Sensor Maximum Temperature Status Registers (`TS_MTSR1` and `TS_MTSR2`). These registers allow software to read the current temperature of each DTS, determine the highest temperature reached during a period of time, and cause an interrupt when the temperature reaches a programmable temperature. System software should mark the real address pages containing the thermal sensor registers as hypervisor privileged.

`TS_CTSR1` and `TS_CTSR2` contain the encoding for the current temperature of each DTS. `TS_MTSR1` and `TS_MTSR2` contain the encoding for the maximum temperature reached for each sensor from the time of the last read of these registers. See the *Cell Broadband Engine Registers* specification for a definition of the encoding. The length of a sample period is controlled by the `SenSampTime` configuration field, as described in the *Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide*.

### 15.2.5 Thermal Sensor Interrupt Registers

The thermal sensor interrupt registers control the generation of a thermal-management interrupt to the PPE. This set of registers consists of the Thermal Sensor Interrupt Temperature Registers (`TS_ITR1` and `TS_ITR2`), the Thermal Sensor Interrupt Status Register (`TS_ISR`), the Thermal Sensor Interrupt Mask Register (`TS_IMR`), and the Thermal Sensor Global Interrupt Temperature Register (`TS_GITR`).

TS_ITR1, TS_ITR2, and TS_GITR contain the encoding for the temperature that causes a thermal-management interrupt to the PPE. See the *Cell Broadband Engine Registers* specification for a definition of the encoding.

When the temperature encoding in TS_CTSR1 and TS_CTSR2 for a sensor is greater than or equal to the corresponding sensor's interrupt temperature encoding in TS_ITR1 and TS_ITR2, the corresponding status bit (TS_ISR[Sx]) is set. When the temperature encoding in TS_CTSR1 and TS_CTSR2 for any sensor is greater than or equal to the global interrupt temperature encoding in TS_GITR, the corresponding status bits (TS_ISR[Gx]) are set.

If any TS_ISR[Sx] bit is set and the corresponding mask bit (TS_IMR[Mx]) is also set, a thermal-management interrupt signal is asserted to the PPE. If any TS_ISR[Gx] bit is set and the corresponding mask bit (TS_IMR[Cx]) is also set, a thermal-management interrupt signal is asserted to the PPE. To clear the interrupt condition, privileged software should clear any corresponding mask bits in TS_IMR to '0'.

To enable a thermal-management interrupt, privileged software must ensure that the temperature is below the interrupt temperature for the corresponding sensors and then perform the following sequence; enabling an interrupt when the temperature is not below the interrupt temperature can result in an immediate thermal-management interrupt's being generated:

1. Write a '1' to the corresponding status bit in TS_ISR.
2. Write a '1' to the corresponding mask bit in TS_IMR.

#### 15.2.5.1 *Thermal Sensor Interrupt Temperature Registers*

The Thermal Sensor Interrupt Temperature Registers (TS_ITR1 and TS_ITR2) contain the interrupt temperature level for the sensors located in the SPEs, PPE, and adjacent to the linear thermal diode. The encoded interrupt temperature levels in this register are compared to the corresponding interrupt temperature encoding in TS_CTSR1 and TS_CTSR2. The results of these comparisons are used to generate a thermal-management interrupt. Each sensor's interrupt temperature level is independent.

#### 15.2.5.2 *Thermal Sensor Global Interrupt Temperature Register*

In addition to the independent interrupt temperature levels set in TS_ITR1 and TS_ITR2, the Thermal Sensor Global Interrupt Temperature Register (TS_GITR) contains a second interrupt temperature level. This level applies to all sensors in the CBEA processors. The encoded global interrupt temperature level in this register is compared to the current temperature encoding for each sensor. The results of these comparisons are used to generate a thermal-management interrupt.

The intent of the global interrupt temperature is to provide an early indication of a temperature rise in the CBEA processors. Privileged software and the system controller can use this information to start actions to control the temperature, for example, increasing the fan speed, rebalancing the application software across units, and so on.

#### 15.2.5.3 *Thermal Sensor Interrupt Status Register*

The Thermal Sensor Interrupt Status Register (TS_ISR) identifies which sensors meet the interrupt conditions. An interrupt condition refers to a particular condition that each TS_ISR bit has that, when met, makes it possible for an interrupt to occur. An actual interrupt is only presented to the PPE if the corresponding mask bit is set in the Thermal Sensor Interrupt Mask Register (TS_IMR).

TS_ISR contains three sets of status bits—the Digital Sensor Global Threshold Interrupt status bit (TS_ISR[Gx]), the Digital Sensor Threshold Interrupt status bit (TS_ISR[Sx]), and the Digital Sensor Global Below Threshold Interrupt status bit (TS_ISR[Gb]).

Hardware sets TS_ISR[Sx] when the temperature encoding for a sensor in TS_CTSR1 and TS_CTSR2 is greater than or equal to the corresponding sensor's interrupt temperature encoding in TS_ITR1 and TS_ITR2 and the corresponding direction bit TM_IMR[Bx] = '0'. Also, hardware sets TS_ISR[Sx] when the temperature encoding for a sensor in TS_CTSR1 and TS_CTSR2 is below the corresponding sensor's interrupt temperature encoding in TS_ITR1 and TS_ITR2 and the corresponding direction bit TM_IMR[Bx] = '1'.

Hardware sets TS_ISR[Gx] when any participating sensor's current temperature is greater than or equal to that of TS_GITR and TS_IMR[BG] is cleared to '0'. The individual TS_ISR[Gx] bits indicate which individual sensors meet these conditions.

Hardware sets TS_ISR[Gb] when all of the participating sensors in TS_IMR[Cx] have a current temperature below that of TS_GITR and the TS_IMR[BG] is set to '1'. Because all participating sensors must have a current temperature below that of TS_GITR, there is only one status bit (TS_ISR[Gb]) for a global below-threshold interrupt condition.

Once a status bit (TS_ISR[Sx], [Gx], or [Gb]) is set to '1', this state is maintained until reset to '0' by privileged software. Privileged software resets a status bit to '0' by writing a '1' to the corresponding bit in TS_ISR.

#### 15.2.5.4 *Thermal Sensor Interrupt Mask Register*

The Thermal Sensor Interrupt Mask Register (TS_IMR) contains two fields for individual sensors and multiple fields for global interrupt conditions. An interrupt condition refers to a particular condition that each TS_IMR bit has that, when met, makes it possible for an interrupt to occur. An actual interrupt is only presented to the PPE if the corresponding mask bit is set.

The two TS_IMR Digital Thermal Threshold Interrupt fields for individual sensors are TS_IMR[Mx] and TS_IMR[Bx]. The TS_IMR[Mx] mask bits prevent an interrupt status bit from generating a thermal-management interrupt to the PPE. The TS_IMR[Bx] directional bits set the temperature direction for the interrupt condition above or below the corresponding temperature in TS_ITR1 and TS_ITR2. Setting TS_IMR[Bx] to '1' sets the temperature for the interrupt condition to be below the corresponding temperature in TS_ITR1 and TS_ITR2. Clearing TS_IMR[Bx] to '0' sets the temperature for the interrupt condition to be equal to or above the corresponding temperature in TS_ITR1 and TS_ITR2.

The TS_IMR fields for the global interrupt conditions are TS_IMR[Cx], TS_IMR[BG], TS_IMR[Cgb], and TS_IMR[A]. The TS_IMR[Cx] mask bits prevent global threshold interrupts and select which sensors participate in the global below threshold interrupt condition. The TS_IMR[BG] directional

bit selects the temperature direction for the global interrupt condition. The TS_IMR[Cgb] mask bit prevents global below threshold interrupts. TS_IMR[A] asserts an attention to the system controller.

Setting TS_IMR[BG] to '1' sets a temperature range for the global interrupt condition to occur when the temperatures of all the participating sensors set in TS_IMR[Cx] are below the global interrupt temperature level. Clearing TS_IMR[BG] to '0' sets a temperature range for the global interrupt condition to occur when the temperature of any of the participating sensors is greater than or equal to the corresponding temperature in TS_GITR. If TS_IMR[A] is set to '1', an attention is asserted when any TS_IMR[Cx] bit and its corresponding status bit (TS_ISR[Gx]) are both set to '1'. Also, an attention is asserted when TS_IMR[Cgb] and TS_ISR[Gb] are both set to '1'.

A thermal-management interrupt is presented to the PPE when any TS_IMR[Mx] bit and its corresponding status bit (TS_ISR[Sx]) are both set to '1'. A thermal-management interrupt is also generated when any TS_IMR[Cx] bit and its corresponding status bit (TS_ISR[Gx]) are both set to '1'. Also, a thermal-management interrupt is presented to the PPE when TS_IMR[Cgb] and TS_ISR[Gb] are both set to '1'.

### 15.2.6 Dynamic Thermal-Management Registers

The dynamic thermal-management registers contain parameters for controlling the execution throttling of the PPE or an SPE. This set of registers contains the Thermal-Management Control Registers (TM_CR1 and TM_CR2), the Thermal-Management Throttle Point Register (TM_TPR), the Thermal-Management Stop Time Registers (TM_STR1 and TM_STR2), the Thermal-Management Throttle Scale Register (TM_TSR), and the Thermal-Management System Interrupt Mask Register (TM_SIMR).

TM_TPR sets the throttle point for the sensors. Two independent throttle points can be set in TM_TPR[ThrottlePPE/ThrottleSPE]—one for the PPE and one for the SPEs. Also contained in this register are temperature points for exiting throttling and stopping the PPE or SPEs. Execution throttling of the PPE or an SPE starts when the temperature is equal to or above the throttle point. Throttling ceases when the temperature drops below the temperature to exit throttling (TM_TPR[EndThrottlePPE/EndThrottleSPE]). If the Temperature reaches the full throttle or stop temperature (TM_TPR[FullThrottlePPE/FullThrottleSPE]), the execution of the PPE or SPE is stopped. TM_CR1 and TM_CR2 are used to control the throttling behavior.

TM_STR1, TM_STR2, and TM_TSR are used to control the frequency and amount of throttling. When the temperature reaches the throttle point, the corresponding PPE or SPE is stopped for the number of clocks specified by the corresponding scale value in the TM_TSR. The PPE or SPE is then allowed to run for the number of clocks specified by the run value in the TM_STR1 and TM_STR2 registers times the corresponding scale value. This sequence, shown below, continues until the temperature falls below the exit throttling (TM_TPR[EndThrottlePPE/EndThrottleSPE]).

TM_SIMR is used to select which interrupts exit throttling of the PPE while the interrupt is pending.

A simplified throttling process for the SPEs is illustrated in the following example. The value TM_Config[MinStopSPE] comes from the configuration ring. All other values in the example are from Thermal-Management MMIO registers.

```
if  (TS_CTSR1[Cur(x)] ≥ TM_TPR[Throttle]) {
       /* where x is the PPE or SPE and 0 ≤ x ≤ 7) */
```

```
    While (TS_CTSR1[Cur(x)] ≥ TM_TPR[EndThrottleSPE]) {
        if (TS_CTSR1[Cur(x)] ≥ TM_TPR[FullThrottleSPE]){
            stop SPE(x)
        } else {
            spcstopunit = (TM_Config[MinStopSPE] * TM_TSR[ScaleSPE])
            stop SPE(x) for (spcstopunit * TM_STR1[StopCore(x)]) cycles
            run SPE(x) for (spcstopunit * (32 - TM_STR1[StopCore(x)])) cycles
        }
    }
}
run SPE(x)
```

The process for throttling the PPE is similar to that for the SPE. The following is the routine for PPE throttling:

```
if  (TS_CTSR2[Cur(8)] ≥ TM_TPR[Throttle]) & (No_Interrupts_Pending) {
       /* No_Interrupts_Pending depends on the setting of TM_SIMR */
    While (TS_CTSR2[Cur(8)] ≥ TM_TPR[EndThrottlePPE]) {
        if (TS_CTSR2[Cur(8)] ≥ TM_TPR[FullThrottlePPE]){
            /* TM_TPR[FullThrottlePPE] should be set at a high enough temperature to
             * avoid stopping PPE
             */
            stop PPE
        } else {
            PPEstopunit = (TM_Config[MinStopPPE] * TM_TSR[ScalePPE])
            stop PPE for (PPEstopunit * TM_STR2[StopCore(8)]) cycles
            run PPE for (PPEstopunit * (32 - TM_STR2[StopCore(8)])) cycles
        }
    }
}
run PPE
```

The preceding code sample shows a simplified throttling process. The actual hardware throttling implementation must also account for updates of the thermal-management and thermal-sensor registers, interrupts, and interface protocols.

Privileged software should set the full throttling temperature `TM_TPR[FullThrottlePPE]`) high enough to avoid the condition where hardware stops the PPE. Stopping the PPE prevents an interrupt from being processed.

#### 15.2.6.1 *Thermal-Management Control Registers*

The Thermal-Management Control Registers (`TM_CR1` and `TM_CR2`) set the throttling mode for the PPE or each SPE independently. The control bits are split between two registers. There are five different modes that can be set for the PPE or each SPE independently:

- Dynamic throttling is disabled (including the Core Stop Safety).
- Normal operation (dynamic throttling and the Core Stop Safety are enabled).
- PPE or SPE is always throttled (Core Stop Safety is enabled).

- Core Stop Safety is disabled (dynamic throttling is enabled and the Core Stop Safety is disabled).
- PPE or SPE is always throttled and Core Stop Safety disabled.

Privileged software should set control bits to normal operation for PPE or SPEs that are running applications or operating systems. If the PPE or an SPE is not running application code, privileged software should set the control bits to disabled. The “PPE or SPE is always throttled” modes are intended for application development. These modes are useful to determine if the application can operate under an extreme throttling condition. Allowing the PPE or an SPE to execute with either the dynamic throttling or Core Stop Safety disabled should only be permitted when privileged software actively manages the thermal events.

#### 15.2.6.2 *Thermal-Management System Interrupt Mask Register*

The Thermal-Management System Interrupt Mask Register (`TM_SIMR`) controls which PPE interrupts cause the thermal-management logic to temporarily stop throttling the PPE. Throttling is temporarily suspended for both threads while the interrupt is pending, regardless of the thread targeted by the interrupt. When the interrupt is no longer pending, throttling can resume if throttle conditions still exist. Throttling of the SPEs is never exited based on a system interrupt condition. The PPE interrupt conditions that can override a throttling condition are:

- External
- Decrementer
- Hypervisor Decrementer
- System Error
- Thermal Management

#### 15.2.6.3 *Thermal-Management Throttle Point Register*

The Thermal-Management Throttle Point Register (`TM_TPR`) contains the encoded temperature points at which execution throttling of the PPE or an SPE begins and ends. This register also contains encoded temperature points at which the PPE or SPE execution is fully throttled.

The values in this register are used to set three temperature points for changing between the three thermal-management states: Normal Run (N), PPE or SPE Throttled (T), and PPE or SPE Stopped (S). Independent temperature points are supported for the PPE and the SPEs.

When the encoded current temperature of a sensor in `TS_CTSR` is equal to or greater than the throttle temperature (`ThrottlePPE/ThrottleSPE`), execution throttling of the corresponding PPE or SPE begins, if enabled. Execution throttling continues until the encoded current temperature of the corresponding sensor is less than the encoded temperature to end throttling (`EndThrottlePPE/EndThrottleSPE`). As a safety measure, if the encoded current temperature is equal to or greater than the full throttle point (FullThrottlePPE/FullThrottleSPE), the corresponding PPE or SPE is stopped.

#### 15.2.6.4 *Thermal-Management Stop Time Registers*

The Thermal-Management Stop Time Registers (TM_STR1 and TM_STR2) control the amount of throttling applied to a specific PPE or SPE in the thermal-management throttled state. The values in this register are expressed as a percentage of time that the PPE or an SPE is stopped versus the time that it is run (CoreStop(x)/32). The actual number of clocks (NClks) that the PPE or an SPE stops and runs is controlled by the Thermal-Management Scale Register (TM_TSR).

#### 15.2.6.5 *Thermal-Management Throttle Scale Register*

The Thermal-Management Throttle Scale Register (TM_TSR) controls the actual number of cycles that the PPE or an SPE stops and runs during the thermal-management throttle state. The values in this register are multiples of a Configuration-Ring setting, TM_Config[MinStopSPE]. The actual number of stop and run cycles is calculated by the following equation:

```
SPE Run and Stop Time:
----------------------

SPE_StopTime = (TM_STR1[StopCore(x)] * TM_Config[MinStopSPE]) * TM_TSR[ScaleSPE]
SPE_RunTime = (32 - TM_STR1[StopCore(x)]) * TM_Config[MinStopSPE]) * TM_TSR[ScaleSPE]

PPE Run and Stop Time:
----------------------

PPE_StopTime = (TM_STR2[StopCore(8)] * TM_Config[MinStopPPE]) * TM_TSR[ScalePPE]
PPE_RunTime = (32 - TM_STR2[StopCore(8)]) * TM_Config[MinStopPPE]) * TM_TSR[ScalePPE]
```

***Note:*** *The run and stop times can be altered by interrupts and privileged software writing various thermal-management registers.*

# 16. Performance Monitoring

The Cell Broadband Engine Architecture (CBEA) processors[1] provide extensive performance-monitoring facilities that assist performance analysis, as well as provide application-optimized and system-optimization features that include:

- Debugging, analyzing, and optimizing processor-architecture features
- Profiling the behavior of the memory hierarchy and the interaction of multiple address spaces, as well as tuning system and application algorithms to optimize scheduling, partitioning, and structuring for tasks and data
- Real-time application-tuning by monitoring bandwidth use and other resource-management behavior

The performance-monitoring facilities can count and log data on over 400 types of internal events. Up to eight types of events can be monitored concurrently. The facilities support the monitoring of classes of instructions selected by an instruction-matching mechanism, the random selection of instructions for detailed monitoring, and count start/stop event pairs that exceed a selected time-out threshold.

The facilities are designed for tuning a broad range of software, including:

- Numerically intensive floating-point applications
- Fixed-point applications
- System software

The facilities give clear visibility to the details of instruction execution, loads and stores, the behavior of caches throughout the CBEA processors, the entire virtual-memory architecture—including effective-to-real address translation (ERAT), translation-lookaside buffers (TLBs), data-prefetch streaming, and the effects of large-page support—and the activity of the element interconnect bus (EIB), memory interface controller (MIC), I/O interface controller (IOC), and several other CBEA processor components that can have large effects on application performance.

The performance monitor facility provides eight 16-bit counters. Counters can be paired (combined) to function as four 32-bit counters for count values to be collected per measurement interval. The interval is timed by a 32-bit programmable interval timer. Performance information can be obtained on signals from either a single island or multiple islands. The performance monitor has the following notable features:

- Counts the number of cycles active or inactive over an interval, or the number of positive or negative transitions over an interval.
- The interval is programmable through an interval timer.
- Performance-monitor signals from the PowerPC Processor Element (PPE) can be masked according to the state of the PPE (for example, hypervisor mode).
- Performance-monitor conditions, such as an interval timeout, can be programmed to cause an interrupt.

The performance monitor counters can be frozen until a selected start count qualifier occurs. The counters can continue to count up until the occurrence of a stop count qualifier. The counters can also be frozen when a counter overflows. An initial count other than ‘0’ can also be applied to the

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

counters to give them a programmable threshold. When the performance monitor counters are frozen, the interval timer is also frozen. See the *Cell Broadband Engine Registers* document for additional details.

## 16.1 How It Works

Access to the registers used in performance monitoring requires hypervisor privilege. The performance-counting facilities share most of the resources used for hardware debug of the CBEA processors. Because performance monitoring tries to maximize the reuse of the debug resources, some debug facilities are reinterpreted for performance-monitoring use. For example, the “trigger bus” and the “event bus” have specific meanings for debug, but for performance monitoring the two buses are regarded as extra buses that are used to route performance-monitoring events.

The basic flow of performance monitoring goes as follows: Each unit in the processor generates various signals (per-event, per-cycle, or both). These signals are selected by writing certain memory-mapped I/O (MMIO) registers. Each signal is connected to one or more bus lines. The bus lines are routed to individual counters. The counters are then connected to a trace array which records the counts periodically.

The debug bus is the transport mechanism for delivering signals from the various CBEA processor units to the performance monitor. This bus is shared by the performance monitor and the trace-logic analyzer. The performance monitor is disabled when the trace-logic analyzer is enabled. Signals to be counted are identified with tags, which indicate whether the signal is useful for counting cycles or edges.

Performance signals are sent to the performance counters by means of the 128-bit debug bus (+ 11 special bits). The 128-bit bus is divided into four words, and most events can be sent to one of two different words of the bus. For example, the signals for a PowerPC processor unit (PPU) thread instruction-unit events can go on either word 2 or word 3, the signals for a memory flow controller (MFC) L2-cache events can go on word 0 or word 1, and the signals for an IOC event can go on word 0 or word 2. Debug-specific signals are reinterpreted for performance monitoring and can use the extra 11 bits normally used for debug.

## 16.2 Events (Signals)

The performance-monitoring facilities can monitor many types of events from all major units in the processor. See *Appendix C Performance Monitor Signals* on page 793 for the full list of verified events (also called signals).

## 16.3 Performance Counters

There are eight 16-bit counters (which can be combined to make four 32-bit counters) and a trace-array collection and storage mechanism. The trace-array mechanism uses both internal and external storage resources. The counters can be set up to count events from either PPU thread, several of the PowerPC processor storage subsystem (PPSS) internal logic blocks, any

of the eight synergistic processor units (SPUs), any of the eight MFCs and their internal logic blocks, the element interconnect bus (EIB), the memory interface controller (MIC), the bus interface controller (BIC), and the I/O interface controller (IOC).

Each counter can monitor one event type at a time, and the events are of three kinds:

- *Cycles*—The cycles in which an event is active or inactive
- *Event Edges*—The edge transitions between two event states. The counters can be frozen until a specified trigger occurs before they begin counting the required event
- *Event Cycles*—The cycles until a specified trigger-event occurs, at which times the counters are frozen

One set of counters can be configured as 16-bit *count qualifiers*, such that the start and stop count of the counting can be specified. This makes it possible to control when the counting of an event starts or stops.

After the counters are set up and begin recording events, the counters can be read interactively by a user program. For example, a user program running in a logical partition can make a system call to its supervisor software, which will then make a hypervisor call that reads the appropriate registers and returns the collected data to the user program.

For details, see the documentation for your tools that are layered on the CBEA processor performance monitoring features.

## 16.4 Trace Array

The trace array contains 1,024 128-bit entries. It is used to store the count values at the end of each interval. When the trace array is filled, a PPE interrupt can be set to occur. The data can be transferred from the trace array to main storage either by simple PPE MMIO reads (128-bits at a time) or by DMA transfers. Count data can also be transferred to external memory to capture thousands of intervals worth of performance data.

Counted events are automatically saved into the trace array according to a specified interval. All performance counters are reset to ‘0’ after each trace-array logging interval. A mechanism is provided for monitoring the operation of the interval timer.

# 17. SPE Channel and Related MMIO Interface

## 17.1 Introduction

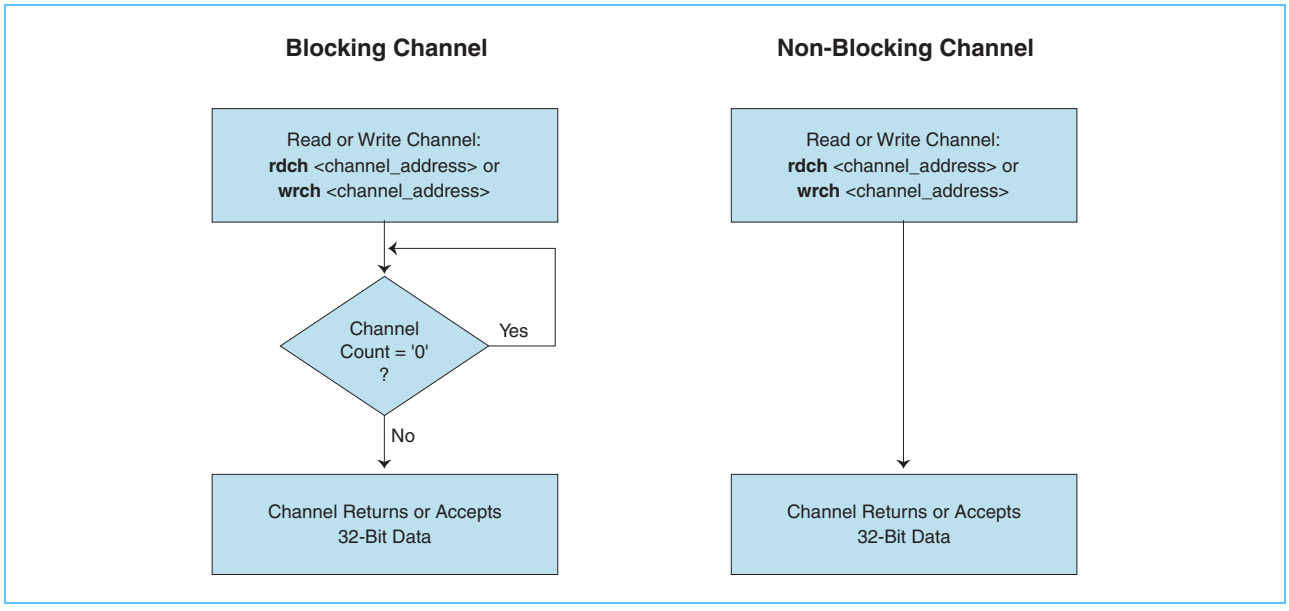
### 17.1.1 An SPE's Use of its Own Channels

A Synergistic Processor Element (SPE) communicates with the PowerPC Processor Element (PPE) and other SPEs and devices through its channels. Channels are unidirectional interfaces for sending and receiving variable-size (up to 32-bit) messages with the PPE and other SPEs, and for sending commands (such as direct memory access [DMA] transfer commands) to the SPE's associated memory flow controller (MFC). Each SPE has its own set of channels.

SPE software accesses channels with three instructions: read channel (**rdch**), write channel (**wrch**), and read channel count (**rchcnt**). These instructions enqueue MFC commands into the SPE's MFC SPU command queue for the purpose of initiating DMA transfers, querying DMA and synergistic processor unit (SPU) status, sending mailbox and signal-notification messages for synchronization of tasks, and accessing auxiliary resources such as the SPE's decrementer.

The channels are architecturally defined as blocking or nonblocking. When SPE software reads or writes a nonblocking channel, the operation executes without delay. However, when SPE software reads or writes a blocking channel, the SPE might stall for an arbitrary length if the associated channel count (which is its remaining capacity) is '0'. In this case, the SPE will remain stalled until the channel count becomes '1' or more, as shown in *Figure 17-1*, or the SPE is interrupted. The stalling mechanism allows an SPE to minimize the power consumed by message-based synchronization of tasks, because fewer logic gates are switching during the stall.

*Figure 17-1. SPE Access to Its Own Channels*



### 17.1.2 Access to Channel Functions by the PPE and other SPEs

Software on the PPE and other devices (including other SPEs) can gain access to some of the facilities an SPU can access through its channel interface—including the passing of mailbox and signal-notification messages—by accessing associated memory-mapped I/O (MMIO) registers in the main-storage effective address (EA) space. Unlike an SPE's accesses to its channels, MMIO accesses are never blocked.

MMIO accesses enqueue MFC commands into the SPE's MFC proxy command queue—a different queue than the MFC SPU command queue used by the SPE, and independent thereof. Although the PPE can use this means of message-passing for synchronization of tasks, such as storage, the PPE more commonly uses memory-barrier instructions (*Section 20* on page 561) available in the *PowerPC Architecture* instruction set.

### 17.1.3 Channel Characteristics

Each SPE channel is either read-only or write-only and is either blocking or nonblocking. Each channel has a corresponding capacity (maximum message entries), and count (remaining capacity). This channel count decrements whenever a channel instruction (**rdch** or **wrch**) is issued to the corresponding channel, and the count increments whenever an operation associated with that channel completes. Thus, a channel count of '0' means "full" for write-only channels and "empty" for read-only channels. A blocking channel will cause the SPE to stall when the SPE reads or writes the channel and the channel count is '0'.

*Table 17-1* summarizes the characteristics of the four channel types.

*Table 17-1. SPE Channel Types and Characteristics*

| Channel Type | Valid Instructions | Blocking | Description |
|---|---|---|---|
| Read | **rdch**<br>**rchcnt** | Nonblocking | **rdch** returns without delay; **rdch** will never stall.<br>**rchcnt** always returns '1'. |
| | | Blocking | **rdch** returns only if the channel count is not '0'; **rdch** stalls if queue empty.<br>**rchcnt** returns number of items waiting in queue. |
| Write | **wrch**<br>**rchcnt** | Nonblocking | **wrch** always puts data without delay; **wrch** will never stall.<br>**rchcnt** always returns '1'. |
| | | Blocking | **wrch** puts data only if the channel count is not '0'; **wrch** stalls if queue is full.<br>**rchcnt** returns number of queue entries available to accept data. |

The *Synergistic Processor Unit Instruction Set Architecture* specifies that a **rdch** or **wrch** instruction to an invalid channel causes the SPE to stop on or after the instruction. However, the Cell/B.E. and PowerXCell 8i implementations differ from the SPU ISA in that **rdch** instructions to reserved channels or valid write channels return zeros and **wrch** instructions to reserved channels or valid read channels have no effect.

Reading or writing an unimplemented (reserved) channel does not cause an illegal channel instruction interrupt. The **rchcnt** instruction responds differently depending on the type of channel it is directed at:

- **rchcnt** returns '0' when directed to an unimplemented (reserved) channel.
- **rchcnt** returns '1' when directed to an implemented, nonblocking channel.
- **rchcnt** returns the capacity (write-only channel) or occupancy (read-only channel) for an implemented, blocking channel.

### 17.1.4 Channel Summary

The key features of the SPE channel operations include:

- All operations on a given channel are unidirectional (they can be only read or write operations for a given channel, not bidirectional).
- Accesses to channel-interface resources through MMIO addresses do not stall.
- Channel operations are done in program order.
- Channel read operations to reserved channels return zeros.
- Channel write operations to reserved channels have no effect.
- Reading of channel counts on reserved channels returns '0'.
- Channel instructions use the 32-bit preferred slot in a 128-bit transfer.

*Table 17-2* on page 450 lists the SPE channels, and their corresponding MMIO registers, that each SPE supports. The EA column under the MMIO Register heading gives the offset in hexadecimal from the base address for the SPE. For further details, see the *Cell Broadband Engine Architecture* document.

The function of channels whose channel mnemonic in *Table 17-2* begins with the prefix `SPU_` is confined to the SPE; channels whose channel mnemonic begins with the prefix `MFC_` access the SPE's MFC for operations outside the SPE, such as DMA transfers, synchronization, mailbox messages, and signal notification.

Most channels are 32 bits wide; only two are 16 bits wide. The *C/C++ Language Extensions for Cell Broadband Engine Architecture* document defines intrinsics for reading and writing channels using 32-bit (scalar `unsigned int`) or 128-bit (vector) operands. When setting a 16-bit channel, only the least-significant 16 bits of a 32-bit scalar operand are used.

*Table 17-2. SPE Channels and Associated MMIO Registers* (Sheet 1 of 3)

| SPE Channel # | Name | Channel Interface[1] (MFC SPU Command Queue) | | | | | MMIO Register Interface[1] (MFC Proxy Command Queue) | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Offset From Base | Mnemonic | Max. Entries | R/W | Width (bits) |
| **SPU Events** | | | | | | | | | | | |
| 0 | SPU Read Event Status | SPU_RdEventStat | 1 | yes | R | 32 | — | — | — | — | — |
| 1 | SPU Write Event Mask | SPU_WrEventMask | 1 | no | W | 32 | — | — | — | — | — |
| 2 | SPU Write Event Acknowledgment | SPU_WrEventAck | 1 | no | W | 32 | — | — | — | — | — |
| **SPU Signal Notification** | | | | | | | | | | | |
| 3 | SPU Signal Notification 1 | SPU_RdSigNotify1 | 1 | yes | R | 32 | x‘1400C’ | SPU_Sig_Notify_1 | 1 | R/W | 32 |
| 4 | SPU Signal Notification 2 | SPU_RdSigNotify2 | 1 | yes | R | 32 | x‘1C00C’ | SPU_Sig_Notify_2 | 1 | R/W | 32 |
| — | SPU Configuration[2,9] | — | — | — | — | — | x‘04078’ | SPU_Cfg | 1 | R/W | 64 |
| 5 | Reserved | — | — | — | — | — | — | — | — | — | — |
| 6 | Reserved | — | — | — | — | — | — | — | — | — | — |
| **SPU Decrementer** | | | | | | | | | | | |
| 7 | SPU Write Decrementer | SPU_WrDec | 1 | no | W | 32 | — | — | — | — | — |
| 8 | SPU Read Decrementer | SPU_RdDec | 1 | no | R | 32 | — | — | — | — | — |
| **MFC Multisource Synchronization** | | | | | | | | | | | |
| 9 | MFC Write Multisource Synchronization Request | MFC_WrMSSyncReq | 1 | yes | W | 32 | x‘00000’ | MFC_MSSync | 1 | R/W | 64 |
| 10 | Reserved | — | — | — | — | — | — | — | — | — | — |
| **SPU and MFC Read Mask** | | | | | | | | | | | |
| 11 | SPU Read Event Mask | SPU_RdEventMask | 1 | no | R | 32 | — | — | — | — | — |
| 12 | MFC Read Tag-Group Query Mask | MFC_RdTagMask | 1 | no | R | 32 | — | — | — | — | — |
| **SPU State Management** | | | | | | | | | | | |
| 13 | SPU Read Machine Status | SPU_RdMachStat | 1 | no | R | 32 | — | — | — | — | — |
| 14 | SPU Write State Save-and-Restore | SPU_WrSRR0 | 1 | no | W | 32 | — | — | — | — | — |
| 15 | SPU Read State Save-and-Restore | SPU_RdSRR0 | 1 | no | R | 32 | — | — | — | — | — |
| **MFC Command Parameters** | | | | | | | | | | | |
| 16 | MFC Local-Storage Address | MFC_LSA | 1 | no | W | 32 | x‘03004’ | MFC_LSA | 1 | W | 32 |
| 17 | MFC Effective Address High | MFC_EAH | 1 | no | W | 32 | x‘03008’ | MFC_EAH | 1 | W | 32 |
| 18 | MFC Effective Address Low or List Address | MFC_EAL | 1 | no | W | 32 | x‘0300C’ | MFC_EAL | 1 | W | 32 |

*Table 17-2. SPE Channels and Associated MMIO Registers* (Sheet 2 of 3)

| SPE Channel # | Name | Channel Interface[1] (MFC SPU Command Queue) | | | | | MMIO Register Interface[1] (MFC Proxy Command Queue) | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Offset From Base | Mnemonic | Max. Entries | R/W | Width (bits) |
| 19 | MFC Transfer Size or List Size | MFC_Size | 1 | no | W | 16[3] | x'03010' | MFC_Size | 1 | W | 16 H |
| 20 | MFC Command Tag Identification | MFC_TagID | 1 | no | W | 16[3] | | MFC_Tag | 1 | W | 16 L |
| 21 | MFC Class ID and Command Opcode | MFC_Cmd | 16 | yes | W | 32 | x'03014' | MFC_ClassID_CMD | 8 | W | 32 |
| — | MFC Command Status[4] | — | — | — | — | — | | MFC_CMDStatus | 1 | R | 32 |
| — | MFC Queue Status[5] | — | — | — | — | — | | MFC_QStatus | 1 | R | 32 |
| **MFC Tag Status** | | | | | | | | | | | |
| 22 | MFC Write Tag-Group Query Mask | MFC_WrTagMask | 1 | no | W | 32 | x'0321C' | Prxy_QueryMask | 1 | R/W | 32 |
| 23 | MFC Write Tag Status Update Request[6] | MFC_WrTagUpdate | 1 | yes | W | 32 | — | — | — | — | — |
| — | Proxy Tag-Group Query-Type[6] | — | — | — | — | — | x'03204' | Prxy_QueryType | 1 | R/W | 32 |
| 24 | MFC Read Tag-Group Status | MFC_RdTagStat | 1 | yes | R | 32 | x'0322C' | Prxy_TagStatus | 1 | R | 32 |
| 25 | MFC Read List Stall-and-Notify Tag Status[7] | MFC_RdListStallStat | 1 | yes | R | 32 | — | — | — | — | — |
| 26 | MFC Write List Stall-and-Notify Tag Acknowledgment[7] | MFC_WrListStallAck | 1 | no | W | 32 | — | — | — | — | — |
| 27 | MFC Read Atomic Command Status[8] | MFC_RdAtomicStat | 1 | yes | R | 32 | — | — | — | — | — |
| **SPU Mailboxes** | | | | | | | | | | | |
| 28 | SPU Write Outbound Mailbox | SPU_WrOutMbox | 1 | yes | W | 32 | x'04004' | SPU_Out_Mbox | 1 | R | 32 |
| 29 | SPU Read Inbound Mailbox | SPU_RdInMbox | 4 | yes | R | 32 | x'0400C' | SPU_In_Mbox | 4 | W | 32 |
| 30 | SPU Write Outbound Interrupt Mailbox[9] | SPU_WrOutIntrMbox | 1 | yes | W | 32 | x'04000' | SPU_Out_Intr_Mbox | 1 | R | 64 |

*Table 17-2. SPE Channels and Associated MMIO Registers* (Sheet 3 of 3)

| SPE Channel # | Name | Channel Interface[1] (MFC SPU Command Queue) | | | | | MMIO Register Interface[1] (MFC Proxy Command Queue) | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Offset From Base | Mnemonic | Max. Entries | R/W | Width (bits) |
| — | SPU Mailbox Status[10] | — | — | — | — | — | x'04014' | SPU_Mbox_Stat | 1 | R | 32 |
| 31:127 | Reserved[11] | — | — | — | — | — | — | — | — | — | — |

1. Because the channel and MMIO command queues are separate queues, a MMIO register has a function that is similar to its corresponding channel, but the MMIO register does not affect the operation of that channel, which is a separate function.
2. On the MMIO register interface, writing the SPU_Cfg register configures the signal-notification to operate in either OR mode or overwrite mode. On the SPU channel interface, no such configuration function is available.
3. When setting a 16-bit channel using a *C/C++ Language Extensions for Cell Broadband Engine Architecture* intrinsic, only the least-significant 16 bits of a 32-bit scalar unsigned int operand are used.
4. On the MMIO register interface, reading the MFC_CMDStatus register returns information about the success or failure of queuing a DMA command. On the SPU channel interface, comparable information can be obtained by writing a command to the MFC_Cmd channel, which will block until a queue entry is available.
5. On the MMIO register interface, reading the MFC_QStatus register returns the number of queue entries available. On the SPU channel interface, comparable information can be obtained by issuing a channel-count (**rchcnt**) instruction.
6. The MFC_WrTagUpdate channel controls the circumstances under which the status in the MFC_RdTagStat channel is updated. The Prxy_QueryType MMIO register performs a comparable function.
7. Only an SPU can issue a DMA list command.
8. Only an SPU can issue MFC atomic commands.
9. Access to this MMIO register is available only to privileged PPE software.
10. On the MMIO register interface, reading the SPU_Mbox_Stat MMIO register returns the number of available mailbox entries. On the SPU channel interface, comparable information can be obtained by issuing a channel-count (**rchcnt**) instruction.
11. The *Synergistic Processor Unit Instruction Set Architecture* supports 128 channels, based on the 7-bit field in the channel command.

### 17.1.5 Channel Instructions

The SPU instruction set defines three instructions to access channels. *Table 17-3* lists the assembler instruction mnemonics and their corresponding C-language intrinsics. The assembler instructions are described in *Table 17-3*; the intrinsics are described in *Section B.2.3* on page 787.

*Table 17-3. SPE Channel Instructions*

| Function | Assembler Instruction | C-Language Intrinsic | Description |
|---|---|---|---|
| Read Channel | **rdch** Rt, chnum | **d = spu_readch**(chnum)<br>**d = spu_readchqw**(chnum) | Data from the channel is loaded into the general-purpose register (GPR) or C-language variable. |
| Write Channel | **wrch** chnum, Ra | **(void) spu_writech**(chnum, a)<br>**(void) spu_writechqw**(chnum, a) | Data from the GPR or C-language variable is stored to the channel. |
| Read Channel Count | **rchcnt** Rt, chnum | **d = spu_readchcnt**(chnum) | The channel count is loaded into the GPR or C-language variable. |

### 17.1.6 Channel Capacity and Blocking

For a nonblocking read channel, a read-channel instruction will execute without delay. Similarly, for a nonblocking write channel, a write-channel instruction will execute without delay. For a blocking channel, however, a read-channel or write-channel instruction can block and cause a delay (SPE stall) of arbitrary length. If a read-channel instruction is directed at a channel connected to an empty input queue (channel count is '0'), the SPE will stall until valid data appears in the queue or the SPE is interrupted. Similarly, if a write-channel instruction is directed at a channel connected to a full output queue (channel count is '0'), the SPE will stall until the queue is able to accept the write data or the SPE is interrupted.

An SPE uses channels to interact with producers and consumers of data that are inherently asynchronous to SPE software. Stalling on a channel read or write has the benefit of minimizing the power consumed by instruction execution, and it also reduces SPE software complexity. This might be useful for cases in which SPE software has no other work to perform until the channel capacity increases.

To avoid stalling on access to a blocking channel, SPE software can read the channel count to determine the available channel capacity; if the read-channel-count instruction returns '0', a read-channel or write-channel instruction would stall.

Reading the channel count might be useful when SPE software has work to perform while waiting for space on the channel to become available (for channel writes) or data on the channel to become available (for channel reads). In this case, SPE software can be structured to intermittently check channel counts (polling) or the SPE interrupt facility can be enabled to implement true asynchronous event handling. Many of the channels have a corresponding event that can be enabled to cause an asynchronous interrupt, and each event interrupt can be enabled independently, as described in *Section 18 SPE Events* on page 471.

## 17.2 SPU Event-Management Channels

There are four SPU event-management channels:

- SPU Read Event Status (`SPU_RdEventStat`)—channel 0
- SPU Write Event Mask (`SPU_WrEventMask`)—channel 1
- SPU Write Event Acknowledgment (`SPU_WrEventAck`)—channel 2
- SPU Read Event Mask (`SPU_RdEventMask`)—channel 11

These are used to control and monitor event reporting. The bit assignments and definitions are the same for all four channels. In addition, there is an internal Pending Event Register that is hidden from SPE software but visible to privileged PPE software.

For details about the SPU event-management channels, see *Section 18.2 Events and Event-Management Channels* on page 472.

## 17.3 SPU Signal-Notification Channels

There are two SPU signal-notification channels, one to read each of the two signal-notification registers. A signal is a short message from outside the SPU (that is, from the PPE, another SPE, or another system device) that can be from one to 32 bits long.

A device outside the SPU sends a signal-notification message to the SPU by writing to the main-storage address of an MMIO register in the SPU's MFC. The signal is latched in the MMIO register, and the SPU executes a read-channel (**rdch**) instruction to get the signal value. An SPU can send a signal-notification message to another SPU with its special send-signal instructions (for example, **sndsig**).

For details about the SPU signalling channels, including programming examples, see *Section 19.7 Signal Notification* on page 551.

## 17.4 SPU Decrementer

Each SPU contains a 32-bit decrementer implemented as a down-counter. All decrementers in the SPUs and PPE count down at the same rate (*Section 13.3.2 SPE Decrementers* on page 390). An SPU decrementer is accessed through two channels, the SPU Write Decrementer Channel (SPU_WrDec) and the SPU Read Decrementer Channel (SPU_RdDec).

A decrementer event becomes pending when the value in the decrementer changes from '0' to negative (the most-significant bit changes from '0' to '1'); see *Section 18 SPE Events* on page 471 for details. *Section 18.10.1* on page 506 gives examples of how a program can use an SPU decrementer event.

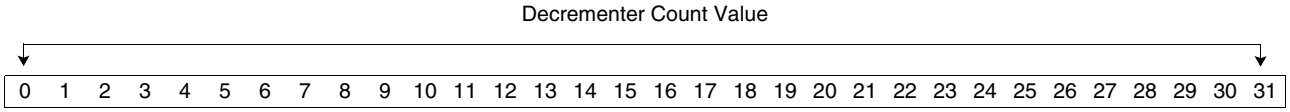
The SPU decrementer can run if the Dh (decrementer halt) bit in the privileged MFC Control Register (MFC_CNTL) is cleared to '0' by PPE software. To start a stopped decrementer, SPE software writes to the SPU Write Decrementer Channel (SPU_WrDec). To stop a running decrementer, SPE software acknowledges the decrementer event when the decrementer event is disabled. The decrementer can also be stopped by setting the MFC_CNTL[Dh] bit to '1'. For additional details and code examples, see *Section 13.3.2 SPE Decrementers* on page 390.

### 17.4.1 SPU Write Decrementer Channel

The value written to the SPU Write Decrementer Channel (SPU_WrDec) sets the starting value of the 32-bit decrementer. When SPE software writes SPU_WrDec with a write-channel (**wrch**) instruction, the decrementer starts running (counting down) if it was not already running. The value loaded into the decrementer determines the time lapse between the **wrch** instruction and the decrementer event becoming pending.

The decrementer event becomes pending when the most-significant bit of the decrementer changes from a '0' to a '1' (the value changes from '0' to negative). If the value loaded into the decrementer causes a change from '0' to '1' in the most-significant bit, the event becomes pending immediately. Setting the decrementer to a value of '0' results in an event after a single decrementer count.

The field assignments and definitions for SPU write decrementer are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | Decrementer Count Value | Decrementer count value. |

### 17.4.2 SPU Read Decrementer Channel

The value read from the SPU Read Decrementer Channel (`SPU_RdDec`) is the current value of the 32-bit decrementer. Reading the decrementer count has no effect on the accuracy of the decrementer (a read does not affect the count-down rate).

The field assignments and definitions for SPU read decrementer are identical to those of the SPU write decrementer (see *Section 17.4.1* on page 454).

## 17.5 MFC Write Multisource Synchronization Request Channel

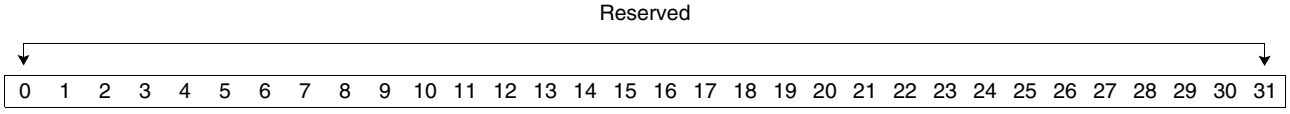
As described in *Section 20.1.5* on page 577, the CBEA processors contain multiple address and communication domains, and the MFC multisource synchronization facility must be used to ensure cumulative ordering across all address domains.

The value written to the MFC Write Multisource Synchronization Request Channel (`MFC_WrMSSyncReq`) causes the MFC to start tracking outstanding write transfers to the MFC. The value written to `MFC_WrMSSyncReq` is ignored, but SPE software should write ‘0’ to ensure compatibility with future architectural revisions.

When software writes to `MFC_WrMSSyncReq`, the channel count is cleared to ‘0’ to indicate an outstanding request is active. When the synchronization is complete (all writes to the MFC that were started before the multisource synchronization request are done), the channel count is set to ‘1’. When the channel count changes from ‘0’ to ‘1’, the multisource synchronization event becomes pending (see *Section 18* on page 471).

When SPE software writes to this channel and the channel count is ‘0’, the SPU stalls until the write transfers being tracked by the currently active multisource synchronization request are complete.

The field assignments and definitions for MFC Write Multisource Synchronization Request are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | Reserved | Reserved. |

See *Section 20.1.5 MFC Multisource Synchronization Facility* on page 577 for a description of multisource synchronization. Graphic flowcharts of the possible software sequences are given in *Figure 20-2* on page 579, *Figure 20-3* on page 581, and *Figure 20-4* on page 582.

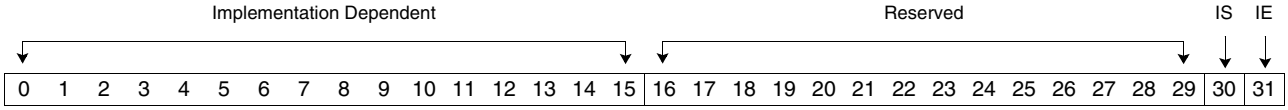
*Section 18.6.2 Procedure for Handling the Multisource Synchronization Event* on page 483 described details for handling a multisource synchronization event.

## 17.6 SPU Read Machine Status Channel

The value read from the SPU Read Machine Status Channel (SPU_RdMachStat) contains two status bits: the isolation status (IS) and the SPU interrupt-enable status (IE). If the IS bit is set to '1', the SPU is operating in the isolated state; if the IS bit is '0', the SPU is not isolated.

If the IE bit is set to '1', any SPU event that is both enabled and pending will cause the SPU to take an interrupt. If the IE bit is '0', the SPU will not take an interrupt regardless of event status. SPE software can set and clear the IE bit by executing an indirect branch with the E feature bit or the D feature bit set, respectively. For more information, see *Section 18* on page 471.

The field assignments and definitions for SPU Read Machine Status are shown in the following table.



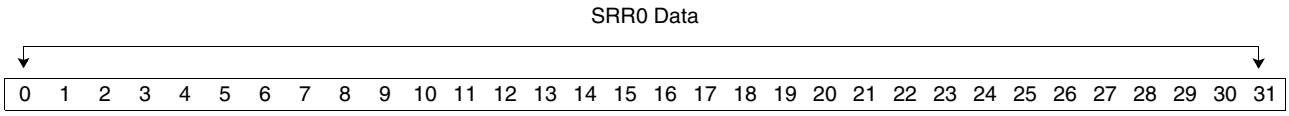

| Bits | Field Name | Description |
|---|---|---|
| 0:15 | Implementation Dependent | |
| 16:29 | Reserved | Cleared to zeros. |
| 30 | IS | Isolation status.<br>0 not isolated.<br>1 isolated. |
| 31 | IE | SPU interrupt enable status.<br>0 interrupts disabled.<br>1 interrupts enabled. |

## 17.7 SPU Write State Save-and-Restore Channel

The value written to the SPU Write State Save-and-Restore Channel (SPU_WrSRR0) sets the return address used by the interrupt-return (**iret**) instruction. The most common reason to write SPU_WrSRR0 is to restore interrupt state when nested interrupts are supported by software. For more information about interrupts and nested interrupts, see *Section 18* on page 471 and *Section 9 PPE Interrupts* on page 239.

SPE software should not write to this channel when interrupts are enabled; doing so can result in the contents of SRR0 being indeterminate. After software writes to this channel and before it executes any instruction that depends on the SRR0 value (such as an **iret** instruction), software must execute the **syncc** instruction (**sync** with the channel feature bit set).

The field assignments and definitions for SPU Write State Save-and-Restore are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | SRR0 Data | State save/restore register 0 data. |

## 17.8 SPU Read State Save-and-Restore Channel

The value read from the SPU Read State Save-and-Restore Channel (SPU_RdSRR0) is the current contents of the SRR0 register, which is the return address saved when the SPU takes an interrupt. The most common reason to read SPU_RdSRR0 is to save interrupt state when nested interrupts are supported by software. For more information about interrupts and nested interrupts, see *Section 18* on page 471 and *Section 9 PPE Interrupts* on page 239.

The field assignments and definitions for SPU Read State Save-and-Restore are identical to those for SPU Write State Save-and-Restore, shown in *Section 17.7* on page 456.

## 17.9 MFC Command Parameter Channels

The channels used by SPE software to enqueue an MFC command to the MFC SPU command queue are described in the following sections. There are six channels involved in setting up an MFC command as shown in *Table 17-4*.

*Table 17-4. MFC Command Parameter and Command Opcode Channels*

| Channel Mnemonic | Channel Number | Description |
|---|---|---|
| MFC_LSA | 16 | MFC Local Storage Address |
| MFC_EAH | 17 | MFC Effective Address High |
| MFC_EAL | 18 | MFC Effective Address Low or List Address |
| MFC_Size | 19 | MFC Transfer Size or List Size |
| MFC_TagID | 20 | MFC Command Tag Identification |
| MFC_Cmd | 21 | MFC Command Opcode and MFC Class ID |

The first five channels listed in *Table 17-4* are referred to as MFC command parameter channels; these channels determine the characteristics of an MFC command. If the SPU MFC command queue has capacity to accept the command, it is enqueued when the command opcode and class ID are written to MFC Command Opcode and MFC Class ID, channel 21. The MFC Command Opcode is the low halfword, and the MFC Class ID logical channel is the high halfword, of a single 32-bit value that is written to channel 21 to complete the command-construction sequence.

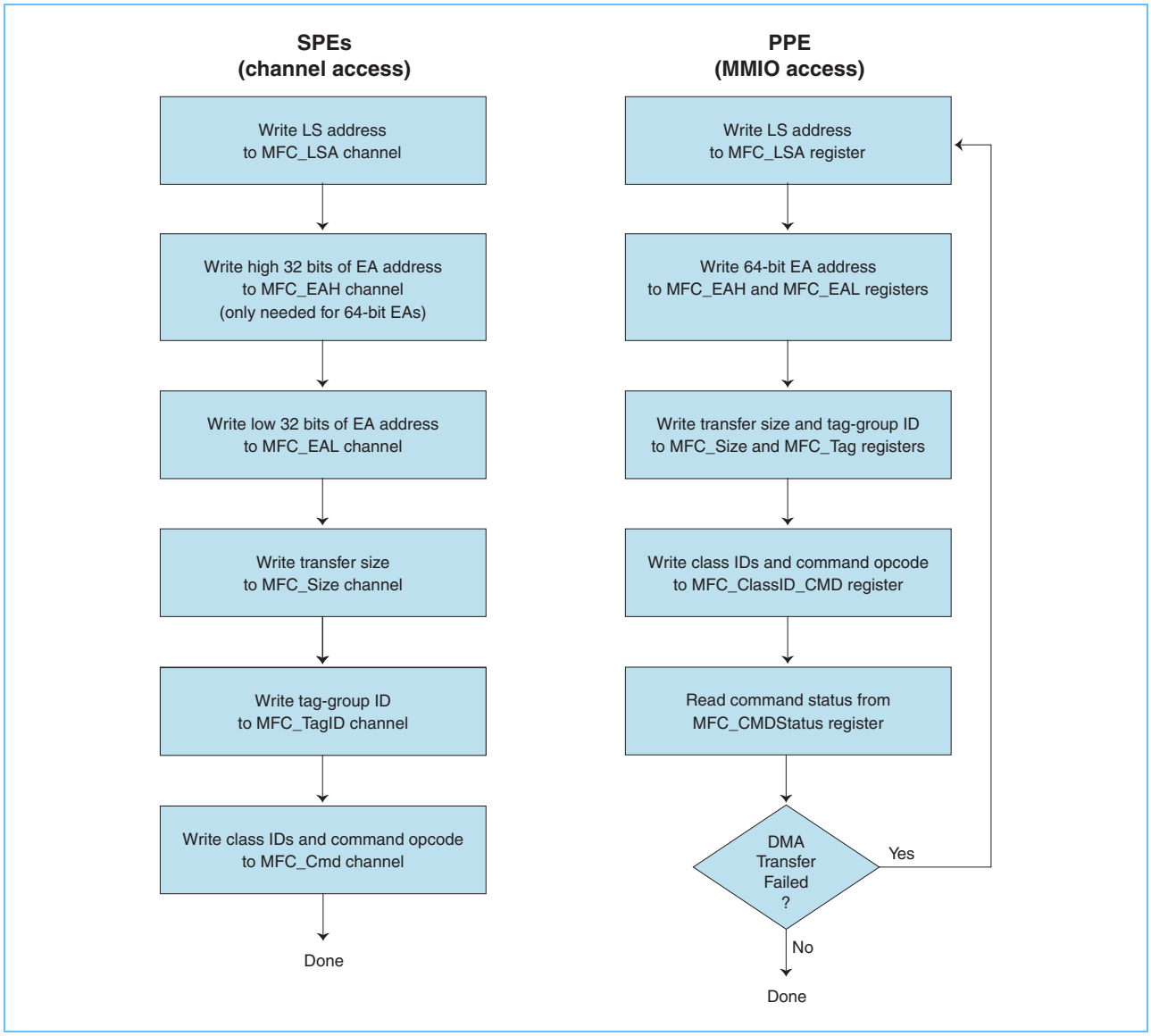
Whenever SPE software writes to MFC Command Opcode, the MFC attempts to enqueue a complete MFC command. Thus, the parameter channels must be written before MFC Command Opcode is written. The only exception is MFC Effective Address High, which need not be written if the default value of '0' is acceptable. SPE software must set all other command parameters before it writes MFC Command Opcode; otherwise, improper operation of the MFC command queue might result.

The parameter channels can be written in any order so long as MFC Command Opcode and MFC Class ID are written last, and if a parameter channel is written more than once, the last value written is used in the queued command.

The MFC command parameter channels (the first five listed in *Table 17-4* on page 457) are nonblocking and do not have channel counts associated with them. Performing a read channel count (**`rchcnt`**) instruction on any of these channels returns '1'. The MFC Command Opcode Channel is blocking and has a maximum channel count of 16. Performing a **`rchcnt`** instruction on this channel returns the remaining capacity (number of empty queue entries) in the MFC command queue.

*Figure 17-2* on page 459 shows the sequences in which MFC commands are issued by SPE software and by PPE software. For additional details about the MFC commands, including DMA-command programming steps and example code, see *Section 19.2 MFC Commands* on page 514.

*Figure 17-2. Sequences for Issuing MFC Commands*




### 17.9.1 MFC Local Storage Address Channel

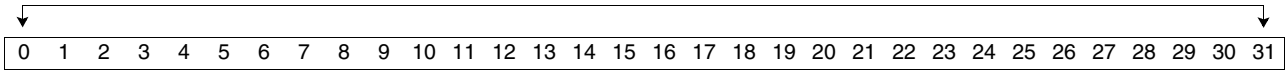
The value written to the MFC Local Storage Address Channel (`MFC_LSA`) sets the SPU local storage (LS) address (LSA) for the MFC command being formed. This address is used as the source (for a **`put`** operation) or destination (for a **`get`** operation) of the MFC transfer as defined in the MFC command. For more information, see *Section 19* on page 513.

If the LSA is unaligned, MFC command queue processing is suspended, and an MFC DMA alignment interrupt is raised to the PPE. To be considered aligned, the four least significant bits of the LS address must match the least-significant four bits of the effective address (MFC Effective Address Low or List Address Channel (see page 460)).

For best performance on transfers of 128 or more bytes, the MFC LSA and the MFC Effective Address Low (EAL) should have bits 25 through 31 cleared to ‘0’, that is, the address should be aligned on a 128-byte boundary.

The field assignments and definitions for MFC LSA are shown in the following table.

MFC Local Storage Address




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | MFC Local Storage Address | MFC local storage address. |

### 17.9.2 MFC Effective Address High Channel

The value written to the MFC Effective Address High Channel (`MFC_EAH`) sets the most-significant 32 bits of the 64-bit effective address for the MFC command being formed. If MFC translation is enabled by the PPE (Relocate bit in the MFC State Register is set to ‘1’), the MFC translates effective addresses into real addresses as described in *Book III* of the *PowerPC Architecture*.

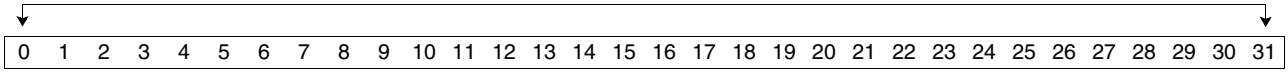
MFC Effective Address High defaults to ‘0’; this is the only MFC command parameter channel that SPE software need not write before it writes to the MFC Command Opcode Channel. If software does not write MFC Effective Address High, the effective address of the queued MFC command will be in the lowest 4 GB of the effective address space.

If the address is invalid due to a segment fault, a mapping fault, or other address violation, MFC command queue processing is suspended, and an interrupt is raised to the PPE.

The MFC checks the validity of the effective address during transfers. Partial transfers can be performed before the MFC encounters an invalid address and raises the interrupt to the PPE.

The field assignments and definitions for MFC Effective Address High are shown in the following table.

High Word of 64-bit Effective Address (Optional)




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | High Word of 64-bit Effective Address | High word of the 64-bit effective address. |

### 17.9.3 MFC Effective Address Low or List Address Channel

The value written to the MFC Effective Address Low or List Address Channel (`MFC_EAL`) sets either the least-significant 32 bits of the 64-bit effective address for the MFC command being formed or the LS address of the list of list elements for the MFC DMA list command being formed. If MFC translation is enabled by the PPE (Relocate bit in the MFC State Register is set to ‘1’), the MFC translates effective addresses into real addresses as described in *Book III* of the *PowerPC Architecture*.

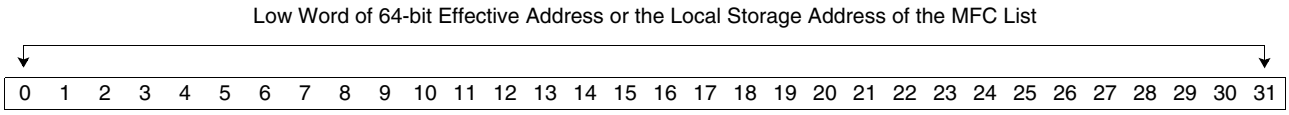
For transfer sizes less than 16 bytes, the MFC Effective Address Low must be naturally aligned (bits 28 through 31 must provide natural alignment based on the transfer size). For transfer sizes of 16 bytes or greater, the MFC Effective Address Low must be aligned to at least a 16-byte boundary (bits 28 through 31 must be '0').

For best performance on transfers of 128 or more bytes, the MFC Effective Address Low (and the MFC Local Storage Address Channel (see page 459)) should have bits 25 through 31 cleared to '0', that is, the address should be aligned on a 128-byte boundary.

For MFC list commands, the MFC List Address must be aligned on an eight-byte boundary (bits 29 through 31 of the List Address must be '0').

If the address is invalid due to a segment fault, a mapping fault, or other address violation, MFC command queue processing is suspended, and an interrupt is raised to the PPE. The MFC checks the validity of the effective address during transfers. Partial transfers can be performed before the MFC encounters an invalid address and raises the interrupt to the PPE.

The field assignments and definitions for MFC Effective Address Low or List Address are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | Low Word of 64-bit Effective Address, or the Local Storage Address of the MFC List | Low word of the 64-bit effective address or the LS address of the MFC list. |

### 17.9.4 MFC Transfer Size or List Size Channel

The value written to the MFC Transfer Size or List Size Channel (`MFC_Size`) sets either the size of the MFC transfer for the command being formed or the size of the list for an MFC DMA list command being formed. In both cases, the value written specifies the number of bytes and cannot be larger than 16 KB.

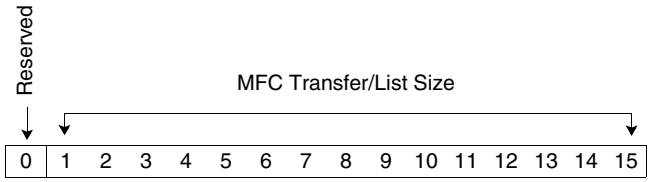
The transfer size can have a value of 1, 2, 4, 8, 16, or a multiple of 16 bytes to a maximum of 16 KB. The source and destination addresses (MFC Local Storage Address Channel (see page 459) and MFC Effective Address Low or List Address Channel (see page 460)) must be naturally aligned to the transfer size for sizes up to 16 bytes. For sizes greater than 16 bytes, alignment on a 16-byte boundary is sufficient.

The List Size of an MFC DMA list command can range from eight bytes up to a maximum of 16 KB. Because each list element has a size of eight bytes, the list can be from one to 2048 elements long. The address of the list (MFC Effective Address Low or List Address Channel (see page 460)) in LS must be aligned on an eight-byte boundary.

If the size is invalid, MFC command queue processing is suspended and an MFC DMA alignment interrupt is raised to the PPE.

Transfers of less than one cache line (128 bytes) should be used sparingly; excessive use of short transfers wastes bus and memory bandwidth. Transfers of less than 16 bytes should only be used to communicate with an I/O device because the setup cost of the transfer is high and bus bandwidth is wasted. When the transfer size is 128 bytes or longer, programmers should align the source and destination addresses on 128-byte boundaries.

The field assignments and definitions for MFC Transfer Size or List Size are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0 | Reserved | Cleared to ‘0’. |
| 1:15 | MFC Transfer/List Size | MFC transfer size or list size. |

### 17.9.5 MFC Command Tag Identification Channel

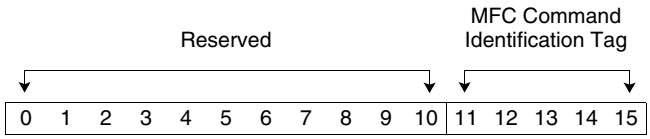
The value written to the MFC Command Tag Identification Channel (`MFC_TagID`) sets the tag value for the command being formed. The identification tag can be any value between 0 and 31. Identification tags have a purely local scope in the hardware.

When two or more MFC commands are enqueued with the same tag value, the commands are in the same tag group. Any number of MFC commands can be tagged with the same identification value.

Tag groups can be formed only within a single MFC command queue. Thus, tags assigned to commands in the MFC SPU command queue are independent of the tags assigned to commands in the MFC proxy command queue.

If the Reserved field (bits 0 through 10) is not cleared to ‘0’, MFC command queue processing is suspended, and an interrupt is raised to the PPE.

The field assignments and definitions for MFC Command Tag Identification are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:10 | Reserved | Cleared to zeros. |
| 11:15 | MFC Command Tag Identification | MFC command tag identification. |

### 17.9.6 MFC Class ID and MFC Command Opcode Channel

The value written to the MFC Class ID and MFC Command Opcode Channel sets the replacement class ID (`RclassID`), the transfer class ID (`TclassID`), and the command opcode for the MFC command that has been formed by previous writes to the MFC command parameter channels (see *Section 17.9* on page 457). In addition, the write to this channel has the side-effect of enqueuing the complete MFC command into the SPU MFC command queue. The write to this channel must be the last channel write during MFC command formation.

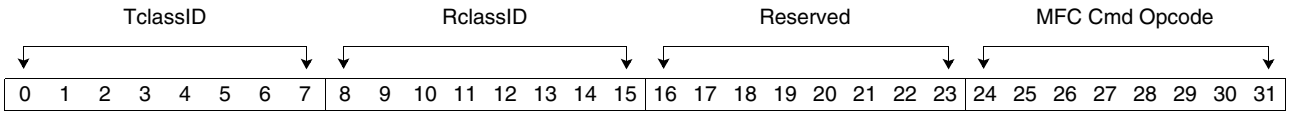
The `RclassID` and `TclassID` are used by the processor and software to improve the overall performance and throughput of the system. The `RclassID` influences L2-cache and translation lookaside buffer (TLB) replacement, as described in *Section 6.3 Replacement Management Tables* on page 154. The `TclassID` influences the allocation of bus bandwidth, as described in *Section 19.2.7 Replacement Class ID and Transfer Class ID* on page 521.

The MFC Command Opcode determines the operation for the MFC command being formed. If the MFC Command Opcode or any of the command parameters is invalid, MFC command-queue processing is suspended and an invalid MFC command interrupt is raised to the PPE. See *Section 19* on page 513 for a complete description of MFC commands.

Software must avoid placing commands in the queue with forward dependencies on newer commands placed in the queue. A string of commands with this type of dependency can create a deadlock, depending on the number of available slots in the MFC queue. Queue depth depends on the CBEA processor version, so software must not be written to require a specific queue depth.

Software can read the channel count of `MFC_Cmd` to determine the number of available queue entries. Because a write to `MFC_Cmd` will stall when the MFC queue is full, software can use the channel count to avoid stalling by waiting for the channel count to become greater than ‘0’.

The field assignments and definitions for MFC Class ID and MFC Command Opcode are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:7 | TclassID | Transfer class identifier. |
| 8:15 | RclassID | Replacement class identifier. |
| 16:23 | Reserved | Should normally be cleared to zeros. Bit 16 set to ‘1’ indicates that the opcode is reserved. |
| 24:31 | MFC Cmd Opcode | MFC command opcode. |

## 17.10 MFC Tag-Group Management Channels

When MFC commands are entered into the command queue, each command is tagged with a 5-bit tag-group identifier, the MFC command tag identifier. The same identifier can be used for multiple MFC commands to create a tag group containing all the commands currently in the queue with the same command tag. Software can use the MFC command tag to check the completion of all queued commands in a tag group. In addition, the MFC command tag is used by

hardware to inform software that an MFC DMA list command has reached an element with its stall-and-notify flag set. Software then uses the command tag to cause the MFC to resume the stalled MFC DMA list command. For more information, see *Section 19* on page 513.

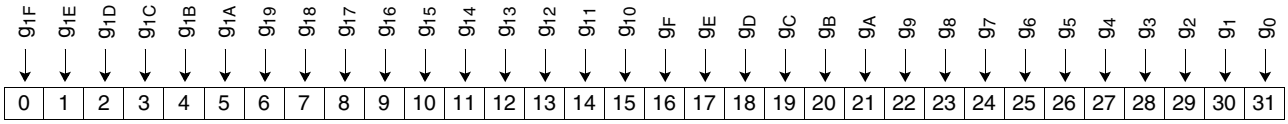
### 17.10.1 MFC Write Tag-Group Query Mask Channel

The value written to the MFC Write Tag-Group Query Mask Channel (`MFC_WrTagMask`) selects the tag groups to be included in a query operation. A query (MFC tag status update request) operation is started by writing to the MFC Write Tag Status Update Request Channel; the status of the query is available in the MFC Read Tag-Group Status Channel.

The value written to this channel is retained by the MFC until changed by a subsequent write. Therefore, the same value can be used for multiple status queries. If this value is changed by software when a query is pending, the results of the query are ambiguous. A pending query should always be cancelled before a modification of this mask. Software cancels a query by writing '0' to the MFC Write Tag Status Update Request Channel.

Software can get the current value in this channel by reading the MFC Read Tag-Group Query Mask Channel.

The field assignments and definitions for MFC Write Tag-Group Query Mask are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | $g_n$ | Tag group "n" select.<br>0 Tag group not in query.<br>1 Tag group is in query. |

### 17.10.2 MFC Read Tag-Group Query Mask Channel

The value read from the MFC Read Tag-Group Query Mask Channel (`MFC_RdTagMask`) is the most recent value written to the MFC Write Tag-Group Query Mask Channel (`MFC_WrTagMask`). Software can read this channel during SPU context save and restore operations and as an alternative to keeping shadow copies of the `MFC_WrTagMask` channel.

The field assignments and definitions for MFC Read Tag-Group Query Mask are identical to those for MFC Write Tag-Group Query Mask, shown in *Section 17.10.1*.

### 17.10.3 MFC Write Tag Status Update Request Channel

The value written to the MFC Write Tag Status Update Request Channel (`MFC_WrTagUpdate`) controls when the MFC tag-group status is updated in MFC Read Tag-Group Status. The MFC Write Tag-Group Query Mask Channel controls which tag groups participate in the update request and therefore which tag groups influence the value read from MFC Read Tag-Group Status.

The value written to MFC Write Tag Status Update determines when MFC Read Tag-Group Status is updated; three update options are possible:

- Status updated immediately after `MFC_WrTagUpdate` is written
- Status updated when any enabled tag group completes after `MFC_WrTagUpdate` is written
- Status updated when all enabled tag groups complete after `MFC_WrTagUpdate` is written

Software must write MFC Write Tag Status Update Request before it reads MFC Read Tag-Group Status.

An MFC Write Tag Status Update request (writing to `MFC_WrTagUpdate`) should be issued after writing a value to MFC Write Tag-Group Query Mask and after enqueuing the MFC commands for the tag groups of interest. If the commands for a tag group are completed before issuing the MFC Write Tag Status Update request thereby satisfying the update status condition, the status is returned without waiting.

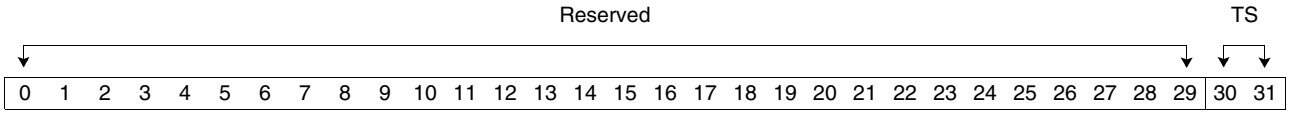
If software reads from MFC Read Tag-Group Status without first writing to MFC Write Tag Status Update Request, this results in a software-induced deadlock.

If software writes to MFC Tag Status Update Request and later needs to cancel the status update request, software should execute the following steps:

1. Request an immediate update status by writing '0' to `MFC_WrTagUpdate`.
2. Read the channel count for `MFC_WrTagUpdate` until a value of '1' is returned.
3. Read from the MFC Read Tag Group Status to discard the unwanted query result and reset the channel count for MFC Read Tag Group Status to '0'.

If software issues two conditional update requests (executes to writes to `MFC_WrTagUpdate`) without an intervening read of the MFC Read Tag Group Status Channel, the value of MFC Read Tag Group Status is unpredictable. To avoid the unpredictable status, software should pair every write to MFC Write Status Update Request with a read from MFC Read Tag-Group Status. The exception is when software cancels a request using the preceding steps.

The field assignments and definitions for MFC Write Tag Status Update Request are shown in the following table.



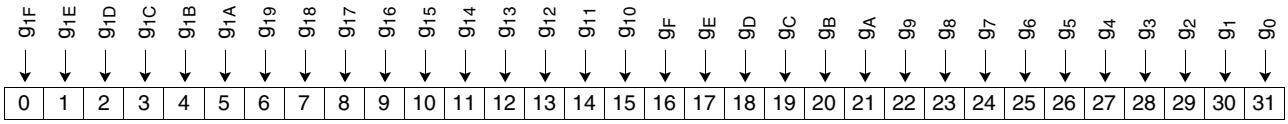

| Bits | Field Name | Description |
|---|---|---|
| 0:29 | Reserved | Reserved. |
| 30:31 | TS | Tag-status update condition:<br>00 MFC_TAG_UPDATE_IMMEDIATE: Update tag status immediately, unconditional.<br>01 MFC_TAG_UPDATE_ANY: Update tag status if or when any enabled tag group completes.<br>10 MFC_TAG_UPDATE_ALL: Update tag status if or when all enabled tag groups complete.<br>11 Reserved. |

### 17.10.4 MFC Read Tag-Group Status Channel

The value read from the MFC Read Tag-Group Status Channel (`MFC_RdTagStat`) reports the status of the tag groups from the last tag-group status update request. Software requests a tag-group status update by writing to MFC Write Tag Status Update Request, and the status reported by MFC Read Tag-Group Status includes only those tag groups enabled by the value that software writes to the MFC Write Tag-Group Query Mask Channel. The bits in the status that correspond to tag groups disabled by the value in MFC Write Tag-Group Query Mask will be '0'.

Software must write a request value to MFC Write Tag Status Update Request before reading from MFC Read Tag-Group Status; failure to do so results in a software-induced deadlock, and only privileged software on the PPE can remove the deadlock condition.

The field assignments and definitions for MFC Read Tag-Group Status are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | $g_n$ | Tag group "n" status.<br>0 Tag group not complete.<br>1 Tag group is complete. |

### 17.10.5 MFC Read List Stall-and-Notify Tag Status Channel

The value read from MFC Read List Stall-and-Notify Tag Status Channel (`MFC_RdListStallStat`) reports which tag groups have an MFC DMA list command in the stall state due to a list element with the stall-and-notify flag set to '1'.

List elements for an MFC list command contain a transfer size, the low 32-bits of an effective address, and a stall-and-notify flag. If the flag is set on a list element, the MFC completes the transfer specified by the list element then stops executing the MFC list command and sets the bit in `MFC_RdListStallStat` corresponding to the tag group of the MFC list command. An MFC list command remains stalled until acknowledged by writing the tag value to the MFC Write List Stall-and-Notify Tag Acknowledgment.

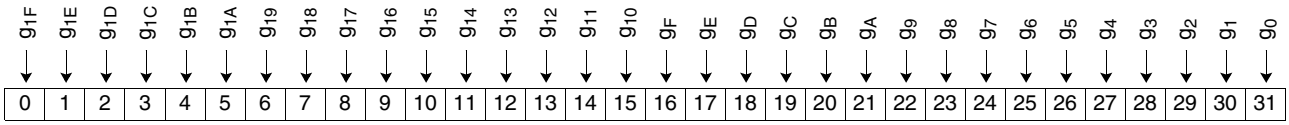
An application program can use the DMA list stall-and-notify capability to monitor the progress of a DMA list command. It can also use stall-and-notify capability when it needs to modify the characteristics (transfer sizes or effective addresses) of list elements that follow the stalled list element. If an application determines that unprocessed list elements should be skipped, it can set the transfer size to '0' in all the elements to be skipped. Software can cancel elements this way because MFC DMA hardware is not allowed to prefetch list elements beyond an element that has its stall-and-notify flag set to '1'.

When software reads MFC Read List Stall-and-Notify Tag Status, all the bits in the status are reset to '0' and the channel count is cleared to '0'. Between reads of the MFC Read List Stall-and-Notify Tag Status, the status accumulates the stall-and-notify status of tag groups as they occur; the channel count for MFC Read List Stall-and-Notify Tag Status never increments above '1'.

If software reads from MFC Read List Stall-and-Notify Tag Status when the channel count is ‘0’, the SPU will stall until the MFC encounters a list element with the stall-and-notify flag set to ‘1’. Thus, software must not read from MFC Read List Stall-and-Notify Tag Status when the channel count is ‘0’ (no bits set in the status) and no outstanding list elements for commands in the MFC queue have the stall-and-notify flag set to ‘1’; if software violates this rule, a software-induced deadlock will occur with the SPU stalled on the read of MFC Read List Stall-and-Notify Tag Status.

Software must also use DMA list commands with the fence or barrier feature bit set if one tag group contains multiple DMA list commands with elements that have a stall-and-notify flag set to ‘1’. Without a fence or barrier to enforce ordering, MFC DMA hardware is free to execute MFC commands (but not list elements) out of order. Thus, without enforced ordering, software cannot be sure which element with the stall-and-notify flag set to ‘1’ has been reached. Software must also implement some mechanism (for example, a linked list) to track DMA list commands and list elements as stalls occur.

The field assignments and definitions for MFC Read List Stall-and-Notify Tag Status are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:31 | $g_n$ | Tag group “n” status.<br>0 Tag group not stalled.<br>1 Tag group stalled. |

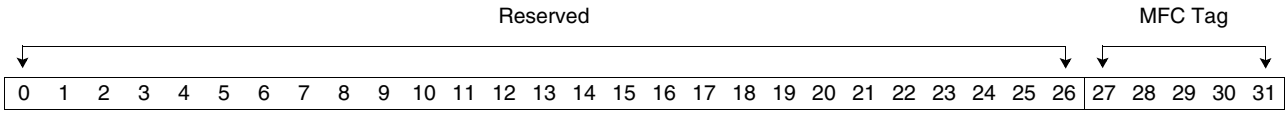
### 17.10.6 MFC Write List Stall-and-Notify Tag Acknowledgment Channel

The value written to the MFC Write List Stall-and-Notify Tag Acknowledgment Channel (`MFC_WrListStallAck`) is the tag group number for a tag group with a stalled DMA list command; writing this tag value to `MFC_WrListStallAck` restarts (acknowledges) the stalled DMA list command in the tag group.

In contrast to most of the other tag group management channels, which have one bit per tag group, MFC Write List Stall-and-Notify Tag Acknowledgment accepts a value corresponding to a tag group number. A tag group number is computed as (31 – tag-group-bit-position), where tag-group-bit-position is the tag group’s bit position in the MFC Read Tag Group Status. MFC Write List Stall-and-Notify Tag Acknowledgment accepts only one tag group number per write, so each write to this channel can restart only one stalled tag group.

If software writes to this channel to acknowledge a tag group that is not stalled due to a stall-and-notify condition, MFC Read List Stall-and-Notify Tag Status will contain an invalid status.

The field assignments and definitions for MFC Write List Stall-and-Notify Tag Acknowledgment are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:26 | Reserved | Reserved. |
| 27:31 | MFC Tag | Tag ID of stalled group that is to be restarted. |

## 17.11 MFC Read Atomic Command Status Channel

The value read from the MFC Read Atomic Command Status Channel (`MFC_RdAtomicStat`) reports the status of the most-recently completed immediate MFC atomic-update command; immediate MFC atomic-update commands are **`getllar`**, **`putllc`**, or **`putlluc`**. This channel does not report any status for the queued **`putqlluc`** MFC atomic-update command. When software reads this channel, its contents are reset to '0'.

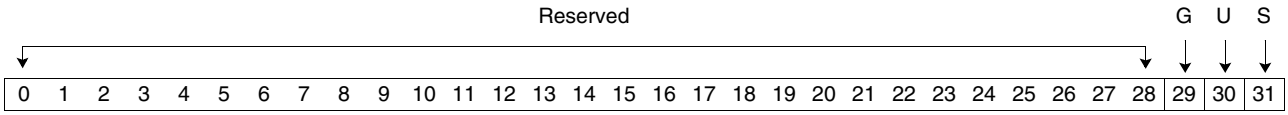
Reading MFC Read Atomic Command Status before issuing one of the immediate MFC atomic-update commands results in a software-induced deadlock.

Software should always read MFC Read Atomic Command Status after every immediate MFC atomic-update command; a read should be paired with each command. If software issues two or more immediate MFC atomic-update commands without intervening reads of the MFC Read Atomic Command Status Channel, the channel will return an incorrect status.

Software can read the channel count for MFC Read Atomic Command Status to determine if the most-recently issued immediate MFC atomic-update command has completed. If the channel count is '0', the command has not completed; if the channel count is '1', the command has completed.

Completion of a subsequent immediate MFC atomic update command overwrites the status of earlier MFC commands.

The field assignments and definitions for MFC Read Atomic Command Status are shown in the following table.




| Bits | Field Name | Description |
|---|---|---|
| 0:28 | Reserved | Reserved. |
| 29 | G | Set if the get lock-line and reserve (**getllar**) command completed. |
| 30 | U | Set if the put lock-line unconditional (**putlluc**) command completed. |
| 31 | S | Put lock-line conditional command (**putllc**).<br>1 Put conditional unsuccessful. The reservation was lost.<br>0 Put conditional successful. |

## 17.12 SPU Mailbox Channels

Mailboxes support the sending and buffering of 32-bit messages between an SPE and other devices, such as the PPE and other SPEs. Each SPE can access three mailbox channels, each of which is connected to a mailbox queue in the SPU's MFC. Two one-entry mailbox channels—the SPU Write Outbound Mailbox and the SPU Write Outbound Interrupt Mailbox—are provided for sending messages from the SPE to the PPE or other device. One four-entry mailbox channel—the SPU Read Inbound Mailbox—is provided for sending messages from the PPE, or other SPEs or devices, to the SPE. Each of the two outbound mailbox channels has a corresponding MMIO register.

An SPE sends a mailbox message by writing the 32-bit message value to either its two outbound mailbox channels. The PPE and other devices can read a message in an outbound mailbox by reading the MMIO register in the SPE's MFC that is associated with the mailbox channel. Likewise, the PPE and other devices send messages to the inbound mailbox by writing the associated MMIO register. For interrupts associated with the SPU Write Outbound Interrupt Mailbox, there is no ordering of the interrupt and previously issued MFC commands.

For details about the mailbox channels, including programming examples, see *Section 19.6 Mailboxes* on page 539.

# 18. SPE Events

This section describes how SPE software can enable and process external events. From the perspective of SPE software, an event is generally an asynchronous change in the state of an resource or unit in a Cell Broadband Engine Architecture (CBEA) processor[1] that is external to that SPE. Each event of potential interest to an SPE has a corresponding bit in the SPE's Event Status Channel. For example, the arrival of a mailbox message from another processor element sets the mailbox bit in the event status and might generate an synergistic processor unit (SPU) interrupt. SPE software can choose to monitor and respond to events synchronously (by polling or stalling) or asynchronously (by enabling the event interrupt).

Each SPE supports an event facility with the capability to mask and unmask events, wait on events, poll for events, and generate interrupts to the SPE when a specific event occurs. If the interrupts are enabled, an occurrence of an unmasked event results in the SPE's interrupt handler being invoked. These interrupts are referred to as "SPU interrupts"; they are unrelated to the PowerPC Processor Element (PPE) interrupts described in *Section 9 PPE Interrupts* on page 239.

## 18.1 Introduction

The main interface between SPE software and the event-tracking hardware is provided by four channels that control event handling and report event status. SPE software deals with events using four basic actions involving these three channels:

- Initialize event handling: write to Event Mask, channel 1.
- Recognize events that have occurred: read Event Status, channel 0.
- Clear events: write Event Acknowledge, channel 2.
- Service events: execute application-specific code.

For a pictorial view of the manner in which events occur, see the Logical Representation of SPU Event Support figure in the *Cell Broadband Engine Architecture* specification. See *Section 17.2 SPU Event-Management Channels* on page 453 for details about the channels.

Typically, an SPE program is interested in only some of the possible SPE events. To enable only the events that are relevant to its operation, SPE software initializes a mask value with event bits set for the relevant events. After software recognizes an event, it writes an acknowledgment value so that the same events can be recognized again.

SPE software can choose to recognize events synchronously or asynchronously. For synchronous event handling, software can either poll for pending events (for example, test a channel count in a loop) or block when no enabled events are pending (for example, stall on a read from an empty channel). For asynchronous event handling, software must enable the event interrupt and provide interrupt handling code. An intermediate approach is to sprinkle Branch Indirect and Set Link if External Data (**bisled**) instructions, either manually or automatically using code-generation tools, throughout application code so that they are executed frequently enough to approximate asynchronous event detection (see *Section 18.5.1* on page 478). This approach has the advantage that interrupts need not be disabled and re-enabled around critical sections of code.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

## 18.2 Events and Event-Management Channels

Each SPE can enable and respond to twelve different external events, listed in *Table 18-2*. There are four event-management channels, with identical bit definitions, that deal with masking, enabling, and reporting events:

- SPU Read Event Status (`SPU_RdEventStat`)—channel 0
- SPU Write Event Mask (`SPU_WrEventMask`)—channel 1
- SPU Write Event Acknowledgment (`SPU_WrEventAck`)—channel 2
- SPU Read Event Mask (`SPU_RdEventMask`)—channel 11

In addition to these four channels, there is an internal Pending Event Register that is hidden from SPE software but visible to privileged PPE software. The internal register and the four event-management channels are described in the following sections.

### 18.2.1 Event Conditions and Bit Definitions for Event-Management Channels

The four event-management channels, plus the internal Pending Event Register, use 32-bit values that defined as shown in *Table 18-2*. This table also explains the condition that leads to each event becoming pending. All of these events that can cause SPU interrupts, as described in *Section 18.3 SPU Interrupt Facility* on page 476.

For more information about conditions that lead to events being raised and the proper protocol for acknowledging and handling events, see *Section 18.6 Event-Specific Handling Guidelines* on page 481.

*Table 18-1. Bit Assignments for Event-Management Channels and Internal Register*

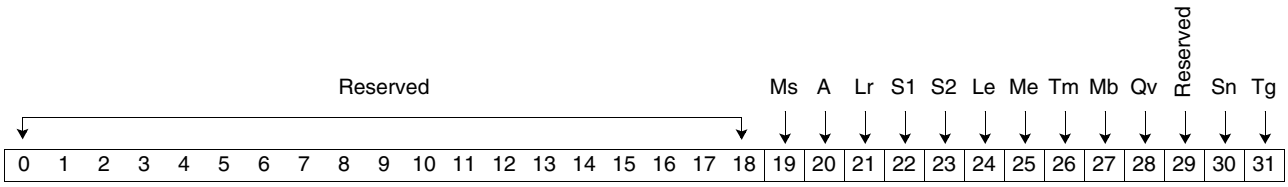



*Table 18-2. Bit Definitions for Event-Management Channels and Internal Register* (Sheet 1 of 2)

| Bit | Field Name | Event Name | Condition for Event To Become Pending | Reason for Condition | Reference |
|---|---|---|---|---|---|
| 0:18 | — | Reserved | | | |
| 19 | Ms | Multisource Sync | Channel count for `MFC_WrMSSyncReq` changes from '0' to '1'. | All of the write transfers directed at the SPE's memory flow controller (MFC) that were started before the multisource synchronization request (the write to `MFC_WrMSSyncReq`) have completed. | *Section 18.6.2* on page 483 |
| 20 | A | Privileged attention | PPE software or other device sets the attention event request bit to '1' in the SPU Privileged Control Register (`SPU_PrivCntl`). | This can be done to control the execution of the SPU. When it is done, a privileged attention event is raised on the SPU. The SPU can reset the condition by acknowledging the event. | *Section 18.6.3* on page 484 |

*Table 18-2. Bit Definitions for Event-Management Channels and Internal Register* (Sheet 2 of 2)

| Bit | Field Name | Event Name | Condition for Event To Become Pending | Reason for Condition | Reference |
|---|---|---|---|---|---|
| 21 | Lr | Lock-line reservation-lost | A lock-line reservation created by a **getllar** MFC command is lost. | The MFC snooped a write to data in the reserved lock-line. | *Section 18.6.4* on page 485 |
| 22 | S1 | Signal-Notification 1 | Channel count for `SPU_RdSigNotify1` changes from '0' to '1'. | Unread signals are pending. | *Section 18.6.5* on page 486 |
| 23 | S2 | Signal-Notification 2 | Channel count for `SPU_RdSigNotify2` changes from '0' to '1'. | | *Section 18.6.6* on page 487 |
| 24 | Le | SPU Write Outbound Mailbox available | Channel count for `SPU_WrOutMbox` changes from '0' to nonzero. | A read from a full SPU Outbound (Interrupt) Mailbox Register has occurred, and the SPU Write Outbound (Interrupt) Mailbox Channel can be written without stalling the SPU. | *Section 18.6.7* on page 488 |
| 25 | Me | SPU Write Outbound Interrupt Mailbox available | Channel count for `SPU_WrOutIntrMbox` changes from '0' to nonzero. | | *Section 18.6.8* on page 489 |
| 26 | Tm | Decrementer | The most-significant bit of the decrementer count changes from '0' to '1' (count changes from '0' to minus-'1'). | The decrementer's elapsed time has expired. If software loads a value that causes the MSb to change from '0' to '1', the event becomes pending immediately. | *Section 18.6.9* on page 489 |
| 27 | Mb | SPU Read Inbound Mailbox Available | Channel count for `SPU_RdInMbox` changes from '0' to nonzero. | A write to an empty SPU Read Inbound Mailbox Register has occurred, and the SPU Read Inbound Mailbox Channel can be read without stalling the SPU. | *Section 18.6.10* on page 491 |
| 28 | Qv | MFC SPU Command Queue Available | Channel count for `MFC_Cmd` changes from '0' nonzero. | An entry in the MFC SPU command queue has become available. | *Section 18.6.11* on page 492 |
| 29 | — | Reserved | | | |
| 30 | Sn | MFC direct memory access (DMA) list command stall-and-notify | Channel count for `MFC_RdListStallStat` changes from '0' to '1'. | A stall-and-notify flag set on a DMA list element has been encountered, the MFC has stopped executing the list, and the MFC has set the bit corresponding to the list's tag group in the `MFC_RdListStallStat` channel. | *Section 18.6.12* on page 492 |
| 31 | Tg | MFC Tag-Group Status Update | Channel count for `MFC_RdTagStat` changes from '0' to '1'. | The tag-group status is currently available by reading the `MFC_RdTagStat` channel. | *Section 18.6.13* on page 494 |

### 18.2.2 Pending Event Register (Internal, SPE-Hidden)

The Pending Event Register records all event occurrences, regardless of whether they are enabled. When an event occurs, the corresponding bit in the Pending Event Register is set to 1. The Pending Event Register has the layout shown in *Table 18-1* on page 472, but it is not directly accessible to SPE software. The contents of the register are, however, visible to privileged PPE software, using the SPU channel access facility described in *Section 12.4.4* on page 362.

Because the Pending Event Register reflects all pending events, regardless of the value of the SPU Write Event Mask, events that are pending but not enabled by the SPU Write Event Mask are nevertheless latched in pending external event. Consequently, if the SPU Write Event Mask is changed to enable a pending event, it will be reported in the SPU Read Event Status Channel

immediately, regardless of how long the event has been pending. For this reason, the initialization code for event handling in SPE software should first clear any pending events (see *Section 18.2.5* on page 475) to prevent the raising of stale or phantom events.

### 18.2.3 SPU Read Event Status

Reading the SPU Read Event Status Channel (`SPU_RdEventStat`) reports events that are both enabled by the SPU Write Event Mask and pending (that is, latched in the Pending Event Register) at the time of the channel read. Any event bit whose corresponding enable is cleared to '0' in the SPU Write Event Mask will read as '0' in the SPU Read Event Status. In other words, when software reads SPU Read Event Status, the value of the SPU Write Event Mask is logically ANDed with the value in the Pending Event Register.

#### 18.2.3.1 *Read Channel and Read Channel Count*

A read-channel-count (**rchcnt**) instruction to the SPU Read Event Status Channel can return either '0' or '1', depending on the channel's state, whereas reading the channel count of any of the other three event-management channels always returns '1'.

A **rchcnt** to SPU Read Event Status returns '0' if no enabled events are pending, and it returns '1' if enabled events have been raised since the last read of the status. A channel count of '1' indicates that at least one event is pending (corresponding bit set in SPU Pending Event) and enabled (corresponding bit set in SPU Write Event Mask). A channel count of '0' indicates no events are both enabled and pending.

Thus, SPE software can poll for pending, enabled events by reading the channel count, or it can execute a **rdch** instruction. A **rdch** when the channel count is '0' causes the SPE to stall until at least one enabled event is pending. This stall behavior allows software to implement a *wait on external event* function. A further benefit is that the SPU uses significantly less power when stalled because most logic gates in the SPU stop switching. To avoid stalling, software can read the channel count before deciding whether to read the channel.

The channel count for SPU Read Event Status is set to '1' for the following conditions:

- An event occurs when the corresponding bit in the SPU Write Event Mask has a value of '1' (the event is enabled when the event occurs).
- At least one event is pending and enabled following a write of the SPU Write Event Mask (an event is already pending when the corresponding bit is set to '1' by a write to the Mask). The previous state of the SPU Write Event Mask has no effect.
- At least one event is both enabled and pending after a write of the SPU Write Acknowledgment Channel (the value written to SPU Write Acknowledgment did not acknowledge all events that were pending and enabled at the time of the write).

#### 18.2.3.2 *Branch Condition*

The channel count of the SPU Read Event Status Channel is used as the condition in the Branch Indirect and Set Link if External Data (**`bisled`**) instruction. If the channel count for SPU Read Event Status is '0', the branch is not taken; if the channel count is '1', the branch is taken. SPE software can execute a **`bisled`** instruction periodically to poll for events, and if **`bisled`** instructions

are executed frequently enough, software can approximate asynchronous event detection. See *Section 18.5.1* on page 478. For more information about the **bisled** instruction, see *Table B-1* on page 772 and the *Synergistic Processor Unit Instruction Set Architecture* document.

#### 18.2.3.3 *SPU Interrupts*

If SPU interrupts are enabled (the IE bit in SPU_RdMachStat is set to '1'), an interrupt is taken when the channel count for SPU Read Event Status changes from '0' to '1', which means that one or more enabled events are pending. The only way to enable or disable interrupts while the SPU is running is to execute an indirect branch with either the D or E feature bit set. The SPU interrupt can be enabled or disabled while the SPU is stopped by setting the I bit in the SPU Next Program Counter Register (SPU_NPC). See the *Synergistic Processor Unit Instruction Set Architecture* document for details.

### 18.2.4 SPU Write Event Mask

An SPE program determines which events it will monitor by setting the appropriate bits in the SPU Write Event Mask Channel (SPU_WrEventMask). The value written to the SPU Write Event Mask determines which pending events will affect the value returned by a read of the SPU Read Event Status Channel. When a bit in the Mask has a value of '1', the corresponding event is said to be enabled. The value written to this channel is retained until a subsequent write changes the value. The current value of the SPU Write Event Mask can be read from the SPU Read Event Mask Channel.

### 18.2.5 SPU Write Event Acknowledgment

Before SPE software services the events reported in SPU Read Event Status, software should write a value to the SPU Write Event Acknowledge Channel (SPU_WrEventAck) to acknowledge (clear) the events that will be processed. The bits set to '1' in the value that is written to SPU Write Event Acknowledge cause the corresponding bits in the hidden Pending Event Register to be cleared to '0'. Thus, after software writes acknowledgment bits, events of the same type can be reported again. Clearing all pending and enabled events has no effect on the channel count for SPU Read Event Status.

All events are recorded in the hidden Pending Event Register, regardless of the SPU Write Event Mask value. A pending, enabled event will continue to be reported in SPU Read Event Status until it is acknowledged—that is, cleared by writing an appropriate value to the SPU Write Event Acknowledgment Channel; each bit position that corresponds to a pending event must be set to '1'.

The bit for a pending event can be cleared in the hidden SPU Pending Event Register with a write to the SPU Event Acknowledgment Channel even if the event is disabled. A disabled event that becomes pending and is subsequently acknowledged (cleared) is not reflected in the value read from the SPU Read Event Status Channel. An event that has been acknowledged is resampled to allow future occurrences of the event to be recorded in the hidden SPU Pending Event Register.

Enabling a currently pending but disabled event by writing an appropriate mask value to the SPU Write Event Mask Channel results in an update of the SPU Read Event Status Channel count (and an SPU interrupt, if enabled).

Acknowledging an event before it occurs results in a software-induced deadlock. Software should be careful in clearing unreported events.

### 18.2.6 SPU Read Event Mask

Software uses the SPU Read Event Mask Channel (SPU_RdEventMask) to read the state of the SPU Write Event Status Mask. The SPU Read Event Mask always returns the data last written to the mask. Instead of keeping software shadow copies of the Event Status Mask, SPE software can read the state of the mask directly. SPE context save and restore operations can also use this channel.

## 18.3 SPU Interrupt Facility

The SPE can monitor event status through the SPU Read Event Status Channel (SPU_RdEventStat). Software can process the event status in the SPU_RdEventStat channel by polling, using the **bisled** instruction or by enabling the SPU interrupt facility. When enabled, the interrupt facility causes the SPE to execute the interrupt handler located at local storage (LS) address x‘0’ when the SPU_RdEventStat channel has nonzero count.

The SPU interrupt can be enabled in one of two ways:

- Setting the E bit in the indirect-branch instruction (see the *Synergistic Processor Unit Instruction Set Architecture* document)
- Setting bit 31 of the SPU Next Program Counter Register (SPU_NPC) before running the SPE

An SPU interrupt is taken if the interrupt is enabled and the count in the SPU_RdEventStat channel has transitioned from ‘0’ to ‘1’. This transition is edge-triggered, based on the ORing of events in the SPU_RdEventStat channel, masked by bits in the SPU Write Event Mask Channel (SPU_WrEventMask), and the acknowledgment bits in the SPU Write Event Acknowledgment Channel (SPU_WrEventAck). For a logical representation of the SPE event support and the descriptions of the SPU_RdEventStat channel, the SPU_WrEventMask channel, and the SPU_WrEventAck channel, see the *Cell Broadband Engine Architecture* document. See the *Synergistic Processor Unit Instruction Set Architecture* document for information about using indirect branches for entering and returning from the interrupt handler.

An interrupt sequence might go through the following steps:

1. Enable interrupts as described previously.
2. The program runs until conditions cause the channel count for the SPU_RdEventStat channel to transition from '0' to '1'; then, the interrupt is taken.
3. Interrupts are acknowledged through internal logic as follows:
   a. Internal logic disables interrupts.
   b. Any blocked channel access is preempted.
   c. The program counter goes to LS address 0.
   d. The return address is stored in SRR0.
4. Read the SPU_RdEventStat channel to see the status of events (the channel count for the SPU_RdEventStat channel bit transitions to '0').

5. Acknowledge all events to be processed (write to the SPU_WrEventAck channel).
6. Process events (for example, mailbox events).
7. Execute an interrupt return (**iret**) instruction, which branches to the return address held in SRR0. The SPU interrupt can be re-enabled using the E bit of the **iret** instruction.

Nested interrupts are supported through reading the SPU Read State Save-and-Restore Channel (SPU_RdSRR0) and writing the SPU Write State Save-and-Restore Channel (SPU_WrSRR0).

The value written to the SPU_WrSRR0 channel is not immediately available to the **iret** instruction. This write must be synchronized by execution of a channel **sync** instruction (with the C bit set in the instruction) before executing an **iret** instruction.

For details about handling SPU interrupts, see *Section 18.5.2 Asynchronous Event Handling Using Interrupts* on page 479.

## 18.4 Interrupt Address Save-and-Restore Channels

Two channels are used for saving and restoring the interrupt return address: SPU Read State Save-and-Restore (SPU_RdSRR0) and SPU Write State Save-and-Restore (SPU_WrSRR0).

### 18.4.1 SPU Read State Save-and-Restore

When an SPE takes an interrupt, the current program counter is stored in the SPU Read State Save-and-Restore Channel (SPU_RdSRR0) and execution continues are address 0. The saved program counter can be acquired by reading SPU_RdSRR0.

SPU interrupts are precise; that is, all instructions in the program flow before the interrupt will have completed execution, and no subsequent instructions will have begun.

### 18.4.2 SPU Write State Save-and-Restore

SPE software can write to the SPU Write State Save-and-Restore Channel (SPU_WrSRR0) to set the return address for an interrupt-return (**iret**) instruction. This channel should only be written when SPU interrupts are disabled, and after it is written, the **syncc** instruction should be executed before the interrupt-return instruction. The **syncc** instruction (**sync** with the channel-feature bit set) makes sure channel writes complete before any subsequent instructions can execute.

### 18.4.3 Nested Interrupts Using SPU Write State Save-and-Restore

The SPE architecture does not have the ability to assign priorities to event interrupts, but nested interrupts can help software approximate prioritized interrupts by allowing a low-priority event interrupt handler to be interrupted by a higher-priority event.

To implement nested interrupts, SPU interrupt handling code can save the SPU Read State Save-and-Restore value before re-enabling interrupts. After handling the initial interrupt, SPU interrupt code must:

1. Disable interrupts.

2. Write to SPU Write State Save-and-Restore to restore the initial interrupt return address.
3. Execute a **syncc** instruction (**sync** with the channel feature-bit set) to force the channel write to complete.
4. Execute an **irete** instruction to return from the initial interrupt and re-enable interrupts.

## 18.5 Event-Handling Protocols

SPE software can be written to handle events either synchronously or asynchronously. Synchronous event handling can be implemented by polling or stalling; asynchronous event handling is implemented with interrupts.

### 18.5.1 Synchronous Event Handling Using Polling or Stalling

SPE software can use polling to recognize an enabled and pending event either by reading the channel count of SPU Read Event Status, channel 0, or by executing the **bisled** instruction.

When no enabled events are pending, the channel count of SPU Read Event Status is '0'; when one or more enabled events are pending, the count is '1'. Thus, SPE event-polling code can implement a loop such as the one shown in the following code fragment:

```
do {
    /* Perform work between checks of the
     * channel count, as appropriate.
     */
} while (!spu_readchcnt(SPU_RdEventStat));
```

A programmer can use the **bisled** instruction or the **spu_bisled** intrinsic to test the SPU Read Event Status channel count. The instruction or intrinsic jumps to the indirect branch target when the channel count is '1' but falls through when the count is '0'.

As shown in *Table 17-2* on page 450, the SPU Read Event Status Channel is a blocking channel. If SPE software executes a **rdch** instruction on the channel, the **rdch** instruction will stall when there are no enabled events pending. This provides a simple way to implement a general "wait on external event" capability as illustrated in the following simple code fragment:

```
unsigned int event;

/* Block until event is available. */
event = spu_readch(SPU_RdEventStat);
```

SPE software that is concerned with only a single type of event, such as a mailbox or signal arrival, can read the associated channel or channel count to achieve stalling or polling behavior if the associated channel can block. Thus, if SPE software is concerned only with events on the SPU Signal Notification 1 Channel, it can poll the channel count for `SPU_RdSigNotify1` with the **rchcnt** instruction, or software can stall by reading `SPU_RdSigNotify1`, a blocking channel, with a **rdch** instruction. If SPE software is concerned only with incoming mailbox events, it can stall by

blocking on a **rdch** from SPU_RdInMbox. Polling or stalling on a specific channel instead of on the SPU Read Event Status Channel can save a small amount of overhead (setting up the event mask and writing the event acknowledgment).

### 18.5.2 Asynchronous Event Handling Using Interrupts

When SPU interrupts are enabled and at least one event is pending in the SPU Read Event Status Channel, SPE hardware takes the following steps to recognize the event interrupt:

- Interrupts are disabled.
- The address of the next instruction that would have executed is saved in Save-Restore Register 0.
- SPE hardware begins executing instructions from address 0 in the SPE's LS.

The only way SPE software can enable and disable interrupts under program control is by using the D and E feature bits in indirect-branch instructions; SPE software cannot write directly to SPE machine status. *Table 18-3* shows the instructions that can be qualified with the D and E feature bits to enable or disable interrupts. See the *Synergistic Processor Unit Instruction Set Architecture* document for details.

*Table 18-3. Indirect-Branch Instructions*

| Mnemonic | Description |
|---|---|
| **bi** | Branch Indirect |
| **bisl** | Branch Indirect and Set Link |
| **bisled** | Branch Indirect and Set Link If External Data |
| **biz** | Branch Indirect If Zero |
| **binz** | Branch Indirect If Not Zero |
| **bihz** | Branch Indirect If Zero Halfword |
| **bihnz** | Branch Indirect If Not Zero Halfword |
| **iret** | Interrupt Return |

These instructions leave the interrupt-enable status unchanged unless either the D or E feature bit is specified (an instruction with both feature bits causes undefined behavior).

When one of these branches is taken with an interrupt feature bit set, the interrupt-enable status change takes effect before the target instruction is executed.

SPE hardware supports one hardwired interrupt vector for all interrupts; this vector is to address 0 in LS. This single interrupt handler at address 0 must determine which events have occurred and process them appropriately.

If an SPE application needs nested interrupt handling, the handler at address 0 must save the interrupt return address from Read Save-Restore Register 0, prepare for a possible nested interrupt (for example, save other volatile state), and enter an interruptible event handler with interrupts enabled (that is, execute an indirect branch with the E feature bit to the interruptible event handler).

When (possibly nested) event processing is complete, an interruptible interrupt handler can return to the instruction stream it interrupted in one of two ways:

- By writing the return address to SPU Write State Save-and-Restore and executing a **syncc** or **iret** instruction pair (**syncc** is needed to force the channel write before **iret** executes).
- By loading the return address into an SPE register and executing a **bie** instruction (or other suitable indirect branch).

In either case, the **iret** or **bi** instruction should have the `E` feature bit set to re-enable event interrupts.

### 18.5.3 Protecting Critical Sections from Interruption

When asynchronous interrupts are enabled, an interrupt can happen at any time while application code is running. Because applications might contain critical sections of code that must not be interrupted, programs must disable interrupts while a critical section is executing.

Critical sections arise in typical SPE application programs. For example, an application that enqueues DMA commands must do so with a sequence of instructions that forms the complete DMA command. When both application and interrupt handler code enqueue DMA commands, the sequence of channel writes that forms the DMA command must not be interrupted; the channel writes for a single DMA command must all complete before the first channel write for the next DMA command happens.

In high-level-language code, the SPE intrinsics `spu_idisable()` and `spu_ienable()` can be used to surround a critical section. In assembler code, indirect branches must be used to disable and enable interrupts. The following code fragments show one way to implement disable and enable operations in assembler; in fact, these are the code sequences in the intrinsics `spu_idisable()` and `spu_ienable()`.

```
idisable:
   ila    $Rt, next_inst  # form address of instruction after bie
   bid    $Rt             # branch to next_inst and disable interrupts
next_inst:

[code inside critical section]

ienable:
   ila    $Rt, next_inst  # form address of instruction after bie
   bie    $Rt             # branch to next_inst and enable interrupts
next_inst:
```

***Note:*** *A branch to the next sequential instruction will not be mispredicted and need not be hinted.*

## 18.6 Event-Specific Handling Guidelines

Most of the semantics of event handling will be application specific and therefore beyond the scope of this document, but all applications should follow general guidelines for event-handling protocol and satisfy the minimum requirements for each possible event. Examples of event handlers with application-specific code are given later in this section (see *Section 18.10.1* on page 506).

For each event, the conditions that lead to it being raised are described in the following sections. One of the necessary conditions for any event to be raised—that the event be enabled in the SPU Write Event Mask—is assumed but not mentioned in the descriptions.

### 18.6.1 Protocol with Multiple Events Enabled

When multiple events are enabled (when the SPU Write Event Mask has more than one bit set to '1'), SPE software should follow a simple procedure to prevent the occurrence of spurious (phantom) events. Before SPE software performs any event-specific actions, it must disable the events that will be processed and then acknowledge them. This allows the MFC hardware to register new occurrences of the events as soon as possible. After all event-specific processing is completed, SPE software restores the event mask. This allows the new occurrences of the processed events to be recognized in the SPU Read Event Status Channel.

The steps of the common protocol, when multiple events are enabled, follow:

1. Save SPU Read Event Status (`SPU_RdEventStat`) in `status`.
2. Save SPU Read Event Mask (`SPU_RdEventMask`) in `mask`.
3. Disable events: write (`mask & ~status`) to SPU Write Event Mask (`SPU_WrEventMask`).
4. Acknowledge events: write `status` to SPU Write Event Acknowledge (`SPU_WrEventAck`).
5. Process events (see *Table 18-4* on page 482 and the following individual event sections).
6. Restore the event mask: write `mask` to SPU Write Event Mask (`SPU_WrEventMask`).

A common implementation of this protocol is to perform steps 1 through 4 and 6 in a first-level event handler routine and to encapsulate the code for handling each specific event in second-level event handling routines. With this organization, the first-level event handler sets up the environment and then calls each second-level event handler routines that corresponds to a bit set to '1' in the event status. This two-level organization is used in the interrupt-driven event handling examples described in *Section 18.7* on page 495, *Section 18.8* on page 501, *Section 18.9* on page 504, and *Section 18.10.1* on page 506.

The following sections describe the recommended protocols for handling specific events. *Table 18-4* on page 482 lists the events along with protocol summaries.

*Table 18-4. SPE Event Handling Protocol Summaries* (Sheet 1 of 2)

| Field Name | Bit | Event Name | Handling Protocol | Reference |
|---|---|---|---|---|
| Ms | 19 | Multisource Synchronization | • Release resource waiting on synchronization. | *Section 18.6.2* on page 483 |
| A | 20 | Privileged Attention | • Perform high-priority response such as stop-and-signal, mailbox communication, or update in-memory status. | *Section 18.6.3* on page 484 |
| Lr | 21 | Lock-Line Reservation Lost | • If lock still being used, re-issue lock and act on new data. | *Section 18.6.4* on page 485 |
| S1 | 22 | SPU Signal-Notification 1 Available | • Read channel count of `SPU_RdSigNotify1`; if '0', handler done.<br>• Read `SPU_RdSigNotify1` and process signal data. | *Section 18.6.5* on page 486 |
| S2 | 23 | SPU Signal-Notification 2 Available | • Read channel count of `SPU_RdSigNotify2`; if '0', handler done.<br>• Read `SPU_RdSigNotify2` and process signal data. | *Section 18.6.6* on page 487 |
| Le | 24 | SPU Outbound Mailbox Available | • Read channel count of `SPU_WrOutMbox`; if '0', handler done.<br>• Write `SPU_WrOutMbox` with mail box data. | *Section 18.6.7* on page 488 |
| Me | 25 | SPU Outbound Interrupt Mailbox Available | • Read channel count of `SPU_WrOutIntrMbox`; if '0', handler done.<br>• Write `SPU_WrOutIntrMbox` with mail box data. | *Section 18.6.8* on page 489 |
| Tm | 26 | SPU Decrementer | • Read `SPU_RdDec`; if negative, this is time delay from event to response (additional time-base ticks since decrementer count expired).<br>• Write `SPU_WrDec` with new timer interval count (optionally reduced by the time delay discovered previously). | *Section 18.6.9* on page 489 |
| Mb | 27 | SPU Read Inbound Mailbox Available | • Read channel count of `SPU_RdInMbox`; if '0', handler done.<br>• Read `SPU_RdInMbox` and process mail box data.<br>• Write `SPU_WrEventAck` with `Mb` bit set (acknowledge each mail box entry).<br>• Repeat. | *Section 18.6.10* on page 491 |

*Table 18-4. SPE Event Handling Protocol Summaries* (Sheet 2 of 2)

| Field Name | Bit | Event Name | Handling Protocol | Reference |
|---|---|---|---|---|
| Qv | 28 | MFC SPU Command Queue Available | • Read channel count of `MFC_Cmd`; if '0', handler done.<br>• Enqueue DMA command with sequence of channel writes.<br>• If no more commands to enqueue, exit.<br>• Write `SPU_WrEventAck` with `Qv` bit set.<br>• Repeat. | *Section 18.6.11* on page 492 |
| Sn | 30 | MFC DMA List Command Stall-and-Notify | • Read `MFC_RdListStallStat` to get stalled tag groups.<br>• Choose a stalled tag group.<br>• Process the stalled tag group.<br>• Write `MFC_WrListStallAck` for the stalled tag group.<br>• Repeat previous three steps until all stalled tag groups processed. | *Section 18.6.12* on page 492 |
| Tg | 31 | MFC Tag-Group Status Update | • Read channel count of `MFC_RdTagStat`; if '0', handler done.<br>• Read `MFC_RdTagStat` to get new tag-group status. | *Section 18.6.13* on page 494 |

### 18.6.2 Procedure for Handling the Multisource Synchronization Event

A multisource synchronization event is the expected response after SPE software writes to MFC Write Multisource Synchronization Request, channel 9 (`MFC_WrMSSyncReq`). SPE software can write any value to start the request, but writing a value of '0' ensures compatibility with future architectural revisions. The occurrence of a multisource synchronization event indicates that all the write transfers directed at the SPE's MFC that were started before the multisource synchronization request (the write to `MFC_WrMSSyncReq`) have completed.

SPE software must use the multisource synchronization facility when cumulative ordering is needed across address domains. Cumulative ordering is the ordering of storage accesses performed by multiple sources (that is, two or more processor elements) with respect to another processor element.

See *Section 20.1.5 MFC Multisource Synchronization Facility* on page 577 for a description of multisource synchronization. Graphic flowcharts of the possible software sequences are given in *Figure 20-2* on page 579, *Figure 20-3* on page 581, and *Figure 20-4* on page 582.

SPE software uses multisource synchronization when it needs data that depends on the completion of a chain of storage operations that other processor elements perform. Thus, when the multisource synchronization event occurs, the action taken by SPE software is application-dependent. Two typical meanings of the event are that (1) an input buffer is completely filled and ready for SPE processing, and (2) that an output buffer is no longer in use and can be reallocated for future use.

To ensure the multisource synchronization event will be raised in response to a specific request, SPE software should perform an initial acknowledgment of the event (write '1' to the `Ms` bit of the `SPU_WrEventAck` channel) before enabling it (write '1' the `Ms` bit of the `SPU_WrEventMask` channel). This initial clearing of a possible stale (phantom) pending multisource synchronization event

need only happen once during SPE application startup, because adherence to the protocols outlined here will ensure that each event occurrence is properly matched with its acknowledgment.

The procedure for handling the multisource synchronization event is as follows:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in a "mask".[2]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Ms]` set to '0'.[2]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Ms]` set to '1'.[2]
4. Perform the application-specific function in response to the completion of a pending multisource synchronization operation.

   This typically indicates that the data in a particular buffer has been completely updated, or that a buffer area is no longer in use.
5. Exit the multisource synchronization event handler.
6. Restore the "mask" by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Ms]` set to '1'.[2]
7. Exit the general event handler.[2]

### 18.6.3 Procedure for Handling the Privileged Attention Event

The privileged attention event can be raised when a '1' is written by the PPE to the SPU Privileged Control Register's Attention Event Required bit. If the event is enabled when this bit is set to '1', the event is raised.

The intended use of the privileged attention event is to inform SPE software that privileged software running on the PPE requires high-priority action. Examples of such actions are an SPE context switch and reporting SPE status to PPE software. Thus, the handler for this event will contain or call the functions that implement the high-priority action.

Because the privileged attention event communicates only one bit of information—the event itself—an application-specific protocol must be implemented if additional information must be passed between the PPE and the SPE. For example, information about the type of attention requested can be passed in a memory location or as a value written to a signal or mailbox channel.

2. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

The procedure for handling the privileged attention event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in “mask”.[3]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[A]` set to ‘0’.[3]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[A]` set to ‘1’.
4. Perform the application-specific function in response to a privileged attention event.

   This can be used to signal that a yield of the SPU is being requested or some other action. An application or operating system-specific response to the privileged attention event should be issued, such as stop and signal, SPU Inbound mailbox write, SPU Outbound Interrupt mailbox write, or an update of a status in system or I/O memory space.
5. Exit the privileged attention event handler.
6. Restore the “mask” by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[A]` set to ‘1’.[3]
7. Exit the general event handler.[3]

### 18.6.4 Procedure for Handling the Lock-Line Reservation Lost Event

This event is raised when a get lock-line and reserve (**getllar**) command is issued, and the reservation is reset due to the modification of data in the same lock line by an outside entity. A lock-line reservation is reset because the snoop hardware in the SPE’s atomic unit (atomic-update cache) detects a data modification by the outside entity. This event will not be raised due to a reservation reset by any local SPE action.

When SPE software executes a **getllar** command, the atomic unit caches the 128-byte line and marks the status of the line as reserved. Typically, the reservation will be reset when SPE software executes a matching **putllc** or **putlluc** to complete an atomic operation (see *Section 20 Shared-Storage Synchronization* on page 561 for a complete description of atomic operations using atomic lock-line commands). When the reservation is reset by a matching **putllc, putlluc,** or **putqlluc** operation, the lock-line reservation lost event will not be raised. If another processor element or device modifies the data in the reserved lock-line, the reservation will be lost and the event will be raised if it is enabled.

The handler for this event should decide if the lock-line reservation should be renewed, and if so, issue a duplicate **getllar** command. There is no need to read the atomic command status to confirm that the reservation was lost; the fact that the event was received means that the reservation was lost.

3. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

The procedure for handling the lock-line reservation lost event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in “mask”.[4]
2. Mask the event by issuing a write channel instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Lr]` set to ‘0’.[4]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Lr]` set to ‘1’.[4]
4. Perform the application-specific function in response to system modification of data in the lock line area.

   This is typically started by checking a software structure in memory to determine if a lock line is still being monitored. If it is still being “waited on,” then the next step typically consists of issuing a **getllar** command to the same lock line area that was modified to obtain the new data and then acting on that data.
5. Exit the lock line reservation lost event handler.
6. Restore the “mask” by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Lr]` set to ‘1’.[4]
7. Exit the general event handler.[4]

***Note:*** *Avoid looping on **getllar** commands. Specifically, in the Cell/B.E. and PowerXCell 8i implementations, avoid looping on more than four different lock lines. This helps to prevent the starvation of writes or low-priority reads, which can result from back-to-back high-priority reads taking precedence in the arbitration scheme. For more information, see the Cell Broadband Engine Architecture document.*

### 18.6.5 Procedure for Handling the Signal-Notification 1 Available Event

The signal-notification 1 available event is raised when another processor or device has written to an empty SPU Signal Notification 1 Register. The event occurrence also means that the channel count has changed from ‘0’ to ‘1’; thus, the SPU can read the value of the signal without stalling.

To account for the possibility of a phantom event, the procedure for handling a signal-notification 1 available event should start by reading the channel count for the SPU Signal-Notification 1 Channel (`SPU_RdSigNotify1`). If the count is ‘0’, the event was a phantom event and no signal data is waiting to be read. If the count is not ‘0’, the handler should read `SPU_RdSigNotify1` to get the signal data and reset the signal data to ‘0’ in preparation for the next signal.

4. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

The procedure for handling the SPU signal-notification 1 available event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in “mask”.[5]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[S1]` set to ‘0’.[5]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[S1]` set to ‘1’.[5]
4. Obtain the channel count by issuing a read channel count (**rchcnt**) instruction to the SPU Signal Notification 1 Channel.
5. If the channel count is ‘0’, skip to step 7.
6. Read the signal data by issuing a read channel (**rdch**) instruction to the SPU Signal Notification 1 Channel (`SPU_RdSigNotify1`).
7. Exit the SPU Signal Notification 1 handler.
8. Restore the “mask” by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[S1]` set to ‘1’.[5]
9. Exit the general event handler.[5]

### 18.6.6 Procedure for Handling the Signal-Notification 2 Available Event

The signal-notification 2 available event is raised when another processor or device has written to an empty SPU Signal Notification 2 Register. The event occurrence also means that the channel count has changed from ‘0’ to ‘1’; thus, the SPU can read the value of the signal without stalling.

To account for the possibility of a phantom event, the procedure for handling a signal-notification 2 available event should start by reading the channel count for the SPU Signal-Notification 2 Channel (`SPU_RdSigNotify2`). If the count is ‘0’, the event was a phantom event and no signal data is waiting to be read. If the count is not ‘0’, the handler should read `SPU_RdSigNotify2`, which will get the signal data and reset the signal data to ‘0’ in preparation for the next signal.

The procedure for handling the SPU signal-notification 2 available event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in “mask”.[5]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[S2]` set to ‘0’.[5]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[S2]` set to a ‘1’.[5]
4. Obtain the channel count by issuing a read channel count (**rchcnt**) instruction to the SPU Signal Notification 2 Channel.
5. If the channel count is ‘0’, skip to step 7.

5. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

6. Read the signal data by issuing a read (**rdch**) channel instruction to the SPU Signal Notification 2 Channel.
7. Exit the SPU Signal Notification 2 handler.
8. Restore the "mask" by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[S2]` set to '1'.[5]
9. Exit the general event handler.[5]

### 18.6.7 Procedure for Handling the SPU Write Outbound Mailbox Available Event

The SPU write outbound mailbox available event is raised when a PPE or other device reads from the SPU Write Outbound Mailbox and the mailbox has available data. When the mailbox data is read by the PPE or device, the channel count for the SPU Write Outbound Mailbox Channel (`SPU_WrOutMbox`) changes from '0' to '1', which means the SPU can write `SPU_WrOutMbox` without stalling.

To account for the possibility of a phantom event, the procedure for handling a SPU write outbound mailbox available event should start by reading the channel count for `SPU_WrOutMbox`. If the count is '0', the event was a phantom event and the mailbox is still full. If the count is not '0', the handler should write `SPU_WrOutMbox` with the next mailbox data to send to the PPE.

The procedure for handling the SPU write outbound mailbox available event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in "mask".[6]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Le]` set to '0'.[6]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Le]` set to '1'.[6]
4. Obtain the channel count by issuing a read channel count (**rchcnt**) instruction to the SPU Write Outbound Mailbox Channel.
5. If the channel count is '0', skip to step 7.
6. Write a new mailbox data entry by issuing a write channel (**wrch**) instruction to the SPU Write Outbound Mailbox Channel.
7. Exit the SPU Outbound Mailbox handler.
8. Restore the "mask" by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Le]` set to '1'.[6]
9. Exit the general event handler.[6]

6. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

### 18.6.8 Procedure for Handling the SPU Write Outbound Interrupt Mailbox Available Event

The SPU write outbound interrupt mailbox available event is raised when a PPE or other device reads from the SPU's PPE interrupt mailbox and the mailbox has available data. When the mailbox data is read by the PPE or device, the channel count for the SPU Write Outbound Interrupt Mailbox Channel (`SPU_WrOutIntrMbox`) changes from '0' to '1', which means the SPU can write `SPU_WrOutIntrMbox` without stalling.

To account for the possibility of a phantom event, the procedure for handling a SPU write outbound interrupt mailbox available event should start by reading the channel count for `SPU_WrOutIntrMbox`. If the count is '0', the event was a phantom event and the mailbox is still full. If the count is not '0', the handler should write `SPU_WrOutIntrMbox` with the next mailbox data to send to the PPE.

The procedure for handling the SPU write outbound interrupt mailbox available event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in the "mask".[7]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Me]` set to '0'.[7]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Me]` set to a '1'.[7]
4. Obtain the channel count by issuing a read channel count (**rchcnt**) instruction to the SPU Write Outbound Interrupt Mailbox Channel.
5. If channel count is '0', skip to step 7.
6. Write a new mailbox data entry by issuing a write channel (**wrch**) instruction to the SPU Write Outbound Interrupt Mailbox Channel.
7. Exit the SPU Outbound Interrupt Mailbox Available handler.
8. Restore the "mask" by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Me]` set to '1'.[7]
9. Exit the general event handler.[7]

### 18.6.9 Procedure for Handling the SPU Decrementer Event

The SPU decrementer event is raised when the decrementer value changes from '0' to negative; that is, the SPU decrementer event is raised when the most-significant bit of the decrementer changes from '0' to '1'.

The decrementer maintains a signed, twos-complement value in a 32-bit down counter. The decrementer counts down at the rate of the PPE time base; all decrementers in the SPUs and the PPE count down at the same rate. The value of the decrementer can be read by SPE software through SPU Read Decrementer (`SPU_RdDec`) and set through SPU Write Decrementer (`SPU_WrDec`).

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7. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

The SPU decrementer can run if the Dh (decrementer halt) bit in the MFC_CNTL register is cleared to '0' by PPE software (see *Section 13 Time Base and Decrementers* on page 381). If the Dh bit is set, SPE software cannot start the decrementer.

After an SPU decrementer event becomes pending, it is possible that a significant amount of time might elapse before the event handler runs. Depending on the needs of the application program, it might be desirable to account for the delay in setting the next time interval to be measured by the decrementer.

To account for the delay, the event handler can read SPU_RdDec to get the current value of the decrementer. If the decrementer has been running during the delay, the value will be the twos-complement of the number of time-base ticks that have elapsed between the time the event was raised and the read of SPU_RdDec. To account for the time lapse in the next decrementer interval, the handler can simply add the required decrementer count to the (negative) value read from SPU_RdDec and write the result to SPU_WrDec. Software will thus set the decrementer to a new count that is adjusted by the number of time-base ticks since the event was raised.

If the protocol given in *Section 18.6.1* on page 481 is followed, however, the decrementer will stop when the acknowledgment step is executed (acknowledging the decrementer event when it is disabled stops the decrementer). In this case, the twos-complement value read from SPU_RdDec as described previously might not account for all the time lapse between the event being raised and the execution of the decrementer event-handler (other handlers might run before the decrementer handler while the decrementer is stopped).

Thus, it might be necessary to either structure interrupt handling to run the decrementer event-handler first or not acknowledge the decrementer event immediately. If the handler follows a protocol that does not cause the decrementer to stop (the event is disabled but not acknowledged immediately or the event is enabled when it is acknowledged), then the twos-complement value read from SPU_RdDec will accurately represent the time delay between the event being raised and the handler running.

To start a stopped decrementer, SPE software writes SPU_WrDec. Thus, when the decrementer event-handler writes a value to SPU_WrDec for another time interval, the decrementer begins counting down immediately if it was stopped.

The procedure for handling the SPU decrementer event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in the "mask".[8]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with SPU_WrEventMask[Tm] set to '0'.[8]
3. Acknowledge the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel (SPU_WrEventAck[Tm] set to '1').
4. Read the decrementer value by issuing a read channel (**rdch**) instruction to the SPU Read Decrementer Channel. If this value is negative, it can be used to determine how much additional time has elapsed from the required interval.

8. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

5. If a new timer event is required, write (**wrch**) a new decrementer value to the SPU Write Decrementer Channel.
6. Exit the SPU decrementer event handler.
7. Restore the "mask" by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Tm]` set to '1'.[8]
8. Exit the general event handler.[8]

### 18.6.10 Procedure for Handling the SPU Read Inbound Mailbox Available Event

The SPU read inbound mailbox event is raised when the PPE or other device writes data to the SPU Read Inbound Mailbox and the mailbox is empty; the data is accessible by the SPE through the SPU Read Inbound Mailbox (`SPU_RdInMBox`). When the `SPU_RdInMBox` channel count changes from '0' to nonzero, the event is raised.

When the PPE or other device writes data to the inbound mailbox, the channel count for the `SPU_RdInMBox` is incremented. The queue depth of the `SPU_RdInMBox` is four, and the maximum channel count is four. It is possible for a PPE or other device to overrun the SPU mailbox which causes the most-recently written value to be lost.

When the SPU mailbox available event handler executes, the SPU mailbox queue can have from one to four valid entries. Because the event is only raised when the `SPU_RdInMBox` channel count changes from '0' to nonzero, the handler must read and process all valid entries to prevent any entries from being ignored for an arbitrarily long period of time

The procedure for handling the SPU read inbound mailbox available event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in "mask".[9]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Mb]` set to '0'.[9]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Mb]` set to '1'.[9]
4. Obtain a channel count by issuing a read channel count (**rchcnt**) instruction to the SPU Read Inbound Mailbox Channel.
5. If the channel count is '0', skip to step 8.
6. Read next mailbox data entry by issuing a read channel (**rdch**) instruction to the SPU Read Inbound Mailbox Channel (`SPU_RdInMbox`).
7. Return to step 3.
8. Exit the SPU inbound mailbox handler.

9. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

9. Restore the “mask” by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Mb]` set to ‘1’.[9]
10. Exit the general event handler.[9]

### 18.6.11 Procedure for Handling the MFC SPU Command Queue Available Event

The MFC SPU command queue available event is raised when the channel count for MFC Class ID and Command Opcode (`MFC_Cmd`) changes from ‘0’ (full) to nonzero (not-full). The occurrence of the event means that the SPU can write `MFC_Cmd` to enqueue a command at least once without stalling.

When the MFC SPU command queue available event handler executes, the MFC SPU command queue might have more than one empty entry. Because the event is only raised when the `MFC_Cmd` channel count changes from ‘0’ to nonzero, the handler can improve efficiency by enqueuing commands until the command queue is full. Thus, software can implement a software queue of MFC commands that is shared by the handler and the main application code.

The procedure for handling the MFC SPU command queue available event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in the “mask”.[10]
2. Mask the event by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Qv]` set to ‘0’.[10]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Qv]` set to ‘1’.
4. Obtain the channel count by issuing a read channel count (**rchcnt**) instruction to the MFC Command Opcode Channel (`MFC_Cmd`).
5. If the channel count is ‘0’, skip to step 8.
6. Enqueue a DMA command to the MFC command queue.
7. If more commands are left to queue, return to step 3.
8. Exit the MFC SPU command queue handler.
9. Restore the “mask” by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Qv]` set to ‘1’.[10]
10. Exit the general event handler.[10]

### 18.6.12 Procedure for Handling the DMA List Command Stall-and-Notify Event

The DMA list command stall-and-notify event is raised when the DMA hardware completes an element of a DMA list command list and the element has the stall-and-notify flag set to ‘1’ (the stall-and-notify flag is bit 0 (the MSb) of the transfer-size word of the element); when the channel count for MFC Read List Stall-and-Notify Tag Status (`MFC_RdListStallStat`) changes from ‘0’ to ‘1’, the event is raised.

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10. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

At this point, the DMA hardware suspends processing of the DMA list and raises the event to give SPE software an opportunity to modify the characteristics (transfer size, starting effective address, stall-and-notify flag) of the remaining elements in the list. See *Section 19 DMA Transfers and Interprocessor Communication* on page 513 for an example of how the stall-and-notify event can be used.

The DMA hardware in the MFC is designed to execute DMA operations with maximum performance. To this end, the DMA hardware can execute the DMA commands from its queues out of order. When the stall-and-notify event is raised, SPE software can query `MFC_RdListStallStat` to find out which tag groups have a stalled DMA list command, but the channel interface provides no information about which DMA list command in the tag group has stalled and no indication of which element in the DMA list command has stalled.

When a DMA list contains multiple elements with the stall-and-notify flag set to '1' or when a tag group has multiple DMA list commands queued that have elements with the stall-and-notify flag set, application software must track stall-and-notify events explicitly. Software can maintain a stall counter for each tag group with elements that can stall, but in addition, the DMA list commands within each such tag group must be explicitly ordered with tag-specific fences or barriers or with the command barrier. Without enforced ordering, DMA hardware can execute DMA commands out of order; if hardware executes commands out of order, software is unable to identify which DMA list command within the tag group is causing a given stall-and-notify event.

When the DMA list command stall-and-notify event handler executes, `MFC_RdListStallStat` might have more than bit set to '1', indicating more than one group with a stalled DMA command. If the handler does not acknowledge and resume all stalled tag groups in response to the event by writing `MFC_WrListStallAck` for each group, the un-acknowledged stalled groups will not cause a subsequent stall-and-notify event and they will remain in the stalled state. Thus, the handler should process and acknowledge all stalled groups reported in `MFC_RdListStallStat`.

Software can adjust the characteristics of a stalled DMA list command in a tag group by changing list-element addresses and transfer sizes for the list elements beyond the element that caused the stall. If application software determines that some list elements should be skipped, software can simply set the list-element transfer sizes to '0'. However, the number of list elements in a queued DMA list command cannot be changed.

The procedure for handling the DMA list command stall-and-notify event is:

1. Perform a read channel (**rdch**) instruction to the SPU Read Event Mask Channel and save the data in the "mask".[11]
2. Mask the event by issuing a write channel instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Sn]` set to '0'.[11]
3. Acknowledge the event by performing a write channel (**wrch**) instruction to the SPU Write Event Acknowledgment Channel with `SPU_WrEventAck[Sn]` set to '1'.[11]
4. Perform a read channel (**rdch**) instruction to the MFC Read List Stall-and-Notify Tag Status Channel `MFC_RdListStallStat[gn]`.

11. When multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

5. Use this information to determine which tag group or tag groups have a DMA List Element in the Stall and Notify state.
6. Perform the application-specific action with respect to each tag group having a stalled DMA list element.

   ***Note:*** *If a DMA list contains multiple list elements having the Stall and Notify flag set, or if a Tag Group has multiple DMA list commands queued with elements having the Stall and Notify flag set, it is essential for the application software to initialize to 0 a tag-group specific stall counter before the DMA list commands are queued for the tag group. In addition, if multiple DMA list commands are queued for a tag group with Stall and Notify elements, ordering must be enforced with tag-specific fences, barriers, or the command barrier. Each time a Stall and Notify status is indicated for a tag group, the corresponding counter should be incremented. Application software can then use this counter to determine at what point in the list the stall has occurred. Application software uses stall and notify to update list element addresses and transfer sizes that follow the list element that has stalled due to dynamically changing conditions. List elements after the stalled list element can be skipped by setting their transfer sizes to 0. However the number of list elements in a queued DMA list command cannot be changed.*
7. Acknowledge and resume each stalled DMA list command by issuing a write channel (**wrch**) instruction to the MFC Write List Stall-and-Notify Tag Acknowledgment Channel (`MFC_WrListStallAck[MFC Tag]`) where the supplied MFC Tag is the encoded Tag ID of the tag group to be resumed.
8. Exit the DMA List Stall and Notify handler.

   ***Note:*** *If application software does not acknowledge all stalled tag groups indicated in the MFC_RdListStallStat[gn] channel, a second stall and notify event does not occur for the unacknowledged tag group.*
9. Restore the “mask” by issuing a write channel (**wrch**) instruction to the SPU Write Event Mask Channel with `SPU_WrEventMask[Sn]` set to ‘1’.[12]
10. Exit the general event handler.[11]

### 18.6.13 Procedure for Handling the Tag-Group Status Update Event

The tag-group status update event is raised when a tag group or groups have completed and resulted in an update of the MFC Read Tag-Group Status (`MFC_RdTagStat`); when the channel count for `MFC_RdTagStat` changes from ‘0’ to ‘1’, the event is raised, and software can read `MFC_RdTagStat` without stalling.

The specific meaning of this event is determined by the values that software writes to the MFC Write Tag Group Query Mask (`MFC_WrTagMask`) and MFC Write Tag Status Update Request (`MFC_WrTagUpdate`). Software writes a mask to `MFC_WrTagMask` to track completion status for the groups with corresponding bits set to ‘1’ in the mask. Then, software writes a request type to `MFC_WrTagUpdate`. If the request type is “any,” then when any of the groups named in the mask complete, `MFC_RdTagStat` is updated and the event is raised. If the request type is “all,” then when

12. If multiple events are enabled, a common handler should be used to save the current event mask, mask all events that are to be handled in one channel write, and acknowledge all events that are to be handled in a single channel write. Then each event specific handler should be invoked to handle the event. The common handler should then restore the current event mask from the saved value and exit. This technique minimizes the generation of spurious events.

all the groups named in the mask complete, `MFC_RdTagStat` is updated and the event is raised. If the request type is “immediate,” then `MFC_RdTagStat` is updated immediately and the event is raised.

The steps in the event-handling protocol follow:

1. Read channel count (**rchcnt**) for `MFC_RdTagStat`.
2. If channel count is ‘0’ then handler done.
3. Read `MFC_RdTagStat`.
4. Perform application-specific processing for each completed tag group.

## 18.7 Developing a Basic Interrupt Handler

Because SPE hardware provides only basic interrupt support, SPE software is responsible for defining most aspects of interrupt handling. Most programmers will want to use high-level languages and tools as much as possible, so the calling conventions and stack management defined by the high-level language development tools should guide the development of the interrupt handling code.

Proper handling of external events on the SPE involves a certain degree of complexity. From a software engineering perspective, this complexity is best managed when a set of system utility functions is provided. To this end, the following items are considered necessary to promote the use of external events (and asynchronous interrupts) on the SPE:

- A set of utilities that safely mask, unmask, and acknowledge a set of external events.
- A first-level interrupt handler (FLIH) that saves and restores resources on entry and exit, and that supports second-level interrupt handlers (SLIHs) that conform to the application binary interface (ABI) of the SPE program. A set of utilities for registering SLIHs should also be included.
- A set of utilities to query, enable, disable, and save interrupt state, thus supporting nested interrupts and SPE critical sections.[13]

Much of this can be absorbed into the default SPE entry function (`crt0.o`), or efficiently implemented as macros or inline assembly functions, with little or no cost to the general SPE program’s LS footprint.

The focus of the description in this section is on developing an interrupt handler that is compatible with function-calling conventions and stack management, as defined in the *SPU Application Binary Interface Specification*. Compatibility with the ABI allows the interrupt handlers to be at least partly written in a high-level language.

### 18.7.1 Basic Interrupt Protocol Features and Design

The example interrupt protocol shown here will consist of a small FLIH coded in assembler language and a small routine for registering SLIHs coded entirely in C. The FLIH respects the register-use and stack-management requirements of the SPE ABI, and the SLIH registration

13. The Linux kernel supports interfaces such as `in_interrupt()`, `local_irq_enable()`, `local_irq_disable()`, `local_irq_save()`, `local_irq_restore()`, and so forth, for these purposes.

code is written in C. Thus, SPE application developers can write SLIH event handlers entirely in C and register these SLIH event handlers with a simple C function call. The example shows that it is possible to have an SPU interrupt mechanism that is easy to use, although a handler written entirely in assembly code would be more efficient.

The example concludes with application code that registers a decrementer event handler and counts ten decrementer events.

### 18.7.2 FLIH Design

The FLIH code resides at address zero in the LS, so it runs every time the SPE takes an interrupt. The tasks of the FLIH are:

- Preserve the interrupted application stack.
- Invoke the SLIH for the pending event.
- Restore the application stack and return to the interrupted application.

The FLIH assembler code follows:

```
.set STACK_SKIP, 125                # max. size needed to protect unallocated stack space
        .extern spu_slih_handlers
        .text
        .align 3
        .global spu_flih
spu_flih:
        stqd        $SP, -(STACK_SKIP+82)*16($SP)    # save back-chain pointer
        stqd        $0, -(STACK_SKIP+80)*16($SP)     # save volatile registers 0, 2-79
        stqd        $2, -(STACK_SKIP+ 2)*16($SP)
        stqd        $3, -(STACK_SKIP+ 3)*16($SP)
        stqd        $4, -(STACK_SKIP+ 4)*16($SP)
        stqd        $5, -(STACK_SKIP+ 5)*16($SP)
        stqd        $6, -(STACK_SKIP+ 6)*16($SP)
        stqd        $7, -(STACK_SKIP+ 7)*16($SP)
        stqd        $8, -(STACK_SKIP+ 8)*16($SP)
        stqd        $9, -(STACK_SKIP+ 9)*16($SP)
        stqd        $10, -(STACK_SKIP+ 10)*16($SP)
              . . .                                  # omitted regs 11-75 for brevity
        stqd        $76, -(STACK_SKIP+76)*16($SP)
        stqd        $77, -(STACK_SKIP+77)*16($SP)
        stqd        $78, -(STACK_SKIP+78)*16($SP)
        stqd        $79, -(STACK_SKIP+79)*16($SP)
        rdch        $6, $SPU_RdEventMask
        stqd        $6, -(STACK_SKIP+1)*16($SP)      # save event mask on stack
        rdch        $3, $SPU_RdEventStat
        andc        $7, $6, $3                       # clear current pending event bits in mask
        wrch        $SPU_WrEventMask, $7             # disable current pending events
        wrch        $SPU_WrEventAck, $3              # acknowledge current pending events
        il          $2, -(STACK_SKIP+82)*16          # stack frame size
        a           $SP, $SP, $2                     # instantiate flih stack frame
next_event:
        clz         $4, $3                           # determine next (left-most) event
```

```
ila         $5, spu_slih_handlers           # address of slih function table
shli        $4, $4, 2                       # make event number a word offset
lqx         $5, $4, $5                      # load four slih function pointers
rotqby      $5, $5, $4                      # rotate slih pointer to preferred slot
bisl        $0, $5                          # branch-and-link to slih function
brnz        $3, next_event                  # more events? ($3 is slih return value)
lqd         $6, 80*16($SP)
wrch         $SPU_WrEventMask, $6           # re-establish previous event mask
lqd         $79, 2*16($SP)                  # restore volatile registers 79-2, 0
lqd         $78, 3*16($SP)
lqd         $77, 4*16($SP)
lqd         $76, 5*16($SP)
        . . .
lqd         $10, 71*16($SP)
lqd         $9, 72*16($SP)
lqd         $8, 73*16($SP)
lqd         $7, 74*16($SP)
lqd         $6, 75*16($SP)
lqd         $5, 76*16($SP)
lqd         $4, 77*16($SP)
lqd         $3, 78*16($SP)
lqd         $2, 79*16($SP)
lqd         $0, 2*16($SP)
lqd         $SP, 0*16($SP)                # restore stack pointer from back chain ptr
irete                                     # Return through SRR0, re-enable interrupts
```

#### 18.7.2.1 *Preserve the Interrupted Application Stack*

The first section of FLIH code allocates an ABI-compliant stack frame. Because the FLIH cannot know the exact context of the application code that has been interrupted, it must assume the worst-case situation, which happens when two conditions are true:

- A leaf function (a function that makes no function calls) has been interrupted.
- The leaf function uses space in its stack frame without actually allocating the frame.

It is legal for a leaf function to use stack space below its caller's stack frame without actually decrementing the stack pointer to allocate a frame. Thus, to be compatible with any interrupt scenario, the FLIH must protect space on the stack equivalent to the maximum size stack frame and preserve registers so that the code for the SLIH, which is written in an ABI-compliant high-level language, can freely use the processor registers and stack according to the ABI rules.

The example FLIH begins by storing the previous stack frame's back-chain pointer at the end of the new stack frame. Next, it saves the return address from R0 (LR, the Link Register) to the correct position in the new stack frame, just above the back-chain pointer. Next, the FLIH saves volatile registers including the registers that might contain parameters for the interrupted function. Then the FLIH sets the SP to the bottom of the new stack frame. At this point, the stack is ready to support calling ABI-compliant functions.

The FLIH also follows the recommended policy of acknowledging all the current pending events before the first event is processed.

#### 18.7.2.2 *Invoke SLIH for Pending Event*

The middle section of the FLIH gets the pending event status, determines the left-most event, loads, aligns the function pointer for the event's corresponding SLIH, and executes a branch-and-link to the SLIH function. Because an ABI-compliant stack frame has been allocated and volatile registers have been saved, the SLIH can be a high-level-language function.

Note that the event status is in R3, which is the first parameter register according to the ABI protocol. Thus, the effect of the FLIH call to the SLIH is equivalent to the C function call:

```
specific_slih (event_status);
```

When the SLIH returns, the FLIH expects the SLIH to have cleared its corresponding event bit in the `event_status` and to have passed back this updated status as the return value; the return value is passed back in R3 according to the SPE ABI. The FLIH uses the return value to determine the next left-most pending event and loops to call the next event's corresponding SLIH.

#### 18.7.2.3 *Restore Application Stack and Return to Interrupted Application*

When all pending events have been handled, the FLIH restores the context of the interrupted code by restoring the saved volatile registers and unwinding the stack. The SP is reset to point to the stack frame of the interrupted code, and the FLIH exits with an **irete** instruction to re-enabled interrupts.

Note that the FLIH and all the SLIH functions are executed with interrupts disabled.

### 18.7.3 SLIH Design and Registering SLIH Functions

SLIH functions are expected to be written as part of an application that is coded in a high-level language; here, the C language is used. The code to register the SLIH functions is also written in C. Shown next is the code for the default SLIH function, the code for the initialization of the SLIH function table, and the function that application code uses to register SLIH functions.

```
#define SPU_EVENT_ID(_mask)(spu_extract(spu_cntlz(spu_promote(_mask, 0)), 0))

static unsigned int spu_default_slih (unsigned int events)
{
    unsigned int mse;

    mse = 0x80000000 >> SPU_EVENT_ID (events); /* get my (left-most) event number */

    return (events & ~mse);                    /* return updated event bits */
}

spu_slih_func  spu_slih_handlers[33] __attribute__ ((aligned (16))) = {
    spu_default_slih, /* event bits 0 through 18: RESERVED, use default slih */
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
```

```
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,

    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih,

    spu_default_slih,
    spu_default_slih,
    spu_default_slih,
    spu_default_slih, /* event bit 19: Multisource Sync */
    spu_default_slih, /* event bit 20: Privilege attention */
    spu_default_slih, /* event bit 21: Lock-line reservation-lost */
    spu_default_slih, /* event bit 22: Signal-notification 1 */
    spu_default_slih, /* event bit 23: Signal-notification 2 */

    spu_default_slih, /* event bit 24: SPU Write Outbound Mailbox available */
    spu_default_slih, /* event bit 25: SPU Write Outbound Interrupt Mailbox avail. */
    spu_default_slih, /* event bit 26: Decrementer */
    spu_default_slih, /* event bit 27: SPE mailbox */
    spu_default_slih, /* event bit 28: DMA-queue */
    spu_default_slih, /* reserved */
    spu_default_slih, /* event bit 30: DMA list command stall-and-notify */
    spu_default_slih, /* event bit 31: SPE tag-status update */

    spu_default_slih, /* extra slih table entry for phantom events */
};

void spu_slih_reg(unsigned int mask, spu_slih_func func)
{
    unsigned int id;

    while (mask)
    {
        id = SPU_EVENT_ID (mask);      /* get next (left-most) event */
        spu_slih_handlers[id] = func;  /* put function pointer into slih table */
        mask &= ~(0x80000000 >> id);   /* clear that event bit */
    }
}
```

#### 18.7.3.1 *Default SLIH Function*

The default SLIH function, `spu_default_slih`, simply satisfies the FLIH protocol that requires each SLIH to clear its corresponding event bit in the status. The event bit is cleared the in the status and the updated status is returned to the caller (the FLIH), but the default SLIH takes no other action on behalf of the event.

#### 18.7.3.2 *Initialization of SLIH Function Table*

The SLIH function table has 32 entries for SLIH functions that correspond to event bits 0 through 31. Each entry in the table is initialized to the default SLIH function, `spu_default_slih()`.

The function table has a 33rd entry, also initialized to `spu_default_slih()`, that will be accessed when an interrupt is taken with no event bits set in the SPU Read Event Status Channel. Such an interrupt is called a “phantom event” interrupt. In this case, the count-leading-zeros instruction (**clz**) in the FLIH code will return the value 32, which will cause the 33rd entry in the table to be used as the function pointer. For a phantom event interrupt, the correct action is to read the SPU Read Event Status Channel, which clears the phantom event, and return from the interrupt. If a phantom event occurs, the code shown previously handles it properly.

#### 18.7.3.3 *Registering an Application Function in SLIH Table*

The last function in the preceding code, `spu_slih_reg()`, is used by an application to register (install) its SLIH functions. Normally, an application will have one function for each event of interest and will call `spu_slih_reg()` with a single-bit mask to install each SLIH function. This `spu_slih_reg()` function is written, however, to allow a single call to register the same function for multiple events. In the extreme, an application can register a single SLIH for all events of interest; this allows the application’s SLIH to supersede the FLIH dispatch and move most aspects of event processing into application code.

### 18.7.4 Example Application Code

With the interrupt-handling infrastructure described previously, only a small amount of programming is needed to enable and process asynchronous events, and no assembler programming is required.

Shown next is a simple application program that sets up an event handler for decrementer events (a decrementer event occurs when the SPU decrementer count makes the transition from ‘0’ to ‘-1’; see *Section 18.10.1 SPU Decrementer Event* on page 506) and waits until ten such events occur before printing an appropriate message.

```
#include <stdio.h>
#include <spu_mfcio.h>
#define DECR_COUNT        10000    /* count 10000 time base ticks per decrementer interrupt */

volatile unsigned int event_count = 0;

unsigned int decr_handler (unsigned int status)
{
        event_count++;                          /* bump global event count */
        spu_writech (SPU_WrDec, DECR_COUNT);    /* reset decrementer */
```

```
        return (status & ~MFC_DECREMENTER_EVENT); /* return updated event bits */
}

int main (void)
{
        spu_slih_reg (MFC_DECREMENTER_EVENT, decr_handler);    /* register the handler */
        spu_writech (SPU_WrEventMask, MFC_DECREMENTER_EVENT);  /* enable decr. event */
        spu_writech (SPU_WrDec, DECR_COUNT);                   /* init decrementer */
        spu_ienable ();                                        /* enable interrupts */
        while (event_count < 10);          /* busy-wait until 10 decrementer events happen */
        printf ("Test successfully completed.\n");
        return (0);
}
```

#### 18.7.4.1 *Decrementer Event Handler*

The decrementer event handler is based on the default SLIH shown previously. In addition to clearing the event bit in the status, the handler increments the global `event_count` variable, re-initializes the decrementer to the start value, and returns the event status with the decrementer-event bit cleared.

#### 18.7.4.2 *Application Code*

The application code for this example demonstrates how simple it is for SPU application programs to setup and use asynchronous events. The main program first registers the SLIH `decr_handler()` in the SLIH function table by calling the `spu_slih_reg()` function defined previously. Next, the decrementer event is enabled by writing a mask to the SPU Write Event Mask Channel. Next, the decrementer is initialized to the start value `DECR_COUNT` and interrupts are enabled. The empty while loop implements a busy-wait until ten decrementer events occur. Finally, the application prints a success message and exits.

In the application, only four lines of code are required to setup interrupt handling including the line to initialize the decrementer starting count. The decrementer SLIH also needs only four lines of code, including re-initializing the starting count and updating the global event counter. A more sophisticated application that uses the decrementer to help implement execution profiling of a running application is described in *Section 18.10.1 SPU Decrementer Event* on page 506.

## 18.8 Nested Interrupt Handling

The interrupt handling design described previously is simple and clean but has the limitation that all pending events must be processed before another event interrupt can be taken. When many events are handled in one invocation of the FLIH, this limitation can lead to long latency between an event occurring and its handler running. Nested interrupt handling removes this limitation and is beneficial when an application must guarantee minimum delay between the occurrence of a high-priority event and handling the event.

For example, it is best to handle DMA-queue interrupts as soon as they occur. CBEA processor system throughput is maximized when all processing elements, including the DMA engines in the MFCs, are active simultaneously. If the mailbox event handler requires lengthy processing and runs with interrupts disabled, a DMA-queue event that occurs shortly after the mailbox handler starts will not be recognized immediately, and, consequently, overall system throughput might be reduced by an idle DMA engine. If nested interrupt handling is implemented, the mailbox handler can be interrupted, which allows immediate handling of a DMA-queue event to minimize DMA idle time.

### 18.8.1 Nested Handler Design

In the basic design described previously, each SLIH was written as a high-level-language function, which means SLIH functions need no modification to support interruptibility. The FLIH code, however, does need some changes to support interruption of SLIH functions.

The implication of interrupting a SLIH function is that the FLIH can run again before a previous FLIH execution runs to completion. The FLIH itself need not be interruptible, but it must save some extra machine context information and re-enable interrupts before transferring control to an interruptible SLIH function. Because the FLIH transfers control to a SLIH before the FLIH finishes all its work, the FLIH must preserve its own context in addition to the context of the interrupted program.

The example code uses a fixed, bit-ordered priority. The code can be modified to handle interrupts in a different priority by bit-sizzling the interrupt events.

### 18.8.2 FLIH Design for Nested Interrupts

The nested FLIH assembler code follows, with new or changed lines of code shown in bold-face type:

```
        .set STACK_SKIP, 125 # max. size needed to protect unallocated stack space
        .extern spu_slih_handlers
        .text
        .align 3
        .global spu_flih
spu_flih:
        stqd        $SP, -(STACK_SKIP+83)*16($SP)   # save back-chain pointer
        stqd        $0, -(STACK_SKIP+81)*16($SP)    # save volatile registers 0, 2-79
        stqd        $2, -(STACK_SKIP+ 2)*16($SP)
        stqd        $3, -(STACK_SKIP+ 3)*16($SP)
        stqd        $4, -(STACK_SKIP+ 4)*16($SP)
        stqd        $5, -(STACK_SKIP+ 5)*16($SP)
        stqd        $6, -(STACK_SKIP+ 6)*16($SP)
        stqd        $7, -(STACK_SKIP+ 7)*16($SP)
        stqd        $8, -(STACK_SKIP+ 8)*16($SP)
        stqd        $9, -(STACK_SKIP+ 9)*16($SP)
        stqd        $10, -(STACK_SKIP+ 10)*16($SP)
                . . .                               # omitted regs 11-75 for brevity
        stqd        $76, -(STACK_SKIP+76)*16($SP)
        stqd        $77, -(STACK_SKIP+77)*16($SP)
        stqd        $78, -(STACK_SKIP+78)*16($SP)
```

```
        stqd        $79, -(STACK_SKIP+79)*16($SP)
        rdch        $3, $SPU_RdSRR0              # get interrupt return address
        stqd        $3, -(STACK_SKIP+80)*16($SP)  # save int. ret. addr. above link reg.
        rdch        $6, $SPU_RdEventMask
        stqd        $6, -(STACK_SKIP+1)*16($SP)   # save event mask on stack
        rdch        $3, $SPU_RdEventStat
        andc        $7, $6, $3                   # clear current pending event bits in mask
        wrch        $SPU_WrEventMask, $7         # disable current pending events
        wrch        $SPU_WrEventAck, $3          # acknowledge current pending events
        il          $2, -(STACK_SKIP+83)*16      # get stack frame size
        a           $SP, $SP, $2                 # instantiate flih stack frame
next_event:
        clz         $4, $3                       # determine next (left-most) event
        ila         $5, spu_slih_handlers        # address of slih function table
        shli        $4, $4, 2                    # make event number a word offset
        lqx        $5, $4, $5                    # load four slih function pointers
        rotqby $5, $5, $4                        # rotate slih pointer to preferred slot
        bisle       $0, $5                       # call slih with interrupts enabled
        brnz        $3, next_event               # more events? ($3 is slih return value)
        ila       $0, next_inst                  # disable interrupts before restoring SRR0
        bid       $0
next_inst:
        lqd        $6, 80*16($SP)
        wrch        $SPU_WrEventMask, $6         # re-establish previous event mask
        lqd        $0, 3*16($SP)
        wrch       $SPU_WrSRR0, $0              # restore interrupt return address

        lqd        $79, 2*16($SP)               # restore volatile registers 79-2, 0
        lqd        $78, 3*16($SP)
        lqd        $77, 4*16($SP)
        lqd        $76, 5*16($SP)
              . . .
        lqd        $10, 71*16($SP)
        lqd        $9, 72*16($SP)
        lqd        $8, 73*16($SP)
        lqd        $7, 74*16($SP)
        lqd        $6, 75*16($SP)
        lqd        $5, 76*16($SP)
        lqd        $4, 77*16($SP)
        lqd        $3, 78*16($SP)
        lqd        $2, 79*16($SP)
        lqd        $0, 2*16($SP)
        lqd        $SP, 0*16($SP)              # restore stack pointer from back chain ptr
        syncc                                  # force channel write to complete
        irete                                  # Return through SRR0, re-enable interrupts
```

The new FLIH, which allows SLIH functions to be interrupted, has only a few minor changes. First, the stack frame that is allocated to protect the context of the interrupted code has been expanded by one quadword. The extra quadword is required to hold the interrupt return address and maintain quadword alignment of the stack. The interrupt return address will be saved on the

stack just above the saved R0 (LR), so the offsets for the back pointer and Link Register in the stack have been adjusted down by one quadword. The interrupt return address is moved from the Save-Restore Register 0 Channel into R0 and then stored onto the stack.

The indirect-branch instruction that calls SLIH functions, **bisl**, has been changed into the version with the `E` feature bit set, **bisle**. Thus, when the first instruction of the SLIH function executes, interrupts will be enabled.

When the SLIH function returns, the FLIH code, as before, looks for more pending events as indicated by bits set to '1' in the updated SPU Read Event Status Channel that is passed back from the SLIH. When all enabled pending events are serviced by calls to SLIH functions, the FLIH context-restore code is executed. The context-restore code is unchanged except for the offsets in the instructions that restore R0 (LR) and the stack pointer (SP), the two new instructions that restore the Save-Restore Register 0 value, and the **syncc** instruction that forces the SPU Write State Save-and-Restore Channel write to complete before the **irete** executes. Before restoring the Save-Restore Register and stack pointer, the interrupts must be disabled. They are re-enabled at end of the FLIH by executing a **irete** instruction.

## 18.9 Using a Dedicated Interrupt Stack

The previous examples of interrupt handling code have used the application stack for temporary storage space and to support calling SLIH functions. However, using the application stack for interrupt handling creates some problems.

One problem is the fact that the FLIH must allocate a very large frame to account for the worst-case situation: interrupting a leaf procedure that is using a maximum-sized stack frame without actually allocating it. The FLIH must allocate more than 125 quadwords—2000 bytes—on the stack even though only a small fraction of that space is likely ever needed.

Another problem is switching contexts when the runtime environment supports multiple processes or multiple threads. When an interrupt handler can be the cause of a process or thread switch, using a separate, dedicated interrupt stack can simplify preserving the context of the suspended thread and restoring the context of the activated thread.

The FLIH code that follows implements a separate interrupt stack; this dedicated stack is used by the FLIH code and SLIH functions. The major changes in the FLIH are to save the application stack pointer and initialize the SP register to the top of the dedicated interrupt stack. Also, this FLIH allocates a much smaller stack frame because the frame need only have space for the volatile registers that must be preserved across calls to SLIH functions.

When the FLIH is finished, it restores the application stack pointer and returns from the interrupt. In a more complete interrupt handler—one that supports thread switching—the application stack pointer would probably be saved in a thread context structure to support a possible interrupt-induced thread switch.

```
.set INTERRUPT_STACK_SIZE, 2048

        .extern spu_slih_handlers
        .text
        .align 3
        .global spu_flih
```

```
spu_flih:
        stqr $SP, main_stack_ptr
        ila $SP, interrupt_stack-82
        stqd $0, 2*16($SP)                # save volatile registers 0, 2-79
        stqd $SP, 0*16($SP)               # save back chain pointer
        stqd $2, 80*16($SP)
        stqd $3, 79*16($SP)
        stqd $4, 78*16($SP)
        stqd $5, 77*16($SP)
        stqd $6, 76*16($SP)
        stqd $7, 75*16($SP)
        stqd $8, 74*16($SP)
        stqd $9, 73*16($SP)
        stqd $10, 72*16($SP)
        . . .                             # omitted regs 11-75 for brevity
        stqd $76, 6*16($SP)
        stqd $77, 5*16($SP)
        stqd $78, 4*16($SP)
        stqd $79, 3*16($SP)
        rdch $6, $SPU_RdEventMask
        stqd $6, 2*16($SP)              # save event mask on stack
        rdch $3, $SPU_RdEventStat
        andc $7, $6, $3                 # clear current pending event bits in mask
        wrch $SPU_WrEventMask, $7       # disable current pending events
        wrch $SPU_WrEventAck, $3        # acknowledge current pending events
next_event:
        clz $4, $3                      # determine next (left-most) event
        ila $5, spu_slih_handlers       # address of slih function table
        shli $4, $4, 2                  # make event number a word offset
        lqx $5, $4, $5                  # load four slih function pointers
        rotqby $5, $5, $4               # rot slih function pointer to preferred slot
        lqd $4, 2*16($SP)               # pass interrupt return address to slih
        bisl $0, $5                     # call slih
        brnz $3, next_event             # more events? ($3 is slih return value)
        lqd $6, 2*16($SP)
        wrch $SPU_WrEventMask, $6       # re-establish previous event mask
        lqd $79, 3*16($SP)              # restore volatile registers 79-2, 0
        lqd $78, 4*16($SP)
        lqd $77, 5*16($SP)
        lqd $76, 6*16($SP)
                . . .
        lqd $10, 72*16($SP)
        lqd $9, 73*16($SP)
        lqd $8, 74*16($SP)
        lqd $7, 75*16($SP)
        lqd $6, 76*16($SP)
        lqd $5, 77*16($SP)
        lqd $4, 78*16($SP)
        lqd $3, 79*16($SP)
        lqd $2, 80*16($SP)
```

```
        lqd $0, 2*16($SP)
        lqr $SP, main_stack_ptr      # restore application stack pointer
        irete                        # return through SRR0 and re-enable interrupts

        .align 4
        .skip INTERRUPT_STACK_SIZE
interrupt_stack:                     # top of private interrupt stack
main_stack_ptr:                      # Place to save main runtime stack pointer
        .skip 16
```

## 18.10 Sample Applications

This section considers a few basic examples of how external events might be used by a Synergistic Processor Element (SPE) program.

A preferred implementation enables asynchronous interrupts and processes specific conditions with SLIHs, which are typically written in a higher-level language like C. Of course **bisled** polling or individual event polling can be used instead, provided that the event status is checked at appropriate intervals.

### 18.10.1 SPU Decrementer Event

The decrementer and the SPU decrementer event have many potential uses, including the watchdog timer function described in *Section 18.10.3.1 Watchdog Timer* on page 508. A few of these are described in the next few sections.

Each SPE controls a 32-bit decrementer. The decrementer value counts down by ‘1’ for every cycle of the PPE time base (see *Section 13.2.3 Time-Base Frequency* on page 383). If the decrementer event is enabled (the `Tm` bit is set in the SPU Read Event Mask), a decrementer event occurs when the value in the decrementer makes the transition from ‘0’ to ‘-1’ (that is, the most-significant bit changes from ‘0’ to ‘1’).

Because the frequency of the time base is a function of the system, SPE software can assume that a specific decrementer value corresponds to a fixed time interval, regardless of the CBEA processor implementation. Thus, software can use the decrementer for time-based scheduling, performance monitoring, and so forth.

If the value loaded into the decrementer causes a change from ‘0’ to ‘1’ in the MSb, an event is signaled immediately. Setting the decrementer to a value of ‘0’ results in an event after a single decrementer interval.

#### 18.10.1.1 *Execution Profiling or Hot-Spot Analysis*

Often, compute-intensive inner loops offer the best opportunities to improve program performance. Programmers can usually identify inner loops by inspection, yet programmer intuition about where to apply optimization effort is subjective and sometimes wrong.

Execution profiling provides objective data about where a program's run time is actually being spent. High-level-language development systems typically have a way to build a version of any program that will gather execution profile information. This type of profiling works well for non-real-time applications, but it has disadvantages due to the extra code that is inserted by the development system to gather the profile data; this extra code changes the real-time behavior of the program by adding to execution time and changing the memory addresses of processor instructions. As a result of the extra profile-gathering code, cache residency and branch alignments can change in ways that affect execution characteristics.

The disadvantages of adding code for profiling can be reduced by gathering profile data through periodic interrupts. When a time interval expires, an interrupt is taken, and the interrupt handler makes note of where the program was executing by incrementing a count associated with the range of PC addresses that contains the interrupt SPU program counter (PC). The handler maintains an array of counts such that the product of the size of the array and the range of PC addresses per count covers the application code of interest. If the interrupt interval is short enough and the application run time long enough, statistical sampling theory holds that the resulting array of counts accurately represents where the application is spending its execution time. The array of counts can be analyzed when the application terminates or even while the application is running by correlating the PC address ranges with high-level-language statements and procedure names.

#### 18.10.1.2 *Garbage Collection*

A handler for the decrementer event can be used to intermittently run a garbage collector for languages supporting such constructs, such as C++ or Java.

#### 18.10.1.3 *Microthreads*

A user-level threading (microthread) environment for the SPE can use the decrementer event to time slice between tasks that share the same SPE.

### 18.10.2 Tag-Group Status Update Event

The tag-group status update event becomes pending when all the MFC commands are complete in at least one unmasked tag group. Because a tag group by definition contains multiple MFC commands, completion of all commands can take a significant amount of time; instead of wasting time waiting for tag-group completion, it makes sense to have the SPU work on other tasks while the MFC autonomously executes the commands in the tag group.

To support such concurrency, SPE software can include an event handler that responds to the tag-group status update event. The handler can set a buffer-complete flag for each tag group reported in the value read from the MFC Read Tag-Group Status Channel. Software can set up a data structure with status bits for each buffer that is currently being processed by the MFC. The structure can have multiple 32-bit status, if necessary, where the assignment of buffer flag bits in the 32-bit words matches the tag-group bit assignments in the 32-bit word of the MFC Read Tag-Group Status Channel. When the tag-group status update event is recognized, the handler can efficiently set and clear flags, such as buffer-complete or buffer-busy flags, with bit-wise logical operations.

### 18.10.3 DMA List Command Stall-and-Notify Event

The purpose of the DMA list command stall-and-notify event is to allow software to set up and start a sequence, or list, of DMA transfers even when the precise characteristics of all list-element transfers is not known at set-up time. By setting the stall-and-notify flag in one or more list elements, the DMA list command can be queued, and the MFC can immediately begin processing some of the DMA transfers.

After the DMA transfer for the first list element in the list with the stall-and-notify flag is completed, the MFC suspends list processing and raises the DMA list command stall-and-notify event. In response to the event, software can re-evaluate conditions affecting the DMA list transfer and adjust the characteristics—transfer size and effective address—of one or more of the remaining list elements. When adjustments are complete, software restarts the MFC processing of the list by acknowledging the event.

When software adjusts list-element characteristics, it can set the stall-and-notify flag in subsequent elements, and the process of responding the event, re-evaluating conditions, and adjusting list-element characteristics can be repeated.

Two example applications of this event-handling procedure follow.

#### 18.10.3.1 *Watchdog Timer*

In real-time applications, such as computing frames in a live-action 3D game, software must respond to strict deadlines. Software might be unable to complete all requested processing in a real-time interval due to unanticipated input data or environment complexity, unusually high input data rate, and so forth. In a 3D game, if the current frame cannot be completed in its allotted time, it might be acceptable to have software drop the current frame or render a low-quality version and begin rendering the next frame.

SPE software can combine the DMA list command stall-and-notify event with a real-time clock maintained by the SPE decrementer to check buffer progress against a hard, real-time deadline.

When software recognizes the DMA list command stall-and-notify event, it can check the value of the decrementer to find out how much time is left before the next upcoming deadline. If software determines that the remainder of the buffer cannot be transferred before the deadline or that all transfers for the buffer will not finish early enough to allow the SPU to complete some buffer-dependent computation, software can simply cancel the remaining list elements by setting their transfer sizes to ‘0’.

When software cancels unneeded list elements in this way, the system might realize two important benefits. The first benefit is that system power consumption might be reduced by preventing the DMA engine from performing needless work that causes driving and switching internal and possibly external buses. The second benefit is that other time-critical tasks that would have been impeded by competition from the remaining list element transfers might finish earlier because bus bandwidth that would have been wasted on needless transfers becomes available for other, productive transfers.

#### 18.10.3.2 *Linked Data Structure Traversal*

Software can cooperate with the MFC DMA engine to search a complex data structure using a novel application of the DMA list command stall-and-notify event.

Consider the task of searching through some portion of an ordered binary tree, and then performing computation on a matching result. In this example, each node in the tree contains a data element, along with pointers to left and right children. In the tree, internal nodes have a valid pointer for either child, while leaf nodes have left and right child pointers equal to the null pointer. Further, for the purposes of this example, assume that the storage for each tree node has been allocated from a heap in some arbitrary order; for example, the tree was created from randomly ordered data.

A straight-forward implementation coded in a high-level language might fetch a parent node and examine its data to determine if it matches the search key. If a match is found, then the search terminates and the compute phase of the program is started. if a match is not found, either the left or right child pointer is followed depending on whether the search key is logically greater than the current node's data. The search terminates when a match is found, or when the traversal can proceed no further (the child pointer to be followed is the null pointer). This approach has the main disadvantage of serializing tree traversal, node fetching, and computation.

One alternative is to have the DMA engine and a handler for the stall-and-notify event cooperate to asynchronously search for a match in the tree while the main program independently computes. The steps in this cooperative process are as follows:

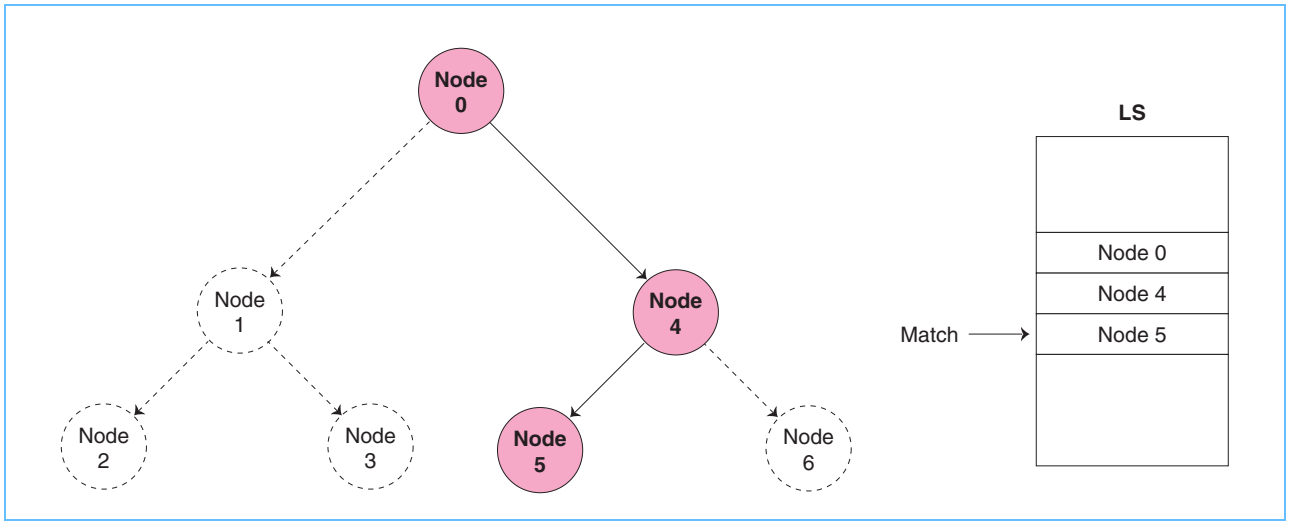
1. Software reserves storage for a DMA list with number of list elements equal to the maximum number of nodes to be searched (the maximum tree depth). Note that individual DMA list elements can be modified after the DMA list command is queued, but the number of elements in the list cannot.
2. For the first element in the DMA list, software sets the effective address to that of the parent node and sets the element's stall-and-notify flag.
3. Software initiates the search operation by enqueuing the DMA list **getl** command and then proceeds to process an independent task; software is assumed to have set up interrupt handling and the DMA list command stall-and-notify event handler at an earlier time.
4. When the interrupt for the stall-and-notify event is taken, a node will have been transferred into SPU LS. The DMA list command stall-and-notify event handler will compare the node's data against the search key, and it terminates the remaining elements in the DMA list—by setting their transfer sizes to '0'—if the data and key match (search success).
5. If the node's data is logically greater than the search key and the left child pointer is not the null pointer, then copy the left child pointer to the effective address of the next element in the DMA list, set the element's stall-and-notify flag, acknowledge the event, and return from the handler to the interrupted program. If the left child pointer is the null pointer, then terminate the remaining elements in the DMA list (search failure).
6. If the node's data is logically less than the search key and the right child pointer is not the null pointer, then copy the right child pointer to the effective address of the next element in the DMA list, set the element's stall-and-notify flag, acknowledge the event, and return from the handler to the interrupted program. If the right child pointer is the null pointer, then terminate the remaining elements in the DMA list (search failure).

Some combination of steps 4-7 will be repeated until the search succeeds or fails. There are at least two benefits to this search method. First, SPU efficiency is improved because all the time spent transferring data into SPU LS can be overlapped with useful SPU computation. Second, a recursive algorithm, which might fail due to a limit on the size of the call stack, is avoided.

*Figure 18-1* on page 510 shows graphically how a search might proceed. First, the algorithm loads node 0 into LS, fails to find a match, and proceeds to the right child. Next, the algorithm loads the right child, node 4, into LS, fails to find a match, and proceeds to the left child. The algorithm loads the left child, node 5 into LS and finds a match. The nodes are loaded into LS autonomously by the DMA engine; checking for a match and setting up the next transfer is done by the event handler. Software performs a minimum of processing while letting the DMA engine do the part of the algorithm it was designed to do well.

In *Figure 18-1*, the nodes are shown being loaded into separate areas of LS due to the limitations of a static drawing. In a real implementation, a single LS area can be used to hold all nodes because only one node is held in LS at a time.

*Figure 18-1. SPE Search of Ordered Binary Tree*




SPE software and MFC hardware can use this algorithm to support multiple simultaneous searches by having the main program initiate several DMA list commands, each with a unique tag group identifier. The DMA List Command Stall-and-Notify handler can maintain a list of active searches and identify which searches have events pending by reading the MFC Read List Stall-and-Notify Tag Status Channel. The handler can resume or terminate each search according to the algorithm described in the preceding paragraphs.

### 18.10.4 MFC SPU Command Queue Available Event

When software has more than sixteen MFC commands to issue, it can issue the first sixteen but then it must wait until space in the queue becomes available to enqueue the remaining commands. If the SPU has no other work to do, it can wait for space by stalling on a write to the MFC Class ID and Command Opcode Channel. If the SPU can proceed with other tasks until space in the MFC queue becomes available, software can set up a handler for the MFC SPU command queue available event so that more MFC commands are enqueued as soon as space is available even though the SPU is attending to other tasks.

Software can set up a queue in LS for MFC commands that are ready for processing. The MFC SPU command queue available event handler takes commands from this queue as MFC queue space becomes available. When the handler finds an empty software queue, it can set a flag that lets the application software know that it must take action to inform the handler that commands

are available to be sent to the MFC queue. The software queue can be made large enough to hold the maximum number of commands generated by application software at any one time. Thus, application software can be written under the assumption that the MFC queue is much larger than it actually is. The software queue and the event handler work together to create this virtual MFC queue.

### 18.10.5 SPU Read Inbound Mailbox Available Event

The SPU read inbound mailbox available event is raised when a message value is written to the inbound mailbox queue. The message can be any 32-bit value, and messages of more than 32 bits can be sent by enqueuing multiple 32-bit words in succession. Software can define a formal message-passing protocol that requires all messages to begin with header word containing message type and length information. The message decoding logic can be encapsulated in the event handler.

For example, consider an application where effective-address and transfer-size information for a DMA command is sent by another processor or device to an SPU through its mailbox. The SPU read inbound mailbox available event handler can read the mailbox data, decode the message, determine that it should handle the message, and initiate a DMA transfer; all this processing can proceed with minimal impact on the SPU as it processes an independent task. Having the handler interact with the MFC tag-group status update event handler and its data structures might increase the efficiency and capability of these asynchronous handlers by allowing tag groups to be used when appropriate.

### 18.10.6 SPU Signal-Notification Available Event

An SPU signal-notification 1 available or SPU signal-notification 2 available events is raised when a signal value is written to a respective signal-notification register in the MFC's memory-mapped I/O (MMIO) address space. Each signal value is a 32-bit word much like a mailbox message value. Unlike mailbox queues, signal notification registers can operate in an OR mode that allows the successive values written to the registers to be accumulated with a logical OR operation.

OR mode allows the single 32-bit signal value to be partitioned so that multiple short messages can be received simultaneously. For example, when software on the SPU and software on another processor participate in a producer-consumer relationship, they need to send and receive simple notices such as a buffer-ready flag. In the simplest case, 32 different flags can be implemented in a single 32-bit signal value.

Because the response times of participants in distributed producer-consumer relationships are typically unpredictable, SPE software can process independent tasks while the SPU signal notification available event handlers respond to signals and perform the simple protocol steps that keep the producer-consumer relationship alive.

### 18.10.7 Lock-Line Reservation Lost Event

The lock-line reservation lost event is raised when the PPE or other device writes to any of the addresses in the lock-line reserved by the **getllar** command. A lock-line reservation lost event handler can be used by an application to receive asynchronous notification that a shared-memory synchronization predicate or condition might have been satisfied.

For example, atomic commands such as the **getllar** command can be executed to reserve a lock-line containing the address of a mutual-exclusion (mutex) lock, also called a semaphore. When the PPE or other device writes to the mutex lock, the lock-line reservation will be lost resulting in the lock-line reservation lost event being raised. In a micro-thread environment, where the SPU implements software multithreading, the event handler can consult a data structure to see that the event allows a previously suspended thread that needs to use the mutex lock to be re-awakened. The handler can interact with the thread-management kernel to move the thread from the suspended list to a list that can be run. A version of the Portable Operating System Interface (POSIX) thread library interfaces `pthread_cond_wait(3)` or `pthread_mutex_lock(3)` for the SPU might use this event to implement efficient user-level threading.

### 18.10.8 Privileged Attention Event

The privileged attention event is raised when the PPE or other device sets the Attention Event Required bit is set in the SPU Privileged Control Register. This event is intended to signal an urgent request for SPU action. The meaning of this event depends on the operating system environment and the organization of resources in the system. One possible meaning for the event is that the SPU has not responded as expected to the passing of a deadline; in this case, the event might signal to SPE software that it must respond in a certain way to prevent the PPE from terminating the SPU tasks.

# 19. DMA Transfers and Interprocessor Communication

## 19.1 Introduction

The Cell Broadband Engine Architecture (CBEA) processors[1] have many attributes of a shared-memory system. The PowerPC Processor Element (PPE) and all Synergistic Processor Elements (SPEs) have coherent access to main storage. But the CBEA processors are not traditional shared-memory multiprocessors. For example, an SPE can execute programs and directly load and store data only from and to its private local storage (LS). In a traditional shared-memory multiprocessor, data communication and synchronization among processors happen at least partially as a side-effect of the fact that all processors use the same shared memory.

Because SPEs lack shared memory, they must communicate explicitly with other entities in the system using three primary communication mechanisms: DMA transfers, mailbox messages, and signal-notification messages. All three communication mechanisms are implemented and controlled by the SPE's memory flow controller (MFC). *Table 19-1* summarizes the three primary mechanisms. This section describes these mechanisms, plus additional MFC mechanisms for synchronizing and otherwise managing DMA transfers and related MFC functions. The channel interface used by an SPE to initiate and monitor these mechanisms is described in *Section 17 SPE Channel and Related MMIO Interface* on page 447.

*Table 19-1. SPE DMA and Interprocessor Communication Mechanisms*

| Mechanism | Description |
|---|---|
| DMA transfers | Used to move data and instructions between main storage and an LS. SPEs rely on asynchronous DMA transfers to hide memory latency and transfer overhead by moving information in parallel with synergistic processor unit (SPU) computation. |
| Mailboxes | Used for control communication between an SPE and the PPE or other devices. Mailboxes hold 32-bit messages. Each SPE has two mailboxes for sending messages and one mailbox for receiving messages. |
| Signal notification | Used for control communication from the PPE or other devices. Signal notification (also called *signaling)* uses 32-bit registers that can be configured for one-sender-to-one-receiver signalling or many-senders-to-one-receiver signalling. |

One type of programming model might rely on the PPE to perform the task of application management by assigning and distributing work to the SPEs. A significant part of this task might be loading main storage with programs and data and then notifying an SPE of available work by either writing to its mailbox or one of its signal-notification registers. After getting the message or signal, the SPE performs a DMA operation to transfer the data and code to its LS. In a variation on this programming model, the PPE might perform the DMA operation and then send a message or signal to the SPE when the DMA operation completes.

After processing the data, an SPE can use another DMA operation to deliver the results to main storage. When the DMA transfer of the SPE results from LS to main storage finishes, the SPE can write a completion message to one of its two outgoing mailboxes that informs the PPE that processing and delivery is complete. If the completion message requires more than 32 bits of information, the SPE can write multiple mailbox messages or use a DMA operation to transfer the

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

long message to main storage where the PPE can read it. Even when a long completion message is transferred with a DMA operation, an outgoing mailbox message can be used to inform the PPE that the message is available.

Although signal notifications and mailbox messages are similar, there are important differences. *Table 19-2* summarizes the differences between using a mailbox and using a signal to communicate.

*Table 19-2. Comparison of Mailboxes and Signals*

| Attribute | Mailboxes | Signals |
|---|---|---|
| Direction | One inbound, two outbound, all accessible through channel interface. | Two inbound accessible through channel interface, but can send signal using MFC send-signal commands. |
| Interrupts | One mailbox can interrupt PPE.<br>Two mailbox-available event interrupts. | Two signal-notification event interrupts. |
| Message Accumulation | No. | Yes: overwrite mode (one-to-one), logical OR mode (many-to-one). |
| Unique SPU Commands | No; programs use channel reads and writes. | Yes, **sndsig**, **sndsigf**, and **sndsigb** for sending signals to other units.<br>See *Section 19.2.3* on page 518. |
| Destructive Read | Reading a mailbox consumes an entry. | Reading a channel resets all 32 bits to ‘0’. |
| Channel Count | Indicates number of available entries. | Indicates waiting signal. |
| Number | Three mailboxes: 4-deep incoming, 1-deep outgoing, 1-deep outgoing with interrupt. | Two signal registers. |

## 19.2 MFC Commands

One of the functions of an MFC is to act as a specialized co-processor for its associated SPU. An MFC has the ability to execute operations from its command set, and it executes them autonomously. The MFC maintains two queues to hold issued commands: the 16-entry MFC SPU command queue and the 8-entry MFC proxy command queue. MFC commands issued by the SPU are buffered only in the MFC SPU command queue; MFC commands issued by the PPE or other system devices are buffered only in the proxy command queue.

An SPU uses its channel interface to send commands to its associated MFC. The SPU can send two types of commands: immediate commands and queueable commands. Queueable commands are entered into the MFC SPU command queue and are processed by the MFC at a time determined by two factors: the MFC execution algorithm and any barrier or fence modifiers attached to commands in the queue. When possible and beneficial, the MFC will execute commands out-of-order. Immediate commands are executed when issued by the SPU and are not entered into a queue.

Other devices in the system, including the PPE, use the MFC’s memory-mapped-I/O (MMIO) registers to send commands to a particular SPU’s associated MFC. The MFC accepts only queueable commands through the MMIO registers. Queueable commands are entered into the MFC proxy command queue, and the MFC processes these commands as it does commands in the MFC SPU command queue: possibly out of order to improve efficiency but respecting barriers and fences. Commands issued through the MMIO registers to the proxy command queue are called proxy commands.

The characteristics of the two command queues are shown in *Table 19-3*. These two queues are independent; the MFC chooses commands to execute from a queue without regard to the activity or conditions in the other queue. Together, these two queues are named the MFC command queues (or queue).

*Table 19-3. MFC Command Queues*

| Queue | Entries | Description |
|---|---|---|
| MFC SPU Command Queue | 16 | For MFC commands sent from the SPU through the channel interface. |
| MFC Proxy Command Queue | 8 | For MFC commands sent from the PPE, other SPUs, or other devices through the MMIO registers. |

Each MFC has a DMA controller that implements DMA transfers. DMA-transfer commands contain both a local storage address (LSA) and an effective address (EA) and thereby initiate transfers between the two storage domains. The MFC converts queued DMA-transfer commands into DMA transfers. Each MFC can maintain and process multiple simultaneous MFC commands and multiple simultaneous DMA transfers.

An MFC can also autonomously perform a sequence of DMA transfers in response to a DMA list command issued by software on its associated SPU.

The majority of MFC commands have names that imply a direction, such as **get**, **put**, and **sndsig**. The data-transfer direction for these MFC commands is always referenced from the perspective of the SPU that is associated with the MFC that executes the commands. Thus, **get** commands transfer data into the LS associated with the MFC; **put** commands transfer data out of the LS associated with the MFC.

Software assigns each MFC queueable command with a 5-bit tag group ID when the command is issued to the MFC. Software can use tag group IDs to monitor the completion of all queued commands in one or more tag groups in a single command queue. SPU software can check the tag-group status by polling, by stalling, or with an asynchronous interrupt.

The MFCs support out-of-order execution of MFC commands. If software needs to maintain order between commands in a queue, it has three ways to do so:

- A command can be issued with its tag-specific barrier feature-bit set to order the command against all preceding and all succeeding commands in the tag group.
- A command can be issued with its tag-specific fence feature-bit set to order the command against all preceding commands in the tag group.
- A separate **barrier** command (*Section 19.2.3* on page 518) can be issued to order the command against all preceding and all succeeding commands in the queue, regardless of tag group.

MFC commands can be classified into two types: DMA commands (*Section 19.2.1*), and synchronization commands (*Section 19.2.3* on page 518). All the DMA commands can be queued; three of the synchronization commands are immediate. The DMA commands can be modified with the suffixes described in *Section 19.2.4* on page 519.

### 19.2.1 DMA Commands

The majority of MFC commands initiate DMA transfers; these are called DMA commands. The basic **get** and **put** DMA commands and all the possible variants are listed in *Table 19-4* and *Table 19-5* on page 517. The command-modifier suffixes for the MFC commands are listed in *Table 19-7* on page 519; these suffixes are used to create the variants shown in *Table 19-4* and *Table 19-5*. Because the LSs of the SPEs and the I/O subsystems are typically mapped into the effective address space, DMA commands can transfer data between an LS and these areas as well.

Regardless of the initiator (SPU, PPE, or other device), DMA transfers move up to 16 KB of data between an LS and main storage. An MFC supports naturally aligned DMA transfer sizes of 1, 2, 4, 8, and 16 bytes and multiples of 16 bytes. For naturally aligned 1, 2, 4, and 8-byte transfers, the source and destination addresses *must* have the same 4 least significant bits.

The performance of a DMA transfer can be improved when the source and destination addresses have the same quadword offsets within a 128-byte cache line. Quadword-offset-aligned transfers generate full cache-line bus requests for every cache line, except possibly the first and last. Transfers that start or end in the middle of a cache line transfer a partial cache line in the first or last bus request, respectively.

The performance of a DMA data transfer when the source and destination addresses have different quadword offsets within a cache line is approximately half[2] that of quadword-aligned transfers, because every bus request is a partial cache-line transfer; in effect, there are two bus requests for each cache line of data.

Peak performance is achieved for transfers in which both the EA and the LSA are 128-byte aligned and the size of the transfer is a multiple of 128 bytes.

Fifteen types of **put** commands move data from LS to main storage. Nine types of **get** commands move data into LS from main storage. For single-transfer DMA code example, see *Section 19.4.1* on page 530.

*Table 19-4. MFC DMA Put Commands* (Sheet 1 of 2)

| Mnemonic | Possible Initiator | | Description |
|---|---|---|---|
| | SPU | PPE | |
| **put** | • | • | Moves data from LS to the effective address. |
| **puts** | | • | Moves data from LS to the effective address and starts the SPU after the DMA operation completes. |
| **putf** | • | • | Moves data from LS to the effective address with fence (this command is locally ordered with respect to all previously issued commands within the same tag group and command queue). |
| **putb** | • | • | Moves data from LS to the effective address with barrier (this command and all subsequent commands with the same tag ID as this command are locally ordered with respect to all previously issued commands within the same tag group and command queue). |
| **putfs** | | • | Moves data from LS to the effective address with fence (this command is locally ordered with respect to all previously issued commands within the same tag group and command queue) and starts the SPU after the DMA operation completes. |

2. For some PowerXCell 8i processor systems configured with large memory, **put** commands are degraded by more than half.

*Table 19-4. MFC DMA Put Commands* (Sheet 2 of 2)

| Mnemonic | Possible Initiator | | Description |
|---|---|---|---|
| | SPU | PPE | |
| **putbs** | | • | Moves data from LS to the effective address with barrier (this command and all subsequent commands with the same tag ID as this command are locally ordered with respect to all previously issued commands within the same tag group and command queue) and starts the SPU after the DMA operation completes. |
| **putl** | • | | Moves data from LS to the effective address using an MFC list. |
| **putlf** | • | | Moves data from LS to the effective address using an MFC list with fence (this command is locally ordered with respect to all previously issued commands within the same tag group and command queue). |
| **putlb** | • | | Moves data from LS to the effective address using an MFC list with barrier (this command and all subsequent commands with the same tag ID as this command are locally ordered with respect to all previously issued commands within the same tag group and command queue). |

*Table 19-5. MFC DMA Get Commands*

| Mnemonic | Possible Initiator | | Description |
|---|---|---|---|
| | SPU | PPE | |
| **get** | • | • | Moves data from the effective address to LS. |
| **gets** | | • | Moves data from the effective address to LS, and starts the SPU after the DMA operation completes. |
| **getf** | • | • | Moves data from the effective address to LS with fence (this command is locally ordered with respect to all previously issued commands within the same tag group and command queue). |
| **getb** | • | • | Moves data from the effective address to LS with barrier (this command and all subsequent commands with the same tag ID as this command are locally ordered with respect to all previously issued commands within the same tag group and command queue). |
| **getfs** | | • | Moves data from the effective address to LS with fence (this command is locally ordered with respect to all previously issued commands within the same tag group), and starts the SPU after the DMA operation completes. |
| **getbs** | | • | Moves data from the effective address to LS with barrier (this command and all subsequent commands with the same tag ID as this command are locally ordered with respect to all previously issued commands within the same tag group and command queue), and starts the SPU after the DMA operation completes. |
| **getl** | • | | Moves data from the effective address to LS using an MFC list. |
| **getlf** | • | | Moves data from the effective address to LS using an MFC list with fence (this command is locally ordered with respect to all previously issued commands within the same tag group and command queue). |
| **getlb** | • | | Moves data from the effective address to LS using an MFC list with barrier (this command and all subsequent commands with the same tag ID as this command are locally ordered with respect to all previously issued commands within the same tag group and command queue). |

### 19.2.2 DMA List Commands

A DMA list is sequence of eight-byte list elements, stored in an SPE's LS, each of which describes a DMA transfer. DMA list commands can be initiated only by SPU programs, not by other devices or by programs running on the PPE. DMA lists are accessed when the MFC executes one of the DMA list commands (**getl**, **getlf**, **getlb**, **putl**, **putlf**, or **putlb**) described in *Table 19-4* on page 516 and *Table 19-5* on page 517.

A DMA list command can specify up to 2048 DMA transfers, each up to 16 KB in length. Thus, a DMA list command can transfer up to 32 MB, which is 128 times the size of the 256 KB LS, more than enough to accommodate future increases in the size of LS. The space required for the maximum-size DMA list is 16 KB (eight bytes for each of 2048 list elements).

DMA list commands are used to move data between a contiguous area in an SPE's LS and possibly noncontiguous area in the effective address space, thus implementing scatter-gather functions between main storage and the LS.

For details and a DMA list code example, see *Section 19.4.4* on page 536.

### 19.2.3 Synchronization Commands

MFC synchronization commands, listed in *Table 19-6*, are used to control the order in which storage accesses are performed. The synchronization commands include four atomic commands (**getllar**, **putllc**, **putlluc**, and **putqlluc**), three send-signal commands (**sndsig**, **sndsigf**, and **sndsigb**), and three barrier commands (**mfcsync**, **mfceieio**, and **barrier**). The MFC can issue up to two outstanding atomic commands, one immediate-type and one **putqlluc**. For details on the four atomic commands, see *Section 20.3 SPE Atomic Synchronization* on page 597.

*Table 19-6. MFC Synchronization Commands*

| Command | Possible Initiator | | Executed Immediately | Description |
|---|---|---|---|---|
| | SPU | PPE | | |
| **getllar** | | • | • | Get lock line and reserve. |
| **putllc** | | • | • | Put lock line conditional. |
| **putlluc** | | • | • | Put lock line unconditional. |
| **putqlluc** | • | • | | Put queued lock line unconditional. |
| **barrier** | • | • | | Barrier type ordering. Ensures ordering of all preceding DMA commands with respect to all commands following the barrier command in the same command queue. The **barrier** command has no effect on the immediate DMA commands: **getllar**, **putllc**, and **putlluc**. |
| **mfceieio** | • | • | | Controls the ordering of **get** commands with respect to **put** commands, and of **get** commands with respect to **get** commands accessing storage that is caching inhibited and guarded. Also controls the ordering of **put** commands with respect to **put** commands accessing storage that is memory coherence required and not caching inhibited. |
| **mfcsync** | • | • | | Controls the ordering of DMA **put** and **get** operations within the specified tag group with respect to other processing units and mechanisms in the system. |
| **sndsig** | • | • | | Write SPU Signal Notification Register in another device. |
| **sndsigf** | • | • | | Write SPU Signal Notification Register in another device, with fence. |
| **sndsigb** | • | • | | Write SPU Signal Notification Register in another device, with barrier. |

### 19.2.4 Command Modifiers

The suffixes listed in *Table 19-7* are included in the mnemonics of DMA commands listed in *Table 19-4* on page 516 and *Table 19-5* on page 517. These suffixes extend or refine the function of a command. For example, a **put** command moves data from LS to the effective address. A **puts** command moves data from LS to an effective address and starts the SPU after the DMA transfer completes. Commands with an **s** suffix can only be issued to the MFC proxy command queue. Commands with an **l** suffix and all the MFC atomic commands can only be issued to the MFC SPU command queue. Commands with an **f** or **b** suffix can be issued to either command queue.

*Table 19-7. MFC Command-Modifier Suffixes*

| Mnemonic | Possible Initiator | | Description |
|---|---|---|---|
| | SPU | PPE | |
| **s** | | • | **Start SPU.** Starts the SPU running at the address in the SPU Next Program Counter Register (SPU_NPC) after the MFC command completes. |
| **f** | • | • | **Tag-specific fence.** Command is locally ordered with respect to all previously issued commands in the same tag group and command queue. |
| **b** | • | • | **Tag-specific barrier.** Command is locally ordered with respect to all previously issued and all subsequently issued commands in the same tag group and command queue. |
| **l** | • | | **List command.** Command processes a list of DMA list elements located in LS. Up to 2048 elements in a list; each list element specifies a transfer of up to 16 KB. |

### 19.2.5 Tag Groups

All DMA commands except **getllar**, **putllc**, and **putlluc** can be tagged with a 5-bit tag group ID. Tag group IDs on commands in the MFC SPU command queue are independent of the tag group IDs on commands in the proxy command queue. Software can ignore the tag group ID capabilities or it can tag one or more commands with the same tag group ID to form a tag group. When multiple commands are in the same tag group, SPU software can use the tag group ID to take advantage of several tag-related features including: checking on the status of a tag group, enabling an interrupt that is raised when one or more tag groups complete, and enforcing an execution order of commands in a tag group.

Programmers can enforce ordering among DMA commands in a tag group with a fence or barrier option by appending an **f**, for fence, or a **b**, for barrier, to the command mnemonic (see *Section 19.2.4* and *Table 19-7*). A fenced command is not executed until all previously issued commands within the same tag group have been performed; commands issued after the fenced command might be executed before the fenced command. A barrier command and all the commands issued after the barrier command are not executed until all previously issued commands in the same tag group have been performed.

#### 19.2.5.1 *Status of Tag Groups Containing MFC Commands*

The SPE uses the following channels, listed in *Table 17-2* on page 450, to query tag groups in the MFC SPU command queue:

- `MFC_WrTagMask`
- `MFC_WrTagUpdate`

- `MFC_RdTagStat`
- `MFC_RdListStallStat`
- `MFC_WrListStallAck`

The PPE uses the following MMIO registers, listed in *Table 17-2* on page 450, to query tag groups in the MFC proxy command queue:

- `Prxy_QueryType`
- `Prxy_QueryMask`
- `Prxy_TagStatus`

The MFC SPU command queue has tag groups 0 to 31, and these are different from MFC proxy command queue tag groups 0 to 31. When one or more commands belonging to a tag group are queued in the MFC, the status of that tag group is not empty. When all commands belonging to a tag group are processed and completed, the status of that tag group becomes empty.

Software can use tag-group identifiers to check the completion status of all queued commands in one or more tag groups. For details and examples, see *Section 19.3.3* on page 525 and *Section 19.4.3* on page 532.

#### 19.2.5.2 *MFC Commands with Tag-Group Dependencies*

The MFC command-execution mechanism checks tag group dependencies in the MFC SPU command queue separately from tag-group dependencies in the MFC proxy command queue; a tag group can have members from only one queue. Tag-group dependencies are determined by the opcode of the command, the command's tag ID, and conditions in the appropriate command queue. Dependencies on still-active commands are cleared as each command is completed.

The **get**, **put**, and **sndsig** commands with **f** (fence) or **b** (barrier) suffixes create a tag-group-specific dependency for the command as compared with other commands in the same tag group and the same command queue.

The **putqlluc** command has a tag-group-specific dependency as compared with other commands in the MFC SPU command queue and a non-tag-specific dependency as compared with other **putqlluc** commands in the MFC SPU command queue.

***Note:*** *The Cell Broadband Engine Architecture specifies that the* ***putqlluc*** *command has an implied tag-specific fence which prevents this command from being issued until all previously issued commands with the same tag have completed. The Cell/B.E. and PowerXCell 8i implementations further add a fence against all other* ***putqlluc*** *commands so that all* ***putqlluc*** *commands are ordered regardless of their tag IDs.*

***Note:*** *The Cell Broadband Engine Architecture specifies that the* ***mfceieio*** *and* ***mfcsync*** *commands are tag-specific. This means that these commands and all subsequent commands in the same tag group are ordered with respect to all previous commands in the queue with the same tag ID. However, the Cell/B.E. and PowerXCell 8i implementations treat these commands as non-tag-specific such that they are ordered with respect to all subsequent and previous commands in the command queue regardless of their tag group.*

### 19.2.6 MFC Command Issue

The MFC supports out-of-order execution of independent MFC commands. The MFC considers commands that have no dependencies to be eligible for issue. The issue mechanism has two slots: slot 0 and slot 1. Each command is assigned to either slot 0 or slot 1 depending on the command type. *Table 19-8* lists the MFC command and their slot assignment.

*Table 19-8. MFC Issue Slot Command Assignments*

| Slot | Instructions |
|---|---|
| Slot 0 | **put, sndsig, putqlluc, sdcrz, sdcrst, sdcrf** |
| Slot 1 | **get, sdcrt, sdcrtst, barrier, mfcsync, mfceieio** |

The MFC issue mechanism alternates issuing commands in slot 0 and slot 1 at the element interconnect bus (EIB) frequency, which is half the SPU frequency. When multiple commands in a slot are eligible, the oldest command in the MFC SPU command queue has priority; if no commands in the MFC SPU command queue are eligible, then the oldest eligible command in the MFC proxy command queue is selected. Thus, commands in the MFC SPU command queue have higher priority than commands in the proxy queue. Sixteen requests can be outstanding.

### 19.2.7 Replacement Class ID and Transfer Class ID

The MFC Class ID and Command Opcode Channel, listed in the MFC Command Parameters group of *Table 17-2* on page 450, can be used to specify both the replacement class ID (`RclassID`) and the transfer class ID (`TclassID`) for each MFC command.

The `TclassID` allows application software to influence the allocation of bus bandwidth for an MFC command. The `TclassID` generates an index into the issue-quota table (IQ0, IQ1, and IQ2 fields in the `MFC_TClassID` register, see *Section 17.9.6* on page 463), which specifies the fraction of bus issue slots to allocate to a transfer class.

The `RClassID` allows privileged software to influence L2-cache and translation lookaside buffer (TLB) replacement for cache misses caused by the MFC command. The `RClassID` is used to generate an index into the replacement management table (RMT), which specifies which members in a congruence class are eligible for replacement. See *Section 6.3 Replacement Management Tables* on page 154 for details.

The MFC Class ID Register performs the same function for commands issued to the MFC by the PPE or the SPU. This register is used to influence cache replacement and bus bandwidth allocation associated with an SPU's MFC, and it has no effect on resources associated with the PPE.

The default class ID ('0') is used for all undefined or invalid class IDs. An invalid class ID does not generate an exception.

#### 19.2.7.1 *TClassID Enabled*

By default, the transfer class ID facility is disabled and the issue mechanism does not use the `TclassID` value of an MFC command. However, performance of the MFC can be improved if privileged software enables `TclassID` in the privileged `MFC_TClassID` register. Enabling the facility allows software to manage the bus and bandwidth allocation for requests to different types of targets, such as on-chip storage, off-chip storage, and I/O.

When TclassID is enabled, the issue mechanism enables round-robin selection between TClassID0, TClassID1, and TClassID2 eligible MFC commands for both slot 0 and slot 1. Slot 0 and slot 1 have separate round-robin tokens that track the class of the last command issued in that slot. When slot alternation (described later) is disabled for a class, that class is skipped by the slot 1 round-robin mechanism, and all entries are issued by the slot 0 round-robin mechanism.

Each TclassID has a quota, specified in the IQ0, IQ1, and IQ2 fields of the MFC_TClassID register, for the number of outstanding requests to the synergistic bus interface (SBI). The sum of these quotas should not exceed sixteen (the depth of the SBI queue). The quota for each class is used to prevent the issue mechanism from selecting an MFC command in that class when the quota is reached. When the number of outstanding requests for a class is below the quota for that class, the issue mechanism can select eligible commands from the class.

Slot alternation of issues between slot 0 and slot 1 can be enabled separately for each TclassID using the SA0, SA1, and SA2 fields of the MFC_TClassID register. Enabling slot alternation for a class allows commands from that class to issue every cycle but alternating between slot 0 and slot 1 at the EIB frequency. Disabling slot alternation for a class forces commands from that class to issue only in slot 0 every 4th cycle at the EIB frequency.

The suggested uses for transfer class IDs are:

- Transfer Class ID 0—For MFC commands that bypass the token request, such as LS to LS (EA translated to LS address) transfers or on-chip MMIO transfers with slot alternation enabled (SA0 = 0) to allow **put** and **get** commands to execute at the same time
- Transfer Class ID 1—For off-chip memory-access MFC commands with slot alternation disabled (SA1 = 1) to reduce the penalties associated with switching the direction of the bidirectional memory interface
- Transfer Class ID 2—For I/O-access MFC commands with slot alternation enabled (SA2 = 0) to allow **put** and **get** commands to execute at the same time

#### 19.2.7.2 *TClassID Disabled*

By default, the transfer class ID facility is disabled, and the issue mechanism does not consider the transfer class ID when it makes issue decisions. When the transfer class ID facility is disabled, the IQ0 (issue quota 0) and SA0 (slot alternation 0) fields in MFC_TClassID affect all MFC commands. IQ0 is sixteen by default to match the depth of the SBI queue. If software lowers the value of IQ0, the maximum number of outstanding DMA requests to the SBI is reduced accordingly. SA0 is 0 by default to allow the issue mechanism to alternate issuing commands in slot 0 and slot 1. If software sets SA0 to 1, this alternation of slots is disabled, and all MFC commands are issued only in slot 0 every fourth cycle at the EIB frequency.

### 19.2.8 DMA-Command Completion

When an MFC command from the SPU is complete, the channel count for MFC_Cmd is incremented to indicate that an entry in the MFC SPU command queue has become available. When an MFC command from the PPE is complete, the MFC_QStatus register (*Table 19-9* on page 523) is updated. Command completion also updates Prxy_TagStatus register (*Table 17-2* on page 450) if the command's Tag Group is empty.

## 19.3 PPE-Initiated DMA Transfers

The PPE or other devices can initiate DMA transfers between main storage and an SPE's LS by accessing the MMIO command-parameter registers maintained by the target SPE's MFC. These MMIO registers are listed in *Table 19-9* on page 523.

*Table 19-9. MFC Command-Parameter Registers for PPE-Initiated DMA Transfers*

| Offset From Base | MMIO Register | Max. Entries | R/W | Width (bits) | Functions |
|---|---|---|---|---|---|
| x'03004' | MFC_LSA | 1 | W | 32 | Specifies the LS address of the DMA transfer. |
| x'03008' | MFC_EAH | 1 | W | 32 | Specifies the high-order half of the EA address of the DMA transfer. |
| x'0300C' | MFC_EAL | 1 | W | 32 | Specifies the low-order half of the EA address of the DMA transfer. |
| x'03010' | MFC_Size | 1 | W | 16 H | Specifies the size of the DMA transfer. |
| | MFC_Tag | 1 | W | 16 L | Specifies an identifier for the DMA-transfer command. |
| x'03014' | MFC_ClassID_CMD | 8 | W | 32 | Specifies the opcode, replacement class ID (RclassID), and transfer class ID (TclassID) for the DMA-transfer command. |
| | MFC_CMDStatus | 1 | R | 32 | Returns information about the success or failure of queuing a DMA command. |
| | MFC_QStatus | 1 | R | 32 | Returns the number of queue entries available. |

### 19.3.1 MFC Command Issue

The PPE issues MFC commands by writing parameters to these registers and writing a command opcode to the `MFC_ClassID_CMD` register. To enqueue an MFC command through the MFC Command Parameter Registers, the registers must be written and read in the following sequence:

1. Write the `MFC_LSA` register. The `MFC_LSA` register must be written with a single 32-bit store.
2. Write the `MFC_EAH` and `MFC_EAL` registers. The `MFC_EAH` register defaults to '0' and need not be written for effective addresses less than 4 GB. Software can write `MFC_EAH` and `MFC_EAL` as two 32-bit stores or one 64-bit store.[3]
3. Write the `MFC_Size` and `MFC_Tag` Parameters. The `MFC_Size` parameter occupies the upper 16 bits of a 32-bit word; `MFC_Tag` occupies the lower 16 bits. `MFC_Size` and `MFC_Tag` can be written with a single 32-bit store or they can be written along with `MFC_ClassID_CMD` in a single 64-bit store.[3]
4. Write the `MFC_ClassID` and `MFC_Cmd` Parameters. The `MFC_ClassID` parameter occupies the upper 16 bits of the `MFC_ClassID_CMD` word; the `MFC_Cmd` parameter occupies the lower 16 bits of a 32-bit word. `MFC_ClassID` and `MFC_Cmd` can be written with a single 32-bit store or they can be written along with the `MFC_Size` and `MFC_Tag` parameters in a single 64-bit store.[3]
5. Read the `MFC_CMDStatus` Register. The read of the `MFC_CMDStatus` register is a 32-bit load.

3. Except when specified explicitly as in this instance, 64-bit access to an address range that includes a 32-bit MMIO register is not allowed.

6. If there was insufficient space in the queue, wait until there is space and retry the sequence starting at step 1.
7. If a sequence error occurred, retry sequence starting at step 1.

The least-significant two bits of the command-status value returned from the read of the `MFC_CMDStatus` register indicate success or error when attempting to enqueue an MFC command, as shown in *Table 19-10*.

*Table 19-10. MFC_CMDStatus Return Values*

| Least-Significant Two Bits | Command Status |
|---|---|
| 0 | The enqueue was successful. |
| 1 | A sequence error occurred while enqueuing the DMA transfer (for example an interrupt occurred, then another DMA transfer was started within an interrupt handler). In this case, the MFC-command enqueue must be restarted from the beginning of the command issue protocol. |
| 2 | The enqueue failed due to insufficient space in the command queue. |
| 3 | Both of these errors occurred. |

*Figure 17-2* on page 459 shows a flowchart of the sequence in which MFC commands are issued by PPE software.

#### 19.3.1.1 *Example: A PPE-Initiated DMA Transfer*

The following C-language program fragment shows how code running on the PPE can initiate a 16 KB DMA transfer from EA x‘10000000’ to LS address x‘500’ using a tag ID of 5.

```
void *ps = get_ps();
unsigned int ls = 0x500;
unsigned int long long ea = 0x10000000;
unsigned int size = 0x4000;
unsigned int tag = 5;
unsigned int classid = 0;
unsigned int cmd = MFC_GET_CMD;
unsigned int cmd_status;

do {
    *((volatile unsigned int *)(ps + MFC_LSA)) = ls;
    *((volatile unsigned long long *)(ps + MFC_EAH)) = ea;
    *((volatile unsigned int *)(ps + MFC_Size)) = (size << 16) | tag;
    *((volatile unsigned int *)(ps + MFC_ClassID)) = (classid << 16) | cmd;
    /*
     * Read MFC_CMDStatus to enqueue command and check enqueue success.
     */
    cmd_status = *((volatile unsigned int *)(ps + MFC_CMDStatus)) & 0x3;

} while (cmd_status); /* Attempt to enqueue until success */
```

### 19.3.2 MFC Command-Queue Control Registers

The PPE can control and monitor an SPE by using the MFC command-queue control registers, which are listed in *Table 19-11*. These MMIO registers are in the SPE problem-state memory region of each SPE along with the MFC Command Parameter Registers.

*Table 19-11. MFC Command-Queue MMIO Registers for PPE-Initiated Commands*

| Offset From Base | MMIO Register | Max. Entries | R/W | Width (bits) | Functions |
|---|---|---|---|---|---|
| x‘03014’ | MFC_CMDStatus | 1 | R | 32 | Returns information about the success or failure of queuing a DMA command. |
| | MFC_QStatus | 1 | R | 32 | Returns the number of queue entries available. |
| x‘0321C’ | Prxy_QueryMask | 1 | R/W | 32 | Specifies the tag groups to be included in a query operation.<br>Returns the current mask value. |
| x‘03204’ | Prxy_QueryType | 1 | R/W | 32 | Specifies a tag-group query-completion condition.<br>Returns query status. A return value of ‘0’ means the query request is complete. |
| x‘0322C’ | Prxy_TagStatus | 1 | R | 32 | Returns the status of the tag groups enabled in the Prxy_QueryMask register. |

### 19.3.3 DMA-Command Issue Status and Errors

The progress of PPE-initiated DMA commands can be monitored in various ways including checking the completion status of tag groups and checking the number of entries available in the MFC proxy command queue. PPE-initiated DMA commands can also fail due to exceptional conditions that include address-translation faults, command errors, and alignment errors. These exceptional conditions are detected and handled by privileged software executing on the PPE.

#### 19.3.3.1 *DMA Tag-Group Completion*

Each of the 32 tag groups (*Section 19.2.5* on page 519) is assigned a bit in the 32-bit `Prxy_QueryMask` register. When software needs to check the completion status of one or more tag groups, it sets the `Prxy_QueryMask` register with a value that has the corresponding bit set for each DMA tag group of interest. Tag group 31 is assigned the most-significant bit of `Prxy_QueryMask`, and tag group 0 is assigned the least-significant bit.

Software can use three basic protocols to determine the completion status of one or more tag groups:

- Poll the Proxy Tag-Group Status Register
- Poll the Proxy Tag-Group Query-Type Register
- Wait for a tag-group completion interrupt

These protocols are described in the following sections.

*Poll Proxy Tag-Group Status Register*

The basic procedure to poll for the completion of an MFC command or group of MFC commands using the Proxy Tag-Group Status Register is as follows:

1. Issue the MFC commands to the MFC proxy command queue.
2. Set the Proxy Tag-Group Query-Mask Register to the groups of interest.
3. Issue an **eieio** instruction before reading the Proxy Tag-Group Status Register to ensure the effects of all previous stores complete.
4. Read the Proxy Tag-Group Status Register.
5. If the value is nonzero, at least one of the tag groups of interest has completed. If polling for all the tag groups of interest to complete, XOR the tag group status value with the tag group query mask. A result of '0' indicates that all groups of interest are complete.
6. Repeat steps 4 and 5 until the tag groups of interest are complete.

The following C-language program fragment shows how to poll for tag group completion of a single tag group. This example assumes that MFC commands have already been issued with a tag-group identifier of 5.

```
void *ps = get_ps();
unsigned int tag_mask = 1 << 5;
unsigned int tag_status;

*((volatile unsigned int *)(ps + Prxy_QueryMask)) = tag_mask;

__asm__("eieio"); /* force write to Prxy_QueryMask to complete */

do {
    tag_status = *((volatile unsigned int *)(ps + Prxy_TagStatus));
} while (!tag_status);
```

The following C-language program fragment shows how to poll for tag group completion of any tag group among many. This example assumes that MFC commands have already been issued with a tag ID of 5, 14, and 31.

```
void *ps = get_ps();
unsigned int tag_mask = (1<<5)|(1<<14)|(1<<31);
unsigned int tag_status;

*((volatile unsigned int *)(ps + Prxy_QueryMask)) = tag_mask;

__asm__("eieio"); /* force Prxy_QueryMask write to complete */

do {
    tag_status = *((volatile unsigned int *)(ps + Prxy_TagStatus));
} while (!tag_status);
```

The following C-language program fragment shows how to poll for tag group completion of all tag groups among many. This example assumes that MFC commands have already been issued with a tag id of 5, 14, and 31.

```
void *ps = get_ps();
unsigned int tag_mask = (1<<5)|(1<<14)|(1<<31);
unsigned int tag_status;

*((volatile unsigned int *)(ps + Prxy_QueryMask)) = tag_mask;

__asm__("eieio");

do {
    tag_status = *((volatile unsigned int *)(ps + Prxy_TagStatus));
} while (tag_status ^ tag_mask);
```

*Poll Proxy Tag-Group Query-Type Register*

The basic procedure to poll for the completion of an MFC command or group of MFC commands using the Proxy Tag-Group Query-Type Register is as follows:

1. Issue the MFC commands to the MFC proxy command queue.
2. Set the Proxy Tag-Group Query-Mask Register to the groups of interest.
3. Request a tag group query by writing a value to the Proxy Tag-Group Query-Type Register: write '01' to indicate a query for completion of any enabled tag group, or write '10' to indicated a query for completion of all enabled tag groups.
4. Read the Proxy Tag-Group Query-Type Register.
5. If the value is '0', then the requested tag group query condition has been met. If the value is nonzero, then repeat step 4 until it returns a value of '0'.

The following C-language program fragment shows how to poll for tag group completion of any tag group. The code will poll for completion of any of the groups enabled by the tag mask; to poll for completion of all of the groups enabled by the tag mask, a tag group query value of '10' is set. This example assumes that MFC commands have already been issued with a tag id of 5, 14, and 31.

```
void *ps = get_ps();
unsigned int tag_mask = (1<<5)|(1<<14)|(1<<31);
unsigned int query_type = 1; /* Use 2 for "all" query. */
unsigned int query_status;

*((volatile unsigned int *)(ps + Prxy_QueryMask)) = tag_mask;
*((volatile unsigned int *)(ps + Prxy_QueryType)) = query_type;

__asm__("eieio");

do {
    query_status = *((volatile unsigned int *)(ps + Prxy_QueryType));
```

```
} while (query_status);
```

***Note:*** *Waiting for any tag-group completion with a 0 tag mask can result in a hang or deadlock condition.*

*Wait for Tag-Group Completion Interrupt*

The basic procedure to wait for a tag-group completion interrupt is as follows:

1. Enable the tag-group completion interrupt by means of PPE privileged software.
2. Issue the MFC commands to the MFC proxy command queue.
3. Set the Proxy Tag-Group Query-Mask Register to the groups of interest.
4. Request a tag-group query by writing a value to the Proxy Tag-Group Query-Type Register: write ‘01’ to indicate a query for completion of any enabled tag group, or write ‘10’ to indicated a query for completion of all enabled tag groups.
5. Wait for the tag-group completion interrupt to occur.
6. PPE privileged software must reset the interrupt status after the interrupt is taken.

The following C-language program fragment shows how to setup up for a tag-group completion interrupt on completion of any tag group. As with the previous example, the tag group query value can be changed to ‘10’ to interrupt on completion of all tag groups. This example assumes that PPE privileged software has already enabled the tag-group completion interrupt, the application has already enabled an interrupt handler, and that MFC commands have already been issued with a tag id of 5, 14, and 31.

```
void *ps = get_ps();
unsigned int tag_mask = (1<<5)|(1<<14)|(1<<31);
unsigned int query_type = 1; /* Use 2 for all query. */

*((volatile unsigned int *)(ps + Prxy_QueryMask)) = tag_mask;
*((volatile unsigned int *)(ps + Prxy_QueryType)) = query_type;
```

Because interrupts are asynchronous, the application can proceed with other tasks while the MFC hardware monitors the tag-group completion status.

### 19.3.3.2 *MFC Proxy Command Queue Space*

Each MFC has an 8-entry MFC proxy command queue for commands initiated by the PPE. As the PPE submits MFC commands, the queue might become full. When the queue is full, PPE software must wait for space to become available before subsequent commands can be issued.

The MFC Queue Status Register contains the current status of the MFC proxy command queue. The most-significant bit is set when the queue is empty. The least-significant 16 bits indicate the number of available slots in the queue. A value of ‘0’ in this field indicates that the queue if full. Software can use the value from this field to set a loop count for the number of MFC commands to enqueue.

The following C-language program fragment shows how to either wait for the MFC proxy command queue to be empty or wait for a specific number of entries, specified in the `space` variable, to become available.

```
void *ps = get_ps();
unsigned int queue_status;

unsigned int space;
  /*
  * When the requested space is greater than eight
  * wait for MFC proxy command queue to become empty.
  */
if (space > 8)
{
    do {
        queue_status = *((volatile unsigned int *)(ps + MFC_QStatus));
    } while (!(queue_status & 0x80000000));
}
else
{
    do {
        queue_status = *((volatile unsigned int *)(ps + MFC_QStatus));
    } while ((queue_status & 0xFFFF) < space);
}
```

## 19.4 SPE-Initiated DMA Transfers

The SPE can initiate DMA transfers between main storage and its LS by accessing its channels, listed in *Table 19-12*.

*Table 19-12. MFC Command-Parameter Channels for SPE-Initiated DMA Transfers*

| SPE Channel # | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Functions |
|---|---|---|---|---|---|---|
| 16 | MFC_LSA | 1 | no | W | 32 | Specifies the LS address of the DMA transfer. |
| 17 | MFC_EAH | 1 | no | W | 32 | Specifies the high-order half of the EA address of the DMA transfer. |
| 18 | MFC_EAL | 1 | no | W | 32 | Specifies the low-order half of the EA address of the DMA transfer. |
| 19 | MFC_Size | 1 | no | W | 16 | Specifies the size of the DMA transfer. |
| 20 | MFC_TagID | 1 | no | W | 16 | Specifies an identifier for the DMA-transfer command. |
| 21 | MFC_Cmd | 16 | yes | W | 32 | Specifies the opcode, replacement class ID (RclassID), and transfer class ID (TclassID) for the DMA-transfer command. |

### 19.4.1 MFC Command Issue

The SPE issues MFC commands by writing parameters to these channels and then writing a command opcode to the MFC_Cmd channel, using the **wrch** instruction, described in *Section 17.1.5* on page 452. To enqueue a DMA command, the first five parameters listed in the following sequence can be written in any order, but they must all be written before the MFC_Cmd parameter, which must be written last. The one exception is the MFC_EAH parameter, which can be skipped if the default value of '0' is acceptable.

1. Write MFC_LSA channel with the LS address.
2. Write MFC_EAH channel with the high 32 bits of the effective address (defaults to '0' if not written).
3. Write MFC_EAL channel with the low 32-bits of the effective address.
4. Write MFC_Size channel with the transfer size.
5. Write MFC_TagID channel with the tag-group identifier value.
6. Write MFC_Cmd channel with the transfer-class and replacement-class IDs and the command opcode.

*Figure 17-2* on page 459 shows a flowchart of the sequence in which MFC commands are issued by SPE software.

#### 19.4.1.1 *Example: Initiating a DMA Transfer from the SPE*

The following examples show how to initiate a DMA transfer from an SPE.

```
extern void dma_transfer(volatile void *lsa,     // local storage address
                         unsigned int eah,       // high 32-bit effective address
                         unsigned int eal,       // low 32-bit effective address
                         unsigned int size,      // transfer size in bytes
                         unsigned int tag_id,    // tag identifier (0-31)
                         unsigned int cmd);      // DMA command
```

An ABI-compliant assembly-language implementation of the subroutine is:

```
        .text
        .global      dma_transfer
dma_transfer:
        wrch         $MFC_LSA, $3
        wrch         $MFC_EAH, $4
        wrch         $MFC_EAL, $5
        wrch         $MFC_Size, $6
        wrch         $MFC_TagID, $7
        wrch         $MFC_Cmd, $8
        bi           $0
```

A comparable C implementation using the SPU composite intrinsic, **spu_mfcdma64**, is:

```
#include <spu_intrinsics.h>

void dma_transfer(volatile void *lsa, unsigned int eah, unsigned int eal,
                        unsigned int size, unsigned int tag_id, unsigned int cmd)
{
       spu_mfcdma64(lsa, eah, eal, size, tag_id, cmd);
}
```

For an example of enqueuing a set of DMA transfers using a DMA list command, see *Section 19.4.4* on page 536.

### 19.4.2 MFC Command-Queue Monitoring Channels

SPE software can monitor the MFC commands in one or more tag groups with the MFC Tag Status channels, listed in *Table 19-13*. An SPE can monitor external events with the MFC Event channels, listed in *Table 19-14*. The following examples show how to use these channels.

*Table 19-13. MFC Tag-Status Channels for Monitoring SPE-Initiated Commands*

| SPE Channel # | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Functions |
|---|---|---|---|---|---|---|
| 22 | MFC_WrTagMask | 1 | no | W | 32 | Specifies the tag groups to be included in a query operation. |
| 23 | MFC_WrTagUpdate | 1 | yes | W | 32 | Specifies when the MFC tag-group status is updated in MFC_RdTagStat. |
| 24 | MFC_RdTagStat | 1 | yes | R | 32 | Returns the status of the tag groups from the last tag-group status update request. |
| 25 | MFC_RdListStallStat | 1 | yes | R | 32 | Returns the tag groups that have an MFC DMA list command in the stall state. |
| 26 | MFC_WrListStallAck | 1 | no | W | 32 | Specifies the tag group number for a tag group with a stalled DMA list command. Writing this tag value to MFC_WrListStallAck restarts (acknowledges) the stalled DMA list command in the tag group. |
| 27 | MFC_RdAtomicStat | 1 | yes | R | 32 | Returns the status of the most-recently completed immediate MFC atomic-update command (the immediate MFC atomic-update commands are getllar, putllc, or putlluc). |

*Table 19-14. MFC Channels for Event Monitoring and Management*

| SPE Channel # | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Functions |
|---|---|---|---|---|---|---|
| 0 | SPU_RdEventStat | 1 | yes | R | 32 | Returns events that are both enabled and pending. |
| 1 | SPU_WrEventMask | 1 | no | W | 32 | Specifies which pending events will affect the value returned by a read of the SPU_RdEventStat channel. |
| 2 | SPU_WrEventAck | 1 | no | W | 32 | Specifies the pending events to be acknowledged (cleared). |

### 19.4.3 DMA Command Issue Status and Errors

The progress of SPU-initiated DMA commands can be monitored in various ways including checking the completion status of tag groups and checking the number of entries available in the MFC SPU command queue. SPU-initiated DMA commands can also fail due to exceptional conditions that include address-translation faults, command errors, and alignment errors. These exceptional conditions are detected and handled by privileged software executing on the PPE.

#### 19.4.3.1 *DMA Tag-Group Completion*

Each MFC command is tagged with a 5-bit tag-group ID. The same identifier can be used for multiple MFC commands. The set of all commands with the same identifier is defined as a tag group. Software can use this tag-group identifier to check the completion status of all queued commands in one or more tag groups. When one or more tag group completes, an interrupt can be raised in the SPU, if enabled by software.

Each of the 32 tag groups is assigned a bit in the 32-bit `MFC_WrTagMask` register. When software needs to check the completion status of one or more tag groups, it sets the `MFC_WrTagMask` register with a value that has the corresponding bit set for each DMA tag group of interest. Tag group 31 is assigned the most-significant bit of `MFC_WrTagMask`, and tag group 0 is assigned the least-significant bit.

Software can use three basic protocols to determine the completion status of one or more commands using tag groups:

- Poll the tag-group update status.
- Wait for tag-group update event.
- Enable interrupt on tag-group update event and process other tasks.

These protocols are described in the following sections.

*Poll Tag-Group Update Status*

The basic procedure for polling the completion of an MFC command or group of commands is:

1. Set the MFC Write Tag-Group Query Mask, `MFC_WrTagMask`.
2. Write ‘0’ to `MFC_WrTagUpdate` to request immediate tag status update.

3. Read MFC_RdTagStat to find out if the tag groups have completed.
4. Repeat steps 2 and 3 as necessary.

The following assembler-language program fragment illustrates the preceding procedure.

```
# Inputs:
#   $0 contains tag mask
# Outputs:
#   $1 contains completed tag status

# Set tag group mask. Once written, the tag mask is retained
# until changed by another write to the tag mask.
    wrch $MFC_WrTagMask, $0

# Set up for immediate tag status update.
    il $1, 0

repeat:
    wrch $MFC_WrTagUpdate, $1
    rdch $1, $MFC_RdTagStat
    brz $1, repeat
```

The following C-language program fragment performs the same function as the preceding assembler-language fragment. The C-language intrinsic **spu_mfcstat(t)** writes the t argument to MFC_WrTagUpdate and then returns the value read from MFC_RdTagStat.

```
#include <spu_intrinsics.h>
#include <spu_mfcio.h>

unsigned int tag_id = 0;
unsigned int tag_mask = 1 << tag_id;

spu_writech(MFC_WrTagMask, tag_mask);

do {
} while (!spu_mfcstat(MFC_TAG_UPDATE_IMMEDIATE)); /* poll for immediate update */
```

*Wait for Tag-Group Update*

The basic procedure for waiting on a conditional tag event (one or more tag-group completions) is:

1. Set the MFC Write Tag-Group Query Mask, MFC_WrTagMask.
2. Write one (MFC_TAG_UPDATE_ANY) to MFC_WrTagUpdate to request completion status for any tag in the mask; write two (MFC_TAG_UPDATE_ALL) if requesting completion status for all tags in the mask.
3. If only waiting for conditional update, read MFC_RdTagStat; SPU will stall until the requested completion status is true.

4. If waiting for any one of multiple events, unmask the MFC tag-group status update event in SPU_WrEventMask and read SPU_RdEventStat; the SPU will stall until an event is raised. See *Section 18* on page 471 for additional methods of detecting and handling SPU external events.

The following assembler-language program fragment illustrates this procedure.

```
# Inputs:
#   $0 contains tag mask
# Outputs:
#   $1 contains completed tag status

# Set tag group mask
    wrch $MFC_WrTagMask, $0

# 0x1 for any tag, 0x2 for all tags.
    il $1, 0x1

# Wait for conditional tag status update (stall the SPU).
    wrch $MFC_WrTagUpdate, $1
    rdch $1, $MFC_RdTagStat
```

***Note:*** *Waiting for any tag-group completion with a 0 tag mask can result in a hang or deadlock condition.*

The following C-language program fragment performs the same function as the preceding assembler-language fragment. The C-language intrinsic **spu_mfcstat(t)** writes the t argument to MFC_WrTagUpdate and then returns the value read from MFC_RdTagStat. If the tag status is not ready, the SPU stalls on the read of MFC_RdTagStat.

```
#include <spu_intrinsics.h>
#include <spu_mfcio.h>

unsigned int tag_id = 0;
unsigned int tag_mask = 1 << tag_id;

spu_writech(MFC_WrTagMask, tag_mask);

/* Wait for all ids in tag group to complete (stall the SPU) */
spu_mfcstat(MFC_TAG_UPDATE_ALL);
```

#### 19.4.3.2 *MFC SPU Command Queue Space*

Each MFC has a 16-entry MFC SPU command queue for commands initiated by the SPE. As SPE software submits MFC commands, the queue can become full. When the queue is full, SPE software must wait for space to become available before subsequent commands can be issued.

Software can use three basic protocols to wait for space to become available in the MFC SPU command queue:

- Poll for queue space.

- Wait for (stall until) queue space is available.
- Wait for (stall until) an external event signals space is available.

*Poll for Queue Space*

The basic procedure for polling for available space in the command queue is:

1. Read the channel count for `MFC_Cmd`.
2. Repeat until the required amount of queue space is available.

The following assembler-language example illustrates this procedure.

```
# Outputs:
#   $0 contains available DMA queue space.

repeat:
    rchcnt $0, $MFC_Cmd
    brz $0, repeat
```

The following C-language program fragment performs the same function as the preceding assembler-language fragment.

```
#include <spu_intrinsics.h>
#include <spu_mfcio.h>

do {
} while (!spu_readchcnt(MFC_Cmd));
```

*Wait for Queue Space*

The basic procedure for waiting for available space in the command queue is:

1. If only waiting for space in the MFC SPU command queue, simply enqueue the MFC command as usual. The SPU will block (stall) on the channel write to `MFC_Cmd` until space becomes available.
2. If waiting for any one of multiple events, enable the MFC SPU command queue available event in `SPU_WrEventMask` and read from `SPU_RdEventStat`. The SPU will stall until an event is raised.

When SPE software has no independent task to perform until MFC SPU command queue space is available, the assembler and C language code sequences are the same as for regular MFC command issue. When software is waiting on MFC SPU command queue space and other events, it reads from `SPU_RdEventStat` and stalls until an event is raised. Software then checks the events that are pending and enabled; if MFC SPU command queue space is available, software can enqueue more MFC commands.

If SPE software has independent tasks to perform while waiting for queue space to become available, software can enable interrupts for asynchronous event handling. See *Section 18* on page 471 for a complete description of event handling and SPU interrupts.

### 19.4.4 DMA List Command Example

*Section 19.2.2* on page 518 summarizes the basics of DMA lists. A DMA list is a sequence of *list elements* that, together with an initiating DMA list command, specifies a sequence of DMA transfers between a single area of LS and possibly discontinuous areas in main storage. Such lists are stored in an SPE's LS, and the sequence of transfers is initiated with a DMA list command such as **getl** or **putl**. DMA list commands can only be issued by programs running on an SPE, but the PPE or other devices can create and store the lists in an SPE's LS. DMA lists can be used to implement scatter-gather functions between main storage and the LS.

Each DMA transfer specified by a list element can transfer up to 16 KB of data, and the list can have up to 2048 (2 K) list elements. Each list element contains a transfer size, the low half of an effective address, and a stall-and-notify bit which can be used to suspend list execution after transferring a list element whose stall-and-notify bit is set.

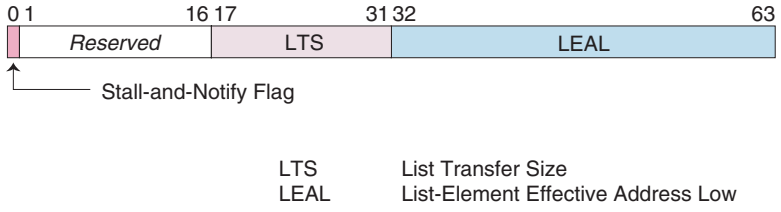
The DMA list command specifies one LS starting address for the entire DMA list, and the data in the LS is accessed sequentially with a minimum step of one quadword. Each list element specifies a new starting address in the effective-address space, but the LS address for a list element starts where the last list element left off. The one exception to this rule is that when a list element in a DMA list contains a short transfer size of 1, 2, 4, or 8 bytes, the LS address increments as needed to start the next list element on a quadword boundary; thus, for a short transfer size, some data in LS will be skipped between list elements. For transfer sizes that are a multiple of a quadword, no data is skipped.

#### 19.4.4.1 *Creating the List*

Software creates the list and stores it in the LS. Lists must be stored in the LS on an 8-byte boundary. The form of a list element is *{LTS, LEAL}*. The first word (LTS) is the list transfer size, the most-significant bit of which serves as an optional *stall-and-notify* flag. The second word (LEAL) is the list element's 32-bit low-order EA. List elements are processed sequentially, in the order they are stored. If the stall-and-notify flag is set for a list element, the MFC will stop processing the DMA list after performing the transfer for that element until the SPE program clears the DMA list command stall-and-notify event from the SPU Read Event Status Channel. This gives programs an opportunity to modify subsequent list elements before they are processed by the MFC.

*Figure 19-1* shows the format of a DMA list element.

*Figure 19-1. DMA List Element*



#### 19.4.4.2 *Initiating the Transfers Specified in the List*

After the list is stored in the LS, the execution of the list is initiated by a DMA list command, such as **getl** or **putl**, from the SPE whose LS contains the list. DMA list commands, like single-transfer DMA commands, require that parameters be written to the MFC Command Parameter channels in the manner described in *Section 19.4.1* on page 530. However, a DMA list command requires two different types of parameters then those required by a single-transfer DMA command:

- *MFC_EAL:* This parameter must be written with the *starting local storage address (LSA) of the list*, rather then with the effective address low (EAL). (The EAL is specified in each list element, although it is named List-Element Effective Address Low, LEAL, as shown in *Figure 19-1*.)
- *MFC_Size:* This parameter must be written with the *size of the list*, in bytes, rather then the transfer size. (The transfer size is specified in each list element as the List Transfer Size (LTS) as shown in *Figure 19-1*.) The list size is equal to the number of list elements, multiplied by the size of the transfer-element structure (8 bytes).

The starting LSA and the EA-high (EAH) are specified only once, in the DMA list command that initiates the transfers. The LSA is internally incremented based on the amount of data transferred by each list element. However, if the starting LSA for each list element does not begin on a 16-byte boundary, then hardware automatically increments the LSA to the next 16-byte boundary.

For list elements smaller than 16 bytes, the least-significant four bits of the transfer-element LSA are equivalent to the least-significant four bits of its EAL.

The EAL for each list element is in the 4 GB area defined by EAH. DMA list element transfers cannot cross the 4 GB area defined by the EAH. If a DMA list element contains an effective address and a transfer size that would result in violation of this rule, the DMA is halted at the 4 GB boundary, and an MFC exception is signalled.

#### 19.4.4.3 *Programming Example*

This C-language sample program creates a DMA list and, in the last line, uses an **spu_mfcdma32** intrinsic to issue a single DMA list command (**getl**) to transfer a main-storage region into LS.

```
/* dma_list_sample.c - SPU MFC-DMA list sample code.
 *
 * This sample defines a list-element data structure, which
 * contains the element's transfer size and low-order 32 bytes of the effective
 * address. Also defined in the structure, but not used by this sample,
 * is the DMA list stall-and-notify bit, which can be used to indicate
 * that the MFC should suspend list execution after transferring a list
 * element whose stall-and-notify bit is set.
 */

#include <spu_mfcio.h>

struct dma_list_elem {
    union {
    unsigned int all32;
```

```
    struct {
       unsigned int stall    : 1;
       unsigned int reserved : 15;
       unsigned int nbytes   : 16;
       } bits;
    } size;
    unsigned int ea_low;
};

struct dma_list_elem list[16] __attribute__ ((aligned (8)));

void get_large_region(void *dst, unsigned int ea_low, unsigned int nbytes)
{
    unsigned int i = 0;
    unsigned int tagid = 0;
    unsigned int listsize;

    /* get_large_region
     *    Use a single DMA list command request to transfer
     *    a "large" memory region into LS. The total size to
     *    be copied may be larger than the MFC's single element
     *    transfer limit of 16kb.
     */

    if (!nbytes)
    return;

    while (nbytes > 0) {
       unsigned int sz;

       sz = (nbytes < 16384) ? nbytes : 16384;
       list[i].size.all32 = sz;
       list[i].ea_low = ea_low;

       nbytes -= sz;
       ea_low += sz;
       i++;
    }


/* Specify the list size and initiate the list transfer */

    listsize = i * sizeof(struct dma_list_elem);
    spu_mfcdma32((volatile *)dst, (unsigned int) &list[0], listsize, tagid,
        MFC_GETL_CMD);
}
```

## 19.5 Performance Guidelines for MFC Commands

- Minimize small transfers. Transfers of less than one cache line consume bus bandwidth equivalent to a full cache-line transfer.
- Align source and destination addresses of large transfers (128 bytes of larger) to a cache-line boundary.
- Have SPEs Initiate DMA transfers. Let the SPEs pull their data instead of PPE pushing the data to the SPE. This is beneficial for four reasons:
  a. There are eight times more SPEs than PPEs.
  b. An MFC SPU command queue is twice as deep (16 entries) as the PPE's MFC proxy command queue (eight entries).
  c. Consumer-managed transfers are easier to synchronize.
  d. The number of cycles required to initiate a DMA transfer from the SPE is smaller than the number to initiate from the PPE.
- Avoid PPE pre-accesses to large data sets, so that most SPE-initiated DMA transfers come from main storage rather than the PPE's L2 cache. DMA transfers from main storage have high bandwidth with moderate latency, whereas transfers from the L2 have moderate bandwidth with low latency.
- Minimize the use of synchronizing and data-ordering commands.

## 19.6 Mailboxes

Mailboxes support the sending and buffering of 32-bit messages between an SPE and other devices, such as the PPE and other SPEs. Each SPE can access three mailbox channels, each of which is connected to a mailbox register in the SPU's MFC. Two one-entry mailbox channels—the SPU Write Outbound Mailbox and the SPU Write Outbound Interrupt Mailbox—are provided for sending messages from the SPE to the PPE or other device. One four-entry mailbox channel—the SPU Read Inbound Mailbox—is provided for sending messages from the PPE, or other SPEs or devices, to the SPE. Each of the two outbound mailbox channels has a corresponding MMIO register that can be accessed by the PPE or other devices.

*Table 19-15* through *Table 19-17* give details about the mailbox channels and their associated MMIO registers.

*Table 19-15. Mailbox Channels and MMIO Registers*

| SPE Channel # | Name | Channel Interface | | | | | MMIO Register Interface | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Offset From Base | Mnemonic | Max. Entries | R/W | Width (bits) |
| 28 | SPU Write Outbound Mailbox | SPU_WrOutMbox | 1 | yes | W | 32 | x‘04004’ | SPU_Out_Mbox | 1 | R | 32 |
| 29 | SPU Read Inbound Mailbox | SPU_RdInMbox | 4 | yes | R | 32 | x‘0400C’ | SPU_In_Mbox | 4 | W | 32 |
| 30 | SPU Write Outbound Interrupt Mailbox[1] | SPU_WrOutIntrMbox | 1 | yes | W | 32 | x‘04000’ | SPU_Out_Intr_Mbox | 1 | R | 64 |
| — | SPU Mailbox Status | — | — | — | — | — | x‘04014’ | SPU_Mbox_Stat | 1 | R | 32 |

1. Access to this MMIO register is available only to privileged PPE software.

*Table 19-16. Functions of Mailbox Channels*

| Channel Interface | SPU Read or Write | Functions |
|---|---|---|
| SPU_WrOutMbox | W | Writes message data to the outbound mailbox. |
| SPU_RdInMbox | R | Returns the next message data from the inbound mailbox |
| SPU_WrOutIntrMbox | W | Writes message data to the outbound interrupt mailbox. |

*Table 19-17. Functions of Mailbox MMIO Registers*

| MMIO Register | PPE Read or Write | Functions |
|---|---|---|
| SPU_Out_Mbox | R | Returns the message data from the corresponding SPU outbound mailbox. |
| SPU_In_Mbox | W | Writes message data to the SPU inbound mailbox. |
| SPU_Out_Intr_Mbox[1] | R | Returns the message data from the corresponding SPU outbound interrupt mailbox. |
| SPU_Mbox_Stat | R | Returns the number of available mailbox entries. |

1. Access to the SPU_Out_Intr_Mbox MMIO register is available only to privileged PPE software.

### 19.6.1 Reading and Writing Mailboxes

Data written by an SPE program to one of the outgoing mailboxes using an SPE write-channel (**wrch**) instruction is available to any processor element or device that reads the corresponding MMIO register in the main-storage space. Data written by a device to the SPU Read Inbound Mailbox using an MMIO write is available to an SPE program by reading that mailbox with a read-channel (**rdch**) instruction.

An MMIO read from either of the outbound mailboxes, or a write to the inbound mailbox, can be programmed to raise an SPE event, which in turn, can cause an SPU interrupt. A **wrch** instruction to the SPU Write Outbound Interrupt Mailbox can also be programmed to cause an interrupt to the PPE or other device, depending on interrupt routing (see *Section 9.6 Direct External Interrupts* on page 265 and *Section 9.8 SPU and MFC Interrupts Routed to the PPE* on page 280).

Each time a PPE program writes to the four-entry inbound mailbox, the channel count for the SPU Read Inbound Mailbox Channel increments. Each time an SPE program reads the mailbox, the channel count decrements. The inbound mailbox acts as a first-in-first-out (FIFO) queue; SPE software reads the oldest data first. If the PPE program writes more than four times before the SPE program reads the data, then the channel count stays at '4', and the fourth location contains the last data written by the PPE. For example, if the PPE program writes five times before the SPE program reads the data, then the data read is the first, second, third, and fifth messages that were written. The fourth message that was written is lost.

### 19.6.2 Mailbox Blocking

Mailbox operations are blocking operations for an SPE: an SPE write to an outbound mailbox that is already full stalls the SPE until a entry is cleared in the mailbox by a PPE read. Similarly, a **rdch** instruction—but not a read-channel-count (**rchcnt**) instruction—from an empty inbound mailbox is stalled until the PPE writes to the mailbox. That is, if the channel count is '0' for a blocking channel, then an **rdch** or **wrch** instruction for that channel causes the SPE to stall until the channel count changes from '0' to nonzero.

To prevent stalling, SPE software should read the channel count associated with the mailbox before deciding whether to read or write the mailbox channel. This stalling behavior for the SPE does not apply to the PPE; if the PPE sends a message to the inbound mailbox and the mailbox is full, the PPE will not stall.

Because a **wrch** instruction will stall when it tries to send a value to a full outbound mailbox, SPE software cannot over-run an outbound mailbox. All outbound messages must be read by some entity outside the SPU before space is made available for more outbound messages. In contrast, the SPU Read Inbound Mailbox can be over-run by an outside entity because the MMIO write used for this purpose does not stall. When the PPE or other device writes to a full inbound mailbox, the last value written to it is overwritten and is lost. See the code sample in *Section 19.7.6.1* on page 553 for a safeguard to avoid this situation.

### 19.6.3 Dealing with Anticipated Messages

There are at least three ways to deal with anticipated mailbox messages:

- SPE software uses an **rdch** instruction to read the inbound mailbox channel, which will block until a mailbox message arrives.
- SPE software uses a **rchcnt** instruction on the mailbox channel, which will return the count ('0' or nonzero). If the count is '0', software can work on other tasks and check again later.
- SPE software enables interrupts to respond to mailbox events. Software can work on other tasks and need never check because the mailbox event interrupt handlers will be activated when mailbox status changes.

### 19.6.4 Uses of Mailboxes

Mailbox message values are intended to communicate messages up to 32 bits in length, such as buffer completion flags or program status. In fact, however, they can be used for any short-data-transfer purpose, such as sending of storage addresses, function parameters, command parameters, and state-machine parameters.

Mailboxes are useful, for example, when the SPE places computational results in main storage via DMA. After requesting the DMA transfer, the SPE waits for the DMA transfer to complete and then writes to an outbound mailbox to notify the PPE that its computation is complete.

If the SPE sends a mailbox message after waiting for a DMA transfer to complete, this ensures only that the SPE's LS buffers are available for reuse. It does not guarantee that data has been coherently written to main storage. The SPE might solve this problem by issuing an **mfcsync** command before notifying the PPE. But doing so is inefficient. Instead, the preferred method is to have the PPE receive the notification and then issue an **lwsync** instruction before accessing any of the resulting data.

Alternatively, an SPE can notify the PPE that it has completed computation by writing, via DMA, such a notification to main storage, from which the PPE can read the notification. (This is sometimes referred to as a *writeback DMA command*, although no such DMA command is defined.) In this case, the data and the writeback must be ordered. To ensure ordering, an `mfceieio` command must be issued between the data DMA commands and the notification to main storage.

Although the mailboxes are primarily intended for communication between the PPE and the SPEs, they can also be used for communication between an SPE and other SPEs, processors, or devices. For this to happen, privileged software needs to allow one SPE to access the mailbox register in another SPE by mapping the target SPE's problem-state area into the EA space of the source SPE. If software does not allow this, then only atomic operations and signal notifications are available for SPE-to-SPE communication.

### 19.6.5 SPU Outbound Mailboxes

The MFC provides two one-entry mailbox channels—the SPU Write Outbound Mailbox and the SPU Write Outbound Interrupt Mailbox—for sending messages from the SPE to the PPE or other device that can read from the main-storage space.

#### 19.6.5.1 *SPU Write Outbound Mailbox Channel*

The value written to the SPU Write Outbound Mailbox Channel (`SPU_WrOutMbox`) is entered into the outbound mailbox in the MFC if the mailbox has capacity to accept the value. If the mailbox can accept the value, the channel count for `SPU_WrOutMbox` is decremented by '1'.

If the outbound mailbox is full, the channel count will read as '0'. If SPE software writes a value to `SPU_WrOutMbox` when the channel count is '0', the SPU will stall on the write. The SPU remains stalled until the PPE or other device reads a message from the outbound mailbox by reading the MMIO address of the mailbox. When the mailbox is read through the MMIO address, the channel count is incremented by '1'.

#### 19.6.5.2 *SPU Write Outbound Interrupt Mailbox Channel*

The value written to the SPU Write Outbound Interrupt Mailbox Channel (`SPU_WrOutIntrMbox`) is entered into the outbound interrupt mailbox if the mailbox has capacity to accept the value. If the mailbox can accept the message, the channel count for `SPU_WrOutIntrMbox` is decremented by '1', and an interrupt is raised in the PPE or other device, depending on interrupt enabling and routing. There is no ordering of the interrupt and previously issued MFC commands.

The raising of an interrupt depends on privileged PPE software having enabled the interrupt for this channel, as described in *Section 9.6 Direct External Interrupts* on page 265 and *Section 9.8 SPU and MFC Interrupts Routed to the PPE* on page 280. After handling the PPE-side interrupt, if enabled, privileged PPE software performs an MMIO read of the SPU Write Outbound Interrupt Mailbox Register to get the value written by the SPU. In response to the MMIO read, the channel count is incremented by '1'. This is the simplest use of the SPU Write Outbound Interrupt Mailbox.

If the outbound interrupt mailbox is full, the channel count will read as '0'. If SPE software writes a value to `SPU_WrOutIntrMbox` when the channel count is '0', the SPU will stall on the write. The SPU remains stalled until the PPE or other device reads a mailbox message from the outbound interrupt mailbox by reading the MMIO address of the mailbox. When this is done, the channel count is incremented by '1'.

The field assignments and definitions for SPU Write Outbound Interrupt Mailbox are identical to those for the SPU Write Outbound Mailbox, shown in *Section 19.6.5.1* on page 542.

#### 19.6.5.3 ***SPU Side***

*Writing SPU Write Outbound Mailbox Data*

SPE software can write to the SPU Write Outbound Mailbox Channel to put a mailbox message in the SPU Write Outbound Mailbox. This write-channel instruction will return immediately if there is sufficient space in the SPU write outbound mailbox queue to hold the message value. If there is insufficient space, the write-channel instruction will stall the SPU until the PPE reads from this mailbox.

The same behavior applies to the SPU Write Outbound Interrupt Mailbox on the SPU side, the only difference is the channel number.

The following assembler-language program fragment shows how to write to an SPU Write Outbound Mailbox.

```
# To write to the SPU Write Outbound Interrupt Mailbox
# instead of the SPU Write Outbound Mailbox, simply replace
# SPU_WrOutMbox with SPU_WrOutIntrMbox in the
# following example.
# Input:
#        $1 contains the mailbox value to be written

wrch $SPU_WrOutMbox, $1
```

The following C-language program fragment shows how to write to the SPU Write Outbound Mailbox.

```
unsigned int mb_value;

spu_writech(SPU_WrOutMbox, mb_value);
```

*Waiting to Write SPU Write Outbound Mailbox Data*

To avoid an SPU stall condition, the SPU can use the read-channel-count instruction on the SPU Write Outbound Mailbox Channel to determine if the queue is empty before writing to the channel. If the read-channel-count instruction returns '0', the SPU write outbound mailbox queue is full. If the read channel-count instruction returns a nonzero value, the value indicates the number of free entries in the SPU write outbound mailbox queue. When the queue has free entries, the SPU can write to this channel without stalling the SPU.

The same behavior described in the preceding paragraph also applies to the SPU write outbound interrupt mailbox queue.

The following assembler-language program fragment shows how to poll SPU Write Outbound Mailbox or SPU Write Outbound Interrupt Mailbox.

```
# To write to the SPU Write Outbound Interrupt Mailbox
# instead of the SPU Write Outbound Mailbox, simply replace
# SPU_WrOutMbox with SPU_WrOutIntrMbox in the
# following example.
# Input:
#        $1 contains the mailbox value to be written

repeat:
    rchcnt $2, $SPU_WrOutMbox
    brz $2, repeat

    wrch $SPU_WrOutMbox, $1
```

The following C-language program fragment shows how to poll SPU Write Outbound Mailbox or SPU Write Outbound Interrupt Mailbox.

```
/*
 * To write the value 1 to the SPU Write Outbound Interrupt Mailbox instead
 * of the SPU Write Outbound Mailbox, simply replace SPU_WrOutMbox
 * with SPU_WrOutIntrMbox in the following example.
 */

unsigned int mb_value;

do {
    /*
     * Do other useful work while waiting.
     */
} while (!spu_readchcnt(SPU_WrOutMbox));

spu_writech(SPU_WrOutMbox, mb_value);
```

The SPU can also avoid a stall by enabling the SPU event facility, described in *Section 18* on page 471. Using events involves the following steps:

1. Enable the PPE core mailbox available event or the PPE core interrupt mailbox available event in SPU Write Event Mask.
2. Monitor for this event. Software can poll with a read-channel-count instruction on SPU Read Event Status; software can block (stall) with a read-channel instruction on the SPU Read Event Status Channel; software can enable SPU interrupts and use interrupt handlers to service mailbox events as they are encountered asynchronously.
3. When the event is detected, it must be acknowledged by writing the corresponding bit to the SPU Write Event Acknowledgment Channel. A typical handler for this event performs the acknowledgment and also writes a value to SPU Write Outbound Mailbox or SPU Write Outbound Interrupt Mailbox.

The following assembler-language program fragment shows how to poll for or block on an SPU write outbound mailbox available event. To block, instead of poll, simply delete the following repeat loop. To write to the SPU Write Outbound Interrupt Mailbox instead of the SPU Write Outbound Mailbox, simply replace `SPU_WrOutMbox` with `SPU_WrOutIntrMbox` and use the appropriate event bit for SPU Write Outbound Interrupt Mailbox in the following example.

```
# Inputs:
#   $2 contains SPU Write Outbound Mailbox event bit
#   $5 contains the 32-bit value to be written to the SPU Write Outbound Mailbox

    il   $2, 0x80
    wrch $SPU_WrEventMask, $2           # Write event mask

repeat:
    rchcnt $3, $SPU_RdEventStat
    brz $3, repeat                      # Wait for event to occur

    rdch $4, $SPU_RdEventStat           # Get event bit
    wrch $SPU_WrEventAck, $2            # Acknowledge event

    wrch $SPU_WrOutMbox, $5             # Send the mailbox message
```

The following C-language program fragment shows how to poll for or block on an SPU write outbound mailbox available event. To block, instead of poll, simply delete the following do/while loop. To write to the SPU Write Outbound Interrupt Mailbox instead of the SPU Write Outbound Mailbox, simply replace `SPU_WrOutMbox` with `SPU_WrOutIntrMbox` and use the appropriate event bit for SPU Write Outbound Interrupt Mailbox in the following example.

```
#define MBOX_AVAILABLE_EVENT  0x00000080
unsigned int event_status;
unsigned int mb_value;

spu_writech(SPU_WrEventMask, MBOX_AVAILABLE_EVENT);

do {
    /*
     * Do other useful work while waiting.
     */
```

```
} while (!spu_readchcnt(SPU_RdEventStat));

event_status = spu_readch(SPU_RdEventStat);          /* read status */

spu_writech(SPU_WrEventAck, MBOX_AVAILABLE_EVENT);   /* acknowledge event */

spu_writech(SPU_WrOutMbox, mb_value);                /* send mailbox message */
```

***Note:*** *The preceding example does not fully accommodate the occurrence of a phantom event. For details on handling phantom events, see Section 18.6.1 on page 481.*

#### 19.6.5.4 *PPE Side*

Before PPE software can read data from one of the SPU Write Outbound Mailboxes, it must first read the Mailbox Status Register to determine that unread data is present in the SPU Write Outbound Mailbox or SPU Write Outbound Interrupt Mailbox; otherwise, stale or undefined data might be returned.

To determine that unread data is available in the SPU Write Outbound Mailbox, PPE software reads the Mailbox Status Register and extracts the count value from the `SPU_Out_Mbox_Count` field. If the count is nonzero, then at least one unread value is present. If the count is '0', PPE software should not read the SPU Write Outbound Mailbox Register because it will get incorrect data and should poll the Mailbox Status Register.

The following C-language example shows how PPE software can read from the SPU Write Outbound Mailbox of an SPE. The example assumes the existence of a operating system call, like **get_ps**, that returns a base address for the problem state area, and some defines, like `SPU_Mbox_Stat` and `SPU_Out_Mbox`, that specify offsets into the problem state area corresponding to the `SPU_Mbox_Stat` and `SPU_Out_Mbox` registers, respectively.

```
void *ps = get_ps();
unsigned int mb_status;
unsigned int new;
unsigned int mb_value;

do {
    mb_status = *((volatile unsigned int *)(ps + SPU_Mbox_Stat));
    new = mb_status & 0x000000FF;
} while ( new == 0 );

mb_value = *((volatile unsigned int *)(ps + SPU_Out_Mbox));
```

SPE software typically uses the SPU Write Outbound Interrupt Mailbox when PPE software needs to be interrupted to detect mailbox data from an SPU. For details on using privileged operations to enable, process, and reset the PPE interrupt associated with the SPU Write Outbound Interrupt Mailbox, see *Section 9 PPE Interrupts* on page 239.

As for the SPU Write Outbound Mailbox, a nonzero value in the `SPU_Out_Intr_Mbox_Count` field of the Mailbox Status Register indicates that unread data is present in the SPU Write Outbound Interrupt Mailbox. The same code example provided previously for the SPU Write Outbound

Mailbox can be used for reading the SPU Write Outbound Interrupt Mailbox, after changes have been made for the location of the SPU_Out_Intr_Mbox_Count field in the Mailbox Status Register and the appropriate offset is used for the SPU Write Outbound Interrupt Mailbox Register.

### 19.6.6 SPU Inbound Mailbox

The MFC provides one mailbox for a PPE to send information to an SPU: the SPU Read Inbound Mailbox. This mailbox has four entries; that is, the PPE can have up to four 32-bit messages pending at a time in the SPU Read Inbound Mailbox.

#### 19.6.6.1 *SPU Read Inbound Mailbox Channel*

If the SPU Read Inbound Mailbox Channel (SPU_RdInMbox) has a message, the value read from the mailbox is the oldest message written to the mailbox. If the inbound mailbox is empty, the SPU_RdInMbox channel count will read as ‘0’.

If SPE software reads from SPU_RdInMbox when the channel count is ‘0’, the SPU will stall on the read. The SPU remains stalled until the PPE or other device writes a message to the mailbox by writing to the MMIO address of the mailbox.

When the mailbox is written through the MMIO address, the channel count is incremented by ‘1’. When the mailbox is read by the SPU, the channel count is decremented by '1'.

#### 19.6.6.2 *PPE Side*

An important difference between the SPU Read Inbound Mailbox and the SPU Write Outbound Mailboxes is that the SPU Read Inbound Mailbox can be overrun by a PPE. A PPE writing to the SPU Read Inbound Mailbox will not stall when this mailbox is full. When a PPE overruns the SPU Read Inbound Mailbox, mailbox message data will be lost.

There is a simple mechanism for the PPE to use to avoid overruns. The fields of the SPU Mailbox Status Register, SPU_Mbox_Stat, shown in *Table 19-18*, contains an SPU_In_Mbox_Count field that contains the number of available entries in the SPU Read Inbound Mailbox. To avoid overruns, the PPE should read the SPU_Mbox_Stat register and send no more mailbox values than indicated in the SPU_In_Mbox_Count field.

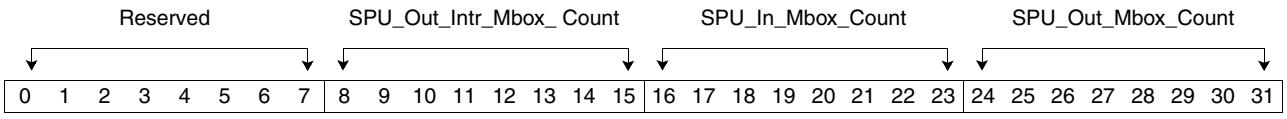



*Table 19-18. Fields of Mailbox Status Register (SPU_Mbox_Stat)*

| Bits | Field Name | Description |
|---|---|---|
| 0:7 | Reserved | Set to zeros. |
| 8:15 | SPU_Out_Intr_Mbox_ Count | Number of valid entries in the SPU Write Outbound Interrupt Mailbox. |
| 16:23 | SPU_In_Mbox_Count | Number of available entries in the SPU Read Inbound Mailbox. |
| 24:31 | SPU_Out_Mbox_Count | Number of valid entries in the SPU Write Outbound Mailbox. |

The following C-language program fragment shows how to write four 32-bit messages to the SPU Read Inbound Mailbox of a particular SPU from the PPE. The code avoids overrunning the SPU Read Inbound Mailbox. The example assumes the existence of a operating system call, like **get_ps**, that returns a base address for the problem state area, and some defines, like `SPU_Mbox_Stat`, and `SPU_In_Mbox`, that specify offsets into the problem state area corresponding to the `SPU_Mbox_Stat` and `SPU_In_Mbox` registers, respectively.

```
void *ps = get_ps();
unsigned int j,k = 0;
unsigned int mb_status;
unsigned int slots;
unsigned int mb_value[4] = {0x1, 0x2, 0x3, 0x4};
do {
    /*
     * Poll the Mailbox Status Register until the SPU_In_Mbox_Count
     * field indicates there is at least one slot available in the
     * SPU Read Inbound Mailbox.
     */
    do {
        mb_status = *((volatile unsigned int *)(ps + SPU_Mbox_Stat));
        slots = (mb_status & 0x0000FF00) >> 8;
    } while ( slots == 0 );

    for (j=0; j<slots && k < 4; j++) {
        *((volatile unsigned int *)(ps + SPU_In_Mbox)) = mb_value[k++];
    }
} while ( k < 4 );
```

#### 19.6.6.3 *SPU Side*

*Reading SPU Read Inbound Mailbox Data*

SPE software can use a read-channel instruction on the SPU Read Inbound Mailbox Channel to read the contents of its SPU Read Inbound Mailbox. This channel read will return immediately if any data written by the PPE is waiting in the SPU Read Inbound Mailbox. This read-channel instruction will cause the SPU to stall if the SPU Read Inbound Mailbox is empty.

The following assembler-language program fragment shows how to read the SPU Read Inbound Mailbox:

```
# Outputs:
#        $1 contains the mailbox value read
rdch $1, $SPU_RdInMbox
```

The following C-language program fragment also shows how to read the SPU Read Inbound Mailbox. Assume that the spu_ macro names have been defined previously.

```
unsigned int mb_value;

mb_value = spu_readch(SPU_RdInMbox);
```

*Waiting for SPU Read Inbound Mailbox Data*

To avoid an SPU stall condition, SPE software can use the read-channel-count instruction on the SPU Read Inbound Mailbox Channel to determine if the mailbox is empty. If the read-channel-count instruction returns a '0', the mailbox is empty and a read-channel instruction would stall. If the read channel-count instruction returns a nonzero value, the value indicates the number of waiting mailbox message values in the SPU Read Inbound Mailbox that can be read without stalling the SPE.

The following assembler-language program fragment shows how to poll an SPU Read Inbound Mailbox read channel.

```
# Outputs:
#  $2 contains the SPU Read Inbound Mailbox value

repeat:
   rchcnt $1, $SPU_RdInMbox
   brz $1, repeat

rdch $2, $SPU_RDInMbox
```

The following C-language program fragment shows how to poll an SPU Read Inbound Mailbox read channel.

```
unsigned int mb_value;

do {
    /*
     * Do other useful work while waiting.
     */
} while (!spu_readchcnt(SPU_RdInMbox));

mb_value = spu_readch(SPU_RdInMbox);
```

The SPU can also avoid a stall by enabling the SPU event facility, described in *Section 18* on page 471. Using events involves the following steps:

1. Enable the SPU inbound mailbox available event in SPU Write Event Mask.
2. Monitor for this event. Software can poll with a read-channel-count instruction on SPU Read Event Status; software can block (stall) with a read-channel instruction on the SPU Read

Event Status Channel; software can enable SPU interrupts and use interrupt handlers to service mailbox events as they are encountered asynchronously.

3. When the event is detected, it must be acknowledged by writing the corresponding bit to the SPU Write Event Acknowledgment Channel. A typical handler for this event performs the acknowledgement and also reads one or more inbound mailbox values from the SPU Read Inbound Mailbox Channel.

The following assembler-language program fragment shows how to poll or block for an SPU read inbound mailbox available event. To block, instead of poll, simply delete the following repeat loop. Only one mailbox value is read to simplify this example, although more than one might have been indicated by the SPU Read Inbound Mailbox Channel count.

```
# Inputs:
#   $2 contains SPU Read Inbound Mailbox event bit
# Outputs:
#   $6 contains one SPU Read Inbound Mailbox value

    il     $2, 0x10
    wrch   $SPU_WrEventMask, $2      # Enable only inbound mailbox event

repeat:
    rchcnt $3, $SPU_RdEventStat
    brz    $3, repeat                # Wait for event

    rdch   $4, $SPU_RDEventStat      # Read event status
    wrch   $SPU_WrEventAck, $2       # Use the value to acknowledge the event
    rdch   $6, $SPU_RdInMbox         # Read first available message
```

The following C-language program fragment shows how to poll or block for an SPU read inbound mailbox available event. To block, instead of poll, simply delete the following do/while loop. All available mailbox messages are placed in the `mb_value` array, and the `mb_count` is set to indicate the number of messages read.

```
#define MFC_SPU_In_Mbox_Event  0x00000010
unsigned int event_status;
unsigned int mb_count, i;
unsigned int mb_value[4];

spu_writech(SPU_WrEventMask, MFC_SPU_In_Mbox_Event);

do {
   /*
    * Do other useful work while waiting.
    */
} while (!spu_readchcnt(SPU_RdEventStat));

event_status = spu_readch(SPU_RdEventStat);

spu_writech(SPU_WrEventAck, event_status);
```

```
mb_count = spu_readchcnt(SPU_RdInMbox);
for (i=0; i< mb_count; i++) {
    mb_value[i] = spu_readch(SPU_RdInMbox);
}
```

## 19.7 Signal Notification

SPE signal-notification channels are connected to inbound registers (into the SPE). The PPE, other SPEs, and other devices use the signal notification registers to send information, such as a buffer-completion synchronization flag, to an SPE. An SPE has two 32-bit signal-notification registers, each of which has a corresponding MMIO register that can be written with signal-notification data.

When SPE software reads a signal-notification channel, hardware clears the channel atomically. In contrast, a read in the MMIO main-storage space does not clear a signal-notification register. SPE software can use polling or blocking when waiting for a signal to appear, or it can set up interrupts to catch signals as they appear asynchronously.

The PPE sends a signal-notification message to the SPE by writing to a MMIO register in the SPE's MFC. The signal is latched in the MMIO register and the SPU signal value by executing a read-channel (**rdch**) instruction.

One SPU can send a signal-notification message to another SPU using one of three special MFC commands: **sndsig**, **sndsigf**, and **sndsigb**. All of these commands are implemented in the same manner as a DMA **put** command, with the effective address of an MMIO register as the destination. In fact, a DMA **put** command can be used to perform exactly the same function as a send-signal command; the send-signal commands are defined in the architecture to support possible future Cell Broadband Engine Architecture implementations that optimize signal-notification performance.

### 19.7.1 SPU Signalling Channels

There are two SPU signal-notification channels, one to read each of the two signal-notification MMIO registers. A signal can be from one bit to 32 bits long. *Table 19-19* through *Table 19-21* list the signal-notification channels and associated MMIO registers. An SPU read from a signal-notification channel will be stalled when no signal is pending at the time of the read.

*Table 19-19. Signal-Notification Channels and MMIO Registers*

| SPE Channel # | Name | Channel Interface | | | | | MMIO Register Interface | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Offset From Base | Mnemonic | Max. Entries | R/W | Width (bits) |
| 3 | SPU Signal Notification 1 | SPU_RdSigNotify1 | 1 | yes | R | 32 | x'1400C' | SPU_Sig_Notify_1 | 1 | R/W | 32 |
| 4 | SPU Signal Notification 2 | SPU_RdSigNotify2 | 1 | yes | R | 32 | x'1C00C' | SPU_Sig_Notify_2 | 1 | R/W | 32 |
| — | SPU Configuration | — | — | — | — | — | x'04078' | SPU_Cfg | 1 | R/W | 64 |

*Table 19-20. Functions of Signal-Notification Channels*

| Channel Interface | SPU Read or Write | Functions |
|---|---|---|
| SPU_RdSigNotify1 | R | Returns signal-notification data. |
| SPU_RdSigNotify2 | R | Returns signal-notification data. |

*Table 19-21. Functions of Signal-Notification MMIO Registers*

| MMIO Register | PPE Read or Write | Functions |
|---|---|---|
| SPU_Sig_Notify_1 | R/W | Reads or writes signal-notification data. |
| SPU_Sig_Notify_2 | R/W | Reads or writes signal-notification data. |
| SPU_Cfg[1] | R/W | Reads or writes the configuration (either OR mode or overwrite mode) of the SPU signal notification registers. |

1. Access to this MMIO register is available only to privileged PPE software.

The two signal-notification channels each have one entry. Thus, the value returned from a read-channel-count (**rchcnt**) instruction indicates the presence or absence of an available signal value; if the channel-count value is '1', a signal is available; if the value is '0', no signal is available. When the channel count is '0' for a signal-notification channel, a read-channel (**rdch**) instruction directed at that channel will stall the SPU until a signal is available.

### 19.7.2 Uses of Signaling

Like mailboxes, signal-notification channels are useful when the SPE places computational results in main storage via DMA. After requesting the DMA transfer, the SPE waits for the DMA transfer to complete and then sends a signal to notify the PPE that its computation is complete. In this case, waiting for the DMA transfer to complete only ensures that the SPE's LS buffers are available for reuse and does not guarantee that data has been coherently written to main storage, as described in *Section 19.6.4 Uses of Mailboxes* on page 541.

### 19.7.3 Mode Configuration

Signal-notification registers can be configured independently by PPE software to operate in either *OR mode* or *overwrite mode*. The SPU Configuration (`SPU_Cfg`) MMIO register is used for this purpose.

In OR mode, the MFC combines all the values written to a signal-notification register using a logical OR operation until the SPU reads the register through the corresponding channel. When the SPU reads the register, the register is reset to '0'. OR mode allows multiple signals to be accumulated in a single signal-notification register. Each signal producer is assigned a subset of the 32 bits in the signal-notification register; this allows the signal producers to send their signals at any time and independently or other signal producers. When SPE software reads the signal-notification register, it becomes aware of all the signals that have been sent since the most-recent read of the register. This mode of signalling can be referred to as *many-to-one signalling*.

In overwrite mode, each value written to a signal-notification register overwrites the value in register. This mode of signalling can be referred to as *one-to-one signaling*. In the case of one-to-one signaling, there is typically no substantial difference in performance between signaling and using an inbound mailbox.

Because the channel count for signal-notification registers can be at most '1', SPE software cannot determine how many writes to the signal-notification register have taken place. If software needs to be made aware immediately of a write to a signal-notification register, it can enable interrupts and configure the event facility to interrupt when signals arrive (see *Section 18* on page 471).

### 19.7.4 SPU Signal Notification 1 Channel

The value read from the SPU Signal Notification 1 Channel (`SPU_RdSigNotify1`) is the 32-bit value of signal register 1. Reading this channel atomically resets to '0' any bits in the corresponding signal register that were set to '1'. A read from this channel when no signals are pending stalls the SPU until a signal is sent. If no signals are pending, a read-channel-count instruction for this channel returns '0'; if unread signals are pending, it returns one.

### 19.7.5 SPU Signal Notification 2 Channel

The value read from the SPU Signal Notification 2 Channel (`SPU_RdSigNotify2`) is the 32-bit value of signal register 2. Reading this channel atomically resets to '0' any bits in the corresponding signal register that were set to '1'. A read from this channel when no signals are pending stalls the SPU until a signal is sent. If no signals are pending, a read-channel-count instruction for this channel returns '0'; if unread signals are pending, it returns one.

### 19.7.6 Sending Signals

#### 19.7.6.1 *From the PPE*

PPE software and other system devices send a signal to an SPE by performing a 32-bit MMIO write to the effective address of the required signal-notification register. This write increments the count of the corresponding channel to one.

When the target signal-notification register is in overwrite mode, a PPE application can send multiple signals to an SPU without overrunning the signal facility by using nonzero signal values and verifying that the previous value has been read by the SPU before writing to the signal-notification register. PPE software verifies that the previous value has been read by reading from the MMIO address of the signal-notification register and ensuring that the MFC has reset the value to '0'; a value of '0' means SPE software read the previous value.

When the target signal-notification register is in OR mode, a PPE application can send multiple signals to an SPU without overrunning the signal facility by designating a unique bit field in the signal value to represent a particular signaller. Then, in a manner similar to the preceding, PPE software should verify that the bit field for the signaller is '0' before initiating another signal.

The following C-language program fragment shows how the PPE can write to a signal-notification register of a particular SPE. The code avoids overrunning the signal notification facility. Assume that the `SPU_Sig_Notify_1` and `SPU_Sig_Notify_2` macro names, representing offsets into the

problem-state region, have been defined previously. This is an example of overwrite-mode signalling, and assumes that privileged PPE software has already configured the SPU Signal Notification 1 Register in overwrite mode.

```
void *ps = get_ps();
unsigned int prev_sn_value;
unsigned int sn_value = 0xF0;

do {
    /* Do some other useful work while waiting. */
    prev_sn_value = *((volatile unsigned int *)(ps + SPU_Sig_Notify_1));
} while ( prev_sn_value != 0 );

*((volatile unsigned int *)(ps + SPU_Sig_Notify_1)) = sn_value; // Write signal value
```

#### 19.7.6.2 *From an SPE*

SPE software can send a signal to another SPE or some other device by issuing an MFC send-signal command (*Section 19.2.3* on page 518) and specifying the effective address of the target SPU signal-notification channel and a signal value. The command increments the channel count of the target SPU's signal-notification channel to one.

When the target signal-notification register is in overwrite mode, an SPU application can send multiple signals to a target SPU without overrunning the signal facility by using nonzero signal values and verifying that the previous value has been read by the target SPU before performing an MFC send-signal command. SPE software verifies that the previous value has been read by by performing an MFC **get** command from the effective address of the target SPU signal-notification register and ensuring that it has been reset to '0' by a channel read on the target SPU.

When the target signal-notification register is in OR mode, an SPU application can send multiple signals to another SPU without overrunning the signal facility by designating a unique bit field in the signal value to represent a particular signaller. Then, in a manner similar to the preceding, SPE software should ensure that the bit field for the signaller is '0' before initiating another signal.

*Send-Signal Command*

The following ABI-compliant assembler-language program fragment shows how a source SPE can write to a signal-notification register in a target SPE using an MFC **sndsig** command. SPE software first performs an MFC **get** command to confirm that the target signal-notification register is '0' before writing a new signal value. This is an overwrite-mode signalling example, and assumes that privileged software has already configured the SPU Signal Notification 1 Register in the target SPE to operate in overwrite mode.

```
# Inputs:
#   $3 contains 64-bit effective address of signal register
#   $4 contains 32-bit value to be written
#   $5 contains tag id, between [0..31]
#

    .set MFC_GET_CMD, 0x40
```

```
    .set MFC_SNDSIG_CMD, 0xA0

send_signal:
    ila  $6, 4                  # size of signal register
    ila  $7, signal_word        # LSA of signal word
    rotqbyi $8, $3, 4           # EAL of signal register

    ila  $9, 1                  # initialize tag mask
    shl  $9, $9, $5
    wrch $MFC_WrTagMask, $9

    # Read signal notification register and wait until 0
repeat:
    wrch $MFC_LSA, $7
    wrch $MFC_EAH, $3
    wrch $MFC_EAL, $8
    wrch $MFC_Size, $6
    wrch $MFC_TagID, $5
    ila  $9, MFC_GET_CMD
    wrch $MFC_Cmd, $9           # enqueue command
    ila  $9, 2
    wrch $MFC_WrTagUpdate, $9
    rdch $9, $MFC_RdTagStat     # wait for DMA get command to complete
    lqa  $9, signal_qword
    rotqbyi $9, $9, 12
    brz  $9, repeat

    # Place the requested data into the LS to be DMAed
    rotqbyi $9, $4, -12
    stqa $9, signal_qword
    dsync

    # Write the specified data to the signal notification register
    wrch $MFC_LSA, $7
    wrch $MFC_EAH, $3
    wrch $MFC_EAL, $8
    wrch $MFC_Size, $6
    wrch $MFC_TagID, $5
    ila  $9, MFC_SNDSIG_CMD
    wrch $MFC_Cmd, $9           # enqueue command

    bi    $0                    # return to caller

    # LS data buffer used for DMA operations
    .align 4
signal_qword:
    .skip 12
signal_word:
    .skip 4
```

Here is a C-language subroutine that is comparable to the preceding example.

```
#include <spu_intrinsics.h>
#include <spu_mfcio.h>

void send_signal(unsigned long long ea, unsigned int value, unsigned int tag)
{
  volatile vector unsigned int signal_qword;
  unsigned int eal, eah;

  spu_writech(MFC_WrTagMask, 1 << tag);

  /* Read signal notification register and wait until 0
   */
  do {
    mfc_get((volatile void *)(&signal_qword) + 12, ea, 4, tag, 0, 0);
    spu_mfcstat(MFC_TAG_UPDATE_ALL);
  } while (spu_extract(signal_qword, 3));

  /* Place the requested data into the LS to be DMAed
   */
  signal_qword = spu_promote(value, 3);
  spu_dsync();

  mfc_sndsig((volatile void *)(&signal_qword) + 12, ea, tag, 0, 0);
}
```

## 19.7.7 Receiving Signals

### 19.7.7.1 *Reading Signals*

SPE software uses a read-channel instruction on the signal-notification channel of interest to receive a 32-bit signal value. This read-channel instruction will return immediately, reset any set bits in the signal-notification register, and reset the channel count to ‘0’ if the associated signal-notification register has a waiting unread signal value. Otherwise, the read-channel instruction will cause the SPU to stall until a write to the signal-notification register happens.

The following assembler-language program fragment shows how to read an SPU signal.

```
# Outputs:
#        $3 contains the signal control value read

rdch     $3, $SPU_RdSigNotify1
```

Here is a C-language program fragment that is comparable to the preceding example.

```
#include <spu_intrinsics.h>
unsigned int sn_value;

sn_value = spu_readch(SPU_RdSigNotify1); // Read SPU Signal Notification 1
```

#### 19.7.7.2 *Waiting for Signals*

To avoid an SPU stall, SPE software can use the read-channel-count instruction on the signal channel of interest to determine if a signal has been sent. If the read-channel-count instruction returns '0', no signal is pending. If the read-channel-count instruction returns a nonzero value, an unread signal is waiting, and the SPU can read the signal channel without stalling.

*Poll For Signals*

The following assembler-language program fragment shows how an SPU can poll one of its signal-notification channels.

```
# Outputs:
#         $3 contains the signal control word 1

repeat:
   rchcnt $3, $SPU_RdSigNotify1
   brz    $3, repeat

rdch      $3, $SPU_RdSigNotify1
```

Here is a C-language program fragment that is comparable to the preceding example.

```
#include <spu_intrinsics.h>
unsigned int sn_value;

do {
    /* Do other useful work while waiting. */
} while (!spu_readchcnt(SPU_RdSigNotify1));

sn_value = spu_readch(SPU_RdSigNotify1); // Read the SPU Signal Notification 1 Channel
```

*Poll or Block for Signal-Notification Events*

The SPU can also avoid a stall by enabling the SPU event facility, described in *Section 18* on page 471. Using events involves the following steps:

1. Enabling an event by setting the Signal Notification 1 Available event bit or the Signal Notification 2 Available event bit in the SPU Write Event Mask Channel.

2. Monitor for this event. Software can poll with a read-channel-count instruction on SPU Read Event Status; software can block (stall) with a read-channel instruction on the SPU Read Event Status Channel; software can enable SPU interrupts and use interrupt handlers to service signal-notification events as they are encountered asynchronously.
3. When the event is detected, it must be acknowledged by writing the corresponding bit to the SPU Write Event Acknowledgment Channel. A typical handler for this event performs the acknowledgement and also reads from the appropriate SPU Signal Notification Channel.

The following assembler-language program fragment shows how to poll for or block on signal-notification events. To block instead of poll, simply delete the repeat loop.

```
# Outputs:
#   $3 contains the signal notification register contents
#   $4 contains the event that occurred

receive_signal:
    ila    $3, 0x200
    wrch $SPU_WrEventMask, $3       # Write event mask

repeat:           # Wait for event to occur
    rchcnt $3, $SPU_RdEventStat
    brz    $3, repeat

    rdch   $4, $SPU_RdEventStat     # Read the events that occurred
    wrch   $SPU_WrEventStat, $4     # Ack the events received
    rdch   $3, $SPU_RdSigNotify1    # Read the signal notification register
```

Here is a C-language program fragment that is comparable to the preceding example. To block, instead of poll, simply delete the following do/while loop.

```
#include <spu_intrinsics.h>
#include <spu_mfcio.h>

unsigned int event_status;
unsigned int sn_value;

spu_writech(SPU_WrEventMask, MFC_SIGNAL_NOTIFY_1_EVENT);  // Enable event

do {
    /* Do other useful work while waiting. */
} while (!spu_readchcnt(SPU_RdEventStat));

event_status = spu_readch(SPU_RdEventStat);              // Read event status

spu_writech(SPU_WrEventAck, event_status);               // Ack. event

sn_value = spu_readch(SPU_RdSigNotify1);                  // Read the signal value
```

### 19.7.8 Differences Between Mailboxes and Signal Notification

For a comparison of the differences between the mailbox and signal-notification facilities, see *Table 19-2* on page 514.

# 20. Shared-Storage Synchronization

This section treats the subject of synchronization in two aspects. The first involves storage-access and instruction-execution ordering for shared-storage environments. The second involves the use of atomic operations to create semaphores, mutex locks, and other synchronization primitives for shared-storage environments.

Synchronization should be done with care because it affects performance. For example, one key to successful parallelization strategies is the minimizing of synchronizing events between the computational elements.

This section describes only nonprivileged (user) synchronization facilities and scenarios. Some operations on privileged registers and memory-management facilities also require synchronization. For details, see *Section 4 Virtual Storage Environment* on page 79, *Section 9 PPE Interrupts* on page 239, *Section 11 Logical Partitions and a Hypervisor* on page 331, and *PowerPC Operating Environment Architecture, Book III*.

## 20.1 Shared-Storage Ordering

This section describes storage-access and instruction-execution ordering for shared-storage environments. The Cell Broadband Engine Architecture (CBEA) processors[1] contain multiple shared-storage domains—the main-storage domain, eight local storage (LS) address domains, and eight local channel domains—as illustrated in *Figure 1-2* on page 47.

### 20.1.1 Storage Model

Although the CBEA processors execute instructions in program order, they load and store data using a *weakly consistent* storage model. That is, the order in which any processor element (PowerPC Processor Element [PPE] or Synergistic Processor Element [SPE]) performs storage accesses, the order in which those accesses are performed with respect to another processor element or mechanism, and the order in which those accesses are performed in main storage might be different. This storage model allows storage accesses to be reordered dynamically, which provides an opportunity for improved overall performance and reduced effect of memory latency on instruction throughput. Therefore, in a memory environment that is shared by other threads of execution or I/O devices, this weakly consistent model requires that programs explicitly order accesses to storage if they want stores to occur in the program order.

To ensure that accesses to shared storage (main storage, which includes external access to the LSs of SPEs) are performed in program order, software must place memory-barrier instructions between storage accesses, as required by the *Cell Broadband Engine Architecture* (CBEA). Programs that do not contain such memory-barrier instructions in appropriate places might be executed correctly for a given implementation or execution environment but will fail if the execution environment is changed or is executed on another implementation of the CBEA.

The term *storage access* means an access to main storage caused by a load, a store, a direct memory access (DMA) read, or a DMA write. There are two orders to consider:

- *Order in which Instructions are executed*—Instructions can be executed in-order or out-of-order. For in-order machines, such as the CBEA processors, each instruction completes

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

before the next instruction starts. In general, from the viewpoint of a program, it appears that the instructions are executed in the order specified by the program. An out-of-order machine does not guarantee that the processor element executes instructions in program order. A program cannot detect that a sequence of instructions is executed in an order different than specified in the program.

- *Order in which storage accesses are performed*—The order in which shared-storage accesses are performed might be different from both program order and the order in which the instructions that caused the accesses are executed. Consider a program that contains a sequence of loads from locations A, B, and C, in that order. The processor element might execute these instructions in a different order (for example, B, A, C) and perform the accesses caused by the instructions in yet a different order (for example, C, B, A).

The CBEA defines four storage-control attributes that are maintained by the operating system in the page-table entries (see *Section 4.2.6 Paging* on page 87):

- *Write-Through Required (W attribute)*—When an access is designated as write-through, if data is in a cache (such as one of the PPE's L1 or L2 caches), a store operation updates both the cache copy of the data and the main-storage location of the data.
- *Caching-Inhibited (I attribute)*—When an access is designated as caching-inhibited, the access is completed by referencing only the location in main storage, bypassing the cache.
- *Memory-Coherence Required (M attribute)*—When an access is designated as memory-coherent, a hardware indication is sent to the rest of the system indicating that the access is global. Other processor elements affected by the access must acknowledge and respond appropriately to the access.
- *Guarded (G attribute)*—Memory is marked as guarded when areas of main storage must be protected from accesses not directly dictated by the program. For example, the guarded attribute can be used to prevent out-of-order, speculative load, or prefetch operations from occurring with peripheral devices; such accesses can produce undesirable results.

Main storage has the memory-coherence required attribute. SPE LSs and I/O devices that are mapped to the effective-address space typically have the cache-inhibited and guarded attributes.

Several means of enforcing an ordering of storage accesses are provided to allow programs to share storage with other programs or threads of execution, or with other mechanisms such as I/O devices. The appropriate mechanism depends upon the subsystem wanting to enforce storage ordering and the storage-control attributes of the storage being accessed.

*Table 20-1* on page 563 summarizes the effects on address and communication domains of the synchronization instructions, commands, and facilities that are described in the sections immediately following this table. Gray shading in a table cell means that the instruction, command, or facility has no effect on the referenced domain.

*Table 20-1. Effects of Synchronization on Address and Communication Domains[1]*

| Issuer | Instruction, Command, or Facility | Main-Storage Domain | | | LS Domain[2] | | | Channel Domain[3] |
|---|---|---|---|---|---|---|---|---|
| | | Accesses by PPE | | Accesses by All Other Processor Elements and Devices | Accesses by Issuing SPU | Accesses by Issuing SPU's MFC | Accesses by All Other Processor Elements and Devices | Accesses by Issuing SPU |
| | | Issuing Thread | Both Threads | | | | | |
| PPU | **sync**[4] | all accesses | | | | Unreliable. Use MFC Multisource Synchronization Facility[5] | | |
| | **lwsync**[6] | accesses to memory-coherence-required locations | | | | | | |
| | **eieio** | accesses to caching-inhibited and guarded locations | | accesses to caching-inhibited and guarded locations | | Unreliable. Use MFC Multisource Synchronization Facility[5] | | |
| | **isync** | instruction fetches | | | | | | |
| SPU | **sync** | | | | all accesses | | | all accesses |
| | **dsync** | | | | load and store accesses | all accesses | | |
| | **syncc** | | | | all accesses | | | |
| MFC | **mfcsync** | | | all accesses | | Unreliable. Use MFC Multisource Synchronization Facility[5] | | |
| | **mfceieio** | | | accesses to caching-inhibited and guarded locations | | | | |
| | **barrier** | | | all accesses | | | | |
| | **<f>, <b>** | | | all accesses for the tag group | | | | |
| | MFC Multisource Synchronization Facility | | | all accesses | all accesses | | | |

1. Gray shading in a table cell means that the instruction, command, or facility has no effect on the referenced domain.
2. The LS of the issuing SPE.
3. The channels of the issuing SPE.
4. This is the PowerPC **sync** instruction with L = '0'.
5. These accesses can exist only if the LS is mapped by the PPE operating system to the main-storage space. This can only be done if the LS is assigned caching-inhibited and guarded attributes.
6. This is the PowerPC **sync** instruction with L = '1'.

### 20.1.2 PPE Ordering Instructions

The PPE supports *barrier instructions* for ordering storage accesses and instruction execution. The *PowerPC Architecture* defines two types of barrier instructions:

- Storage Barriers—the **sync**, **lwsync**, **ptesync**, and **eieio** instructions
- Instruction Barrier—the **isync** instruction

These instructions can be used between storage-access instructions to define a memory barrier that divides the instructions into those that precede the barrier instruction and those that follow it. The storage accesses caused by instructions that precede the barrier instruction are performed before the storage accesses caused by instructions that follow the barrier instruction.

The PPE supports one additional synchronization instruction, **tlbsync**, that is privileged. For details, see *Section 4 Virtual Storage Environment* on page 79.

#### 20.1.2.1 *The PPE sync Instruction*

The PPE **sync** instruction—which differs from the synergistic processor unit (SPU) **sync** instruction (*Section 20.1.3.1* on page 569)—ensures that all instructions preceding the **sync** appear to have completed before the **sync** instruction completes, and that no subsequent instructions are initiated until after the **sync** instruction completes. This does not mean that the previous storage accesses have completed before the **sync** instruction completes.

The **sync** instruction is also known as the *heavyweight* **sync** instruction, as opposed to the lightweight **lwsync** instruction and the **ptesync** instruction. The opcode of the **sync** instruction has its L bit cleared to ‘0’, whereas the **lwsync** instruction is the **sync** instruction with L = ‘1’ and the **ptesync** instruction is the **sync** instruction with L = ‘2’. See the *PowerPC Operating Environment Architecture, Book III* for details.

The **sync** instruction orders all main-storage accesses, regardless of the main-storage memory attributes. For example, it orders a caching-inhibited load with respect to a caching-enabled store. However, only processor elements operating in the same logical partition (see *Section 4 Virtual Storage Environment* on page 79) will accept the **sync**; processor elements in other logical partitions ignore it.

In addition, the **sync** instruction does the following:

- Prevents store-combining to storage that is caching-inhibited or write-through-required.
- Prevents load-combining for storage that is caching-inhibited.
- Is cumulative—that is, all processor elements in the system will see the same barrier sequence that is performed and seen by the processor element that executes the **sync** instruction (the PPE).

The **sync** instruction can be used, for example, to ensure that the results of all stores into a data structure, caused by store instructions executed in a critical section of a program, are seen by other processor elements before the data structure is seen as unlocked.

The functions performed by the **sync** instruction normally take a significant amount of time to complete, and the time required to execute **sync** can vary from one execution to another, so this instruction should be used with care. The **eieio** instruction might be more appropriate than **sync** for many cases (see *Table 20-3* on page 567 and *Table 20-4* on page 568).

#### 20.1.2.2 *The lwsync Instruction*

The **lwsync** instruction (also known as the *lightweight* **sync** instruction) creates the same barrier as the **sync** instruction for storage accesses that have the memory-coherence-required attribute. However, the **lwsync** instruction does not create a barrier for storage accesses that have either the caching-inhibited or write-through-required attribute.

Unlike the **sync** instruction, the **lwsync** instruction orders only the main-storage accesses of the PPE. It has no effect on the main-storage accesses of other processor elements in the system. Thus, the effects of the **lwsync** instruction are not cumulative.

The **lwsync** instruction should be used when ordering is required only for coherent memory, because **lwsync** executes faster than **sync**.

#### 20.1.2.3 *The eieio Instruction*

The **eieio** instruction (enforce in-order execution of I/O) ensures that all main-storage accesses caused by instructions proceeding the **eieio** have completed, with respect to main storage, before any main-storage accesses caused by instructions following the **eieio**.

The **eieio** instruction orders loads and stores in the following two sets, which are ordered separately in the following sequence:

1. Loads and stores to storage that is both caching-inhibited and guarded, and stores to storage that is write-through-required. All accesses in this set are ordered as a single set—that is, there is not one order for loads and stores to caching-inhibited and guarded memory and another order for stores to write-through-required memory.
2. Stores to storage that is memory-coherence-required but that is neither caching-inhibited nor write-through-required.

In addition, the **eieio** instruction does the following:

- Prevents store-combining to storage that is both caching-inhibited and guarded, or to storage that is write-through-required
- Prevents load-combining for storage that is caching-inhibited
- Is cumulative (see *Section 20.1.2.1 The PPE sync Instruction* on page 564) for storage that is neither caching-inhibited nor write-through-required

The *PowerPC Architecture* specifies that stores to storage that is both caching-inhibited and guarded are performed in program order, so **eieio** is needed for such storage only when loads must be ordered with respect to stores or with respect to other loads, or when load and store combining operations must be prevented.

An **eieio** instruction issued on one PPE thread has no architected effect on the other PPE thread. However, there is a performance effect because all load or store operations (cacheable or noncacheable) on the other thread are serialized behind the **eieio** instruction.

The **eieio** instruction does not order accesses with differing storage attributes. For example, if an **eieio** is placed between a caching-enabled store and a caching-inhibited store, the access might still be performed in an order different than specified by the program.

The **eieio** instruction is intended for use in managing shared data structures, in accessing memory-mapped I/O, and in preventing load or store combining. For the first use, the shared data structure and the lock that protects it must be altered only by stores that are in the same **eieio** ordering set as described previously. For the second use, **eieio** can be thought of as placing a barrier into the stream of memory accesses issued by a processor element, such that any given memory access appears to be on the same side of the barrier to both the processor element and the I/O device.

#### 20.1.2.4 *The isync Instruction*

The **isync** instruction ensures that all PPE instructions proceeding the **isync** are completed before **isync** is completed. An **isync** (L = 0 or 1) instruction causes issue stall and blocks all other instructions from both PPE threads until the **isync** instruction completes.

Prefetched instructions on the issuing thread are discarded; those on the other PPE thread are unaffected. The **isync** instruction only flushes based on thread. The instructions from the other thread which are already fetched or dispatched are not refetched after the **isync** completes.

This instruction is most often used in conjunction with self-modifying PPE code. After an instruction is modified, an **isync** instruction must be issued before execution of the modified instruction. The **isync** instruction might also be used during context-switching when the memory management unit (MMU) translation rules are being changed.

The **isync** instruction is *context synchronizing*, as opposed to the **sync** instruction which is *execution synchronizing*. See the *Glossary* on page 835 and the Synchronization section of *PowerPC Operating Environment Architecture, Book III* for a further information.

#### 20.1.2.5 *Summary of Storage-Barrier Instruction Use*

*Table 20-2* on page 567 summarizes the functions of the storage-barrier instructions with respect to storage attributes. In this table, “yes” (and “no”) mean that the instruction performs (or does not perform) a barrier function on storage with the related attribute.

*Table 20-2. Storage-Barrier Function Summary*

| Storage Attribute or Function | Loads or Stores | **sync** | **lwsync** | **eieio** |
|---|---|---|---|---|
| memory-coherence-required | loads | yes | yes | no |
| | stores | | | yes, if neither caching-inhibited nor write-through-required[1] |
| caching-inhibited | loads and stores | | no | yes, if both caching-inhibited and guarded[2] |
| guarded | loads and stores | | | |
| write-through-required | stores | | | yes[2] |
| | loads | not applicable | | |
| combining prevented | loads | for caching-inhibited | | |
| | stores | for (a) caching-inhibited, and for (b) write-through-required | | for (a) caching-inhibited and guarded, and for (b) write-through-required |
| cumulative[3] | loads and stores | yes | no | yes, if neither caching-inhibited nor write-through-required |

1. Second ordering: stores to storage that is memory-coherence-required but that is neither caching-inhibited nor write-through-required.
2. First ordering: loads and stores to storage that is both caching-inhibited and guarded, and stores to storage that is write-through-required
3. All processor elements in the system will see the same barrier sequence that is performed and seen by the processor element that executes the instruction.

*Table 20-3* and *Table 20-4* summarize the use of the storage-barrier instructions for two common types of main-storage memory:

- *System Memory*—Memory that has the memory-coherence-required attribute.
- *Device Memory*—Memory that has caching-inhibited and guarded attributes, which is typical of memory-mapped I/O devices. (The mapping of SPEs' LS to main storage is generally caching-inhibited but not guarded.)

In these tables, "yes" (and "no") mean that the instruction performs (or does not perform) a barrier function on the related storage sequence, "recommended" means that the instruction is the preferred one, "not recommended" means that the instruction will work but is not the preferred one, and "not required" and "no effect" mean the instruction has no effect.

*Table 20-3. Storage-Barrier Ordering of Accesses to System Memory*

| Storage-Access Instruction Sequence | **sync** | **lwsync** | **eieio** |
|---|---|---|---|
| load-*barrier*-load | yes | recommended | no affect |
| load-*barrier*-store | yes | recommended | no affect |
| store-*barrier*-load | yes | no | no affect |
| store-*barrier*-store | yes | recommended | not recommended |

*Table 20-4. Storage-Barrier Ordering of Accesses to Device Memory*

| Storage-Access Instruction Sequence | **sync** | **lwsync** | **eieio** |
|---|---|---|---|
| load-*barrier*-load | yes | no affect | yes |
| load-*barrier*-store | yes | no affect | yes |
| store-*barrier*-load | yes | no affect | yes |
| store-*barrier*-store | not required [1] | no affect | not required [1] |

1. Two stores to caching-inhibited storage are performed in the order specified by the program, regardless of whether they are separated by a barrier instruction or not.

The following rules apply to ordering operations for cacheable and caching-inhibited locations:

- The ordering of cacheable loads is not guaranteed. If ordering is required, software needs to issue a **sync** or **lwsync** instruction.
- Caching-inhibited loads are implemented as guarded operations. This means that the PPE does not issue a new load until the data for the older load is returned. There can only be one outstanding caching-inhibited load at any time.
- Caching-inhibited instruction fetches are treated as nonguarded operations. The PPE does not maintain ordering for caching-inhibited, nonguarded operations.

### 20.1.3 SPU Ordering Instructions

An LS can experience asynchronous interaction from the following streams that access it:

- Instruction fetches by the local SPU
- Data loads and stores by the local SPU
- DMA transfers by the local memory flow controller (MFC) or another SPE's MFC
- Loads and stores in the main-storage space by other processor elements or devices

With regard to an SPU, the CBEA processors' in-order execution model[2] guarantees only that SPU instructions that access that SPU's LS appear to be performed in program order with respect to that SPU. These accesses might not appear to be performed in program order with respect to external accesses to that LS or with respect to the SPU's instruction fetch. Thus, in the absence of external LS writes, an SPU load from an address in its LS returns the data written by that SPU's most-recent store to that address. However, an instruction fetch from that address does not necessarily return that data. The SPU may buffer and otherwise reorder its LS accesses. Instruction fetches, loads, and stores can access the LS in any order. The SPU can speculatively read the LS. However, the SPU does not speculatively write the LS; the SPU only writes the LS on behalf of stores required by the program.

The SPU Instruction Set Architecture (ISA) provides three synchronization instructions: synchronize (**sync**), synchronize data (**dsync**), and synchronize channel (**syncc**). These instructions have both coherency and instruction-serializing effects.

2. The in-order execution model is implementation-specific, rather than part of the *Cell Broadband Engine Architecture*.

In the descriptions that follow, the following terms are used:

- "External Read" and "External Write" mean accesses by any DMA transfer (by the local MFC or the MFC associated with another SPE) or any processor element (including the PPE) or other device—other than the SPU that executes the synchronization instruction.
- "SPU Load" and "SPU Store" mean accesses by the SPU that executes the synchronization instruction.

#### 20.1.3.1 *The SPU sync Instruction*

The SPU **sync** instruction—which differs from the PPE **sync** instruction (*Section 20.1.2.1* on page 564)—causes the SPU to wait until all pending store instructions to its LS have completed before fetching the next sequential instruction. It orders instruction fetches, loads, stores, and channel accesses.

After an SPU has executed a **sync** instruction, the SPU's instruction fetches from an LS address return data stored by the most-recent store instruction or external write to that address. In addition, the SPU will have completed all prior loads, stores, and channel accesses and will not have begun execution of any subsequent loads, stores, or channel accesses. At this time, an external read from an LS address returns the data stored by the most-recent SPU store to that address. SPU loads after the **sync** return the data externally written before the moment when the **sync** completes.

The **sync** instruction affects only that SPU's instruction sequence and the coherency of that SPU's fetches, loads, and stores, with respect to actual LS state. The SPU does not broadcast **sync** notification to external devices that access its LS, and therefore the **sync** does not affect the state of external devices.

The SPU's instruction-fetch buffers and pipelines are flushed when it executes the **sync** instruction. The **sync** instruction has no affect on, and does not wait for, DMA **get** operations controlled by the SPU's MFC.

The **sync** instruction must be used before attempting to execute new code that either arrives through DMA transfers or is written with store instructions. This instruction is most often used in conjunction with self-modifying SPE code.

The CBEA processors invalidate the instruction prefetch buffer whenever a mispredicted branch is executed. In such a case, the **sync** instruction can be omitted, but for architecture portability the **sync** instruction should still be used.

#### 20.1.3.2 *The dsync Instruction*

The **dsync** (data synchronization) instruction allows SPE software to ensure that data has been stored in the LS before the data becomes visible to the local MFC or other external devices. It orders loads, stores, and channel accesses but not instruction fetches, and it does not affect any prefetching of instructions that the SPE might have done.

After a **dsync** completes, the SPU will have completed all prior loads, stores, and channel accesses and will not have begun execution of any subsequent loads, stores, or channel accesses. At this time, an external read from an LS address returns the data stored by the most-recent SPU store to that address. SPU loads after the **dsync** return the data written before the moment when the **dsync** completes.

The **dsync** instruction affects only that SPU's instruction sequencing and the coherency of that SPU's loads and stores, with respect to actual LS state. The SPU does not broadcast **dsync** notification to external devices that access its LS, and therefore the **dsync** does not affect the state of the external devices.

DMA transfers can interfere with store instructions and the store buffer. Architecturally, the **dsync** instruction is used to ensure that all store buffers have been flushed to the LS, so that all previous stores to the LS will be seen by subsequent LS accesses. However, the CBEA processors do not require the **dsync** instruction for this purpose.

#### 20.1.3.3 *The syncc Instruction*

The **syncc** (channel synchronization) instruction performs channel synchronization followed by the same synchronization provided by the **sync** instruction. It forces all channel instructions to complete before executing the next SPU instruction. It ensures that the effects on SPU state caused by prior **wrch** instructions are propagated and influence the execution of the following instructions.

Data through a particular channel is never reordered. A channel is either blocking or nonblocking. If a channel is blocking, read data must be present or previous write data must have been consumed before the instruction that accesses the channel can complete. Thus, the hardware synchronizes the channel streams. However, if a channel is nonblocking, then software must manage the synchronization based on the ordering of channel data.

Some channel activity might, as a side effect, alter the SPU operating state. These side effects are not ordered with respect to the SPU pipeline. Such side effects can be synchronized with SPU-instruction execution using the **syncc** instruction. Channel synchronization does not ensure that DMA commands generated as a result of issuing a channel command have completed; it ensures only that such DMA commands have been placed in the MFC command queue.

See *Section 17* on page 447 for details about channels.

#### 20.1.3.4 *Summary of SPU Ordering Instructions*

The SPU synchronization instructions are summarized in *Table 20-5* on page 571. *Table 20-6* on page 571 shows which SPU synchronization instructions are required between LS writes and LS reads to ensure that reads access data written by prior writes.

*Table 20-5. SPU Synchronization Instructions*

| Instruction | Coherency Effects | Instruction Serialization Effects |
|---|---|---|
| **sync** | Ensures that subsequent external[1] reads access data written by prior SPU stores.<br>Ensures that subsequent instruction fetches access data written by prior SPU stores and external[1] writes.<br>Ensures that subsequent SPU loads access data written by external[1] writes. | Forces all access of LS and channels due to instructions before the **sync** to be completed before completion of **sync**.<br>Forces all access of LS and channels due to instructions after the **sync** to occur after completion of the **sync**. |
| **dsync** | Ensures that subsequent external[1] reads access data written by prior SPU stores.<br>Ensures that subsequent SPU loads access data written by external[1] writes. | Forces SPU load and SPU store access of LS due to instructions before the **dsync** to be completed before completion of **dsync**.<br>Forces read channel operations due to instructions before the **dsync** to be completed before completion of the **dsync**.<br>Forces SPU load and SPU store access of LS due to instructions after the **dsync** to occur after completion of the **dsync**.<br>Forces read-channel and write-channel operations due to instructions after the **dsync** to occur after completion of the **dsync**. |
| **syncc** | Ensures that subsequent external[1] reads access data written by prior SPU stores.<br>Ensures that subsequent instruction fetches access data written by prior SPU stores and external[1] writes.<br>Ensures that subsequent SPU loads access data written by external[1] writes.<br>Ensures that subsequent instruction processing is influenced by all internal execution states modified by previous **wrch** instructions. | Forces all access of LS and channels due to instructions before the **syncc** to be completed before completion of **syncc**.<br>Forces all access of LS and channels due to instructions after the **syncc** to occur after completion of the **syncc**. |

1. By any DMA transfer (from the local MFC or a nonlocal MFC), the PPE, or other device—other than the SPU that executes the synchronization instruction.

*Table 20-6. Synchronization Instructions for Accesses to an LS*

| Writer | Reader | | |
|---|---|---|---|
| | SPU Instruction Fetch | SPU Load | External[1] Read |
| SPU Store | **sync** | nothing required | **dsync** |
| External[1] Write | **sync** | **dsync** | not applicable |

1. By any DMA transfer (from the local MFC or a nonlocal MFC), the PPE, or other device—other than the SPU that executes the synchronization instruction.

### 20.1.3.5 *Ordering and Synchronization of External Accesses to the MFC and LS*

Privileged software on the PPE can make the SPE's LS and MFC resources accessible to the PPE or other devices in the main-storage space. If two accesses are made to two different addresses, no ordering is maintained between the accesses unless the PPE **eieio** or **sync** instruction is used between the two accesses. (The **eieio** instruction can be used because LS mappings to main storage have the caching-inhibited attribute.)

### 20.1.4 MFC Ordering Mechanisms

The MFC of each SPE supports synchronization commands and command options that control the order in which DMA storage accesses are performed by that MFC. These synchronization commands are of two kinds:

- *Barrier Commands*—The **mfcsync**, **mfceieio**, and **barrier** commands order storage accesses made through the MFC with respect to all other MFCs, processor elements, and other devices in the system. The CBEA specifies that the **mfcsync** and **mfceieio** commands are tag-specific, but the Cell/B.E. and PowerXCell 8i processors treat all three barrier command identically, having no tag-specific effects.
- *Fence or Barrier Command Options*—These are options of the **get**, **put**, and **sndsig** commands. The options order storage accesses and signaling only for a specific, local tag group and MFC command queue.

For an overview of MFC commands, command queues, and tag groups, see *Section 19.2* on page 514.

#### 20.1.4.1 *The mfcsync Command*

The **mfcsync** command is similar in operation to the PPE **sync** instruction, described in *Section 20.1.2.1* on page 564. The CBEA specifies that the **mfcsync** command controls the order in which MFC commands within a specified tag group for that MFC are executed with respect to storage accesses by all other processor elements and devices in the system. The **mfcsync** command and all subsequent commands in the same tag group are ordered with respect to all previous commands in the queue with the same tag ID.

Although the CBEA specifies the **mfcsync** command as having tag-specific effects, the Cell/B.E. and PowerXCell 8i processors treat this command as not having tag-specific effects and as behaving the same as the **barrier** command, described in *Section 20.1.4.3* on page 573.

The C-intrinsic syntax for an SPE issuing an **mfcsync** command with the `tag_id` tag group is:

```
spu_writech(MFC_TagID, tag_id);
spu_writech(MFC_Cmd, 0xCC);
```

To obtain the best performance across the widest range of implementations, use the **mfcsync** command only when its entire functionality is needed. Use the **mfceieio** command, the **barrier** command, or fence **<f>** or barrier **<b>** options of other commands, if any of these are sufficient for the ordering needs.

If the storage accesses between two DMA operations that access storage having different storage attributes need to be strongly ordered, the fence **<f>** or barrier **<b>** options of the **get** and **put** commands can do this. However, these options are insufficient to guarantee the ordering function with respect to other processor elements or devices. In these cases, an **mfcsync** command must be issued between commands involving storage with different storage attributes to provide the required ordering function.

#### 20.1.4.2 *The mfceieio Command*

The **mfceieio** command is similar in operation to the PPE **eieio** instruction, described in *Section 20.1.2.3* on page 565. The CBEA specifies that the **mfceieio** command controls the order in which DMA commands within a specified tag group for that MFC are executed with respect to storage accesses by all other processor elements and devices in the system, when the storage being accessed has the attributes of caching-inhibited and guarded (typical for I/O devices).

Although the CBEA specifies the **mfceieio** command as having tag-specific effects, the Cell/B.E. and PowerXCell 8i processors treat this command as not having tag-specific effects and as behaving the same as the **barrier** command, described in *Section 20.1.4.3*.

The **mfceieio** command only orders:

- **get** commands with respect to other **get** commands from storage that is caching-inhibited and guarded
- **get** commands with respect to **put** commands accessing storage that is caching-inhibited and guarded
- **put** commands with respect to **put** commands accessing storage that is memory-coherence-required and is not caching-inhibited

The C-intrinsic syntax for issuing an **mfceieio** command with the `tag_id` tag group is:

```
spu_writech(MFC_TagID, tag_id);
spu_writech(MFC_Cmd, 0xC8);
```

The **mfceieio** command is intended for use in managing shared data structures, in performing memory-mapped I/O, and in preventing load and store combining in main storage. To obtain the best performance across the widest range of implementations, use the fence **<f>** or barrier **<b>** options of other commands, if any of these are sufficient for the ordering needs.

#### 20.1.4.3 *The barrier Command*

The **barrier** command orders all subsequent MFC commands with respect to all MFC commands preceding the **barrier** command in the DMA command queue, independent of tag groups. The **barrier** command will not complete until all preceding commands in the queue have completed. After the command completes, subsequent commands in the queue may be started.

The C syntax for issuing a **barrier** command is:

```
spu_writech(MFC_TagID, tag_id);
spu_writech(MFC_Cmd, 0xC0);
```

The setting of the tag ID is optional and need only be provided to assign an ID to the **barrier** command for subsequent tag-completion testing.

In contrast to the **barrier** command, which operates independently of tag-groups, the fence **<f>** or barrier **<b>** options of **get** and **put** commands form of fence or barrier that only affects commands within the same tag group.

#### 20.1.4.4 *Tag-Specific Ordering Commands*

Local ordering of the MFC commands within a specific tag group can be achieved using the fence **<f>** or barrier **<b>** options of the **get**, **put**, or **sndsig** commands. Local ordering only ensures ordering of the MFC commands with respect to that particular MFC command queue and tag group (tag-groups cannot span multiple MFCs).

The storage system may complete requests in an order different then the order in which they are issued, depending on the storage attributes. However, caching-inhibited storage accesses are ordered in the order in which they are issued. *Table 20-7* lists the tag-specific ordering commands.

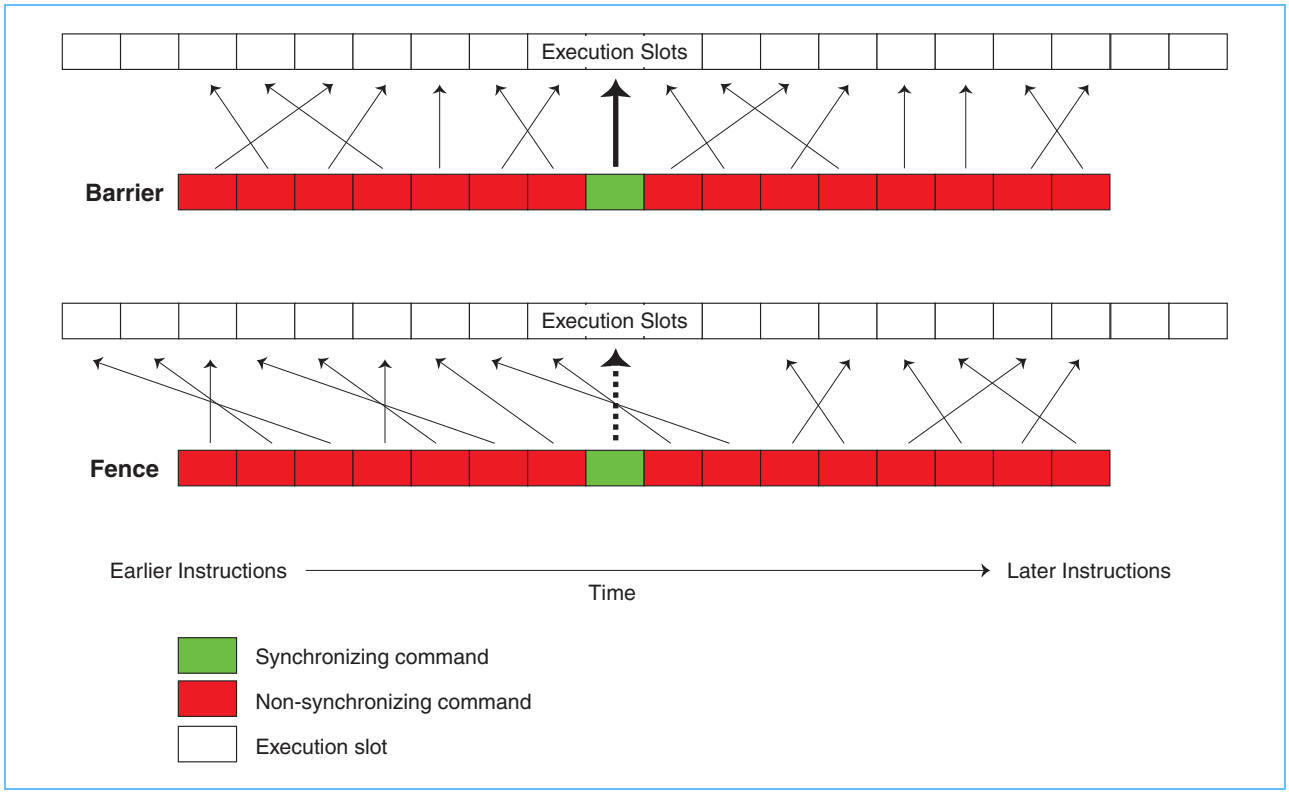
*Table 20-7. Tag-Specific Ordering Commands*

| Option | Commands |
|---|---|
| barrier | **getb, getbs, getlb, putb, putbs, putrb, putlb, putrlb, sndsigb** |
| fence | **getf, getfs, getlf, putf, putfs, putrf, putlf, putrlf, sndsigf** |

The fence form ensures that all previous MFC commands of the same tag group are issued prior issuing the fenced command. The barrier form ensures that all previous MFC commands in the same tag group are issued before issuing the barrier command, and that none of the subsequent MFC commands in the same tag group are issued before the barrier command.

*Figure 20-1* on page 575 illustrates the behavior of barriers and fences. In this example, the row of white boxes represents command-execution slots, in real-time, in which the DMA commands (red and green boxes) might execute. Each DMA command is assumed to transfer the same amount of data, thus, all boxes are the same size. The arrows show how the DMA hardware, using out-of-order execution, might execute the DMA commands over time. So, a barrier does not allow the DMA hardware to reach across the barrier to find eligible commands to execute. A fence, in contrast, allows the DMA hardware to reach across the fence in one direction to find eligible DMA commands to run.

*Figure 20-1. Barriers and Fences*




*Storage-Access Consistency for DMA-Command Combinations*

The fence and barrier options provide stronger storage-access consistency for some combinations of DMA commands:

- A **put(l)** command followed by a **put(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are memory-coherence-required and neither write-through-required nor caching-inhibited, has a storage-order effect at least as strong as two stores on the PPE separated by a **lwsync** instruction.
- A **get(l)** command followed by a **get(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are memory-coherence-required and neither write-through-required nor caching-inhibited, has a storage-order effect at least as strong as two loads on the PPE separated by a **lwsync** instruction.
- A **put(l)** command followed by a **put(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attribute is caching-inhibited, has a storage-order effect at least as strong as two stores on the PPE separated by a **sync** instruction.
- A **get(l)** command followed by a **get(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are caching-inhibited and not guarded, has a storage-order effect at least as strong as two loads on the PPE separated by a **sync** instruction.
- A **get(l)** command followed by a **get(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are caching-inhibited and guarded, has a storage-order effect at least as strong as two loads on the PPE separated by an **eieio** instruction.

- A **put(l)** command followed by a **get(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are caching-inhibited and not guarded, has a storage-order effect at least as strong as a store followed by a load on the PPE separated by a **sync** instruction.
- A **get(l)** command followed by a **put(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are memory-coherence-required and neither write-through-required nor caching-inhibited, has a storage-order effect at least as strong as a load followed by a store on the PPE separated by a **lwsync** instruction.
- A **get(l)** command followed by a **put(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are caching-inhibited and not guarded, has a storage-order effect at least as strong as a load followed by a store on the PPE separated by a **sync** instruction.
- A **get(l)** command followed by a **put(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are caching-inhibited and guarded, has a storage-order effect at least as strong as a load followed by a store on the PPE separated by an **eieio** instruction.
- A **put(l)** command followed by a **put(l)** command with either the fence **<f>** or barrier **<b>** option, where the storage attributes are caching-inhibited and guarded, are always performed in program order with respect to any processor element or mechanism.

*Common Fenced-Option Use*

The most common use of the fenced DMA is when writing back notifications. Consider an example in which an SPE computes some data, writes the computed data back to main storage, then writes a notification that the data is available. This notification might be to any type memory (main memory, I/O memory, a memory-mapped I/O [MMIO] register, a signal-notification register, or another SPE's mailbox). To ensure ordering of the DMA write and the notification, the notification can be sent using a fenced DMA command, so that the notification is not sent until all previous DMA commands of the group are issued.

A fence option might also be useful when multiple DMA commands are needed to load an SPU program and to start its execution. In this example, one DMA command is used to load the initialized data segment and a second DMA command with both the SPU start ("**s**") suffix (causing execution to start after the DMA transfer completes) and the tag-specific fence ("**f**") suffix is used to load the code (the "text" segment). If the two commands have the same tag ID, the fence ensures that the load of the data segment is completed before loading the text data and before starting the SPU program execution. Without the fence, the second DMA command might complete and might start the SPU program before the data segment is loaded.

*Common Barrier-Option Use*

A barrier option might be useful when a buffer read takes multiple commands and must be performed before writing the buffer, which also takes multiple commands. In this example, the commands to read the buffer can be queued and performed in any order. Using the barrier-form for the first command to write the buffer allows the commands used to write the buffer to be queued without waiting for an MFC tag-group event on the read commands. (The barrier-form is only required for the first buffer-write command.) If the buffer read and buffer write commands have the same tag ID, the barrier ensures that the buffer is not written before being read. If

multiple commands are used to read and write the buffer, using the barrier option allows the read commands to be performed in any order and the write commands to be performed in any order, which provides better performance but forces all reads to finish before the writes start.

Barrier commands are also useful when performing double-buffered DMA transfers in which the data buffers used for the input data are the same as the output data buffers. Consider the following example for such a double-buffered environment:

```
int i;
i = 0;
read buffer 0
while (more buffers) {

    read buffer i^1
    wait for buffer i
    compute buffer i
    write buffer i
    i = i^1;
}
wait buffer i
compute buffer i
write buffer i
```

At the bottom of the loop, data is written from the same buffer that is read into immediately at the top of the next loop iteration. It is critical that the writes complete before the reads are started. Therefore, the first read (at the top of the loop) should be a barrier read. (For examples of using double-buffering techniques to overlap DMA transfers with computation on an SPU, see *Section 24.1.2* on page 692.)

### 20.1.5 MFC Multisource Synchronization Facility

The PPE **sync** instruction (*Section 20.1.2.1* on page 564) provides cumulative ordering, such that all threads in the system see the same barrier sequence that is performed and seen by the PPE, one of whose two threads executes the **sync** instruction. However, **sync** is only cumulative with respect to the main-storage domain. The CBEA processors contain multiple address and communication domains—the main-storage domain, eight local LS-address domains, and eight local channel domains—as illustrated in *Figure 1-2* on page 47. To ensure cumulative ordering across all address domains, the MFC multisource synchronization facility must be used.

Standard PowerPC storage-ordering rules apply to storage accesses performed by one processor element or device with respect to another processor element or device. Ordering of storage accesses performed by multiple sources (that is, two or more processor elements or devices) with respect to another processor element or device is referred to as cumulative ordering, as described in the *PowerPC Virtual Environment Architecture, Book II*. Cumulative ordering is ensured when all accesses are performed in the main-storage (effective-address) domain and the proper synchronization instructions are performed. Cumulative ordering cannot be ensured when accesses are performed from both the main-storage and the LS-address domains.

The MFC multisource synchronization facility addresses this cumulative-ordering need by providing two independent multisource synchronization-request methods:

- The MFC Multisource Synchronization Register, which allows the PPE or other processor elements or devices to control synchronization from the main-storage domain
- The MFC Write Multisource Synchronization Request Channel, which allows an SPE to control synchronization from its LS-address and channel domain

Each of these two synchronization-request methods ensures that all write transfers to the associated MFC are sent and received by the MFC before the MFC synchronization-request is completed. This facility does not ensure that read data is visible at the destination when the associated MFC is the source.

The channel and register supporting the MFC multisource synchronization facility are shown in *Table 20-8*. For details about channels and their associates MMIO registers, see *Section 17* on page 447.

*Table 20-8. MFC Multisource Synchronization Facility Channel and MMIO Register*

| SPE Channel # | Name | Channel Interface | | | | | MMIO Register Interface | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|
| | | Mnemonic | Max. Entries | Blocking | R/W | Width (bits) | Offset From Base | Mnemonic | Max. Entries | R/W | Width (bits) |
| 9 | MFC Write Multisource Synchronization Request | MFC_WrMSSyncReq | 1 | yes | W | 32 | x‘00000’ | MFC_MSSync | 1 | R/W | 64 |

These two synchronization-request methods operate independently. A synchronization request through the MMIO register has no effect on synchronization requests through the channel, and vice versa.

All outstanding requests that are in an SPE’s snoop-write queue when a Multisource Synchronization Request is received are completed before reporting the completion of the Multisource Synchronization Request.

#### 20.1.5.1 *MFC Multisource Synchronization Register*

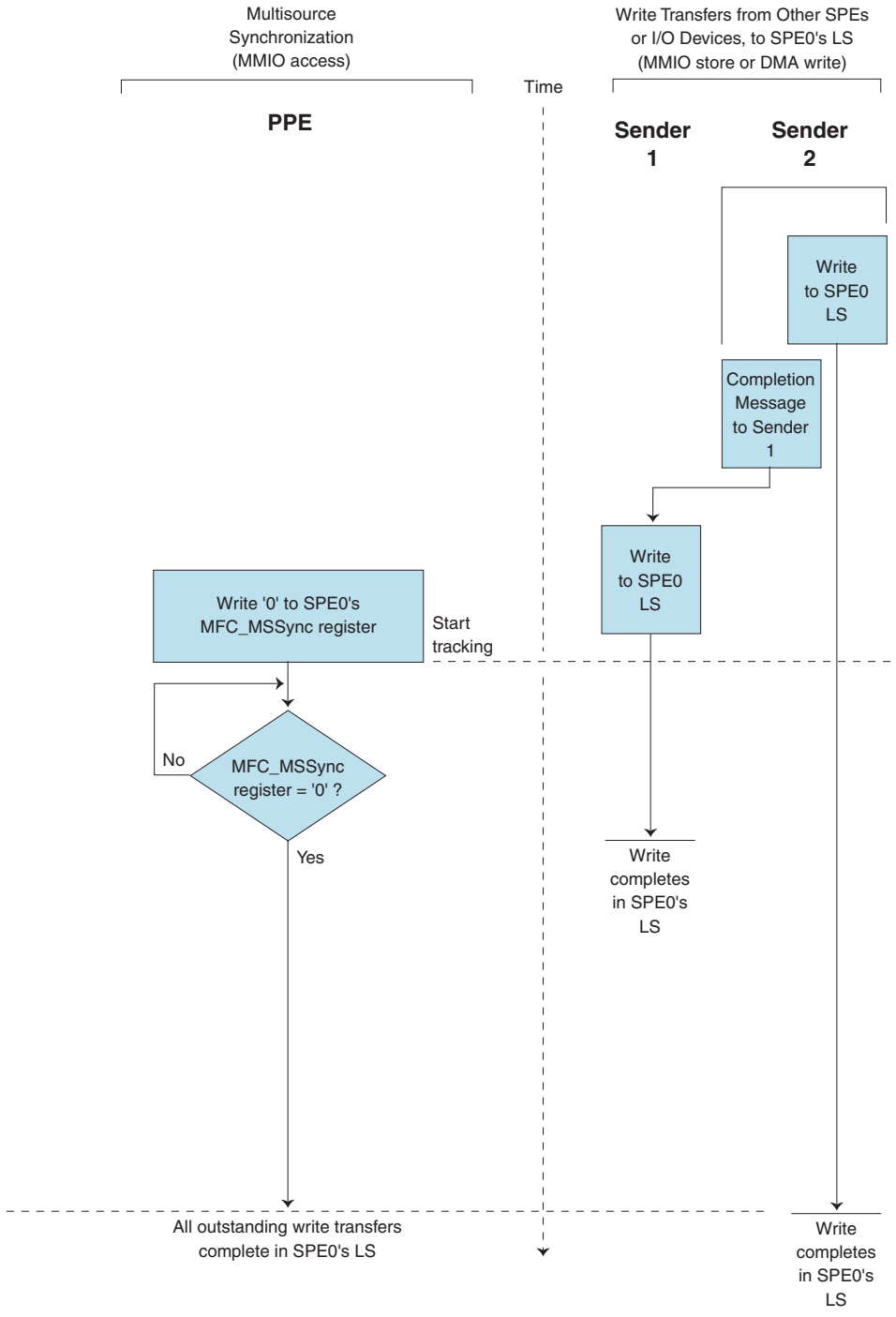
The MFC Multisource Synchronization Register (`MFC_MSSync`) allows processor elements or other devices to control ordering from the main-storage domain using the MMIO register at problem-state offset x‘00000’. Writing any value to this register causes the MFC to track all pending transfers targeted at the associated SPE received before the `MFC_MSSync` write. A read of the `MFC_MSSync` register will return a value of ‘1’ if any transfers being tracked are still outstanding. When all the transfers are being tracked are complete, a read of this register will return a value of ‘0’.

To use the MMIO facility, a program must perform the following steps:

1. Write any value to the `MFC_MSSync` register.
2. Poll the `MFC_MSSync` register until a value of ‘0’ is read.

*Figure 20-2* on page 579 shows the sequence of events for PPE control of multisource synchronization.

*Figure 20-2. PPE Multisource Synchronization Flowchart*




Use of this register is required for swapping context on an SPE. After stopping the SPU, privileged software on the PPE must prevent any new transfers from being initiated to the SPE by unmapping the associated resources. Next, privileged software must use the `MFC_MSSync` register

to ensure the completion of all outstanding transfers. This has the side-effect of ensuring that the count of the MFC Write Multisource Synchronization Request Channel, described in *Section 20.1.5.2*, is ‘1’ and that the SPU Read Event Status (*Section 18.6.2 Procedure for Handling the Multisource Synchronization Event* on page 483) is updated.

***Note:*** *Frequent use of a single MFC multisource synchronization facility by two or more processor elements or devices can result in a livelock condition. The livelock occurs when a read from a processor element or from a device never returns ‘0’ due to a synchronization request from other processor elements or devices.*

### 20.1.5.2 *MFC Write Multisource Synchronization Request Channel*

MFC Write Multisource Synchronization Request Channel (`MFC_WrMSSyncReq`) allows an SPE to control ordering from its LS-address and channel domain using its channel 9. Writing any value (‘0’ is recommended) to this channel causes the MFC to track all pending transfers targeted at itself that are received before the channel write. A second write to this channel will result in the SPE stalling until the outstanding transfers being tracked by the first write are complete; then, the second synchronization request will receive its response. A read of the channel count will return ‘0’ if any transfers being tracked are still outstanding. When all the transfers are being tracked are complete, a read of this channel will return a value of ‘1’.

To use the channel facility, a program must perform the following steps:

1. Write any value (‘0’ is recommended) to the `MFC_WrMSSyncReq` channel.
2. Wait for the `MFC_WrMSSyncReq` channel to become available. This can be accomplished by one of the following means:
   - Non-Event-Based:
     - Initiate a second multisource synchronization request by writing to the `MFC_WrMSSyncReq` channel. This method will block, waiting for a previous request to complete, without having to poll.
     - Poll the `MFC_WrMSSyncReq` channel count for the count to be set back to ‘1’. The `MFC_WrMSSyncReq` channel has a maximum count of ‘1’.
   - Event-Based:
     - Read the `SPU_RdEventStat` channel. This method will block, waiting for a multisource synchronization event, indicating that all outstanding writes have completed.
     - Poll the `SPU_RdEventStat` channel count for the count to be set back to ‘1’.
     - Wait for an event interrupt, then determine the cause.
     - Loop periodically on execution of a **bisled** instruction. This instruction uses the channel count of the `SPU_RdEventStat` channel for the condition in the branch, as described in *Section 18.2.3.2 Branch Condition* on page 474.

*Figure 20-4* on page 582 shows the sequence of events for SPE control of multisource synchronization using monitoring methods unrelated to events. *Figure 20-4* on page 582 shows the sequence of events for SPE control of multisource synchronization using event-based monitoring methods.

*Figure 20-3. SPE Multisource Synchronization Flowchart (Non-Event-Based)*

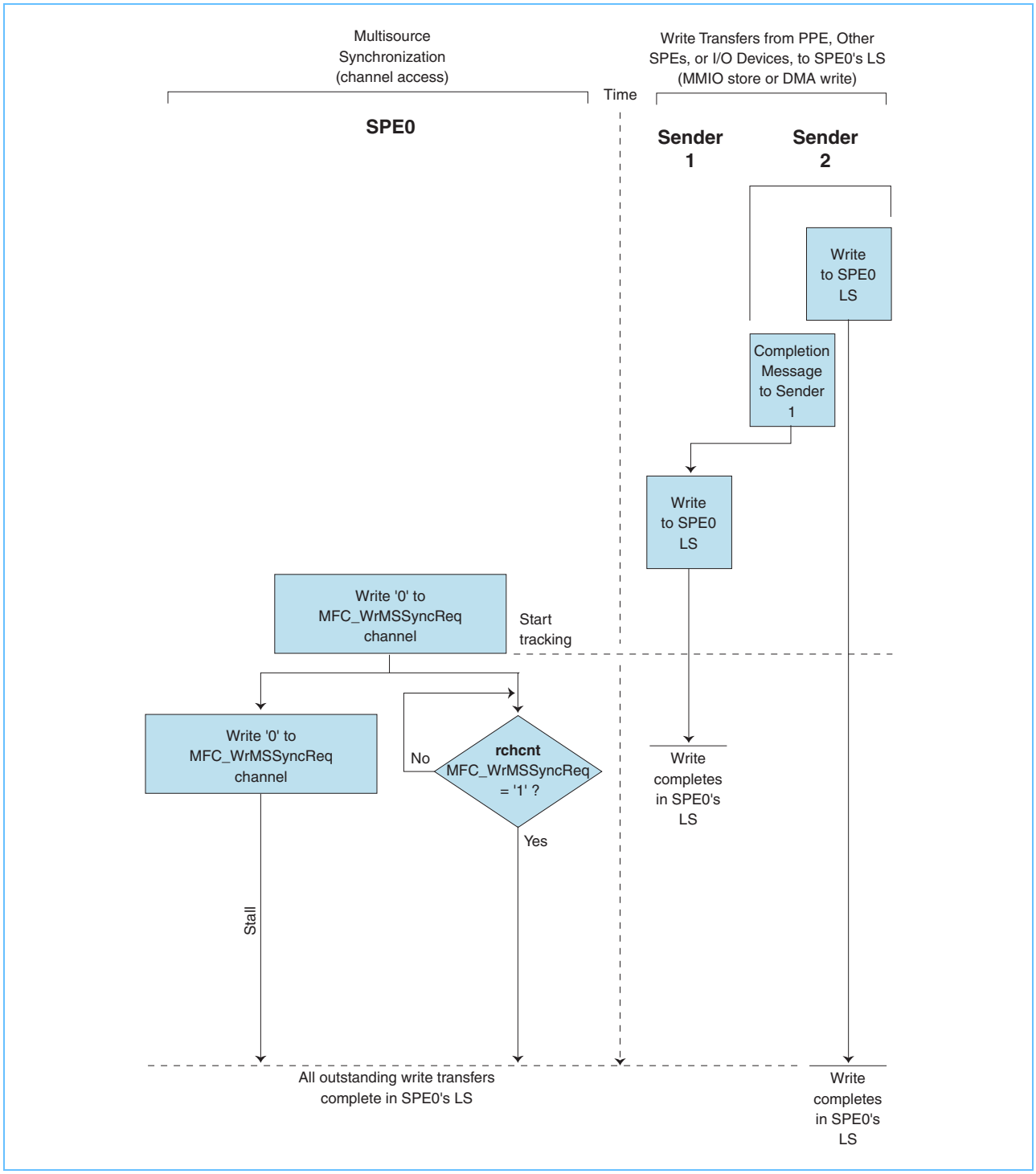
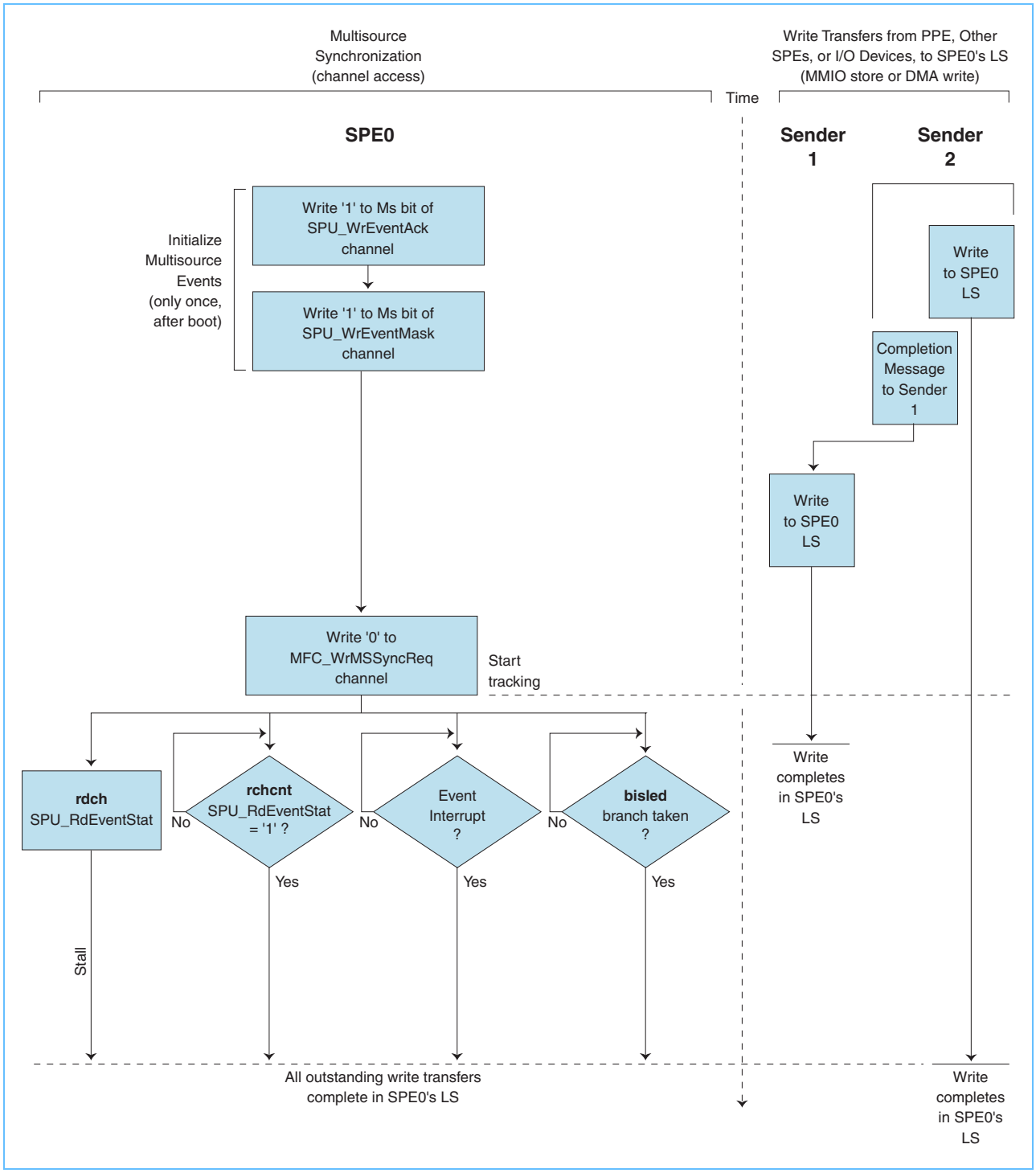


*Figure 20-4. SPE Multisource Synchronization Flowchart (Event-Based)*




For details about the procedures for handling a multisource synchronization event, see *Section 18.6.2* on page 483.

#### 20.1.5.3 *Multisource Synchronization Examples*

To best understand when ordering across multiple domains is required and how to use each of these facilities, consider the following examples of an I/O device storing a value of ‘1’ to an LS area, X, that is aliased in main storage, followed by a store of the value ‘2’ to a main-storage area, Y. The PPE reads main-storage area Y and stores a value of ‘3’ to another location, Z, in the LS area aliased in main storage. Multisource synchronization is required to guarantee that the value of ‘1’ that is stored in location X is ordered, with respect to the SPE, before the value of ‘3’ that is stored to location Z.

The following three examples illustrate the use of MMIO and channel multisource synchronization. The third example illustrates the use of the MMIO facility when starting an SPE. Without these facilities, the SPE is not guaranteed to receive a value of ‘1’ when reading its LS location, X.

Throughout these examples, recall that I/O transfers are always in-order.

*Example 1: MFC Multisource Synchronization Register*

This example demonstrates MMIO-initiated SPE multisource synchronization:

- *I/O Device:*
  a. Stores a value of ‘1’ to LS-aliased X.
  b. Stores a value of ‘2’ to main-storage location Y.
- *PPE:*
  a. Loops loading from main storage location Y until a value of ‘2’ is obtained.
  b. Stores any value to `MFC_MSSync` register.
  c. Polls `MFC_MSSync` register until a ‘0’ is obtained.
  d. Stores a value of ‘3’ to LS-aliased Z.
- *SPE:*
  a. Loops loading from LS location Z until a value of ‘3’ is obtained.
  b. Loads from LS location X, guaranteed to read a value of ‘1’

This example is provided for explanation purposes only. It is not intended as a recommendation for the use of polling loops.

*Example 2: MFC Write Multisource Synchronization Request Channel*

This example demonstrates channel-initiated SPE multisource synchronization:

- *I/O Device:*
  a. Stores a value of ‘1’ to LS-aliased X.
  b. Stores a value of ‘2’ to main-storage location Y.
- *PPE:*
  a. Loops loading from main storage location Y until a value of ‘2’ is obtained.
  b. Stores a value of ‘3’ to LS-aliased Z.

- *SPE:*
  a. Writes a ‘0’ to the MFC Write Multisource Synchronization Request (`MFC_WrMSSyncReq`) channel.
  b. Loads from LS location Z and obtains a 3.
  c. Wait for an SPE multisource synchronization event to occur by one of the following means:
     — Waiting for the channel count of the `MFC_WrMSSyncReq` channel to become ‘1’.
     — Waiting for a multisource synchronization request event.
  d. Loads from LS location X, guaranteed to read a value of ‘1’.

*Example 3: MFC Multisource Synchronization Register*

This example demonstrates MMIO-initiated SPE multisource synchronization when starting an SPE:

- *I/O Device:*
  a. Stores a value of ‘1’ to LS-aliased X.
  b. Interrupts PPE.
- *PPE:*
  a. Receives interrupt from I/O Device.
  b. Stores any value to `MFC_MSSync` register.
  c. Polls `MFC_MSSync` register until a ‘0’ is obtained.
  d. Starts an SPE program by writing a value of ‘1’ to the SPE Run Control Register.
- *SPE:*
  a. Begins execution as a result of step d.
  b. Loads from LS location X, guaranteed to read a value of ‘1’.

### 20.1.6 Scenarios for Using Ordering Mechanisms

This section includes some examples showing how the storage-ordering facilities of the CBEA processor can be used.

#### 20.1.6.1 *PPE-to-SPE Communications*

When the PPE is used as an application controller, managing and distributing work to the SPEs, the PPE typically loads main storage with the data to be processed, and then the PPE notifies the SPE by writing to either the PPE-to-SPE mailbox or one of the SPE’s signal-notification registers.

To make this feasible, it is important that the data storage be visible to the SPE before receiving the work-request notification. To ensure guaranteed ordering, a **lwsync** storage barrier instruction must be issued by the PPE between the final data store in memory and the PPE write to the SPE mailbox or signal-notification register.

#### 20.1.6.2 *SPE-to-PPE Communications*

Consider a case in which the SPE places computational results into main storage, waits for the DMA transfers to complete, and then either writes to an SPE-to-PPE mailbox or performs a *stop and signal* instruction (**stop** or **stopd**) to notify the PPE that its computation is complete. Waiting for the DMA transfers to complete only ensures that the LS buffers are available for SPE reuse; it does not guarantee that data has been placed into main storage. In this case, the SPE might issue an **mfcsync** command before notifying the PPE. However, doing so is inefficient. Instead, the preferred method is to have the PPE receive the notification and then issue an **lwsync** instruction before accessing any of the resulting data.

Alternatively, the SPE can perform a writeback to main storage to notify the PPE that the SPE's computation is complete. A writeback is a flag written to main storage that notifies the PPE that a specific event (for example, data computation complete) has occurred. In this case, the data and the writeback must be ordered. To ensure ordering, the writeback must either use the fence **<f>** or barrier **<b>** option.

## 20.2 PPE Atomic Synchronization

This section describes the use of PPE atomic operations to create semaphores and mutex locks for synchronization of storage or other functions among applications. *Section 20.3* on page 597 describes the comparable subject for SPE atomic operations.

### 20.2.1 Atomic Synchronization Instructions

The PPE atomic synchronization instructions include the **lwarx**, **ldarx**, **stwcx.**, and **stdcx.** instructions.

The load and reserve instructions (**lwarx**, **ldarx**) load the addressed value from memory and then set a reservation on an aligned unit of real storage (called a *reservation granule*) containing the address. The CBEA processor reservation granule is 128 bytes, corresponding to the size of a PPE cache line. These instructions control the order in which memory operations are completed with respect to asynchronous events, and the order in which memory operations are seen by other processor elements or memory access mechanisms.

A subsequent store conditional instruction (**stwcx**, **stdcx**) to this address verifies that the reservation is still set on the granule before carrying out the store. If the reservation does not exist, the instruction completes without altering storage. If the store is performed, bit 2 of CR0 is set to '1'; otherwise, it is cleared to '0'. The processor element clears the reservation by setting another reservation or by executing a conditional store to any address. Another processor element may clear the reservation by accessing the same reservation granule. A pair of load-and-reserve and store-conditional instructions permits the atomic update of a single aligned word or doubleword (only in 64-bit implementations) in memory.

A compiler that directly manages threads may use these instructions for in-line locks and to implement wait-free updates using primitives similar to compare and swap. Because locked or synchronized operations in multiprocessor systems are complex, these operations are typically exposed only through calls to appropriate runtime library functions.

The **lwarx**, **ldarx**, **stwcx.**, and **stdcx.** instructions, when the load and store versions of the instructions are used together, permit atomic update of a storage location by implementing atomic primitives such as fetch-and-increment, fetch-and-decrement, and test-and-set. A programmer can choose to implement a lock, a condition variable, or a semaphore with these instructions. Examples of semaphore operations can be found in *Section 20.2.2* on page 587.

The **lwarx** instruction must be paired with an **stwcx.** instruction, and the **ldarx** instruction must be paired with an **stdcx.** instruction, with the same effective address specified by both instructions of the pair. The only exception is that an unpaired **stwcx.** or **stdcx.** instruction to any effective address can be used to clear any reservation held by the processor element. The conditional store is performed based upon the existence of a reservation established by the preceding **lwarx** or **ldarx** instruction. If the reservation exists when the store is executed, the store is performed and a bit is set in the Condition Register (CR). If the reservation does not exist when the store is executed, the target memory location is not modified and a bit is cleared in the CR.

The **lwarx**, **ldarx**, **stwcx.**, and **stdcx.** instructions allow software to read a semaphore, compute a result based on the value of the semaphore, store the new value back into the semaphore location only if that location has not been modified since it was first read, and determine if the store was successful. If the store was successful, the sequence of instructions from the read of the semaphore to the store that updated the semaphore appears to have been executed atomically (that is, no other processor element or mechanism modified the semaphore location between the read and the update), thus providing the equivalent of a real atomic operation. However, in reality, other processor elements might have read from the location during this operation.

At most one reservation exists simultaneously on any processor element. The address associated with the reservation can be changed by a subsequent **lwarx** or **ldarx** instruction. The conditional store is performed based upon the existence of a reservation established by the preceding **lwarx** or **ldarx** instruction.

A reservation held by the processor element is cleared (or may be cleared, in the case of the fourth and fifth items in the following list) by one of the following:

- The processor element holding the reservation executes another **lwarx** or **ldarx** instruction. This clears the first reservation and establishes a new one.
- The processor element holding the reservation executes any **stwcx.** or **stdcx.** instruction regardless of whether its address matches that of the **lwarx** or **ldarx**.
- Some other processor element executes a store or **dcbz** (data cache block clear to zero) to the same reservation granule (128 bytes), or modifies a referenced or changed bit in the same reservation granule.
- Some other processor element executes a **dcbtst** (data cache block touch for store), **dcbst** (data cache block store), or **dcbf** (data cache block flush) to the same reservation granule. Whether the reservation is cleared is undefined.
- Some other mechanism modifies a memory location in the same reservation granule (128 bytes). If the processor element holding the reservation modifies a Reference (R) or Change (C) bit in the same reservation granule, the lose of reservation is undefined.

Exceptions do not clear reservations. However, when a exception occurs, a **stwcx.** or **stdcx.** instruction might need to be issued to ensure that a **lwarx** or **ldarx** in the interrupted program is not paired with a **stwcx.** or **stdcx** in the new program.

See the *PowerPC Architecture* for descriptions of the atomic synchronization instructions. See *Section 20.2* on page 585 for additional memory synchronization instructions.

***Note:*** *The PowerPC Architecture is likely to be changed in the future to permit the reservation to be lost if a* ***dcbf*** *(data cache block flush) instruction is executed on the processor element holding the reservation. Therefore* ***dcbf*** *instructions should not be placed between a load and reserve instruction and the subsequent store conditional instruction.*

The following points provide general information about the **lwarx** and **stwcx.** instructions:

- In general, the **lwarx** and **stwcx.** instructions should be paired, and the same effective address (EA) should be used for both instructions. The only exception is that an unpaired **stwcx.** instruction to any (scratch) effective address can be used to clear any reservation held by the processor element.
- It is acceptable to execute an **lwarx** instruction for which no **stwcx.** instruction is executed.
- To increase the likelihood that forward progress is made, it is important that looping on **lwarx** or **stwcx.** pairs be minimized. For example, this can be achieved by testing the old value before attempting the store; were the order reversed, more **stwcx.** instructions might be executed, and reservations might more often be lost between the **lwarx** and the **stwcx.** instructions.
- The manner in which **lwarx** and **stwcx.** are communicated to other processor elements and mechanisms, and between levels of the memory subsystem within a given processor element, is chip-implementation-dependent.
- In a multiprocessor, livelock (a state in which processor elements interact in a way such that no processor element makes progress) is possible if a loop containing an **lwarx** and **stwcx.** pair also contains an ordinary store instruction for which any byte of the affected memory area is in the reservation granule of the reservation.
- The reservation granularity is 128 bytes. The **lwarx**, **ldarx**, **stwcx.**, and **stdcx.** instructions require the effective address to be aligned.

### 20.2.2 PPE Synchronization Primitives

The following examples demonstrate how the PPE's atomic memory synchronization instructions, **lwarx** and **stwcx**, can be used to achieve commonly used synchronization primitives. These include atomic operations such as addition, mutexes, and condition variables. Equivalent synchronization primitives for the SPE are provided in *Section 20.2.3 SPE Synchronization Primitives* on page 590.

Additional sample PPE synchronization primitives can be found in the *PowerPC Virtual Environment Architecture, Book II*. All of the synchronization primitive examples require that all synchronization variables reside in memory marked with the memory coherency required storage control attribute (see *Section 4.2.6.7 WIMG-Bit Storage Control Attributes* on page 91 for details).

#### 20.2.2.1 *Atomic Addition*

The atomic addition primitive atomically adds a value to a word in memory. GPR3 contains the pointer to the word in memory, GPR4 contains the value to be added to the memory word. The previous value is returned in GPR5.

```
atomic_add:
    lwarx    r5, 0, r3# load memory word with reservation
    add      r0, r5, r4# add value
    stwcx.   r0, 0, r3# store result if still reserved
    bne-     atomic_add# retry if reservation was lost
```

#### 20.2.2.2 *Mutex Lock and Unlock*

A mutex can be used ensure mutual exclusive access to a resource or object. It is used to cause other threads to wait will while the thread holding the mutex executes code in a critical section.

The atomic synchronization instructions can be used to construct such a mutex. The mutex consists of a memory variable, pre-initialized to 0, and at least two methods, lock and unlock, that operate on the mutex variable.

The following code sample demonstrates a spin lock. The spin lock waits indefinitely until the lock can be acquired. Ideally, if the lock cannot be acquired, control should be yielded to other tasks and retried at a later time.

In this example, GPR3 contains the pointer to the lock variable.

```
spin_lock:
     lwarx    r4, 0, r3# load lock variable with reservation
     cmpwi    r4, 0# is the mutex already locked
     bne-     spin_lock# keep waiting if it is
     li       r4, 1#
     stwcx.   r4, 0, r3# set the lock variable
     bne-     spin_lock# retry if store conditional failed
     isync    # delay subsequent instruction til lock acquired
```

To unlock the mutex, a simple store of zero to the mutex variable need only be done.

#### 20.2.2.3 *Condition Variables*

A condition variable is used by a thread to make itself wait until an expression involving shared data attains a particular state. Condition variable can be used to implement other synchronization constructs such as counting semaphores.

The three primary operations on condition variables are condition wait, condition signal, and condition broadcast. This sample condition variable implementation uses a 32-bit word, pointed to by GPR3, that is atomically accessed. The word is partitioned into two half words in which the least significant half word is incremented for each conditional wait request. While the most significant halfword contains a count of the number of threads signaled. The initial value of the condition variable word is zero.

Condition wait, atomically increments the “waiting halfword”, then waits until it is signaled, that is when the “signaled” (high) halfword equals or is larger then the incremented waiting halfword. A Portable Operating System Interface (POSIX) standard condition wait function also will include a mutex unlock at the start of the condition wait and a mutex lock at the end of the condition wait. These have been omitted from the sample for brevity sake.

```
cond_wait:                          # r3 = pointer to condition variable
  _retry:
    lwarx       r5,0,r3             # atomically fetch the condition variable
    addi        r0,r5,1             # increment the waiting count
    rlwinm      r6,r5,16,16,31      # extract signaling count for cond var
    clrlwi      r7,r0,16            # extract incremented waiting count
    rlwimi      r0,r5,0,0,15
    stwcx.      r0,0,r3             # store the incremented condition variable
    bne-        _retry              # retry increment if reservation was lost
    subf        r5,r6,r7            # delta = abs(waitint_cnt-signaling_cnt)
    subf        r7,r7,r6
    cmpwi       r5,0
    bge         _wait_for_signal
    mr          r5,r7
    UNLOCK MUTEX
                                    # r5 = delta, r6 = signaling count
  _wait_for_signal:                 # atomically wait until signaled
    lwarx       r0,0,r3             # re-fetch the condition variable
    rlwinm      r0,r0,16,16,31      # extract current signaling count
    subf        r7,r6,r0            # compute num signaled since wait start
    subf        r0,r0,r6
    cmpwi       r7,0
    bge         _skip_next
    mr          r7,r0
  _skip_next:
    cmpw        r5,r7               # if num signaled is less than delta
    bgt-        _wait_for_signal    # keep waiting
    LOCK MUTEX
```

Condition signal notifies one thread waiting on the condition variable. This is accomplished by atomically incrementing the “signaled” halfword if the waiting count is not equal to the signaled count. This will signal the oldest waiting thread to discontinue waiting.

```
cond_signal:
    lwarx       r4,0,r3             # load cond variable with reservation
    rlwinm      r5,r4,16,16,31      # compare signaled and waiting counts
    clrlwi      r6,r4,16
    cmplw       r5,r6
    beq         done                # do nothing if waiting == signaled
    addis       r4,r4,1             # increment signaled (high) halfword
    stwcx.      r4, 0, r3           # store result if still reserved
    bne-        cond_signal         # retry if reservation was lost
  _done:
```

All currently waiting threads can be signaled by making a broadcast request. A broadcast, is implemented by atomically copying the “waiting count” (low halfword) of the condition variable to the “signaled count” (high halfword) of the condition variable.

```
cond_broadcast:
    lwarx      r4, 0, r3         # load condition variable with reservation
    rlwimi     r4, r4,16,0,15    # copy low halfword to high halfword
    stwcx.     r4, 0, r3         # store result if still reserved
    bne-       cond_broadcast    # retry if reservation was lost
```

### 20.2.3 SPE Synchronization Primitives

The following examples demonstrate how the MFC synchronization commands, **getllar** and **putllc**, can be used to perform synchronization primitives that are compatible with equivalent PPE primitives demonstrated in *Section 20.2.2 PPE Synchronization Primitives* on page 587. The example primitives include atomic addition, mutexes, and condition variables.

#### 20.2.3.1 *Atomic Addition*

The atomic addition primitive atomically adds a value to a word in memory and returns the previous value. Because the granularity of the MFC synchronization commands is a cache line, this function can be easily extended to atomically alter an entire 128 byte block.

"ea_ptr" is a 64 bit effective address pointer to a naturally aligned 32-bit word in which "addend" is it to be added.

```
#include <spu_intrinsics.h>

int atomic_add_return(unsigned long long ea_ptr, int addend)
{
    int old_word;
    unsigned int offset, status;
    unsigned int ea_lo, ea_hi;
    volatile char buf[256], *buf_ptr;
    volatile int *word_ptr;

    // Determine the offset to the word within its cache line. Align
    // the effective address to a cache line boundary.
    ea_hi = (unsigned int)ea_ptr >> 32;
    ea_lo = (unsigned int)ea_ptr;
    offset = ea_lo & 0x7F;
    ea_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    word_ptr = (volatile int *)(buf_ptr + offset);

    do {
        // Get, with reservation, the cache line containing the word to be
    // atomically added.
        spu_mfcdma64(buf_ptr, ea_hi, ea_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);
```

```
    // Fetch original value and add with addend.
    old_word = *word_ptr;
    *word_ptr = old_word + addend;

    // Put, conditionally the cache line containing the modified word.
        spu_mfcdma64(buf_ptr, ea_hi, ea_lo, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;

    // Retry if the reservation was lost.
    } while (status);

    return (old_word);
}
```

#### 20.2.3.2 *Mutex Lock and Unlock*

A mutex can be used ensure mutual exclusive access to a resource or object. It is used to cause other threads to wait will while the thread holding the mutex executes code in a critical section.

The MFC synchronization commands can be used to construct such a mutex. The mutex consists of a memory variable, pre-initialized to 0, and at least two methods, lock and unlock, that operate on the mutex variable.

The following code sample demonstrates a spin lock. The spin lock waits indefinitely until the lock can be acquired. The lock-line reservation lost event is used to initiate a low power stall if the mutex is currently held by another party. The input parameter, `mutex_ptr`, is a 64-bit effective address of the naturally aligned lock word.

```
#include <spu_intrinsics.h>

void lock(unsigned long long mutex_ptr)
{
    unsigned int offset, status, events, mask;
    unsigned int mutex_lo, mutex_hi;
    volatile char buf[256], *buf_ptr;
    volatile int *lock_ptr;

    // Determine the offset to the mutex word within its cache line. Align
    // the effective address to a cache line boundary.
    mutex_hi = (unsigned int)mutex_ptr >> 32;
    mutex_lo = (unsigned int)mutex_ptr;
    offset = mutex_lo & 0x7F;
    mutex_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    lock_ptr = (volatile int *)(buf_ptr + offset);

    // Setup for use possible use of lock line reservation lost events.
```

```
    // Detect and discard phantom events.
    mask = spu_readch(SPU_RdEventMask);
    spu_writech(SPU_WrEventMask, 0);
    if (spu_readchcnt(SPU_RdEventStat)) {
        spu_writech(SPU_WrEventAck, spu_readch(SPU_RdEventStat));
    }
    spu_writech(SPU_WrEventMask, MFC_LLR_LOST_EVENT);

    do {
        // Get, with reservation, the cache line containing the mutex lock
    // word.
        spu_mfcdma64(buf_ptr, mutex_hi, mutex_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    if (*lock_ptr) {
        // The mutex is currently locked. Wait for the lock line
        // reservation lost event before checking again.
        events = spu_readch(SPU_RdEventStat);
        spu_writech(SPU_WrEventAck, events);

        status = MFC_PUTLLC_STATUS;
    } else {
        // The mutex is not currently locked. Attempt to lock.
        *lock_ptr = 1;

        // Put, conditionally, the cache line containing the lock word.
        spu_mfcdma64(buf_ptr, mutex_hi, mutex_lo, 128, 0, MFC_PUTLLC_CMD);
        status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;
    }

    // Retry if the reservation was lost.
    } while (status);

    // Restore the event mask
    spu_writech(SPU_WrEventMask, mask);
}
```

To unlock the mutex, a simple word store of zero to the mutex variable need only be done. This can be accomplished by either atomically writing zero to the mutex word, or alternatively, performing a DMA put to the mutex variable.

```
#include <spu_intrinsics.h>

void unlock(unsigned long long mutex_ptr)
{
    unsigned int mask;
    volatile vector unsigned int zero = (vector unsigned int){0};
    volatile void * word_ptr;
    unsigned int mutex_lo, mutex_hi;
```

```
    // Determine the offset to the word within its quadword so that our
    // source and destination addresses have the same sub-quadword alignment.
    mutex_hi = (unsigned int)mutex_ptr >> 32;
    mutex_lo = (unsigned int)mutex_ptr;

    word_ptr = (volatile void *)(&zero) + (mutex_lo & 0xF);

    spu_mfcdma64(word_ptr, mutex_hi, mutex_lo, 4, 0, MFC_PUT_CMD);

    // Wait until the DMA put completes before returning.
    mask = spu_readch(MFC_RdTagMask);
    spu_writech(MFC_WrTagMask, 1 << 0);
    (void)spu_mfcstat(MFC_TAG_UPDATE_ALL);
    spu_writech(MFC_WrTagMask, mask);
}
```

#### 20.2.3.3 *Condition Variables*

A condition variable is used by a thread to make itself wait until an expression involving shared data attains a particular state. Condition variable can be used to implement other synchronization constructs such as counting semaphores.

The three primary operations on condition variables are condition wait, condition signal, and condition broadcast. This sample condition variable implementation uses a 32-bit word that is atomically accessed. The word is partitioned into two half words counters in which the least significant half word is incremented for each conditional wait request. While the most significant half word contains a count of the number of threads signaled. The initial value of the condition variable word is zero.

Condition wait atomically increments the “waiting halfword”, then waits until it is signaled, that is when the “signaled” (high) halfword equals or is larger then the incremented waiting halfword. A POSIX standard condition wait function also will include a mutex unlock at the start of the condition wait and a mutex lock at the end of the condition wait.

The inputs `cond_ptr` and `mutex_ptr` are 64-bit effective address to the 32-bit condition variable and mutex lock word, respectively. These are assumed to be naturally (4 byte) aligned.

```
#include <spu_intrinsics.h>

void cond_wait(unsigned long long cond_ptr, unsigned long long mutex_ptr)
{
    signed short delta, cur_delta;
    unsigned short signaled_cnt, waiting_cnt;
    unsigned int offset, events, mask, status;
    unsigned int cond_lo, cond_hi;
    volatile char buf[256], *buf_ptr;
    volatile unsigned short *cv_ptr;

    // Determine the offset to the condition variable word within its
    // cacheline. Align the effective address to a cacheline boundary.
```

```
cond_hi = (unsigned int)cond_ptr >> 32;
cond_lo = (unsigned int)cond_ptr;
offset = cond_lo & 0x7F;
cond_lo &= ~0x7F;

// Cache line align the local stack buffer.
buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
cv_ptr = (volatile unsigned short *)(buf_ptr + offset);

// Setup for use possible use of lock line reservation lost events.
// Detect and discard phantom events.
mask = spu_readch(SPU_RdEventMask);
spu_writech(SPU_WrEventMask, 0);
if (spu_readchcnt(SPU_RdEventStat)) {
    spu_writech(SPU_WrEventAck, spu_readch(SPU_RdEventStat));
}
spu_writech(SPU_WrEventMask, MFC_LLR_LOST_EVENT);
// Increment the waiting halfword.
do {
    // Get, with reservation, the cache line containing the condition
    // variable word.
    spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    // Increment the low halfword - waiting count.
    signaled_cnt = *(cv_ptr+0);
    waiting_cnt = *(cv_ptr+1) + 1;
    *(cv_ptr+1) = waiting_cnt;

    // Put, conditionally, the cache line containing the condition variable.
    spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;
} while (status);

// Unlock the mutex
unlock(mutex_ptr);

delta = waiting_cnt - signaled_cnt;
if (delta < 0) delta = -delta;

// Wait until a signal is received. I.E., high halfword of the condition
// variable is greater than or equal to the waiting_cnt fetched when
// the wait was established with the condition variable.

while (1) {
    // Get, with reservation, the cache line containing the condition
    // variable word.
    spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);
```

```
            cur_delta = *(cv_ptr+0) - signaled_cnt;
            if (cur_delta < 0) cur_delta = -cur_delta;

            if (cur_delta < delta) {
                // No signal received yet. Wait until the lock line reservation
                // is lost before checking again.
                events = spu_readch(SPU_RdEventStat);
                spu_writech(SPU_WrEventAck, events);
            } else {
                // Signal received. Done waiting.
                break;
            }
        }

        // Restore the event mask
        spu_writech(SPU_WrEventMask, mask);

        lock(mutex_ptr);
    }
```

Condition signal notifies one thread waiting on the condition variable. This is accomplished by atomically incrementing the “signaled” halfword count if the waiting count is not equal to the signaled count. This will signal the oldest waiting thread to discontinue waiting.

```
#include <spu_intrinsics.h>

void cond_signal(unsigned long long cond_ptr)
{
    unsigned int offset, status;
    unsigned int cond_lo, cond_hi;
    volatile char buf[256], *buf_ptr;
    volatile unsigned short *cv_ptr;

    // Determine the offset to the word within its cache line. Align
    // the effective address to a cache line boundary.
    cond_hi = (unsigned int)cond_ptr >> 32;
    cond_lo = (unsigned int)cond_ptr;
    offset = cond_lo & 0x7F;
    cond_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    cv_ptr = (volatile unsigned short *)(buf_ptr + offset);

    do {
        // Get, with reservation, the cache line containing the condition
    // variable word.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);
```

```
    // If there is anyone waiting on the condition variable, increment
    // the waiting count. Else, do nothing.
    if (*(cv_ptr+0) != *(cv_ptr+1)) {
        *cv_ptr++;

        // Conditionally put the lock line containing the adjusted
        // signaled halfword.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_PUTLLC_CMD);
        status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;

    } else {
        status = 0;
    }
    } while (status);
}
```

All currently waiting threads can be signaled by making a broadcast request. A broadcast, is implemented by atomically copying the “waiting” count (low halfword) of the condition variable to the “signaled” count (high halfword) of the condition variable.

```
#include <spu_intrinsics.h>

void cond_broadcast(unsigned long long cond_ptr)
{
    unsigned int offset, status;
    unsigned int cond_lo, cond_hi;
    volatile char buf[256], *buf_ptr;
    volatile unsigned short *cv_ptr;

    // Determine the offset to the word within its cache line. Align
    // the effective address to a cache line boundary.
    cond_hi = (unsigned int)cond_ptr >> 32;
    cond_lo = (unsigned int)cond_ptr;
    offset = cond_lo & 0x7F;
    cond_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    cv_ptr = (volatile unsigned short *)(buf_ptr + offset);

    do {
        // Get, with reservation, the cache line containing the condition
    // variable word.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    // Copy the wait count into the signaled count.
    *(cv_ptr+0) = *(cv_ptr+1);
```

```
    // Conditionally put the lock line containing the adjusted
    // signaled halfword.
    spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;
    } while (status);
}
```

## 20.3 SPE Atomic Synchronization

This section describes the use of MFC atomic commands to create locks for synchronization of storage or other functions among applications. *Section 20.2* on page 585 describes a comparable subject for PPE atomic operations.

The MFC for each SPE provides two main functions for the SPE: moving data between the SPE's LS and main storage, and providing synchronization between the SPE and other processor elements in the system. The MFC has an atomic unit that supports synchronization. In the Cell/B.E. and PowerXCell 8i implementations, the atomic unit has a buffer that holds six 128-byte cache lines, and it dedicates four of its six cache lines to providing synchronization functions with other processing units in the system. The atomic unit does this by processing atomic MFC commands and by maintaining cache coherence when the MFC receives snoop requests from other processor elements in the system. Even though there are four cache lines dedicated to synchronization, only one reservation is maintained in each atomic unit. Four cache lines are implemented in a nested fashion (up to four deep) to provide high performance in situations where atomic updates are used for lock acquisition.

### 20.3.1 MFC Commands for Atomic Updates

Each MFC supports commands for atomic operations. The *get and reserve* (**getllar**) and *put conditional* (**putllc**) commands provide essentially the same functions as the PowerPC **lwarx** and **stwcx** instructions, described in *Section 20.2* on page 585. Together, these MFC commands permit atomic update of a main-storage location. In addition, the MFC supports the *put unconditional* **(putlluc**) and the *put unconditional in queued* (**putqlluc**) commands.

The **getllar**, **putllc**, and **putlluc** commands are not tagged, and they are executed immediately instead of being queued behind other DMA commands. Attempts to issue these three commands before a previous **getllar, putllc,** or **putlluc** command has completed result in an error. The **putqlluc** command enables additional **putlluc** commands to be queued while the previous command executes. These four atomic-update commands can only be issued by the SPU to the MFC SPU command queue.

*Table 20-9* lists the MFC commands for atomic updates.

*Table 20-9. MFC Commands for Atomic Updates* (Sheet 1 of 2)

| Command | Opcode | Description |
|---|---|---|
| **getllar** | x'D0' | Get lock-line and create a reservation (executed immediately) |

*Table 20-9. MFC Commands for Atomic Updates* (Sheet 2 of 2)

| Command | Opcode | Description |
|---|---|---|
| **putllc** | x'B4' | Put lock-line conditional on a reservation (executed immediately) |
| **putlluc** | x'B0' | Put lock-line unconditional (executed immediately) |
| **putqlluc** | x'B8' | Put lock-line unconditional (queued form) |

#### 20.3.1.1 *The Get Lock-Line and Reserve Command—getllar*

The get lock-line and reserve (**getllar**) command is similar to the PowerPC **lwarx** instruction, with the exception of the size and destination of the load. The **getllar** command loads 128 bytes (the PPE cache-line size) into the SPE's LS and sets a reservation visible in main storage. Issuing the **getllar** command requires an available slot in the MFC SPU command queue. This command is issued immediately, is not queued behind other commands, and has no associated tag. An attempt to issue this command before a previous **getllar**, **putllc**, or **putlluc** command has completed results in an error.

A read from the MFC Read Atomic Command Status Channel (`MFC_RdAtomicStat`) must be performed after issuing the **getllar** command, before issuing another **getllar, putllc,** or **putlluc** command.

#### 20.3.1.2 *The Put Lock-Line Conditional Command—putllc*

The put lock-line conditional (**putllc**) command is similar to the operation of the PowerPC **stwcx** instruction, with the exception of the size and source of the store-conditional. This is a **put** command of 128 bytes that is conditioned on two things:

- The existence of the reservation created by the **getllar**
- Whether the same storage location is specified by both commands

The **putllc** command is a conditional-store operation. The store will not be successful if no reservation for the same address has been made, or if the reservation has been lost. A read of the MFC Read Atomic Command Status Channel (`MFC_RdAtomicStat`) is required to determine the success or failure of this command.

#### 20.3.1.3 *The Put Lock-Line Unconditional Command—putlluc*

The put lock-line unconditional (**putlluc**) command is similar to the operation of the **putllc** command, but the store for the **putlluc** is always performed. The **putlluc** command store is not dependent on the existence of a previously-made reservation. The store size of the **putlluc** command is 128 bytes.

#### 20.3.1.4 *The Put Queued Lock-Line Unconditional Command—putqlluc*

The put queued lock-line unconditional (**putqlluc**) command is functionally equivalent to the put lock-line unconditional (**putlluc**) command. The difference between the two commands is the order in which the commands are performed and how completion is determined. The **putlluc** is performed immediately (not queued), and the **putqlluc** is placed into the MFC SPU command queue along with other DMA commands. Because this command is queued, it executes independently of any pending immediate **getllar**, **putllc**, or **putlluc** commands.

To determine if the **putqlluc** command is complete, software must wait for a tag-group completion. The **putqlluc** command requires a tag input and has an implied tag-specific fence. This command cannot be issued unless all previously issued commands with the same tag have completed.

Multiple **putqlluc** commands may be issued and pending in the MFC SPU command queue. All the command-queue ordering rules apply to the **putqlluc** command.

The **putqlluc** command allows software to queue the release of a lock behind commands accessing storage associated with the lock. For proper operation, the **putqlluc** command must either be within the same tag group as the commands accessing the storage, or the barrier command must be used to ensure ordering.

### 20.3.2 The MFC Read Atomic Command Status Channel

The MFC Read Atomic Command Status Channel (`MFC_RdAtomicStat`), channel 27, contains the status of the last-completed immediate atomic update DMA command (**getllar**, **putllc**, or **putlluc**). A read from this channel before issuing an atomic command results in a software-induced deadlock. For details about the channel interface, see *Section 17* on page 447.

Software can read the count associated with this channel to determine if an atomic DMA command has completed. A value of '0' is returned if an atomic DMA command has not completed. A value of '1' is returned if an atomic DMA command has completed and the status is available by reading this channel. A read from the `MFC_RdAtomicStat` channel should always follow an atomic DMA command. Performing multiple atomic DMA commands without an intervening read from the `MFC_RdAtomicStat` channel might result in an incorrect status.

This channel is read-blocking, with a maximum count of '1'. The contents of this channel are cleared when read. Completion of a subsequent atomic DMA command will overwrite the status of prior atomic DMA commands. *Table 20-10* shows the content of this channel.

*Table 20-10. MFC Read Atomic Command Status Channel Contents*

| Bits | Field Name | Description |
|---|---|---|
| 0:28 | Reserved | Reserved. |
| 29 | G | Set if the **getllar** command completed. |
| 30 | U | Set if the **putlluc** command completed. |
| 31 | S | Put lock-line conditional command status:<br>1 Put conditional unsuccessful. The reservation was lost.<br>0 Put conditional successful. |

### 20.3.3 Avoiding Livelocks

The MFC atomic synchronization commands have the benefit of allowing an SPE program to participate in atomic shared-memory updates with other processor elements in the system, including the PPE. However, care must be taken when applying these commands. Incorrect use of atomic synchronization commands can lead to large performance penalties or livelocks.

Serialize programs as infrequently as possible to maximize parallelism, per Amdahl's law[3]. Try to partition problems statically, and only use atomic operations. Livelock is a common problem if atomic commands are not used, and used carefully. The following sections illustrate how to avoid livelocks.

#### 20.3.3.1 *Lock-Line Reservation Lost Event*

Use lock-line reservation lost events instead of spin loops. This event is triggered when a get lock-line and reserve (**getllar**) command is issued and the reservation is reset due to modification of the data in the same lock-line by an outside entity. It will not be set due to a reservation reset by a local action.

Using the notification of this event allows the program to accomplish other tasks while waiting for an external modification to the lock-line data. If no other task is available, software can perform a read from the SPU Read Event Status Channel (`SPU_RdEventStat`) to put the SPE into a low-power state until the lock-line data has been modified. A lock-line reservation lost event can be used to implement a sleeping lock rather than spin locks.

For details about handling lock-line reservation lost events, see *Section 18.6.4* on page 485.

Here is an example of how to use atomic commands with lock-line reservation lost events:

```
       * In this example, ea and lsbuf are 128-byte aligned */

/* Disable all events */
spu_writech (SPU_WrEventMask, 0);

/* Clear phantom events */
if (spu_readchcnt (SPU_RdEventStat)) spu_readch (SPU_RdEventStat);
spu_writech(SPU_WrEventAck, MFC_LLAR_LOST_EVENT_MASK);

/* set event mask */
spu_writech (SPU_WrEventMask, MFC_LLAR_LOST_EVENT_MASK);

do {
  /* Atomically fetch the synchronization variable
   */
  spu_mfcdma32(lsbuf, ea, 128, 0, MFC_GETLLAR_CMD);
  (void)spu_readch (MFC_RdAtomicStat);

  if (lsbuf[0]) {
    /* Non-zero synchronization variable, wait until
     * reservation is lost
     */
    event = spu_readch(SPU_RdEventStat);
    spu_writech (SPU_WrEventAck, MFC_LLAR_LOST_EVENT_MASK);
    status = MFC_PUTLLC_STATUS;
```

3. If N is the number of processor elements, and s is the amount of time spent by one processor element on a serial parts of a program and p is the amount of time spent by one processor element on parts of the program that can be done in parallel, then the speedup is given by Speedup = 1/(s + p/N). So, speedup is inversely proportional to s, the amount of time spent in the serial parts of a program.

```
  } else {
    /* Zero synchronization variable. Atomically set it.
     */
    lsbuf[0] = 1;
    spu_mfcdma32(lsbuf, ea, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch (MFC_RdAtomicStat);
  }
} while (status & MFC_PUTLLC_STATUS);

/* Clean up any remnant events, as needed */
spu_writech (SPU_WrEventMask, 0);
if (spu_readchcnt (SPU_RdEventStat)) {
  event = spu_readch(SPU_RdEventStat);
  spu_writech (SPU_WrEventAck, MFC_LLAR_LOST_EVENT_MASK);
}
```

#### 20.3.3.2 *Nesting Atomic Primitives*

Avoid nesting more than four atomic primitives. The MFC commands provide a 4-line cache for data involved in atomic updates used for lock acquisition. Nesting more than four atomic primitives might degrade performance.

### 20.3.4 Synchronization Primitives

The following examples demonstrate how the MFC's synchronization commands, **getllar** and **putllc**, can be used to perform synchronization primitives that are compatible with equivalent PPE primitives demonstrated in *Section 20.2.2* on page 587. The example primitives include atomic addition, mutexes, and condition variables.

#### 20.3.4.1 *Atomic Addition*

The atomic addition primitive atomically adds a value to a word in memory and returns the previous value. Because the granularity of the MFC synchronization commands is a cache line, this function can be easily extended to atomically alter an entire 128-byte cache line.

`ea_ptr` is a 64-bit effective-address pointer to a naturally aligned 32-bit word, in which `addend` is it to be added.

```
#include <spu_intrinsics.h>

int atomic_add_return(unsigned long long ea_ptr, int addend)
{
    int old_word;
    unsigned int offset, status;
    unsigned int ea_lo, ea_hi;
    volatile char buf[256], *buf_ptr;
    volatile int *word_ptr;

    // Determine the offset to the word within its cache line. Align
    // the effective address to a cache line boundary.
```

```
    ea_hi = (unsigned int)ea_ptr >> 32;
    ea_lo = (unsigned int)ea_ptr;
    offset = ea_lo & 0x7F;
    ea_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    word_ptr = (volatile int *)(buf_ptr + offset);

    do {
        // Get, with reservation, the cache line containing the word to be
    // atomically added.
        spu_mfcdma64(buf_ptr, ea_hi, ea_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    // Fetch original value and add with addend.
    old_word = *word_ptr;
    *word_ptr = old_word + addend;

    // Put, conditionally the cache line containing the modified word.
        spu_mfcdma64(buf_ptr, ea_hi, ea_lo, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;

    // Retry if the reservation was lost.
    } while (status);

    return (old_word);
}
```

#### 20.3.4.2 *Mutex Lock and Unlock*

A mutex can be used ensure mutual-exclusive access to a resource or object. It is used to cause other threads to wait will while the thread holding the mutex executes code in a critical section.

This mutex code sample consists of a memory variable, pre-initialized to ‘0’, and at least two methods, lock and unlock, that operate on the mutex variable. The sample demonstrates a spin lock. The spin lock waits indefinitely until the lock can be acquired. The lock-line reservation lost event is used to initiate a low-power stall if the mutex is currently held by another party. The input parameter, `mutex_ptr`, is a 64-bit effective address of the naturally aligned lock word.

```
#include <spu_intrinsics.h>

void lock(unsigned long long mutex_ptr)
{
    unsigned int offset, status, events, mask;
    unsigned int mutex_lo, mutex_hi;
    volatile char buf[256], *buf_ptr;
    volatile int *lock_ptr;
```

```
    // Determine the offset to the mutex word within its cache line. Align
    // the effective address to a cache line boundary.
    mutex_hi = (unsigned int)mutex_ptr >> 32;
    mutex_lo = (unsigned int)mutex_ptr;
    offset = mutex_lo & 0x7F;
    mutex_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    lock_ptr = (volatile int *)(buf_ptr + offset);

    // Setup for use possible use of lock line reservation lost events.
    // Detect and discard phantom events.
    mask = spu_readch(SPU_RdEventMask);
    spu_writech(SPU_WrEventMask, 0);
    if (spu_readchcnt(SPU_RdEventStat)) {
        spu_writech(SPU_WrEventAck, spu_readch(SPU_RdEventStat));
    }
    spu_writech(SPU_WrEventMask, MFC_LLR_LOST_EVENT);

    do {
        // Get, with reservation, the cache line containing the mutex lock
    // word.
        spu_mfcdma64(buf_ptr, mutex_hi, mutex_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    if (*lock_ptr) {
        // The mutex is currently locked. Wait for the lock line
        // reservation lost event before checking again.
        events = spu_readch(SPU_RdEventStat);
        spu_writech(SPU_WrEventAck, events);

        status = MFC_PUTLLC_STATUS;
    } else {
        // The mutex is not currently locked. Attempt to lock.
        *lock_ptr = 1;

        // Put, conditionally, the cache line containing the lock word.
        spu_mfcdma64(buf_ptr, mutex_hi, mutex_lo, 128, 0, MFC_PUTLLC_CMD);
        status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;
    }

    // Retry if the reservation was lost.
    } while (status);

    // Restore the event mask
    spu_writech(SPU_WrEventMask, mask);
}
```

To unlock the mutex, a simple word store of zero to the mutex variable is done. This can be accomplished by either atomically, by writing zero to the mutex word, or by performing a DMA **put** to the mutex variable.

```
#include <spu_intrinsics.h>

void unlock(unsigned long long mutex_ptr)
{
    unsigned int mask;
    volatile vector unsigned int zero = (vector unsigned int){0};
    volatile void * word_ptr;
    unsigned int mutex_lo, mutex_hi;

    // Determine the offset to the word within its quadword so that our
    // source and destination addresses have the same sub-quadword alignment.
    mutex_hi = (unsigned int)mutex_ptr >> 32;
    mutex_lo = (unsigned int)mutex_ptr;

    word_ptr = (volatile void *)(&zero) + (mutex_lo & 0xF);

    spu_mfcdma64(word_ptr, mutex_hi, mutex_lo, 4, 0, MFC_PUT_CMD);

    // Wait until the DMA put completes before returning.
    mask = spu_readch(MFC_RdTagMask);
    spu_writech(MFC_WrTagMask, 1 << 0);
    (void)spu_mfcstat(MFC_TAG_UPDATE_ALL);
    spu_writech(MFC_WrTagMask, mask);
}
```

#### 20.3.4.3 *Condition Variables*

A condition variable is used by a thread to make itself wait until an expression involving shared data attains a particular state. Condition variables can be used to implement other synchronization constructs, such as counting semaphores.

The three primary operations on condition variables are condition wait, condition signal, and condition broadcast. This sample condition-variable implementation uses a 32-bit word that is accessed atomically. The word is partitioned into two halfword counters, in which the least-significant halfword is incremented for each conditional-wait request while the most-significant halfword contains a count of the number of threads signaled. The initial value of the condition-variable word is zero.

Condition wait atomically increments the “waiting halfword”, then waits until it is signaled—that is, when the “signaled” (high) halfword equals or is larger then the incremented waiting halfword. A POSIX standard condition wait function also will include a mutex unlock at the start of the condition wait and a mutex lock at the end of the condition wait.

The inputs `cond_ptr` and `mutex_ptr` are 64-bit effective address to the 32-bit condition variable and mutex lock word, respectively. These are assumed to be naturally (4-byte) aligned.

```
#include <spu_intrinsics.h>

void cond_wait(unsigned long long cond_ptr, unsigned long long mutex_ptr)
{
    signed short delta, cur_delta;
    unsigned short signaled_cnt, waiting_cnt;
    unsigned int offset, events, mask, status;
    unsigned int cond_lo, cond_hi;
    volatile char buf[256], *buf_ptr;
    volatile unsigned short *cv_ptr;

    // Unlock the mutex
    unlock(mutex_ptr);

    // Determine the offset to the condition variable word within its
    // cacheline. Align the effective address to a cache line boundary.
    cond_hi = (unsigned int)cond_ptr >> 32;
    cond_lo = (unsigned int)cond_ptr;
    offset = cond_lo & 0x7F;
    cond_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    cv_ptr = (volatile unsigned short *)(buf_ptr + offset);

    // Setup for use possible use of lock line reservation lost events.
    // Detect and discard phantom events.
    mask = spu_readch(SPU_RdEventMask);
    spu_writech(SPU_WrEventMask, 0);
    if (spu_readchcnt(SPU_RdEventStat)) {
        spu_writech(SPU_WrEventAck, spu_readch(SPU_RdEventStat));
    }
    spu_writech(SPU_WrEventMask, MFC_LLR_LOST_EVENT);

    // Increment the waiting halfword.
    do {
        // Get, with reservation, the cache line containing the condition
    // variable word.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    // Increment the low halfword - waiting count.
    signaled_cnt = *(cv_ptr+0);
    waiting_cnt = *(cv_ptr+1) + 1;
    *(cv_ptr+1) = waiting_cnt;

    // Put, conditionally, the cache line containing the condition variable.
    spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;
    } while (status);
```

```
    delta = waiting_cnt - signaled_cnt;
    if (delta < 0) delta = -delta;

    // Wait until a signal is received. I.E., high halfword of the condition
    // variable is greater than or equal to the waiting_cnt fetched when
    // the wait was established with the condition variable.

    while (1) {
        // Get, with reservation, the cache line containing the condition
    // variable word.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    cur_delta = *(cv_ptr+0) - signaled_cnt;
    if (cur_delta < 0) cur_delta = -cur_delta;

    if (cur_delta < delta) {
        // No signal received yet. Wait until the lock line reservation
        // is lost before checking again.
        events = spu_readch(SPU_RdEventStat);
        spu_writech(SPU_WrEventAck, events);
    } else {
        // Signal received. Done waiting.
        break;
    }
    }

    // Restore the event mask
    spu_writech(SPU_WrEventMask, mask);

    lock(mutex_ptr);
}
```

Condition signal notifies one thread waiting on the condition variable. This is accomplished by atomically incrementing the “signaled” halfword count if the waiting count is not equal to the signaled count. This will signal the oldest waiting thread to discontinue waiting.

```
#include <spu_intrinsics.h>

void cond_signal(unsigned long long cond_ptr)
{
    unsigned int offset, status;
    unsigned int cond_lo, cond_hi;
    volatile char buf[256], *buf_ptr;
    volatile unsigned short *cv_ptr;

    // Determine the offset to the word within its cache line. Align
    // the effective address to a cache line boundary.
    cond_hi = (unsigned int)cond_ptr >> 32;
```

```
    cond_lo = (unsigned int)cond_ptr;
    offset = cond_lo & 0x7F;
    cond_lo &= ~0x7F;

    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    cv_ptr = (volatile unsigned short *)(buf_ptr + offset);

    do {
        // Get, with reservation, the cache line containing the condition
    // variable word.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    // If there is anyone waiting on the condition variable, increment
    // the waiting count. Else, do nothing.
    if (*(cv_ptr+0) != *(cv_ptr+1)) {
        *cv_ptr++;

        // Conditionally put the lock line containing the adjusted
        // signaled halfword.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_PUTLLC_CMD);
        status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;

    } else {
        status = 0;
    }
    } while (status);
}
```

All currently waiting threads can be signaled by making a broadcast request. A broadcast is implemented by atomically copying the “waiting” count (low halfword) of the condition variable to the “signaled” count (high halfword) of the condition variable.

```
#include <spu_intrinsics.h>

void cond_broadcast(unsigned long long cond_ptr)
{
    unsigned int offset, status;
    unsigned int cond_lo, cond_hi;
    volatile char buf[256], *buf_ptr;
    volatile unsigned short *cv_ptr;

    // Determine the offset to the word within its cache line. Align
    // the effective address to a cache line boundary.
    cond_hi = (unsigned int)cond_ptr >> 32;
    cond_lo = (unsigned int)cond_ptr;
    offset = cond_lo & 0x7F;
    cond_lo &= ~0x7F;
```

```
    // Cache line align the local stack buffer.
    buf_ptr = (char *)(((unsigned int)(buf) + 127) & ~127);
    cv_ptr = (volatile unsigned short *)(buf_ptr + offset);

    do {
        // Get, with reservation, the cache line containing the condition
    // variable word.
        spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_GETLLAR_CMD);
    (void)spu_readch(MFC_RdAtomicStat);

    // Copy the wait count into the signaled count.
    *(cv_ptr+0) = *(cv_ptr+1);

    // Conditionally put the lock line containing the adjusted
    // signaled halfword.
    spu_mfcdma64(buf_ptr, cond_hi, cond_lo, 128, 0, MFC_PUTLLC_CMD);
    status = spu_readch(MFC_RdAtomicStat) & MFC_PUTLLC_STATUS;
    } while (status);
}
```

# 21. Parallel Programming

Programming the Cell Broadband Engine Architecture (CBEA) processors[1] requires an understanding of parallel programming. Traditional computing platforms contain a single processor, which computes a single thread of control. High-performance computing platforms contain many processors, with potentially many threads of control. Parallel programming has become the default in many fields where immense amounts of data needs to be processed as quickly as possible: oil exploration, automobile manufacturing, pharmaceutical development and in animation and special effects studios.

These widely disparate tasks, and the algorithms for tackling those tasks, showcase the many different styles of parallel programming. Some tasks are data-centric and algorithms for working on them fit neatly into the single instruction, multiple data (SIMD) model. Others are characterized by distinct chunks of distributed programming, and these algorithms rely on good communication models among subtasks. This section reviews different styles of parallel programming on the CBEA processors and describes ways to implement those different strategies. The section covers topics of interest to both application programmers and compiler writers.

## 21.1 Challenges

The key to parallel programming is to locate *exploitable concurrency* in a task. The basic steps for parallelizing any program are:

- Locate concurrency.
- Structure the algorithms to exploit concurrency.
- Tune for performance.

The major challenges are:

- Data dependencies.
- Overhead in synchronizing concurrent memory accesses or transferring data between different processor elements and memory might exceed any performance improvement.
- Partitioning work is often not obvious and can result in unequal units of work.
- What works in one parallel environment might not work in another, due to differences in bandwidth, topology, hardware synchronization primitives, and so forth.

## 21.2 Patterns of Parallel Programming

There are several ways to express parallelism: programming language features (intrinsics and pragmas), languages with explicit support for parallelism (such as Unified Parallel C [UPC] and Java), tools, and application programming interfaces (APIs) for using parallel-programming libraries. But with all these ways, important basic programming structures can be used to create parallel programs with good performance. These programming structures include types of algorithms and ways to synchronize access to data shared among concurrent tasks. This section describes these programming structures.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

### 21.2.1 Terminology

The terms listed in *Table 21-1* commonly appear in textbooks and web sites focused on parallel programming. These terms describe concepts and properties that are important for efficient parallel programming.

*Table 21-1. Parallel-Programming Terminology* (Sheet 1 of 2)

| Term | Definition |
|---|---|
| Amdahl's Law | Potential speedup = 1 / 1 - P, where P is the fraction of code that can be parallelized. |
| Asynchronous | In an asynchronous programming model, different processor elements execute different tasks without needing to synchronize with other processor elements or tasks. |
| Atomic | An atomic operation is uninterruptable. In parallel programming, this can mean an operation (or set of instructions) that has been protected by synchronization methods. |
| Bandwidth | The number of bytes per second that can be moved across a network. A program is bandwidth-limited when it generates more data-transfer requests than can be accommodated by the network. |
| Busy wait | Using a timed loop to create a delay. This is often used to wait until some condition is satisfied. Because the device is executing instructions in the loop, no other work can be done until the loop ends. |
| Critical section | A critical section is a piece of code that can only be executed by one task at a time. It typically terminates in fixed time, and another task only has to wait a fixed time to enter it. Some synchronization mechanism is required at the entry and exit of the critical section to ensure exclusive use. A semaphore is often used. |
| Deadlock | A deadlock occurs when two or more tasks or processor elements are stalled, waiting for each other to perform some action such as release a shared resource. |
| Latency | Latency is a time delay between the moment something is initiated, and the moment one of its effects begins. In parallel programming, this often refers to the time between the moment a message is sent and the moment it is received. Programs that generate large numbers of small messages are latency-bound. |
| Load balance | Distributing work among processor elements so that each processor element has roughly the same amount of work. |
| Monitor | A software monitor consists of:<br>• A set of procedures that allow interaction with a shared resource.<br>• A mutual-exclusion lock.<br>• The variables associated with the shared resource.<br>• A monitor invariant that defines the conditions needed to avoid race conditions.<br>A monitor procedure takes the lock before doing anything else, and holds it until it either finishes or waits for a condition. If every procedure guarantees that the invariant is true before it releases the lock, then no task can ever find the resource in a state that might lead to a race condition. |
| Mutual exclusion | Mutual exclusion (often abbreviated to *mutex*) algorithms are used in parallel programming to avoid the concurrent use of nonshareable resources by pieces of computer code called critical sections.<br>When several processor elements share memory, an indivisible test-and-set of a flag is used in a tight loop to wait until the other processor element clears the flag. This test-and-set ensures that when the code enters the critical region, the flag is set. When the code leaves the critical region, it clears the flag. In a spin lock or busy wait mutex, the wait loop terminates when the test finds that the flag is not set, and the wait continues if the flag is set. |
| Race condition | A race condition occurs when the output exhibits a dependence (typically unexpected) on the relative timing of events. The term originates with the idea of two signals racing each other to influence the output first. The term race condition also refers to an error condition where the output of a program changes as the scheduling of (multiple) processor elements varies. |

*Table 21-1. Parallel-Programming Terminology* (Sheet 2 of 2)

| Term | Definition |
|---|---|
| Semaphore | A semaphore is a protected variable (or abstract data type) and constitutes the classic method for restricting access to shared resources (for example, shared memory) in a parallel programming environment. There are two operations on a semaphore: V (sometimes called up) and P (sometimes called down). The V operation increases the value of the semaphore by one. The P operation decreases the value of the semaphore by one. The V and P operations must be atomic operations. |
| Starvation | Starvation occurs when a task tries to access some resource but is never granted access to that resource. |
| Synchronous | Coordinated in time among tasks, such as when one task notifies a second task about some state change, which the second task receives when it is polling for such notification. |
| Synchronization | Synchronization enforces constraints on the ordering of events occurring in different processor elements. |
| Task | A unit of execution. |
| Thread | A fundamental unit of execution in many parallel programming systems. One process can contain several (or dozens or hundreds) of threads. |

### 21.2.2 Finding Parallelism

The first, vital step in parallelizing any program is to consider where there might be exploitable concurrency. Time spent analyzing the program and its algorithms and data structures will be repaid many-fold in the implementation and coding phase. In other words, by no means immediately start to code a program to take advantage of this or that parallel programming model. Spend time understanding the data flow, data dependencies, and functional dependencies.

The most important question is: Will the anticipated speedup from parallelizing a program be greater than the effort to parallelize a program, which includes any overhead for synchronizing different tasks or access to shared data?

The second question is: Which parts of the program are the most computationally intensive? It is worthwhile to do initial performance analysis on typical data sets, to be sure the hot spots in the program are being targeted.

When you know which parts of the program can benefit from parallelization, you can consider different patterns for breaking down the problem. Ideally, you can identify ways to parallelize the computationally-intensive parts:

- Break down the program into tasks that can execute in parallel.
- Identify data that is local to each subtask.
- Group subtasks so that they execute in the correct order.
- Analyze dependencies among tasks.

In this context, a *task* is a unit of execution that is enough to keep one processor element busy for a significant amount of time, and that performs enough work to justify any overhead for managing data dependencies. Key elements to examine are:

- Function calls
- Loops
- Large data structures that might be operated on in chunks

Questions to ask about data-sharing include:

- Do all processor elements update shared data?
- Does one task need another's data?
- Are there any race conditions, in which data might be read either before or after it is written in another task?
- Is there too much synchronization overhead to justify the parallelism?
- Are interactions among tasks synchronous or asynchronous?
- Are the processing advantages worth the data-communication costs, including latency, bandwidth, and amount of data that needs to be transferred?

Main organizing principles include:

- By tasks (including control flow between tasks and the amount of work in different tasks)
- By data decomposition
- By flow of data

### 21.2.3 Strategies for Parallel Programming

#### 21.2.3.1 *Task Parallelism*

Some programs are embarrassingly parallel. For example, ray-tracing routines fit well into task-parallel designs. Such programs have, inside them, many tasks that are completely independent. These programs can be easily mapped onto standard parallel programming models.

The two essential questions are:

- Are the tasks independent enough that managing dependencies does not consume too much time?
- Can the tasks be load-balanced among the CBEA processor elements?

#### 21.2.3.2 *Data Partitioning*

Some programs operate on data that can be broken down into chunks that can be operated on independently. This approach makes sense when the most computationally-intensive part of the program is centered on manipulating a large data structure. Use this approach when the chunks of data are operated on in the same way. Arrays of data, and recursive data structures like trees, are often good candidates for this style of data decomposition.

#### 21.2.3.3 *Task Grouping*

Sometimes groups of tasks—several sets of operations on data—can be pulled together to form the basis of parallelism. Combining a group of tasks can create enough work to justify assigning it to a single processor element. This approach can also reduce the amount of synchronization between different groups, with each group being scheduled to execute as a single task. This approach also makes sense when one task shares data dependencies with another task. Those tasks might be candidates for grouping.

#### 21.2.3.4 *Divide and Conquer*

A divide and conquer algorithm works by recursively breaking down a problem into two or more sub-problems of the same (or related) type, until these become simple enough to be solved directly. The solutions to the sub-problems are then combined to give a solution to the original problem. This strategy is often used in mathematical problems, such as fast Fourier transform (FFT) or Cholesky matrix decomposition. The parallelism is easy to use: because the sub-problems are solved independently, they can be computed in parallel.

#### 21.2.3.5 *Pipelining*

Sometimes the overall computation involves feeding data through a series of operations. Graphics applications often fall into this category. The key here is to express clearly any ordering constraints—which operations must occur before any other. Pipelining can be expressed at different levels: at the SIMD or vector-processing level, as well as at a higher, algorithmic level. The key to parallelism is assigning each stage of the pipeline to a different processor element, and managing data flow among the processor elements.

#### 21.2.3.6 *Event Parallelism*

Event parallelism concerns a group of independent tasks that are interacting, but in a somewhat irregular fashion. This style of programming often relies on asynchronous communication, in which a task sends an event but does not wait for a response. In a shared-memory system, the program might use a queue to represent message-passing among the tasks. This requires safe, concurrent access to mutex variables. One challenge is avoiding deadlock, in which one task waits for an event that will never arrive. Another challenge is load-balancing the tasks across processor elements.

#### 21.2.3.7 *Master and Subordinate*

The master and subordinate approach is similar to task parallelism. One master processor element might farm out chunks of works to subordinate processor elements and wait for these subordinate tasks to return. This approach works well when you have variable and unpredictable workloads, and simple parallel loops are not enough. The master task is in charge of load-balancing among the several subordinate tasks. A shared queue is a good way to manage the workload. This approach has good scalability if the number of tasks greatly exceeds the number of subordinate processor elements.

#### 21.2.3.8 *Fork and Join*

The fork and join approach works well when the number of parallel tasks varies as the program runs, and structures such as parallel loops are not powerful enough to contain each task's execution. Tasks are created dynamically: that is, they are *forked*. When tasks finish, they return their status to the parent: that is, they are *joined*. A one-to-one mapping of tasks onto processor elements often works well, but this is dependent on the number of tasks and the number of processor elements.

## 21.3 Steps for Parallelizing a Program

Now that the common strategies for parallel programming have been introduced, this section suggests ways in which you can apply these strategies to a program.

### 21.3.1 Step 1: Understand the Problem

- *Performance Objectives*—Is the objective to perform the computation as fast as possible, using all available resources? Or is absolute performance not critical, and is efficiency more important because the algorithm must be run simultaneously with other problems?
- *Type of Computation*—Does the problem contain small kernels of computation that can be partitioned and loaded onto separate Synergistic Processor Elements (SPEs)? Also, because of the heterogeneous architecture, it is important to understand what code sections must be performed on the PowerPC Processor Element (PPE) (such as system calls) or might be faster on a PPE (such as code dominated by unpredictable branches), and what code sections are likely to be faster on an SPE.
- *Data Types*—Is the fundamental data type fixed-point or floating-point? What is the precision—byte, halfword, word, or doubleword? The answers to these questions might determine which computational resources can be used. For example, double-precision floating-point is not supported by the PPE's vector/SIMD multimedia extension instructions, although it is supported by the PPE's standard *PowerPC Architecture* instructions and the synergistic processor unit (SPU) instructions. (The SPEs, however, are optimized for single-precision floating-point operations.)
- *Data-Access Patterns*—Is data randomly accessed or sequentially accessed? If randomly accessed, what is the size of the data accesses? Small data accesses can negatively affect the memory bandwidth so that alternative data organizations might be warranted. Understanding data-access patterns includes understanding what data is read-only, write-once, accessed by a single task in a particular parallel region, or accessed only by a specific subset of tasks.

### 21.3.2 Step 2: Choose Programming Tools and Technology

- *System Configuration*—Is this application going to be deployed on a networked cluster of CBEA processors, on a shared-memory system of CBEA processors, on a single CBEA processor, or on only one (or a few) CBEA processor processing elements?
- *Availability*—What tools are available? Because the Cell Broadband Engine Architecture (CBEA) is a new architecture, new tools are being developed all the time.
- *Quality and Reliability*—Are the tools stable and of high-enough quality for the development needs?
- *Programmer Skills*—Are the programmers experienced in the use of tools and languages? Can the project schedule absorb the learning curve associated with using new programming tools or languages?
- *Schedule Objectives*—How much time is available to develop the software?
- *Performance Objectives*—What are the system performance objectives? Can software be quickly developed now and optimized over time?

### 21.3.3 Step 3: Develop High-Level Parallelization Strategy

The key to maximizing system performance on a CBEA processor is choosing the optimal parallelization strategy. Depending on the system configuration, multiple strategies might be required—for example, a high-level parallelization strategy and a moderate-level parallelization strategy. A high-level strategy is useful when partitioning the problem across systems without shared-memory access. A moderate-level strategy is useful when partitioning the problem across processing elements that have high-speed, shared-memory access.

The key to nearly all successful parallelization strategies is *minimizing synchronizing events* between the computational elements. See *Section 20 Shared-Storage Synchronization* on page 561. One effective method for the CBEA processors is to partition the data for independent computation across parallel SPE threads of execution. Each SPE thread operates as an independent stream processor—streaming in blocks of input data, performing the necessary computation, and streaming out the results. Multibuffering techniques (see *Section 24.1.2* on page 692) can be employed to ensure simultaneous computation and data transfer.

The work (data) can be allocated among the SPEs in one of three methods:

- *Algorithmically Assigned*—For example, if *w* work is partition among *n* SPEs, then each SPE is assigned *w*/*n* amount of work. This solution works well for simple algorithms in which the computation patterns are deterministic.
- *Self-Managed*—Making SPEs autonomously arbitrate for work is effective at load-balancing parallel computation when computation is not easily predictable.
- *Mastered*—In this method, one of the processing elements (typically the PPE) serves as a master distributing work among the available SPEs. Work is communicated to the SPEs either by using a work queue in main storage or by using one of the hardware communication features—mailboxes (*Section 19.6* on page 539) or signal-notification registers (*Section 19.7* on page 551).

There is a set of problems in which an SPE's local storage (LS) is insufficient to hold both the working data set and all the code simultaneously. In these situations, it might be effective to stream the code to the data by sequencing a set of code modules over a single set of LS data.

### 21.3.4 Step 4: Develop Low-Level Parallelization Strategy

Language extensions, like vector intrinsics, and pragmas, like OpenMP's `parallel do`, can be used to direct the compiler in its SIMDization phase. Identify data structures that can be operated on using vector instructions, and annotate the code accordingly. See *Section 22* on page 629 for more information about SIMD Programming, and see *Section 21.4.1* on page 618 for details on programmer-directed SIMDization.

### 21.3.5 Step 5: Design Data Structures for Efficient Processing

After the basic low-level parallelization strategy is identified, the data structures must be developed for efficient processing. This includes:

- *Format*—Organize structures and arrays according to the SIMD format chosen in step 4 (*Section 21.3.4*). Consider padding SOA-formatted data (*Section 22.1.1.1* on page 630) to an even multiple of quadwords, so that partial quadword processing need not be used, thus improving performance.

- *Quadword Align*—Quadword-align arrays and structures to ensure correct alignment for transferring data from main storage to an LS, and to ensure correct alignment for LS loads and stores. Some data structures may also be aligned on a cache-line boundary to maximize bus bandwidth.
- *ABI Differences*—Pointers and long variables may be of different sizes, depending upon whether the PPE is compiled according the 32-bit application binary interface (ABI) or 64-bit ABI. For example, the 32-bit ABI may be ILP32 (that is, 32-bit integers, longs, and pointers). The 64-bit ABI may be LP64 (that is, 64-bit longs and pointers). The SPEs are always ILP32, with DMA support for 32-bit and 64-bit effective-address pointers.
- *Synchronization Variables in Independent Cache Lines*—Place synchronization variables in their own, independent 128-byte cache lines to minimize atomic-reservation thrashing.
- *Optimize Layout*—Tailor layout of shared data structures for parallel computation—that is, to provide spatial locality. For example, structure data so that (as much as possible) data items used locally within a processor are not interspersed with data items shared by multiple processor elements.

When data structures are shared, one or more tasks might modify the data, and one or more tasks might access the modified data. Explicitly managing shared data structures is one of the most error-prone aspects of parallel programming. Try to reduce the size of the critical section. Use a mutual-exclusion protocol to access the shared data: mutex locks, synchronization blocks, semaphores, or nested locks. Consider partially replicating shared data, so that many tasks can update the data without needing to synchronize with one another. But do not prevent concurrent operations on shared data that could safely have taken place at the same time. Also, prevent bottlenecks, in which several tasks are waiting to access the same areas of shared data at the same time.

### 21.3.6 Step 6: Iterate and Refine

The next step is to evaluate system performance and determine the effectiveness of the parallelization strategy. Various performance-analysis techniques can be employed to identify the performance bottlenecks. Based upon these findings, the strategies devised in step 3 (*Section 21.3.3* on page 615) and step 4 (*Section 21.3.4* on page 615) might need to be revisited.

### 21.3.7 Step 7: Fine-Tune

After the basic parallelization strategy is established in the preceding steps, the software should be fine-tuned. Generally, this entails fine-tuning only the critical code sections identified by either programmer knowledge or hot-spot analysis. Programmer knowledge of PPE and SPE programming, as well as previous experience with the code-development tools, will influence the degree of additional fine-tuning that is necessary.

Fine-tuning might include one or more of the following considerations:

- *Branches*—Eliminate branches by exploiting the select-bits instruction (see *Section 24.3.2* on page 700). Reduce mispredicted branches by using feedback-directed optimization, software branch-prediction tables, or programmer-directed branch prediction.
- *Inline*—Inline frequently-called functions, to eliminate call-linkage overhead and provide additional instructions for improved scheduling and latency-hiding. See *Section 24.3.1* on page 700. Most compilers give users control over what and how much to inline.

- *Loops*—Unroll loops to reduce loop iterations and provide additional instructions for improved instruction scheduling. See *Section 24.3.1* on page 700. Loop-unrolling might not prove fruitful for algorithms in which dependencies exist between loop iterations. Be careful of the interaction of manual loop-unrolling with compiler SIMDization; check the results after compiling. Also, loop unrolling might not always be beneficial, due to constrained LS size in an SPE.
- *False Dependencies*—Restructure code to eliminate false-dependency stalls. False-dependency stalls are instruction dependencies that exist only because of the order in which the operations were scheduled.
- *Dual-Issue*—Change instruction mix to improve dual-issuing. Frequently, there are multiple ways an operation can be performed. The multiple solutions might use instructions on different execution pipelines. Choosing the solution that uses instructions executed on an underused pipeline can result in improved performance. Explicit instruction control can be achieved when programming in intrinsics.
- *Memory-Bank Access*—Refine data organization to ensure better memory-bank access patterns. Many algorithms have regular, power-of-two, data-access patterns. Because memory banks are organized on powers of two, this can result in multiple processor elements accessing the same memory banks at the same time. This, in turn, results in less-then-theoretical memory bandwidth. Offsetting data buffers or algorithmically changing the access patterns can improve performance on algorithms that are memory bound.
- *Inefficient Operations*—Eliminate inefficient operations based upon known input constraints. For example, integer multiply can be optimized if one or more of the operands is known to be 16 bits or less. See *Section 24.7* on page 716.

## 21.4 Levels of Parallelism in the CBEA Processors

The CBEA processors provide a foundation for many levels of parallelization. Starting from the lowest, fine-grained parallelization—SIMD processing—up to the highest, course-grained parallelization—networked multiprocessing—the CBEA processors provide many opportunities for concurrent computation. The levels of parallelization include:

- SIMD processing
- Dual-issue superscalar microarchitecture
- Multithreading
- Multiple execution units with heterogeneous architectures and differing capabilities
- Shared-memory multiprocessing
- Networked distributed-memory multiprocessing

The shared-memory model is key to data transfer and program execution across the PPE and the SPEs. Compilers and other utilities can mediate the partitioning of code and data and orchestrate any data movement implied by this partitioning. A variety of compiler techniques can be used to exploit the performance potential of the SPEs and to enable the multilevel heterogeneous parallelism supported by the CBEA processors. For example, see *Section 21.5.3* on page 621 and *Section 21.5.4* on page 623.

### 21.4.1 SIMD Parallelization

*Section 22* on page 629 describes SIMD parallelization techniques. Both the PPE and the SPEs are capable of SIMD computation. Compilers can support programmer-directed parallelism by means of programmer-inserted intrinsics and pragmas. Compilers can also provide auto-SIMDization that generates vector instructions from scalar source code for both the PPE and the SPEs. Auto-SIMDization minimizes the overhead due to misaligned data streams and can be tailored to handle many of the data structures and code structures found in multimedia applications.

### 21.4.2 Superscalar Parallelization

The PPE and SPEs have multiple, parallel execution units and are capable of executing two instructions per clock. The PPE execution units are described in *Section 2.1* on page 52. The SPE execution units are described in *Section 3.1* on page 65. Dual-issue success depends upon the instructions being issued, their address, and the state of the system during execution.

A dual-issue instruction pair consists of two instructions on an aligned doubleword address. *Section A.5* on page 760 describes the dual-issue rules for the PPE's vector/SIMD multimedia extension instructions. *Section B.1.3* on page 779 describes the dual-issue rules for SPE instructions.

### 21.4.3 Hardware Multithreading

The PPE supports two simultaneous threads of execution in hardware, so the PPE can be viewed as a two-way multiprocessor with shared dataflow. This gives PPE software the effective appearance of two independent processing units. The performance of the two threads is limited, however, because they must share resources such as the L1 and L2 caches and they must reside in the same logical partition. For more details on multithreading, see *Section 10 PPE Multithreading* on page 299.

Programmers typically think of multithreading in terms of application programs and the thread libraries they may use. CBEA processor thread libraries can enable multithreading at the application level.

### 21.4.4 Multiple Execution Units

Each of the nine processor elements provides independent computation and can be considered as asymmetric threads of execution. All processor elements have access to the coherent main storage for shared-memory multiprocessing. The SPE mailboxes and SPU signal notification registers support parallel-processing message-passing.

Compilers can parallelize and partition a single source program across the PPE and SPEs, guided by user directives. Compilers can also optimize data transfer, allowing a single SPE to process data that far exceeds the LS's capacity. The SPE can also use code that exceeds the size of LS. Compilers can schedule necessary DMA transfers so that they overlap ongoing computation.

In addition to the nine processor elements, each of the eight memory flow controllers (MFCs) serves as independent DMA controllers capable of supporting up to 16 concurrent transfers for a total of 128 such concurrent transfers. The large number of concurrent transfers, in conjunction with a uniform distribution of memory transfers across both memory channels and all memory banks, makes near-theoretical peak memory throughput possible.

### 21.4.5 Multiple CBEA Processors

Multiple-instruction, multiple-data (MIMD) parallelization can be supported on one CBEA processor or on multiple CBEA processors. This section reviews two common methods of interconnecting multiple CBEA processors.

#### 21.4.5.1 *Shared-Memory Processing*

Two CBEA processors can be directly connected by means of the Cell Broadband Engine interface (BEI) unit, which manages data transfers between the element interconnect bus (EIB) and I/O devices as shown in *Figure 7-1* on page 161. The BEI supports two Rambus FlexIO I/O interfaces. One of the two interfaces supports only a noncoherent I/O Interface (IOIF) protocol, which is suitable for I/O devices. The other interface is software-selectable between the noncoherent interface and the fully coherent Cell Broadband Engine interface (BIF) protocol—the EIB's native internal protocol—which coherently extends the EIB to another device that can be another CBEA processor.

*Figure 7-2* on page 163 shows a 2-way shared-memory processing configuration, and *Figure 7-3* on page 164 shows a 4-way configuration. Configurations with more than two CBEA processors require a BIF-protocol switch, and such switches can be connected together to support very large multiprocessing configurations.

The BIF protocol supports shared, addressable I/O memory. I/O addresses can be optionally translated from effective addresses into real addresses using the I/O address-translation mechanism described in *Section 7.4* on page 176.

The programming of shared-memory systems, whether on one CBEA processor or multiple CBEA processors, is often done using either OpenMP, Pthreads, or UPC. See *Section 21.5.3* on page 621 for details.

#### 21.4.5.2 *Distributed-Memory Processing*

Multiple processors and processor complexes can be loosely clustered in a distributed-memory configuration. Such a system consists of CBEA processors, memory associated with each CBEA processor, and an interconnection network. Because such systems lack shared memory between processors, data that is needed by more than one CBEA processor must be explicitly sent from processor to processor through the interconnection network.

Distributed-memory processing systems are typically programmed in one of three methods:

- Using a proprietary distributed multiprocessing framework.
- Using the Message Passing Interface (MPI) standard. MPI is widely available and was designed for high-performance communication on both massively parallel architectures and clustered distributed-memory systems.

- Using grid computing services, consisting of a set of standards and protocols. For example, the Open Grid Services Architecture (OGSA) enables communication across heterogeneous, geographically dispersed environments.

## 21.5 Tools for Parallelization

Given an understanding of the patterns and steps of parallelization, and an understanding of what parallelism is available on a CBEA processor, how can parallelism be implemented efficiently? This section describes some tools for creating parallel programs.

### 21.5.1 Language Extensions: Intrinsics and Directives

A broad set of intrinsic extensions for C and C++ is defined in the *C/C++ Language Extensions for Cell Broadband Engine Architecture* specification and summarized in *Appendix B.2* on page 784. Compiler intrinsics function as fast shorthand for expressing assembly-level instructions. Compilers that support the intrinsics will likely also support vector data types. Each data type in C or C++ can be prepended by the keyword, `vector`, to form a 128-bit vector of two doublewords, four words, eight halfwords, or 16 bytes.

Intrinsics are very useful when you are manually SIMDizing your code. If you are coding at this level, be sure to place small structures that must be accessed by mutual exclusion in their own cache line, so they can be operated on by the intrinsic equivalents of the MFC atomic commands, **getllar**, **putllc**, **putlluc** and **putqlluc** (*Section 19.2.1 DMA Commands* on page 516). However, it might not be appropriate to use intrinsics in all cases. For example, if you are relying on the compiler to perform SIMDization, it is best to let the compiler select the appropriate assembly-level instructions.

The major advantage of using intrinsics is that they provide full control over the handling of data alignment and the selection of SIMD instructions, yet still provide the benefits from the compiler's other optimizations such as loop unrolling, function inlining, scheduling, and register allocation. In addition, you can normally rely on the compiler to generate scalar code, such as address, branch, and loop code. Such code is error-prone when coded manually in assembly language. Furthermore, it is much easier to modify code that uses intrinsics than it is to maintain assembly code.

In addition to the hundreds of intrinsics at your disposal, a compiler may provide the special directives that are summarized in *Appendix B.2.5* on page 791. For example, branch prediction can be helped with feedback-directed optimization:

```
int __builtin_expect(int exp, int value)
```

This instructs the compiler that `exp` is expected to equal `value`.

The `__align_hint` directive enables compiler analysis and optimization for pointers. It provides compilers with information that is needed for auto-SIMDization:

```
__align_hint(ptr, base, offset)
```

Here, `ptr` points to data with the base alignment of `base` and an offset from `base` of `offset`.

The most important tool for efficient, highly optimized low-level assembly code is your compiler. See its documentation for a description of compiler optimizations that you can enable for your program.

### 21.5.2 Compiler Support for Single Shared-Memory Abstraction

Manually partitioning code and data among the PPE and the SPEs, and manually managing transfers of code and data between main storage and SPEs' LSs, are great burdens on a programmer. A programmer is accustomed to viewing a computer system as possessing a single addressable memory in which all program data reside.

In the CBEA processors, the LSs, which are directly addressable by their respective SPUs, are memories separate from the vastly larger main storage. Each SPE can initiate DMA transfers between its LS and main storage. The compiler, under certain conditions, can also initiate DMA transfers by explicitly inserting DMA commands. Moreover, there are many optimizations that the compiler can perform to optimize these data transfers, especially when memory references are regular. The compiler can attempt to abstract the concept of separate memories by allocating SPE program data in main storage and automatically managing the movement of this data between its home location and a temporary location in an LS. If this movement is done to satisfy the demands of an executing SPE program, and if the resulting buffers are organized in such a way as to permit reuse, this is referred to as a *software cache*. See *Section 21.5.4* on page 623 for details about software caches.

### 21.5.3 OpenMP Directives

Many compilers support some set of OpenMP directives. OpenMP is a parallel programming model for shared-memory systems. Pioneered by Silicon Graphics, Inc. (SGI), it is now embraced by most of the computing industry and is exhaustively documented at http://www.openmp.org. The book *Parallel Programming in OpenMP*, by Chandra et al., is another good source, as is *Parallel Programming with OpenMP* at http://www.osc.edu/hpc/training/openmp.

OpenMP directives are instructions to the compiler to perform tasks in parallel. In Fortran, they are expressed as source-code comments; in C and C++, they are expressed as a `#pragma`. All OpenMP directives in C and C++ are prefixed by, `#pragma omp`. A typical OpenMP directive looks like this one, which says that individual iterations of the loop are independent and can be executed in parallel:

```
#pragma omp parallel for
    for (int i = 0; i < 1000; i++)
        big_calc(x[i]);
```

OpenMP includes a small set of runtime library routines and environment variables. These routines are designed to be thread-safe. The environment variables can be used to control thread behavior such as spin, yielding, and thread binding. Also, each thread created by OpenMP has its own stack that is off the heap. The compiler often has flags to control the default stack size for each thread. Without this functionality, a program might easily run out of memory.

The basic execution model for OpenMP is the fork and join model. OpenMP provides two parallel control constructs. The first control construct is the `parallel` directive. It encloses a block of code and causes new threads of execution to be spun off to execute that block of code. The second control construct is a mechanism that divides work among a group of already-existing parallel threads.

Variables can be declared `shared`, which means they have a single storage location during the execution of a parallel section. In contrast, a variable can be declared to be `private`. It will have multiple copies: one within the execution context of each thread. There are different flavors of `private`, including `threadprivate`, `firstprivate`, and `lastprivate`. The latter two include an element of shared storage behavior. Finally, a variable can be declared to be a `reduction` variable. Reduction variables have both private and shared storage behavior. A common example is the final summation of temporary local variables at the end of a parallel construct.

OpenMP provides mutual exclusion and event synchronization as ways for multiple OpenMP threads to communicate with each other. The directives `critical` and `barrier` can be used.

OpenMP allows programmers to specify loops, loop nests, or code sections that may be executed in parallel. This, of course, relies heavily on the programmer doing the initial data and task analysis, to make sure that sections marked for parallelization are safe to parallelize. Although OpenMP is a shared memory model, it can also be implemented for distributed memory systems. In the case of the CBEA processors, an SPE's LS is not shared memory, but a compiler can simulate shared memory for SPEs using a combination of a software-managed cache (*Section 21.5.4* on page 623) and DMA-transfer optimizations (*Section 24.1.2* on page 692).

An OpenMP implementation typically uses a runtime library. The runtime library includes functions for initialization, work distribution, and synchronization of data, as well as control flow. The compiler can insert calls to runtime library functions appropriate to the OpenMP directives contained in the source code. The OpenMP *master thread* runs on the PPE and uses the runtime library to distribute work to the SPEs. The master thread itself partakes in all work-sharing constructs. This thread also handles all operating system service requests. The PPE runtime library can include the facility to spawn new SPE threads and terminate them. When a new SPE thread is spawned, it can continuously loop, waiting for the PPE to assign work items to it. A work item specifies a handle that determines the function to execute, any input parameters, and if needed, the current address of the stack in main storage.

Such a runtime library requires communication between the PPE and the SPEs to coordinate execution. SPEs can use explicit DMA commands to read work items assigned to them from a circular queue that is shared with the PPE. The *Cell Broadband Engine Architecture* also includes more efficient communication channels in the form of signal notification and mailbox facilities. The PPE can use asynchronous signals to inform an SPE that there is work available, or that it should terminate. The SPEs can use the mailbox facility to update the PPE on the status of their execution.

Several factors complicate the implementation of OpenMP on the CBEA processors. There are two distinct instruction sets. Certain code sections may be executed on both the PPE and the SPEs. Because the SPEs are not designed to run an operating system, certain code sections such as system calls must execute on the PPE. The compiler must identify data in main storage that is accessed in SPE code sections, and it must insert DMA commands at appropriate points to transfer this data to and from the SPE LSs. The SPE LSs are limited to 256 KB, and this space

contains both the code and data that is executing on the SPE. The compiler might need to split SPE code into multiple partitions if it does not fit in the limited SPE LS (see *Section 14* on page 397 for some alternatives).

The challenge for the compiler is to efficiently manage all this complexity while still enabling the performance potential of the CBEA processors. To do so, the compiler must:

- Minimize the amount of data that needs to be transferred using DMA commands.
- Align some data in main storage for efficient DMA transfers.
- Issue DMA commands early in the code to overlap computation and communication.
- Carefully choose the size of code partitions and individual data transfers.

For details about OpenMP in the context of SIMD programming, see *Section 22.4.1* on page 674.

Here are some things a programmer can do to get better performance when using a compiler that supports OpenMP:

- Choose intelligent values for OpenMP attributes like scheduling and number of threads.
- Minimize use of pointers in parallel code, to enable more precise compiler analysis.
- Lay out shared data structures carefully, paying attention to locality of data in the context of parallel execution.
- Use alignment and other intrinsics to guide compiler analysis.

### 21.5.4 Compiler-Controlled Software Cache

An SPE's local storage (LS), which is filled from main storage using software-initiated DMA transfers, can be regarded as a software-managed cache (or simply *software cache)*. Although most processors reduce latency to memory by using hardware caches, an SPE uses its DMA-filled LS. The approach provides a high degree of control, but it is advantageous only if the DMA transfer-size is sufficiently large and the DMA command is issued well before the data is needed, because the latency and instruction overhead associated with DMA transfers exceeds the latency of servicing a cache miss on the PPE.

When compiling SPE code, the compiler can identify references to data in main storage that have not been optimized using explicit DMA transfers, and the compiler can insert code to invoke the software-managed cache mechanism before each such reference. The call to the software cache takes a main-storage address and returns an address in the SPE LS that contains the corresponding data. The SPE code is generated by a compiler code generator that is different from the one that generates PPE code. Symbols that refer to main storage will have space allocated to them when the PPE code is compiled and linked. The SPE code does not have access to these addresses until it is linked with the PPE code at the very end. During compilation, the compiler can create new symbols—local table of contents (TOC) entries—in the SPE code as placeholders for the main storage addresses. The correct values for these symbols can be filled in when the PPE and SPE code sections are linked together.

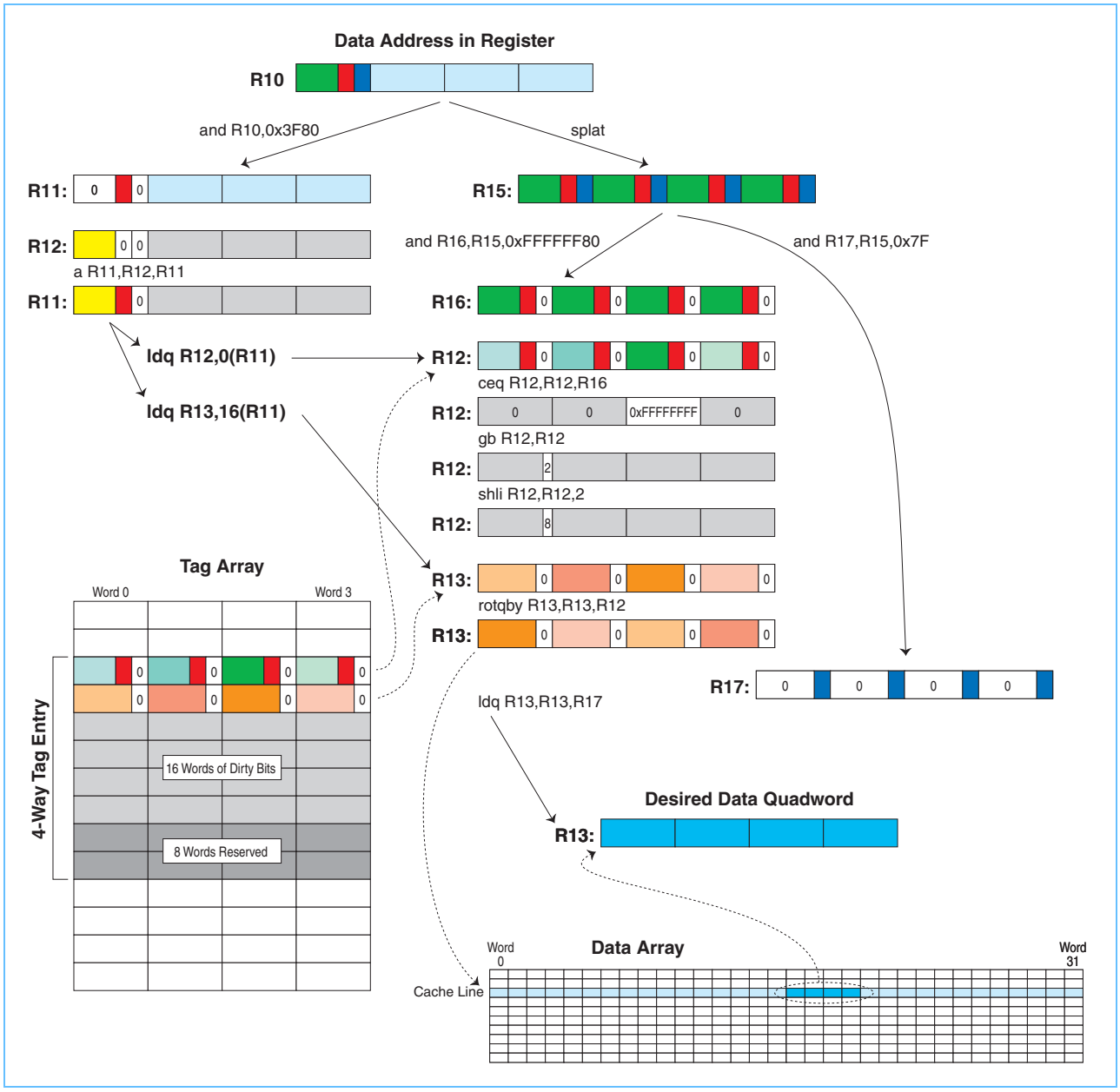
To support the software cache, a separate directory is maintained in each SPE. The compiler, using interprocedural analysis, replaces a subset of load and store instructions with instructions that explicitly look up the effective address of the data in a directory. If the directory lookup indicates that the data is present in the LS, the address of the requested variable cached in the LS is computed and the load or store proceeds using this LS location. Otherwise, a miss-handler

subroutine is invoked and the requested data is transferred from main storage. As is typical in hardware caches, the directory is updated and typically another line is selected for eviction, to make room for the new data.

*Figure 21-1* on page 625 illustrates a software-cache lookup. Starting with the address of the required variable in the left-most (or preferred) slot of an SPE register, instructions are executed to produce the address of the appropriate directory entry in the Tag Array. This task consists of masking out all bits in the address except for those that are used to index into the Tag Array. The cache-line size for the CBEA processors is 128 bytes. By also making the directory entry 128 bytes long, the index obtained after masking does not need to be shifted right and multiplied, and it can be used directly as an offset into the Tag Array.

This example shows a 64 KB 4-way cache, so there are also 128 lines in each way, and the index is thus also 7 bits. To this offset we add the base address of the Tag Array, and using this address we load two consecutive quadwords from the array. The first contains four tags, for the four ways of the cache, and the second contains the LS addresses of the four lines in the Data Array. Independent of this computation, the data address is replicated in all four slots of another register. This register is then masked twice to produce, in the first instance, four identical tags, and in the second, a register containing the offset in the line of the required data. These two instruction sequences are independent and may be scheduled for parallel execution. We now do a SIMD compare of the loaded tags and the replicated tag from our data address. If any of the comparisons contains the tag we are seeking, then the 128-bit result of the compare will be nonzero, and we can use this to test for a hit. If the result is zero, we have a miss and the miss handler (not shown in the figure) is invoked.

*Figure 21-1. Software-Cache Lookup*




The call to the miss handler need not be expanded by the compiler until very late, so it does not appear as a call to the compiler optimizations. This allows the common case, which we assume to be a hit, to suffer no penalty due to the call setup, but it also means that the miss handler cannot use the system-register conventions. If there is a hit, as shown in *Figure 21-1*, we can use the result of the compare to determine the rotation needed to place the correct line address in the preferred slot, making it ready to be used as a LS address. Finally, we add the line offset to the line address, and we load the required data. Store processing proceeds in much the same way,

but additional instructions are required to set the dirty bits. These bits, which are stored in the Tag Array entry, are used to ensure correct execution when multiprocessing (“correct execution” here means ensuring that updates are correctly reflected in main storage).

It takes roughly 12 instructions to process a cache hit, and a similar number to set dirty bits for a store, but because some of these instructions are dual-issued, and there is often other independent work that can be simultaneously scheduled, the cost in cycles is not high. Still, it is clearly important to attempt to reduce the number of cache lookups inserted when compiling a program.

This software cache can be used for either a serial or a parallel program. There is some additional cost involved in supporting parallel execution, because we must keep track of which bytes in a cache line have been modified to support multiple writers of the same line. If the compiler knows that no data will be shared between SPEs (for example, in the case of a serial SPE application), the additional cost of maintaining these dirty bits can be avoided.

Because the compiler must deal with the potential for aliases between cached data and data obtained by other means (for example by means of explicit prefetch in the compiler’s tiler), the miss handler is required to take more care in choosing a line to be evicted than is normally the case with a hardware implementation.

### 21.5.5 Compiler and Runtime Support for Code Partitioning

Because the limited LS of each SPE must accommodate both code and data, there is always a possibility that a single SPE object will be too large to fit in the LS. Compilers can provide a code-partitioning technique to reduce the impact of the LS limitations on the program’s code segment. This approach can be used standalone with an SPE compiler or by programmers choosing to manually partition their algorithms. When using OpenMP with single source compilation, the code partitioning can be integrated with the data software cache to allow for the execution of large functions with large data sets to run seamlessly across multiple SPEs.

In one approach to code partitioning, the compiler is used to automatically divide the SPE program into multiple partitions. The basic unit of partitioning is a function. Just as with data in the software-cache approach, the home locations of code partitions are in main storage. These SPE code partitions are overlaid during linking. That is, they are all assigned to the same starting virtual address. SPE code can then fit into a virtual-address space equal to the size of the largest code partition. Because the compiler determines how code is partitioned, it controls the partition size, and thus controls the amount of space used in the LS for code.

The overlaid partitions cannot execute at the same time. This implies that if code in one partition calls a function in another partition, the two partitions need to be swapped in and out of the LS at the point of the function call and return. To run a partitioned program, such partition transitions must be handled properly, and this can be done collaboratively by the compiler and a runtime partition manager.

When the compiler partitions SPE code, it can also reserve a small portion of the SPE LS for the runtime partition manager. The reserved memory can be divided into two segments: one to hold the continuously resident partition manager, and the other to hold the active code partition. The code partitions need to be relocatable, which implies that function calls should not use absolute addressing. The partition manager is responsible for loading partitions from their home location in main storage into LS during an interpartition function call or an interpartition return. The compiler

can modify the original SPE program to replace each interpartition call with a call to the partition manager. The partition manager also modifies the return address on the stack before branching to the called function to ensure that control returns to the partition manager first.

When an interpartition call is directed through the partition manager, a function pointer and arguments to the function are passed to the partition manager. The partition manager must determine which partition contains the called function. For this purpose, the compiler can assign an index to all partitions it creates, and it can encode the corresponding index in the function pointer that is passed back to the partition manager. An SPE pointer is 32 bits, but because the LS is 256 KB, only 18 bits are used. Of the 14 unused bits, 13 bits can be used for the partition index (the most-significant bit can be used to indicate special handling for calls to certain library functions).

For example, an interpartition call to function `foo` in partition 3 can be transformed from `foo(arg1, arg2, ...)` to `call_partition_manager(3<<18 | foo, arg1, arg2, ...)`. The partition manager fetches the correct partition using the partition index, and transfers control to the proper location within the partition by using the lower 18 bits of the function pointer.

### 21.5.6 Thread Library

Most CBEA processor software systems include a thread library, and many such libraries follow the Portable Operating System Interface (POSIX) definition, documented as *POSIX threads explained* at http://www-128.ibm.com/developerworks/linux/library/l-posix3/index.html. A thread library can be used to create, manage, and terminate threads of execution, manage access to shared data using mutexes and condition variables, and distribute and manage workloads among threads.

Although programming directly with threads is powerful, it puts a burden on the programmer. You are responsible not only for protecting access to all shared data, but also partitioning and allotting chunks of work among threads. You need to explicitly manage a fork and join model or some other programming model. A program that explicitly uses threads can be a challenge to debug, unless you have access to a debugger that has been taught to track threads of execution.

Excellent books are available on this subject. *Programming with Threads*, by Steve Kleiman is a good place to start. *Threads Primer: A Guide to Multithreaded Programming*, by Bill Lewis is a gentle introduction. There are also many good Web sites.

# 22. SIMD Programming

## 22.1 SIMD Basics

Both the PowerPC Processor Element (PPE) and the Synergistic Processor Elements (SPEs) support parallel processing of Single Instruction Multiple Data (SIMD) vector elements. A *vector* is an instruction operand containing a set of data elements packed into a one-dimensional array, as shown in *Figure 22-1*. In the Cell Broadband Engine Architecture (CBEA) processors[1], the vector elements can be fixed-point (integer) or floating-point values. Almost all vector/SIMD multimedia extension and synergistic processor unit (SPU) instructions operate on vector operands. Vectors are also called *SIMD operands* or *packed operands*.

*Figure 22-1. A Vector with Four Elements*


As the SIMD name implies, this style of programming allows one instruction to be applied to the multiple data elements of a vector in parallel. In this way, SIMD processing exploits data-level parallelism. SIMD programming is prevalent in multimedia, graphics-intensive stream processing (such as gaming), and high performance computing—basically, in any compute-intensive application. The CBEA processors are designed to operate efficiently on SIMD code, so programmers benefit from learning how to efficiently exploit data parallelism in their programs and how to take advantage of compiler optimizations for SIMD code.

Support for 128-bit-wide SIMD operations is pervasive in the CBEA processors. In the PPE, these operations are supported by the 32-entry vector register file, vector/SIMD multimedia extensions to the PowerPC instruction set, and C/C++ intrinsics for the vector/SIMD multimedia extensions, as summarized in *Section 2* on page 51. In the SPEs, SIMD operations are supported by the 128-entry vector register file, SPU instruction set, and C/C++ intrinsics, as summarized in *Section 3* on page 65. In both the PPE and SPEs, the vector registers hold multiple data elements as a single vector. The data paths and registers supporting SIMD operations are 128 bits wide, corresponding to four full 32-bit words. This means that four 32-bit words can be loaded into a single register, and, for example, added to four other words in a different register in a single operation. Similar operations can be performed on vector operands containing 16 bytes, 8 halfwords, or (for the SPEs) 2 doublewords.

In a traditional vector processor, multiple vector elements (typically up to 64) are processed one after another, in a pipelined fashion. Registers are very wide—each one holding, for example, 64 vector elements of 64 bits each—but there are only a small number of such registers (for example, eight). These processors have elaborate memory instructions, such as *gather data* and *scatter data*, that use a vector register as an index into another array. For example, two vector loads can be used to fetch the values of `a[b[0...63]]`, as detailed below:

```
r1 = vector load b[0..63];
r2 = indirect vector load into array of a, using r1 as index
```

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

In contrast, a SIMD processor, such as an SPE, stores fewer vector elements in each register but it has far more of such registers. The memory unit of a SIMD processor is more similar to the memory unit of a scalar load-store processor than to the memory unit of a traditional vector processor. The most important distinction is that, in the newer SIMD processors, memory accesses must be aligned on 16-byte boundaries, and loads and stores operate only on 16 bytes at a time. Architectures that support only loads and stores that are aligned to the vector-register length, such as the CBEA processors, are called architectures with *alignment constraints*. An *aligned reference* means that the required data resides at an address that is a multiple of a vector register's 16-byte width.

This section describes some basic techniques of SIMD application programming and the reasons behind it—such as techniques used by compilers to extract parallelism from SIMD application programs. The material is intended for both application programmers and compiler writers. Some material is of particular interest to compiler writers. Nevertheless, an understanding of the compiler-related concepts is very useful for application programmers who want to get the best performance from their programs.

### 22.1.1 Converting Scalar Data to SIMD Data

Depending on performance requirements and code size constraints, advantages can be gained by properly grouping the data that represent the elements of SIMD vectors.

When scalar code for an SPE is fed to a compiler, the compiler's back-end code generator will typically expand a scalar load as read-aligned and a scalar store as a read-modify-write code sequence. This results in large code size with less than optimal performance. When the code is SIMDized (manually by the programmer or automatically by the compiler), such expansion of scalar code does not occur. The expansion of scalar code also does not occur in the case of the PPE, because the PPE has full-featured scalar units

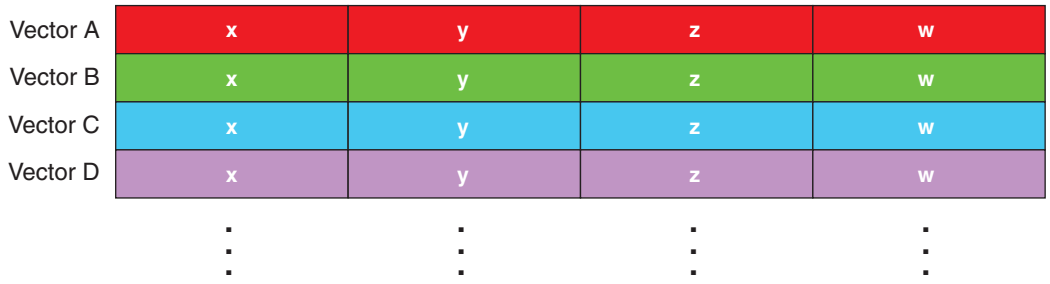
#### 22.1.1.1 *Organizing Data*

Although there are several methods of organizing data, two forms are commonly used for SIMD data: an *array-of-structures (AOS)* form and a *structure-of-arrays (SOA)* form. Both are described in the following sections. For explanation purposes, 3-D homogeneous vertex coordinates will be used as the basic data type. Such a vertex consists of four floating-point components, x, y, z, and w.

*Array Of Structures*

The AOS form—also known as the *vector-across* or *vec-across* form—is shown in *Figure 22-2* on page 631. In this form, the vertex coordinates are structures, and several such structures represent an array of vertices. All components of a single vertex have the same color in *Figure 22-2*.

*Figure 22-2. Array-of-Structures (AOS) Data Organization*




In the AOS form, a set of vertices can be represented as an array of `structs`:

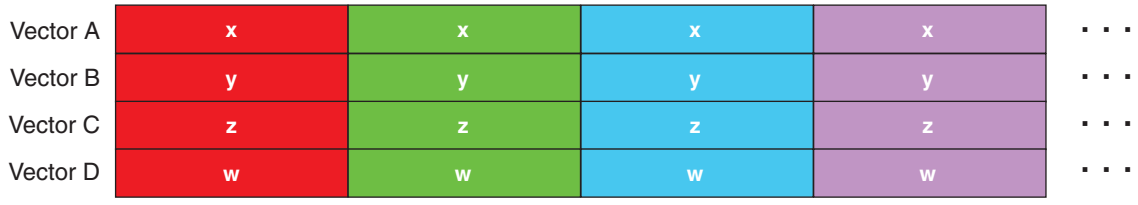
```
union {
    struct _coordinate{
        float x, y, z, w;
    } coordinate;
    vector float vertex;
} array[];
```

This `union` specifies four single-precision floating-point values `x`, `y`, `z`, and `w`, which can be treated either as individual components, accessible one at a time, or as part of the packed vector, named `vertex`, accessible in a single vector operation.

*Structure Of Arrays*

The SOA form—also known as the *parallel-array* form—is shown in *Figure 22-3*. In this form, which is orthogonal to the AOS form, the components of the vertices are separated into independent arrays. These arrays are arranged so that each vector contains a single component from four independent vertices. As in *Figure 22-2*, all elements of a single vertex have the same color in *Figure 22-3*.

*Figure 22-3. Structure-of-Arrays (SOA) Data Organization*




In the SOA form, a set of vertices can be represented as a `struct` of arrays:

```
struct {
    float x[];
    float y[];
    float z[];
```

```
    float w[];
} vertices;
```

If data is stored in the AOS form, the **spu_shuffle** intrinsic can be used to transform it to the SOA form, and vice versa.

*Converting Between AOS and SOA Organization*

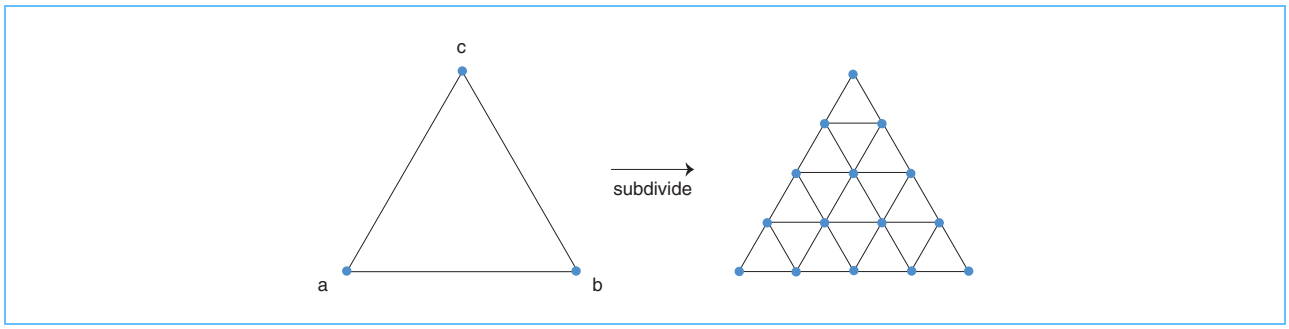
Converting between the AOS vectors (A, B, C, D) shown in *Figure 22-2* on page 631 and the SOA vectors (A, B, C, D) shown in *Figure 22-3* entails performing a 4 × 4 matrix transpose. An efficient transpose can be accomplished using only two shuffle patterns, as follows:

```
vector unsigned char hi = ((vector unsigned char){
0x00,0x01,0x02,0x03,0x10,0x11,0x12,0x13,0x04,0x05,0x06,0x07,0x14,0x15,0x16,0x17});

vector unsigned char lo = ((vector unsigned char){
0x08,0x09,0x0A,0x0B,0x18,0x19,0x1A,0x1B,0x0C,0x0D,0x0E,0x0F,0x1C,0x1D,0x1E,0x1F});

tmp1 = spu_shuffle(A, C, hi);
tmp2 = spu_shuffle(A, C, lo);
tmp3 = spu_shuffle(B, D, hi);
tmp4 = spu_shuffle(B, D, lo);

A = spu_shuffle(tmp1, tmp3, hi);
B = spu_shuffle(tmp1, tmp3, lo);
C = spu_shuffle(tmp2, tmp4, hi);
D = spu_shuffle(tmp2, tmp4, lo);
```

### 22.1.1.2 *Triangle-Subdivision Example*

As an example of SIMDizing an algorithm, consider subdividing a single triangle into multiple smaller triangles, as shown in *Figure 22-4*. This is known as *triangle subdivision*. The original triangle is defined by the three vertices: `a`, `b`, and `c`. Two levels of subdivision results in a subdivided triangle consisting of 16 smaller triangles defined by 15 vertices.

*Figure 22-4. Triangle Subdivision*



For this algorithm, there are several methods of exploiting SIMD parallelization in the generation of the subdivided vertices. Reviewing these methods should inspire programmers to consider alternate, nontraditional, methods that can result in more efficient solutions, in both development time and execution performance.

*Method 1: Evaluate One Vertex at a Time Using AOS*

In the AOS data-organization form (*Figure 22-2* on page 631), the coordinates for each vertex are packed into a vector. The three vertices that describe the original triangle in *Figure 22-4* on page 632 can be represented as three vectors of four elements each. Each vertex is then evaluated one at a time. Performance can be improved by unrolling the loops and evaluate several vertices at a time. The additional instructions within the loop can be interleaved, manually or by the compiler, so that dependent-instruction latency can be hidden.

Although the AOS form is a natural fit for 3-D vertices exploiting SIMD methods, it is typically not the most efficient. Efficiency is compromised when the 3-D vertices have only three components. In addition, geometric operations such as dot products and cross products are not ideal. Dot products on the AOS form generate scalars, and cross products require vertex-component reordering.

*Method 2: Evaluate One Vertex at a Time Simultaneously for Four Triangles Using SOA*

Another way to subdivide the triangle is to evaluate one subdivision vertex at a time, over four independent triangles. In this approach, which uses the SOA data-organization form (*Figure 22-3* on page 631), each vector is populated with *independent* data, so a program can be developed as though the algorithm were operating on scalar data. Each vertex component—for example, the `x` component of each of four triangle vertices—is stored in a single vector.

Think of the data as if it were scalar, and the vectors were populated with independent data across the vector. This differs from the AOS form, in which the four values of each vertex are stored in one vector. *Figure 22-3* on page 631 shows the use of SIMD vectors to represent the four single-precision floating-point values `x`, `y`, `z`, and `w` for one vertex of four triangles. Not only are the data types the same across the vector, but now their data interpretation is the same. Depending on the algorithm, software might execute more efficiently with this SOA data organization than with the AOS organization shown in *Figure 22-2* on page 631.

Assuming that the algorithm can be independently and simultaneously performed on multiple data objects, this approach often is the most efficient.

*Method 3: Evaluate Four Vertices at a Time Using SOA*

A third approach is to evaluate four vertices at a time for a single subdivided triangle. In this case, the vertex component data is replicated across the vectors. Unique weighting factors are maintained in each element of the vector, so that four subdivided vertices can be computed in parallel. However, inefficiency can result when the number of subdivision vertices is not a multiple of four. In the case of the two levels of subdivision shown in *Figure 22-4* on page 632, three iterations of four vertices and one iteration of three vertices are performed.

### 22.1.2 Approaching SIMD Coding Methodically

Two of the principal objectives of SIMD programming are:

- Extract parallelism from both loops and basic blocks.
- Satisfy the constraints of data alignment[2], data-size conversions, and data stride[3].

These objectives can be achieved in both SPE code and PPE vector/SIMD multimedia extension code.

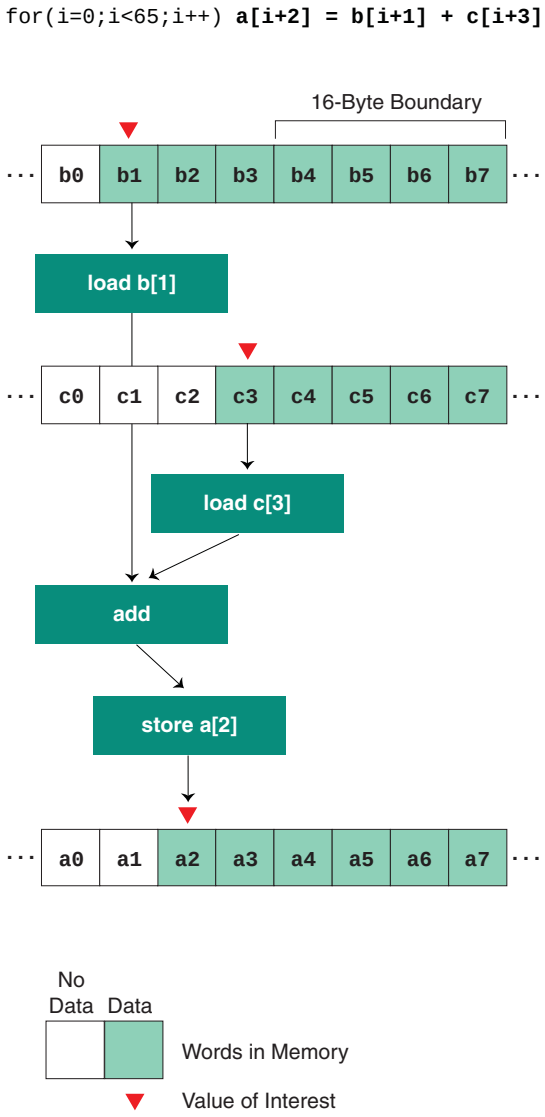
Consider the execution of the following loop:

```
for (i=0;i<65;i++) a[i+2] = b[i+1] + c[i+3]
```

This loop can be performed in a traditional sequential way as shown in *Figure 22-5* on page 635. Here, the first pair of 32-bit operands, `b[1]` and `c[3]`, are loaded from their respective addresses in memory using standard (nonvector) load instructions. Then, they are added together, and the result, `a[2]`, is stored using a standard (nonvector) store instruction. This produces a correct result, but the operands and the result are offset from one another, relative to the alignment of vectors of which they could be a part, so the method cannot be SIMDized without some realignment of the operands.

2. Alignment on 16-byte boundaries.
3. A *data stride* is a memory access pattern in which each element in a list is accessed sequentially.

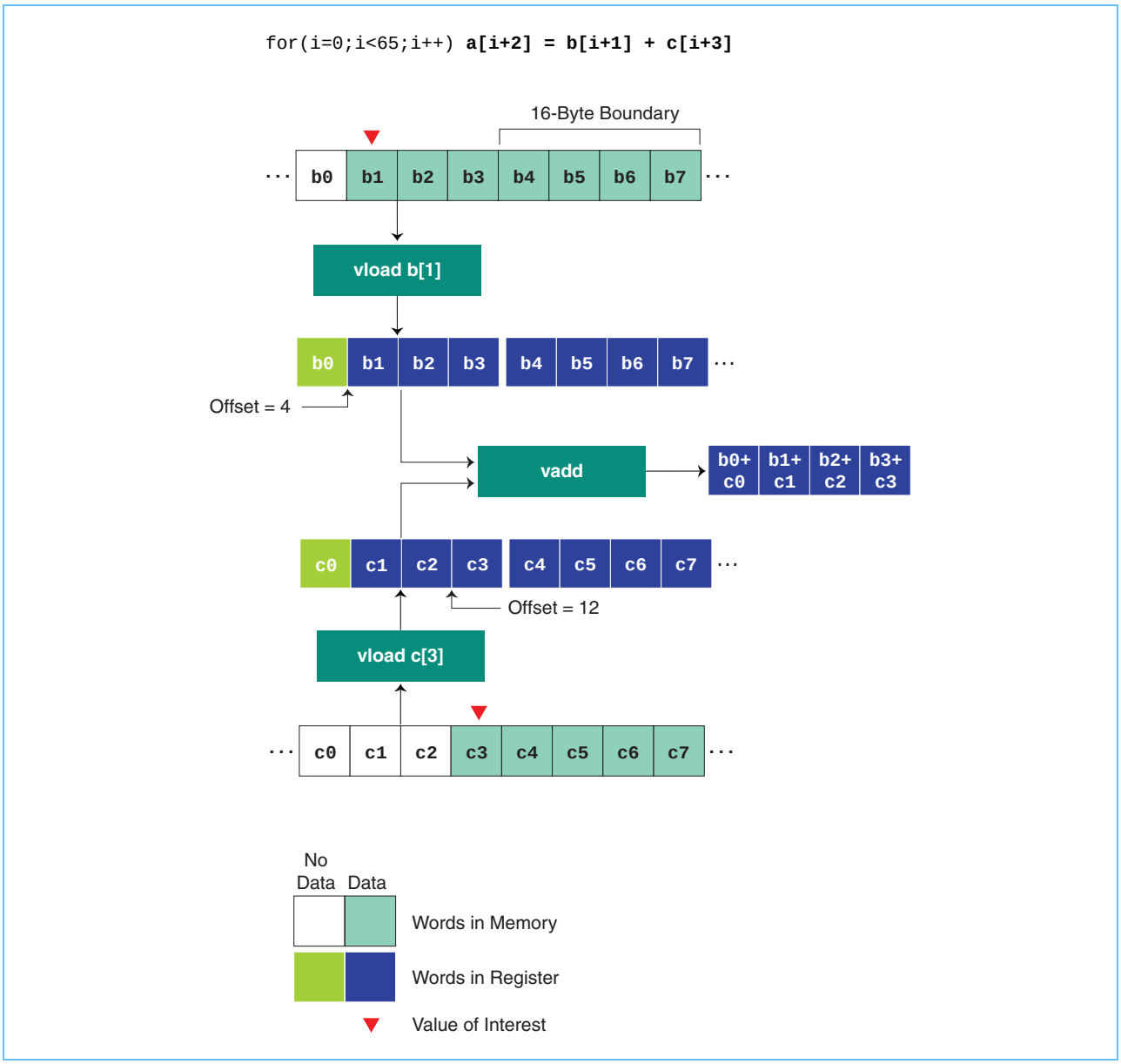
*Figure 22-5. Traditional Sequential Implementation of a Loop*



### 22.1.2.1 *Simple, Incorrect SIMDization*

A simple, but incorrect, SIMD implementation of the loop is illustrated in *Figure 22-6*. Here, vector-load instructions (called **vload** in this example) load the complete four-element vectors containing the `b[i]` and `c[i]` words. Then, a vector-add instruction sums the four pairs of vector elements in parallel. However, the result is not correct because the alignments of the `b[i]` and `c[i]` vector elements within the vector registers do not match—the `b[1]` element is added to `c[1]` rather than to `c[3]`.

*Figure 22-6. Simple, Incorrect SIMD Implementation of a Loop*

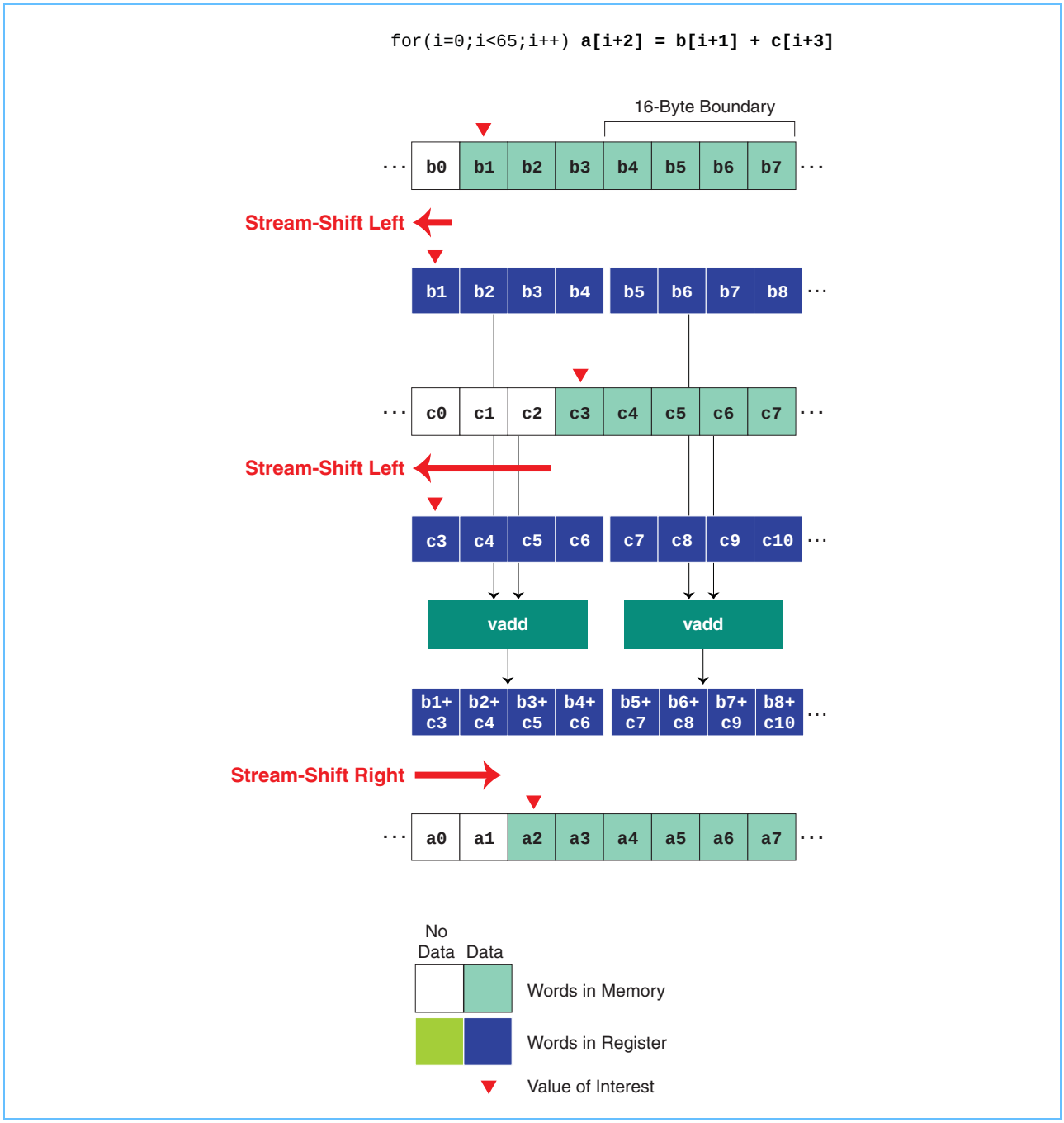


#### 22.1.2.2 *Simple, Correct SIMDization*

A simple, and correct, SIMD implementation of the loop is illustrated in *Figure 22-7* on page 638. It uses operations, called *stream-shift operations*, which shift memory loads or stores the right or left to align vectors within registers. A *stream* is a sequence of contiguous memory locations that are accessed by a memory reference throughout the lifetime of a loop (also called a *memory stream*), or a sequence of contiguous register values that are produced by an operation over the lifetime of a loop (also called a *register stream)*.

In the *Figure 22-7* example, each memory stream is shifted left, by different amounts, to place the vector elements, `b[1]` and `c[3]`, at corresponding register locations. A vector-add instruction sums the four pairs of elements in parallel, and these additions are repeated for subsequent vectors. A final stream-shift operation, to the right, is used prior storing the results so that the first result, `a[2]`, is stored in the expected memory location.

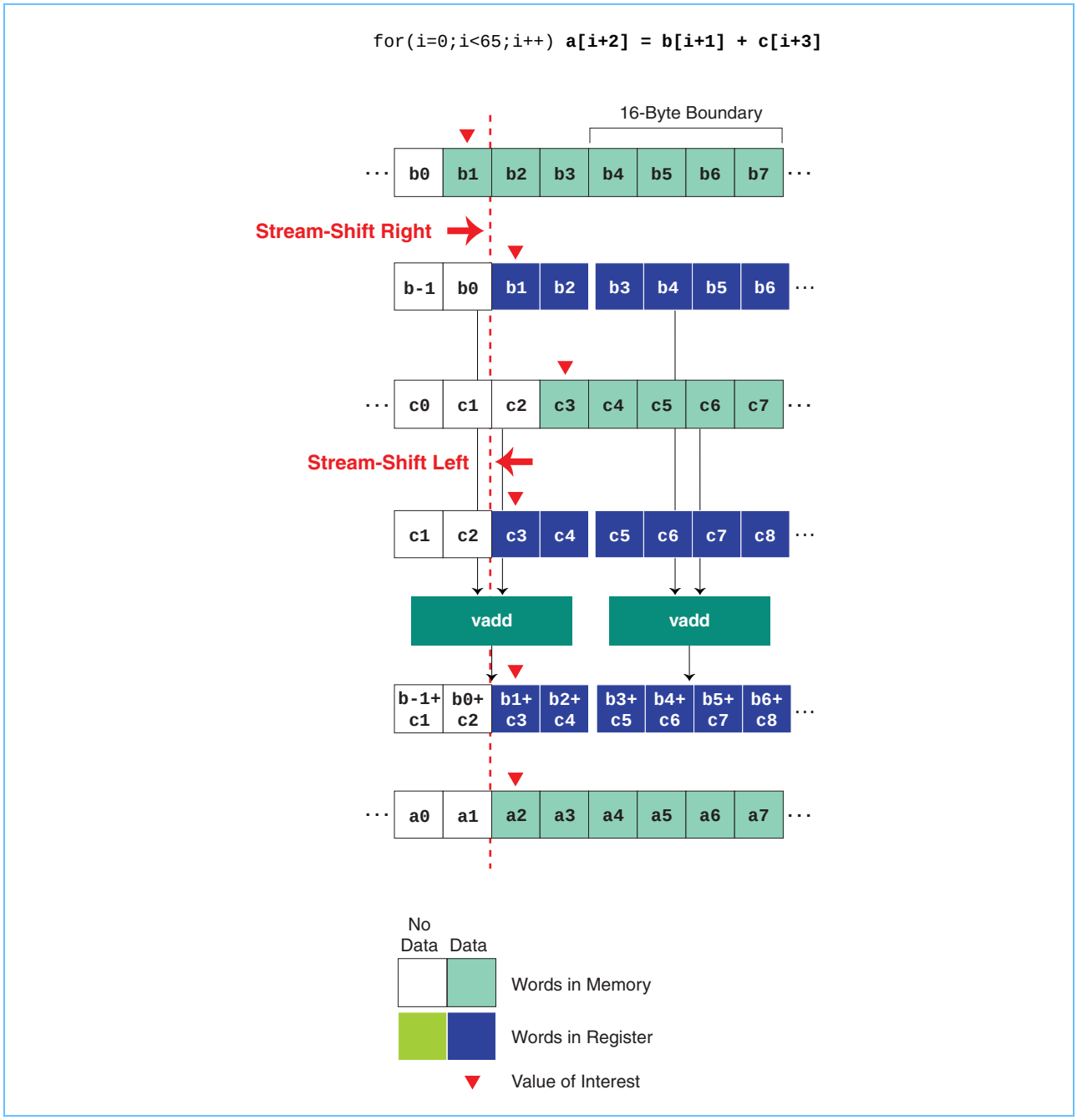
*Figure 22-7. Simple, Correct SIMD Implementation of a Loop*



### 22.1.2.3 *Optimized, Correct SIMDization*

An optimized, and correct SIMD implementation of the loop is illustrated in *Figure 22-8*. Here, the `b[1]` memory stream is shifted right and the `c[3]` memory stream is shifted left, so that both the `b[1]` and `c[3]` elements align at the same register position from which the result can be stored back to memory without stream-shifting. The advantage of this method, compared with that in *Figure 22-7* on page 638, is the elimination of the last stream-shift.

*Figure 22-8. Optimized, Correct SIMD Implementation of a Loop*



#### 22.1.2.4 *SIMDization Code Structure for a Loop*

The basic code structure for the SIMDization of a loop is:

```
<SIMDization prolog>
for(i=0; i<65; i+=4)
   <SIMDization loop body>
<SIMDization epilog>
```

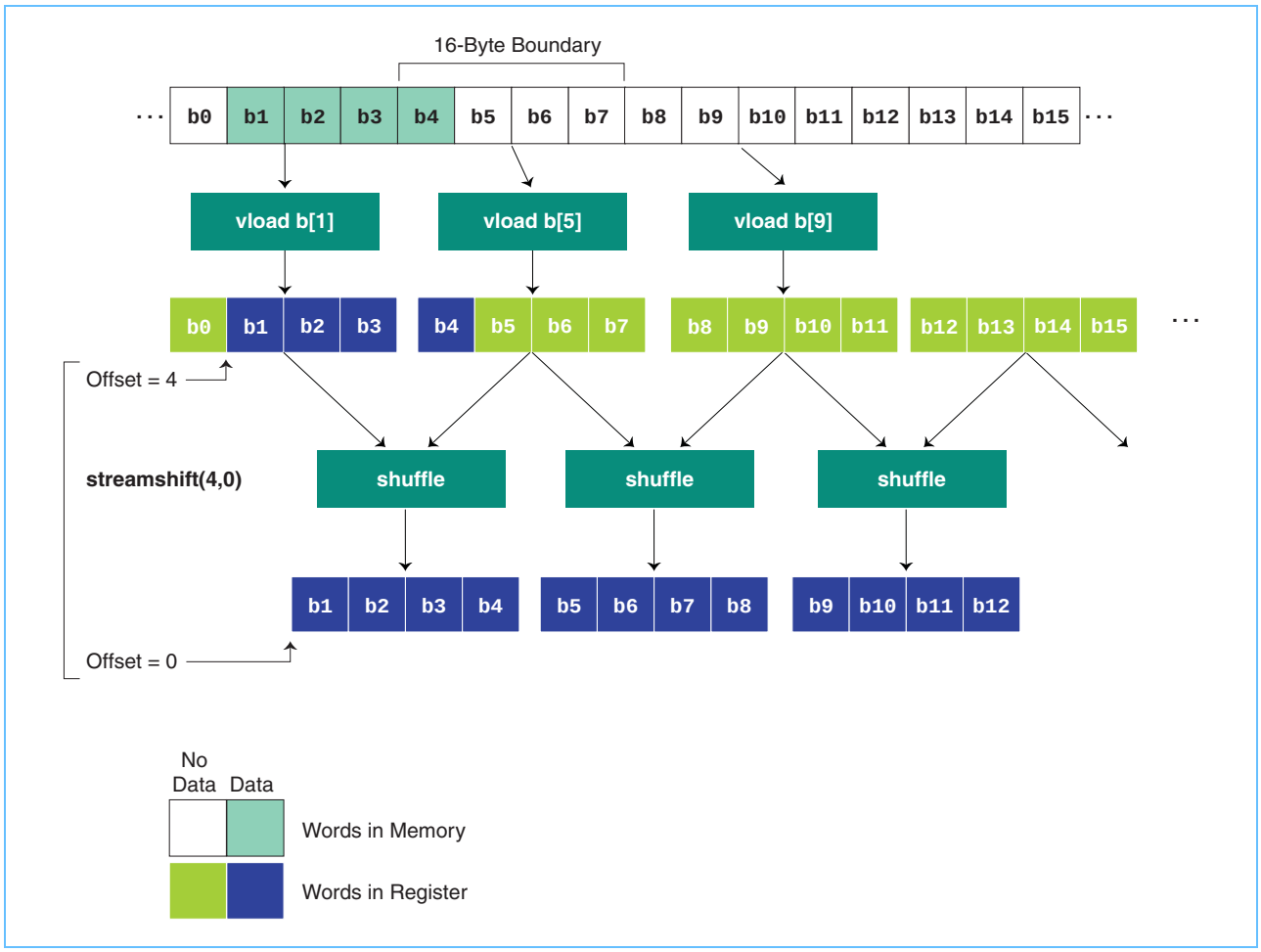
The *SIMDization prolog* (*Figure 22-10* on page 642) is used if the data to be stored is misaligned. The *SIMDization epilog* (*Figure 22-11* on page 643) is used, depending on the requirements for store alignment and the loop's trip-count.

#### 22.1.2.5 *Implementing Stream-Shift Operations on Data*

*Figure 22-9* on page 641 shows how to implement stream-shift operations using shuffle operations (the **shufb** instruction for an SPE, or the **vperm** instruction for the PPE). The operations on `b[i]` require a stream of vectors, starting at memory address `b[1]`, in vector registers. To obtain this, the vector that contains `b[1]`—the vector that begins at address `b[0]` in memory—is first loaded into a vector register with a vector-load instruction, followed by similar loads of subsequent vectors. After the loads are complete, shuffle operations are performed on the contents of the registers so that the `b[1]` element in the first vector is shifted from register offset 4 to offset 0. This operation is denoted in the figure as `streamshift(4,0)`, indicating that offset 4 is moved to offset 0. The pattern for either the **vperm** or **shufb** operation for `streamshift(4,0)` is: (x'04050607', x'08090A0B', x'0C0D0E0F', x'10111213').

Similar stream-shift operations are performed on the vector stream containing `c[3]`. Then, the adds and a store are performed, as shown in *Figure 22-8* on page 639.

*Figure 22-9. Implementing Stream-Shift Operations for a Loop*

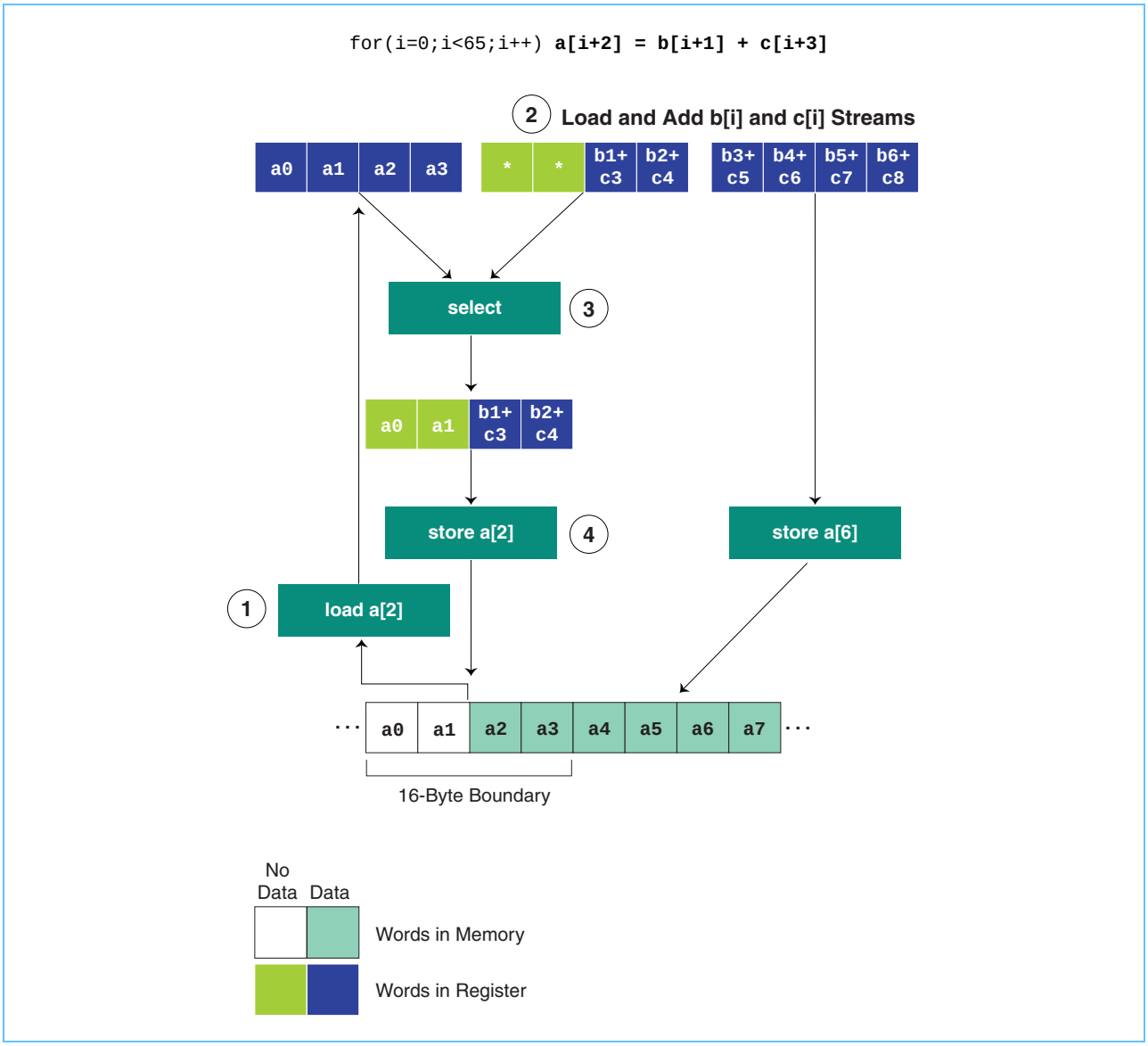


### 22.1.2.6 *The Prolog*

*Figure 22-10* shows how to implement a SIMDization prolog, which is used if the data to be stored is misaligned:

1. The vector that contains `a[2]`—the vector that begins at address `a[0]` in memory—is loaded.
2. The `b[i]` and `c[i]` streams are loaded and added as in *Figure 22-8* on page 639.
3. The most-significant two elements of the `a[i]` vector and the least-significant two elements of the first add result are selected.
4. The result of the select is stored, starting at the memory address of `a[0]`.

*Figure 22-10. SIMDization Prolog*

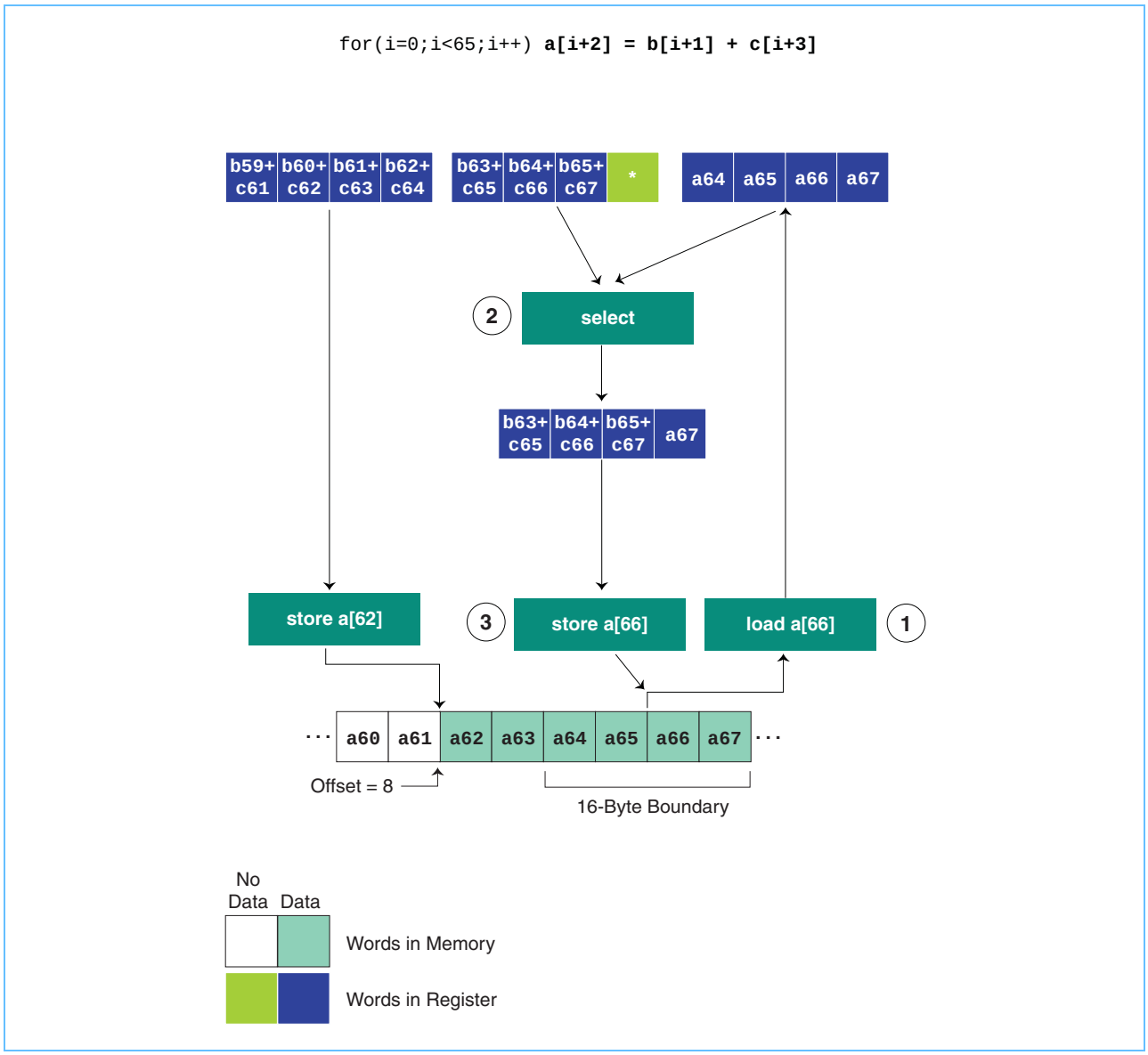


### 22.1.2.7 *The Epilog*

*Figure 22-11* shows how to implement a SIMDization epilog, which depends on the requirements for store alignment and the loop's trip-count. In this example, the trip count is 66. Thus, the last-computed vector result has only three of its four elements with valid data. To handle this:

1. The `a[i]` vector containing the last element, `a[66]`, is loaded.
2. This `a[i]` vector and the three valid result elements are selected.
3. The vector containing the `a[66]` result is stored.

*Figure 22-11. SIMDization Epilog*



### 22.1.2.8 *Machine-Independent Pseudocode*

The machine-independent pseudocode result of the preceding SIMDization operations—before software pipelining of adjacent loads—follows. The code is based on the following assumptions:

- All operations are on vectors.
- Loads and stores are normalized to truncated addresses.
- For SPE application, the **splice**, **shiftpairl**, and **shiftpairr** operations are data-reorganization operations that map to the **spu_shuffle** or **spu_sel** intrinsics.

```
i = 0;
a[i] = splice(a[i], shiftpairr(b[i - 4], b[i], 4) + shiftpairl(c[i - 4], c[i],4), 8);
i=4;
do {
   a[i] = shiftpairr(b[i-4], b[i], 4) + shiftpairl(c[i-4], c[i],4);
   i = i + 4;
} while (i < 65-4);
a[i] = splice(shiftpairr(b[i-4], b[i],4) + shiftpairl(c[i-4], c[i], 4), a[i],12));
```

### 22.1.2.9 *Typical SPE Code*

A typical example of actual code for an SPE—after software pipelining of loads, loop normalization, and address truncation—follows.

```
extern vector unsigned int a[], b[], c[];            // quadword aligned

int i;
vector unsigned int tc, tb, oldSPCopy0, oldSPCopy1;

vector unsigned char patt_b = (vector unsigned char){12,13,14,15,16,17,18,19,
                                                     20,21,22,23,24,25,26,27};
vector unsigned char patt_c = (vector unsigned char){4,5,6,7,8,9,10,11,
                                                     12,13,14,15,16,17,18,19};

tc = c[1];
tb = b[0];

a[0] = spu_sel(a[0],spu_add(spu_shuffle(b[-1],tb,patt_b),
             spu_shuffle(c[0], tc,patt_c)), (vector unsigned int)spu_maskb(255));

i = 0;
do {
   oldSPCopy0 = tc;
   oldSPCopy1 = tb;

   tc = c[i+2];
   tb = b[i+1];
```

```
    a[i+1] = spu_add(spu_shuffle(oldSPCopy1,tb,patt_b),
                     spu_shuffle(oldSPCopy0,tc,patt_c));

    i = i + 1;
} while (i < 15);

a[16] = spu_sel(spu_add(spu_shuffle(tb,b[16],patt_b),
           spu_shuffle(tc,c[17],patt_c)), a[16],(vector unsigned int)spu_maskb(15));
```

### 22.1.3 Coding for Effective Auto-SIMDization

Before going further into the subject of SIMD programming, here are some general application-programming guidelines for writing code that can be effectively SIMDized by a compiler—a method called *auto-SIMDization*: Details about why these methods work well are described in subsequent sections.

#### 22.1.3.1 *Organize Loops*

Organize loops so that they can be auto-SIMDized. This might not be easy, but it is possible. Such loops should be inner-most loops. They should use the `for` loop construct, not the `while` loop construct; the `while` loop construct is not SIMDizable.

*Constant Trip Count*

A constant trip count[4] is best. But in any case, the trip count should be more than three times the number of elements per vector. Loops with long trip count results in more efficient SIMDization. Generally, a trip count of more than three times the number of elements per vector is sufficient for SIMDization. But shorter loops might also be SIMDizable, depending on their alignment and the number of statements in the loop. For example, a loop such as:

```
for(i=0; i<4;i++) a[i] = xxx;
```

This loop has a trip count of four and an array, `a[i]`, of `int`. It can be auto-SIMDized if `a[0]` is aligned.

The reason that the trip count should be more than three times the number of elements per vector is because there must be at least one iteration in the loop for auto-SIMDization to work. Because the prolog and epilog can each require one iteration, three iterations of the loop might be required if the loop has more than two statements and the compiler knows nothing about the alignment.

*Declare Constant Loop Bounds*

Use `#define` to declare constant loop bounds. Doing so increases the reliability that the compiler will use the correct constant.

4. “Trip count” is the number of iterations in a loop.

*Use the Select Operation for Conditional Branches*

Loops that contain if-then-else statements might not always be SIMDizable. However, if the if-then-else statement has a common left-hand variable and can be expressed using the C-language `:?` (colon question-mark) operator, the code will SIMDize using the *select bits* instruction. Here is an example of a realizable data-parallel select operation:

```
for (i = 0; i < VL; i++)
    if (a[i] > b[i])
        m[i] = a[i]*2;
    else
        m[i] = b[i]*3;
```

### 22.1.3.2 *Organize Data in Memory*

Lay out data in memory so that operations on it can be easily SIMDized.

- Use stride-one accesses. Non-stride-one accesses[5] are less efficiently SIMDized, if at all. Random or indirect accesses are not SIMDizable.
- Lay out the data to maximize aligned accesses.
- Use arrays; do not do your own pointer arithmetic in the application to access large data structures. Use global arrays that are statically declared. Arrays are aligned to 16-byte boundaries.
- If it is not possible to use aligned data, use the `alignx` directive to indicate to the compiler what the alignment is. For example:

  ```
  #pragma alignx(16, p[i+n+1]);
  ```

- If you know arrays are disjoint, use the disjoint directive to indicate to the compiler that the arrays specified by the pragma are not overlapping:

  ```
  #pragma disjoint (*ptr_a, b)
  #pragma disjoint (*ptr_b, a)
  ```

### 22.1.3.3 *Organize Algorithms*

Structure algorithms to reduce dependencies:

- Loops with inherent dependences are not SIMDizable.
- Avoid using pointers when not absolutely necessary. Potential aliasing can confuse the compiler.

5. A stride of one refers to a memory access pattern in which each element in a list is accessed sequentially.

#### 22.1.3.4 *Use Appropriate Computations*

Avoid computations that are not SIMDizable, including:

- Function calls are not SIMDizable. Use inline functions or macros in place of function calls. If function calls are kept, be sure to enable inlining by the compiler and possibly add an inlining directive to make sure that it happens.
- Operations that do not map onto native vector instructions are less SIMDizable. In general, all operations except load, branch, and hint-for branch are capable of being mapped.

## 22.2 Auto-SIMDizing Compilers

Compilers that optimize for traditional vector processors are called *vectorizing compilers*. To make the distinction between traditional vector processing and SIMD processing, the terms *SIMDization* and *SIMDize* are used to refer to vectorization on a SIMD processor.

A compiler that automatically merges scalar data into a parallel-packed SIMD data structure is called an *auto-vectorizing* compiler. Such compilers must handle all the high-level language constructs and therefore do not always produce optimal code. Although programming with intrinsics gives programmers a direct and powerful way to access the SIMD parallelization of SPU operations, this kind of programming can be challenging and time-consuming. In such cases, use of an auto-vectorizing compiler is generally less tedious and more productive than hand-coding, except perhaps for very critical kernel loops. But even in the latter cases, an auto-vectorizing compiler can often achieve performance that is nearly as good as hand-coding with intrinsics.

Although this section, and the remainder of this section, are intended primarily for compiler writers, application programmers can benefit from understanding the challenges and limitations facing a SIMDizing compiler, so that their application programs can avoid coding styles that are not readily SIMDizable by a compiler.

To understand the context of this explanation, consider the following example of SIMDizing a loop. The original, scalar loop multiplies two arrays, term by term. In the SIMDized version, the arrays are assumed to remain scalar outside of the subroutine, `vmult`.

```
/* Scalar version */
int mult(float *array1, float *array2, float *out, int arraySize) {
    int i;
    for (i = 0; i < arraySize; i++) {
        out[i] = array1[i] * array2[i];
    }
    return 0;
}


/* SIMDized version for an SPE */
int vmult(float *array1, float *array2, float *out, int arraySize) {

    /* This code assumes that the arrays are quadword-aligned. */
    /* This code assumes that the arraySize is divisible by 4. */

    int i;
```

```
    vector float *varray1 = (vector float *) (array1);
    vector float *varray2 = (vector float *) (array2);
    vector float *vout = (vector float *) (out);

    for (i = 0; i < arraySize/4; i++) {
        /* spu_mul is an intrinsic that multiplies vectors */
        vout[i] = spu_mul(varray1[i], varray2[i]);
    }

    return 0;
}
```

### 22.2.1 Motivation and Challenges

Compiling for the SPU's SIMD operations most closely resembles the work done with traditional loop-based vectorization, pioneered for vector supercomputers. However, architectural constraints for accessing memory and performing vector operations present new challenges.

#### 22.2.1.1 *Misaligned Data*

One of the most significant challenges of SIMDization is dealing with misaligned data. Because the SPU hardware supports loads and stores only on quadword (16-byte) alignment, SIMDizing loops that operate on misaligned data accesses leads to incorrect results. Traditionally, a valid vectorization is constrained only by dependences: whether or not data in the loop depends on other data being computed first. SIMDization must satisfy three additional alignment conditions:

- The alignment of data in memory dictates the byte-offset in its destination vector register when it is loaded.
- Vector operations must operate on data that are at the same byte-offset in the input vector registers.
- Data being stored must be at the same byte-offset in the vector register as the memory alignment of the store address.

#### 22.2.1.2 *Scatter-Gather*

*Scatter-gather* refers to a technique for operating on sparse data, using an index vector. A *gather* operation takes an *index vector* and loads the data that resides at a base address added to the offsets in the index vector. A *scatter* operation stores data back to memory, using the same index vector.

The CBEA processor SIMD Architecture does not directly support scatter-gather in hardware. Some processors do support scatter-gather directly in hardware, but this capability comes at a great cost to the memory subsystem. Because the CBEA processors do not support scatter-gather in hardware, the best way to extract SIMD parallelism is to combine operations on data in adjacent memory addresses into vector operations.

The CBEA processors support only stride-one loads and stores. Only data stored consecutively in memory is loaded or stored. Array accesses such as those represented by `a[2i+1]`, where `i` is an iteration index, result in poor memory-access patterns. Random indirect accesses are also a potential performance problem.

The CBEA processor architectural approach has the advantage of providing very efficient support for the most common types of memory access, while still providing relatively inexpensive support for rarer types of memory access, such as scatter-gather. This architectural approach enables the CBEA processors to operate at faster clock speeds and higher compute density than would be the case if they supported scatter-gather in hardware.

#### 22.2.1.3 *Packed Data*

Because SIMD vectors consist of packed data elements, a compiler writer must keep track of the number of data elements in each register. A packed vector refers to the fact that a 16-byte register can hold 16 1-byte values, eight 2-byte values, four 4-byte values, or two 8-byte values—for example, `char`, `short`, `int`, or `long long`, signed or unsigned. The number of values in a register changes when data is converted from one data type to another, such as when a vector of 1-byte `char` values is converted to a vector of 2-byte `short` values. In this case, the original vector expands to become two vectors.

#### 22.2.1.4 *Mix of Data Types*

The mix of data types in multimedia programs can present problems. Traditional vectorizable programs contain loops with a single data type. In contrast, multimedia programs contain a rich mix of data types, and converting from one type to another is common.

Data is often stored as 1-byte `char`, but operated on as 2-byte `short`. Without taking care, it is easy to use the whole computation bandwidth by operating on eight `short` values at a time, but waste half the memory bandwidth by only storing eight `char` values at a time. A compiler writer must efficiently handle mixed data types in a loop to maximize both computation and memory bandwidth.

#### 22.2.1.5 *Nonuniform Instruction Set*

Not all data types are equally supported by the SPU instruction set. For example, vectors of 16-bit integers are the best supported. There is weak support for `vector long long` (vectors of 64-bit integers). There is slow throughput for `vector double` (vectors of double-precision floating-point numbers). 16-bit vector multiply is supported, but 32-bit vector multiply is not. There is no support for complex-number arithmetic.

#### 22.2.1.6 *Parallelism Program Scope*

Although long-running loops were traditionally targets for optimization on vector processors, the relatively short vectors in SIMD Architectures make SIMD programming attractive even for short, unstructured computations. In many multimedia programs, SIMD parallelism can be extracted from computations among the x, y, and z coordinates of a vertex. Multimedia programs often contain manually unrolled loops and loops with a short trip count. A compiler writer must be able to extract SIMD parallelism from various program scopes, such as a single basic block, an inner-most loop, or a loop nest.

The CBEA processors provide plenty of support for complex-number arithmetic, such as the ability to perform SIMD and shuffle/permute operations. Most of the CBEA processor hardware support has been abstracted away from special operations (such as complex-number arithmetic) to general operations on SIMD code—(for example, from `chars` to `double`). Thus, although there is no dedicated hardware for complex-number arithmetic, this results is higher overall performance for most SIMD applications.

#### 22.2.1.7 *SIMDization Interaction with Other Optimizations*

SIMD functional units are more constrained than scalar functional units, SIMD vectors are relatively short, and multimedia programs are complex applications. Because of these considerations, the interaction between SIMDization and other compiler optimizations can be quite involved. For example, there is a trade-off between SIMDization and instruction-level parallelism (ILP). Without taking care, a compiler can SIMDize a loop but degrade performance.

Addressing each of these challenges is difficult in itself, but they can and do appear in actual applications all at once. The following code fragment exhibits several of these problems: misaligned access, length conversion, short loops and available SIMD parallelism across nested loops:

```
short input[], coef[];
for (i=0; i < NInputs; i++) {
    int sum = 0;
    for (j=0; j < 16; j++)
        sum += input[i+j] * coef[j];
    output[i] = sum;
}
```

### 22.2.2 Examples of Invalid and Valid SIMDization

*Section 22.1.2 Approaching SIMD Coding Methodically* on page 634 presents an example of invalid and valid SIMDization methods for a loop. This section presents a similar example but with a bit more detail.

Consider the following code fragment, in which integer arrays `a`, `b`, and `c` are quadword-aligned:

```
for (i = 0; i < 20; i++) {
       a[i+3] = b[i+1] + c[i+2];
}
```

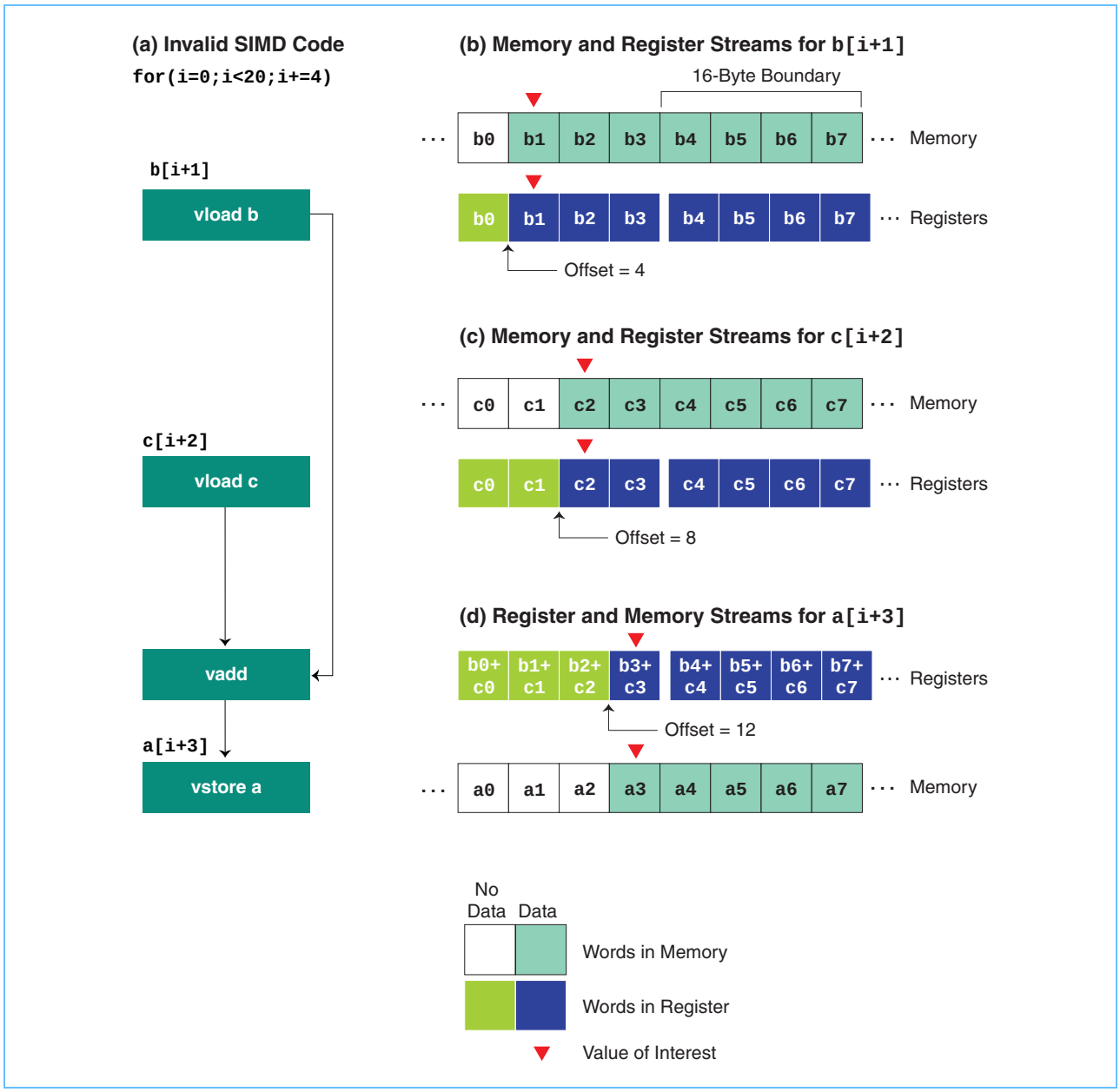
Because there is no loop-carried dependence, the loop can be easily SIMDized for traditional vector processors that have no alignment constraints. However, such SIMDized code is invalid for SIMD units, such as the CBEA processors, that support only aligned loads and stores.

*Figure 22-12* on page 651 shows the incorrect execution of this loop when it is SIMDized for traditional vector processors on hardware with alignment constraints. Consider the first iteration of the SIMDized loop in *Figure 22-12(a)*, especially the values of the expression `a[3]=b[1]+c[2]` that are highlighted with red triangles in items (b) through (d). The `vload b[1]` operation loads vector `b[0],...,b[3]` with b[1] at byte-offset 4 in its vector register. The `vload` operation is a

compiler-internal function, described in *Section 22.3.4.1 Streams* on page 658. The operation `vload(addr(i))` loads a vector from a stride-one memory reference described by `addr(i)`. This produces a register stream whose byte-offset is the alignment of `addr(i)`.

Similarly, the `vload c[2]` operation loads `c[0],...,c[3]` with `c[2]` at byte offset 8, as shown in *Figure 22-12* (c). Adding these two vector registers yields the values `b[0]+c[0], ..., b[3]+c[3]`, as illustrated in *Figure 22-12* (d). This is not the result specified by the original `b[i+1]+c[i+2]` computation.

*Figure 22-12. Invalid SIMDization on Hardware with Alignment Constraints*

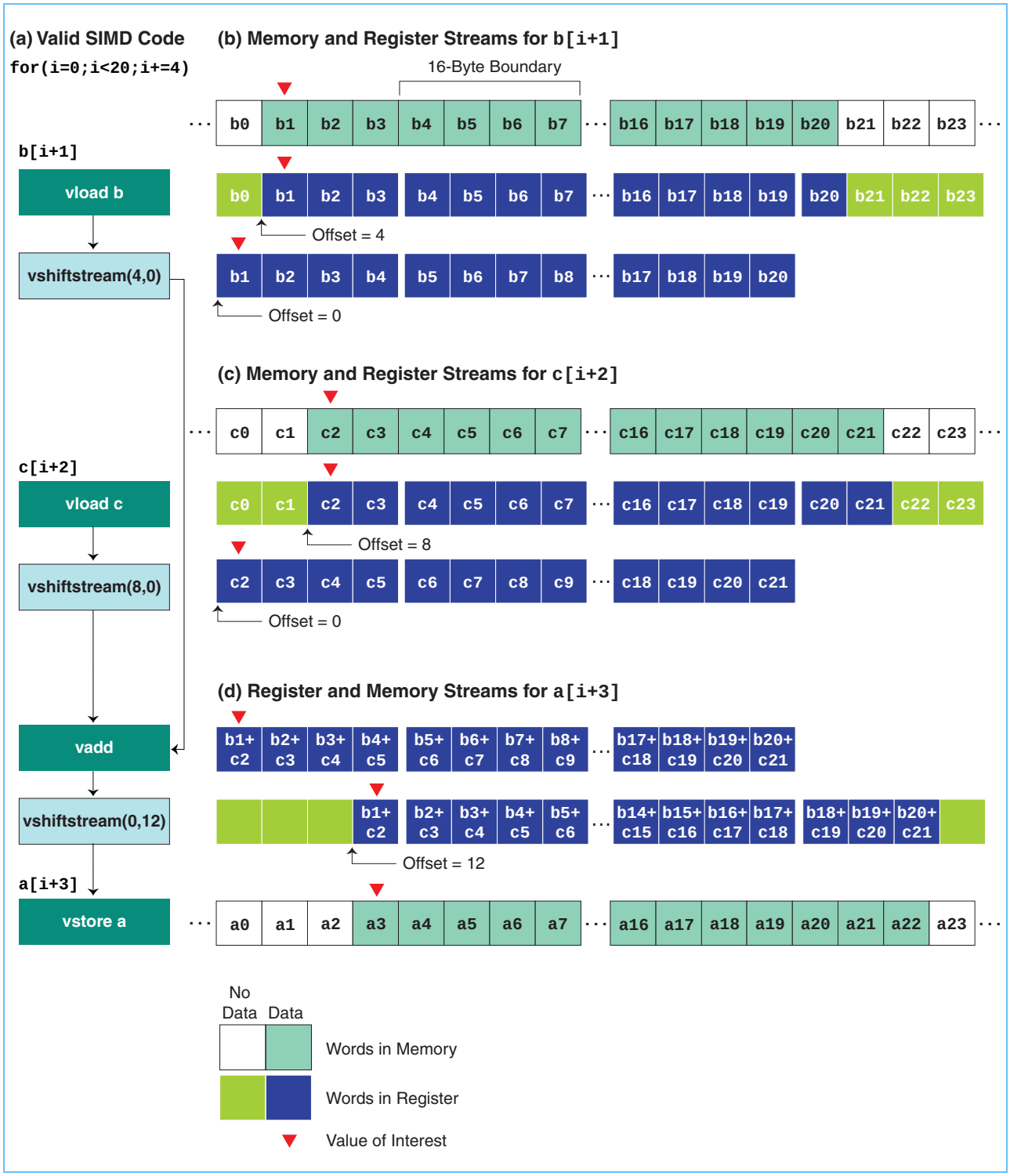


A valid vectorization for traditional vector units is only constrained by dependences. This is no longer sufficient when SIMDizing for architectures with alignment constraints. A valid SIMDization must satisfy these conditions:

- The alignment of data in memory dictates the byte-offset in its destination vector register when it is loaded.
- Vector operations must operate on data that are at the same byte-offset in the input vector registers.
- Data being stored must be at the same byte-offset in the vector register as in the memory alignment of the store address.

*Figure 22-13* on page 653 shows one example of a valid SIMDization for the loop with misaligned references. A compiler can automatically generate data reorganization instructions to align data in registers to satisfy these alignment constraints. Leading and trailing **vstore** operations require store Prolog and Epilog, as described in sections *Section 22.1.2.6* on page 642 and *Section 22.1.2.7* on page 643.

*Figure 22-13. An Example of a Valid SIMDization on SIMD Unit with Alignment Constraints*



## 22.3 SIMDization Framework for a Compiler

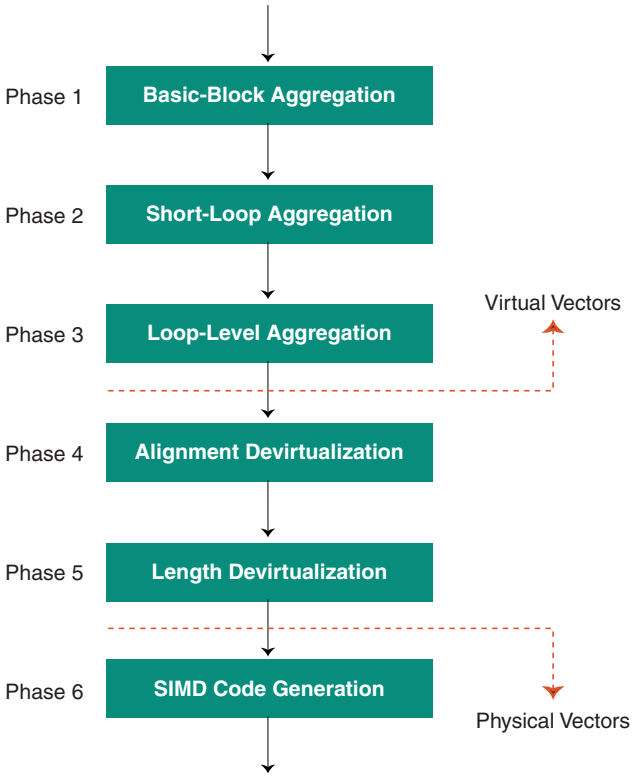
This section describes how a compiler can provide a SIMDization framework using a set of SIMDization algorithms. Although the intended audience is primarily compiler writers, the description might also be interesting to application programmers. It explains techniques for efficiently handling alignment and multiple-sourcing of SIMD data, and the conversion techniques that can be applied.

A SIMDization framework is presented that SIMDizes loops primarily with stride-one memory accesses. A stride of one refers to a memory access pattern in which each element in a list is accessed sequentially. However, the SIMDization framework described accomplishes more than SIMDizing the obvious candidates. Some of the features of the framework include:

- SIMD parallelism can be extracted from different program scopes, including basic block, inner-most loops, and inner-most loop nests.
- SIMDization can result in mixed scalar and SIMD code.
- Loops may contain multiple statements.
- Data inside the loops may have different data types.
- Loops may convert data from one type to another.
- Data inside the loops may have mixed lengths.
- Loops that are manually unrolled are SIMDized.
- Operations across vectors (reductions) can be SIMDized.
- Loops may contain induction, reduction, and private variables.
- Both compile-time alignment and runtime alignment can be performed.
- Interprocedural alignment analysis is used to obtain accurate alignment information.
- Programmer-inserted directives describe alignment of data (for example, the `_align_hint` intrinsic specified by the *C/C++ Language Extensions for Cell Broadband Engine Architecture* document).
- Redundant sign-extensions are minimized, to reduce unnecessary packing and unpacking overhead.
- Programmer-inserted directives indicate which loops can be parallelized.
- Diagnostic information is provided on SIMDization failure.

The approach uses *virtual vectors* to carry out the early SIMD transformations. These vectors have no alignment constraints and can have any length. Compilers gradually devirtualize virtual vectors to physical vectors by later SIMD transformations. At that stage, architectural constraints, such as alignment and actual vector length, are taken into account. *Figure 22-14* on page 655 shows the six phases of the SIMDization process.

*Figure 22-14. SIMDization Framework Based on Virtual Vectors*




SIMDization is broken into six phases. Phases 1, 2, and 3 operate on virtual vectors (basic-block aggregation, short-loop aggregation, loop-level aggregation). Phases 4 and 5 transition the virtual vectors to physical vectors (alignment and length devirtualization). Phase 6 operates on physical vectors (SIMD instruction scheduling and code generation). The first three aggregation phases combine operations on data at adjacent memory into vector operations. Such contiguous memory access patterns are found in two kinds of places:

- At the basic-block level, adjacent memory accesses like `a.y` and `a.z`, or `b[i]` and `b[i+1]`, can be combined.
- At the loop level, stride-one accesses like `c[i]` across consecutive loop iterations can be combined.

Unfortunately, these two types of accesses are often not compatible. For example, although `d[i].x` and `d[i].y` are adjacent accesses at the basic-block level, they are not stride-one accesses at the loop level.

### 22.3.1 Phase 1: Basic-Block Aggregation

Basic-block aggregation extracts SIMD parallelism within a basic block by packing isomorphic computation on adjacent memory accesses into vector operations. That is, when data is being operated on in the same way, but the data is not already in a vector, the data is moved into a vector so that it can later be SIMDized. This is possible only if the data is stored in adjacent addresses in memory.

Such parallelism is often found in unrolled loops, whether the loops are unrolled manually by the programmer or automatically by the compiler. It is also common in graphic applications that manipulate adjacent fields of aggregates, such as the subdivision computation shown here:

```
for (i=0; i < n; i++) {
    t = triangles[i];
    v[i].x=w0 * t.p[0].x + w1 * t.p[1].x + w2 * t.p[2].x;
    v[i].y=w0 * t.p[0].x + w1 * t.p[1].x + w2 * t.p[2].y;
    v[i].z=w0 * t.p[0].x + w1 * t.p[1].x + w2 * t.p[2].z;
    v[i].w=w0 * t.p[0].x + w1 * t.p[1].x + w2 * t.p[2].w;
}
```

In this example, the last three computations can be packed into one vector statement. This is because they are isomorphic, and also operate on adjacent fields in memory (`x, y, z, w`). The first statement in this loop (`t = triangles[i]`) is an aggregate copy that can be SIMDized into a vector copy.

A loop SIMDized at the basic-block level can be further SIMDized at the loop-level. For example, if `x, y, z` and `w` are 16-bit integers, both basic-block and loop-level transformations are needed to extract enough SIMD parallelism for a 16-byte vector.

### 22.3.2 Phase 2: Short-Loop Aggregation

Short-loop aggregation eliminates inner loops that have short trip counts. SIMDizable short loops are collapsed into vector operations. The next loop in the original loop nest becomes the new inner-most loop. This expands the scope of subsequent SIMDization transformations.

For example, consider a simplified loop nest, which negatively indexes the array, `dp`:

```
for (j=40; j ≤ 120; j++) {
    sum = 0;
    for (i=0; i < 40; i++) {
        sum += wt[i] * dp[i-j];
    }
    L_result[j] = sum;
}
```

After short-loop aggregation, the inner loop is SIMDized into a single statement operating on a vector of 40 elements, because its trip count is 40. See *Section 22.3.4.1 Streams* on page 658 for a definition of `reduct`.

```
for (j=40; j ≤ 120; j++) {
    sum = 0;
    sum = sum + reduct(+, wt[0:39] * dp[-j:39-j]);
    L_result[j] = sum;
}
```

### 22.3.3 Phase 3: Loop-Level Aggregation

Loop-level aggregation extracts SIMD parallelism across loop iterations. Computations on stride-one accesses across loop iterations are combined into vector operations. The loop is blocked by `B`, defined such that the byte length of each virtual vector is a multiple of the physical vector length. Loop-blocking refers to rewriting the loops into blocks of size `B`, to make each block accessible for later SIMDization.

In loop-level aggregation, each statement is packed with itself across consecutive iterations. Each final vector length must be a multiple of 16 bytes. A scalar `b[i]` and a vector `a[i].x,y,z` are treated alike, regardless of alignment.

Continuing with the example in short-loop aggregation (*Section 22.3.2 Phase 2: Short-Loop Aggregation* on page 656), the remaining `j` loop is SIMDized. The variable `sum` is a private variable (a variable used only in this loop), `L_result[j]` is a stride-one access, and `reduct` operates over a vector. Assuming that all the data is `int`, the blocking factor is computed as:

```
B = 16/GCD(16, sizeof(int)) = 16/4 = 4
```

where `GCD` is the Greatest Common Divisor—the largest integer that divides into both of the integer operands. SIMDization of `sum` and `L_result[j]` is straightforward. `reduct` is more complicated because its data is not stride-one across the `j`-loop. Specifically, `wt[0:39]` is now loop-invariant (its stride is 0). `dp[-j:39-j]` has a stride of -1/40. The goal is to SIMDize `reduct` into a 4-way `parallel-reduct` (see *Section 22.3.4.1 Streams* on page 658 for a definition of `parallel_reduct`, `elem_splat`, and `elem_slide`).

Because `wt[0:39]` has stride 0, to transform the reduction into a 4-way parallel reduction, each of its elements needs to be replicated 4 times. A compiler-internal operation, `elem_splat(wt[0:39, 4)`, can be used for this.

For `dp[-j:39-j]`, the four vectors to be reduced in parallel are `dp[-j:39-j], ..., dp[-j-3:30-j-3]`. The corresponding elements of the four data streams form a packed vector among themselves. To produce a 4-way parallel reduction, each element of the original vector needs to be expanded by applying a sliding window of -4 elements: `elem_slide(dp[-j:39-j], -4)`. The code after loop-level SIMDization is:

```
for (j=40; j ≤ 120; j+=4) {
    vsum = (0, 0, 0, 0);
    vsum = vsum + parallel_reduct(+, 4, elem_splat(wt[0:39])
        *elem_slide(dp[-j:39-j], -4));
    L_result[j:j+3] = vsum;
}
```

### 22.3.4 Phase 4: Alignment Devirtualization

The key to automatically SIMDizing loops with arbitrary misalignment is to provide automatically generating data-reorganization operations that align data in registers. This data reorganization satisfies the alignment constraints stated previously: that vector operations can only operate on aligned data, and that loads and stores to memory must be properly aligned.

As illustrated in *Figure 22-14* on page 655, the top-level algorithm for SIMDizing a loop has six phases:

- *Phases 1, 2, and 3:* SIMDize the loop as if there are no alignment constraints.
- *Phase 4:* Insert data-reorganization operations to satisfy actual alignment constraints. This phase creates the data-reorganization graph:
  - Insert the data-reorganization operations.
  - Apply optimizations to reduce the number of data-reorganization operations.
- *Phase 5:* Length devirtualization.
- *Phase 6:* Generate SIMD code. In this phase, take care of runtime alignment, unknown loop bounds, multiple misalignments, and multiple statements.

Although some alignments are known at compile-time, others are known only at runtime. In fact, runtime alignment is pervasive in applications either because it is part of the algorithms used, or because it is an artifact of the compiler's inability to extract alignment information from complex applications.

The alignment framework incorporates *length conversions*, which are conversions between data of different sizes. Length conversions are pervasive in multimedia applications, where mixed integer types are often used. Supporting length conversion greatly improves the number of SIMDizable loops.

This loop contains misaligned accesses:

```
for (i = 0; i < 100; i++) {
    a[i+2] = b[i+1] + c[i+3];
}
```

#### 22.3.4.1 *Streams*

A *stream* is a sequence of contiguous memory locations that are accessed by a memory reference throughout the lifetime of a loop. This is also called a *memory stream*. By analogy, a stream is also a sequence of contiguous registers that are produced by an operation over the lifetime of a loop. This is also called a *register stream*.

Vector operations in a loop can be viewed as operations on streams. For example, a vector load consumes a stream of memory and produces a stream of registers. The stream is a way to view values throughout the lifetime of a loop.

An important property of a stream is its *stream offset*. This is defined as the byte-offset of the first required value in the first register of a stream. In the presence of misaligned streams, data-reorganization operations must be inserted to enforce required stream offsets. There are several such compiler-internal operations, including:

- `vshiftstream(stream, inputOffset, outputOffset)`—Shift `stream`, the register stream, from stream offset `inputOffset` to stream offset `outputOffset`. For example, `vshiftstream(stream, 4, 0)` shifts the stream to the left 4 bytes, and `vshiftstream(stream, 0, 12)` shifts the stream to the right 12 bytes. Another way to think of this: `vshiftstream` takes an input register stream whose offset is `inputOffset`, and generates a register stream whose offset is `outputOffset`.
- `vshiftpair(v1, v2, length)`—Create a double-length vector by concatenating vectors `v1` and `v2`, and select the bytes `length, length + 1, length + 2, ... length + V - 1`, where `V` is the vector length.
- `vsplice(v1, v2, length)`—Concatenate the first length bytes of `v1` with the last (`V - length`) bytes of `v2`.
- `prependW(Stream, x)`—Prepend `x` bytes to the beginning of `Stream` of `W`-byte wide registers.
- `skipW(Stream, x)`—Skip the first `x` bytes of `Stream` of `W`-byte wide registers.
- `reduct(op, V(n, t))`—Perform operation `op` over all n elements of vector `V` and generate a single result. This is called a reduction.
- `parallel_reduct(op, x, V(n, V(x, t)))`—Perform operation `op` in parallel `x` times, over `n` distinct values. The data in the vector is interleaved: the first `x` values are the first elements in the of the `x` reductions, the next x values are the second elements in each of the reductions, and so on. The parallel reduction generates a vector of `x` reduction values.
- `vsplat(x)`—Replicate the loop invariant x across a register stream.
- `elem_splat(V(n, t), x)`—Generate a new `V(n, V(x,t))` vector where each of the values in the original vector appears `x` times, consecutively.
- `elem_slide(V(n, t), w)`—Apply a sliding window of `w` elements along the original vector.
- `vload(addr(i))`—Load a vector from a stride-one memory reference `addr(i)`. This produces a register stream whose byte-offset is the alignment of `addr(i)`.
- `vstore(addr(i), src)`—Store a vector stream from `src` to a stride-one reference at `addr(i)`.
- `vop(src1, ..., srcN)`—Take registers `src1` through `srcN` as inputs for a vector operation, and produce an output register stream. Input register streams must have matching stream offsets.
- `vpack(Stream, f)`—Pack `Stream` by the factor `f`. For example, converting from a 4-byte data type to a 2-byte data type is a stream pack with a factor of 2.
- `vunpack(Stream, f)`—Unpack `Stream` by the factor `f`. For example, converting from a 4-byte data type to an 8-byte data type is a stream unpack with a factor of 2.

In this context, *shifting left* and *shifting right* refer to data-reorganization operations that operate on a stream of consecutive registers, not to traditional arithmetic or logical shift operations.

#### 22.3.4.2 *Zero-Shift Policy*

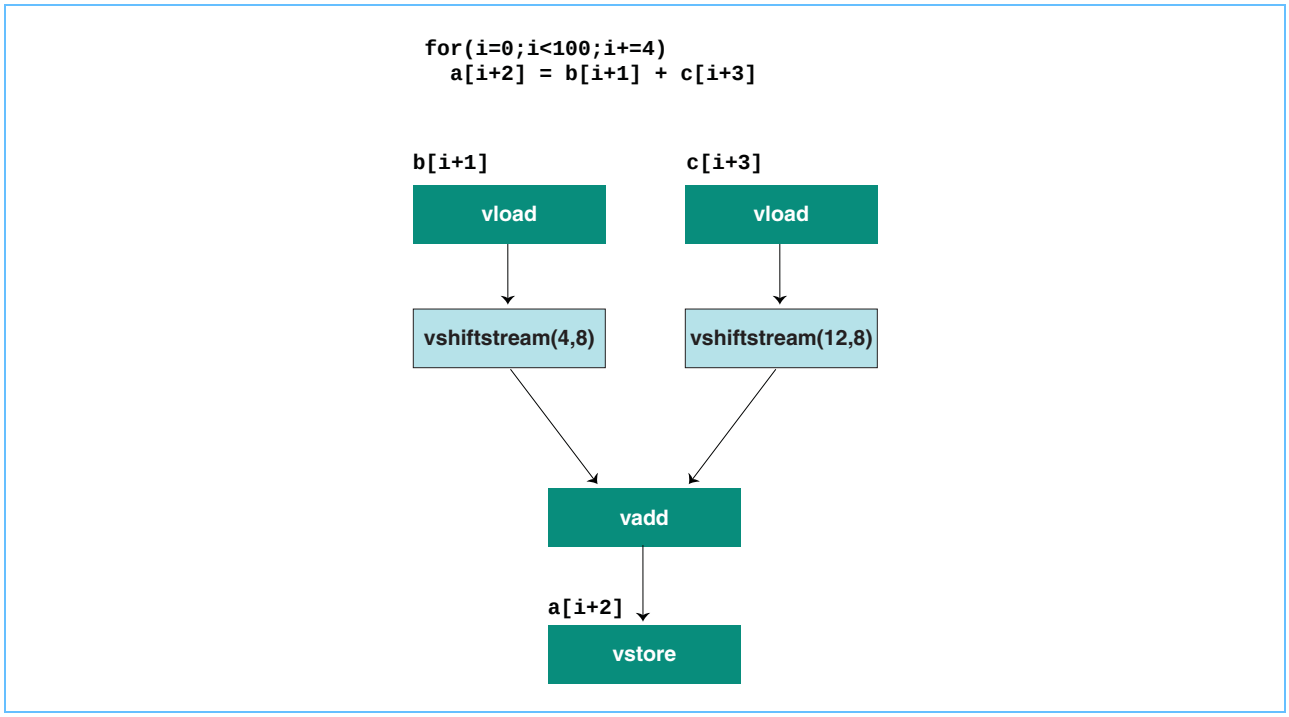
The Zero-Shift Policy is the least-optimized way to reorganize misaligned data. The idea is to shift each misaligned register stream to a stream offset of 0 immediately after it is loaded from memory. Then, each register stream is shifted to the alignment of the store address just before it is stored back to memory. Zero-Shift Policy inserts one `vshiftstream` operation for each misaligned memory stream. Zero-Shift Policy is used for runtime alignment. *Figure 22-7* on page 638 shows an example of the Zero-Shift Policy.

### 22.3.4.3 *Eager-Shift Policy*

The Eager-Shift Policy shifts a misaligned load stream directly to the alignment of the store, rather than to 0 as in the Zero-Shift Policy. This policy requires that alignments of loads and stores be known at compile-time. It can be more efficient than Zero-Shift Policy, lowering the number of `vshiftstream` operations needed. Eager-Shift Policy is applicable only to compile-time alignments, due to code-generation problems: the code sequence for shifting left or right is different.

*Figure 22-15* shows the flow of the Eager-Shift policy, and *Figure 22-8* on page 639 illustrates the policy. It lowers the total number of stream-shift operations from 3 to 2, compared with the Zero-Shift Policy.

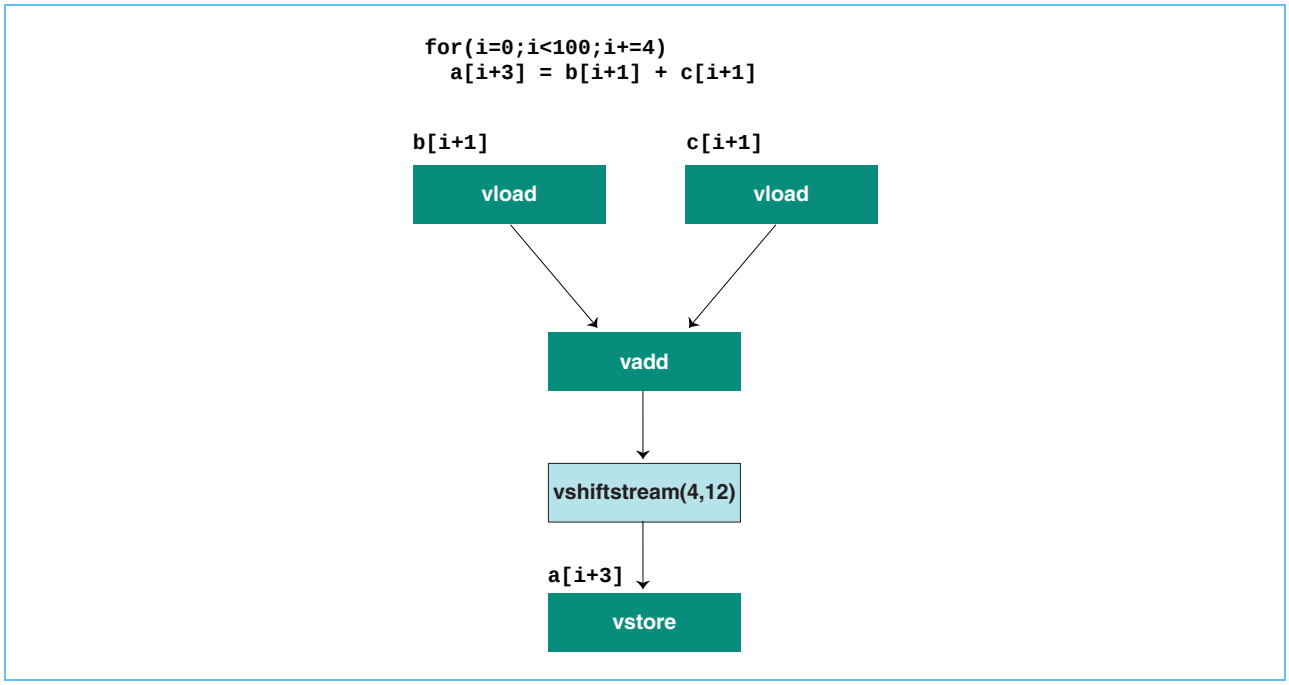
*Figure 22-15. Eager-Shift Policy*




### 22.3.4.4 *Lazy-Shift Policy*

The Lazy-Shift Policy improves upon the Eager-Shift Policy by delaying stream shifts as long as possible. For example, two relatively aligned vectors can be added by first adding them and then shifting the result, instead of first shifting them to the alignment of the store and then adding them. This helps reduce the number of `vshiftstream` operations. *Figure 22-16* on page 661 shows the lazy-shift policy. This policy exploits the case in which `b[i+1]` and `c[i+1]` are relatively aligned, and can be safely operated on as is. Only the result of the add operation needs to be shifted to match the alignment of the store.

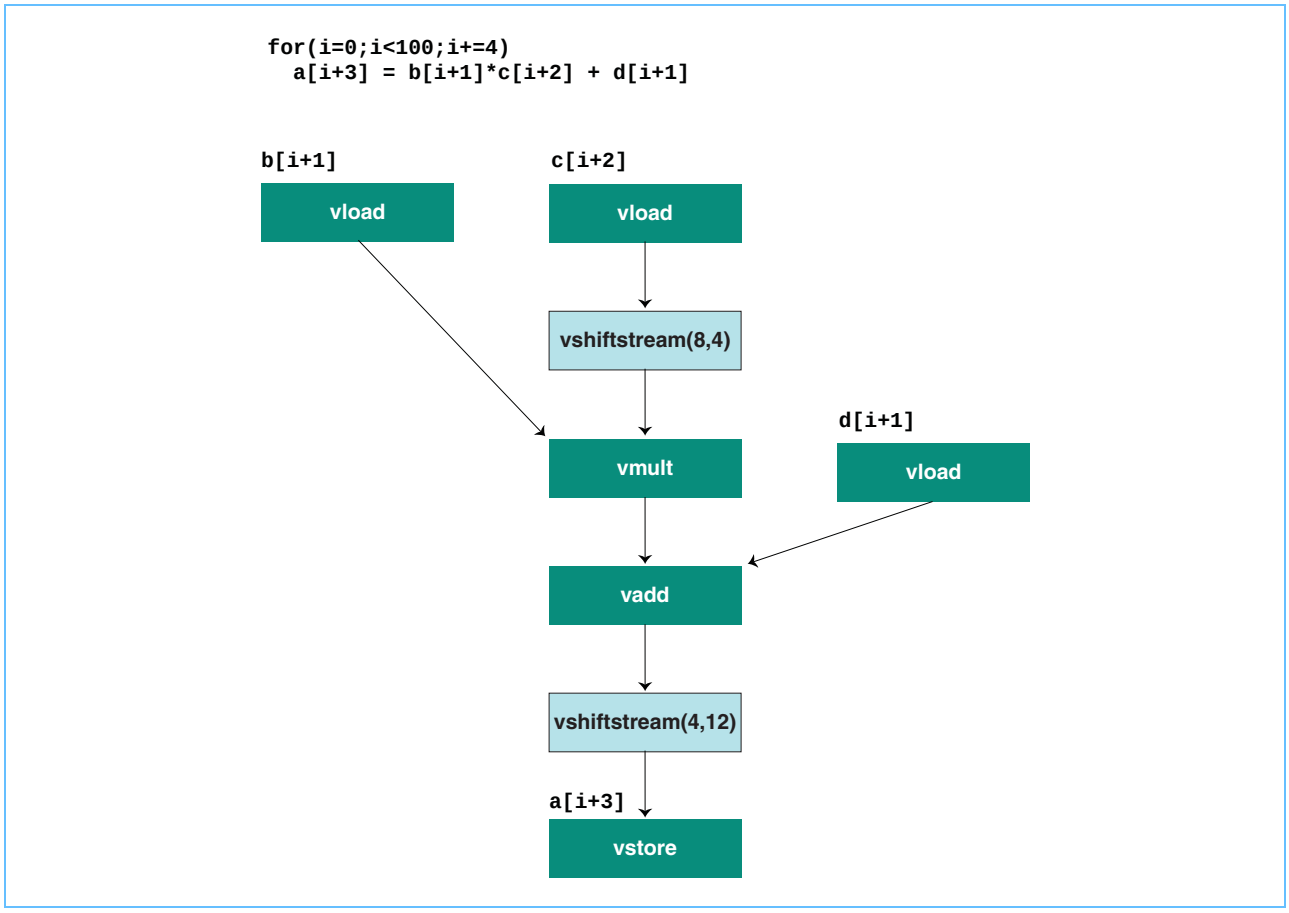
*Figure 22-16. Lazy-Shift Policy*




#### 22.3.4.5 *Dominant-Shift Policy*

The Dominant-Shift Policy reduces the number of `vshiftstream` operations the most. It is most effective if applied after the Lazy-Shift Policy. The offset that is dominant (most often used) among a set of streams is chosen as the offset to shift to. For example, if 4 is the offset used most often, then 4 is the dominant offset, and chosen as the shift target. In *Figure 22-17* on page 662, the dominant offset is the stream offset 4. Shifting the stream to this offset decreases the number of `vshiftstreams` from 4 (for the Zero-Shift Policy) to 2.

*Figure 22-17. Dominant-Shift Policy*




### 22.3.4.6 *Length Conversion*

A length conversion results from any operation that converts a value of one length to a value of another length. The most common length conversions occur between data types of different sizes. Such data conversions are common in multimedia workloads because multimedia data, such as pixels or colors, are stored in compact forms like `short` or `char`, but are operated on as integers or floats to reduce rounding errors. The data-reorganization operations, `pack` and `unpack`, are used to handle data of different lengths. All four shift policies can be applied to expression trees with length conversion. The data-reorganization graph created by phase 4 (*Section 22.3.4* on page 658) is a typical compiler expression tree augmented with data-reorganization operations.

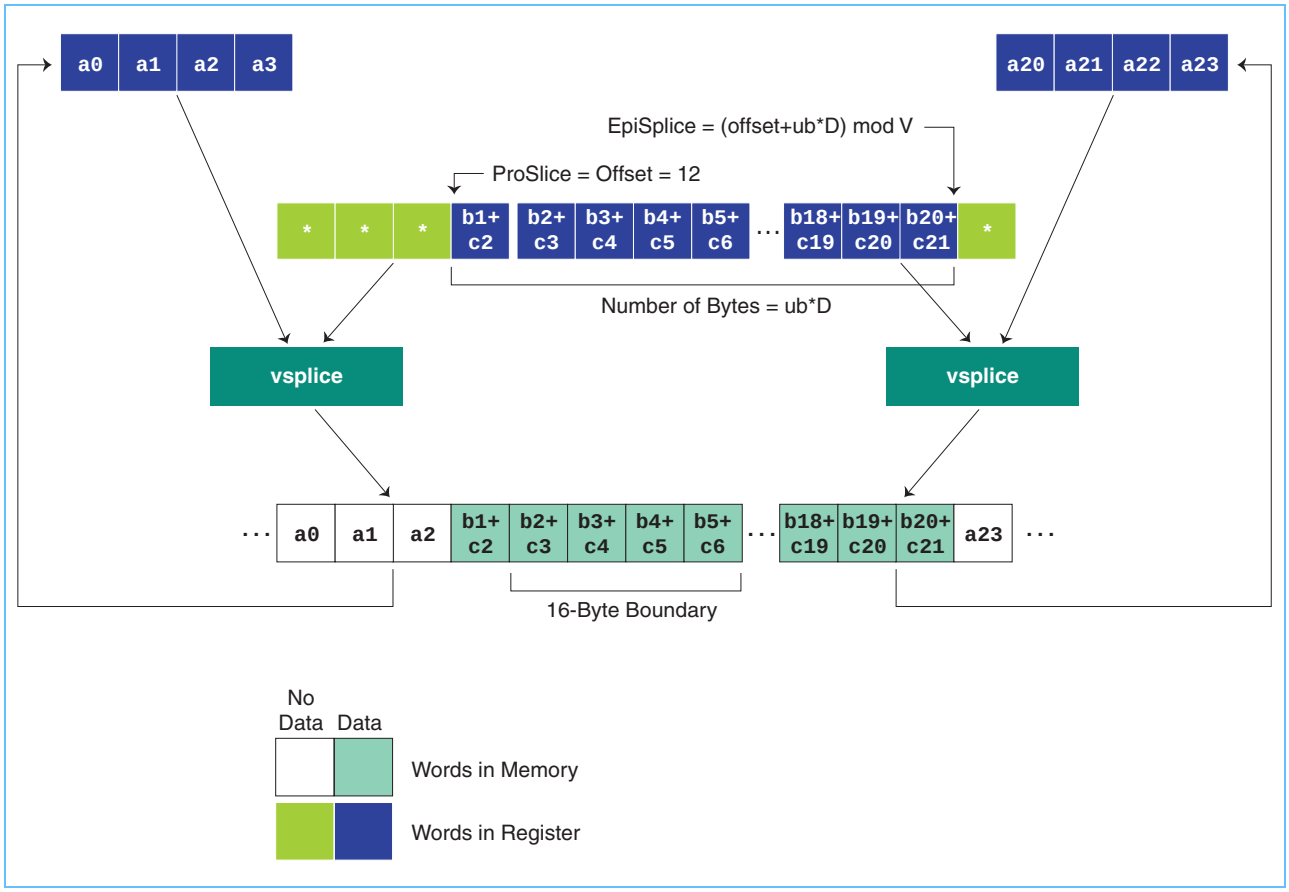
### 22.3.4.7 *Handling a Partial Store*

The first and last iterations of a loop often need to be handled separately, in a prolog (*Section 22.1.2.6* on page 642) and epilog (*Section 22.1.2.7* on page 643).

Partial stores are implemented by loading the original value before the store, splicing it with the newly computed value, and then storing the spliced value back into memory, using the `vsplice` operation. *Figure 22-18* shows how partial stores are implemented by first loading the original value before the store. Then, this value is spliced with the newly computed value. Finally, the spliced value is stored back into memory, using the compiler-internal function, `vsplice`. At the end of phase 4, the vectors are aligned, and their lengths are multiples of the physical vector length.

The vector/SIMD multimedia extension store vector left and right instructions (**lvlx[l]**, **lvrx[l]**, **stvlx[l]**, and **stvrx[l]**) can also be useful for simplifying partial-vector stores.

*Figure 22-18. Implementing Partial Vector Store*




### 22.3.5 Phase 5: Length Devirtualization

In Length Devirtualization, virtual vectors are first flattened to vectors of primitive types. Then, operations on the virtual vectors are either mapped to operations on multiple physical vectors, or rewritten as operations on scalar values. The decision of whether to map a virtual vector onto physical vectors or onto scalars is based on the length of the virtual vector and other heuristics.

At the end of phase 5, all vectors have become physical vectors. They are aligned and operate on 16-byte data.

### 22.3.6 Phase 6: SIMD Code Generation and Instruction Scheduling

#### 22.3.6.1 *Code Generation*

The Code Generation phase maps the SIMDized operations, including the data-reorganization operations, to SIMD instructions specific to the target platform (either vector/SIMD multimedia extension or SPE). Code generation takes the augmented expression tree (the expression tree plus data-reorganization operations) as input, and maps the generic stream-shift operations to native SIMD permutation instructions. *Figure 22-19* on page 665 shows the basic mechanism to shift a stream with offset 4 to offset 0. In this case, the stream is shifted left by 4 bytes. The compiler-internal function, `vshiftstream`, is mapped to a sequence of shift operations across pairs of vector registers. These shift operations are called `vshiftpair` operations; `vshiftpair` is another compiler-internal function.

The code-generation and earlier phases of the framework also tackle four difficult problems:

- Runtime alignment
- Unknown loop bounds
- Multiple misalignments
- Multiple statements (a loop body that contains more than one statement)

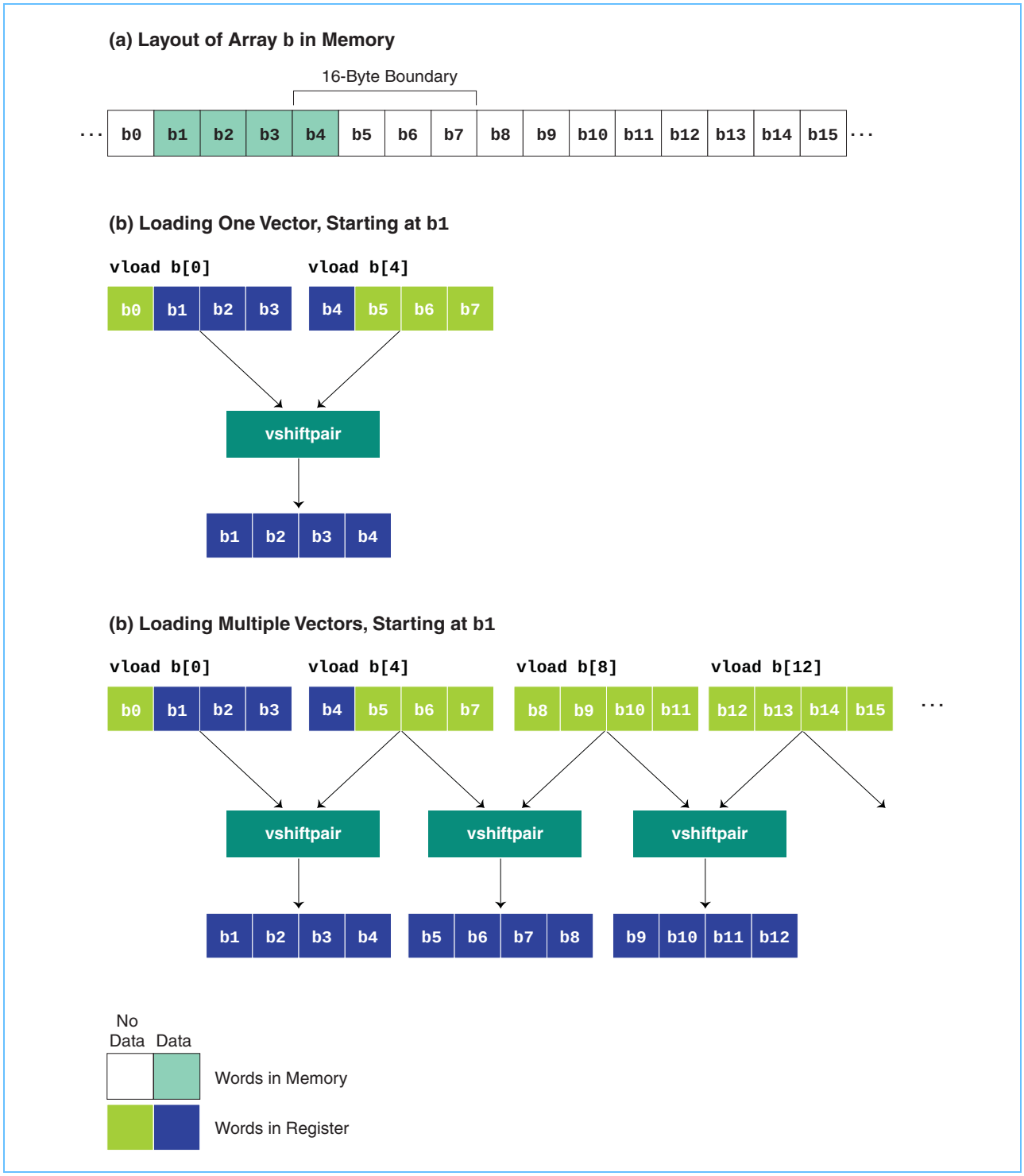
When memory alignment cannot be known at compile-time, it is referred to as *runtime alignment*. Runtime alignment can be inherent to a program, or it can result when the program uses complex pointer arithmetic. To handle runtime alignment, transform the original memory streams and computation streams into equivalent streams that can be shifted to arbitrary alignments. An arbitrary stream-shift, denoted by `vshiftstream(stream, left, right)`, can be converted to a left stream-shift to the required offset. The derived stream starts exactly `left` bytes before the first value of `stream`.

Here are some guidelines for applying data-reorganization operations for runtime alignment:

- Lazy-Shift, Eager-Shift, and other policies can be applied to runtime alignment.
- When propagating stream offsets, that is, propagating the alignment of the store to the load during an Eager-Shift, the stream offset needs to be scaled up or down. The stream offset is scaled by the packing factor.
- Avoid placing a stream-shift that causes an unpack between a register stream and a memory stream, with no other stream-shifts in between.
- When possible, place stream-shifts at the shorter end of a data conversion. It takes fewer machine instructions to shift a stream of shorter length.

### 22.3.7 SIMDization Example: Multiple Sources of SIMD Parallelism

*Figure 22-19. Implementing Stream Shift*




This section demonstrates how to extract parallelism at the basic-block and loop levels. The original code sample, which follows, is a simple loop with four statements. This code is followed by

three other samples—basic-block aggregation, loop aggregation, and length devirtualization. These code samples are truncated, so that they highlight the data elements on which we focus.

The focus is only on the stores generated by each of the loop's statements, which are illustrated in *Figure 22-20* on page 667, although the compiler would take all memory references into account. Each 4-byte store is represented by a box in *Figure 22-20*. The sets of boxes alternative in color to distinguish consecutive iterations of the original loop.

*Original Code:*

```
for (i = 0; i < n; i++) {
    a[i].x =
    a[i].y =
    a[i].z =
    b[i] =
}
```

*Step 1: Basic-Block Aggregation:*

```
for (i = 0; i < n; i++) {
    (a[i].x,y,z)   =
    b[i]           =
}
```

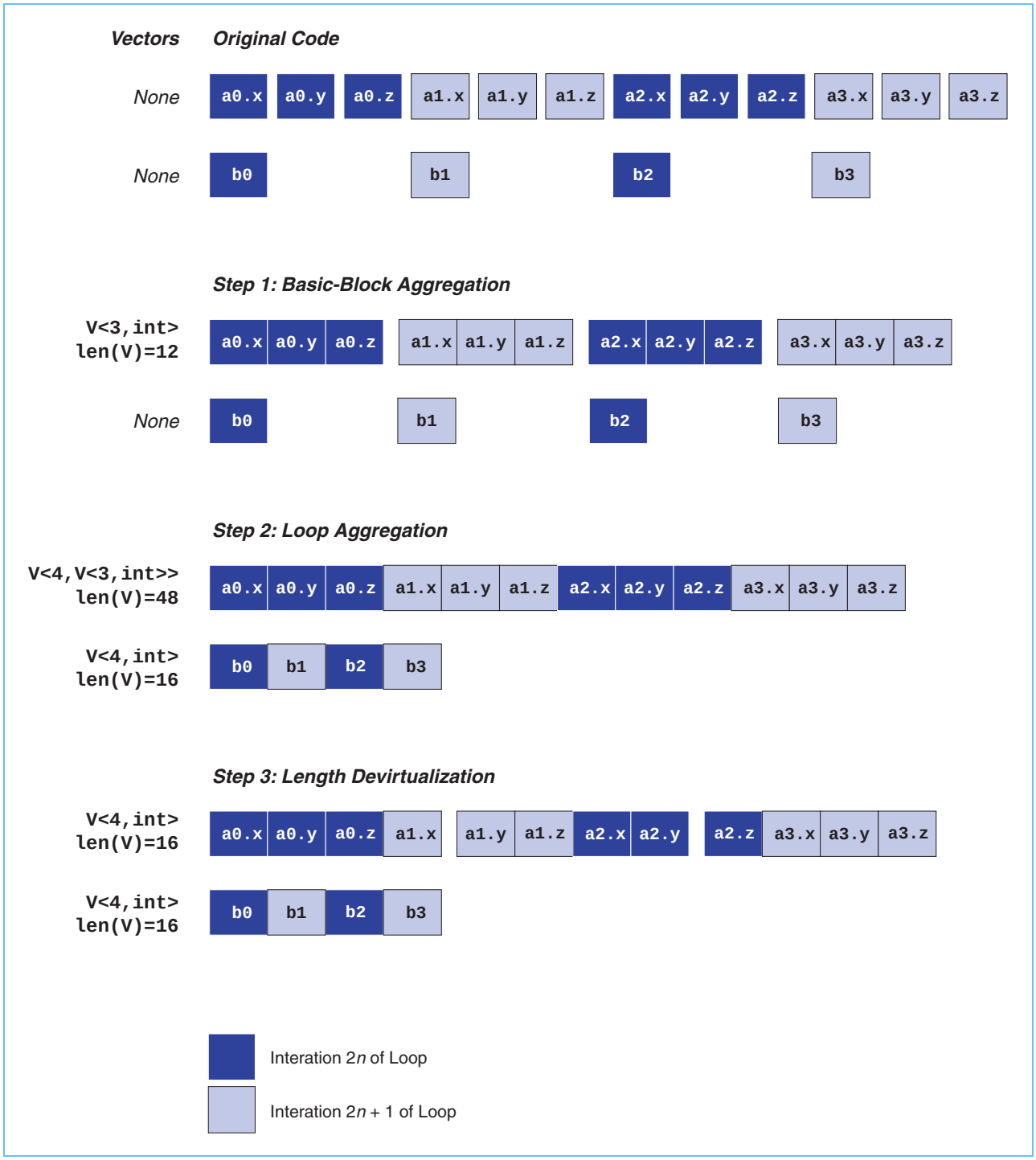
*Step 2: Loop Aggregation:*

```
for (i = 0; i < n; i+=4) {
    (a[i].x...a[i+3].z)   =
    (b[i].......b[i+3])   =
}
```

*Step 3: Length Devirtualization:*

```
for (i = 0; i < n; i+=4) {
    (a[i].x.....a[i+1].x) =
    (a[i+1].y...a[i+2].y) =
    (a[i+2].z...a[i+3].z) =
    (b[i].......b[i+3])   =
}
```

*Figure 22-20. SIMDization Example with Mixed Stride-One and Adjacent Accesses*




The first step is basic-block aggregation. The compiler recognizes that the first three stores (`a[i].x a[i].y, a[i].z`) are adjacent in memory. Assuming the right-hand sides of the three statements are isomorphic, the compiler aggregates the three statements into a vector of three integers stored into `a[i].x, y, z`. The `b[i]` statement remains unchanged.

The second step is loop-level aggregation. Recognizing the vector store `a[i].x,y,z` and the element store `b[i]` as stride-one accesses, the compiler aggregates these accesses across loop iterations. It treats the new vector `a[i].x,y,z` statement no differently than any other statements in the loop. The only difference between the `a[i].x,y,z` and `b[i]` statements is that the first one generates a 12-byte value but the second one generates a 4-byte value.

During loop-level aggregation, the compiler extracts SIMD parallelism from the smallest number of consecutive iterations, while ensuring that each vector in the loop has a length that is a multiple of the physical vector length: 16 bytes. The optimal blocking factor is 4, because it aggregates 4 of the 12-byte `a{i}.x,y,z` vectors into a new compound vector of 48 bytes. Four of the 4-byte `b[i]` values aggregate into a new vector of 16 bytes.

An alignment-devirtualization step might be needed. This step is necessary when arrays `a` and `b` are not aligned. Because both `a[i].x ... a[i+3].z` and `b[i] ... b[i+3]` are stride-one accesses, their alignment constraints are handled no differently than any stride-one access of native data types. For simplicity, our example and *Figure 22-20* on page 667 assume both vectors are aligned. But suppose that one of them were not; for example. suppose a[i].y is 16-byte aligned instead of `a[i].x`. Then, the alignment devirtualization phase logically skews the computation associated with vector `a` to the right by one value. As a result, the blocked loop effectively computes and stores the `a[i].y ... a[i+4].x` instead of `a[i].x ... a[i+3].z`.

The third step (in this example) is length devirtualization. The long `a[i].x ... a[i+3].z` vector is broken down to three physical vectors. Each of the three physical vectors mixes original stores from two consecutive iterations. At this stage, all vectors have become physical vectors, meaning that they are aligned and operate on 16-bytes of data.

In this example, the basic-block and loop-level aggregation steps each handle some aspect of SIMDizing the loop, but neither can SIMDize the loop alone:

- On one hand, basic-block SIMDization combines the otherwise non-stride-one accesses `a[i].x, a[i].y, a[i].z` into stride-one accesses that can be subsequently exploited in loop-level aggregation.
- On the other hand, loop-level SIMDization aggregates the unaligned vector of 12-bytes, `a[i].x,y,z` produced by basic-block aggregation into an aligned vector of a multiple of the physical vector length.

This is achieved transparently within the proposed SIMDization framework without expensive loop restructuring such as loop distribution, loop rerolling and collapsing, and loop unrolling.

### 22.3.8 SIMDization Example: Multiple Data Lengths

There are two types of loops that mix data of different lengths—for example, `char`, `short int`, and `float`. The first type has homogenous statements: statements that each operate on data of uniform length. Some of the homogenous statements do operate on types of different lengths. The second type has statements that operate on distinct data lengths within a unique statement. This code example demonstrates such heterogeneous statements:

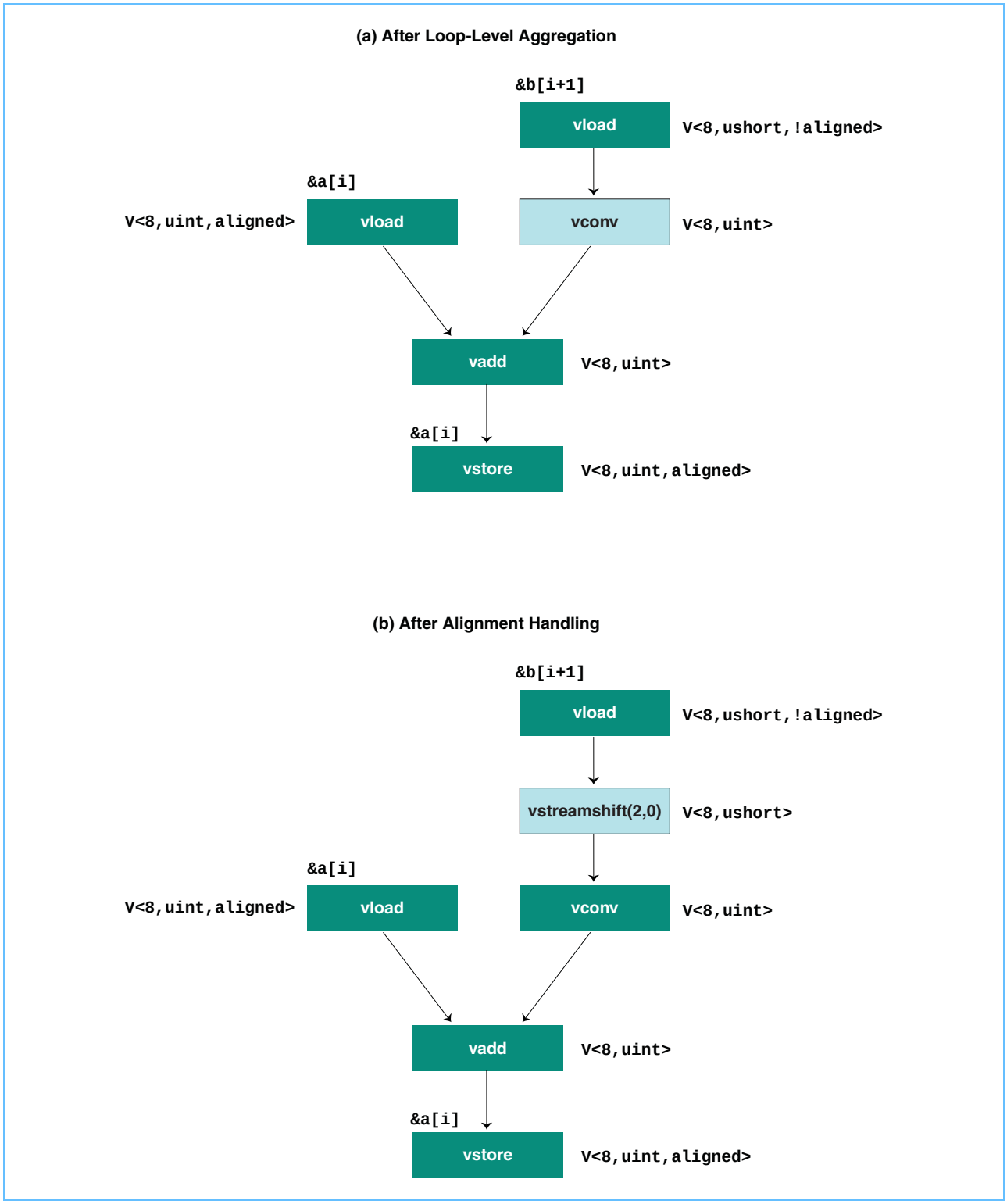
```
uint a[N];
ushort b[N];
for (int i = 0; i < N; i++) {
    a[i] += (uint) b[i+1];
}
```

Such heterogeneous statements often require conversion operations between data of different lengths. This is referred to as *length conversion*. Loops that require length conversion among heterogeneous statements are more constraining than loops that contain homogenous statements. The complexity is due to the fixed-length of SIMD vectors. For example, a SIMDized length conversion from 1-byte length to 2-byte length consumes one vector and produces two vectors. Consuming one vector but producing multiple vectors (or a fraction of a vector, if the conversion is the other way around) is unnatural for most compiler intermediate representations. The example that follows focuses on the more challenging length conversion among heterogeneous statements.

#### 22.3.8.1 *Alignment Handling and Alignment Devirtualization*

*Figure 22-21* on page 670 shows the expression tree for the preceding code after the first three phases (basic-block, short-loop, and loop-level aggregation).

*Figure 22-21. Length Conversion After Loop-Level SIMDization*



In the tree representation, each node is labeled with the tuple `V(n, t, aligned)` to represent the number of elements, the type, and the alignment properties of each vector. Because `aligned` only affects vector load and store semantics, it is specified only for memory nodes.

For alignment devirtualization, each array accessed in the loop must be stride-one to form a stream of contiguous memory accesses across iterations. The alignment-devirtualization algorithm is mainly concerned with code generation problems using stream-shift operations with length conversion.

In our example, assume the base addresses `a` and `b` are aligned. Assume the streams represented by `a[i]` have an offset of 0 and the streams represented by `b[i+1]` have an offset of 2. These streams are therefore relatively misaligned. *Figure 22-21(b)* shows the alignment tree after alignment handling. The node `vstreamshift(2, 0)` represents a pseudo operation (a compiler-internal function) that shifts the entire input stream of `b[i+1]` with an offset of 2 to 0, to match the offset of the `a[i]` stream.

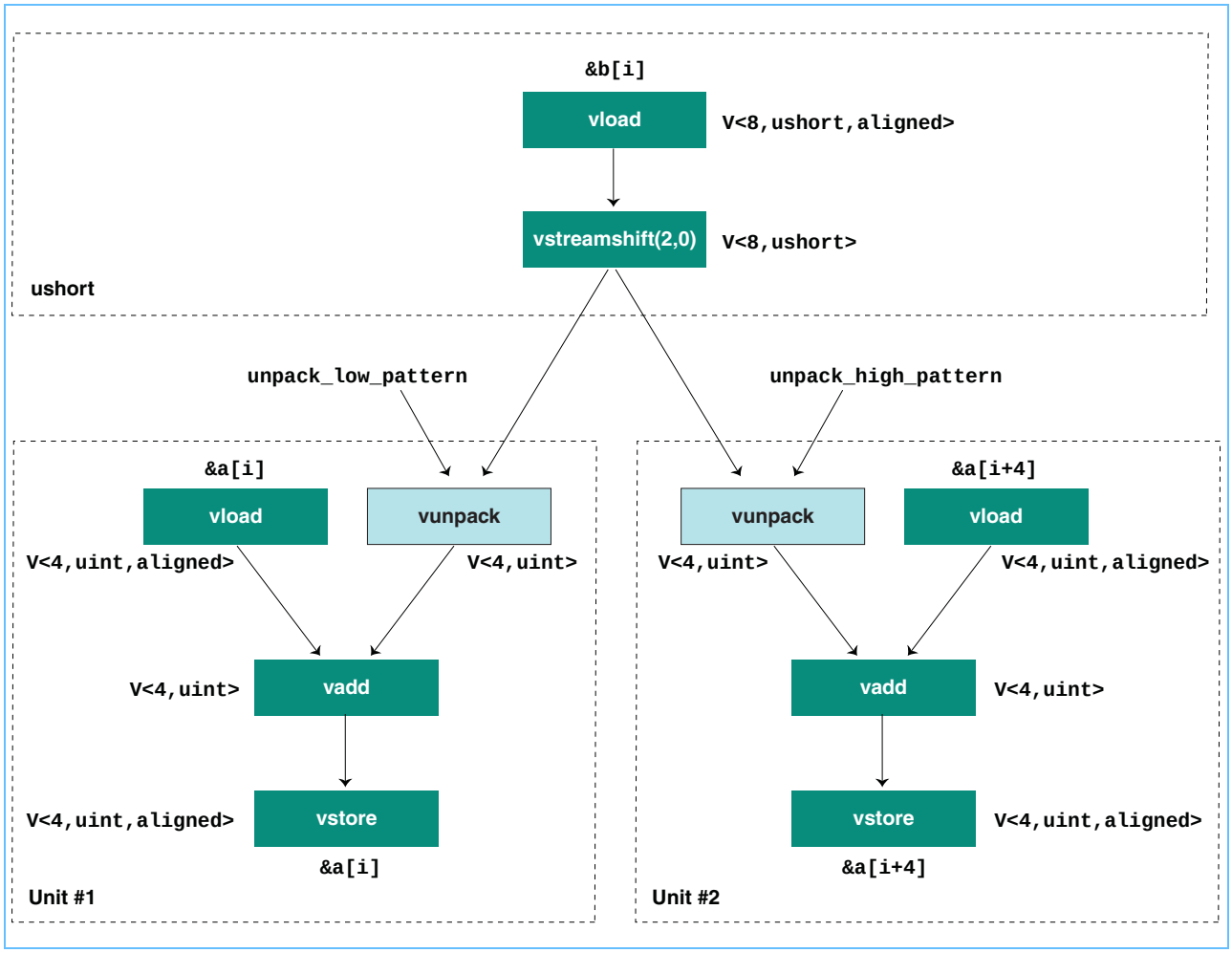
#### 22.3.8.2 *Length Devirtualization*

For most vector operations, devirtualization simply involves breaking a vector into multiple physical vectors. Length conversion, however, requires special handling because a devirtualized length conversion might consume one physical vector and produce multiple physical vectors, or vice versa. The two generic operations, `pack` and `unpack`, are used to express such data reorganization:

- `pack(v1, v2, vpattern)`—This operation produces a physical vector by selecting bytes from physical vector `v1` and `v2` according to the byte pattern specified by physical vector `vpattern`.
- `unpack(v, vpattern)`—This operation produces a physical vector by selecting bytes from physical vector `v` according to the byte pattern specified by physical vector `vpattern`.

In our example, the conversion operation unpacks a vector of `short` (16-bit integers) into two vectors of `int` (32-bit integers). *Figure 22-22* on page 672 shows the expression trees after devirtualization. This SIMDized code consists of three distinct parts. The top part represents the computations done with the variables of type `short`, namely loading from memory. The bottom two parts represent the computations done with the variables of type `int`, with the left and right bottom parts each consuming half the data produced by the top part.

*Figure 22-22. Length Conversion After Devirtualization*




### 22.3.8.3 *SIMD Code Generation*

The generic `pack` and `unpack` operations are mapped to instructions of the target platform. Integrating conversion-handling into the SIMDization framework is made much easier because stride-one accesses are preserved in early phases of SIMDization. As shown in *Figure 22-22*, when devirtualizing the `short` to `int` conversion, the stride-one access `a[i]` in the original loop becomes two non-stride-one vector accesses, `a[i:i+3]` and `a[i+4:i+7]`, in the SIMDized loop. Because stride-one accesses are critical both to aggregation phases and alignment handling, by maintaining the stride-one accesses until devirtualization, the early SIMDization phases can process length-conversion in the same way that other regular operations are processed.

### 22.3.9 Vector Operations and Mixed-Mode SIMDization

The concept of virtual vectors and devirtualization—phase 4 in *Section 22.3.4* on page 658, and phase 5 in *Section 22.3.5* on page 663—can be extended to vector operations. During early phases of SIMDization, a vector operation can be SIMDized even if there is no direct support in the underlying SIMD hardware. When these virtual vector operations are devirtualized, they can be replaced by:

- A single SIMD instruction
- A sequence of SIMD instructions
- A library call
- A sequence of scalar operations

This process is called *mixed-mode SIMDization*. A loop that contains computation that can be executed on the scalar units, or on the SIMD unit, or both, is a candidate for mixed-mode SIMDization. For example, on the PPE's vector/SIMD multimedia extension unit, double-precision floating-point instructions are executed in scalar units, but saturate-arithmetic is executed on the SIMD (VXU) unit. Most fixed-point instructions can execute on both units. Because the VXU unit is more constrained than the scalar unit, it can be more efficient to compute in scalar mode if accesses are misaligned or not stride-one.

Loop distribution is one way in which heterogeneous loops (loops containing a mixture of scalar and vector computation) can be handled. Loop distribution might cause large loops be distributed into smaller ones. More loops with shorter loop bodies works against finding instruction-level parallelism, because it decreases the size of basic blocks. Overly aggressive loop distribution is one of the major causes of SIMDization-related performance degradation.

Mixed-mode SIMDization avoids the problem of loop distribution affecting ILP, and might help ILP as the scheduler can use the SIMD and scalar units (and registers) at the same time. It is not always straightforward to decide which mix of SIMD and scalar computation is best. Both hardware constraints and resource constraints of the currently executing program need to be taken into account. In addition, the VXU and scalar units can be used for the same operation at the same time.

Virtualizing vector operations also provides a way to change transformations made earlier, if during later analysis, when more information is available, it becomes clear that those transformations did not improve the code. For example, during basic-block SIMDization, non-stride-one memory accesses can be packed into vectors, hoping that these vectors become stride-one after loop-level aggregation. But if that does not happen and the packing overhead is high, the compiler can undo the overly optimistic packing decision by devirtualizing the vector operation back to scalar operations.

## 22.4 Other Compiler Optimizations

### 22.4.1 OpenMP

A compiler can support OpenMP directives and routines. OpenMP is a platform-neutral and language-neutral way for programmers to insert parallelization directives in source code. OpenMP is defined by a consortium of universities and vendors, and the specification is found at http://www.openmp.org. See *Section 21 Parallel Programming* on page 609 for a description of how to use OpenMP constructs in programs run on the CBEA processors.

Most of the OpenMP directives are orthogonal to SIMDization, and they should not interfere with SIMDization. Some OpenMP directives do parallelize loops in the way that SIMDization does, and there can be some side effects. If the OpenMP directives are processed first, the loop structures can be perturbed in a way that prevents SIMDization. One of the OpenMP directives is the *parallel for* construct. This is used to parallelize a simple loop. For example:

```
void demo(int n, float *a, float *b) {
   int i;

#pragma omp parallel for

   for (i = 1; i < n; i++)  /* i is private, by definition */
      b[i] = (a[i] + a[i-1]) / 2.0;
}
```

It is most important for SIMDization that OpenMP not corrupt the alignment information. When compile-time alignment is available for arrays before OpenMP directives are processed, this compile-time alignment should remain in force after processing OpenMP directives. Whenever possible, keep array subsections in multiples of 16 bytes.

### 22.4.2 Subword Data Types

Both 2-byte and 1-byte data types are called *subword* data types. The C programming language data types `char` and `short` are 1 and 2 bytes, respectively. From the architecture's perspective, not all data types are equally supported. For example, the SPU Instruction Set Architecture (ISA) provides 16-bit multiply only, so 32-bit multiply is implemented in software using 16-bit multiplies. Subword data types are commonly used in multimedia and games programs, so the SIMD units are designed to provide better support for subword data types.

The C programming language requires that subword data types be promoted to a full word before being operated on. A word is 32-bits or 4-bytes; the C data type `int` corresponds to a word. C allows a programmer to declare a subword variable, in terms of alignment, size and layout in memory—but there is no subword arithmetic. When a `short` or `char` is loaded into a register, the values are converted into 32-bit integers by sign-extensions. Without additional optimizations, there is no way to generate subword instructions directly from scalar C source programs.

Subword arithmetic optimizations bridge the gap between C's inability to specify subword arithmetic and the underlying hardware's strength in supporting subword operations. There are two steps:

- *Eliminate Redundant Partial-Copy Operations (from sign extensions).* This relies on partial-copy propagation. Because subword instructions use only a partial region of a register and sign-extension instructions copy only a partial region of a register, redundant sign-extensions are readily removed. In addition, this step performs general constant folding. Constant folding refers to the compiler precalculating constant expressions.
- *Generate Subword Arithmetic Operations.* This converts integer arithmetic operations to the equivalent arithmetic operation on narrower-width data types. This produces more efficient code when the ISA provides subword instructions, but not the equivalent instructions on 32-bit data types. In addition, because this optimization narrows the operating data, it exposes more opportunity for eliminating redundant sign extensions.

### 22.4.3 Backend Scheduling for SPEs

For SPEs, there are additional code-generation issues that must be handled for both automatically SIMDized instructions and scalar instructions. Instruction scheduling occurs just before the compiler emits the final assembly or object code. The scheduling process consists of two closely interacting subtasks: scheduling and bundling. While typical schedulers deal with resources and latencies, a highly optimizing scheduler takes into account constraints that are expressed in numbers of instructions. For example, the hint for branch (HBR) instruction cannot be more than 256 instructions away from its target branch, and should be no closer than 8 instructions. (See *Section 24.3.3 Branch Hints* on page 701.) The scheduling is complicated by the fact that the number of instructions in a scheduling unit is known only after the bundling subtask has completed.

The bundler's main role is to ensure that each pair of instructions that are expected to be dual-issued satisfies the SPE's instruction-issue constraints. The processor will dual-issue independent instructions only when the first instruction uses the even pipeline and resides in an even word address, and the second instruction uses the odd pipeline and resides in an odd word address. After instruction ordering is set by the scheduling subtask, the bundler can only impact the even or odd word address of instructions by inserting `nop` instructions into the instruction stream.

Another important task of the bundler is to prevent instruction-fetch starvation. Recall that a single local memory port is shared by the instruction fetch mechanism and the processor's memory instructions. As a result, a large number of consecutive memory instructions can stall instruction fetching. With 2.5 instruction-fetch buffers reserved for the fall-through path, the SPE can run out of instructions in as few as 40 dual-issued cycles. When a buffer is empty, there might be as few as 9 cycles for issuing an instruction-fetch request to still hide its full 15-cycle latency. Because the refill window is so small, the bundling process must keep track of the status of each buffer and insert explicit `ifetch` instructions when a starvation situation is detected.

The refill window is even smaller after a correctly-hinted branch that is taken, because there is only one valid buffer after a branch. This compares to 2.5 buffers for the fall-through path. Instruction starvation is prevented only when all instructions in the buffer are useful. That is, the branch target must point to an address that is a multiple of 16 instructions, which is the alignment constraint of the instruction-fetch unit. This constraint is enforced by introducing additional `nop` instructions before a branch target, to push it to the next multiple of 16 instructions. These compiler heuristics are successful in scavenging any idle issue slots, so that `nop` instructions can be inserted without performance penalties.

A final concern of the bundling process is to make sure that there is a sufficient number of instructions between a branch hint and its branch instruction. This constraint is due to the fact that a hint is only successful if its target branch address is computed before that branch enters the instruction-decode pipeline. The bundler adds extra `nop` instructions when the schedule did not succeed in interleaving a sufficient number of independent instructions between a hint and its branch.

A unified scheduling and bundling phase produces best performance. The compiler can preserve the cycle scheduling approach, in which each nondecreasing cycle of the schedule is filled in turn. The exception is that the bundling process may retroactively insert `nop` or `ifetch` instructions. When scheduling the next cycle in the scheduling unit, the compiler checks to determine whether an `ifetch` instruction is required. If it is, an `ifetch` instruction is inserted, and the scheduling resource model is updated. The normal scheduling process then continues. When no additional instructions can be placed in the current cycle, the compiler checks to determine whether `nop` instructions must be inserted in previous cycles to enable dual-issuing.

### 22.4.4 Interacting with Typical Optimizations

SIMDization should occur after aggressive dependence analysis and high-level loop transformations. These loop transformations enhance SIMDization by removing loop-carried dependences among inner-most loops:

- Loop interchange, which swaps an inner (nested) loop with an outer loop. This is done when the inner loop appears to be the candidate for SIMDization.
- Loop distribution, which splits off those parts that cannot be SIMDized.
- Loop versioning, a technique for hoisting an individual array index exception check outside a loop by providing two copies of the loop: the safe loop and the unsafe (original) loop.

SIMDization can be followed by other loop transformations, including:

- Loop fusion.
- Loop unrolling, which removes needless copy operations.
- Predictive commoning, a special form of common subexpression elimination that exploits data reuse among consecutive loop iterations.

Some optimizations are not suitable for SIMDization:

- Loop rerolling converts basic-block parallelism to loop-level parallelism by rerolling statements into loops. The drawback is that a new inner-most loop is added to the loop nest. Because loop-level SIMDization happens at the inner-most loop, rerolling does not expand the scope of SIMDization. In addition, rerolling requires a loop counter, which is not always known.
- Loop collapsing rerolls statements into a loop and then collapses it to the inner-most surrounding loop. Loop collapsing exploits parallelism at both the basic-block level and the loop level. However, it is more constraining because it requires the original inner-most loop to be perfectly nested in its surrounding loop, and it requires all its statements to be rerolled.
- Loop unroll-and-pack converts loop-level parallelism to basic-block parallelism by unrolling the loop and packing the computation in a basic block with SIMD instructions. Because SIMDization is limited to a basic block in this case, it is very inefficient to handle misalignment

and reduction. In addition, packing algorithms are more expensive and randomized, and sometimes perform less well, than loop-level SIMDization based on dependence vectors.

# 23. Vector/SIMD Multimedia Extension and SPU Programming

Both the PowerPC Processor Element (PPE) and the Synergistic Processor Elements (SPEs) support very similar sets of vector instructions. The PPE processes single instruction, multiple data (SIMD) operations in the vector/SIMD multimedia extension unit within its PowerPC processor unit (PPU). The operations are those of the vector/SIMD multimedia extensions to the PowerPC instruction set. The SPEs process SIMD operations in their synergistic processor unit (SPU). The operations are those of the SPU instruction set. Both instruction sets operate on 128-bit SIMD vectors and both support many of the same types of SIMD operations.

This section summarizes key similarities between these instructions and intrinsics, with the goal of identifying strategies for developing code that is portable between the PPE and the SPEs.

## 23.1 Architectural Differences

The objectives of the vector/SIMD multimedia extension and SPU architectures are essentially the same—to perform well on data-intensive processing, such as graphics pipelines, stream processing (encoding, decoding, encryption, decryption), and physical modeling. The architectural resources, however, that each architecture brings to bear differ in some important respects. Some of the major SIMD-support differences between the PPE and SPE architectures are summarized in *Table 23-1*.

*Table 23-1. PPE and SPE SIMD-Support Comparison*

| Feature | PPE | SPE |
|---|---|---|
| Number of processing elements | 1 | 8 |
| Modes supported | user and supervisor | user only |
| Number of SIMD registers | 32 (128-bit) | 128 (128-bit) |
| Organization of register files | separate fixed-point, floating-point, and SIMD registers | unified SIMD registers |
| Load latency | variable (cached) | fixed |
| Addressability | $2^{64}$-byte main storage | 256 KB local storage (LS)<br>$2^{64}$-byte main storage via DMA |
| Memory architecture | 2-level caching | software-controlled LS |
| SIMD instruction set | general SIMD, supported by PowerPC scalar and control instructions | SIMD only, optimized for single-precision floating-point, 16-bit fixed-point, and 32-bit fixed-point |
| Single-precision floating-point SIMD | IEEE 754-1985 and SPE-compatible graphics-rounding mode supported | extended range[1] |
| Double-precision floating-point SIMD | not supported | IEEE 754-1985 supported |
| Doubleword fixed-point SIMD | not supported | supported |
| Interelement communication facility | memory mapped I/O (MMIO) | channel |

1. See *Section 3.1.4* on page 70.

The PPE excels at control processing. Unlike the SPEs, which supports only problem (user) state, the PPE supports problem (user) state, privileged (supervisor) state, and hypervisor state. All of the PPE instructions—standard PowerPC and vector/SIMD multimedia extensions—can be intermixed in a single program. The PowerPC and vector/SIMD multimedia extensions architectures are also supported by well-developed tools for debug and performance analysis, although this advantage is changing as new tools become available for the SPE. Perhaps most importantly, there is a substantial quantity of existing PowerPC and vector/SIMD multimedia extension code.

The eight SPEs provide a significant boost in execution resources for both scalar and SIMD processing, as compared to the single PPE. SPU instructions provide excellent support for single-precision floating-point graphics applications. Several SPU instructions are available for eliminating inefficient execution penalties typically associated with branch execution. Blocking channels for communication with the memory flow controller (MFC) or other parts of the system external to the SPU, provide an efficient mechanism for waiting on the completion of external events without the need for polling or interrupts, both of which burn power needlessly.

### 23.1.1 Registers

#### 23.1.1.1 *PPE Registers Used by Vector/SIMD Multimedia Extension Instructions*

The PPE has 32 Vector Registers (VRs) for holding vector/SIMD multimedia extension instruction operands, as described in *Section 2.3* on page 54. It also has 32 General-Purpose Registers (GPRs) for holding fixed-point scalar operands and 32 Floating-Point Registers (FPRs) for holding floating-point scalar operands.

In addition to these instruction-operand register files, the PPE also has three control and status registers that are affected by vector/SIMD multimedia extension operations—a 32-bit Vector Status and Control Register (VSCR), a 32-bit Vector Save Register (VRSAVE), and a 32-bit (8-element) Condition Register (CR). The VSCR register has two defined bits—non-Java-mode and saturation [SAT]. The VRSAVE register may be used in accordance with the application binary interface (ABI) for saving and restoring vector/SIMD multimedia extension state across context switches. The CR6 field of the CR register is useful for branch control; the bit is set or cleared according to the result of a vector/SIMD multimedia extension compare instruction that is of the recording form (a compare instruction that has a dot suffix on the mnemonic).

The VSCR[SAT] bit is set if saturation (an inexact result) occurs during a saturating SIMD operation. To use this bit, the contents of the VSCR can be copied to a VR (using the **mfvscr** instruction), bits other than the SAT bit can be cleared in the VR (**vand** with a constant), the result can be compared to zero so as to set CR6 (**vcmpequb.**), and a PowerPC conditional branch instruction can test CR6.

The PPE's Floating-Point Status and Control Register (FPSCR) is not affected by vector/SIMD multimedia extension floating-point operations. In general, no status bits are set to reflect the results of such operations, except that VSCR[SAT] bit can be set by the Vector Convert To Fixed-Point Word instructions.

#### 23.1.1.2 *SPE Registers*

Each SPE has 128 registers in one unified register file that holds all types of operands and instructions. In addition, the SPE also has a 128-bit Floating-Point Status and Control Register (FPSCR) that is updated after every floating-point operation to record information about the result and any associated exceptions.

The SPE architecture does not have a condition register comparable to the PPE's CR register. Instead, SPE comparison operations set result fields in the result vector that are either 0 (false) or -1 (true) and that are the same width as the operand elements being compared. These results can be used for bitwise masking, select instructions, or conditional branches.

The SPU instruction set includes many conditional branch instructions. An SPE conditional branch requires moving the vector element containing the branch condition into the preferred slot using a quadword rotate, performing a compare, and branching based on the result of the compare. By comparison, a vector/SIMD multimedia extension conditional branch requires a compare, a mask, a compare of CR6, and then a conditional branch, which is more expensive than an SPE's conditional branch.

### 23.1.2 Data Types

For the most part, vector/SIMD multimedia extension instructions operate on 128-bit vectors whose elements are byte, halfword, or word data types. The only area in which they do not is in the sourcing of a pointer from a GPR and writing a predicate[1] result to a GPR.

To this list of data types, the SPU architecture adds the doubleword data type. *Table 23-2* summarizes the supported vector data types for each architecture.

*Table 23-2. Vector/SIMD Multimedia Extension and SPU Vector Data Types*

| Data Type | Vector/SIMD Multimedia Extension | SPU |
|---|---|---|
| vector char (signed and unsigned) | yes | yes |
| vector bool char | yes | no[1] |
| vector short (signed and unsigned) | yes | yes |
| vector bool short | yes | no[1] |
| vector pixel | yes | no |
| vector int (signed and unsigned) | yes | yes |
| vector bool int | yes | no[1] |
| vector float | yes | yes |
| vector long long (signed and unsigned) | no | yes |
| vector double | no | yes |

1. The *C/C++ Language Extensions for Cell Broadband Engine Architecture* document does not define specific boolean vector types. Boolean vectors are supported as unsigned vectors.

1. A predicate is Boolean-logic term denoting a logical expression that determines the state of some variables. All the vector/SIMD multimedia extension predicates return their results to a GPR. They do this by performing a vector compare and then moving and masking CR6 to a GPR.

Among the important differences are:

- Only the vector/SIMD multimedia extension instructions supports pixel vectors.
- Only the SPU instructions support doubleword vectors.

The SPUs quadword data type is excluded from the list because it is a nonspecific vector data type instead of a specific vector data type. The quadword data type is used exclusively as an operand in specific intrinsics—those which have a one-to-one mapping with a single assembly-language instruction.

### 23.1.3 Instruction-Set Differences

The vector/SIMD multimedia extension and SPU instructions were architected to integrate seamlessly, with an emphasis on frequent and important data types. Thus, they have many similarities, in addition to their use of 128-bit vector operands. Both architectures support a large collection of fixed-point operations and 16-bit multiplies. Both support a shuffle or permute operation on each of 32 source bytes, although the vector/SIMD multimedia extensions also support several merge and replicate operations. Both support many rotates and shifts, but there are approximately twice as many choices among the SPU instructions.

Some details of their distinctive features follow. For summary lists of the instructions, see *Section A.3* on page 748 and *Section B.1* on page 771.

#### 23.1.3.1 *Vector/SIMD Multimedia Extension Instructions*

The vector/SIMD multimedia extensions provide these operations that are not directly supported by the SPU instruction set:

- Saturating math
- Sum-across
- $Log_2$ and $2^x$
- Ceiling and floor.
- Complete byte instructions

Vector/SIMD multimedia extensions support several vector-pack and unpack operations, although comparable operation can be performed with SPU instructions. Vector/SIMD multimedia extension compare instructions have an optional recording form that is useful for creating conditional-branch tests, as described in *Section 23.1.1.1* on page 680.

The CBEA processors support instructions with a graphics rounding mode. This mode allows programs written with vector/SIMD multimedia extension instructions to produce floating-point results that are equivalent in precision to those written in the SPU instruction set. In this mode, as in the SPU environment, the default rounding mode is r*ound to zero*, denormals are treated as zero, and there are no infinities or NaNs. For details, see *Section A.3.3* on page 752.

Mechanisms are available in both the PPE and SPE to hide memory-load latency. The PPE PowerPC instructions include data-cache block touch instructions that can cause data to be prefetched into the cache, plus cache-line locking instructions to prevent prefetched data from being cast out. However, the advantages of prefetching data are limited by the small number of

registers with which to operate on the data. The SPEs, in contrast, support DMA transfers in parallel with computation, and their 128 registers provide ample resources with which to operate on data. Thus, SPU loads might be better suited for deterministic real-time applications.

#### 23.1.3.2 *SPU Instructions*

The SPU instruction set provides these operations that are not directly supported by a single vector/SIMD multimedia extension instruction:

- Immediate operands
- Double-precision floating-point
- Sum of absolute difference
- Count ones in bytes
- Count leading zeros
- Equivalence
- AND complement
- OR complement
- Extend sign
- Gather bits
- Form select mask
- Integer multiply and accumulate
- Multiply subtract
- Multiply float
- Shuffle byte special conditions
- Carry and borrow generate
- Sum bytes across
- Extended shift range

SPU single-precision floating-point instructions execute at approximately twice the speed of vector/SIMD multimedia extension single-precision floating-point, and the SPU instructions—unlike the vector/SIMD multimedia extension instructions—also support double-precision vector operations.

The SPU Instruction Set Architecture (ISA) defines compare instructions for setting masks. These masks can be used in three-operand select instructions to create efficient conditional assignments to avoid difficult-to-predict branches. The SPU also supports hint for branch instructions (*Section 24.3.3* on page 701) that avoid branch penalties on taken branches when the branch can be predicted sufficiently early.

## 23.2 Porting SIMD Code from the PPE to the SPEs

It is sometimes easier to write SIMD programs by writing them first for the PPE than porting them to the SPEs. This approach postpones some SPE-related considerations—such as the size of the SPE's LS, data movements, and debug—until after the port. This approach can also allow partitioning of the work into simpler steps on the SPEs. Nevertheless, it is important to begin considering code partitioning for multiple SPUs early in the development of program architecture.

After vector/SIMD multimedia extension code is working properly on the PPE, a strategy for parallelizing the algorithm across multiple SPEs can be implemented. This is often, but not always, a data-partitioning method. The effort might involve converting from vector/SIMD multimedia extension intrinsics to SPU intrinsics, adding data-transfer and synchronization constructs, and tuning for performance. It might be useful to test the impact of various techniques, such as DMA double-buffering, loop unrolling, branch elimination, alternative intrinsics, number of SPEs, and so forth (see *Section 24* on page 691 for some examples). Debugging tools such as the static timing-analysis tool and the IBM Full System Simulator for the Cell Broadband Engine are available to assist this effort.

Alternatively, experienced CBEA processor programmers might prefer to skip the vector/SIMD multimedia extension coding phase and go directly to SPU programming. In some cases, SIMD programming can be easier on an SPE than the PPE because of the SPE's unified register file.

The earlier sections in this tutorial describe the vector/SIMD multimedia extension and SPU programming environments and some of their differences. Armed with knowledge of these differences, one can devise a strategy for developing code that is portable between the PPE and the SPEs. The strategy one should employ depends upon the type of instructions to be executed, the variety of vector data types, and the performance objectives. Solutions span the range of simple macro translation to full functional mapping.

### 23.2.1 Code-Mapping Considerations

There are several challenges associated with mapping code designed for one instruction set and compiled for another instruction set. These including performance, unmappable constructs, limited size of LS, and equivalent precision, as described in the following sections.

#### 23.2.1.1 *Performance*

Simplistic mapping of low-level intrinsics might not deliver the full performance potential of the SPE, depending upon the intrinsics used. Understanding the dynamic range of the mapping's operands—such as whether the range of operands is known to be restricted—can reduce the performance impact of simple mapping.

For example, consider mapping the **spu_shuffle** intrinsics (**shufb** instruction) to **vec_perm** (**vperm** instruction). If it was known than the shuffle-pattern bytes are all in the range x'0' through x'1F', then **spu_shuffle** can be directly mapped to **vec_perm**. Otherwise, the mapping requires several instructions.

#### 23.2.1.2 *Unmappable Constructs*

Differences in the processing of intrinsics make simple translation of certain intrinsics unmappable. The unmappable SPU intrinsics include:

- **`stop`** and **`stopd`**
- Conditional halt
- Interrupt enable and disable
- Move to and from status control and special purpose registers
- Channel instructions
- Branch on external data

#### 23.2.1.3 *Limited Size of LS*

Vector/SIMD multimedia extension programs mapped to SPU programs might not fit within the LS of the SPE, either because the program is initially too big or because mapping expands the code.

#### 23.2.1.4 *Equivalent Precision*

The SPU instruction set does not fully implement the IEEE 754 single-precision floating-point standard[2] (default rounding mode is r*ound to zero*, denormals are treated as zero, and there are no infinities or NaNs). Therefore, floating-point results on an SPE might differ slightly from floating-point results using the PPE's PowerPC instruction set. In addition, all estimation intrinsics (for example, ceiling, floor, reciprocal estimate, reciprocal square root estimate, exponent estimate, and log estimate) do not have equivalent accuracy on the SPU and PPE PowerPC instruction sets.

However, the instructions in the PPE's vector/SIMD multimedia extension have a *graphics rounding mode* (enabled by default) that is not available in the standard PowerPC Architecture. This mode is a special feature of the *Cell Broadband Engine Architecture*. It allows programs written with vector/SIMD multimedia extension instructions to produce floating-point results for add, subtract, multiply, multiply-add, and multiply-subtract that are equivalent in precision to those written in the SPU instruction set (but that differ from the estimate instructions). In this vector/SIMD multimedia extension graphics rounding mode, as in the SPU environment, the default rounding mode is r*ound to zero*, denormals are treated as zero, and there are no infinities or NaNs. For details on the graphics rounding mode, see *Appendix A.3.3 Graphics Rounding Mode* on page 752.

### 23.2.2 Simple Macro Translation

For many programs, it is possible to use a simple macro translation strategy for developing code that is portable between the vector/SIMD multimedia extension and SPU instruction sets. The keys to simple macro translation are described in the sections that follow.

2. The vector/SIMD multimedia extension instructions also do not fully implement the IEEE standard. The instruction sets are architected this way to better support applications such as graphics and real-time code. For details, see the *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual*.

#### 23.2.2.1 *Use a Compatible Vector-Literal Construction Format*

The vector/SIMD multimedia extension and SPU instruction sets support two styles of constructing literal vectors: curly brace (preferred style for SPE) and parenthesis (deprecated AltiVec style). Some compilers support both styles. A set of construction macros can be used to insulate programs from any differences in the tools.

There are multiple methods of constructing literal vectors, and all tools might not support the same construction format. The key differences in the methods are syntax (the deprecated AltiVec *parenthesis style* or the preferred *curly-braces style*) and handling of unspecified elements. The AltiVec style requires either a single element or all elements be specified. If only a single element is specified, then all elements of the vector are initialized with the specified value. The curly braces style allows one to specify any number of vector elements. The vector/SIMD multimedia extension and SPU instruction sets and most compilers support both syntax styles. Unspecified elements are cleared to '0'.

A simple set of construction macros can be used to insulate programs from any differences in the tools and assure portability of vector construction. For example:

```
#ifdef ALTIVEC_STYLE
    /* deprecated legacy support */
    #define VECTOR_UINT(a,b,c,d) \
        ((vector unsigned int)(a,b,c,d))
#else // CURLY_BRACE_STYLE
    /* preferred CBEA format */
    #define VECTOR_UINT(a,b,c,d) \
        ((vector unsigned int){a,b,c,d})
#endif
```

#### 23.2.2.2 *Use Single-Token Vector Data Types*

The *C/C++ Language Extensions for Cell Broadband Engine Architecture* document specifies a set of single-token vector data types, listed in *Table B-8* on page 784. Because these are single-token, the data types can be easily redefined by a preprocessor to the required target processor. Additional single-token data types have been standardized for the unique vector/SIMD multimedia extension data types. *Table 23-3* lists the additional data types. Also, see *Table 23-2* on page 681.

*Table 23-3. Additional Vector/SIMD Multimedia Extension Single-Token Data Types*

| Vector Data Type | Single-Token Data Type |
|---|---|
| vector bool char | vec_bchar16 |
| vector bool short | vec_bshort8 |
| vector bool int | vec_bint4 |
| vector pixel | vec_pixel8 |

#### 23.2.2.3 *Use Intrinsics that Map One-to-One*

Regardless of the technique used to provide portability, performance will be optimized if the operations map one-to-one between vector/SIMD multimedia extension intrinsics and SPU intrinsics. The SPU intrinsics that map one-to-one with vector/SIMD multimedia extension (except specific intrinsics) are shown in *Table 23-4*. The vector/SIMD multimedia extension intrinsics that map one-to-one with SPU are shown in *Table 23-5*.

*Table 23-4. SPU Intrinsics with One-to-One Vector/SIMD Multimedia Extension Mapping*

| SPU Intrinsic | Vector/SIMD Multimedia Extension Intrinsic | For Data Types |
|---|---|---|
| spu_add | vec_add | vector operands only, no scalar operands |
| spu_genc | vec_addc | all |
| spu_and | vec_and | vector operands only, no scalar operands |
| spu_andc | vec_andc | all |
| spu_avg | vec_avg | all |
| spu_cmpeq | vec_cmpeq | vector operands only, no scalar operands |
| spu_cmpgt | vec_cmpgt | vector operands only, no scalar operands |
| spu_convtf | vec_ctf | limited scale range (5 bits) |
| spu_convts | vec_cts | limited scale range (5 bits) |
| spu_convtu | vec_ctu | limited scale range (5 bits) |
| spu_madd | vec_madd | float only |
| spu_mulhh | vec_mule | all |
| spu_mulo | vec_mulo | halfword vector operands only, no scalar operands |
| spu_nmsub | vec_nmsub | float only |
| spu_nor | vec_nor | all |
| spu_or | vec_or | vector operands only, no scalar operands |
| spu_re | vec_re | all |
| spu_rl | vec_rl | vector operands only, no scalar operands |
| spu_rsqrte | vec_rsqrte | all |
| spu_sel | vec_sel | all |
| spu_sub | vec_sub | vector operands only, no scalar operands |
| spu_genb | vec_subc | vector operands only, no scalar operands |
| spu_xor | vec_xor | vector operands only, no scalar operands |

*Table 23-5. Vector/SIMD Multimedia Extension Intrinsics with One-to-One SPU Mapping* (Sheet 1 of 2)

| Vector/SIMD Multimedia Extension Intrinsic | SPU Intrinsic | For Data Types |
|---|---|---|
| vec_add | spu_add | halfwords, words, and floats only (not bytes) |
| vec_addc | spu_genc | all |
| vec_and | spu_and | all |

*Table 23-5. Vector/SIMD Multimedia Extension Intrinsics with One-to-One SPU Mapping* (Sheet 2 of 2)

| Vector/SIMD Multimedia Extension Intrinsic | SPU Intrinsic | For Data Types |
|---|---|---|
| vec_andc | spu_andc | all |
| vec_avg | spu_avg | unsigned chars only |
| vec_cmpeq | spu_cmpeq | all |
| vec_cmpgt | spu_cmpgt | all |
| vec_ctf | spu_convtf | all |
| vec_cts | spu_convts | all |
| vec_ctu | spu_convtu | all |
| vec_madd | spu_madd | all |
| vec_mulo | spu_mulo | halfwords only (not bytes) |
| vec_nmsub | spu_nmsub | all |
| vec_nor | spu_nor | all |
| vec_or | spu_or | all |
| vec_re | spu_re | all |
| vec_rl | spu_rl | halfwords and words only (not bytes) |
| vec_rsqrte | spu_rsqrte | all |
| vec_sel | spu_sel | all |
| vec_sub | spu_sub | halfwords, words, and floats only |
| vec_subc | spu_genb | all |
| vec_xor | spu_xor | all |

### 23.2.3 Full Functional Mapping

Because intrinsics have function-call syntax, code written for a specific functional unit can be mapped using inline functions. The keys to full functional mapping are:

- *C++ Function Overloading*—Because generic-intrinsic mappings are polymorphic (that is, one generic intrinsic maps to one of several instructions), C++ function overloading must be used to assure correct mapping.
- *Map Nonsupported Data Types to Unique Data Types*—Using C++ function overloading to map a generic intrinsic to a specific function requires that its parameter types be unique for every specific function. The vector/SIMD multimedia extensions support vector data types that are not valid on an SPU, and vice-versa. Therefore, these data types must be mapped to a supported type on the target processor element. *Table 23-6* on page 689 recommends a mapping that guarantees that the overloading function will uniquely select the correct mapped function without having to mandate that new data types be supported.

*Table 23-6. Unique SPU-to-Vector/SIMD Multimedia Extension Data-Type Mappings*

| SPU Data Type | Vector/SIMD Multimedia Extension Data Type |
|---|---|
| vector unsigned long long | vector bool char |
| vector signed long long | vector bool short |
| vector double | vector bool int |
| vector unsigned short | vector pixel[1] |

1. There are insufficient unique vector types on the SPU to provide a unique mapping for the *vector pixel* data type. This causes an overloading conflict when attempting to support the vec_unpackl and vec_unpackh intrinsics for both *vector signed short/vector bool short* and *vector pixel* types. This conflict can be resolved by using conditional compilation selection of the required mapping. Simultaneous mapping in a single object is not possible without further language extensions.

- *Use Intrinsics that Map One-to-One*—Use intrinsics that map one-to-one, as described in *Section 23.2.2.3* on page 687.

### 23.2.4 Code-Portability Typedefs

To improve code portability between PPE and SPU programs, single-token typedefs can be provided for the vector keyword data types. An example of such typedefs is shown in *Table B-8* on page 784. Because they are single-token, they serve as class names for extending generic intrinsics for mapping to-and-from vector/SIMD multimedia extension intrinsics and SPU intrinsics.

### 23.2.5 Compiler-Target Definition

To support the development of code that can be conditionally compiled for multiple targets, such as SPU and PPU, the *C/C++ Language Extensions for Cell Broadband Engine Architecture* specifies that compilers must define `__SPU__` when code is being compiled for the SPU. As an example, the following code supports misaligned quadword loads on both the SPU and PPU. The `__SPU__` define is used to conditionally select which code to use. The code that is selected will be different depending on the processor target.

```
vector unsigned char load_qword_unaligned(vector unsigned char *ptr)
{
    vector unsigned char qw0, qw1, qw;
#ifdef __SPU__
    unsigned int shift;
#endif
    qw0 = *ptr;
    qw1 = *(ptr+1);
#ifdef __SPU__
    shift = (unsigned int)(ptr) & 15;
    qw = spu_or(spu_slqwbyte(qw0, shift, spu_rlmaskqwbyte(qw1, (signed)(shift - 16)));
#else /* PPU */
    qw = vec_perm(qw0, qw1, vec_lvsl(0, ptr));
#endif
    return (qw);
}
```

# 24. SPE Programming Tips

This section provides general advice and code examples about how to write efficient code for the Synergistic Processor Element (SPE). Specific suggestions and examples are given for overlapping DMA and computation, eliminating and predicting branches, loop unrolling and pipelining, vectorizing a loop, and transformations and table lookups. References to more detailed descriptions of these and other topics are also given.

Among the most general guidelines for synergistic processor unit (SPU) coding are these:

- *Intrinsics*—Use intrinsics to achieve machine-level control without the need to write assembly-language code. Understand how the intrinsics map to assembly instructions and what the assembly instructions do. See *Section B* on page 771 for a summary of the instruction set and its C-language intrinsics, and see the *C/C++ Language Extensions for Cell Broadband Engine Architecture* document for complete details.
- *Overlap DMA Transfers with Computation*—Data transfers and computation can proceed in parallel by using multibuffering techniques. In addition, use the SPEs rather than the PowerPC Processor Element (PPE) to initiate most DMA transfers. See *Section 24.1*.
- *Fetch and Issue Rules*—An understanding of the fetch and issue rules will maximize your dual-issue rate. See *Section 24.2* on page 698.
- *Design for Limited Local Storage*—Each local storage (LS) is limited to 256 KB for program, stack, local data structures, and DMA buffers. Many optimization techniques put pressure on this limited resource. As such, all optimizations might not be possible for a given application. Often it is possible to reduce LS pressure by dynamically managing the program store using code overlays.
- *Branch Instructions*—Eliminate branches or reduce branch misprediction using hint-for branch instructions. See *Section 24.3* on page 699.
- *Loop Optimization*—Unroll and pipeline loops. See *Section 24.4* on page 709.
- *Integer Multiplication*—Use the SPU's 16 x16 bit multiplier to best advantage by avoiding 32-bit integer multiplies. See *Section 24.7* on page 716.
- *Table Lookups and Permutations*—Use the shuffle-bytes instruction for efficient table lookups and reorganization of operands. See *Section 24.6* on page 712.
- *Scalar Code*—Scalar (sub-quadword) loads and stores require multiple instructions and have long latencies. *Section 24.8* on page 716 has advice about efficient use of scalar code.

For a general overview of the SPE architecture, see *Section 3* on page 65.

## 24.1 DMA Transfers

This section builds on the basic description of DMA and DMA list transfer methods given in *Section 19.4* on page 529. It describes how to overlap DMA transfers with computation by using double-buffering and multi-buffering, plus techniques to improve DMA performance.

### 24.1.1 Initiating DMA Transfers from SPEs

Performance is typically best if you initiate DMA transfers from the SPEs that need the transfers rather than from the PPE. There are eight SPEs and only one PPE. The PPE can have only four in-flight data references (plus two for L2-cache loads and stores). Each of the eight SPEs, in contrast, can have up to 16 in-flight DMA transfers. Because of these considerations, the SPEs can initiate many more DMA transfers than the PPE in a given period of time.

### 24.1.2 Overlapping DMA Transfers and Computation

To hide data-access latencies, use double-buffering or multi-buffering techniques. You can either double-buffer the data (typical) or double-buffer the code, depending upon the code and data sizes and data-access patterns. This section outlines basic strategies for overlapping DMA transfers with computation on an SPU.

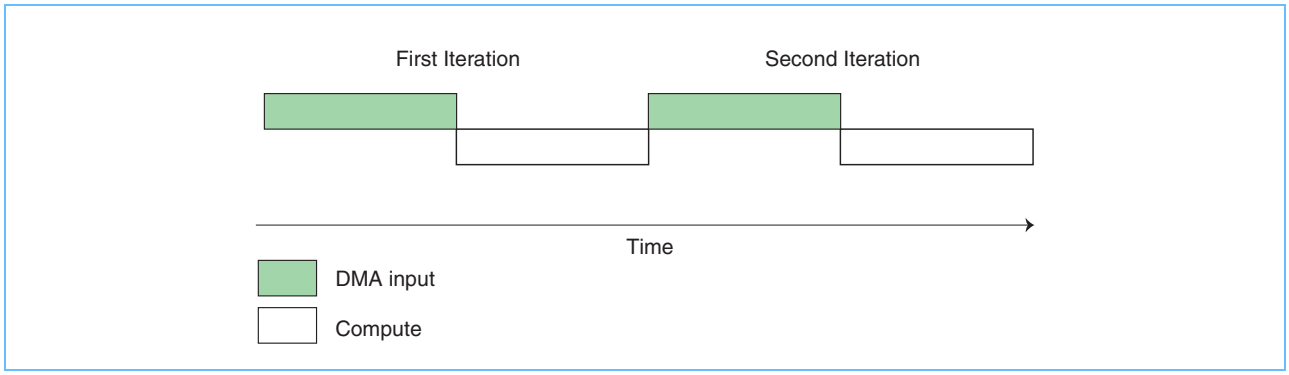
#### 24.1.2.1 *Double Buffering*

Consider a simple SPU program that performs the following steps:

1. DMA incoming data from effective address (EA) in main storage to LS buffer B.
2. Wait for the transfer to complete.
3. Compute on data in buffer B.
4. Repeat as necessary.

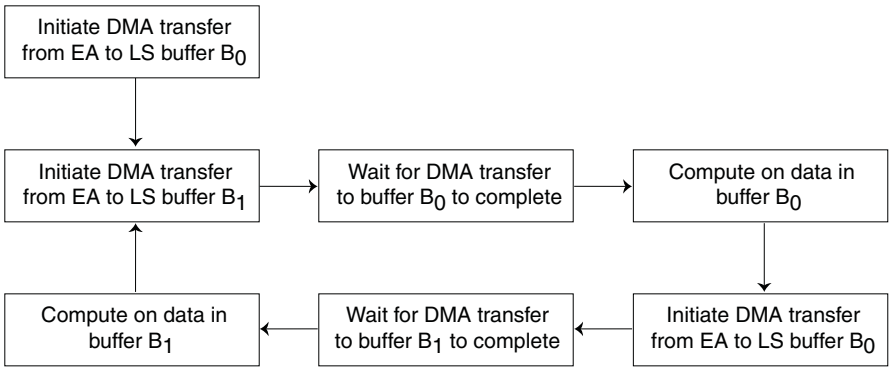
This sequence has no overlap between data transfer and computation. If considered over time, the time graph might look like *Figure 24-1*.

*Figure 24-1. Serial Computation and Transfer*




When these steps are known to iterate more than once, performance might be improved by allocating two LS buffers, $B_0$ and $B_1$, and overlapping computation on one buffer with data transfer for the next. This technique is known as *double-buffering*. *Figure 24-2* on page 693 shows a flow diagram for this double buffering scheme. Double buffering is a form of *multibuffering,* which is the method of using multiple buffers in a circular queue to overlap computation and data transfer.

*Figure 24-2. DMA Transfers Using a Double-Buffering Method*

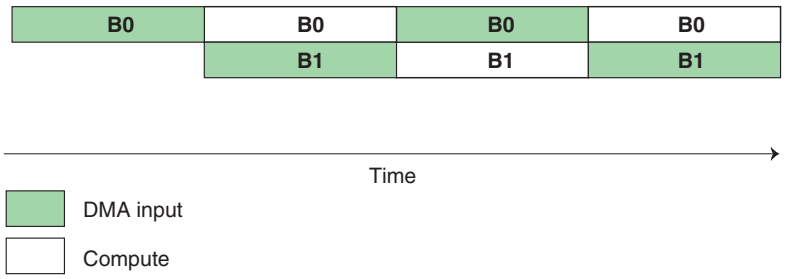



The double-buffering sequence is:

1. Initiate DMA transfer of incoming data from EA to LS buffer $B_0$.
2. Initiate DMA transfer of incoming data from EA to LS buffer $B_1$.
3. Wait for transfer of buffer $B_0$ to complete.
4. Compute on data in buffer $B_0$.
5. Initiate DMA transfer of incoming data from EA to LS buffer $B_0$.
6. Wait for transfer of buffer $B_1$ to complete.
7. Compute on data in buffer $B_1$.
8. Repeat steps 2 through 7 as necessary.

As the new time graph in *Figure 24-3* shows, computation and data transfers proceed concurrently once the first DMA for $B_0$ completes.

*Figure 24-3. Parallel Computation and Transfer*




Double-buffering can be achieved on the SPU by applying tag-group identifiers (*Section 19.2.5* on page 519). Tag-group ID 0 is applied to all transfers involving $B_0$ (steps 1, 3, and 5) and tag-group ID 1 is applied to all transfers involving $B_1$ (steps 2 and 6).

The key to double-buffering on the SPU lies in properly setting the tag-group mask, such that the program only waits for the particular buffer of interest. To wait for $B_0$, software sets the tag-group mask to include only tag ID 0 and requests conditional tag-status update (step 3). To wait for $B_1$, software sets the tag-group mask to include only tag ID 1 and again requests conditional tag-status update (step 6).

The following C-language program fragment shows how to initiate a buffer transfer.

```
#include <spu_intrinsics.h>
#include <spu_mfcio.h>

volatile void *B[2]; /* Pointers to LS Buffers */

/* Initiate transfer using LS buffer B[idx] */
static inline void xfer(unsigned int ea, unsigned int size, unsigned int id)
{
    spu_mfcdma32(B[idx], ea, size, idx, MFC_GET_CMD);
}
```

The following C-language program fragment shows how to wait for a buffer transfer to complete.

```
/* Wait for B[idx] transfer to complete. */
static inline void wait_xfer(unsigned int idx)
{
    unsigned int tag_mask = (1 << idx);

    spu_writech(MFC_WrTagMask, tag_mask);
    spu_mfcstat(MFC_TAG_UPDATE_ALL);
}
```

Both incoming or outgoing data transfers can cause bottlenecks for application performance. It might therefore be necessary to implement buffering schemes for both input and output.

The following C-language program fragment shows one way to implement double-buffering. The code starts the first DMA transfer, then enters the loop. The loop starts the next transfer and waits on the first one. When the first completes, the code calls the `use_data()` function to process the fresh buffer of DMA data. The code then toggles the buffer index and loops to start the next transfer. The code waits for the oldest transfer to complete, and then calls `use_data()` to process the fresh data. The process repeats until all the data has been transferred and processed.

Assume that the `MFC_` macro names have been defined previously.

```
/* Example C code demonstrating double buffering using buffers B[0] and B[1].
 * In this example, an array of data starting at the effective address
 * eahi|ealow is DMAed into the SPU's local storage in 4 KB chunks and processed
 * by the use_data subroutine.
 */

#include <spu_intrinsics.h>
#include <spu_mfcio.h>
```

```
#define BUFFER_SIZE 4096
volatile unsigned char B[2][BUFFER_SIZE] __attribute__ ((aligned(128)));

void double_buffer_example(unsigned int ea, int buffers)
{
    int next_idx, idx = 0;

    // Initiate first DMA transfer
    spu_mfcdma32(B[idx], ea, BUFFER_SIZE, idx, MFC_GET_CMD);

    ea += BUFFER_SIZE;
    while (--buffers) {
        next_idx = idx ^ 1;
        spu_mfcdma32(B[next_idx], ea, BUFFER_SIZE, idx, MFC_GET_CMD);
        ea += BUFFER_SIZE;
        spu_writech(MFC_WrTagMask, 1 << idx);
        (void)spu_mfcstat(MFC_TAG_UPDATE_ALL); // Wait for previous transfer done
        use_data(B[idx]);                       // Use the previous data
        idx = next_idx;
    }
    spu_writech(MFC_WrTagMask, 1 << idx);
    (void)spu_mfcstat(MFC_TAG_UPDATE_ALL);    // Wait for last transfer done
    use_data(B[idx]);                         // Use the last data
}
```

#### 24.1.2.2 *Multibuffering*

Data buffering can be extended to use more than two buffers, provided that enough space is available in the LS. If more than two buffers are involved, the technique is known as multi-buffering.

Building on the concepts from the previous section, a basic recipe for multibuffered data transfers on the SPU is:

1. Allocate multiple LS buffers, $B_0..B_n$.
2. Initiate transfers for buffers $B_0..B_n$. For each buffer $B_i$, apply tag group identifier i to transfers involving that buffer.
3. Beginning with $B_0$ and moving through each of the buffers in round robin fashion:
   - Set tag group mask to include only tag i, and request conditional tag status update.
   - Compute on $B_i$.
   - Initiate the next transfer on $B_i$.

This algorithm waits for and processes each $B_i$ in round-robin order, regardless of when the transfers complete with respect to one another. In this regard, the algorithm uses a strongly ordered transfer model. Strongly ordered transfers are useful when the data must be processed in a known order.

An alternative is to set the tag-group mask to include all outstanding buffer transfers, and then request conditional tag-status update for any tags that have completed. The buffers can be processed when their transfers complete, rather than in the order that their transfers were initiated. This represents a weakly ordered transfer model. Weakly ordered transfers might be useful when the data can be processed in any order.

#### 24.1.2.3 *Shared I/O Buffers*

This section examines a special case of multibuffering in which the storage for the incoming buffer is reused for subsequent outbound transfers. This allows the program to compute on input data from a buffer, store the results back to that same buffer, and then initiate an outbound transfer to write the results back to main storage. The advantage of using shared buffers is that LS memory requirements can be reduced.

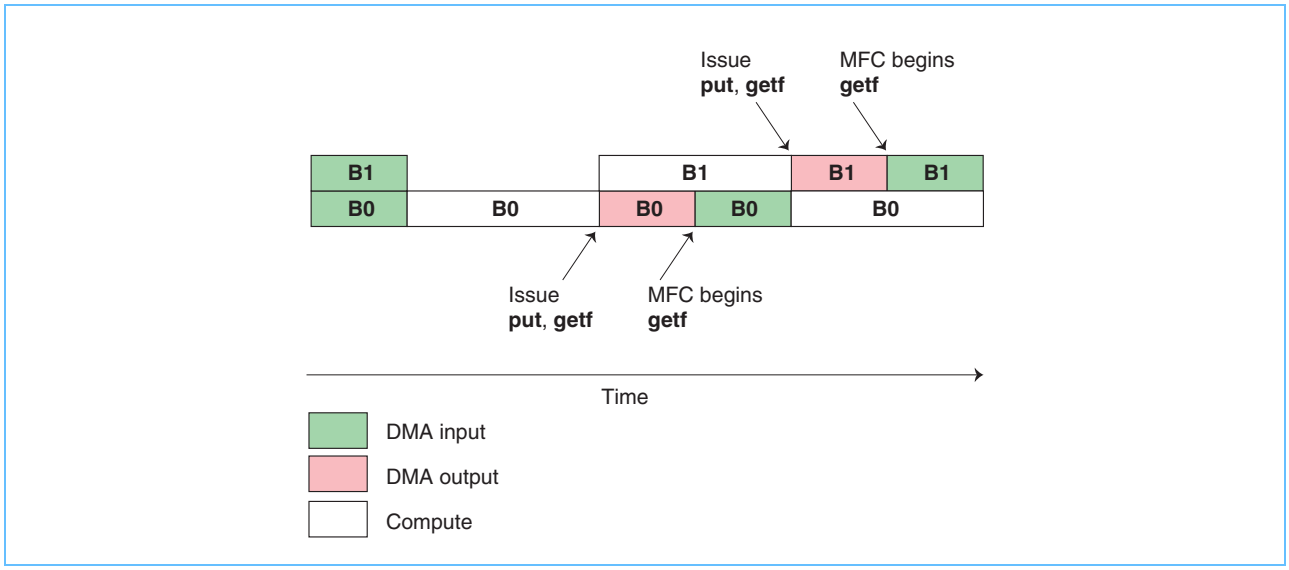
Shared I/O buffers can be particularly useful when a task has roughly a one-to-one correspondence between the amount of incoming and outgoing data. Vertex transformation (or more generally vertex shading) is an example of a task in which there is a one-to-one correspondence between incoming and outgoing data. Image convolution is another example.

However, an ordering dependency is introduced when sharing buffers for both input and output: previous outbound transfers need to complete before subsequent incoming transfers can be initiated on the same buffer. Such an ordering dependency can lead to increased program complexity, particularly when multiple incoming and outgoing transfers are in flight.

The memory flow controller (MFC) supports fenced variants of all the **get** and **put** commands (*Section 20.1.4.4* on page 574). The fence attribute causes a command to be locally ordered with respect to all previously issued commands within the same tag group. This is useful for shared buffers.

*Figure 24-4* on page 697 show how a fenced **get** command works. The ordering provided by the fence ensures that the **get** is not initiated until the **put** from the same buffer has completed. By using fenced **get** commands, an SPU program does not need to wait for outbound transfers to complete. Both parallelism and program simplicity can be maintained.

*Figure 24-4. Shared I/O Buffers with Fenced **get** Command*




#### 24.1.2.4 *Guidelines for Multibuffering Data*

To use multibuffering effectively for data, follow these rules for DMA transfers on the SPE:

- Use multiple LS buffers.
- Use unique DMA tag IDs, one for each LS buffer.
- Use fenced or barriered commands to order the DMA transfers within a tag group.
- Use separate barrier commands to order DMA transfers within the MFC command queue.
- Keep array elements a power-of-2 size to avoid multiplication when indexing.

### 24.1.3 DMA Transfers and LS Accesses

When using DMA buffers, declare the DMA buffers as volatile to ensure that buffers are not accessed by SPU load or store instructions until after DMA transfers have completed.

The *C/C++ Language Extensions for Cell Broadband Engine Architecture* document specifies that channel commands are ordered with respect to volatile-memory accesses. The DMA commands specify LS addresses as volatile, void pointers. By declaring all DMA buffers as volatile, it forces all accesses to these buffers to be performed (that is, they cannot be cached in a register) and ordered.

Depending on the algorithm and buffer access, this can result in additional memory accesses (loads and stores) and additional constraints on the compiler's optimizing scheduler that need not exist. The solution to this problem is to either:

- Load DMA-buffer contents into nonvolatile (typically local stack) variables, and reference these variables.

- Declare the DMA buffers as nonvolatile, and perform DMA transfers and synchronization in noninlined functions. This orders the DMA transfers and memory accesses by separating the two by a function-call boundary.

When coding DMA transfers, exploit DMA list transfers whenever possible. See *Section 19.2.2 DMA List Commands* on page 518 for details. See *Section 3.1.1.3* on page 68 for LS access priorities.

## 24.2 SPU Pipelines and Dual-Issue Rules

The SPU has two pipelines, even (pipeline 0) and odd (pipeline 1), into which it can issue and complete up to two instructions per cycle, one in each of the pipelines. Whether an instruction goes to the even or odd pipeline depends on its instruction type, which is related to the execution unit that performs the function. Each execution unit is assigned to one of the two pipelines.

To obtain high performance from the pipeline:

- Design for balanced pipeline use. Typically, the algorithm dictates the instruction mix. However, there might be multiple ways to achieve the same computational results. Choosing the one that achieves balanced pipeline use will often result in improved performance.
- Unroll loops and interleave computation to hide latency (reduce dependency stalls) and improve dual-issue rates.

*Table 24-1* and *Table 24-2* provide an overview of:

- *Pipeline*—The pipeline on which the instructions are issued.
- *Stalls*—The number of additional cycles before another instruction of the same type can be issued. For example, double-precision floating-point operations on the first implementation of the CBEA, the Cell/B.E. processor, have a 6-cycle stall. Therefore, for back-to-back double-precision floating-point operations, the second operation will be issued at least 7 cycles after the first operation.
- *Latency*—The number of instructions before the result is available.

*Table 24-1. Pipeline 0 Instructions and Latencies*

| Pipeline-0 Instructions | Latency (clocks) | Stall (clocks) |
|---|---|---|
| Single-precision floating-point operations | 6 | 0 |
| Integer multiplies, convert between floating-point and integer, interpolate | 7 | 0 |
| Immediate loads, logical operations, integer add and subtract, signed extend, count leading zeros, select bits, carry and borrow generate | 2 | 0 |
| Double-precision floating-point operations | | |
| Cell/B.E. processor | 7 | 6 |
| PowerXCell 8i processor | 9 | 0 |
| Element rotates and shifts | 4 | 0 |
| Byte operations (count ones, absolute difference, average, sum) | 4 | 0 |

*Table 24-2. Pipeline 1 Instructions and Latencies*

| Pipeline-1 Instructions | Latency (clocks) | Stall (clocks) |
|---|---|---|
| Shuffle bytes, quadword rotates, and shifts | 4 | 0 |
| Gather, mask, generate insertion control | 4 | 0 |
| Estimate | 4 | 0 |
| Loads | 6 | 0 |
| Branches | 4 | 0 |
| Channel operations, move to/from special purpose registers (SPRs) | 6 | 0 |

The SPU issues all instructions in program order according to the pipeline assignment. Each instruction is part of a doubleword-aligned instruction-pair called a *fetch group*. A fetch group can have one or two valid instructions. This means that the first instruction in the fetch group is from an even word address, and the second instruction from an odd word address. The SPU processes fetch groups one at a time, continuing to the next fetch group when the current fetch group becomes empty. An instruction becomes issueable when register dependencies are satisfied and there is no structural hazard (resource conflict) with prior instructions or LS contention due to DMA or error-correcting code (ECC) activity. See *Section 3.1.1.3* on page 68 for LS access priorities.

Dual-issue occurs when a fetch group has two issueable instructions in which the first instruction can be executed on the even pipeline and the second instruction can be executed on the odd pipeline. If a fetch group cannot be dual-issued, but the first instruction can be issued, the first instruction is issued to the proper execution pipeline and the second instruction is held until it can be issued. A new fetch group is loaded after both instructions of the current fetch group are issued.

For details on the SPU fetch and issue rules, see *Section B.1.3* on page 779.

## 24.3 Eliminating and Predicting Branches

The SPU hardware assumes sequential instruction flow. A branch instruction has the potential of disrupting the assumed sequential flow. Correctly predicted branches execute in one cycle, but a mispredicted branch (conditional or unconditional) incurs a penalty of 18 to 19 cycles, depending on the address of the branch target. Considering the typical SPU instruction latency of two-to-seven cycles, mispredicted branches can seriously degrade program performance. The branch instructions also restrict a compiler's ability to optimally schedule instructions by creating a barrier on instruction reordering.

The most effective method of reducing the impact of branches is to eliminate them using three primary methods—inlining, unrolling, and predication. The second-most effective method of reducing the impact of branches is to use the hint-for branch instructions. If software speculates that the instruction branches to a target path, a branch hint is provided. If a hint is not provided, software speculates that the branch is not taken (that is, instruction execution continues sequentially).

### 24.3.1 Function-Inlining and Loop-Unrolling

Function-inlining and loop-unrolling are two techniques often used to increase the size of *basic blocks* (sequences of consecutive instructions without branches), which increases scheduling opportunities.

Function-inlining eliminates the two branches associated with function-call linkage. These include the *branch and set link* (such as **brasl**) for function-call entry, and the *branch indirect* (such as **bi**) for function-call return. Loop-unrolling eliminates branches by decreasing the number of loop iterations. Loop unrolling can be manual, programmer-directed, or compiler-automated. Typically, branches associated with looping are inexpensive because they are highly predictable. However, if a loop can be fully unrolled, then all branches can be eliminated—including the final nonpredicted branch. For more about loop unrolling, see *Section 24.4* on page 709.

Over-aggressive use of inlining and loop unrolling can result in code that reduces the LS space available for data storage or, in the extreme case, is too large to fit in the LS.

### 24.3.2 Predication Using Select-Bits Instruction

The *select-bits* (**selb**) instruction is the key to eliminating branches for simple control-flow statements (for example, *if* and *if-then-else* constructs). An *if-then-else* statement can be made branchless by computing the results of both the *then* and *else* clauses and using *select bits* (**selb**) to choose the result as a function of the conditional. If computing both results costs less than a mispredicted branch, then there are additional savings.

For example, consider the following simple if-then-else statement:

```
unsigned int a, b, d;
. . .
if (a > b) d += a;
else       d += 1;
```

This code sequence when directly converted to an SPU instruction sequence without branch optimizations looks like:

```
    clgt   cc, a, b
    brz    cc, else
then:
    a      d, d, a
    br     done
else:
    ai     d, d, 1
done:
```

Using the *select bits* instruction, this simple conditional becomes:

```
clgt   cc, a, b                  /* compute the greater-than condition */
a      d_plus_a, d, a            /* add d + a */
ai     d_plus_1, d, 1            /* add d + 1 */
selb   d, d_plus_1, d_plus_a, cc /* select proper result */
```

This example shows:

- Both branches were eliminated, and the correct result was placed in `d`.
- New registers were needed to maintain potential values of `d` (`d_plus_a` and `d_plus_1`). This does not put significant pressure on the register file because the register file is so large and life of these variables is very short.
- The rewritten code sequence is smaller.
- The latency of the operations permits the scheduler to cover most of the cost of computing both conditions. Further scheduling these instructions with those before and after this code sequence will likely improve performance even further.

Here is an example of using the select bits with C intrinsics. This code fragment uses the `spu_cmpgt`, `spu_add`, and `spu_sel` intrinsics to eliminate conditional branches.

```
vector unsigned int va, vb, vd;

va = spu_promote(a, 0);
vb = spu_promote(b, 0);
vd = spu_promote(d, 0);

vd = spu_sel(spu_add(vd, 1), spu_add(vd, va), spu_cmpgt(a, b));

d = spu_extract(vd, 0);
```

***Note:*** *Most optimizing compilers will exploit the* ***selb*** *instructions for small conditionals.*

### 24.3.3 Branch Hints

General-purpose processors have typically addressed branch prediction by supporting hardware look-asides with branch history tables (BHT), branch-target address caches (BTAC), or branch-target instruction caches (BTIC). The SPU supports branch prediction through a set of *hint-for branch* (HBR) instructions (**hbr**, **hbra**, and **hbrr**) and a branch-target buffer (BTB). These instructions support efficient branch processing by allowing programs to avoid the penalty of taken branches.

Although the effects of hint-for branch instructions are not defined by the *Synergistic Processor Unit Instruction Set Architecture*, the hint-for branch instructions provide three kinds of advance knowledge about future branches:

- Address of the branch target (that is, where will the branch take the flow of control)
- Address of the actual branch instruction (known as the *hint-trigger address*)

- Prefetch schedule (when to initiate prefetching instructions at the branch target)

Hint-for branch instructions have no program-visible effects. They provide a hint to the SPU about a future branch instruction, with the intention that the information be used to improve performance by prefetching the branch target.

If software provides a branch hint, software is speculating that the instruction branches to the branch target. If a hint is not provided, software speculates that the instruction does not branch to a new location (that is, it stays inline). If speculation is incorrect, the speculated branch is flushed and refetched. It is possible to sequence multiple hints in advance of multiple branches. As with all programmer-provided hints, care must be exercised when using branch hints because, if the information provided is incorrect, performance might degrade.

#### 24.3.3.1 *Hint-for Branch Instructions*

The SPE hint-for branch instructions are shown in *Table 24-3*. There are immediate and indirect forms for this instruction class. The location of the branch is always specified by an immediate operand in the instruction.

*Table 24-3. Hint-for Branch Instructions*

| Instruction | Description |
|---|---|
| **hbr** s11, ra | Hint for branch (r-form). Hint that the instruction addressed by the sum of the address of the current instruction and the signed extended 11-bit value, s11, will branch to the address contained in word element 0 of register ra. This form is used to hint function returns, pointer function calls, and other situations that give rise to indirect branches. |
| **hbra** s11, s18 | Hint for branch (a-form). Hint that the instruction addressed by the sum of the address of the current instruction and the signed extended 11-bit value, s11, will branch to the address specified by the sign extended, 18-bit value s18. |
| **hbrr** s11, s18 | Hint for branch relative. Hint that the instruction addressed by the sum of the address of the current instruction and the signed extended 11-bit value, s11, will branch to the address specified by the sum of the address of the current instruction and sign extended, 18-bit value s18. |

#### 24.3.3.2 *Hint-for Branch Fetch*

Hint-for branch instructions load the SPU's branch-target buffer (BTB) with a branch target address and an address of a branch instruction (the hint-trigger address). After loading, the BTB monitors the instruction stream as it goes into the issue stage of the pipeline. When the address of the instruction going into issue matches the hint-trigger address, the hint is triggered, and the SPU speculates to the target address.

Branch hints arbitrate for the LS as though they were loads, except that they require three idle DMA cycles before they can issue. See *Section 3.1.1.3* on page 68 for details about LS access priority.

#### 24.3.3.3 *Hint Trigger*

The BTB is invalid at SPU program start. After a branch hint is loaded, the hint remains valid either until it is replaced by another hint or is invalidated by a **sync** instruction. Hints are disabled for the first instruction pair sent to issue after a flush.

If software's objective is to load the hint early enough so that the target is issued in the cycle after the branch, the branch hint should precede the branch by at least eleven cycles plus four instruction-pairs.

#### 24.3.3.4 *Hint Stall and Pipelined Hint Mode*

It is expected that most hints are used for end-of-loop branches or paired with an upcoming branch. It is also possible to sequence multiple hint in advance of multiple branches in a pipeline fashion.

After the BTB and the trigger are loaded, they remain available for triggering until either another hint is loaded or a synchronizing event occurs (**sync** or start). In particular, the BTB is not cleared if an instruction-sequencing error[1] occurs. Therefore, it might be appropriate to hoist hint instructions from certain simple loops into loop-initialization code. The following types of loops are probably not candidates for hint hoisting:

- Loops that execute in interrupt-enabled mode
- Loops that contain **sync** instructions
- Loops that feature a likely-taken branch in the loop body

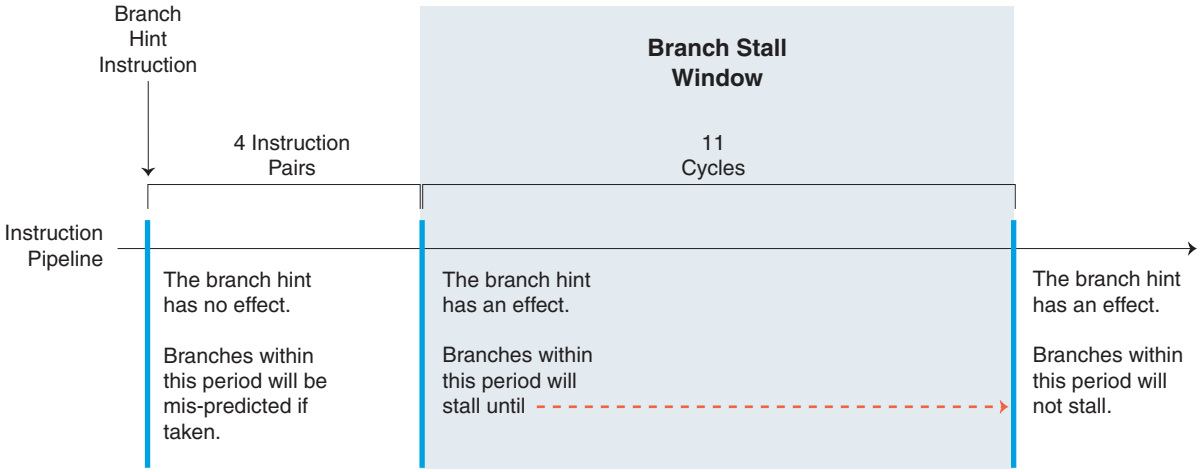
If possible, a hint should precede the branch by at least eleven cycles plus four instruction pairs. More separation between the hint and branch will probably improve the performance of applications on future SPU implementations. Hints that are too close to the branch do not affect the speculation after the branch. Hint stall has been added to reduce the number of instructions required between the hint and the branch, so that the hinted target follows the branch through issue. If there are at least four instruction pairs between the hint instruction and the branch instruction (based on the branch-address index in the hint instruction), the SPU enters hint stall. Hint stall does not stall the hint; rather, it holds the branch instruction in or before the stage of the pipeline in which triggering occurs. The branch is held there until the hint trigger is loaded. When the hint is ready for trigger, hint stall is released, proper address-matching occurs, and the target follows the branch through the issue pipeline.

Hint stall is set up by hints whose branches are at least eight, but less than 32, instructions ahead. These hints are said to be stallable. Issue stalls if the hinted branch advances to or just before the stage of the pipeline in which triggering occurs, and the BTB has not yet been loaded by the hint instruction.

*Figure 24-5* on page 704 shows the branch stall window with respect to when branch hints have an effect.

1. An instruction-sequencing error occurs when instructions are executed in a different order other than expected. The SPU expects instructions to be executed sequentially unless modified by hint trigger.

*Figure 24-5. Branch Stall Window*




If another stallable hint is issued in the window between the first stallable hint and the issuance of the original hinted branch, the SPU disables the hint stall and enters pipelined hint mode. When pipelined hint mode is entered, the hint stall is negated and no further hint stall is generated for the duration of the pipelined hint mode. Pipelined hint mode lasts for 16 cycles after the last stallable hint is issued in pipelined hint mode. Hints executed during pipelined hint mode are written to the hint buffer in order, and are available for trigger from the time they are written until the time the next hint is written. Performance enhancements obtained by using pipelined hint mode might not be portable to future implementations of SPU.

Four instruction pairs and one cycle must separate the hint from the branch in order for the branch to be predicted to the taken path. If the hint is closer than this to the branch, hint stall does not occur. To use hint stall, the instruction sequence must be padded with no operation (**nop**) and no operation-load (**lnop**) instructions. If hint stall occurs, the speculation might still be incorrect. In this case, the incorrect speculation is flushed when it gets to branch resolution.

#### 24.3.3.5 *Rules for Using Branch Hints*

The following general rules apply to the *hint for branch* (HBR) instructions:

- An HBR instruction should be placed at least 11 cycles followed by four instruction pairs before the branch instructions being hinted by the HBR instruction. In other words, an HBR instruction must be followed by at least 11 cycles of instructions, followed by eight instructions aligned on an even address boundary. More separation between the hint and branch will probably improve the performance of applications on future SPU implementations.
- If an HBR instruction is placed too close to the branch, then a hint stall will result. This results in the branch instruction stalling until the timing requirement of the HBR instruction is satisfied.
- If an HBR instruction is placed closer to the hint-trigger address than four instruction pairs plus one cycle, then the hint stall does not occur and the HBR is not used.
- Only one HBR instruction can be active at a time. Issuing another HBR cancels the current one.

- An HBR instruction can be moved outside of a loop and will be effective on each loop iteration as long as another HBR or **sync** instruction is not executed.
- The HBR instruction must be placed within -256 to +255 instructions of the branch instruction.
- The HBR instruction only affects performance.

The HBR instructions can be used to support multiple strategies of branch prediction. These include:

- *Static Branch Prediction*—Prediction based upon branch type or displacement (*Section 24.3.4 Program-Based Branch Prediction*), and prediction based upon profiling or linguistic hints (*Section 24.3.5 Profile or Linguistic Branch-Prediction* on page 706).
- *Dynamic Branch Prediction*—Software caching of branch-target addresses (*Section 24.3.6 Software Branch-Target Address Cache* on page 707), and using control flow to record branching history (*Section 24.3.7 Using Control Flow to Record Branch History* on page 708).

### 24.3.4 Program-Based Branch Prediction

Statically predicting branches based upon the program can improve branch prediction. Ball and Larus (1993) developed a set of heuristics that performed well for a large and diverse set of programs. Even though these programs are not necessarily appropriate candidates for SPU execution, the heuristic are still valid.

*Table 24-4* proposes static branch predictions for various types of program constructs. For example, “always” means that the branch should always be predicted taken, and “never” means that branch should never be predicted taken.

*Table 24-4. Heuristics for Static Branch Prediction*

| Program Feature | Static Prediction |
|---|---|
| Unconditional branch | always |
| Conditional branch | never |
| Loop-iteration branch | always |
| Loop-termination branch | never |
| Call and return | always |
| Pointer comparison | false |
| Less than zero | false |
| Greater than zero | true |
| Floating-point equal | false |
| Blocks containing store instructions[1] | avoid |
| Path length[1] | shortest |
| Short-circuit evaluation | avoid |

1. “Blocks containing store instructions” and “Path length” refer to the properties of code in the predicted branch target.

For the path-length heuristic in *Table 24-4* on page 705, if prediction is correct, branch penalty is small, but if prediction is incorrect, the mispredict branch penalty is small compared to the computation cost of the executed path.

For the short-circuit evaluation heuristic in *Table 24-4*, a common optimization performed by compilers on complex conditionals is to terminate (short-circuit) the evaluation as soon as possible. Consider the following example:

```
C Source Conditional:
    if ((a > b) && (c > d))
Assembly:
    clgt    cc_ab, a, b
    clgt    cc_cd, c, d
    and     cc, cc_ab, cc_cd
    brz     cc, not_true
        . . .
not_true:
```

Short-circuit evaluation introduces a branch on the first condition, so the second need not be evaluated:

```
Assembly w/ Short-Circuit Evaluation:
    clgt    cc_ab, a, b
    brz     cc_ab, not_true
    clgt    cc_cd, c, d
    brz     cc_cd, not_true
        . . .
not_true:
```

This technique is not well-suited for effective branch hints on the SPU because there is generally not adequate runway for hinting all the branches. As such, this optimization technique should probably be avoided.

### 24.3.5 Profile or Linguistic Branch-Prediction

A common approach to generating static branch prediction is to use expert knowledge that is obtained either by feedback-directed optimization techniques or using linguistic hints supplied by the programmer.

There are many arguments against profiling large bodies of code, but most SPU code is not like that. SPU code tends to be well-understood loops. Thus, obtaining realistic profile data should not be time-consuming. Compilers should be able to use this information to arrange code so as to increase the number of fall-through branches (that is, conditional branches not taken). The information can also be used to select candidates for loop unrolling and other optimizations that tend to unduly consume LS space.

Programmer-directed hints can also be used effectively to encourage compilers to insert optimally predicted branches. Even though there is some anecdotal evidence that programmers do not use them very often, and when they do use them, the result is wrong, this is likely not the case for SPU programmers. SPU programmers generally know a great deal about performance and will be highly motivated to generate optimal code.

The *SPU C/C++ Language Extension* specification defines a compiler directive mechanism for branch prediction (*Section B.2.5 Compiler Directives* on page 791). The `__builtin_expect` directive allows programmers to predicate conditional program statements. The following example demonstrates how a programmer can predict that a conditional statement is false (a is not larger than b).

```
if(__builtin_expect((a>b),0))
  c += a;
else
  d += 1;
```

Not only can the `__builtin_expect` directive be used for static branch prediction, it can also be used for dynamic branch prediction. The return value of `__builtin_expect` is the value of the `exp` argument, which must be an integral expression. For dynamic prediction, the `value` argument can be either a compile-time constant or a variable. The `__builtin_expect` function assumes that `exp` equals `value`.

A static-prediction example might look like this:

```
if (__builtin_expect(x, 0)) {
    foo(); /* programmer doesn't expect foo to be called */
}
```

A dynamic-prediction example might look like this:

```
cond2 = .../* predict a value for cond1 */
...
cond1 = ...
if (__builtin_expect(cond1, cond2)) {
    foo();
}
cond2 = cond1;/* predict that next branch is the same as the previous */
```

Compilers may require limiting the complexity of the expression argument because multiple branches can be generated. When this situation occurs, the compiler must issue a warning if the program's branch expectations are ignored.

### 24.3.6 Software Branch-Target Address Cache

The HBR instructions can also be used to implement a software version of a branch-target address cache (BTAC). Consider the following example in which a loop contains a conditional branch:

```
    lqr    i, iterations
loop:
     . . .
    brz    cc, skip      // condition variable cc
    . . .
skip:
    ai     i, i, -1      // decrement loop counter
    brnz   i, loop
```

The conditional branch (the **brz** instruction) can be predicted as follows:

```
    lqr    i, iterations
    ila    hint, skip    // initial hint - skip condition
loop:
    hbr    br_addr, hint // hint based upon previous iteration
     . . .
    ila    hint, skip    // assume hint for next iteration
br_addr:
    brz    cc, skip
    ila    hint, br_addr+4// correct hint for next iteration
    . . .
skip:
    ai     i, i, -1
    brnz   i, loop
```

In this example, the `hint` variable is used to store the previous iteration's branch target, under the assumption that the previous iteration's branch result is a good predictor of the next iteration's branch. It is possible to extend this concept and keep a history of branch results to be used as the predictor.

### 24.3.7 Using Control Flow to Record Branch History

Another technique for improving branch prediction is to replicate loops with its branches reversed and taken branches being vectored to the other copy of the loop. For example, consider the following if-then clause within a loop:

```
     op1
loop:
     op2
    brz    cc, label1
     op3
label1:
     op4
    brnz   cntr, loop
     op5
```

This loop can be rewritten into two tightly coupled loops as follows:

```
LOOP 1 (branch not taken)    LOOP 2 (branch taken)
-------------------------    ---------------------

    op1
loop:                        loop2:
    op2                          op2
   brz    cc, label3           brnz   cc, label2
label2:                      label3:
    op3                          op4
label1:                      brnz   cntr, loop2
    op4                      br     label4
   brnz   cntr, loop
label4:
```

This optimization is effectively a 1-bit branch predictor for the conditional. The left loop (LOOP 1) is used when predicting the conditional as true. The right loop (LOOP 2) is used when predicting the conditional as not true. The main problem with this technique is that code can grow significantly as it is applied to more complex and larger code sequences. Judicious application of this technique is recommended.

## 24.4 Loop Unrolling and Pipelining

Loops are the foundation in nearly all programs, especially streaming applications. If the number of loop iterations is a small constant, then consider removing the loop altogether. Otherwise, consider unrolling the loop if the loop is relatively independent (that is, an iteration of the loop is not dependent upon the previous iteration). The SPU has a large register file and significant loop unrolling can be accomplished before register spill occurs. Register spilling occurs when the instantaneous number of active variables exceeds the size of the register file. Unrolled loops provide additional computation for the optimizer to improve issue rates and reduce dependency stalls.

When unrolling loops, additional local variables might be required to eliminate false dependencies amongst the loop variables. Failure to eliminate these false dependencies can prevent unrolled loops from being interleaved by the compiler.

***Note:*** *Applications using auto-vectorization technology should not explicitly remove, unroll, or pipeline loops. This only adds complexity to the auto-vectorization process.*

Consider a sample workload, called `xformlight`, that performs basic graphics vertex-processing, four vertices at a time. The vertex processing includes 4 × 4 vertex transformation, perspective division, one local light computation with OpenGL-style specularity, color conversion, clamping, and RGBA color packing. *Table 24-5* on page 710 shows how loop unrolling affects the performance of this workload. These values are typical of many computational workloads.

*Table 24-5. Loop Unrolling the `xformlight` Workload*

| Metric | Unroll Factor | | | |
|---|---|---|---|---|
| | 1 | 2 | 4 | 8 |
| Normalized Performance | 1.00 | 1.52 | 1.73 | 1.66 |
| Cycles Per Instruction (CPI) | 1.35 | 0.91 | 0.76 | 0.67 |
| Dual Issue Rate | 3.3% | 19.8% | 34.3% | 35.8% |
| Dependency Stalls | 27.2% | 5.9% | 0.9% | 1.5% |
| Registers Used | 78 | 103 | 128 | 128 |
| Text Size (bytes) | 768 | 1344 | 3076 | 5252 |

Most loops have the same basic structure, as shown in *Figure 24-6*. Per iteration they load input data, perform computation, and finally store the results. Because loads, stores, quadword rotates, and shuffles execute on pipeline 1, and most computation instructions execute on pipeline 0, loops can be software pipelined to improve dual-issue rates by computing at the same time as loading and storing data. *Figure 24-7* on page 711 shows how software pipelining functions, and *Table 24-5* shows some representative results of software pipelining the unrolled `xformlight` workload.

*Figure 24-6. Basic Loop*

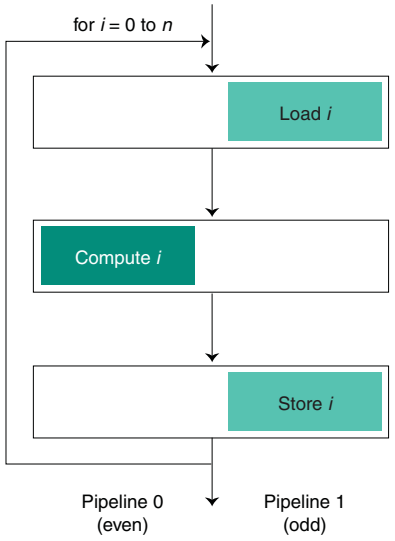


*Figure 24-7. Software-Pipelined Loop*

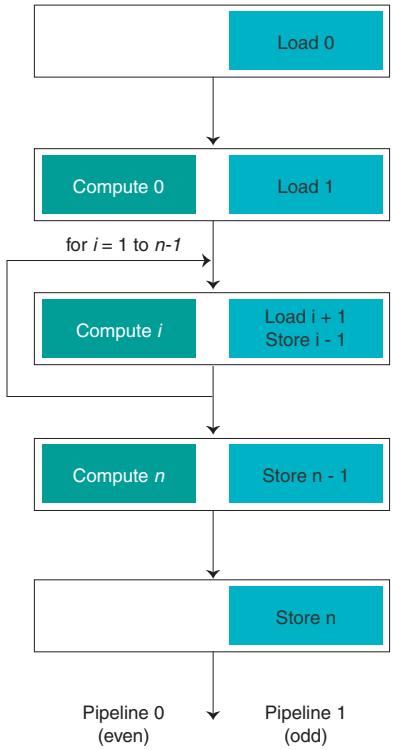



*Table 24-6. Software Pipelining the Unrolled `xformlight` Workload*

| Metric | Unroll Factor | |
|---|---|---|
| | 2 | 4 |
| Normalized Performance | 1.82 | 1.83 |
| Cycles Per Instruction (CPI) | 0.75 | 0.69 |
| Dual Issue Rate | 36.7% | 47.8% |
| Dependency Stalls | 1.3% | 0.5% |
| Resource Conflicts | 1.0% | 0.5% |
| Registers Used | 114 | 128 |
| Text Size (bytes) | 2436 | 3468 |
| Speedup versus nonpipelined | 1.16 | 1.08 |

## 24.5 Offset Pointers

The PPE's PowerPC instruction set supports *load/store with update* instructions. These instructions allow one to sequentially index through an array without the need of additional instructions to increment the array pointer. The SPU does not support this instruction form. Instead, one should exploit the SPU d-form load instructions by specifying small, literal array offsets from the base array pointer.

For example, consider the following PPE code that exploits the PowerPC *store with update* instruction:

```
#define FILL_VEC_FLOAT(_q, _data)   *(vector float)(_q++) = _data;

FILL_VEC_FLOAT(q, x);
FILL_VEC_FLOAT(q, y);
FILL_VEC_FLOAT(q, z);
FILL_VEC_FLOAT(q, w);
```

The same code can be modified for SPU execution as follows:

```
#define FILL_VEC_FLOAT(_q, _offset, _data) *(vector float)(_q+(_offset)) = _data;

FILL_VEC_FLOAT(q, 0, x);
FILL_VEC_FLOAT(q, 1, y);
FILL_VEC_FLOAT(q, 2, z);
FILL_VEC_FLOAT(q, 3, w);
q += 4;
```

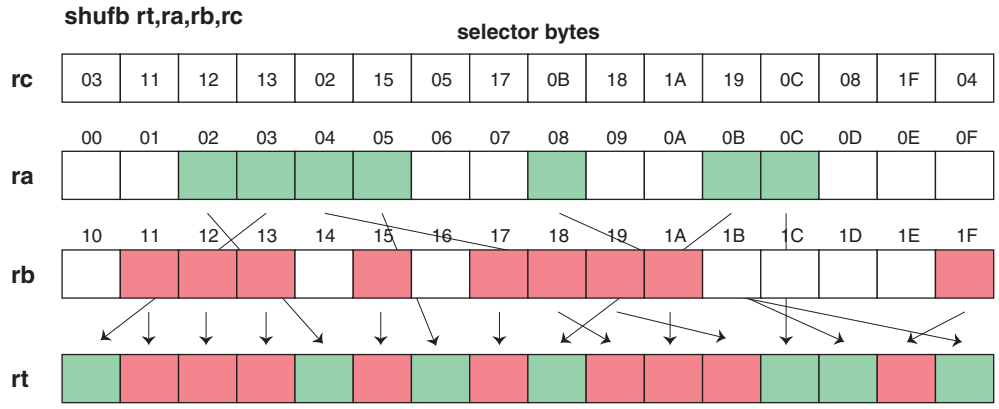
## 24.6 Transformations and Table Lookups

### 24.6.1 The Shuffle-Bytes Instruction

The SPU's **shufb** (shuffle bytes) instruction, which is comparable to the **vperm** instruction for the PPE, supports shuffles and permutations that can be used for many purposes, including table lookups, transformations between array-of-structure (AOS) and structure-of-array (SOA) vector forms, and operand alignment such as endian-reversal.

*Figure 24-8* on page 713 shows an example of how the **shufb** (shuffle bytes) instruction works. Each of the vectors, **ra**, **rb**, **rc**, and **rt**, contain sixteen byte elements. Each byte of **rc** is used to select a byte from either **ra**, **rb**, or a constant (0, x'80', or x'FF'). For example, an **rc** byte of x'03' will select byte element 3 of **ra**, and an **rc** byte of x'15' will select byte element 5 of **rb**. The results are placed in the corresponding bytes of **rt**.

*Figure 24-8. Shuffle Bytes Instruction*




A compiler can use the **shufb** instruction to perform a (sub-quadword) scalar store, using a shuffle pattern generated by one of the *Generate Controls for Insertion* instructions (**cbd**, **cbx**, **cdd**, **cdx**, **chd**, **chx**, **cwd**, **cwx**).

### 24.6.2 Fast SIMD 8-Bit Table Lookups

Many applications require table lookups. An example that uses 8-bit table lookups is the Advanced Encryption Standard (AES) byte-substitution (*ByteSub)* transformations. See the *Advanced Encryption Standard (AES)*, Federal Information Processing Standards Publications, FIPS PUB 197, 2001, and http://csrc.nist.gov/CryptoToolkit/aes/rijndael/.

Here is an example of an AES ByteSub transformation:

```
unsigned char s_box[256];
unsigned char state[16];
for (i=0; i<16; i++) {
    state[i] = s_box[state[i]];
}
```

Because table lookups involve an indirection, most vector processors do not support single instruction, multiple data (SIMD) table lookups. Instead, programmers often must fall back to scalar operations, which lack the acceleration found in parallel operations.

#### 24.6.2.1 *Scalar Table Lookup*

Assuming an AES state block is stored in a 128-bit quadword vector and the `s_box` array is either a 256-entry byte array of unknown alignment or a 16-entry quadword array, a scalar SPU implementation of the AES ByteSub transformation can be implemented as follows:

```
   ila     s_box, s_box_address    # initialize s_box array pointer
   il      cnt, 16                 # initialize loop counter
next_byte:                         # for each byte in the state quadword
```

```
andi    idx, state, 0xFF        #  strip off other bytes in the word
lqx     new_state, s_box, idx   # fetch sbox array quadword indexed by state byte
cbx     insert, s_box, idx      # generate shuffle pattern for next instruction
shufb   state, new_state, state, insert  # insert sbox lookup into state quadword
ai      cnt, cnt, -1            #  decrement loop counter
rotqbyi state, state, 1         #  rotate state quadword for next byte
brnz    cnt, next_byte          #  branch if more bytes in the quadword
```

The preceding example uses approximately seven instructions per byte-lookup. By completely unrolling the loop, the number of instructions per byte-lookup can be reduced to five. Similar results can be obtained using the PPE's vector/SIMD multimedia extension unit.

#### 24.6.2.2 *Exploiting the Shuffle-Bytes Operation*

The key to improving the performance of 8-bit table lookups on an SPE is the **shufb** instruction, described in *Section 24.6.1* on page 712. This instruction is capable of performing 16 simultaneous 8-bit table lookups in a 32-entry table. Assuming the table is quadword-aligned, the following SPU instructions perform 16 simultaneous table lookups in a 32-entry byte table:

```
lqa    table0_15, table + 0
lqa    table16_31, table + 16
andbi  idx, state, 0x1F
shufb  state, table0_15, table16_31, idx
```

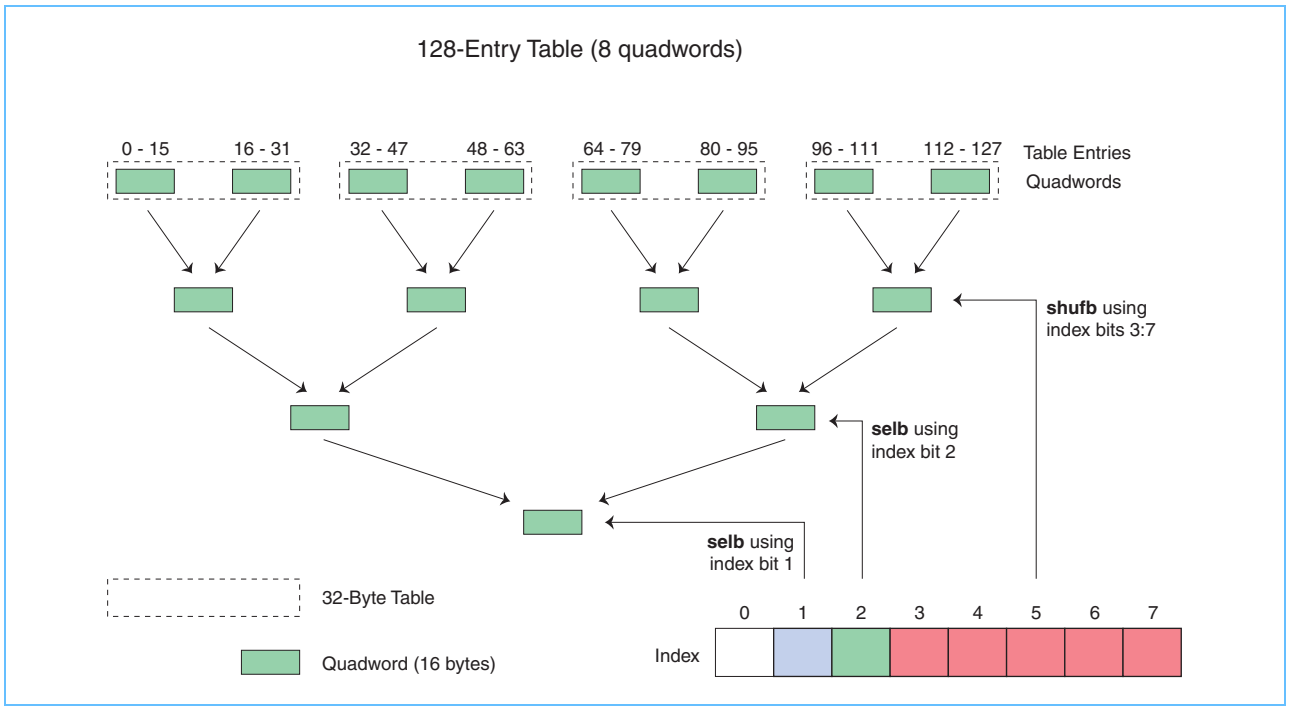
Most real-life examples require tables larger than 32 entries. Such larger tables can be handled by performing a binary-tree pruning process on a series of 32-entry table lookups. The pruning is performed using successively more significant bits in the table index.

For example, *Figure 24-9* on page 715 shows the operations required to perform a 128-entry table lookup. In this case, the 128 entries of the table are loaded, one quadword at a time, into eight quadwords (quadwords 0 through 7). The five least-significant bits are used as the shuffle pattern to perform a byte-for-byte select from the four 32-entry subtables. Each successive bit of the table index is used as a selector (using the **selb** instruction, *Section 24.3.2* on page 700) to choose the correct subtable value. This is repeated for each of the remaining most-significant bits of the table index, as follows:

```
vector unsigned char idx;
vector unsigned char tbl[8];
vector unsigned char tbl0_1, tbl2_3, tbl4_5;
vector unsigned char tbl6_7, tbl0_3, tbl4_7;
vector unsigned char lsb3_7, bit3, bit2;
vector unsigned char result;
lsb3_7 = spu_and(idx, 0x1F);
tbl0_1 = spu_shuffle(tbl[0],tbl[1],lsb3_7);
tbl2_3 = spu_shuffle(tbl[2],tbl[3],lsb3_7)
tbl4_5 = spu_shuffle(tbl[4],tbl[5],lsb3_7);
tbl6_7 = spu_shuffle(tbl[6],tbl[7],lsb3_7);
```

```
bit2 = spu_cmpeq(spu_and(idx,0x20),0x20);
tbl0_3 = spu_sel(tbl0_1,tbl2_3,bit2);
tbl4_7 = spu_sel(tbl4_5,tbl6_7,bit2);
bit1 = spu_cmpgt(idx,0x3F);
result = spu_sel(tbl0_3,tbl4_7,bit1);
```

*Figure 24-9. 128-Entry Table Lookup Example*




### 24.6.2.3 *Performance Characterization*

The previous sections have compared a scalar implementation of a table lookup with a SIMD implementation. Both implementations use quadword data-loads and loop-unrolling to maximize dual-issue rates and minimize dependencies. The performance results are shown in *Table 24-7* on page 716. The results clearly show the performance benefit of SIMD table lookups. Despite the fact that dual-issue efficiency is compromised in the SIMD version, the SIMD results show performance gains of 3.7 to 14.3 times that of scalar lookups, and code reductions of 4.6 to 27 times that of scalar lookups. The scalar implementation has higher dual-issue rates because the number of instructions and instruction mix provides more opportunities to achieve higher dual-issue rates. This example demonstrates that a small CPI does not always corollate to good performance.

*Table 24-7. Performance Comparison*

| Technique | Table Size | Code Size[1] | Cycles/Instruction | Cycles/Byte Looked Up |
|---|---|---|---|---|
| scalar | 1 to 255 | 1760 | 0.536 | 3.58 |
| SIMD | 1 to 32 | 64 | 1.07 | 0.250 |
| SIMD | 1 to 64 | 128 | 0.774 | 0.375 |
| SIMD | 1 to 128 | 224 | 0.691 | 0.594 |
| SIMD | 1 to 255 | 384 | 0.653 | 0.969 |

1. Inner-loop instructions only, expressed in bytes.

## 24.7 Integer Multiplies

The SPU contains only a 16 x16 bit multiplier. Therefore, to perform a 32-bit integer multiply, it takes five instructions—three 16-bit multiplies and two adds to accumulate the partial products.

To avoid extraneous multiply cycles, observe the following rules:

- Make array-element size a power of 2. This avoids multiplication when indexing.
- If the operands are less than 16 bits in size, cast them to `unsigned short` before multiplication. This takes maximum advantage of the multiplier.
- Always cast constants, because they have an implicit type of `int`.
- To avoid inadvertent introduction of signed extends and masks due to casting, consider introducing a macro to explicitly perform an integer multiply whose operands are 16-bits or less. For example:

```
#define MULTIPLY(a, b)\
    (spu_extract(spu_mulo((vector unsigned short)spu_promote(a,0),\
    (vector unsigned short)spu_promote(b, 0)),0))
```

## 24.8 Scalar Code

### 24.8.1 Scalar Loads and Stores

The SPU only loads and stores a quadword at a time. As such, scalar (sub-quadword) loads and stores require multiple instructions and have long latencies. This is due to the fact that scalar loads must be rotated into the preferred scalar slot (*Figure 3-6* on page 77), and stores require a read-modify-write operation to insert the scalar into the quadword being stored. This overhead is demonstrated by the following scalar load and store sample:

```
void add1(int *p) { *p += 1; }

ASSEMBLY

add1:
```

```
lqd        $4, 0(p)         # load quadword containing int pointed to by p
rotqby $5, $4, p            # rotate int into preferred scalar slot
ai         $5, $5, 1        # add 1 to int
cwd        $6, 0(p)         # generate shuffle pattern for scalar word insertion
shufb      $4, $5, $4, $6   # insert scalar into quadword
stqd       $4, 0(p)         # store updated quadword back to local storage
```

There are several strategies for making scalar code (code that is not appropriate for vectorization) more efficient. These include:

- Change the scalars to quadword vectors. This might seem wasteful, however, if you consider the three extra instructions associated with loading and storing scalars, this trade-off can actually have a positive impact on code size.
- Cluster scalars into groups and load multiple scalars at a time using a quadword memory access. Manually extract or insert the scalars on an as-needed basis. This will eliminate redundant loads and stores.

These strategies have been applied to an implementation of RC4, a character base encryption algorithm that uses a 256-byte dynamic state table. The basic algorithm is:

```
/* cipher msg_in to msg_out */
unsigned char idx1, idx2, *msg_in, *msg_out;
...
for (i=0; i<msg_len; i++)  {
    idx2 = state[++idx1] + idx2;
    SWAP(state[idx1], state[idx2]);
    msg_out[i] = msg_in[i] ^ (state[idx1] + state[idx2])
}
```

Because each iteration of the algorithm is dependent upon previous iteration, it cannot be parallelized. However, by exploiting the strategies outlined previously, significant performance enhancements can still be achieved, as shown in *Table 24-8*. The 256-byte state table is expanded into a 256-quadword state table in which all 16 elements of the unsigned character vector contain the same byte value. The input and output messages are loaded and stored a quadword at a time to eliminate the extraneous loads and stores and their respective latencies. These changes resulted in an 86% performance speedup.

*Table 24-8. RC4 Scalar Improvements*

| Implementation | Instructions | Cycles | CPI | speedup |
|---|---|---|---|---|
| original | 540120 | 723245 | 1.34 | - |
| optimized | 265794 | 388457 | 1.46 | 1.86 |

### 24.8.2 Promoting Scalar Data Types to Vector Data Types

The **spu_promote** and **spu_extract** SPU intrinsics are provided to efficiently promote scalars to vectors, or vectors to scalars. These intrinsics are listed in *Table 24-9*. When promoting a scalar operand to a vector operand, be sure to promote the scalar to the *preferred scalar slot* (*Figure 3-6* on page 77) to avoid the need for additional instructions. The preferred scalar slot for byte, halfword, word, and doubleword operations is shown in *Table 24-10*.

*Table 24-9. Intrinsics for Changing Scalar and Vector Data Types*

| Instruction | Description |
|---|---|
| d = spu_promote(a, element) | Promote a scalar to a vector. |
| d = spu_extract(a, element) | Extract a vector element from its vector. |

*Table 24-10. Preferred Scalar Slot*

| Scalar Data Type | Byte Address of Preferred Scalar Slot |
|---|---|
| Byte | 3 |
| Halfword | 1 |
| Word | 0 |
| Doubleword | 0 |

## 24.9 Unaligned Loads

On the SPE, shift-left and shift-right instructions must be used to perform an unaligned quadword vector load:

```
vector unsigned char load_vec_unaligned(unsigned char *ptr)
{
    vector unsigned char qw0, qw1, qw;
    unsigned int shift;

    shift = (unsigned int)(ptr) & 15;
    qw0 = *((vector unsigned char *)(ptr));
    qw1 = *((vector unsigned char *)(ptr+16));
    qw = spu_or(spu_slqwbyte(qw0, shift),
               spu_rlmaskqwbyte(qw1, (signed int)(shift-16)));
    return (qw);
}
```

In contrast, unaligned loads using the PPE's vector/SIMD multimedia extensions use the load-left (**lvlx**) and load-right (**lvrx**) instructions. Unaligned stores use the store-left (**stvlx**) and store-right (**stvrx**) instructions. Two examples are:

```
# Rb contains a pointer to the misaligned quadword.
# Vt contains the loaded quadword vector.
lvlx               Vhi, 0, Rb           # load left-most portion
```

```
addi                Rb, Rb, 16
lvrx                Vlo, 0, Rb            # load right-most portion
vor                 Vt, Vhi, Vlo          # combine the two portions

# Rb contains a pointer to the misaligned quadword.
# Vs contains the quadword to be stored.
stvlx                Vs, 0, Rb            # store the left-most portion
addi                 Rb, Rb, 16
stvrx                Vs, 0, Rb            # store the right-most portion
```

# Appendixes, Glossary, and Index

This section includes:

# Appendix A. PPE Instruction Set and Intrinsics

## A.1 PowerPC Instruction Set

An overview of the PowerPC Processor Element (PPE) PowerPC instructions and data types is given in *Section 2.4* on page 57. The sections that follow summarize key points of the instructions. For a complete description, see the *PowerPC Architecture* documents.

### A.1.1 Data Types

*Section 2.4.1* on page 57 contains a list of the PowerPC data types.

### A.1.2 PPE Instructions

*Table A-1* summarizes the PPE instruction performance characteristics. The entries in this table are organized by the functional units that execute the instructions. Some considerations for interpreting the table follow:

- An instruction can be executed in one or more execution units.
- Latency is the issue-to-issue latency for a dependent instruction. For example, if an **add** issues in cycle 20 and the soonest a dependent **xor** can issue is cycle 22, the latency for **add** is listed as 2. There are multiple latencies for some instructions when the instruction writes more than one type of register. In general, latency will be one cycle larger than use penalty.
- Throughput refers to the maximum sustained rate at which the PPE can execute instructions of the type noted in the absence of any dependencies and assuming infinite caches. It is shown as instructions per cycle (IPC).
- Some fields are labelled “N/A” to indicate that a more elaborate description is required (and is best understood by taking a broader view of the machine as a whole). This includes complex microcoded instructions. Instructions that do not modify a register are also labelled “N/A.”
- In the Notes column, “MC” indicates instructions that are microcoded. “CSI” indicates a context-synchronizing instruction.
- The Execution Units column indicates BRU: branch unit, IU: instruction unit, FPU: floating-point unit, FXU: fixed-point unit, and LSU: load and store unit.

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 1 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **b ba bl bla** | Branch | BRU | N/A | 1 per latency cycle | | Unconditional branches always “predicted” correctly. |
| **bc bca bcl bcla** | Branch conditional | BRU | 1 (lr / ctr) | 1 per latency cycle | | Branch direction predicted at the top of the pipeline. May update the link stack. |
| **bclr bclrl** | Branch conditional to link register | BRU | 1 (lr / ctr) | 1 per latency cycle | | Branch direction and target address predicted at the top of the pipeline. May use the link stack. |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 2 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **bcctr bcctrl** | Branch conditional to count register | BRU | 1 (lr)<br>2 (ctr) | 1 per latency cycle | | Branch direction and target address predicted at the top of the pipeline. May use the link stack. |
| **sc rfid** | System call<br>Return from interrupt | IU | N/A | N/A | CSI | |
| **crand cror crxor crnand crnor crnandc creqv crorc** | CR logical operations | BRU | 1 | 1 per latency cycle | | |
| **mcrf** | Move Condition Register (CR) field | BRU | 1 | 1 per latency cycle | | |
| **lbz lbzx lhz lhzx lwz lwzx ld ldx lhbrx lwbrx** | Load | LSU | 2 | 1 per latency cycle | | 5-cycle latency for **load** followed by dependent **load**. 2-cycle latency for **load** followed by a dependent FXU operation. |
| **lha lhax lwa lwax** | Load algebraic | LSU, FXU | See *Table A-4* on page 737 | See note 2 | MC | |
| **lbzu lhzu lwzu ldu** | Load with update | LSU, FXU | 2 (RT)<br>2 (RA) | 1 per latency cycle | | Hardware breaks into basic **load** and an **add**. |
| **lbzux lhzux lwzux ldux** | Load and zero with update indexed | LSU, FXU | 2 (RT)<br>2 (RA) | 1 per latency cycle | | Hardware breaks into a basic **load** and an **add**. |
| **lhau** | Load algebraic with update | LSU, FXU | See *Table A-4* on page 737 | See note 2 | MC | |
| **lhaux lxaux** | Load algebraic with update indexed | LSU, FXU | See *Table A-4* on page 737 | See note 2 | MC | |
| **stb sth stw std** | Store | LSU, FXU | N/A | 1 per latency cycle | | |
| **stbx sthx stwx stdx** | Store indexed | LSU, FXU | N/A | 1 per latency cycle | | |
| **stbu sthu stwu stdu** | Store with update | LSU, FXU | 2 for updated register | 1 per latency cycle | | Hardware breaks into basic **store** and an **add**. |
| **sthbrx stwbrx** | Store byte-reversed indexed | LSU, FXU | N/A | 1 per latency cycle | | |
| **stbux sthux stwux stdux** | Store with update indexed | LSU, FXU | 2 for updated register | 1 per latency cycle | | Hardware breaks into basic **store** and an **add**. |
| **lmw** | Load multiple word | LSU | See *Table A-4* on page 737 | 1 register load per latency cycle after startup | MC | Microcode generates an inline sequence of basic loads. |
| **stmw** | Store multiple word | LSU, FXU | See *Table A-4* on page 737 | 1 register store per cycle after startup | MC | Microcode generates inline sequence of basic stores. |
| **lswi** (naturally aligned) | Load string word immediate | LSU | See *Table A-4* on page 737 | 1 register store per cycle after startup | MC | Microcode assumes natural alignment and generates inline sequence of basic loads. |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 3 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **lswx** (naturally aligned) | Load string word indexed | LSU | See *Table A-4* on page 737 | 1 register store per cycle after startup | MC | Microcode assumes natural alignment and generates inline sequence of basic loads |
| **stswi** (naturally aligned) | Store string word immediate | LSU, FXU | See *Table A-4* on page 737 | 1 register store per cycle after startup | MC | Microcode assumes natural alignment and generates inline sequence of basic stores. |
| **stswx** (naturally aligned) | Store string word indexed | LSU, FXU | See *Table A-4* on page 737 | 1 register store per cycle after startup | MC | Microcode assumes natural alignment and generates inline sequence of basic stores. |
| **lswi lswx stswi stswx** (unaligned) | Load or store string word immediate<br>Load or store string word indexed | LSU | See *Table A-4* on page 737 | See note 2 | MC | A string instruction is first decoded and issued assuming natural alignment. At execution, the LSU notes that it is unaligned and causes a machine flush.<br>As the string instruction goes through microcode the second time, it is broken up in a way that takes the misalignment into account. |
| **lwarx ldarx** | Load and reserve indexed | LSU | N/A | N/A | | Forced to miss data L1 cache.<br>One outstanding **lwarx** in system at a time. |
| **stwcx. stdcx.** | Store conditional indexed | LSU, FXU | N/A | N/A | | Must establish coherency block ownership before completing the instruction (other stores do not have to do this). Can take anywhere from 10 cycles to hundreds of cycles depending upon the state of the coherency block in the memory hierarchy. |
| **addi addis add subf neg** | Add<br>Subtract<br>Negate | FXU | 2 (gpr) | 1 per latency cycle | | |
| **add. subf. subfic neg. subfe** | Add<br>Subtract<br>Negate | FXU | 2 (gpr)<br>1 (cr) | 1 per latency cycle | | |
| **addco subfco addeo subfeo addmeo subfmeo addzeo subfzeo nego addo addze subfo addc subfc addic subic adde addme subfme subfze** | Add<br>Add carrying<br>Add extended<br>Subtract<br>Subtract carrying<br>Subtract extended<br>Negate | FXU | 2 (gpr)<br>2 (xer) | 1 per latency cycle | | |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 4 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **addic. adde. addze. subfe. addme. subfme. subfze. addo. subfo. subfeo. addeo. addmeo. subfmeo. addzeo. subfzeo. addc. addco. subfc. subfco.** | Add<br>Add carrying<br>Add extended<br>Subtract<br>Subtract carrying<br>Subtract extended | FXU | 2 (gpr)<br>1 (cr)<br>2 (xer) | 1 per latency cycle | | |
| **mulli** | Multiply immediate | FXU | 6 | 1 per 6 latency cycles | | Not pipelined in the FXU. Causes a 6-cycle stall after issue. |
| **mullw mulhw mulhwu mullwo** | Multiply word | FXU | 9 | 1 per 9 latency cycles | | Not pipelined in the FXU. Causes a 9-cycle stall after issue. |
| **mulld mulhd mulhdu mulldo** | Multiply doubleword | FXU | 15 | 1 per 15 latency cycles | | Not pipelined in the FXU. Causes a 15-cycle stall after issue. |
| **mullw. mulld. mulhd. mulhw. mulhdu. mulhwu.** | Multiply recording | FXU | See *Table A-3* on page 735 | See note 2 | MC | Not pipelined in the FXU. Microcode breaks into baseline operation and a compare. |
| **mullwo. mulldo.** | Multiply recording | FXU | See *Table A-3* on page 735 | See note 2 | MC | Not pipelined in the FXU. Microcode breaks into baseline operation and a compare. |
| **divd divdu divdo divduo** | Divide | FXU (one) | 10 - 70 | 1 per 10 to 1 per 70 latency cycles | | The performance is determined by the number of bits required to represent the result.<br>PowerPC processor unit (PPU) cycles equal:<br>`((1 setup) + (ceil ((rb leading digits - ra leading digits)/2) + 1 iterations) + (1 fixup)) × 2`<br>word minimum = 10, maximum = 38 cycles<br>doubleword minimum = 10, maximum = 70 cycles<br>Overflow cases will complete in 10 cycles. |
| **divw divwu divwo divwuo** | Divide word | FXU (one) | 10 - 38 | 1 per 10 to 1 per 38 latency cycles | | Not pipelined in the FXU. The XU uses complex queue and pipeline stalls until write back.<br>See the performance notes for **divd**. |
| **divd. divw. divdu. divwu.** | Divide recording | FXU (one) | See *Table A-3* on page 735 | See note 2 | MC | Not pipelined in the FXU. Microcode breaks into a divide and a compare.<br>The XU uses complex queue and pipeline stalls until write back. |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 5 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **divdo. divwo. divduo. divwuo.** | Divide recording | FXU (one) | See *Table A-3* on page 735 | See note 2 | MC | Not pipelined in the FXU. Microcode breaks into a divide and a compare. See the performance notes for **divd**. |
| **cmpi cmp cmpli cmpl** | Compare | FXU | 1 | 1 per latency cycle | | |
| **tdi twi td tw** | Trap | FXU | See note 2 | 1 per latency cycle if a trap does not occur; otherwise, see note 2 | | Causes a program interrupt. |
| **ori oris xori xoris and or xor nand nor eqv andc orc** | Logical | FXU | 2 (gpr) | 1 per latency cycle | | |
| **andi. andis. and. or. xor. nand. nor. eqv. andc. orc. nego.** | Logical recording | FXU | See *Table A-3* on page 735 | See note 2 | MC | |
| **extsb**<br>**extsh**<br>**extsw** | Extend sign | FXU | 2 | 1 per latency cycle | | |
| **extsb.**<br>**extsh.**<br>**extsw.** | Extend sign recording | FXU | See *Table A-3* on page 735 | See note 2 | MC | |
| **cntlzd cntlzw** | Count leading zeros | FXU | 2 (gpr) | 1 per latency cycle | | |
| **cntlzd. cntlzw.** | Count leading zeros recording | FXU | See *Table A-3* on page 735 | See note 2 | MC | |
| **rldicl rldic rldicr rlwinm rldimi rlwimi** | Rotate left immediate | FXU | 2 (gpr) | 1 per latency cycle | | |
| **rldicl. rldicr. rldic. rlwinm. rldcl rldcl. rldcr rldcr. rlwnm rlwnm. rldimi. rlwimi.** | Rotate left recording<br>Rotate left clear | FXU | See *Table A-3* on page 735 | See note 2 | MC | |
| **sld sld. slw slw. srd srd. srw srw. srad srad. sraw sraw. sradi. srawi.** | Shift left<br>Shift right | FXU | See *Table A-3* on page 735 | See note 2 | MC | |
| **sradi srawi** | Shift right algebraic immediate | FXU | 2 (gpr) | 1 per latency cycle | | |
| **mtspr(xer)** | Move to XER | FXU (one) | See note 2 | See note 2 | | |
| **mtspr(lr) mtspr(ctr)** | Move to Link Register (LR)<br>Move to Count Register (CTR) | BRU | 1 | 1 per latency cycle | | |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 6 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **mtspr** (others)<br>**mtmsr mtmsrd** | Move to SPR<br>Move to MSR | Either FXU or BRU (depends on SPR) | Varies based on SPR | Varies based on SPR | | **mtmsrd** L = '0' is microcoded. **mtmsrd** L = '1' is not.<br>LR and CTR have 1-cycle latency and issue 1 per cycle. |
| **mtcrf** | Move to CR fields | FXU | See *Table A-3* on page 735 | See note 2 | MC | Bit 11 = '0' |
| **mtcrf, mtocrf** | | FXU | 1 | 1 per latency cycle | | Bit 11 = '1' |
| **mcrxr** | Move to CR from Fixed-Point Exception Register (XER) | BRU | 1 | 1 per latency cycle | | |
| **mfcr** | Move from CR | BRU | ~34 | ~1 per 34 latency cycles | | |
| **mfcrf** | | BRU | ~34 | ~1 per 34 latency cycles | | Bit 11 = '1'. This form is nonpipelined. |
| **mftb** | Move from time base | FXU (one) | ~40 | ~2 to 40 latency cycles | | |
| **mfspr(lr)**<br>**mfspr(ctr)** | | BRU | 1 | 1 per latency cycle | | |
| **mfspr** (others)<br>**mfmsr** | Move from Special Purpose Register (SPR)<br>Move from Machine State Register (MSR) | Either FXU or BRU (depends on SPR) | Varies based on SPR | Varies based on SPR | | Other SPRs are handled by the FXU and are nonpipelined. The pipeline stalls until writeback. |
| **lfs lfsx lfd lfdx** | Load floating-point | LSU, FPU Load | 1 | 1 per latency cycle | | Instructions are issued to both the FPU and the LSU units; only executes in the LSU unit |
| **lfsu lfsux lfdu lfdux** | Load floating-point with update | LSU, FXU, FPU Load | 1 (fpr)<br>2 (gpr) | 1 per latency cycle | | |
| **stfs stfsx stfd stfdx** | Store floating-point | LSU, FPU Store | See note 2 | 1 per latency cycle | | Instructions are dispatched to both the FPU and the LSU units |
| **stfsu stfsux stfdu stfdux** | Store floating-point with update | LSU, FPU, FXU | 2 (gpr) | 1 per latency cycle | | |
| **stfiwx** | Store floating-point as integer indexed | LSU, FPU | See note 2 | 1 per latency cycle | | |
| **fmr fneg fabs fnabs fadd fadds fsub fsubs fmul fmuls** | Floating-point move<br>Floating-point absolute<br>Floating-point negative<br>Floating-point negative absolute<br>Floating-point add<br>Floating-point subtract<br>Floating-point multiply | FPU | 10 | 1 per latency cycle | | |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 7 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **fmr. fneg. fabs. fnabs. fadd. fadds. fsub. fsubs. fmul. fmuls.** | Floating-point move recording<br>Floating-point absolute recording<br>Floating-point negative recording<br>Floating-point negative absolute recording<br>Floating-point add recording<br>Floating-point subtract recording<br>Floating-point multiply recording | FPU | 10 (fpr)<br>+17 (cr) | 1 per latency cycle | | The processor flushes if a CR-using operation is issued while one of these instructions is in flight. |
| **fmadd fmadds fmsub fmsubs fnmadd fnmadds fnmsub fnmsubs** | Floating-point multiply-add<br>Floating-point multiply-subtract | FPU | 10 | 1 per latency cycle | | |
| **fmadd. fmadds. fmsub. fmsubs. fnmadd. fnmadds. fnmsub. fnmsubs.** | Floating-point multiply-add recording<br>Floating-point multiply-subtract recording | FPU | 10 (fpr)<br>+17 (cr) | 1 per latency cycle | | The processor flushes if a CR-using operation is issued while one of these instructions is in flight. |
| **fdiv fdivs** | Floating-point divide (IEEE) | FPU | 74 | 1 per 74 latency cycles | | |
| **fdiv fdivs** | Floating-point divide (non-IEEE) | FPU | 56 | 1 per 56 latency cycles | | Non-IEEE results are accurate to within 1 or 2 units in last place (ULPs) of the correctly rounded IEEE result, depending on rounding mode: 1 ULP for to-nearest and to-zero rounding modes (the most common), 2 ULP for the other modes. |
| **fsqrt fsqrts** | Floating-point square root (IEEE) | FPU | 84 | 1 per 84 latency cycles | | Nonpipelined in the FPU. |
| **fsqrt fsqrts** | Floating-point square root (non-IEEE) | FPU | 66 | 1 per 66 latency cycles | | Non-IEEE results are accurate to within 1 or 2 ULPs of the correctly rounded IEEE result, depending on rounding mode: 1 ULP for to-nearest and to-zero rounding modes (the most common), 2 ULP for the other modes. |
| **fres** | Floating-point reciprocal estimate | FPU | 10 | 1 per latency cycle | | |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 8 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **frsqrte** | Floating-point reciprocal square-root estimate | FPU | 10 | 1 per latency cycle | | |
| **fdiv. fdivs. fsqrt. fsqrts. fres. frsqrte.** | Floating-point divide recording<br>Floating-point square root recording<br>Floating-point reciprocal estimate recording | FPU | same as above +17 (cr) | Same as nonrecording instructions with an additional 17 cycles for CR | | The processor flushes if a CR-using operation is issued while one of these instructions is in flight. |
| **frsp** | Floating-point round to single-precision | FPU | 10 | 1 per latency cycle | | |
| **fctid fctidz fctiw fctiwz fcfid** | Floating-point convert to integer<br>Floating-point convert from integer | FPU | 10 | 1 per latency cycle | | |
| **frsp. fctid. fctidz. fctiw. fctiwz. fcfid.** | Floating-point convert to integer recording<br>Floating-point convert from integer recording | FPU | same as above +17 (cr) | 1 per latency cycle | | The processor flushes if a CR-using operation is issued while one of these instructions is in flight. |
| **fcmpu fcmpo** | Floating-point compare | FPU | 1 | 1 per latency cycle | | |
| **fsel** | Floating-point select | FPU | 10 | 1 per latency cycle | | |
| **fsel.** | Floating-point select recording | FPU | same as above +1 (cr) | 1 per latency cycle | | |
| **mffs mffs.** | Move from Floating-Point Status and Control Register (FPSCR) | FPU | 11-28 | 1 per 28 latency cycles | | These instructions are stalled at VQ8 until all older vector scalar unit (VSU) operations are complete. |
| **mcrfs** | Move to CR from FPSCR | FPU | See *Table A-3* on page 735 | See note 2 | MC | |
| **mtfsfi mtfsfi. mtfsf mtfsf. mtfsb0 mtfsb0. mtfsb1 mtfsb1.** | Move to FPSCR | FPU | 1 | 1 per latency cycle | | |
| **sync** | Synchronize | LSU | See note 2 | See note 2 | | The IU holds at issue until all queues and pipelines have drained. After issue, the **sync** forces previous stores to finish into the cache or memory hierarchy; that is, out of the store queues. |
| **lwsync** | Lightweight **sync** | LSU | See note 2 | See note 2 | | The IU holds at issue until all queues and pipelines have drained. After issue, the **lwsync** forces previous stores to finish into the cache or memory hierarchy; that is, out of the store queues.<br>Still broadcasts a **sync** transaction onto the element interconnect bus (EIB) (but does not block) |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 9 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **ptesync** | | LSU | N/A | N/A | | The IU holds at issue until all queues and pipelines have drained. After issue, the **ptesync** forces previous stores to finish into the cache or memory hierarchy; that is, out of the store queues.<br>Still broadcasts sync transaction onto the EIB (but does not block). |
| **sc rfi rfid** | System call<br>Return from interrupt | IU | See note 2 | See note 2 | CSI | |
| **eieio** | Enforce in-order execution of I/O | LSU | See note 2 | See note 2 | | |
| **isync** | Instruction synchronize | LSU | See note 2 | See note 2 | CSI | The IU holds at issue until all queues and pipelines have drained. Issued to the IU, and the IU performs a flush (N+1) when complete. |
| **icbi** | Instruction-cache block invalidate | LSU | See note 2 | See note 2 | | After the LSU generates and translates the effective address (EA), the **icbi** looks like a snooped **icbi** to the instruction fetcher. |
| **dcbt dbtst** | Data-cache block touch | LSU | See note 2 | See note 2 | | |
| **dcbz** | Instruction-cache block zero | LSU | See note 2 | See note 2 | | Invalidates the L1 cache line on its way to the L2. Allocation and zero function occur at the L2 cache. |
| **dcbst** | Data-cache block store | LSU | See note 2 | See note 2 | | |
| **dcbf** | Data-cache block flush | LSU | See note 2 | See note 2 | | |
| **slbie** | Segment lookaside buffer (SLB) invalidate entry | LSU | See note 2 | See note 2 | | Causes class-based and thread-based invalidate in both the I-ERAT and the D-ERAT |
| **slbia** | SLB invalidate all | LSU | See note 2 | See note 2 | | Fully invalidates the SLB, the I-ERAT, and D-ERAT |
| **tlbie** | Translation lookaside buffer (TLB) invalidate entry | LSU | See note 2 | See note 2 | | Causes index-based invalidate in both the I-ERAT and the D-ERAT<br>Broadcast onto the EIB. |
| **tlbiel** | TLB invalidate entry local | LSU | See note 2 | See note 2 | | Causes index-based invalidate in both the I-ERAT and the D-ERAT<br>Is not broadcast onto the EIB. |
| **tlbsync** | TLB synchronize | LSU | See note 2 | See note 2 | | |

*Table A-1. PowerPC Instructions by Execution Unit* (Sheet 10 of 10)

| Mnemonic | Description | Execution Unit[1] | Latency (cycles)[2] | Throughput (IPC) | Notes[3] | Comments |
|---|---|---|---|---|---|---|
| **slbmte** | SLB move to entry | LSU | See *Table A-3* on page 735 | See note 2 | MC | |
| **slbmfev slbmfee** | SLB move from entry virtual segment ID (VSID)<br>SLB move from entry effective segment ID (ESID) | LSU | See *Table A-3* on page 735 | See note 2 | MC | |

1. The execution units are BRU: branch unit, IU: instruction unit, FPU: floating-point unit, FXU: fixed-point unit, LSU: load and store unit.
2. Parameters with this note require elaborate descriptions and are best understood by taking a broader view of the machine as a whole. This might include complex microcoded instructions or instructions that do not modify a register.
3. MC: Microcode, CSI: Context Synchronizing Instruction. A minimum of 11 cycles is required before the first instruction is received from the microcode ROM, so microcoded instructions should be avoided if possible. See *Section A.1.3* on page 733 for details about microcode.

*Table A-2. Storage Alignment for PowerPC Instructions*

| Operand | | | Alignment | | |
|---|---|---|---|---|---|
| Type | Size (bytes) | Byte Alignment | Within 8-Byte Block | Crosses 8-Byte Boundary | Crosses 32-Byte Boundary[1] |
| Integer load[2] | 1, 2, 4, 8 | any | optimal | optimal | poor (to microcode) |
| Integer store[2] | 1, 2, 4, 8 | any | optimal | optimal | poor (to microcode) |
| Floating-point load | 4, 8 | not word | poor (alignment interrupt) | poor (alignment interrupt) | poor (alignment interrupt) |
| Floating-point load | 4, 8 | word | optimal | optimal | poor (to microcode) |
| Floating-point store | 4, 8 | not word | poor (alignment interrupt) | poor (alignment interrupt) | poor (alignment interrupt) |
| Floating-point store | 4, 8 | word | optimal | optimal | poor (to microcode) |
| **lmw**, **stmw** | any multiple of 4 bytes (word) | any | poor (to microcode) | poor (to microcode) | poor (to microcode) |
| load string word, store string word | any | any | poor (to microcode) | poor (to microcode) | poor (to microcode) |
| Caching-inhibited load (not from microcode) | 1, 2, 4, 8 | natural alignment | optimal | optimal | N/A |
| Caching-inhibited load (not from microcode) | 1, 2, 4, 8 | not natural alignment | poor (alignment interrupt) | poor (alignment interrupt) | poor (alignment interrupt) |
| Caching-inhibited store (not from microcode) | 1, 2, 4, 8 | natural alignment | optimal | optimal | N/A |
| Caching-inhibited store (not from microcode) | 1, 2, 4, 8 | not natural alignment | poor (alignment interrupt) | poor (alignment interrupt) | poor (alignment interrupt) |
| Caching-inhibited load or store from microcode | 1, 2, 4, 8 | any | poor (alignment interrupt) | poor (alignment interrupt) | poor (alignment interrupt) |

1. Operations crossing 4 KB, 64 KB, or segment boundaries behave the same as crossing a 32-byte boundary.
2. Byte-reversed loads and stores have no special alignment characteristics.

### A.1.3 Microcoded Instructions

Instructions that are either too complex or require too many system resources to implement (such as load string instructions) are microcoded**.** This means that they are split into several simpler instructions (microwords).

***Note:*** *A minimum of 11 cycles is required before the first instruction is received from the microcode ROM, so microcoded instructions should be avoided if possible.*

Most microcoded instructions are decoded into two or three simple PowerPC instructions, and they can be avoided in most cases. The microcoded instructions are typically decomposed into an integer and a load or store operation, with a dependency between them. Although most microcoded PowerPC instructions are decoded into only a few simple instructions, it is important to keep in mind that there are typically dependencies between the internal operations of the microcode, which generate stalls at the issue stage. Replacing the microcoded instructions with PowerPC instructions not only avoids stalling but also gives more latitude in scheduling instructions to avoid stalls, as well as potentially improving multithreaded performance.

Some microcoded instructions are more complex than just a few PowerPC instructions, and some instructions are only microcoded in certain conditions, so replacement might not be possible. Instructions that are always decoded into microcode are referred to as unconditionally microcoded. Instructions that are microcoded only under specific conditions are referred to as conditionally microcoded. The only instructions that are conditionally microcoded are unaligned loads and stores that would not be microcoded if they were aligned.

A microcoded instruction is a single instruction and must execute atomically. Therefore, an asynchronous interrupt cannot be taken while a microcoded instruction is executing. This is another reason why microcoded instruction should be avoided.

### A.1.3.1 *Unconditionally Microcoded instructions*

Instructions that are difficult to implement in hardware or are infrequently executed can be split into microcoded instructions. Instructions that are always microcoded can be summarized by the following list:

- Shifts and rotates that read the shift amount from a register instead of from an immediate field
- Load or store algebraic
- Load or store strings and multiples
- Several condition register (CR) recording instructions (“Rc” = ‘1’)
- Other instructions

*Table A-3* on page 735 describes the unconditionally microcoded instructions, except load and store instructions. *Table A-4* on page 737 describes the unconditionally microcoded load and store instructions. In both tables, the Latency column indicates the latency to access the first instructions. The Number Of Microwords column indicates the number of microwords generated, not the number of cycles required to execute the microwords, because the exact latency will depend on adjacent instructions.

The largest class of microcoded instructions is “CR recording” instructions, which set a value in the Condition Register (CR). These are composed of two instructions: the base operation followed by a dependent compare instruction. Because the compare is dependent on a general-purpose register (GPR) target produced by the arithmetic logical unit (ALU), the total amount of time taken is (`microcode-ROM access + 1 for the base instruction + 1 dependency stall + 1 compare instruction`). Some microcoded instructions, such as load and store string and load and store multiple, are decomposed into many more operations.

Other things to note:

- If replacing a microcoded **or.** instruction with an **or** instruction followed by a **cmp**, be sure to avoid inadvertently changing the priority of a thread. See *Section 10 PPE Multithreading* on page 299 for details.
- The load and store string, and the load and store multiple, instructions can be replaced with a sequence of PowerPC load or store instructions, which can be an advantage if misses are expected in the data effective-to-real address translation (D-ERAT) buffer or the L1 data cache, and if interrupt latency is important.
- The load string instruction will first attempt to use load word instructions to move the data. If the access would cross a 32-byte boundary when it accesses the L1 data cache, the load will be flushed and refetched and will proceed byte-by-byte. The store string instruction behaves similarly.

*Table A-3. Unconditionally Microcoded Instructions (Except Loads and Stores)* (Sheet 1 of 3)

| Mnemonic | Class | Latency (cycles) | Number of Microwords | Comment |
|---|---|---|---|---|
| **and.** | CR recording | 11 | 2 | Record instructions are all handled the same way. The “root” instruction is issued followed by the cmpi_x instruction. |
| **andc.** | CR recording | 11 | 2 | |
| **andi.** | CR recording | 11 | 2 | The nonrecord form used in the microcode sequence is only available to microcode. |
| **andis.** | CR recording | 11 | 2 | The nonrecord form used in the microcode sequence is only available to microcode. |
| **nand.** | CR recording | 11 | 2 | |
| **nor.** | CR recording | 11 | 2 | |
| **nego.** | CR recording | 11 | 2 | |
| **or.** | CR recording | 11 | 2 | |
| **orc.** | CR recording | 11 | 2 | |
| **xor.** | CR recording | 11 | 2 | |
| **cntlzd.** | CR recording | 11 | 2 | |
| **cntlzw.** | CR recording | 11 | 2 | |
| **divd.** | CR recording | 11 | 2 | |
| **divdo.** | CR recording | 11 | 2 | |
| **divdu.** | CR recording | 11 | 2 | |
| **divduo.** | CR recording | 11 | 2 | |
| **divw.** | CR recording | 11 | 2 | |
| **divwo.** | CR recording | 11 | 2 | . |
| **divwu.** | CR recording | 11 | 2 | |
| **divwuo.** | CR recording | 11 | 2 | |
| **eqv.** | CR recording | 11 | 2 | |
| **extsb.** | CR recording | 11 | 2 | |
| **extsh.** | CR recording | 11 | 2 | |

*Table A-3. Unconditionally Microcoded Instructions (Except Loads and Stores)* (Sheet 2 of 3)

| Mnemonic | Class | Latency (cycles) | Number of Microwords | Comment |
|---|---|---|---|---|
| **extsw.** | CR recording | 11 | 2 | |
| **mulhd.** | CR recording | 11 | 2 | |
| **mulhdu.** | CR recording | 11 | 2 | |
| **mulhw.** | CR recording | 11 | 2 | |
| **mulhwu.** | CR recording | 11 | 2 | |
| **mulld.** | CR recording | 11 | 2 | |
| **mulldo.** | CR recording | 11 | 2 | |
| **mullw.** | CR recording | 11 | 2 | |
| **mullwo.** | CR recording | 11 | 2 | |
| **nego.** | CR recording | 11 | 2 | |
| **rldcl** | indirect rotate | 11 | 4 | All indirect shift and rotate instructions are handled using the same technique. First the mt_shr is issued, followed by two no-ops for delay, followed by the “root” instruction (that is, rldcl_sh). |
| **rldcl.** | CR recording | 11 | 5 | |
| **rldcr** | indirect rotate | 11 | 4 | |
| **rldcr.** | CR recording | 11 | 5 | |
| **rldic.** | CR recording | 11 | 2 | |
| **rldicl.** | CR recording | 11 | 2 | |
| **rldicr.** | CR recording | 11 | 2 | |
| **rldimi.** | CR recording | 11 | 2 | |
| **rlwimi.** | CR recording | 11 | 2 | |
| **rlwinm.** | CR recording | 11 | 2 | |
| **rlwnm** | indirect rotate | 11 | 4 | |
| **rlwnm.** | CR recording | 11 | 5 | |
| **sld** | indirect shift | 11 | 4 | |
| **sld.** | CR recording | 11 | 5 | |
| **slw** | indirect shift | 11 | 4 | |
| **slw.** | CR recording | 11 | 5 | |
| **srad** | indirect shift | 11 | 4 | |
| **srad.** | CR recording | 11 | 5 | |
| **sradi.** | CR recording | 11 | 2 | |
| **sraw** | indirect shift | 11 | 4 | |
| **sraw.** | CR recording | 11 | 5 | |
| **srawi.** | CR recording | 11 | 2 | |
| **srd** | indirect shift | 11 | 4 | |
| **srd.** | CR recording | 11 | 5 | |

*Table A-3. Unconditionally Microcoded Instructions (Except Loads and Stores)* (Sheet 3 of 3)

| Mnemonic | Class | Latency (cycles) | Number of Microwords | Comment |
|---|---|---|---|---|
| **srw** | indirect shift | 11 | 4 | |
| **srw.** | CR recording | 11 | 5 | |
| **slbmte** | miscellaneous | 11 | 2 | |
| **slbmfev** | miscellaneous | 11 | 2 | |
| **slbmfee** | miscellaneous | 11 | 2 | |
| **mtcrf** (bit 11 = '0') | miscellaneous | 11 | 9 | |
| **mcrfs** | miscellaneous | 11 | 2 | |
| **mtmsrd** L = '0' | miscellaneous | 11 | 4 | **mtmsrd** is microcoded if and only if the L field of the instruction is 0. The microcode routine is required for synchronization. The sequence is: **sync** L = '0'; **mtmsrd** L = '0'; **sync** L = '0'; **isync;** |
| **mtmsr** L = '0' | miscellaneous | 11 | 4 | **mtmsr** is microcoded if and only if the L field of the instruction is 0. The microcode routine is required for synchronization. The sequence is: **sync** L = '0'; **mtmsr** L = '0'; **sync** L = '0'; **isync;** |

*Table A-4. Unconditionally Microcoded Loads and Stores* (Sheet 1 of 2)

| Mnemonic | Class | Latency (cycles) | Number of Microwords [1, 2] | Comment |
|---|---|---|---|---|
| **lha** | load algebraic | 11 | 7 | Handled by byte. |
| **lhau** | load algebraic | 11 | 8 | Handled by byte. |
| **lhaux** | load algebraic | 11 | 8 | Handled by byte. |
| **lhax** | load algebraic | 11 | 8 | Handled by byte. |
| **lmw** | load multiple | 11 | (2 + 1 × words) | This instruction is broken down into a series of load words. |
| **lswi** | load string/optimized | 10 | By word: (1 × words + 2 × bytes) By byte: (2 × bytes) | Optimized instruction [3] |
| **lswx** | load string/optimized | By word: 11 By byte: 7 | By word: 4 + (1 × words + 2 × bytes) By byte: 4 + (2 × bytes) | Optimized instruction [3] |
| **lwa** | load algebraic | 11 | 13 | Handled by byte. |
| **lwaux** | load algebraic | 11 | 12 | Handled by byte. |
| **lwax** | load algebraic | 11 | 12 | Handled by byte. |
| **stmw** | store multiple | 11 | (2 + 1 × words) | Broken into a series of store words. |

*Table A-4. Unconditionally Microcoded Loads and Stores* (Sheet 2 of 2)

| Mnemonic | Class | Latency (cycles) | Number of Microwords [1, 2] | Comment |
|---|---|---|---|---|
| **stswi** | store string/optimized | 10 | By word: (1 × words + 2 × bytes) By byte: (2 × bytes) | Optimized instruction [3] |
| **stswx** | store string/optimized | 7 | By word: 4 + (1 × words + 2 × bytes) By byte: 4 + (2 × bytes) | Optimized instruction [3] |

1. A microcode load or store operation can access an 8-bit byte, indicated as “by byte”, or a 32-bit word, indicated as “by word”.
2. “Words” means the number of words.
3. The instruction is first broken down into a series of load-word instructions (odd bytes are handled by byte). If this does not cause an alignment exception, then the instruction is complete. If an alignment exception occurs, the first attempt is flushed. When the instruction is returned to microcode it is then handled a byte at a time. Odd bytes, if any, are defined as the remainder of string_count / 4. For store instructions, it is a series of store words.

### A.1.3.2 *Conditionally Microcoded Instructions*

The L1 data cache is physically implemented with 32-byte sectors. Thus, a conditionally microcoded load or store instruction that attempts to perform an L1 data-cache access that crosses a 32-byte boundary must be split into several instructions. When one of these misaligned loads or stores first attempts to access the L1 data cache, the misalignment is detected and the pipeline (*Section A.6* on page 762) is flushed when the instruction reaches the EX7 stage. The flushed load or store is then refetched, converted to microcode at the decode stage, and split into the appropriate loads or stores, as well as any instructions needed to merge the values together into a single register.

Doubleword integer loads that cross a 32-byte alignment boundary are first attempted as two word-sized loads or stores. If these still cross the 32-byte boundary, they are flushed and attempted again at byte granularity. The word and halfword integer loads behave similarly.

Doubleword floating-point loads and stores that are aligned to a word boundary, but not to a doubleword boundary, are handled in microcode. If these loads or stores are not word aligned, or if they cross a virtual page boundary, a PowerPC alignment interrupt is taken.

Integer loads and stores that are misaligned but do not cross a 32-byte boundary are not converted into microcode and will have the same performance characteristics as aligned loads and stores.

All of the conditionally microcoded instructions are loads and stores. The details of these are shown in *Table A-5*.

*Table A-5. Conditionally Microcoded Instructions* (Sheet 1 of 3)

| Mnemonic | Class | Latency (cycles) | Number of Microwords[1] | Comment |
|---|---|---|---|---|
| **ld** | load optimized | 11 | By word: 6 By byte: 18 | Optimized instruction[2] |

*Table A-5. Conditionally Microcoded Instructions* (Sheet 2 of 3)

| Mnemonic | Class | Latency (cycles) | Number of Microwords[1] | Comment |
|---|---|---|---|---|
| **ldbrx** | load | 11 | 19 | Optimized instruction[2] |
| **ldu** | load optimized | 11 | By word: 9<br>By byte: 21 | Optimized instruction[2] |
| **ldux** | load optimized | 11 | By word: 7<br>By byte: 19 | Optimized instruction[2] |
| **ldx** | load optimized | 11 | By word: 7<br>By byte: 19 | Optimized instruction[2] |
| **lfd** | load | 11 | 5 | Handled by word or handled by byte |
| **lfdu** | load | 11 | 6 | Handled by word or handled by byte |
| **lfdux** | load | 11 | 6 | Handled by word or handled by byte |
| **lfdx** | load | 11 | 6 | Handled by word or handled by byte |
| **lhbrx** | load | 11 | 7 | Handled by byte. |
| **lhz** | load | 11 | 6 | Handled by byte |
| **lhzu** | load | 11 | 7 | Handled by byte |
| **lhzux** | load | 11 | 7 | Handled by byte |
| **lhzx** | load | 11 | 7 | Handled by byte |
| **lwbrx** | load | 11 | 11 | Handled by byte |
| **lwz** | load | 11 | 10 | Handled by byte |
| **lwzu** | load | 11 | 11 | Handled by byte |
| **lwzux** | load | 11 | 11 | Handled by byte |
| **lwzx** | load | 11 | 11 | Handled by byte |
| **std** | store optimized | 11 | By word: 6<br>By byte: 18 | Optimized instruction[2] |
| **stdbrx** | store | 11 | 18 | New instruction |
| **stfdux** | store | 11 | 5 | Handled by word or handled by byte |
| **stdu** | store optimized | 11 | By word: 9<br>By byte: 21 | Optimized instruction[2] |
| **stdux** | store optimized | 11 | By word: 6 | Optimized instruction[2] |
| **stdx** | store optimized | 11 | By byte: 18 | Optimized instruction[2] |
| **stfd** | store | 11 | 3 | Handled by word |
| **stfdu** | store | 11 | 5 | Handled by word |
| **stfdx** | store | 11 | 5 | Handled by word |
| **sth** | store | 11 | 5 | Handled by byte |
| **sthbrx** | store | 11 | 6 | Handled by byte |

*Table A-5. Conditionally Microcoded Instructions* (Sheet 3 of 3)

| Mnemonic | Class | Latency (cycles) | Number of Microwords[1] | Comment |
|---|---|---|---|---|
| **sthu** | store | 11 | 7 | Handled by byte |
| **sthux** | store | 11 | 7 | Handled by byte |
| **sthx** | store | 11 | 6 | Handled by byte |
| **stw** | store | 11 | 4 | Handled by byte |
| **stwbrx** | store | 11 | 10 | Handled by byte |
| **stwu** | store | 11 | 11 | Handled by byte |
| **stwux** | store | 11 | 10 | Handled by byte |

1. A microcode load or store operation can access an 8-bit byte, indicated as “by byte”, or a 32-bit word, indicated as “by word”.
2. The instruction is first broken down into two load halfword instructions. If this does not cause a misalignment, then the instruction is complete. If a misalignment occurs, the first attempt is flushed. When the instruction is returned to microcode, it is then handled a byte at a time. For store instructions, it is a series of store words.

## A.2 PowerPC Extensions in the PPE

This section describes the differences between the instruction-set architecture of the PPE and version 2.0.2 of the *PowerPC Architecture*, as defined in the *PowerPC Architecture, Books I, II, and III, version 2.02* and further described in the *PowerPC Microprocessor Family: The Programming Environments for 64-Bit Microprocessors*. These differences include:

- New PowerPC instructions added in the PPE
- New meaning to existing PowerPC instructions in the PPE
- Optional PowerPC instructions implemented by the PPE
- PowerPC instructions not implemented by the PPE
- Endian support

### A.2.1 New PowerPC Instructions

The PPE implements the following new instructions, relative to version 2.02 of the *PowerPC Architecture*:

- **ldbrx**—Load Doubleword Byte Reverse Indexed X-form
- **sdbrx**—Store Doubleword Byte Reverse Indexed X-form

Details follow starting on the next page.

# Load Doubleword Byte Reverse Indexed X-form

**ldbrx rt,ra,rb**

| 0 | 1 | 1 | 1 | 1 | 1 | RT | | | | | RA | | | | | RB | | | | | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | / |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 30 | 31 |

/ ... denotes a reserved field.

```
if RA = 0 then    b ← 0
else              b ← (RA)
EA ←      b +tea (RB)
RT ←      MEM(EA+tea7, 1) || MEM(EA+tea6, 1) || MEM(EA+tea5, 1) || MEM(EA+tea4, 1) ||
          MEM(EA+tea3, 1) || MEM(EA+tea2, 1) || MEM(EA+tea1, 1) || MEM(EA, 1)
```

Let the effective address (EA) be the sum (RA|0)+$_{tea}$(RB). Bits 0:7 of the doubleword in storage addressed by EA are loaded into RT$_{56:63}$. Bits 8:15 of the word in storage addressed by EA are loaded into RT$_{48:55}$. Bits 16:23 of the word in storage addressed by EA are loaded into RT$_{40:47}$. Bits 24:31 of the word in storage addressed by EA are loaded into RT$_{32:39}$. Bits 32:39 of the word in storage addressed by EA are loaded into RT$_{24:31}$. Bits 40:47 of the word in storage addressed by EA are loaded into RT$_{16:23}$. Bits 48:55 of the word in storage addressed by EA are loaded into RT$_{8:15}$. Bits 56:63 of the word in storage addressed by EA are loaded into RT$_{0:7}$.

That is, the bytes in memory aligned location as:

```
ABCDEFGH
```

will be stored in the GPR as:

```
HGFEDCBA
```

**Special Registers Altered**:

- None

# Store Doubleword Byte Reverse Indexed X-form

**sdbrx rt,ra,rb**

| 0 | 1 | 1 | 1 | 1 | 1 | RT | | | | | RA | | | | | RB | | | | | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | / |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 30 | 31 |

/ ... denotes a reserved field.

```
if RA = 0 then    b ← 0
else              b ← (RA)
EA ← b +tea (RB)
MEM(EA, 8) ← (RS)56:63 || (RS)48:55 || (RS)40:47 || (RS)32:39 || (RS)24:31 || (RS)16:23 ||
(RS)8:15 || (RS)0:7
```

Let the effective address (EA) be the sum (RA|0)+$_{tea}$(RB). $(RS)_{56:63}$ are stored into bits 0:7 of the doubleword in storage addressed by EA. $(RS)_{48:55}$ are stored into bits 8:15 of the doubleword in storage addressed by EA. $(RS)_{40:47}$ are stored into bits 16:23 of the doubleword in storage addressed by EA. $(RS)_{32:39}$ are stored into bits 24:31 of the doubleword in storage addressed by EA. $(RS)_{24:31}$ are stored into bits 32:39 of the doubleword in storage addressed by EA. $(RS)_{16:23}$ are stored into bits 40:47 of the doubleword in storage addressed by EA. $(RS)_{8:15}$ are stored into bits 48:55 of the doubleword in storage addressed by EA. $(RS)_{0:7}$ are stored into bits 56:63 of the doubleword in storage addressed by EA.

That is, bytes in GPR as:

```
ABCDEFGH
```

will be stored in a memory aligned location as:

```
HGFEDCBA
```

**Special Registers Altered**:

- None

## A.2.2 Implementation-Dependent Interpretation of PowerPC Instructions

### A.2.2.1 *No-Op Forms of OR and ORI Instructions*

The PPE supports a subset of the program priorities defined in the *PowerPC Operating Environment Architecture, Book III*, Section 3.4.1. The PPE also adds specific forms of the no-op instruction for performance enhancements.

For details about the program priorities and the no-op forms of the **or** and **ori** instructions that support them, see *Section 10.6.2.2 nop Instructions that Change Thread Priority and Dispatch Policy* on page 320.

### A.2.2.2 *Logical Partitioning (LPAR)*

The PPE supports 32 different logical partitions (LPARs), each of which is a virtual machine managed by the hypervisor. The Logical Partitioning Control Register (LPCR) has a new bit field [52] named Mediated Interrupt (MER). The associated interrupt is described in *Section 9 PPE Interrupts* on page 239. Values in the Real Mode Limit Select (RMLS) field [34:37] have the meanings given in *Table A-6*. All reserved values are treated as 1 MB.

*Table A-6. Summary of Real Mode Limit Select (RMLS) Values*

| RMLS(0:2) | Effective Address Limit |
|---|---|
| 0000 | 256 GB |
| 0001 | 16 GB |
| 0010 | 1 GB |
| 0011 | 64 MB |
| 0100 | 256 MB |
| 0101 | 16 MB |
| 0110 | 8 MB |
| 0111 | 128 MB |
| 1000 | Reserved |
| 1001 | 4 MB |
| 1010 | 2 MB |
| 1011 | 1 MB |
| 11-- | Reserved |

### A.2.2.3 *TLB Invalidation*

Software must use the **tlbie** or **tlbiel** instructions to invalidate entries in the translation lookaside buffer (TLB). For these instructions, the PPE supports bits [22:(63 – *p*)] of the RB source register, where *p* is the page size. This support results in a selective invalidation in the TLB, based on VPN[38:(79 – *p*)] and the page size. Thus, any entry in the TLB with matching VPN[38:(79 – *p*)] and page size will be invalidated. For small pages (*p* = 12), the Effective to Real Address Trans-

lation (ERAT) is also invalidated as a result of **tlbie** or **tlbiel** based on VPN[63:67], which is equivalent to EA[47:51]. Thus, any entry in the ERAT table with matching EA[47:51] is invalidated.

The PPE adds new fields to the RB register of the **tlbiel** instruction that are not currently defined in the *PowerPC Architecture*. These include the Large Page Selector (LP) and Invalidation Selector (IS) bits.

*Table A-7* gives details of the implementation of the IS field. Bit 1 of the IS field is ignored. Bit 0 of the IS field is provided in RB[52] of the **tlbiel** instruction.

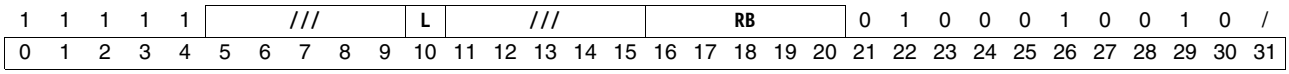
*Table A-7. Summary of Supported IS Values in **tlbiel***

| IS Field | Behavior |
|---|---|
| 00 | The TLB is as selective as possible when invalidating TLB entries. The invalidation match criteria is VPN[38:79-p], L, LP, and logical partition id (LPID). |
| 01 | Reserved. Implemented the same as IS = ‘00’. |
| 10 | Reserved. Implemented the same as IS = ‘11’. |
| 11 | The TLB does a congruence class (index-based) invalidate. All entries in the TLB matching the index of the virtual page number (VPN) supplied will be invalidated. |

The modified instruction definition for **tlbiel** specific to the PPE is shown on the next page.

# TLB Invalidate Entry Local X-form

**tlbiel rb,l**




/, //, ///, ... denotes a reserved field.

```
inval_sel ← (RB)52
if inval_sel = 0 then
       if L = 0 then pg_size ← 4 KB
       else if L = 1 then
           LP ← (RB)51
       if LP = 0 then
           pg_size = large page size 1 selected by HID6[LB]16:17
       else if LP = 1 then
           pg_size = large page size 2 selected by HID6[LB]18:19
       p ← log_base_2(pg_size)
       for each TLB entry
           if (entry_VPN38:79-p = (RB)22:63-p) & (entry_pg_size = pg_size) then
               TLB entry ← invalid
   else if inval_sel = 1 then
       if entry_pg_size = 4 KB then
           if (entry_VPN52:55 / entry_VPN60:63 || entry_VPN64:67) = (RB)44:51 then
               TLB entry ← invalid
       else if entry_pg_size = 64 KB then
           if (entry_VPN52:55 / entry_VPN56:59 || entry_VPN60:63) = (RB)44:51 then
               TLB entry ← invalid
       else if entry_pg_size = 1 MB then
           if entry_VPN52:59 = (RB)44:51 then
               TLB entry ← invalid
       else if entry_pg_size = 16 MB then
           if entry_VPN48:55 = (RB)44:51 then
               TLB entry ← invalid
```

Let the invalidation selector be $(RB)_{52}$. If the invalidation selector is 0, then let all TLB entries that have the following properties be made invalid on the processor which executes this instruction:

- The entry translates a virtual address for which $VPN_{38:79-p}$ is equal to $(RB)_{22:63-p}$.
- The page size of the entry matches the page size specified by the L and LP field of the instruction.

If the invalidation selector is 1, then let all TLB entries that are in the congruence class of the TLB corresponding to $(RB)_{44:51}$ be made invalid on the processor which executes this instruction.

Because the ERAT table does not support large pages, any **tlbie** or **tlbiel** instruction to a large page (L = 1) will cause a complete invalidation of all entries in the ERAT table. For a **tlbie** or **tlbiel** instruction to a small page (L = 0), any I-ERAT or D-ERAT entries that have VPN[63:67] match RB[47:51] of the instruction will be invalidated.

***Note:*** *It is possible for a* ***tlbiel*** *with L = '0' and IS = '1' to invalidate a TLB entry in a congruence class of the TLB but not invalidate the same translation in the I-ERAT or D-ERAT table. It is software's responsibility in this case to make sure that ERAT coherency is maintained. Software can set L = 1 to avoid this (with the performance consequence of invalidating the entire ERAT table).*

The **tlbie** instruction register transfer language (RTL) is similar except that IS = '0' is assumed, and the instruction is performed on all processors belonging to the same partition as the processor issuing the **tlbie** instruction.

Because a **tlbiel** instruction to a large page causes the ERAT to be flushed, software may issue a **tlbiel** to a large page that is not currently in use (unmapped in the page table) to cause an ERAT flush. This is useful for invalidating ERAT entries during a partition context switch. The IS field of **tlbiel** should be set to '00' for this operation.

For more information about the TLB, see *Section 4 Virtual Storage Environment* on page 79.

### A.2.3 Optional PowerPC Instructions Implemented

The following optional PowerPC user-mode instructions are implemented in the PPE:

#### A.2.3.1 *Book I Optional Instructions Implemented*

- **fsqrt(.)**—Floating-Point Square Root
- **fsqrts(.)**—Floating-Point Square Root Single
- **fres(.)**—Floating Reciprocal Estimate Single A-form
- **frsqrte(.)**—Floating-Point Reciprocal Square Root Estimate A-form
- **fsel(.)**—Floating-Point Select
- **mtocrf**—Move To One Condition Register Field XFX-form
- **mfocrf**—Move From One Condition Register Field XFX-form

#### A.2.3.2 *Book III Optional Instructions Implemented*

- **tlbie**—TLB Invalidate Entry (large and small page)
- **tlbiel**—Processor local form of TLB Invalidate Entry (large and small page)
- **tlbsync**—TLB Synchronize
- **mtmsr**—Move to Machine State Register (32-bit)

### A.2.4 PowerPC Instructions Not Implemented

The following PowerPC instructions are not implemented in the PPE:

#### A.2.4.1 *Book I Unimplemented Instructions*

The following instructions are not implemented in the PPE:

- **popcntb**—Population Count Bytes
- **fre(.)**—Floating Reciprocal Estimate A-form
- **frsqrtes(.)**—Floating Reciprocal Square Root Estimate Single A-form

The following instruction is obsolete and thus not implemented:

- **mcrxr**—Move from Condition Register to XER Register

An attempt to execute one of these instructions will result in an illegal instruction-type exception.

#### A.2.4.2 *Book II Unimplemented Instructions*

The optional external control facility is not implemented in the PPE. Instructions associated with this facility are likewise not implemented. These include:

- **eciwx**—External Control In Word indeX
- **ecowx**—External Control Out Word indeX

An attempt to execute one of these instructions will result in an illegal instruction-type exception.

#### A.2.4.3 *Book III Unimplemented Instructions*

The External Access Register (EAR) is not implemented in the PPE. Instructions associated with the EAR are likewise not implemented. These include:

- **eciwx**—External Control In Word indeX
- **ecowx**—External Control Out Word indeX
- **tlbia**—TLB Invalidate All
- **mtsr**—Move to Segment Register
- **mtsrin**—Move to Segment Register Indirect
- **mfsr**—Move from Segment Register
- **mfsrin**—Move from Segment Register Indirect

An attempt to execute one of these instructions will result in an illegal instruction-type exception.

### A.2.5 Endian Support

The PPE does not support little-endian mode; it only supports big-endian mode. Writing to the Little-Endian mode (LE) and Interrupt Little-Endian mode (ILE) bits in the Machine State Register has no effect.

## A.3 Vector/SIMD Multimedia Extension Instructions

An overview of the vector/SIMD multimedia extension instructions is given in *Section 2.5* on page 59. The sections that follow summarize key points of the instructions. For a complete description, see the *PowerPC Microprocessor Family: Vector/SIMD Multimedia Extension Technolgy Programming Environments Manual*.

### A.3.1 Data Types

*Section 2.5.2* on page 61 contains a list of the vector/SIMD multimedia extension data types.

### A.3.2 Vector/SIMD Multimedia Extension Instructions

*Table A-8* summarizes the vector/SIMD multimedia extension instructions, showing their pipeline, latency, and throughput. Some comments on the table:

- An instruction can be executed in one or more execution units, as shown in *Table A-8*.
- Latency is the issue-to-issue latency for a dependent instruction. There are multiple latencies for some instructions when the instruction writes more than one type of register. In general, latency will be one cycle larger than use penalty.
- Throughput refers to the maximum sustained rate at which the PPE can execute instructions of the type noted in the absence of any dependencies and assuming infinite caches. It is shown as instructions per cycle (IPC).
- Some fields are labelled “N/A” (not applicable) to indicate that a more elaborate description is required. This includes complex microcoded instructions and instructions that do not modify a register.
- In the Execution Units column, LSU is the load and store unit, and VXU is the vector/SIMD multimedia extension unit.

*Table A-8. Vector/SIMD Multimedia Extension Instructions* (Sheet 1 of 4)

| Instruction | Description | Execution Unit[1] | Latency | Throughput (IPC) | Comments |
|---|---|---|---|---|---|
| **lvebx lvehx lvewx lvlx lvlxl lvrx lvrxl lvsl lvsr lvxl** | Load vector indexed | VXU load, LSU | 2 | 1 | |
| **stvebx stvehx stvewx stvlx stvlxl stvrx stvrxl stvx stvxl** | Store vector indexed | VXU store, LSU | N/A | 1 | |
| **dst dstt dstst dss dssall** | Nop | | | | Not implemented (these instructions are no longer part of the architecture. They will be treated as no-ops). |
| **vpkpx vpkshss vpkshus vpkswss vpkuwus vpkuhum vpkuhus vpkuwum vpkuwus** | Vector pack | permute | 4 | 1 | |
| **vupkhpx vupkhsb vupkhsh vupklpx vupklsb vupklsh** | Vector unpack | permute | 4 | 1 | |

*Table A-8. Vector/SIMD Multimedia Extension Instructions* (Sheet 2 of 4)

| Instruction | Description | Execution Unit[1] | Latency | Throughput (IPC) | Comments |
|---|---|---|---|---|---|
| **vmrghb vmrghh vmrghw vmrglb vmrglh vmrglw** | Vector merge | permute | 4 | 1 | |
| **vspltb vsplth vspltw vspltisb vspltish vspltisw** | Vector splat | permute | 4 | 1 | |
| **vperm** | Vector permute | permute | 4 | 1 | |
| **vsldoi** | Vector shift left double by octet immediate | permute | 4 | 1 | |
| **vslo vsro** | Vector shift left by octet | permute | 4 | 1 | |
| **vaddubm vadduhm vadduwm vaddubs vadduhs vadduws vaddcuw vsububm vsubuhm vsubuwm vsububs vsubuhs vsubuws vsubsbs vsubshs vsubsws vsubcuw vsubsws vaddsbs vaddshs vaddsws** | Vector add<br>Vector subtract | simple | 4 | 1 | |
| **vavgub vavguh vavguw vagvsb vagvsh vavgsw** | Vector average | simple | 4 | 1 | |
| **vand vor vxor vandc vnor** | Vector logical | simple | 4 | 1 | |
| **vsel** | Vector select | simple | 4 | 1 | |
| **vrlb vrlh vrlw vslb vslh vslw vsl vsrb vsrh vsrw vsr vsrab vsrah vsraw** | Vector rotate<br>Vector shift | simple | 4 | 1 | |
| **vcmpgtub vcmpgtsb vcmpgtuh vcmpgtsh vcmpgtuw vcmpgtsw vcmpgtfp** | Vector compare | simple | 4 | 1 | |
| **vcmpgtub. vcmpgtsb. vcmpgtuh. vcmpgtsh. vcmpgtuw. vcmpgtsw. vcmpgtfp.** | Vector compare recording | simple | 4 | 1 | |
| **vcmpequb vcmpequh vcmpequw vcmpeqfp** | Vector compare equal to | simple | 4 | 1 | |
| **vcmpequb. vcmpequh. vcmpequw. vcmpeqfp.** | Vector compare equal to recording | simple | 4 | 1 | |

*Table A-8. Vector/SIMD Multimedia Extension Instructions* (Sheet 3 of 4)

| Instruction | Description | Execution Unit[1] | Latency | Throughput (IPC) | Comments |
|---|---|---|---|---|---|
| **vcmpbfp**<br>**vcmpgefp** | Vector compare bounds<br>Vector greater than or equal to | simple | 4 | 1 | |
| **vcmpbfp.**<br>**vcmpgefp.** | Vector compare bounds recording<br>Vector greater than or equal to recording | simple | 4 | 1 | |
| **vmaxub vmaxuh vmaxuw vmaxsb vmaxsh vmaxsw vmaxfp** | Vector maximum | simple | 4 | 1 | |
| **vminub vminuh vminuw vminsb vminsh vminsw vminfp** | Vector minimum | simple | 4 | 1 | |
| **mtvscr** | Move to the Vector Status and Control Register (VSCR) | simple | N/A | N/A | Stalls during issue until all older instructions are complete. Then, after issuing from the vector scalar unit issue queue, all younger vector scalar unit instructions stall until complete. |
| **mfvscr** | Move from VSCR | simple | N/A | N/A | Stalls during issue until all older instructions are complete. |
| **vaddfp vsubfp vmaddfp vnmsubfp** | Vector add single-precision<br>Vector subtract single-precision<br>Vector multiply-add single-precision<br>Vector negative multiply-subtract single-precision | float | 12 | 1 | |
| **vrefp vrsqrtefp** | Vector reciprocal estimate single-precision<br>Vector reciprocal square root estimate single-precision | estimate | 14 | 1 | |
| **vlogefp vexptefp** | Vector log base-2 estimate floating-point<br>Vector 2 raised to the exponent estimate floating-point | float | 12 | 1 | |

*Table A-8. Vector/SIMD Multimedia Extension Instructions* (Sheet 4 of 4)

| Instruction | Description | Execution Unit[1] | Latency | Throughput (IPC) | Comments |
|---|---|---|---|---|---|
| **vrfin vrfiz vrfip vrfim**<br>**vcfpsxws**<br>**vcfpuxws**<br>**vcuxwfp vcsxwfp** | Vector round<br>Vector convert | float | 12 | 1 | |
| **vmuloub vmulouh**<br>**vmulosb vmulosh**<br>**vmuleub vmuleuh**<br>**vmulesb vmulesh**<br>**vmhaddshs**<br>**vmhraddshs**<br>**vmladduhm**<br>**vmsumubm**<br>**vmsummbm**<br>**vmsumuhm**<br>**vmsumuhs**<br>**vmsumshm**<br>**vmsumshs**<br>**vsum4ubs**<br>**vsum4sbs**<br>**vsum4shs**<br>**vsum2sws**<br>**vsumsws** | Vector multiply<br>Vector multiply-add<br>Vector multiply-sum | complex | 9 | 1 | |

1. The vector/SIMD multimedia extension and FPU instructions are categorized into two groups. The first group, VSU type 1, includes vector/SIMD multimedia extension simple, vector/SIMD multimedia extension float, and FPU arithmetic instructions. The second group, VSU type 2, includes vector/SIMD multimedia extension load, vector/SIMD multimedia extension store, vector/SIMD multimedia extension permute, FPU load, and FPU store instructions.

*Table A-9. Storage Alignment for Vector/SIMD Multimedia Extension Instructions*

| Operand | | | Alignment | | |
|---|---|---|---|---|---|
| Type | Size (bytes) | Byte Alignment | Within 8-Byte Block | Crosses 8-Byte Boundary | Crosses 32-Byte Boundary[1] |
| VXU load (except **lvlx[l]**, **lvrx[l]**) | 16[3] | any | optimal | optimal | N/A[3] |
| VXU store (except **stvlx[l]**, **stvrx[l]**) | 16[3] | any | optimal | optimal | N/A[3] |
| VXU load or store left (**lvlx[l]**, **stvlx[l]**) | any | any | optimal | optimal | N/A[4] |
| VXU load or store right (**lvrx[l]**, **stvrx[l]**) | any | not quadword aligned | optimal | optimal | N/A[4] |
| VXU load or store right (**lvrx[l]**, **stvrx[l]**) | any | quadword aligned[2] | optimal | optimal | N/A[4] |

1. Operations crossing 4 KB, 64 KB, or segment boundaries behave the same as crossing a 32-byte boundary.
2. For Load Vector Right and Store Vector Right that are quadword aligned, no attempt is made to access storage. On a load, the data returned is zero; on a store, the operation becomes a no-op.
3. Regular vector/SIMD multimedia extension unit (VXU) operations are assumed to be aligned to the 16-byte boundary.
4. Due to the way these operations are defined, they do not cross a 32-byte boundary (data is accessed up to a16-byte boundary).

### A.3.3 Graphics Rounding Mode

The vector/SIMD multimedia extension instructions have a *graphics rounding mode* (enabled by default) that allows programs written with vector/SIMD multimedia extension instructions to produce floating-point results that are equivalent in precision to those written in the synergistic processor unit (SPU) instruction set. In this vector/SIMD multimedia extension mode, as in the SPU environment, the default rounding mode is *round to zero*, denormals are treated as zero, and there are no infinities or NaNs.

When the `HID1[grap_mode]` bit is set to 1, the VXU operates in graphics rounding mode.

The following set of rules apply to graphics rounding mode:

- The VXU operates in non-Java mode.
  - Denormal inputs are flushed to zero.
  - If the result is an underflow, the result will be zero.
  - If the infinitely precise unrounded result would be denormal, the result will be zero.
- Except where noted, the rounding mode is round toward zero.
- Infinity and NaN inputs are operated upon as normal (extended-range) numbers.
- The positive overflow boundary is moved from x‘7F80_0000’ to x‘7FFF_FFFF’.
- The negative overflow boundary is moved from x‘FF80_0000’ to x‘FFFF_FFFF’.
- If the result is greater than the new positive_overflow_boundary, then the result will be x‘7FFF_FFFF.’

- If the result is less than the new negative_overflow_boundary, then the result will be x‘FFFF_FFFF’.

The following subsections detail graphics rounding mode behavior for individual instruction classes.

#### A.3.3.1 *Math Instructions (Multiply-Add, Multiply-Subtract, Add, Subtract)*

Because there are no infinity or NaN inputs, there is not a NaN result for `infinity - infinity` or for infinity × 0. The result of infinity or NaN multiplied by zero is zero.

#### A.3.3.2 *Floating-Point Compares*

NaN and Infinity inputs are treated as large contiguous numbers, and they are correctly compared. Comparisons of +0 to -0 are the same as those for non-graphics-rounding mode.

#### A.3.3.3 *Reciprocal Estimate*

The result of a reciprocal estimate is the same as in non-Java mode, except that NaN and Infinity inputs are interpreted as large numbers and translated to correct mathematical estimates.

#### A.3.3.4 *Reciprocal Square-Root Estimate*

The instruction takes the absolute value of the input, then returns a result that is equivalent to the non-Java mode reciprocal square root estimate of the positive value. NaN and Infinity inputs are interpreted as large numbers and translated to the correct mathematical estimate.

#### A.3.3.5 *Logarithm Base-2 Estimate*

The same results as non-Java mode are generated except that NaN and Infinity inputs are interpreted as large numbers and translated to correct mathematical estimates. In non-Java mode, the log of a negative number is a NaN result. For negative inputs in graphics rounding mode, the result is the log of the absolute value of the input, except when the input is +0 or -0. If the input is +0 or -0, the output is x‘FFFF_FFFF’.

#### A.3.3.6 *Power-of-2 Estimate*

The same results as non-Java mode are generated except that NaN and Infinity inputs are interpreted as large numbers and translated to correct mathematical estimates. The detection for an overflow boundary is the same as for non-Java mode. However, the result is an overflow boundary for graphics rounding mode (provided in *Section A.3.3 Graphics Rounding Mode* on page 752).

#### A.3.3.7 *Scaled Conversion to integer*

Because the non-Java rounding mode is truncated for these instructions, the result is the same as non-Java mode except for NaN and Infinity inputs. NaN and Infinity inputs are too large for the integer range, so they cause the same integer saturation results as non-Java mode for Infinity inputs. Positive NaN inputs give the same result as positive infinity in non-Java mode. Negative NaN inputs give the same result as negative infinity in non-Java mode.

#### A.3.3.8 *Scaled Conversion from integer*

The results are the same as non-Java mode. The rounding mode is round-to-nearest-even. Because the input is an integer, there are no differences for NaN and Infinity inputs.

#### A.3.3.9 *Round to Floating-Point Integer*

Because the rounding mode is part of the instruction, the rounding mode is not suppressed as in the default graphic mode rules. The result is the same as for non-Java mode. There is a form of the instruction that uses the rounding mode truncate.

## A.4 C/C++ Language Extensions (Intrinsics) for Vector/SIMD Multimedia Extensions

An overview of the C/C++ intrinsics for the vector/SIMD multimedia extension instructions is given in *Section 2.6* on page 62. The sections that follow summarize key points of the intrinsics.

### A.4.1 Vector Data Types

The vector/SIMD multimedia extension programming model adds a set of fundamental data types, called *vector types*, as shown in *Table A-10*.

***Note:*** *Because the token, vector, is a keyword in the vector/SIMD multimedia extension data types, it is recommended that the term not be used elsewhere in the program as, for example, a variable name.*

*Table A-10. Vector/SIMD Multimedia Extension Data Types* (Sheet 1 of 2)

| Vector Data Type | Meaning | Values[1] |
|---|---|---|
| vector unsigned char | Sixteen 8-bit unsigned values | 0 ... 255 |
| vector signed char | Sixteen 8-bit signed values | -128 ... 127 |
| vector bool char | Sixteen 8-bit unsigned boolean | 0 (false), 255 (true) |
| vector unsigned short | Eight 16-bit unsigned values | 0 ... 65535 |
| vector unsigned short int | Eight 16-bit unsigned values | 0 ... 65535 |
| vector signed short | Eight 16-bit signed values | -32768 ... 32767 |
| vector signed short int | Eight 16-bit signed values | -32768 ... 32767 |
| vector bool short | Eight 16-bit unsigned boolean | 0 (false), 65535 (true) |
| vector bool short int | Eight 16-bit unsigned boolean | 0 (false), 65535 (true) |

1. The represented values are in decimal (base-10) notation.

*Table A-10. Vector/SIMD Multimedia Extension Data Types* (Sheet 2 of 2)

| Vector Data Type | Meaning | Values[1] |
|---|---|---|
| vector unsigned int | Four 32-bit unsigned values | 0 ... $2^{32}$ - 1 |
| vector signed int | Four 32-bit signed values | $-2^{31}$ ... $2^{31}$ - 1 |
| vector bool int | Four 32-bit unsigned values | 0 (false), $2^{31}$ - 1 (true) |
| vector float | Four 32-bit single precision | IEEE-754 values |
| vector pixel | Eight 16-bit unsigned values | 1/5/5/5 pixel |

1. The represented values are in decimal (base-10) notation.

Introducing fundamental vector data types permits a compiler to provide stronger type-checking and supports overloaded operations on vector types.

### A.4.2 Vector Literals

There are two standard formats of constructing a vector literal. A compiler may support one or both formats. The first format, as specified by the *AltiVec Technology Programming Interface Manual*, is shown in *Table A-11*. This format consists of a parenthesized vector type followed by a parenthesized set of constant expressions. The AltiVec format is deprecated and will often only be supported for existing software.

*Table A-11. AltiVec Vector-Literal Format* (Sheet 1 of 2)

| Notation | Description |
|---|---|
| (vector unsigned char)(unsigned int) | A set of 16 unsigned 8-bit quantities that all have the value specified by the integer. |
| (vector unsigned char)(unsigned int, ..., unsigned int) | A set of 16 unsigned 8-bit quantities specified by the 16 integers. |
| (vector signed char)(signed int) | A set of 16 signed 8-bit quantities that all have the value specified by the integer. |
| (vector signed char)(signed int, ..., signed int) | A set of 16 signed 8-bit quantities specified by the 16 integers. |
| (vector unsigned short)(unsigned int) | A set of 8 unsigned 16-bit quantities that all have the value specified by the integer. |
| (vector unsigned short)(unsigned int, ..., unsigned int) | A set of 8 unsigned 16-bit quantities specified by the 16 integers. |
| (vector signed short)(signed int) | A set of 8 signed 16-bit quantities that all have the value specified by the integer. |
| (vector signed short)(signed int, ..., signed int) | A set of 8 signed 16-bit quantities specified by the 16 integers. |
| (vector unsigned int)(unsigned int) | A set of 4 unsigned 32-bit quantities that all have the value specified by the integer. |
| (vector unsigned int)(unsigned int, ..., unsigned int) | A set of 4 unsigned 32-bit quantities specified by the 16 integers. |
| (vector signed int)(signed int) | A set of 4 signed 32-bit quantities that all have the value specified by the integer. |

*Table A-11. AltiVec Vector-Literal Format* (Sheet 2 of 2)

| Notation | Description |
|---|---|
| (vector signed int)(signed int, ..., signed int) | A set of 4 signed 32-bit quantities specified by the 16 integers. |
| (vector float)(float) | A set of 4 32-bit floating-point quantities that all have the value specified by the float. |
| (vector float)(float, float, float, float) | A set of 4 32-bit floating-point quantities that all have the value specified by the 4 floats. |

The second format for constructing a vector literal is written as a parenthesized vector type followed by a curly-braced set of constant expressions, as shown in *Table A-12*. The elements of the vector are initialized to the corresponding expression. Elements for which no expressions are specified default to zero. This format is the preferred syntax and should be used for all new application software.

*Table A-12. Curly-Brace Vector-Literal Format*

| Notation | Description |
|---|---|
| (vector unsigned char)*{unsigned int}* | A set of 16 unsigned 8-bit quantities that all have the value specified by the integer. |
| (vector unsigned char)*{unsigned int, ..., unsigned int}* | A set of 16 unsigned 8-bit quantities specified by the 16 integers. |
| (vector signed char)*{signed int}* | A set of 16 signed 8-bit quantities that all have the value specified by the integer. |
| (vector signed char)*{signed int, ..., signed int}* | A set of 16 signed 8-bit quantities specified by the 16 integers. |
| (vector unsigned short)*{unsigned int}* | A set of 8 unsigned 16-bit quantities that all have the value specified by the integer. |
| (vector unsigned short)*{unsigned int, ..., unsigned int}* | A set of 8 unsigned 16-bit quantities specified by the 16 integers. |
| (vector signed short)*{signed int}* | A set of 8 signed 16-bit quantities that all have the value specified by the integer. |
| (vector signed short)*{signed int, ..., signed int}* | A set of 8 signed 16-bit quantities specified by the 16 integers. |
| (vector unsigned int)*{unsigned int}* | A set of 4 unsigned 32-bit quantities that all have the value specified by the integer. |
| (vector unsigned int)*{unsigned int, ..., unsigned int}* | A set of 4 unsigned 32-bit quantities specified by the 16 integers. |
| (vector signed int)*{signed int}* | A set of 4 signed 32-bit quantities that all have the value specified by the integer. |
| (vector signed int)*{signed int, ..., signed int}* | A set of 4 signed 32-bit quantities specified by the 16 integers. |
| (vector float)*{float}* | A set of 4 32-bit floating-point quantities that all have the value specified by the float. |
| (vector float)*{float, float, float, float}* | A set of 4 32-bit floating-point quantities that all have the value specified by the 4 floats. |

### A.4.3 Intrinsics

The intrinsics are shown in *Table A-13* on page 757.

*Table A-13. Vector/SIMD Multimedia Extension Intrinsics* (Sheet 1 of 4)

| Intrinsic | Description |
|---|---|
| **Arithmetic** | |
| d = vec_abs(a) | Vector Absolute Value |
| d = vec_abss(a) | Vector Absolute Value Saturated |
| d = vec_add(a,b) | Vector Add |
| d = vec_addc(a,b) | Vector Add Carryout Unsigned Word |
| d = vec_adds(a,b) | Vector Add Saturated |
| d = vec_avg(a,b) | Vector Average |
| d = vec_madd(a,b,c) | Vector Multiply Add |
| d = vec_madds(a,b,c) | Vector Multiply Add Saturated |
| d = vec_max(a,b) | Vector Maximum |
| d = vec_min(a,b) | Vector Minimum |
| d = vec_mladd(a,b,c) | Vector Multiply Low and Add Unsigned Half Word |
| d = vec_mradds(a,b,c) | Vector Multiply Round and Add Saturated |
| d = vec_msum(a,b,c) | Vector Multiply Sum |
| d = vec_msums(a,b,c) | Vector Multiply Sum Saturated |
| d = vec_mule(a,b) | Vector Multiply Even |
| d = vec_mulo(a,b) | Vector Multiply Odd |
| d = vec_nmsub(a,b,c) | Vector Negative Multiply Subtract |
| d = vec_sub(a,b) | Vector Subtract |
| d = vec_subc(a,b) | Vector Subtract Carryout |
| d = vec_subs(a,b) | Vector Subtract Saturated |
| d = vec_sum4s(a,b) | Vector Sum Across Partial (1/4) Saturated |
| d = vec_sum2s(a,b) | Vector Sum Across Partial (1/2) Saturated |
| d = vec_sums(a,b) | Vector Sum Saturated |
| **Rounding And Conversion** | |
| d = vec_ceil(a) | Vector Ceiling |
| d = vec_ctf(a,b) | Vector Convert from Fixed-Point Word |
| d = vec_cts(a,b) | Vector Convert to Signed Fixed-Point Word Saturated |
| d = vec_ctu(a,b) | Vector Convert to Unsigned Fixed-Point Word Saturated |
| d = vec_floor(a) | Vector Floor |
| d = vec_trunc(a) | Vector Truncate |
| **Floating-Point Estimate** | |
| d = vec_expte(a) | Vector Is 2 Raised to the Exponent Estimate Floating-Point |
| d = vec_loge(a) | Vector Log2 Estimate Floating-Point |
| d = vec_re(a) | Vector Reciprocal Estimate |
| d = vec_rsqrte(a) | Vector Reciprocal Square Root Estimate |

*Table A-13. Vector/SIMD Multimedia Extension Intrinsics* (Sheet 2 of 4)

| Intrinsic | Description |
|---|---|
| **Compare** | |
| d = vec_cmpb(a,b) | Vector Compare Bounds Floating-Point |
| d = vec_cmpeq(a,b) | Vector Compare Equal |
| d = vec_cmpge(a,b) | Vector Compare Greater Than or Equal |
| d = vec_cmpgt(a,b) | Vector Compare Greater Than |
| d = vec_cmple(a,b) | Vector Compare Less Than or Equal |
| d = vec_cmplt(a,b) | Vector Compare Less Than |
| **Logical** | |
| d = vec_and(a,b) | Vector Logical AND |
| d = vec_andc(a,b) | Vector Logical AND with Complement |
| d = vec_nor(a,b) | Vector Logical NOR |
| d = vec_or(a,b) | Vector Logical OR |
| d = vec_xor(a,b) | Vector Logical XOR |
| **Rotate and Shift** | |
| d = vec_rl(a,b) | Vector Rotate Left |
| d = vec_round(a) | Vector Round |
| d = vec_sl(a,b) | Vector Shift Left |
| d = vec_sld(a,b,c) | Vector Shift Left Double |
| d = vec_sll(a,b) | Vector Shift Left Long |
| d = vec_slo(a,b) | Vector Shift Left by Octet |
| d = vec_sr(a,b) | Vector Shift Right |
| d = vec_sra(a,b) | Vector Shift Right Algebraic |
| d = vec_srl(a,b) | Vector Shift Right Long |
| d = vec_sro(a,b) | Vector Shift Right by Octet |
| **Load and Store** | |
| d = vec_ld(a,b) | Vector Load Indexed |
| d = vec_lde(a,b) | Vector Load Element Indexed |
| d = vec_ldl(a,b) | Vector Load Indexed Least Recently Used (LRU) |
| d = vec_lvsl(a,b) | Vector Load for Shift Left |
| d = vec_lvsr(a,b) | Vector Load Shift Right |
| vec_st(a,b,c) | Vector Store Indexed |
| vec_ste(a,b,c) | Vector Store Element Indexed |
| vec_stl(a,b,c) | Vector Store Indexed LRU |
| **Pack and Unpack** | |
| d = vec_pack(a,b) | Vector Pack |
| d = vec_packpx(a,b) | Vector Pack Pixel |
| d = vec_packs(a,b) | Vector Pack Saturated |

*Table A-13. Vector/SIMD Multimedia Extension Intrinsics* (Sheet 3 of 4)

| Intrinsic | Description |
|---|---|
| d = vec_packsu(a,b) | Vector Pack Saturated Unsigned |
| d = vec_unpackh(a) | Vector Unpack High Element |
| d = vec_unpackl(a) | Vector Unpack Low Element |
| **Merge** | |
| d = vec_mergeh(a,b) | Vector Merge High |
| d = vec_mergel(a,b) | Vector Merge Low |
| **Permute and Select** | |
| d = vec_perm(a,b,c) | Vector Permute |
| d = vec_sel(a,b,c) | Vector Select |
| **Stream** | |
| vec_dss(a) | Vector Data Stream Stop |
| vec_dssall() | Vector Stream Stop All |
| vec_dst(a,b,c) | Vector Data Stream Touch |
| vec_dstst(a,b,c) | Vector Data Stream Touch for Store |
| vec_dststt(a,b,c) | Vector Data Stream Touch for Store Transient |
| vec_dstt(a,b,c) | Vector Data Stream Touch Transient |
| **Move** | |
| d = vec_mfvscr | Vector Move from Vector Status and Control Register |
| vec_mtvscr(a) | Vector Move to Vector Status and Control Register |
| **Replicate** | |
| d = vec_splat(a,b) | Vector Splat |
| d = vec_splat_s8(a) | Vector Splat Signed Byte |
| d = vec_splat_s16(a) | Vector Splat Signed Half-Word |
| d = vec_splat_s32(a) | Vector Splat Signed Word |
| d = vec_splat_u8(a) | Vector Splat Unsigned Byte |
| d = vec_splat_u16(a) | Vector Splat Unsigned Half-Word |
| d = vec_splat_u32(a) | Vector Splat Unsigned Word |
| **All Predicates** | |
| d = vec_all_eq(a,b) | All Elements Equal |
| d = vec_all_ge(a,b) | All Elements Greater Than or Equal |
| d = vec_all_gt(a,b) | All Elements Greater Than |
| d = vec_all_in(a,b) | All Elements in Bounds |
| d = vec_all_le(a,b) | All Elements Less Than or Equal |
| d = vec_all_lt(a,b) | All Elements Less Than |
| d = vec_all_nan(a) | All Elements Not a Number |
| d = vec_all_ne(a,b) | All Elements Not Equal |
| d = vec_all_nge(a,b) | All Elements Not Greater Than or Equal |

*Table A-13. Vector/SIMD Multimedia Extension Intrinsics* (Sheet 4 of 4)

| Intrinsic | Description |
|---|---|
| d = vec_all_ngt(a,b) | All Elements Not Greater Than |
| d = vec_all_nle(a,b) | All Elements Not Less Than or Equal |
| d = vec_all_nlt(a,b) | All Elements Not Less Than |
| d = vec_all_numeric(a) | All Elements Numeric |
| **Any Predicates** | |
| d = vec_any_eq(a,b) | Any Element Equal |
| d = vec_any_ge(a,b) | Any Element Greater Than or Equal |
| d = vec_any_gt(a,b) | Any Element Greater Than |
| d = vec_any_le(a,b) | Any Element Less Than or Equal |
| d = vec_any_lt(a,b) | Any Element Less Than |
| d = vec_any_nan(a) | Any Element Not a Number |
| d = vec_any_ne(a,b) | Any Element Not Equal |
| d = vec_any_nge(a,b) | Any Element Not Greater Than or Equal |
| d = vec_any_ngt(a,b) | Any Element Not Greater Than |
| d = vec_any_nle(a,b) | Any Element Not Less Than or Equal |
| d = vec_any_nlt(a,b) | Any Element Not Less Than |
| d = vec_any_numeric(a) | Any Element Numeric |
| d = vec_any_out(a,b) | Any Element Out of Bounds |

## A.5 Issue Rules

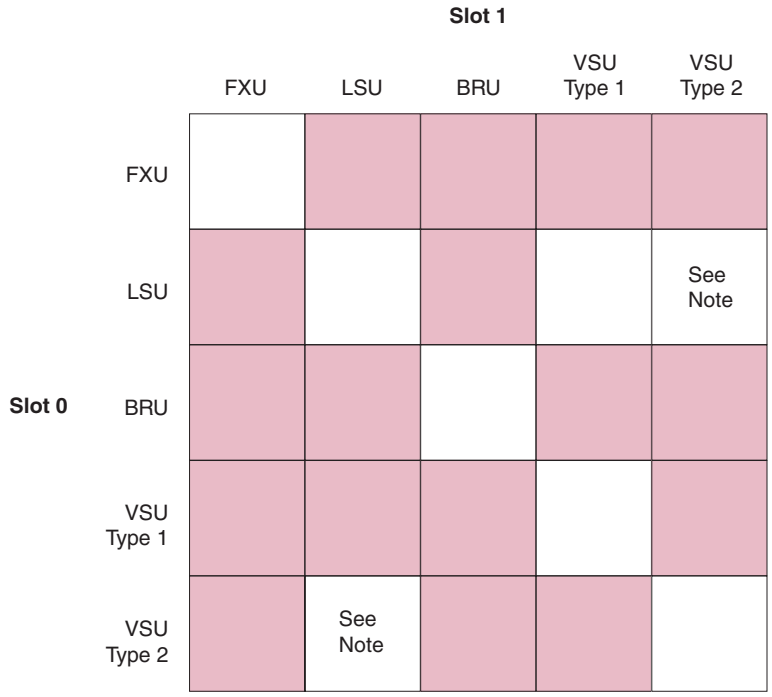
The PPE supports dual-issue of PowerPC and vector/SIMD multimedia extension instructions at both the pipeline stages at which instructions are issued. However, there are several restrictions that limit which instructions can be issued together.

*Figure A-1* on page 761 illustrates valid issue combinations. The instruction in slot 0 (the older instruction) is shown on the left vertical axis and the instruction in slot 1 (the younger instruction) is shown on the top horizontal axis. The pink squares demonstrate valid dual-issue combinations, and the white squares indicate illegal combinations. The vector/SIMD multimedia extension unit (VXU) and floating-point unit (FPU) instructions are categorized into two groups:

- Type 1: VXU simple, VXU complex, VXU floating-point, and FPU arithmetic instructions
- Type 2: VXU load, VXU store, VXU permute, FPU load, and FPU store instructions

Also, some instructions are classified in more than one group. For example, all fixed-point store and load-with-update instructions require both the fixed-point unit (FXU) and the load and store unit (LSU) to execute. In this case, both rows or columns for FXU and LSU must be checked to see if an issue conflict will occur. The classification of each *PowerPC Architecture* instruction can be found in *Table A-1* on page 723.

*Figure A-1. Dual-Issue Combinations*




| | |
|---|---|
| BRU | Branch Unit |
| FPU | Floating-Point Unit |
| FXU | Fixed-Point Unit |
| LSU | Load and Store Unit |
| VSU | Vector Scalar Unit (a combination of VXU and FPU) |
| VXU | Vector/SIMD Multimedia Extension Unit |

**Note:** A VXU Permute instruction in slot 0 can dual issue with an LSU instruction in slot 1.

There are several specific cases that prevent dual instruction issue. These are:

- A dependency stall exists. Input parameters are not yet available.
- The even instruction (in slot 0—the older instruction in a dispatch pair) is a nonpipelined instruction. The nonpipelined instructions include **mfspr**, **mulli**, **mullw**, **mullwo**, **mulhu**, **mulhwu**, **mulld**, **mulldo**, **mulhd**, **mulhdu**, **divd**, **divdu**, **divdo**, **divduo**, **fdiv**, **fdivs**, **fsqrt**, **fsqrts**, and **mtvscr**.
- The even instruction is a context-synchronizing instruction. The context-synchronizing instructions include **sync** (L = 0, 1, or 2), **isync**, **mtctrl**, **mtmsr** and **mtmsrd** (L = 0), **sc**, **rfid**, and **hrfid**.
- One of the instructions is a microcoded instruction. The PPE has two types of microcoded instructions—unconditional and conditional. The unconditional microcoded instructions include shifts and rotates that source the shift amount from a register (as opposed to an immediate field), load/store algebraic instructions, load/store strings and multiples, and sev-

eral CR recording (Rc=1) instructions. The conditional microcoded instructions are all load and store instructions and occur conditionally based upon the alignment of the load or store.

- The FPU is running single-step as a result of precise floating-point exceptions being enabled (machine-state register bits MSR[FE0] or MSR[FE1] being set to 1).
- The issue queue for the vector scalar unit (VSU, a combination of the VXU and FPU) is single-stepping as a result of an internal flush. Internal flushes occur when VSU instructions have denormalized operands or produce denormalized results.
- A VXU floating-point instruction (except for a reciprocal-estimate operation) is issued within two cycles following a reciprocal estimate instruction, **vrefp** or **vrsqrtefp**.
- A **mffs** (move from FPSCR) instruction is issued one cycle following an FPU instruction, other than an FPU load or store.
- No instruction can be issued following a nonpipelined FPU or VXU instruction until the nonpipelined instruction completes.
- No FPU instruction, including loads and stores, can be issued after a **fdiv** or **fsqrt** until the **fdiv** or **fsqrt** completes.
- A younger (slot 1) VSU load or store can be issued in the same cycle as any other VSU type 1 instruction, but an older VXU or FPU load or store cannot.
- Nothing issues while an internal flush condition is resolved. During this time, instructions are issued from the denorm-recycle queue in a single-step fashion.

Additional information can be found in *Section A.7* on page 767.

## A.6 Pipeline Stages

The instruction unit (IU) is responsible for all instruction handling and control. After instructions are issued from the IU, they enter one of the execution units (LSU, FXU, VXU, FPU, or BRU). The BRU is an execution unit, but is considered part of the IU because it deals with instruction branching.

### A.6.1 Instruction-Unit Pipeline

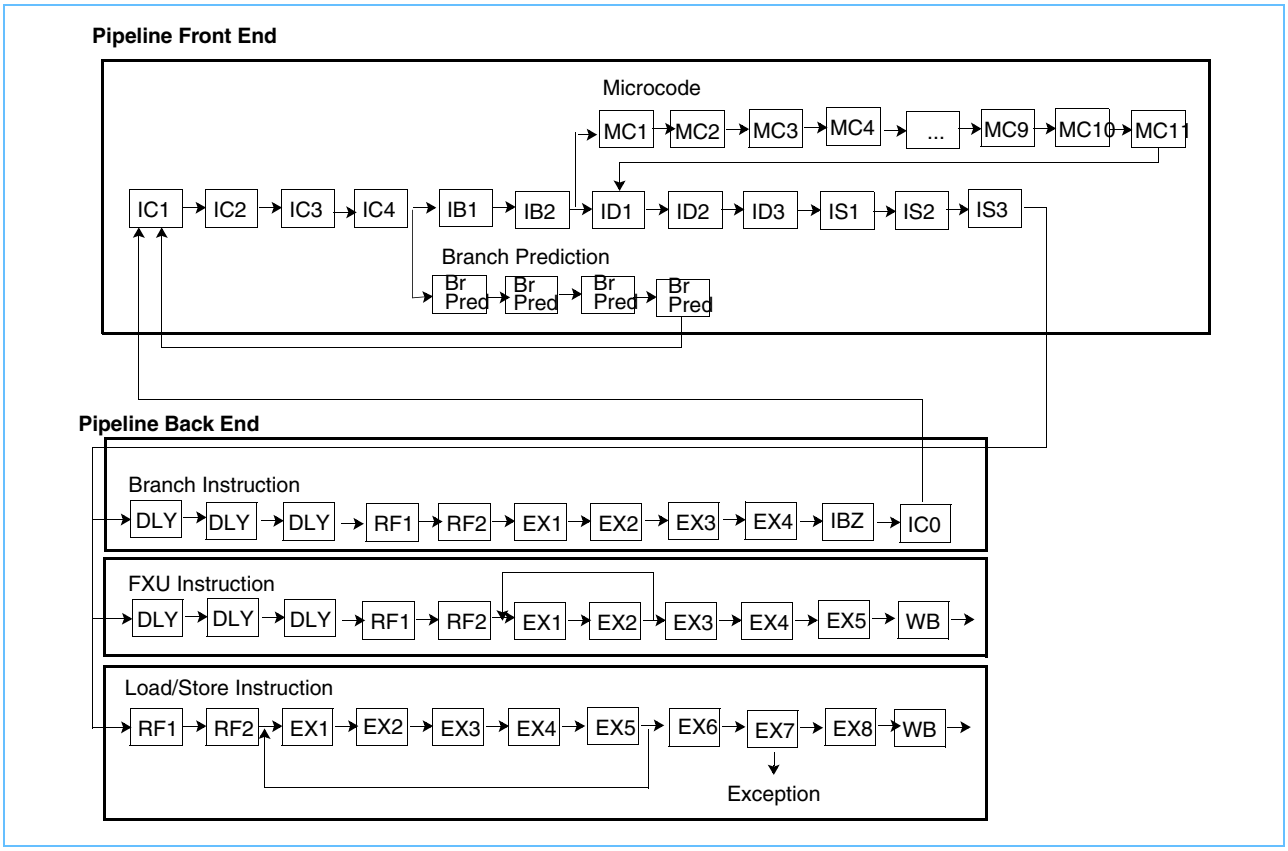
*Figure A-2* on page 763 shows the portion of the IU pipeline that serves of the BRU, FXU, and LSU. The front end of the IU pipeline consists of a 4-cycle cache access: IC1 through IC4. The fetch address is based on either the sequential path or the branch target path. The branch execution stage can correct the instruction fetch with the branch-target address if the branch is mispredicted not taken. The branch execution stage can correct the instruction fetch with the sequential address if the branch is mispredicted taken.

After pipeline stage IC4, there are two instruction-buffer stages: IB1 and IB2. The instruction buffer (IBuf) is a first in, first out (FIFO) queue. It is used to buffer the four instructions fetched from the L1 ICache when there is a downstream stall condition. Instruction-buffer stage IB1 is used to load the IBuf; there is one set of IBufs for each thread. IB2 is used to unload the IBuf and multiplex down to two instructions (this is the instruction dispatch stage). Each thread is given equal priority in dispatch, toggling every other cycle, unless specified otherwise in software (see *Section 10 PPE Multithreading* on page 299). If one thread is not fetching instructions into the IBuf (due to conditions such as a cache miss or thread disabled), then the other thread may have

exclusive access to dispatch. Dispatch also controls the flow of instructions to and from micro-code. Microcode (MC) is used to break an instruction that is difficult to execute into multiple smaller operations.

The stages labeled RF are GPR read cycles. Those labeled EX are execution cycles. Those labeled WB are write-back.

*Figure A-2. BRU, FXU, and LSU Pipeline for PPE*



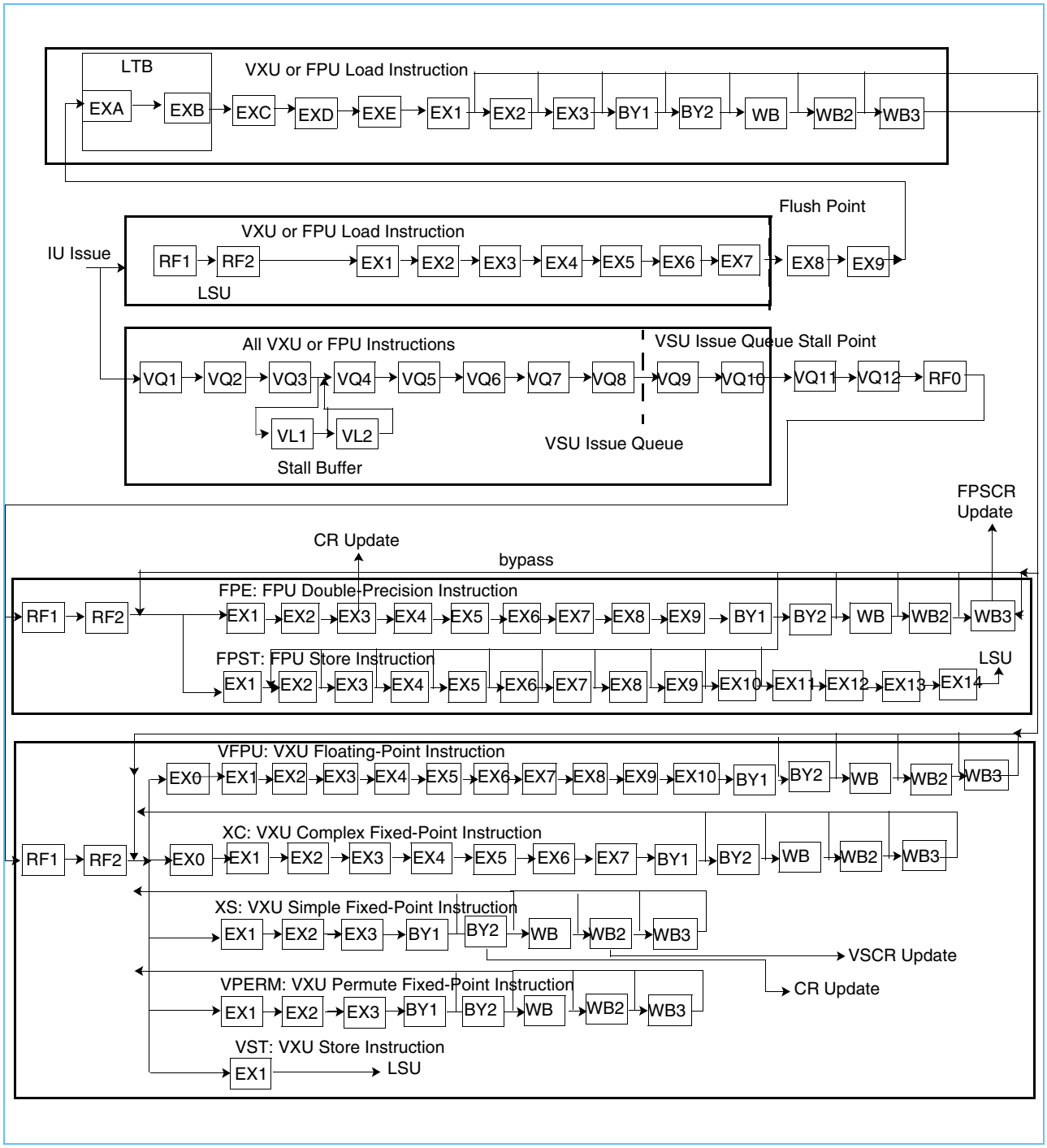

Dispatch grouping occurs based on the effective address bits [60:63] of the instruction. The dispatcher groups instructions with addresses x‘0’ and x‘4’, as well as instructions with addresses x‘8’ and x‘C’. An instruction with address x‘0’ cannot be grouped with an instruction with address x‘8’ or x‘C’. Similarly, an instruction with address x‘4’ cannot be grouped with an instruction with address x‘8’ or x‘C’. In other words, instructions grouped together must be aligned on 8-byte boundaries.

Instruction-decode pipeline stages ID1 through ID3 are used to assemble the internal opcodes and register source and target fields. In addition, dependency checking starts in ID2, which checks for data hazards, such as read-after-write (RAW) or write-after-write (WAW). The issue logic continues in ID3, IS1, and IS2 to create a single stall point at IS2, which is propagated up the pipeline to the IBufs, stalling both threads. The IS2 stall point is driven by data-hazard detection, in addition to resource-conflict detections, among other conditions. The IS2 issue stage determines the appropriate routing of the instructions. Then, they are issued to the execution units in the IS3 cycle. Each instruction in IS2 can be routed to five different issue slots: to the FXU, LSU, BRU, and to two slots in the VSU issue queue, described in the next section.

### A.6.2 Vector/Scalar Unit Issue Queue

*Figure A-3* shows the VSU issue queue (VIQ). VSU instructions are issued from the IU's IS2 stage to the VIQ without checking for dependencies or issue restrictions. The VIQ has a separate dependency-checking facility, used exclusively for VSU instructions. A separate stall point is generated at the bottom of the VIQ in cycle VQ8.

*Figure A-3. VSU (FPU and VXU) Pipeline for PPE*



Instructions are sent in pairs to the VIQ in cycle IS2. If there are no VSU instructions to issue, a bubble propagates down the queue. The queue does not compress the bubbles. In other words, the queue is really a pipeline. The VL1 and VL2 stages of the queue represent an overflow buffer. When a VQ8 stall occurs, this buffer continues to allow instructions to flow for an additional two cycles. The IS2 point will stall two cycles after a VQ8 stall only if there is a VSU instruction that would overflow the VIQ if issued. Otherwise, the IS2 point can continue issuing FXU, LSU, and BRU instructions.

After VQ8, instructions are sent to the VXU or FPU. The VQ9, VQ10, VQ11, and VQ12 stages are transmit cycles required to physically send the data to the execution unit.

FPU and VXU instructions can have denormalized operands as described later in this section. Therefore, a denorm recycle queue is maintained to track each instruction past issue until the VXU or FPU signals that no internal flushes are necessary for the instruction. If no such denormalized condition occurs, the instruction completes and is dropped from the denorm recycle queue. If an internal flush occurs, the instruction and all younger instructions that are past VQ8 are reissued to the VSU in a single-step fashion until the denorm recycle queue is empty. Instructions are issued to the execution units and to the denorm recycle queue in parallel.

VXU and FPU loads are sent to the LSU to access the L1 DCache, in addition to being sent to the VIQ. When data is available for these loads, it is sent to the VXU load target buffer (VLTB) or FPU load target buffer (FLTB) as appropriate for the type of instruction. The load data is then held in these buffers until the instruction is issued from the VIQ. The VLTB and FLTB are 16 entries deep to accommodate a total of 16 loads in flight past the IS2 issue point.

If a VXU or FPU load misses the L1 DCache, it is issued to the VXU and FPU load miss queue (VMQ). It is held in the VMQ until the cache miss is resolved in the LSU. When the miss is resolved, the load is reissued to the VSU.

Internal flushes are caused in the VSU for the following reasons:

- An FPU or VXU floating-point instruction with a denormalized operand is encountered.
- An FPU instruction with a Not a Number (NaN) operand (as defined in *PowerPC User Instruction Set Architecture, Book I*) is encountered.
- An FPU instruction uses a bypassed suppressed result caused by an FPU enabled exception.
- The zero divide or invalid operation exceptions are enabled in the FPU when both were previously disabled.

### A.6.3 Stall and Flush Points

In normal instruction execution, everything will flow in an orderly pipelined fashion. Several conditions exist in which this normal execution is interrupted with a stall or a flush condition. There are three separate stall points for the PPE.

The first stall point occurs just after the instruction buffers in the IB2 stage. This stall point is commonly referred to as a *dispatch block* because it prevents instructions from being dispatched. There are several possible reasons for a dispatch block. The following are two of the more common reasons:

- A special **nop** instruction (see *Section 10.6.2.2 nop Instructions that Change Thread Priority and Dispatch Policy* on page 320).

- The flush logic, whenever an L1 DCache miss dependency (RAW or WAW) or D-ERAT miss occurs; or whenever a Condition Register Field 0 (CR0) source-dependent operation issues while an **stdcx.** or **stwcx.** instruction is pending.

Stalling at this point in the pipeline allows the other thread to continue to dispatch instructions while the current thread is stalled. All other stall points in the pipeline block both threads.

The second stall point occurs in the Issue logic at the IS2 stage. This stall point is activated by the hardware whenever one of the following conditions is met:

- An LSU or FXU instruction dependency occurs.
- A nonpipelined instruction is issued.
- An invalid dual-issue combination is present.
- A context serializing instruction is issued.
- The processor is in single-stepping mode.
- The lower VSU stall point (in cycle VQ8) is active, and a VSU instruction is presently trying to issue to the VIQ.

This stall point is also activated whenever a stall request is received from the LSU. This stall request is always honored if there is a valid instruction at IS2. A stall at IS2 also causes the microcode pipeline to stall.

The third stall point is exclusive to the VIQ (in cycle VQ8). It is activated for the following conditions:

- VXU or FPU dependencies
- VXU or FPU write-port conflicts on the Vector Registers (VRs) or Floating-Point Registers (FPRs), invalid dual-issue combinations
- Nonpipelined instruction issue
- Special single-stepping conditions

This stall point is also activated whenever a stall request is received from the LSU; this stall request is always honored.

A 2-stage buffer separates the third stall point from the second stall point to help decouple the FXU and LSU instruction flow from the VXU and FPU instruction flow. If this buffer fills up and there is an instruction in the IS2 stage that needs to issue to the VIQ, then the second stall point will activate to prevent overflowing the VIQ.

When an instruction is past all stall points, it must either flow to completion or be flushed. A flush is initiated by one of the execution units in the back end of the pipeline for various reasons, such as the occurrence of a dependency on a cache miss or an exception condition. Whenever a flush occurs because of an exception, an L1 DCache-miss dependency, or a D-ERAT address-translation miss, it will be taken at the ninth stage of the execution unit (EX7). All instructions in the machine that have been fetched from the L1 ICache and are younger (in program order) than the flushing instruction will be invalidated.

The VXU and FPU contain an internal flush for instructions that have denormalized operands. (For more information about denormalization, see *PowerPC User Instruction Set Architecture, Book I.)* In this case, all instructions that have issued out of the VIQ and are younger than the denormalized instruction will be invalidated and then reexecuted. To facilitate this, there is a

denormalized instruction queue (DIQ) that tracks all instructions that are issued from the VIQ. When the internal flush condition is signaled, the VIQ is stalled. Instructions in the DIQ are reissued one at a time (each instruction must complete before the next is issued).

## A.7 Compiler Optimizations

The following compiler optimizations apply to all software written for the PPE, using either the PowerPC or vector/SIMD multimedia extension instruction sets. These optimizations will improve the performance of the PPE.

For many additional compiler optimizations, including those specifically for single instruction, multiple data (SIMD) operations, see *Section 22* on page 629.

### A.7.1 Instruction Arrangement

PPE instructions are dispatched in 8-byte aligned pairs. The smaller address of the pair is termed I0 (EA[60:63] = x‘0’ or x‘8’); the larger address of the pair is termed I1 (EA[60:63] = x‘4’ or x‘C’). By convention, I0 is termed “older” than I1 and I1 is termed “younger” than I0.

See *Section A.5* on page 760 for an overview of issue rules and the names of execution units. The following rules apply:

- I0 and I1 should not attempt to use the same execution unit (BRU, LSU, FXU, VSU type 1, or VSU type 2). Doing so causes a 1-cycle stall for I1.
- I0 and I1 should not both access (read or write) the CR or LR registers. Doing so causes a 1-cycle stall for I1.
- If I0 is a load or store instruction and I1 is a any vector/SIMD multimedia extension or FPU instruction other than a load or store and if I1 is not dependent on I0, then the order of I0 and I1 should be reversed. Failure to do so will cause a 1-cycle stall for I1.
- Nonpipelined operations (see *Section A.7.2*) and context-synchronizing instructions (CSI, see *Section A.7.2*) should be placed in I1. If placed in I0, then the I1 instruction cannot dual-issue with the second-issue of the double-issue operation (causing a 1-cycle stall penalty for the I1 instruction.
- The target of a branch instruction should be to an x‘0’ or x‘8’ (I0) address. This allows the dispatcher to avoid an unpaired instruction (resulting in a 1-cycle inefficiency).
- Predicted-taken branches should have an address of x‘4’ or x‘8’ for maximum dispatch efficiency.

### A.7.2 Avoiding Slow Instructions and Processor Modes

For better performance, software should avoid setting MSR[FE0] or MSR[FE1] to ‘1’. When either of these is set, the processor will run in a “single-step” mode whereby each instruction must wait for all older instructions to complete before getting issued.

Because of the 11-cycle penalty for beginning a microcode sequence, it is best for performance if software avoids instructions that go to microcode. See *Table A-3* on page 735, *Table A-4* on page 737, and *Table A-5* on page 738.

The following are context-synchronizing instructions (CSI). These instructions require all older instructions to complete before they are issued, and are therefore best to avoid in performance-critical code:

- **sync** (L = 0, 1, or 2)
- **isync**
- **mtctrl**
- **mtmsr**, **mtmsrd** L = 0
- **sc**
- **rfid**, **hrfid**

The following are nonpipelined instructions. These instructions require all younger instructions to stall until they are completed:

- **mfspr** LR or **mfspr** CTR—inserts a 2-cycle stall.
- all other **mfspr** instructions—inserts a minimum 9-cycle stall.
- **mulli**—inserts a 6-cycle stall.
- **mullw**, **mullwo**, **mulhw**, **mulhwu**—inserts a 9-cycle stall.
- **mulld**, **mulldo**, **mulhd**, **mulhdu**—inserts a 15-cycle stall.
- **divd**, **divdu**, **divdo**, **divduo**—inserts a 10 to 70 cycle stall.
- **divw**, **divwu**, **divwo**, **divwuo**—inserts a 10 to 70 cycle stall.
- **fdiv**, **fdivs**, **fsqrt**, **fsqrts**, **mtvscr**—inserts a minimum 13 cycle stall.

The vector/SIMD multimedia extension and FPU instructions that have either denormalized or underflow (NaN) operands or results should be avoided. If such an instruction is encountered, an internal flush will occur in the VIQ, resulting in a temporary single-step mode. This will cost a minimum of 100 cycles.

### A.7.3 Avoiding Dependency Stalls and Flushes

The following dependencies should be avoided if possible:

- An instruction that reads the XER while a **stwcx.** or **stdcx.** instruction is pending will be flushed.
- An instruction that reads the XER should not be I1 in the same dispatch group or I0 or I1 in the following dispatch group as an XER writing instruction. If so, the issue logic will stall to separate them by a 1-cycle bubble.
- Two instructions that read or write the CTR should not be in the same or back-to-back dispatch groups. If so, the issue logic will stall to separate them by a 1-cycle bubble.
- In general, FXU instructions that read the CR should not access the same CR fields that are in use by vector/SIMD multimedia extension or FPU instructions. Doing so can cause a flush if the FXU instruction issues before the vector/SIMD multimedia extension or FPU instruction has written results to the CR field.
- An FXU instruction that is dependent on an older instruction should not be in the same or the following dispatch group as the older instruction. If so, the issue logic will stall to separate them by a 1-cycle bubble.

- An LSU instruction that is dependent on an older LSU instruction should not be in the same or the following dispatch group as the older instruction. If so, the issue logic will stall to separate them by a 1-cycle bubble.
- An LSU instruction that is dependent on an older FXU instruction should not be in the same or the following four dispatch groups as the older instruction. If so, the issue logic will stall to separate them by a 4-cycle bubble.

  There is an exception for a store instruction that has only data (not an address) dependent on an older FXU instruction. In this case the store should not be in the same or the following dispatch group. If it is, then a 4-cycle bubble will occur (even though it appears that only a 1-cycle bubble should be necessary).

For vector/SIMD multimedia extension and FPU instructions, the following read-after-write (RAW) dependencies should be avoided if possible:

- An instruction dependent on an older FPU arithmetic result should not be in the same or the following nine dispatch groups as the older instruction. If so, the issue logic will stall to separate them by a 9-cycle bubble. The exception is if the dependent instruction is an FPU store that should not be in the same dispatch group as the FPU arithmetic (or a 1-cycle stall will result).
- An instruction dependent on an older vector/SIMD multimedia extension Simple or vector/SIMD multimedia extension Permute result should not be in the same or the following three dispatch groups as the older instruction. If so, the issue logic will stall to separate them by a 3-cycle bubble.
- An instruction dependent on an older vector/SIMD multimedia extension Complex result should not be in the same or the following eight dispatch groups as the older instruction. If so, the issue logic will stall to separate them by a 3-cycle bubble.
- An instruction dependent on an older vector/SIMD multimedia extension Float result should not be in the same or the following eleven dispatch groups as the older instruction. If so, the issue logic will stall to separate them by a 11-cycle bubble.
- An instruction dependent on an older vector/SIMD multimedia extension Estimate result should not be in the same or the following thirteen dispatch groups as the older instruction. If so, the issue logic will stall to separate them by a 13-cycle bubble.

For vector/SIMD multimedia extension and FPU instructions, the following write-after-write (WAW) dependencies should be avoided if possible. Software can reuse the same target for instructions that execute in the same pipeline without penalty. For example, a vector/SIMD multimedia extension Float instruction can reuse the same target register as a previous vector/SIMD multimedia extension Float immediately without any penalty.

- The target of a vector/SIMD multimedia extension Estimate instruction should not be reused within ten dispatch groups.
- The target of a vector/SIMD multimedia extension Float instruction should not be reused within eight dispatch groups.
- The target of a vector/SIMD multimedia extension Complex instruction should not be reused within five dispatch groups.
- An FPU load should not reuse an FPU arithmetic target within six dispatch groups.

For vector/SIMD multimedia extension and FPU instructions, there exist write-port-collisions (WPCs). In general these are difficult to avoid in a precise fashion other than to avoid dispatch combinations similar to the following:

- A vector/SIMD multimedia extension Simple instruction must stall for 1 cycle if a vector/SIMD multimedia extension Complex was issued 5 cycles ago, a vector/SIMD multimedia extension Float was issued 8 cycles ago, or a vector/SIMD multimedia extension Estimate was issued 10 cycles ago.
- A vector/SIMD multimedia extension Complex instruction must stall for 1 cycle if a vector/SIMD multimedia extension Float was issued 2 cycles ago, or a vector/SIMD multimedia extension Estimate was issued 4 cycles ago.
- A vector/SIMD multimedia extension Float instruction must stall for 1 cycle if a vector/SIMD multimedia extension Estimate was issued 2 cycles ago.
- A vector/SIMD multimedia extension Store instruction must stall for 1 cycle if a FPU Store was issued 12 cycles ago.
- An FPU Load instruction must stall for 1 cycle if an FPU Arithmetic instruction was issued 6 cycles ago.

Finally, VSU type 1 instructions should avoid writing the same target register as a VSU type 2 instruction within 12 cycles of each other. This prevents a potential 1-cycle stall that can occur if the writes occur in the same cycle.

### A.7.4 General Recommendations

The following general recommendations can be used to help improve performance:

- Use the data-cache block touch (cache hint) instructions (**dcbt**).
- Use the cache replacement management facility for controlling the L2 cache and TLB replacement policy based on a class identifier (class ID).
- Preload the TLB.
- Use all of the branch-prediction hints given in the branch instructions.
- Avoid load-hit-store[1] dependencies.
- Avoid D-ERAT thrashing by allocating data addresses so that collisions are less likely to occur on 128 KB boundaries (2-way set associative with 32 × 4 KB entries).
- Avoid flush conditions due to filling up the fixed-point unit store queue by having no more than 15 stores outstanding at one time.
- Avoid VSU flush conditions due to filling up the VSU load target buffer by having no more than 15 VSU loads outstanding at one time.

1. The load-hit-store (LHS) condition occurs when a load request is made to store to a similar address in the processor's store queue. Because the store queue in the processor holds stores that have not yet written into the cache, the load must wait in the load miss queue (LMQ) until the store has written its data into the cache. Loads that are marked as LHS can be recycled after all older stores have completed.

# Appendix B. SPU Instruction Set and Intrinsics

The Synergistic Processor Elements (SPEs) support the synergistic processor unit (SPU) instruction set architecture (ISA), and the C/C++ language extensions for this instruction set provide data types and intrinsics to programmers writing in C or C++ when they use the `spu_intrinsics.h` header file. This section summarizes both the ISA and the extensions. Because the functions performed by the intrinsics are closely related to the assembly-language instructions of the SPU instruction set, this overview can be helpful in understanding the functions of the intrinsics.

The SPU ISA operates primarily on single instruction, multiple data (SIMD) vector operands, both fixed-point and floating-point, with support for some scalar operands. The PowerPC Processor Element (PPE) and the SPE both execute SIMD instructions, but the two processors execute different instruction sets, and programs for the PPE and SPEs must be compiled by different compilers.

For a complete description of the SPU instructions, intrinsics, and their use, see:

- *Synergistic Processor Unit (SPU) Instruction Set Architecture*
- *C/C++ Language Extensions for Cell Broadband Engine Architecture*
- *SPU Application Binary Interface Specification*
- *SPU Assembly Language Specification*

## B.1 SPU Instruction Set

The SPU instruction set uses instructions that are 4 bytes long and word-aligned. It supports 16-byte operand accesses between storage and its 128 vector registers. For a brief overview of the SPU instruction set, including the data types, addressing modes and instruction types, see *Section 3.3* on page 76. The sections that follow summarize key points of the instruction set.

### B.1.1 Data Types

*Section 3.3.1* on page 76 contains a list of the SPU data types.

### B.1.2 Instructions

*Table B-1* on page 772 contains a complete list of SPU instructions, ordered alphabetically. The SPU has two pipelines, named even (pipeline 0) and odd (pipeline 1), into which it can issue and complete up to two instructions per cycle, one in each of the pipelines. The pipeline is identified for each instruction in the table, along with the latency and stall cycles. The latency values include all stall cycles.

A stall is the number of cycles, after issuing an instruction of a given type, before another instruction of the same type can be issued. For example, double-precision floating-point operations have a 6-cycle stall. Therefore, for sequential double-precision floating-point operations, the second operation will be issued at least 7 cycles after the first operation.

The SPU ISA supports only 16-bit multiplies, so 32-bit multiplies are implemented in software using 16-bit multiplies.

For details about the instruction issue rules, see *Section B.1.3* on page 779. For more information about the instructions, see the *Synergistic Processor Unit Instruction Set Architecture* specification and the *SPU Assembly Language Specification*. Instructions not supported on specific processor implementations are indicated by 'N/S'.

*Table B-1. SPU Instructions* (Sheet 1 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **a rt, ra, rb** | **Add word**. | 0 | 2 | 0 | 2 | 0 |
| **absdb rt, ra, rb** | **Absolute difference of bytes**. | 0 | 4 | 0 | 4 | 0 |
| **addx rt, ra, rb** | **Add word extended**. | 0 | 2 | 0 | 2 | 0 |
| **ah rt, ra, rb** | **Add halfword**. | 0 | 2 | 0 | 2 | 0 |
| **ahi rt, ra, s10** | **Add halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **ai rt, ra, s10** | **Add word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **and rt, ra, rb** | **And**. | 0 | 2 | 0 | 2 | 0 |
| **andbi rt, ra, s10** | **And byte immediate**. | 0 | 2 | 0 | 2 | 0 |
| **andc rt, ra, rb** | **And with complement**. | 0 | 2 | 0 | 2 | 0 |
| **andhi rt, ra, s10** | **And halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **andi rt, ra, s10** | **And word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **avgb rt, ra, rb** | **Average bytes**. | 0 | 4 | 0 | 4 | 0 |
| **bg rt, ra, rb** | **Borrow generate word**. | 0 | 2 | 0 | 2 | 0 |
| **bgx rt, ra, rb** | **Borrow generate word extended**. | 0 | 2 | 0 | 2 | 0 |
| **bi ra** | **Branch indirect**. | 1 | N/A | 0 | N/A | 0 |
| **bid ra** | **Branch indirect, disable**. | 1 | N/A | 0 | N/A | 0 |
| **bie ra** | **Branch indirect, enable**. | 1 | N/A | 0 | N/A | 0 |
| **bihnz rt, ra** | **Branch indirect if not zero halfword**. | 1 | N/A | 0 | N/A | 0 |
| **bihnzd rt, ra** | **Branch indirect if not zero halfword, disable**. | 1 | N/A | 0 | N/A | 0 |
| **bihnze rt, ra** | **Branch indirect if not zero halfword, enable**. | 1 | N/A | 0 | N/A | 0 |
| **bihz rt, ra** | **Branch indirect if zero halfword**. | 1 | N/A | 0 | N/A | 0 |
| **bihzd rt, ra** | **Branch indirect if zero halfword, disable**. | 1 | N/A | 0 | N/A | 0 |
| **bihze rt, ra** | **Branch indirect if zero halfword, enable**. | 1 | N/A | 0 | N/A | 0 |
| **binz rt, ra** | **Branch indirect if not zero word**. | 1 | N/A | 0 | N/A | 0 |
| **binzd rt, ra** | **Branch indirect if not zero word, disable**. | 1 | N/A | 0 | N/A | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 2 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **binze rt, ra** | **Branch indirect if not zero word, enable**. | 1 | N/A | 0 | N/A | 0 |
| **bisl rt, ra** | **Branch indirect and set link**. | 1 | 4 | 0 | 4 | 0 |
| **bisld rt, ra** | **Branch indirect and set link, disable**. | 1 | 4 | 0 | 4 | 0 |
| **bisle rt, ra** | **Branch indirect and set link, enable**. | 1 | 4 | 0 | 4 | 0 |
| **bisled rt, ra** | **Branch indirect and set link on external data**. | 1 | 4 | 0 | 4 | 0 |
| **bisledd rt, ra** | **Branch indirect and set link on external data, disable**. | 1 | 4 | 0 | 4 | 0 |
| **bislede rt, ra** | **Branch indirect and set link on external data, enable**. | 1 | 4 | 0 | 4 | 0 |
| **biz rt, ra** | **Branch indirect if zero word**. | 1 | N/A | 0 | N/A | 0 |
| **bizd rt, ra** | **Branch indirect if zero word, disable**. | 1 | N/A | 0 | N/A | 0 |
| **bize rt, ra** | **Branch indirect if zero word, enable**. | 1 | N/A | 0 | N/A | 0 |
| **br s18** | **Branch relative**. | 1 | N/A | 0 | N/A | 0 |
| **bra s18** | **Branch absolute**. | 1 | N/A | 0 | N/A | 0 |
| **brasl rt, s18** | **Branch absolute and set link**. | 1 | N/A | 0 | N/A | 0 |
| **brhnz rt, s18** | **Branch if not zero halfword**. | 1 | N/A | 0 | N/A | 0 |
| **brhz rt, s18** | **Branch if zero halfword**. | 1 | N/A | 0 | N/A | 0 |
| **brnz rt, s18** | **Branch if not zero word**. | 1 | N/A | 0 | N/A | 0 |
| **brsl rt, s18** | **Branch relative and set link**. | 1 | N/A | 0 | N/A | 0 |
| **brz rt, s18** | **Branch if zero word**. | 1 | N/A | 0 | N/A | 0 |
| **cbd rt, u7(ra)** | **Generate controls for byte insertion (d-form)**. | 1 | 4 | 0 | 4 | 0 |
| **cbx rt, ra, rb** | **Generate controls for byte insertion (x-form)**. | 1 | 4 | 0 | 4 | 0 |
| **cdd rt, u7(ra)** | **Generate controls for doubleword insertion (d-form)**. | 1 | 4 | 0 | 4 | 0 |
| **cdx rt, ra, rb** | **Generate controls for doubleword insertion (x-form)**. | 1 | 4 | 0 | 4 | 0 |
| **ceq rt, ra, rb** | **Compare equal word**. | 0 | 2 | 0 | 2 | 0 |
| **ceqb rt, ra, rb** | **Compare equal byte**. | 0 | 2 | 0 | 2 | 0 |
| **ceqbi rt, ra, s10** | **Compare equal byte immediate**. | 0 | 2 | 0 | 2 | 0 |
| **ceqh rt, ra, rb** | **Compare equal halfword**. | 0 | 2 | 0 | 2 | 0 |
| **ceqhi rt, ra, s10** | **Compare equal halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **ceqi rt, ra, s10** | **Compare equal word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **cflts rt, ra, scale7** | **Convert floating to signed integer**. | 0 | 7 | 0 | 7 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 3 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **cfltu rt, ra, scale7** | **Convert floating to unsigned integer**. | 0 | 7 | 0 | 7 | 0 |
| **cg rt, ra, rb** | **Carry generate word**. | 0 | 2 | 0 | 2 | 0 |
| **cgt rt, ra, rb** | **Compare greater than word**. | 0 | 2 | 0 | 2 | 0 |
| **cgtb rt, ra, rb** | **Compare greater than byte**. | 0 | 2 | 0 | 2 | 0 |
| **cgtbi rt, ra, s10** | **Compare greater than byte immediate**. | 0 | 2 | 0 | 2 | 0 |
| **cgth rt, ra, rb** | **Compare greater than halfword**. | 0 | 2 | 0 | 2 | 0 |
| **cgthi rt, ra, s10** | **Compare greater than halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **cgti rt, ra, s10** | **Compare greater than word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **cgx rt, ra, rb** | **Carry generate word extended**. | 0 | 2 | 0 | 2 | 0 |
| **chd rt, u7(ra)** | **Generate controls for halfword insertion (d-form)**. | 1 | 4 | 0 | 4 | 0 |
| **chx rt, ra, rb** | **Generate controls for halfword insertion (x-form)**. | 1 | 4 | 0 | 4 | 0 |
| **clgt rt, ra, rb** | **Compare logical greater than word**. | 0 | 2 | 0 | 2 | 0 |
| **clgtb rt, ra, rb** | **Compare logical greater than byte**. | 0 | 2 | 0 | 2 | 0 |
| **clgtbi rt, ra, s10** | **Compare logical greater than byte immediate**. | 0 | 2 | 0 | 2 | 0 |
| **clgth rt, ra, rb** | **Compare logical greater than halfword**. | 0 | 2 | 0 | 2 | 0 |
| **clgthi rt, ra, s10** | **Compare logical greater than halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **clgti rt, ra, s10** | **Compare logical greater than word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **clz rt, ra** | **Count leading zeros**. | 0 | 2 | 0 | 2 | 0 |
| **cntb rt, ra** | **Count ones in bytes**. | 0 | 4 | 0 | 4 | 0 |
| **csflt rt, ra, scale7** | **Convert signed integer to floating**. | 0 | 7 | 0 | 7 | 0 |
| **cuflt rt, ra, scale7** | **Convert unsigned integer to floating**. | 0 | 7 | 0 | 7 | 0 |
| **cwd rt, u7(ra)** | **Generate controls for word insertion (d-form)**. | 1 | 4 | 0 | 4 | 0 |
| **cwx rt, ra, rb** | **Generate controls for word insertion (x-form)**. | 1 | 4 | 0 | 4 | 0 |
| **dfa rt, ra, rb** | **Double-precision floating add**. | 0 | 13 | 6 | 9 | 0 |
| **dfceq rt, ra, rb** | **Double-precision floating compare equal** | 0 | N/S | N/S | 9 | 0 |
| **dfcgt rt, ra, rb** | **Double-precision floating compare greater than** | 0 | N/S | N/S | 9 | 0 |
| **dfcmeq rt, ra, rb** | **Double-precision floating compare magnitude equal** | 0 | N/S | N/S | 9 | 0 |
| **dfcmgt rt, ra, rb** | **Double-precision floating compare magnitude greater than** | 0 | N/S | N/S | 9 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 4 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **dfm rt, ra, rb** | **Double-precision floating multiply**. | 0 | 13 | 6 | 9 | 0 |
| **dfma rt, ra, rb** | **Double-precision floating multiply and add**. | 0 | 13 | 6 | 13 | 6 |
| **dfms rt, ra, rb** | **Double-precision floating multiply and subtract**. | 0 | 13 | 6 | 9 | 0 |
| **dfnma rt, ra, rb** | **Double-precision floating negative multiply and add**. | 0 | 13 | 6 | 9 | 0 |
| **dfnms rt, ra, rb** | **Double-precision floating negative multiply and subtract**. | 0 | 13 | 6 | 9 | 0 |
| **dfs rt, ra, rb** | **Double-precision floating subtract**. | 0 | 13 | 6 | 9 | 0 |
| **dftsv rt, ra, u7** | **Double-precision floating test special value** | 0 | N/S | N/S | 9 | 0 |
| **dsync** | **Synchronize data**. | 1 | N/A | 0[5] | N/A | 0[5] |
| **eqv rt, ra, rb** | **Equivalent**. | 0 | 2 | 0 | 2 | 0 |
| **fa rt, ra, rb** | **Floating add**. | 0 | 6 | 0 | 6 | 0 |
| **fceq rt, ra, rb** | **Floating compare equal**. | 0 | 2 | 0 | 2 | 0 |
| **fcgt rt, ra, rb** | **Floating compare greater than**. | 0 | 2 | 0 | 2 | 0 |
| **fcmeq rt, ra, rb** | **Floating compare magnitude equal**. | 0 | 2 | 0 | 2 | 0 |
| **fcmgt rt, ra, rb** | **Floating compare greater than**. | 0 | 2 | 0 | 2 | 0 |
| **fesd rt, ra** | **Floating extend single to double**. | 0 | 13 | 6 | 13 | 6 |
| **fi rt, ra, rb** | **Floating interpolate**. | 0 | 7 | 0 | 7 | 0 |
| **fm rt, ra, rb** | **Floating multiply**. | 0 | 6 | 0 | 6 | 0 |
| **fma rt, ra, rb, rc** | **Floating multiply and add**. | 0 | 6 | 0 | 6 | 0 |
| **fms rt, ra, rb, rc** | **Floating multiply and subtract**. | 0 | 6 | 0 | 6 | 0 |
| **fnms rt, ra, rb, rc** | **Floating negative multiply and subtract**. | 0 | 6 | 0 | 6 | 0 |
| **frds rt, ra** | **Floating round double to single**. | 0 | 13 | 6 | 13 | 6 |
| **frest rt, ra** | **Floating reciprocal estimate**. | 1 | 4 | 0 | 4 | 0 |
| **frsqest rt, ra** | **Floating reciprocal square root estimate**. | 1 | 4 | 0 | 4 | 0 |
| **fs rt, ra, rb** | **Floating subtract**. | 0 | 6 | 0 | 6 | 0 |
| **fscrrd rt** | **Floating-point status control register read**. | 0 | 13 | 6 | 13 | 6 |
| **fscrwr ra**<br>**fscrwr rt, ra** | **Floating-point status control register write**. | 0 | 7 | 0 | 7 | 0 |
| **fsm rt, ra** | **Form select mask for words**. | 1 | 4 | 0 | 4 | 0 |
| **fsmb rt, ra** | **Form select mask for bytes**. | 1 | 4 | 0 | 4 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 5 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **fsmbi rt, u16** | **Form select mask for byte immediate**. | 1 | 4 | 0 | 4 | 0 |
| **fsmh rt, ra** | **Form select mask for halfwords**. | 1 | 4 | 0 | 4 | 0 |
| **gb rt, ra** | **Gather bits from words**. | 1 | 4 | 0 | 4 | 0 |
| **gbb rt, ra** | **Gather bits from bytes**. | 1 | 4 | 0 | 4 | 0 |
| **gbh rt, ra** | **Gather bits from halfwords**. | 1 | 4 | 0 | 4 | 0 |
| **hbr s11, ra** | **Hint for branch (r-form)**. | 1 | 15 | 0 | 15 | 0 |
| **hbra s11, s18** | **Hint for branch (a-form)**. | 1 | 15 | 0 | 15 | 0 |
| **hbrp** | **Hint for branch, prefetch (r-form)**. | 1 | 15 | 0 | 15 | 0 |
| **hbrr s11, s18** | **Hint for branch relative**. | 1 | 15 | 0 | 15 | 0 |
| **heq ra, rb** | **Halt if equal**. | 0 | N/A | 0 | N/A | 0 |
| **heqi ra, s10** | **Halt if equal immediate**. | 0 | N/A | 0 | N/A | 0 |
| **hgt ra, rb** | **Halt if greater than**. | 0 | N/A | 0 | N/A | 0 |
| **hgti ra, s10** | **Halt if greater than immediate**. | 0 | N/A | 0 | N/A | 0 |
| **hlgt ra, rb** | **Halt if logically greater than**. | 0 | N/A | 0 | N/A | 0 |
| **hlgti ra, s10** | **Halt if logically greater than immediate**. | 0 | N/A | 0 | N/A | 0 |
| **iiretd ra** | **Interrupt return, disable**. | 1 | N/A | 0 | N/A | 0 |
| **il rt, s16** | **Immediate load word**. | 0 | 2 | 0 | 2 | 0 |
| **ila rt, u18** | **Immediate load address**. | 0 | 2 | 0 | 2 | 0 |
| **ilh rt, u16** | **Immediate load halfword**. | 0 | 2 | 0 | 2 | 0 |
| **ilhu rt, u16** | **Immediate load halfword upper**. | 0 | 2 | 0 | 2 | 0 |
| **iohl rt, u16** | **Immediate OR halfword lower**. | 0 | 2 | 0 | 2 | 0 |
| **iret ra** | **Interrupt return**. | 1 | N/A | 0 | N/A | 0 |
| **irete ra** | **Interrupt return, enable**. | 1 | N/A | 0 | N/A | 0 |
| **lnop** | **Nop operation (load)**. | 1 | 0 | 0 | 0 | 0 |
| **lqa rt, s18** | **Load quadword (a-form)**. | 1 | 6 | 0 | 6 | 0 |
| **lqd rt, s14(ra)** | **Load quadword (d-form)**. | 1 | 6 | 0 | 6 | 0 |
| **lqr rt, s18** | **Load quadword instruction relative (a-form)**. | 1 | 6 | 0 | 6 | 0 |
| **lqx rt, ra, rb** | **Load quadword (x-form)**. | 1 | 6 | 0 | 6 | 0 |
| **mfspr rt, spr** | **Move from special purpose register**. | 1 | 6 | 0 | 6 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 6 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **mpy rt, ra, rb** | **Multiply**. | 0 | 7 | 0 | 7 | 0 |
| **mpya rt, ra, rb, rc** | **Multiply and add**. | 0 | 7 | 0 | 7 | 0 |
| **mpyh rt, ra, rb** | **Multiply high**. | 0 | 7 | 0 | 7 | 0 |
| **mpyhh rt, ra, rb** | **Multiply high high**. | 0 | 7 | 0 | 7 | 0 |
| **mpyhha rt, ra, rb** | **Multiply high high and add**. | 0 | 7 | 0 | 7 | 0 |
| **mpyhhau rt, ra, rb** | **Multiply high high unsigned and add**. | 0 | 7 | 0 | 7 | 0 |
| **mpyhhu rt, ra, rb** | **Multiply high high unsigned**. | 0 | 7 | 0 | 7 | 0 |
| **mpyi rt, ra, s10** | **Multiply immediate**. | 0 | 7 | 0 | 7 | 0 |
| **mpys rt, ra, rb** | **Multiply and shift right**. | 0 | 7 | 0 | 7 | 0 |
| **mpyu rt, ra, rb** | **Multiply unsigned**. | 0 | 7 | 0 | 7 | 0 |
| **mpyui rt, ra, s10** | **Multiply unsigned immediate**. | 0 | 7 | 0 | 7 | 0 |
| **mtspr spr, ra** | **Move to special purpose register**. | 1 | 6 | 0 | 6 | 0 |
| **nand rt, ra, rb** | **Nand**. | 0 | 2 | 0 | 2 | 0 |
| **nop**<br>**nop rt** | **Nop operation (execute)**. | 0 | 0 | 0 | 0 | 0 |
| **nor rt, ra, rb** | **Nor**. | 0 | 2 | 0 | 2 | 0 |
| **or rt, ra, rb** | **Or**. | 0 | 2 | 0 | 2 | 0 |
| **orbi rt, ra, s10** | **Or byte immediate**. | 0 | 2 | 0 | 2 | 0 |
| **orc rt, ra, rb** | **Or with complement**. | 0 | 2 | 0 | 2 | 0 |
| **orhi rt, ra, s10** | **Or halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **ori rt, ra, s10** | **Or word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **orx rt, ra** | **Or word across.** | 1 | 4 | 0 | 4 | 0 |
| **rchcnt rt, ch** | **Read channel count**. | 1 | 6[2, 3, 4] | 0 | 6[2, 3, 4] | 0 |
| **rdch rt, ch** | **Read channel**. | 1 | 6[2, 3, 4] | 0 | 6[2, 3, 4] | 0 |
| **rot rt, ra, rb** | **Rotate word**. | 0 | 4 | 0 | 4 | 0 |
| **roth rt, ra, rb** | **Rotate halfword**. | 0 | 4 | 0 | 4 | 0 |
| **rothi rt, ra, s7** | **Rotate halfword immediate**. | 0 | 4 | 0 | 4 | 0 |
| **rothm rt, ra, rb** | **Rotate and mask halfword**. | 0 | 4 | 0 | 4 | 0 |
| **rothmi rt, ra, s6** | **Rotate and mask halfword immediate**. | 0 | 4 | 0 | 4 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 7 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **roti rt, ra, s7** | **Rotate word immediate**. | 0 | 4 | 0 | 4 | 0 |
| **rotm rt, ra, rb** | **Rotate and mask word**. | 0 | 4 | 0 | 4 | 0 |
| **rotma rt, ra, rb** | **Rotate and mask algebraic word**. | 0 | 4 | 0 | 4 | 0 |
| **rotmah rt, ra, rb** | **Rotate and mask algebraic halfword**. | 0 | 4 | 0 | 4 | 0 |
| **rotmahi rt, ra, s6** | **Rotate and mask algebraic halfword immediate**. | 0 | 4 | 0 | 4 | 0 |
| **rotmai rt, ra, s7** | **Rotate and mask algebraic word immediate**. | 0 | 4 | 0 | 4 | 0 |
| **rotmi rt, ra, s7** | **Rotate and mask word immediate**. | 0 | 4 | 0 | 4 | 0 |
| **rotqbi rt, ra, rb** | **Rotate quadword by bits**. | 1 | 4 | 0 | 4 | 0 |
| **rotqbii rt, ra, u3** | **Rotate quadword by bits immediate**. | 1 | 4 | 0 | 4 | 0 |
| **rotqby rt, ra, rb** | **Rotate quadword by bytes**. | 1 | 4 | 0 | 4 | 0 |
| **rotqbybi rt, ra, rb** | **Rotate quadword by bytes from bit shift count**. | 1 | 4 | 0 | 4 | 0 |
| **rotqbyi rt, ra, s7** | **Rotate quadword by bytes immediate**. | 1 | 4 | 0 | 4 | 0 |
| **rotqmbi rt, ra, rb** | **Rotate and mask quadword by bits**. | 1 | 4 | 0 | 4 | 0 |
| **rotqmbii rt, ra, s3** | **Rotate and mask quadword by bits immediate**. | 1 | 4 | 0 | 4 | 0 |
| **rotqmby rt, ra, rb** | **Rotate and mask quadword by bytes**. | 1 | 4 | 0 | 4 | 0 |
| **rotqmbybi rt, ra, rb** | **Rotate and mask quadword by bytes from bit shift count.** | 1 | 4 | 0 | 4 | 0 |
| **rotqmbyi rt, ra, s6** | **Rotate and mask quadword by bytes immediate**. | 1 | 4 | 0 | 4 | 0 |
| **selb rt, ra, rb, rc** | **Select bits**. | 0 | 2 | 0 | 2 | 0 |
| **sf rt, ra, rb** | **Subtract from word**. | 0 | 2 | 0 | 2 | 0 |
| **sfh rt, ra, rb** | **Subtract from halfword**. | 0 | 2 | 0 | 2 | 0 |
| **sfhi rt, ra, s10** | **Subtract from halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **sfi rt, ra, s10** | **Subtract from word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **sfx rt, ra, rb** | **Subtract from word extended**. | 0 | 2 | 0 | 2 | 0 |
| **shl rt, ra, rb** | **Shift left word**. | 0 | 4 | 0 | 4 | 0 |
| **shlh rt, ra, rb** | **Shift left halfword**. | 0 | 4 | 0 | 4 | 0 |
| **shlhi rt, ra, u5** | **Shift left halfword immediate**. | 0 | 4 | 0 | 4 | 0 |
| **shli rt, ra, u6** | **Shift left word immediate**. | 0 | 4 | 0 | 4 | 0 |
| **shlqbi rt, ra, rb** | **Shift left quadword by bits**. | 1 | 4 | 0 | 4 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

*Table B-1. SPU Instructions* (Sheet 8 of 8)

| Instruction | Description | Pipeline | Cell/B.E. Processor | | PowerXCell 8i Processor | |
|---|---|---|---|---|---|---|
| | | | Latency[1] (cycles) | Stalls (cycles) | Latency[1] (cycles) | Stalls (cycles) |
| **shlqbii rt, ra, u3** | **Shift left quadword by bits immediate**. | 1 | 4 | 0 | 4 | 0 |
| **shlqby rt, ra, rb** | **Shift left quadword by bytes**. | 1 | 4 | 0 | 4 | 0 |
| **shlqbybi rt, ra, rb** | **Shift left quadword by bytes from bit shift count**. | 1 | 4 | 0 | 4 | 0 |
| **shlqbyi rt, ra, u5** | **Shift left quadword by bytes immediate**. | 1 | 4 | 0 | 4 | 0 |
| **shufb rt, ra, rb, rc** | **Shuffle bytes**. | 1 | 4 | 0 | 4 | 0 |
| **stop u14** | **Stop and signal**. | 1 | 4 | 0 | 4 | 0 |
| **stopd ra, rb, rc** | **Stop and signal with dependencies**. | 1 | 4 | 0 | 4 | 0 |
| **stqa rt, s18** | **Store quadword (a-form)**. | 1 | 6 | 0 | 6 | 0 |
| **stqd rt, s14(ra)** | **Store quadword (d-form)**. | 1 | 6 | 0 | 6 | 0 |
| **stqr rt, s18** | **Store quadword instruction relative (a-form)**. | 1 | 6 | 0 | 6 | 0 |
| **stqx rt, ra, rb** | **Store quadword (x-form)**. | 1 | 6 | 0 | 6 | 0 |
| **sumb rt, ra, rb** | **Sum bytes into halfword**. | 0 | 4 | 0 | 4 | 0 |
| **sync** | **Synchronize**. | 1 | N/A | 0[5] | N/A | 0[5] |
| **syncc** | **Synchronize channel**. | 1 | N/A | 0[5] | N/A | 0[5] |
| **wrch ch, ra** | **Write channel**. | 1 | N/A[2, 4] | 0 | N/A[2, 4] | 0 |
| **xor rt, ra, rb** | **Xor**. | 0 | 2 | 0 | 2 | 0 |
| **xorbi rt, ra, s10** | **Exclusive or byte immediate**. | 0 | 2 | 0 | 2 | 0 |
| **xorhi rt, ra, s10** | **Exclusive or halfword immediate**. | 0 | 2 | 0 | 2 | 0 |
| **xori rt, ra, s10** | **Exclusive or word immediate**. | 0 | 2 | 0 | 2 | 0 |
| **xsbh rt, ra** | **Extend sign byte to halfword**. | 0 | 2 | 0 | 2 | 0 |
| **xshw rt, ra** | **Extend sign halfword to word**. | 0 | 2 | 0 | 2 | 0 |
| **xswd rt, ra** | **Extend sign word to doubleword**. | 0 | 2 | 0 | 2 | 0 |

1. The latency cycles include all stall cycles.
2. A **rdch** or **wrch** instruction for a blocking channel is not issued unless the channel count for the blocking channel is greater than zero.
3. This is the latency for a nonblocking case.
4. If successive **rdch**, **wrch**, or **rchcnt** instructions are issued within five SPU cycles to the same channel (or to the SPU_RdEventStat Channel, the SPU_WrEventMask Channel, or the SPU_WrEventAck Channel), the second channel instruction is retried (causing a flush).
5. This instruction causes a pipeline flush, so in some cycles the pipeline makes no progress.

## B.1.3 Fetch and Issue Rules

### B.1.3.1 *Fetch*

Instruction fetches load 32 instructions per LS request. Because the LS is single-ported and load and store instruction frequency is likely to drive LS occupancy to high levels, it is important that DMA and instruction-fetch activity transfer as much useful data as possible in each LS request.

There are three types of instruction fetches: flush-initiated fetches, hint-for branch fetches, and inline (sequential) prefetches. Of these, application software can directly influence hint-for branch fetches and it has some indirect control over inline prefetches.

### B.1.3.2 *Issue*

The SPU issues all instructions in program order according to the pipeline assignment. Each instruction is part of a doubleword-aligned instruction pair called a *fetch group*. A fetch group can have one or two valid instructions, but it must be aligned to doubleword boundaries. This means that the first instruction in the fetch group is from an even word address, and the second instruction from an odd word address. The SPU processes fetch groups one at a time, continuing to the next fetch group when the current instruction group becomes empty. An instruction becomes issueable when register dependencies are satisfied and there is no structural hazard (resource conflict) with prior instructions or direct memory access (DMA) or error-correcting code (ECC) activity.

Dual-issue occurs when a fetch group has two issueable instructions in which the first instruction can be executed on the even pipeline and the second instruction can be executed on the odd pipeline. If a fetch group cannot be dual-issued, but the first instruction can be issued, the first instruction is issued to the proper execution pipeline and the second instruction is held until it can be issued. A new fetch group is loaded after both instructions of the current fetch group are issued.

*Table B-2* shows how a short program can be issued by the SPU in the absence of resource conflicts and register dependencies. Even-pipeline instructions are shaded, odd-pipeline instructions are not shaded. The program features four cases of pipeline assignment versus instruction alignment. The eight instructions are issued in seven cycles. Dual-issue occurs only for the first fetch group. The first fetch group consists of an instruction executed by the even pipeline in address x‘0000’ and an instruction executed by the odd pipeline at address x‘0004’.

*Table B-2. Instruction Issue Example*

| Address | Fetch Group in Memory | |
|---|---|---|
| | Even Address | Odd Address |
| x‘0000’ | 1: Even | 2: Odd |
| x‘0008’ | 3: Even | 4: Even |
| x‘0010’ | 5: Odd | 6: Odd |
| x‘0018’ | 7: Odd | 8: Even |

dual issue
(no register dependency)

| Instruction Issue | | Clock Cycles |
|---|---|---|
| Even Pipeline | Odd Pipeline | |
| 1: Even | 2: Odd | 1 |
| 3: Even | Penalty | 2 |
| 4: Even | | 3 |
| | 5: Odd | 4 |
| | 6: Odd | 5 |
| | 7: Odd | 6 |
| 8: Even | | 7 |

Instruction alignment can be optimized to minimize execution time by inserting no operation (NOP) and load NOP (LNOP) instructions. For example, *Table B-3* shows how another short program can be scheduled on an SPU. These eight instructions issue in six cycles. This same program can be issued in five cycles if the instruction alignment is improved.

*Table B-3. Example of Instruction Issue without NOP and LNOP*

| Address | Fetch Group in Memory | | | Instruction Issue | | Clock Cycles |
|---|---|---|---|---|---|---|
| | Even Address | Odd Address | | Even Pipe | Odd Pipe | |
| x'0000' | 1: Even | 2: Odd | dual issue (no register dependency) | 1: Even | 2: Odd | 1 |
| x'0008' | 3: Odd | 4: Even | | — | 3: Odd | 2 |
| x'0010' | 5: Odd | 6: Even | | 4: Even | — | 3 |
| x'0018' | 7: Even | 8: Odd | | — | 5: Odd | 4 |
| x'0020' | | | | 6: Even | — | 5 |
| | | | dual issue (no register dependency) | 7: Even | 8: Odd | 6 |

*Table B-4* on page 781 shows that a NOP added between instructions 2 and 3 maintains dual-issue through the third fetch group, and an added LNOP continues dual-issue for the remainder of the program.

*Table B-4. Example of Instruction Issue Using NOP and LNOP*

| Address | Fetch Group in Memory | | | Instruction Issue | | Clock Cycles |
|---|---|---|---|---|---|---|
| | Even Address | Odd Address | | Even Pipe | Odd Pipe | |
| x'0000' | 1: Even | 2: Odd | dual issue (no register dependency) | 1: Even | 2: Odd | 1 |
| x'0008' | NOP | 3: Odd | dual issue (no register dependency) | NOP | 3: Odd | 2 |
| x'0010' | 4: Even | 5: Odd | dual issue (no register dependency) | 4: Even | 5: Odd | 3 |
| x'0018' | 6: Even | LNOP | dual issue (no register dependency) | 6: Even | LNOP | 4 |
| x'0020' | 7: Even | 8: Odd | dual issue (no register dependency) | 7: Even | 8: Odd | 5 |
| | | | | — | — | 6 |

*Table B-5 Cell/B.E. SPU Execution Pipelines and Result Latency* on page 782 and *Table B-6 PowerXCell 8i SPU Execution Pipelines and Result Latency* on page 783 show how the various operations on the even and odd pipelines are related. Instructions are classified in this table as:

- LS: Load and Store
- HB: Branch Hints
- BR: Branch Resolution
- CH: Channel Interface, Special Purpose Registers
- SH: Shuffle
- LNOP: No Operation (load)
- NOP: No Operation (execute)
- SP: Single-Precision Floating-Point
- DP: Double-Precision Floating-Point
- FI: Floating-Point Integer
- BO: Byte Operations
- FX: Simple Fixed-Point
- WS: Word Rotate and Shift

*Table B-5. Cell/B.E. SPU Execution Pipelines and Result Latency*

| Pipeline | Pipeline Stage/Latency (cycles) | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|
| | M/1 | N/2 | O/3 | P/4 | Q/5 | R/6 | S/7 | T | U | V |
| Even | | FX | | BO, WS | | SP | DP, FI | Pre-writeback | Register file write | Read bypass |
| Odd | | | | BR, SH | | LS, CH | | Pre-writeback | Register file write | Read bypass |

Each pipeline stage has a unique alphabetical identifier such that a later stage has a letter that occurs later in the alphabet. Execution begins in stage M, where inputs to the execution units are launched and latched. No unit produces a result in the first cycle, but FX-class instruction results become available for use by subsequent instructions in the second cycle. When a result becomes available in a particular cycle, it is available for use by any subsequent instruction executed by any unit. Although shorter-latency instructions make their results available for forwarding earlier, their results on a Cell/B.E. processor are staged out to latency 8 (stage T) before the results are sent to the register file. The longest latency of a pipelined result is seven cycles, which allows seven cycles of result-equalization to commit register-file results in program order.

Register-file reads and writes are both 2-cycle operations. Result data is sent to the register file while the write addresses are decoded during the pre-writeback stage. In the next cycle, the register file is written. One final stage of forwarding, read bypass, is necessary before the written result is available from the register file.

Double-precision instructions are performed as two double-precision operations in 2-way SIMD fashion. However, the Cell/B.E. SPU is capable of performing only one double-precision operation per cycle. Thus, the Cell/B.E. SPU executes double-precision instructions by breaking up the SIMD operands and executing the two operations in consecutive instruction slots in the pipeline.

Although double-precision instructions have 13-clock-cycle latencies, only the final seven cycles are pipelined. No other instructions are dual-issued with double-precision instructions, and no instructions of any kind are issued for six cycles after a double-precision instruction is issued.

*Table B-6. PowerXCell 8i SPU Execution Pipelines and Result Latency*

| Pipeline | Pipeline Stage Name/Latency Cycles | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|
| | M/1 | N/2 | O/3 | P/4 | Q/5 | R/6 | S/7 | T/8 | U/9 | V | W | X |
| Even | | FX | | BO, WS | | SP | FI | | DP | Pre-write back | Register file write | Read bypass |
| Odd | | | | BR, SH | | LS, CH | | | | Pre-write back | Register file write | Read bypass |

For the PowerXCell 8i processor, the results are staged out to latency 10 (stage V) before the results are sent to the register file. The longest latency of a pipelined result is nine cycles. Double-precision instructions are fully pipelined and dual-issued.

### B.1.4 Inline Prefetch and Instruction Runout

For inline (sequential) prefetch, the prefetcher uses empty slots on the LS schedule to perform the fetch, rather than using a specific instruction as in a branch hint. It is possible to prevent prefetch with long sequences of load and store instructions. The prefetcher starts looking for slots when one of the inline prefetch buffer pairs becomes empty. At this point, there should be 16 instructions in the predicted-path buffer and another 32 instructions in the other prefetch buffer. If the SPU sustains dual-issue, there should be 24 cycles before the prefetched text is needed for inline speculation.

Prefetch requests require 15 cycles to load the instruction line buffer. This means that the prefetcher has nine cycles in which it must succeed in securing a slot. If the LS is busy during this interval and the prefetch is not started, instruction runout might occur (that is, there might be no instructions in the pipeline, and the SPU becomes idle). The minimum runout delay to reset the fetch state machine is three cycles if the prefetch was late by only one or two cycles. Prefetches late enough to cause runout for more than three cycles incrementally delay the issue of the next instruction by one cycle for every cycle the prefetch is delayed. As the pipeline drains, loads and stores cease, the prefetcher can schedule the prefetch, and execution resumes.

Long sequences of instructions that access the LS every cycle can cause instruction runout. Instruction runout can be prevented by breaking up these sequences with instructions that do not access the LS. The **hbrp** instruction allows programs to ensure that the cycles needed for instruction fetch are not consumed by DMA transfers. In principle, any instruction that does not use the LS, and that does not dual-issue with an instruction that uses the LS, is sufficient.

The following rules apply:

- The empty slot in the LS schedule may be adjacent to a slot used by a line-read DMA transfer or branch hint, and this prevents the prefetch from using this empty slot.
- The issue-control unit can schedule non-LS instructions during a DMA transfer, but this causes the slot that the schedule intended for prefetch to be lost.

The **hbrp** instruction has a P feature bit that alters the behavior of the **hbr** instruction. When this bit is set, **hbrp** arbitrates for the LS as though it were a hint. However, after it is scheduled, it becomes a **nop** instruction. This leaves a hole wide enough for an instruction prefetch. For more information, see:

- *SPU Assembly Language Specification*
- *SPU Application Binary Interface Specification*

## B.2 C/C++ Language Extensions (Intrinsics) for SPU Instructions

An overview of the SPU intrinsics is given in *Section 3.4* on page 77. The sections that follow summarize key points of the intrinsics. For a complete description, see the *C/C++ Language Extensions for Cell Broadband Engine Architecture* document.

### B.2.1 Vector Data Types

The C/C++ language extensions define eleven vector data types, as shown in *Table B-7*. These data types are all 128 bits long and contain from 1 to 16 elements per vector. The `qword` data type is a special 16-byte quadword that is only used as input or output for specific intrinsics (see *Section B.2.3* on page 787).

*Table B-7. Vector Data Types*

| Vector Data Type | Content |
|---|---|
| vector unsigned char | 16 8-bit unsigned chars |
| vector signed char | 16 8-bit signed chars |
| vector unsigned short | 8 16-bit unsigned halfwords |
| vector signed short | 8 16-bit signed halfwords |
| vector unsigned int | 4 32-bit unsigned words |
| vector signed int | 4 32-bit signed words |
| vector unsigned long long | 2 64-bit unsigned doublewords |
| vector signed long long | 2 64-bit signed doublewords |
| vector float | 4 32-bit single-precision floating-point numbers |
| vector double | 2 64-bit double-precision floating-point numbers |
| qword | quadword (16-byte) |

To improve code portability, the `spu_intrinsics.h` header file provides `typedefs` for the vector data types, as shown in *Table B-8*. The `typedefs` might be useful when porting code from the PPE to an SPE.

*Table B-8. Programmer Typedefs for Vector Data Types* (Sheet 1 of 2)

| Vector Data Type | Typedef |
|---|---|
| vector unsigned char | vec_uchar16 |
| vector signed char | vec_char16 |
| vector unsigned short | vec_ushort8 |

*Table B-8. Programmer Typedefs for Vector Data Types* (Sheet 2 of 2)

| Vector Data Type | Typedef |
|---|---|
| vector signed short | vec_short8 |
| vector unsigned int | vec_uint4 |
| vector signed int | vec_int4 |
| vector unsigned long long | vec_ullong2 |
| vector signed long long | vec_llong2 |
| vector float | vec_float4 |
| vector double | vec_double2 |

*Table B-9* shows the size and alignment of the data types.

*Table B-9. Data Type Alignments*

| Data Type | Size | Alignment |
|---|---|---|
| char | 1 | byte |
| short | 2 | halfword |
| int | 4 | word |
| long | 4 | word/doubleword |
| long long | 8 | doubleword |
| float | 4 | word |
| double | 8 | doubleword |
| pointer | 4 | word |
| vector | 16 | quadword |

The `aligned` attribute is used to align data on particular boundaries, as in:

```
float factor _attribute_((aligned (16)));
```

The variable `factor` is aligned on a quadword boundary.

The `__align_hint` intrinsic helps in two ways:

- Enables compiler analysis and optimization for pointers.
- Provides compilers with extra information needed for auto-SIMDization.

The `__align_hint(ptr, base, offset)` intrinsic tells the compiler that the pointer `ptr` points to data with a base alignment of `base` and with the given `offset`. The `base` must be a power of 2. A `base` of zero implies the pointer has no known alignment. The `offset` must be less than the `base`, or zero.

Use `__align_hint()` with care. If it is used with pointers that are not aligned, data can end up straddling quadword boundaries. If you specify alignment incorrectly, your program might be compiled incorrectly (that is, not as you expect) and you might get incorrect results.

One last caveat: although a compiler compliant with the *C/C++ Language Extensions for Cell Broadband Engine Architecture* specification is required to provide the `__align_hint` intrinsic, the compiler is allowed to ignore these hints.

#### B.2.1.1 *Vector Element Ordering*

The SPE intrinsics support big-endian data ordering of elements within a vector data type, as shown in *Figure B-10*.

*Table B-10. Element Ordering for Vector Types*

MSB LSB

| byte 0 | byte 1 | byte 2 | byte 3 | byte 4 | byte 5 | byte 6 | byte 7 | byte 8 | byte 9 | byte 10 | byte 11 | byte 12 | byte 13 | byte 14 | byte 15 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| doubleword 0 | | | | | | | | doubleword 1 | | | | | | | |
| word 0 | | | | word 1 | | | | word 2 | | | | word 3 | | | |
| halfword 0 | | halfword 1 | | halfword 2 | | halfword 3 | | halfword 4 | | halfword 5 | | halfword 6 | | halfword 7 | |
| char 0 | char 1 | char 2 | char 3 | char 4 | char 5 | char 6 | char 7 | char 8 | char 9 | char 10 | char 11 | char 12 | char 13 | char 14 | char 15 |

### B.2.2 Vector Literals

A vector literal is written as a parenthesized vector type followed by a curly-braced set of constant expressions, as shown in *Table B-11*. The elements of the vector are initialized to the corresponding expression. Elements for which no expressions are specified default to 0. Vector literals may be used either in initialization statements or as constants in executable statements.

*Table B-11. Vector-Literal Format*

| Notation | Represents |
|---|---|
| (vector unsigned char) *{unsigned int, ...}* | A set of 16 unsigned 8-bit quantities. |
| (vector signed char) *{signed int, ...}* | A set of 16 signed 8-bit quantities. |
| (vector unsigned short) *{unsigned short, ...}* | A set of 8 unsigned 16-bit quantities. |
| (vector signed short) *{signed short, ...}* | A set of 8 signed 16-bit quantities. |
| (vector unsigned int) *{unsigned int, ...}* | A set of 4 unsigned 32-bit quantities. |
| (vector signed int) *{signed int, ...}* | A set of 4 signed 32-bit quantities. |
| (vector unsigned long long) *{unsigned long long, ...}* | A set of 2 unsigned 64-bit quantities. |
| (vector signed long long) *{signed long long, ...}* | A set of 2 signed 64-bit quantities. |
| (vector float) *{float, ...}* | A set of 4 32-bit floating-point quantities. |
| (vector double) *{double, ...}* | A set of 2 64-bit floating-point quantities. |

An alternate format may also be supported, as shown in *Table B-12*. This form consists of a parenthesized vector type followed by a parenthesized set of constant expressions.

*Table B-12. Alternate Vector-Literal Format* (Sheet 1 of 2)

| Notation | Description |
|---|---|
| (vector unsigned char)(unsigned int) | A set of 16 unsigned 8-bit quantities that all have the value specified by the integer. |
| (vector unsigned char)(unsigned int, ..., unsigned int) | A set of 16 unsigned 8-bit quantities specified by the 16 integers. |
| (vector signed char)(signed int) | A set of 16 signed 8-bit quantities that all have the value specified by the integer. |

*Table B-12. Alternate Vector-Literal Format* (Sheet 2 of 2)

| Notation | Description |
|---|---|
| (vector signed char)(signed int, ..., signed int) | A set of 16 signed 8-bit quantities specified by the 16 integers. |
| (vector unsigned short)(unsigned int) | A set of eight unsigned 16-bit quantities that all have the value specified by the integer. |
| (vector unsigned short)(unsigned int, ..., unsigned int) | A set of eight unsigned 16-bit quantities specified by the 16 integers. |
| (vector signed short)(signed int) | A set of eight signed 16-bit quantities that all have the value specified by the integer. |
| (vector signed short)(signed int, ..., signed int) | A set of eight signed 16-bit quantities specified by the 16 integers. |
| (vector unsigned int)(unsigned int) | A set of four unsigned 32-bit quantities that all have the value specified by the integer. |
| (vector unsigned int)(unsigned int, ..., unsigned int) | A set of four unsigned 32-bit quantities specified by the 16 integers. |
| (vector signed int)(signed int) | A set of four signed 32-bit quantities that all have the value specified by the integer. |
| (vector signed int)(signed int, ..., signed int) | A set of four signed 32-bit quantities specified by the 16 integers. |
| (vector unsigned long long)(unsigned long long) | A set of two unsigned 64-bit quantities that all have the value specified by the long long integer. |
| (vector unsigned long long)(unsigned long long, unsigned long long) | A set of two unsigned 64-bit quantities specified by the two long long integers. |
| (vector signed long long)(signed long long) | A set of two signed 64-bit quantities that all have the value specified by the long long integer. |
| (vector signed long long)(signed long long, signed long long) | A set of two signed 64-bit quantities specified by the two long long integers. |
| (vector float)(float) | A set of four 32-bit floating-point quantities that all have the value specified by the float. |
| (vector float)(float, float, float, float) | A set of four 32-bit floating-point quantities specified by the 4 floats. |
| (vector double)(double) | A set of two 64-bit double-precision quantities that all have the value specified by the double. |
| (vector double)(double, double) | A set of two 64-bit quantities specified by the 2 doubles. |

### B.2.3 Intrinsics

*Table B-13* on page 788 lists the generic and composite intrinsics. Generic intrinsics map to one or more assembly-language instructions, as a function of the type of its input parameters. *Built-ins* are a subset of generic intrinsics that map to more than one SPU instruction. All of the generic intrinsics and built-ins are prefixed by the string, `spu_`. For example, the intrinsic that implements the stop assembly instruction is named `spu_stop`. Many generic intrinsics accept scalars as one of their operands. These correspond to intrinsics that map to instructions with immediate values.

Generic intrinsics are provided for all SPU instructions, *except the following*:

- Branch

- Branch hint (see *Section 24.3* on page 699)
- Interrupt return
- Generate control for insertion (used for scalar stores)
- Constant formation
- Nop
- Memory load and store
- Stop and signal with dependencies (**stopd**)

Composite intrinsics are constructed from a sequence of specific or generic intrinsics. Specific intrinsics are not included in this summary, because they are rarely used.

*Table B-13. SPU Intrinsics* (Sheet 1 of 3)

| Intrinsic | Description |
|---|---|
| **Constant Formation** | |
| d = spu_splats(a) | Replicate scalar a into all elements of vector d |
| **Conversion** | |
| d = spu_convtf(a, scale) | Convert integer vector to float vector |
| d = spu_convts(a, scale) | Convert float vector to signed int vector |
| d = spu_convtu(a, scale) | Convert float vector to unsigned float vector |
| d = spu_extend(a) | Sign extend vector |
| d = spu_rountf(a) | Round double vector to float vector |
| **Arithmetic** | |
| d = spu_add(a, b) | Vector add |
| d = spu_addx(a, b, c) | Vector add extended |
| d = spu_genb(a, b) | Vector generate borrow |
| d = spu_genbx(a, b, c) | Vector generate borrow extended |
| d = spu_genc(a, b) | Vector generate carry |
| d = spu_gencx(a, b, c) | Vector generate carry extended |
| d = spu_madd(a, b, c) | Vector multiply and add |
| d = spu_mhhadd(a, b, c) | Vector multiply high high and add |
| d = spu_msub(a, b, c) | Vector multiply and subtract |
| d = spu_mul(a, b) | Vector multiply |
| d = spu_mulh(a, b) | Vector multiply high |
| d = spu_mulhh(a, b) | Vector multiply high high |
| d = spu_mulo(a, b) | Vector multiply odd |
| d = spu_mulsr(a, b) | Vector multiply and shift right |
| d = spu_nmadd(a, b, c) | Negative vector multiply and add |
| d = spu_nmsub(a, b, c) | Negative vector multiply and subtract |
| d = spu_re(a) | Vector floating-point reciprocal estimate |

*Table B-13. SPU Intrinsics* (Sheet 2 of 3)

| Intrinsic | Description |
|---|---|
| d = spu_rsqrte(a) | Vector floating-point reciprocal square root estimate |
| d = spu_sub(a, b) | Vector subtract |
| d = spu_subx(a, b, c) | Vector subtract extended |
| **Byte Operation** | |
| d = spu_absd(a, b) | Vector absolute difference |
| d = spu_avg(a, b) | Vector average |
| d = spu_sumb(a, b) | Vector sum bytes into shorts |
| **Compare, Branch, and Halt** | |
| d = spu_bisled(func) | Branch indirect and set link if external data |
| d = spu_cmpabseq(a, b) | Vector compare absolute equal |
| d = spu_cmpabsgt(a, b) | Vector compare absolute greater than |
| d = spu_cmpeq(a, b) | Vector compare equal |
| d = spu_cmpgt(a, b) | Vector compare greater than |
| (void) spu_hcmpeq(a, b) | Halt if compare equal |
| (void) spu_hcmpgt(a, b) | Halt if compare greater than |
| d = spu_testsv(a, values) | Vector test special value |
| **Bit and Mask** | |
| d = spu_cntb(a) | Vector count ones for bytes |
| d = spu_cntlz(a) | Vector count leading zeros |
| d = spu_gather(a) | Gather bits from elements |
| d = spu_maskb(a) | Form select byte mask |
| d = spu_maskh(a) | Form select halfword mask |
| d = spu_maskw(a) | Form select word mask |
| d = spu_sel(a, b, pattern) | Select bits |
| d = spu_shuffle(a, b, pattern) | Shuffle bytes of a Vector |
| **Logical** | |
| d = spu_and(a, b) | Vector bit-wise AND |
| d = spu_andc(a, b) | Vector bit-wise AND with complement |
| d = spu_eqv(a, b) | Vector bit-wise equivalent |
| d = spu_nand(a, b) | Vector bit-wise complement of AND |
| d = spu_nor(a, b) | Vector bit-wise complement of OR |
| d = spu_or(a, b) | Vector bit-wise OR |
| d = spu_orc(a, b) | Vector bit-wise OR with complement |
| d = spu_orx(a) | Bit-wise OR word elements |
| d = spu_xor(a, b) | Vector bit-wise exclusive OR |
| **Rotate** | |
| d = spu_rl(a, count) | Element-wise bit rotate left |

*Table B-13. SPU Intrinsics* (Sheet 3 of 3)

| Intrinsic | Description |
|---|---|
| d = spu_rlmask(a, count) | Element-wise bit rotate left and mask |
| d = spu_rlmaska(a, count) | Element-wise bit algebraic rotate and mask |
| d = spu_rlmaskqw(a, count) | Bit rotate and mask quadword |
| d = spu_rlmaskqwbyte(a, count) | Byte rotate and mask quadword |
| d = spu_rlmaskqwbytebc(a, count) | Byte rotate and mask quadword using bit rotate count |
| d = spu_rlqw(a, count) | Bit rotate quadword left |
| d = spu_rlqwbyte(a, count) | Byte rotate quadword left |
| d = spu_rlqwbytebc(a, count) | Byte rotate quadword left using bit rotate count |
| **Shift** | |
| d = spu_sl(a, count) | Element-wise bit shift left |
| d = spu_slqw(a, count) | Bit shift quadword left |
| d = spu_slqwbyte(a, count) | Byte shift quadword left |
| d = spu_slqwbytebc(a, count) | Byte shift quadword left using bit shift count |
| **Control** | |
| (void) spu_idisable() | Disable interrupts |
| (void) spu_ienable() | Enable interrupts |
| (void) spu_mffpscr() | Move from floating-point status and control register |
| (void) spu_mfspr(register) | Move from special purpose register |
| (void) spu_mtfpscr(a) | Move to floating-point status and control register |
| (void) spu_mtspr(register, a) | Move to special purpose register |
| (void) spu_dsync() | Synchronize data |
| (void) spu_stop(type) | Stop and signal |
| (void) spu_sync() | Synchronize |
| **Scalar** | |
| d = spu_extract(a, element) | Extract vector element from vector |
| d = spu_insert(a, b, element) | Insert scalar into specified Vector element |
| d = spu_promote(a, element) | Promote scalar to vector |
| **Channel Control** | |
| d = spu_readch(channel) | Read word channel |
| d = spu_readchqw(channel) | Read quadword channel |
| d = spu_readchcnt(channel) | Read channel count |
| (void) spu_writech(channel, a) | Write word channel |
| (void) spu_writechqw(channel, a) | Write quadword channel |
| **Composite Intrinsics** | |
| spu_mfcdma32(ls, ea, size, tagid, cmd) | Initiate DMA to or from 32-bit effective address |
| spu_mfcdma64(ls, eahi, ealow, size, tagid, cmd) | Initiate DMA to or from 64-bit effective address |
| spu_mfcstat(type) | Read memory flow controller (MFC) tag status |

For further information about the SPU intrinsics, see the *C/C++ Language Extensions for Cell Broadband Engine Architecture* document.

### B.2.4 Inline Assembly

Occasionally, the C/C++ language constructs and intrinsic might not be sufficient to achieve a required low-level programming result. In these situations, inline assembly instructions can be used. The inline assembly syntax must match the AT&T assembly syntax implemented by gcc. The `.balignl` directive may be used within the inline assembly to ensure the known alignment that is needed to achieve effective dual-issue by the hardware.

### B.2.5 Compiler Directives

Like compiler intrinsics, a compiler might provide directives that are very useful for SPE programming. For example, the `restrict` qualifier is well-known in many C/C++ implementations, and it is part of the *C/C++ Language Extensions for Cell Broadband Engine Architecture*. When the `restrict` keyword is used to qualify a pointer, it specifies that all accesses to the object pointed to are done through the pointer. For example:

```
void *memcpy(void * restrict s1, void * restrict s2, size_t n);
```

By specifying `s1` and `s2` as pointers that are restricted, the programmer is specifying that the source and destination objects (for the memory copy) do not overlap.

Another useful directive is `__builtin_expect`. Because branch mispredicts are relatively expensive, `__builtin_expect` provides a way for the programmer to direct branch prediction. This example:

```
int __builtin_expect(int exp, int value)
```

returns the result of evaluating `exp`, and means that the programmer expects `exp` to equal `value`. The `value` can be a constant for compile-time prediction, or a variable used for runtime prediction.

Two more useful directives are the `aligned` attribute, and the `_align_hint` directive. The `aligned` attribute is used to ensure proper DMA alignment, for efficient data transfer. The syntax is the same as in many implementations of gcc:

```
float factor __attribute__((aligned (16)); //aligns “factor” to a quadword
```

The `_align_hint` directive helps compilers auto-vectorize. Although it looks like an intrinsic, it is more properly described as a compiler directive, because no code is generated as a result of using the directive. The example

```
_align_hint(ptr, base, offset)
```

can inform the compiler that the pointer, `ptr`, points to data with a base alignment of `base`, with a byte offset from the base alignment of `offset`. The base alignment must be a power of two. Giving 0 as the base alignment implies that the pointer has no known alignment. The offset must be less than the base, or, zero. The `_align_hint` directive should not be used with pointers that are not naturally aligned.

# Appendix C. Performance Monitor Signals

The performance-monitoring facilities can monitor many types of events (also called signals) from all major units on the Cell Broadband Engine Architecture (CBEA) processors[1]. For an introduction to the performance monitor signals, see *Section 16 Performance Monitoring* on page 443.

## C.1 Selecting Performance Monitor Signals on the Debug Bus

The performance monitor signals are routed on a 128-bit-wide debug bus to the performance monitor registers in the pervasive logic. To enable the signals from an island to reach the debug bus, one or more privilege 1 (hypervisor mode) implementation-specific registers within the unit needs to be set up correctly. By default, a value of zero in the affected bit positions in these registers disables the unit from routing its signals to the debug bus. Any unit that needs to be set up for signal routing must have the appropriate registers explicitly initialized by writing the nonzero values as specified in this document. Typically, the original values are written back to the affected registers after the performance monitoring of the specific island is complete.

The performance monitor can be used to count the number of occurrences of an event (event counting), or the number of cycles that a signal is active or inactive (cycle counting), or in some cases, both. Event counting is performed using the rising edge of a signal, and cycle counting is done using the positive-high level of the signal. Note that the cycles refer to the clock rate at which the island operates. The clock rate is either the core clock rate (NClk) or half of the core clock rate (NClk/2).

In the signal group tables that follow, each signal is designated as a type B, C, E, V, or S. These types define how the signal needs to be interpreted, and are here described:

- **B**—a signal that counts both the number of occurrences and their cumulative length. This signal can present either a single-cycle event or a multicycle event. For this signal type, the rising or falling edge of the signal indicates an occurrence of the event, and the level signal denotes the length of the occurrence. Two successive occurrences on this signal type are separated by at least one cycle.
- **C**—a signal that counts only the cumulative length of all occurrences. A level detector is used to count this signal type. Two successive occurrences of this signal type can occur in back-to-back cycles.
- **E**—a signal that counts only the occurrences. An edge detector is used to count the number of times this event occurs. Two successive events of this type are separated by at least one idle cycle.
- **V**—a signal that lasts for only one cycle even though it represents multicycle events. This signal type is useful for counting the number of occurrences of a signal, but not their duration.
- **S**—a signal that counts single-cycle event occurrences. Because of the single-cycle nature, the length and the number of occurrences are synonymous. Therefore, counting cycles is sufficient to determine the number of occurrences and the length.

The *count_cycles* field of the performance monitor *PMx_control* registers must be explicitly programmed to count either cycles or edges or both, as shown in *Table C-1* on page 794.

1. The Cell Broadband Engine (Cell/B.E.) and PowerXCell 8i processors are collectively called CBEA processors.

*Table C-1. Count_cycles Field of the PMx_control Register*

| Tag | Cycle Counting | Event Counting |
|---|---|---|
| B | Cycle | Edge |
| C | Cycle | N/A |
| E | N/A | Edge |
| V | N/A | Cycle |
| S | Cycle | Cycle |

All memory-mapped input and output (MMIO) addresses referred to in this document are relative to the BE_MMIO_Base address. Some of the registers must be altered with a read-modify-write operation. These are designated “rmw” in the register descriptions.

***Note:*** *The registers listed in this appendix are solely for the setup and use of the performance monitor facility. No further description or details of these registers are provided.*

The CBEA processors consist of several units, also referred to as islands. These include the pervasive unit, the PowerPC Processor Element (PPE), the Synergistic Processor Elements (SPEs), the element interconnect bus (EIB), the Cell Broadband Engine interface (BEI), and the memory interface controller (MIC). The PPE consists of two subunits: the PowerPC processor unit (PPU), and the PowerPC storage subsystem (PPSS). Similarly, each SPE consists of a synergistic processor unit (SPU) and a memory flow controller (MFC). The EIB also contains the token manager. Finally, the BEI consists of the interface controller and the input and output controller. All islands clock at NClk/2 except for the PPU, SPU, and parts of the PPSS, all of which clock at NClk.

Each island can logically OR their performance monitor signals onto the debug bus, allowing the monitoring of multiple islands at the same time. The 128-bit-wide debug bus is logically partitioned into four 32-bit-wide lanes. An island might be capable of routing the signals on one or more of the debug bus lanes. The lane on which the signals are routed is selected by initializing the special purpose registers (SPRs) or the MMIO registers in that unit. Care must be taken *not* to allow more than one island to use the *same* lane on the debug bus at the same time.

*Table C-2* shows the mapping between the words selected for each island and the debug bus lane to which they are routed. The routed-to debug bus lanes become busy when an island is selected to route their performance monitor signals. Again, care must be taken not to allow any island to route their signals onto a lane that is in use by another island at the same time.

In general, any island word 0 can be selected with any other island word 2 for performance monitor counting. Note also that an NClk signal group (PPU, SPU, or PPSS-NClk) can be selected in combination with any NClk/2 signal group providing that they do not use the same lane on the debug bus.

*Table C-2. Lanes Used by Different Islands on the Debug Bus*

| Debug Bus Lane 0 | Debug Bus Lane 1 | Debug Bus Lane 2 | Debug Bus Lane 3 |
|---|---|---|---|
| PPU word 0 | PPU word 1 | | |
| PPSS-NClk word 0 | PPSS-NClk word 1 | | |
| PPSS-NClk/2 word 0 | PPSS-NClk/2 word 1 | PPSS-NClk/2 word 2 | PPSS-NClk/2 word 3 |

*Table C-2. Lanes Used by Different Islands on the Debug Bus*

| Debug Bus Lane 0 | Debug Bus Lane 1 | Debug Bus Lane 2 | Debug Bus Lane 3 |
|---|---|---|---|
| MIC word 0 | MIC word 1 | MIC word 2 | MIC word 3 |
| EIB word 0 | EIB word 1 | EIB word 2 | EIB word 3 |
| BEI word 0 | BEI word 1 | BEI word 2 | BEI word 3 |
| MFC word 0 | MFC word 1 | MFC word 2 | MFC word 3 |
| SPU word 0 phase A | SPU word 0 phase B | SPU word 1 phase A | SPU word 1 phase B |

To use the performance monitor, the signals must pass through all the intervening islands to the pervasive logic where the performance monitor logic is located. This is done by activating the debug bus at *all* the units as described in *Table C-3*. All address offsets are relative to the BE_MMIO_Base address.

*Table C-3. Register Settings to Enable Pass-Through of Debug Bus Signals through Islands*

| Address Offset | Register Short Name | Enable Pass-Through |
|---|---|---|
| x‘50 0858’ | L2_Debug1 | rmw bit 61 = 1 |
| x‘50 0958’ | CIU_DR1 | rmw bit 5 = 1 |
| x‘50 9C98’ | on_ramp_trace | rmw bit 39 = 1 |
| x‘50 A010’ | MBL_Debug | rmw bit 20 = 1 |
| x‘40 03C8’ | SBI_PMCR (MFC0) | rmw bit 38 = 1 |
| x‘40 23C8’ | SBI_PMCR (MFC1) | rmw bit 38 = 1 |
| x‘40 43C8’ | SBI_PMCR (MFC2) | rmw bit 38 = 1 |
| x‘40 63C8’ | SBI_PMCR (MFC3) | rmw bit 38 = 1 |
| x‘40 83C8’ | SBI_PMCR (MFC4) | rmw bit 38 = 1 |
| x‘40 A3C8’ | SBI_PMCR (MFC5) | rmw bit 38 = 1 |
| x‘40 C3C8’ | SBI_PMCR (MFC6) | rmw bit 38 = 1 |
| x‘40 E3C8’ | SBI_PMCR (MFC7) | rmw bit 38 = 1 |

### C.1.1 An Example of Setting up the Performance Monitor in PPSS L2 Mode A

A step-by-step description follows of how to set up the performance monitor logic to count the number of hits and misses in the PPSS level 2 (L2) cache. Hit events are available in bit 0, and miss events are available in bit 1 of the L2 Mode A signal group described in *Section C.3.2 PPSS L2 Cache Controller - Group 1 (NClk/2)* on page 800.

The following example uses performance monitor counters 0 and 1 in 32-bit mode, and assumes that they will not overflow. If there is any possibility of the counters reaching their maximum value, the chosen counters must be set up to either stop counting when they reach the maximum value or to use the trace buffer to store their contents periodically, as defined by the *pm_interval* register. See *Cell Broadband Engine Book IV for DD 3.0, DD 3.1, DD 3.2* and *Cell Broadband Engine Registers* for details about how to do set up these options, and for a description of all the registers used in this example.

1. *Activate the main debug bus through all the units.* This step involves doing a read-modify-write to each of the registers shown in *Table C-3* on page 795.

2. *Enable performance monitor logic in the L2 unit.* This step requires two read-modify-write operations: one to *L2_Perfmon1* and another to *L2_Debug1* as shown in *Section C.3.2 PPSS L2 Cache Controller - Group 1 (NClk/2)* on page 800. For this example, the signals are assumed to be enabled to route onto the debug bus lane 2.
3. *Enable the debug bus for the performance monitor signals from the PPSS L2 island.* This is done by writing the value of x‘0800 0000’ to the *debug_bus_control* register.
4. *Route debug bus lane 2 (or word 2) to the performance monitor counter input multiplexer bits 0-31*. This is done by writing the value of x‘8000 0000’ to the *group_control* register.
5. *Choose the signals to be monitored by performance monitor counters 0 and 1.* The two events counted in this example, L2 hits and L2 misses, are edge-triggered events and are in bits 0 and 1, respectively. Set up the two performance monitor counters to count these signals by writing x‘0140 0000’ to the *pm0_control register,* and x‘0540 0000’ to the *pm1_control* register.
6. *Set the interval timer.* Because the interval timer counts up, write the value x‘0000 0000’ to the *pm_interval* register to count up to $2^{32}$-1 NClks. In this example, the counters are assumed not to overflow.
7. *Initialize the counters*. The two counters pm0_4 and pm1_5 must each be initialized by explicitly writing zeros into them.
8. *Start the performance monitor counting logic at the appropriate time.* This is done by writing the value of x‘8000 0000’ to the *pm_control* register.
9. *Stop the performance monitor counting logic at the appropriate time.* This is done by writing the value of x‘0000 0000’ to the *pm_control* register.

   **Note:** The value written to the *pm_control* register must the same as in step 8, except that bit 0 must be set to zero to disable the counting. For example, if the performance monitor counters are being used in 16-bit mode, the nonzero values in bits 7-10 must be preserved as written in step 8. This step also stops the interval timer.
10. *Read out the values in the counters.* The 32-bit value in the *pm0_4* counter indicates the number of L2 hits, and the 32-bit value in the *pm1_5* counter indicates the number of L2 misses. The number of NClks that elapsed between the enabling and disabling of the performance monitor logic can be found by reading the *pm_interval* register.

## C.2 PowerPC Processor Unit (PPU) Signal Selection

### C.2.1 PPU Instruction Unit

| SPR Address | Register Short Name | Enable Word 0 | Enable Word 1 |
|---|---|---|---|
| 1009 | HID1 | rmw bits 31, 32:35, 40:41, 48, 58, 60:61 = ‘1 1111 01 0 0 00’ | rmw bits 31, 36:39, 40:41, 48, 58, 62:63 = ‘1 1111 01 0 0 01’ |

| Bit | Type | Description |
|---|---|---|
| **Thread 0** | | |
| 0 | V | Branch instruction committed. |
| 1 | E | Branch instruction that caused a misprediction flush is committed. Branch misprediction includes mispredictions of taken or not-taken on a conditional branch and mispredictions of branch target address on bclr[1] and bcctr[1]. |
| 2 | C | Instruction buffer empty. |
| 3 | E | Instruction effective-address-to-real-address translation (I-ERAT) miss. |
| 4 | B | Level 1 (L1) instruction cache miss cycles. Counts the cycles from the miss event until the returned instruction is dispatched or cancelled due to branch misprediction, completion restart, or exceptions (see **Note 1**). |
| 6 | C | Valid instruction is available for dispatch, but the dispatch is blocked (see **Note 2**). |
| 9 | E | Instruction in pipeline stage EX7 causes a flush. |
| 11 | V | Two PowerPC instructions committed. For microcode sequences, only the last microcode operation is counted. Committed instructions are counted two at a time. If only one instruction has committed for a given cycle, this event is not raised until another instruction has been committed in a future cycle. |
| **Thread 1** | | |
| 19 | V | Branch instruction committed. |
| 20 | E | Branch instruction that caused a misprediction flush is committed. Branch misprediction includes mispredictions of taken or not-taken on a conditional branch and mispredictions of branch target address on bclr[1] and bcctr[1]. |
| 21 | C | Instruction buffer empty. |
| 22 | E | I-ERAT miss. |
| 23 | B | L1 Instruction cache miss cycles. Counts the cycles from the miss event until the returned instruction is dispatched or cancelled due to a branch misprediction, a completion restart, or exceptions (see **Note 1**). |
| 25 | C | Valid instruction is available for dispatch, but dispatch is blocked (see **Note 2**). |
| 28 | E | Instruction in pipeline stage EX7 causes a flush. |
| 30 | V | Two PowerPC instructions committed. For microcode sequences only, the last microcode operation is counted. Committed instructions are counted two at a time. If only one instruction has committed for a given cycle, this event is not be raised until another instruction has been committed in a future cycle. |

1. Counting the rising edge of this event is only an approximation of the number of I-cache misses. If counting edges, note that the rising edge does not occur when an older I-cache miss is being resolved while a new I-cache miss occurs. This case can happen when the first miss is speculative and is canceled on a later cycle while a new nonspeculative miss occurs on the redirect fetch address.
2. These performance monitor signals are not asserted or counted in the event of a blocked dispatch if its hardware thread has a higher priority than the other hardware thread.

### C.2.2 PPU Execution Unit (NClk)

| SPR Address | Register Short Name | Enable Word 0 | Enable Word 1 |
|---|---|---|---|
| 1009 | HID1 | rmw bits 31, 32:35, 42:44, 49, 58, 60:61 = ‘1 1111 001 0 0 10’ | rmw bits 31, 36:39, 45:47, 49, 58, 62:63 = ‘1 1111 011 0 0 10’ |

| Bit | Type | Description |
|---|---|---|
| **Thread 0** | | |
| 2 | S | Data effective-address-to-real-address translation (D-ERAT) miss. This event is not speculative. |
| 3 | S | Store request counted at the L2 interface. This counts microcoded PowerPC Processor Element (PPE) sequences more than once (see **Note 1** for exceptions). |
| 4 | S | Load valid at a particular pipe stage. This is speculative because flushed operations are also counted. Counts microcoded PPE sequences more than once. Misaligned flushes might be counted the first time as well. Load operations include all loads that read data from the cache, **dcbt** and **dcbtst**. This event does not include load Vector/single instruction multiple data (SIMD) multimedia extension pattern instructions. |
| 5 | S | L1 D-cache load miss. Pulsed when there is a miss request that has a tag miss but not an effective-address-to-real-address translation (ERAT) miss. This is speculative because flushed operations are counted as well. |
| **Thread 1** | | |
| 18 | S | D-ERAT miss. This event is not speculative. |
| 20 | S | Load valid at a particular pipe stage. This event is speculative because flushed operations are counted as well. Counts microcoded PPE sequences more than once. Additionally, misaligned flushes might also be counted the first time. Load operations include all loads that read data from the cache, **dcbt** and **dcbtst**. This event does not include load Vector/SIMD multimedia extension pattern instructions. |
| 21 | S | L1 D-cache load miss. This event is pulsed when there is a miss request that has a tag miss but not an ERAT miss. This event is speculative because flushed operations are counted as well. |

1. In single-threaded mode, bit 3 counts the stores for thread 0. In multithreaded mode, bit 3 might be incorrect.

## C.3 PowerPC Storage Subsystem (PPSS) Signal Selection

The address offset of each setup register is relative to the BE_MMIO_Base address.

### C.3.1 PPSS Bus Interface Unit (NClk/2)

| Address Offset | Register Short Name | Enable Word 2 |
|---|---|---|
| x‘50 0B58’ | BIU_DbgPerfmon | rmw bits 0:2 = ‘110’ |
| x‘50 0858’ | L2_Debug1 | rmw bit 61 = 1 |

| Bit | Type | Description |
|---|---|---|
| 0 | E | Load from memory flow controller (MFC) memory-mapped input and output (MMIO) space. |
| 1 | E | Stores to MFC MMIO space |
| 6 | E | Request token type (even memory bank numbers 0-14) |
| 15 | E | Receive 8-beat data from the element interconnect bus (EIB). |
| 16 | E | Eight-beat data was sent to the EIB. |
| 17 | E | A command was sent to the EIB (includes retried commands). |
| 18 | C | Number of cycles between a data request and a data grant. |
| 23 | C | The 5-entry noncacheable unit store command queue is not empty. |

### C.3.2 PPSS L2 Cache Controller - Group 1 (NClk/2)

The L2 performance group 1 (mode A) shares the upper half of the L2 trace bus with the other L2 performance group modes. Therefore, only one group is allowed to be active at a time. Do not set more than one bit in L2_Perfmon1[59:62]. The upper half of the L2 trace bus (bits 0 to 31) is also shared with the upper half of the L2 debug mode selection. These debug modes must also be disabled to observe any of the L2 performance groups.

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 0838’ | L2_Perfmon1 | rmw bits 59:63 = ‘10001’ | |
| x‘50 0858’ | L2_Debug1 | rmw bits 43:61, 62 = ‘0000000000000000001 1’ | rmw bits 43:61, 63 = ‘0000000000000000001 1’ |

| Bit | Type | Description |
|---|---|---|
| 0 | E | Cache hit for core interface unit (CIU) loads and stores. |
| 1 | E | Cache miss for CIU loads and stores. |
| 4 | E | CIU load miss. |
| 5 | E | CIU store to Invalid state (miss). |
| 7 | E | Load word and reserve indexed (**lwarx**/**ldarx**) for thread 0 hits Invalid (I) cache state. |
| 14 | E | Store word conditional indexed (**stwcx**/**stdcx**) for thread 0 hits Invalid (I) cache state when reservation is set. |
| 25 | C | All four snoop state machines are busy. |

### C.3.3 PPSS L2 Cache Controller - Group 2 (NClk/2)

The L2 performance group 2 (mode B) shares the upper half of the L2 trace bus with the other L2 performance group modes. Therefore, only one group is allowed to be active at a time. Do not set more than one bit in L2_Perfmon1[59:62]. The upper half of the L2 trace bus (bits 0 to 31) is also shared with the upper half of the L2 debug mode selection. These debug modes also must be disabled to observe any of the L2 performance groups.

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 0838’ | L2_Perfmon1 | rmw bits 59:63 = ‘01001’ | |
| x‘50 0858’ | L2_Debug1 | rmw bits 43:61, 62 = ‘0000000000000000001 1’ | rmw bits 43:61, 63 = ‘0000000000000000001 1’ |

| Bit | Type | Description |
|---|---|---|
| 4 | E | Data line claim (**dclaim**) that received a good combined response. This event includes store/**stcx**/**dcbz** to Shared (S), Shared Last ($S_L$), or Tagged (T) cache state. This event does not include **dcbz** to Invalid (I) cache state (see **Note 1**). |
| 11 | E | **Dclaim** converted into **rwitm**; might still not get to the bus if **stcx** is cancelled (see **Note 2**). |
| 12 | E | Store to Modified (M), Modified Unsolicited ($M_U$), or Exclusive (E) cache state |
| 13 | C | Eight-entry store queue is full. |
| 14 | E | Store dispatched to read and claim (RC) machine is acknowledged. |
| 15 | V | Gatherable store (type = ‘00000’) was received from CIU. |
| 18 | E | Snoop push |
| 19 | E | Send intervention from ($S_L$ \| E) cache state to a destination within the same CBEA processor. |
| 20 | E | Send intervention from (M \| $M_U$) cache state to a destination within the same CBEA processor. |
| 25 | E | Respond with retry to a snooped request due to one of the following conflicts: RC state machine full address, castout congruence class, snoop machine full address, all snoop machines busy, directory lockout, or parity error. |
| 26 | E | Respond with retry to a snooped request because all snoop machines are busy. |
| 27 | E | Snooped response causes a cache state transition from (M \| $M_U$) to (E \| S \| T). |
| 28 | E | Snooped response causes a cache state transition from E to S. |
| 29 | E | Snooped response causes a cache state transition from (E \| SL \| S \| T) to Invalid (I). |
| 30 | E | Snooped response causes a cache state transition from (M \| $M_U$) to I. |

1. A combined response of retry for an atomic **dclaim** (due to a lost reservation) is valid.
2. An atomic **dclaim** that is converted to **rwitm** might not get to the bus because of a lost reservation, but it is still counted.

### C.3.4 PPSS L2 Cache Controller - Group 3 (NClk/2)

The L2 performance group 3 (mode C) shares the upper half of the L2 trace bus with the other L2 performance group modes. Therefore, only one group is allowed to be active at a time. Do not set more than one bit in L2_Perfmon1[59:62]. The upper half of the L2 trace bus (bits 0 to 31) is also shared with the upper half of the L2 debug mode selection. These debug modes also must be disabled to observe any of the L2 performance groups.

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 0838’ | L2_Perfmon1 | rmw bits 59:63 = ‘00101’ | |
| x‘50 0858’ | L2_Debug1 | rmw bits 43:61, 62 = ‘0000000000000000001 1’ | rmw bits 43:61, 63 = ‘0000000000000000001 1’ |

| Bit | Type | Description |
|---|---|---|
| 12 | E | Load and reserve indexed (**lwarx**/**ldarx**) for thread 1 hits an Invalid (I) cache state. |
| 19 | E | Store conditional indexed (**stwcx**/**stdcx**) for thread 1 hits an Invalid (I) cache state when reservation is set. |

### C.3.5 PPSS Noncacheable Unit (NCIk/2)

| Address Offset | Register Short Name | Enable Word 0 |
|---|---|---|
| x‘50 0858’ | L2_Debug1 | rmw bit 61 = 1 |
| x‘50 0A58’ | NCU_DR1 | rmw bit 0 = 1 |
| x‘50 0A60’ | NCU_DR2 | rmw bits 0:6 = ‘0000001’ |

| Bit | Type | Description |
|---|---|---|
| 0 | V | Noncacheable store request received from CIU. This event includes all synchronization operations such as **sync** and **eieio**. |
| 1 | V | **sync** received from CIU. |
| 4 | V | A noncacheable store request was received from the CIU (includes *only* stores). |
| 6 | V | **eieio** received from CIU. |
| 7 | V | **tlbie** received from CIU. |
| 8 | C | A **sync** is at the bottom of the store queue (waiting on *st_done* signal from the bus interface unit [BIU] and *sync_done* signal from L2). |
| 9 | C | An **lwsync** is at the bottom of the store queue (waiting for a *sync_done* signal from the L2). |
| 10 | C | An **eieio** is at the bottom of the store queue (waiting for a *st_done* signal from the BIU and a *sync_done* signal from the L2). |
| 11 | C | A **tlbie** is at the bottom of the store queue (waiting for a *st_done* signal from the BIU). |
| 12 | V | Noncacheable store combined with the previous noncacheable store with a contiguous address |
| 15 | C | All four store-gather buffers are full. |
| 16 | E | A noncacheable load request was received from the CIU (includes instruction and data fetches). |
| 17 | C | The 4-deep store queue is not empty. |
| 18 | C | The 4-deep store queue is full. |
| 19 | C | At least one store gather buffer is not empty. |

## C.4 Synergistic Processor Unit (SPU) Signal Selection

The address offset of each setup register is relative to the BE_MMIO_Base address + (x‘2000’ × Synergistic Processor Element (SPE) number). For example, the address offset of the *SPU_trace_sel* register in SPU 1 is at offset x‘40 3048’ from the BE_MMIO_Base address.

### C.4.1 SPU (NCIk)

| Address Offset | Register Short Name | Enable Word 0 (see **Note 1**) | Enable Word 2 (see **Note 1**) |
|---|---|---|---|
| x‘40 1048’ | SPU_trace_sel | rmw bits 32:63 = x‘0000 1000’ | rmw bits 32:63 = x‘0000 0010’ |
| x‘40 1070’ | SPU_trace_cntl | rmw bit 48 = 1 | |
| x‘40 03C8’ | SBI_PMCR | rmw bits 28:38 = ‘11111111011’ (even SPU) | rmw bits 20:27, 36:38 = ‘11111111 011’ (even SPU) |
| | | rmw bits 20:27, 36:38 = ‘11111111 101’ (odd SPU) | rmw bits 28:35, 36:38 = ‘11111111 101’ (odd SPU) |

| Bit | Type | Description |
|---|---|---|
| 0 | S | A dual instruction is committed. |
| 1 | S | A single instruction is committed. |
| 2 | S | A pipeline 0 instruction is committed. |
| 3 | S | A pipeline 1 instruction is committed. |
| 5 | B | Local storage is busy. |
| 6 | S | A direct memory access (DMA) might conflict with a load or store. |
| 7 | S | A store instruction to local storage is issued. |
| 8 | S | A load instruction from local storage is issued. |
| 9 | S | A floating-point unit exception occurred. |
| 10 | S | A branch instruction is committed. |
| 11 | S | A nonsequential change of the SPU program counter has occurred. This can be caused by branch, asynchronous interrupt, stalled wait on channel, error-correction code (ECC) error, and so forth. |
| 12 | S | A branch was not taken. |
| 13 | S | Branch miss prediction. This count is not exact. Certain other code sequences can cause additional pulses on this signal (see **Note 2**). |
| 14 | S | Branch hint miss prediction. This count is not exact. Certain other code sequences can cause additional pulses on this signal (see **Note 2**). |
| 15 | S | Instruction sequence error |

1. When SPU word 0 is enabled, the speed-converted data is routed from the SPU to debug bus lanes 0 and 1. When SPU word 2 is enabled, the speed-converted data is routed from the SPU to debug bus lanes 2 and 3.
2. Examples of false pulses on misprediction signals include the following items:
   - A blocking channel (count = 0) accessed through **wrch** or **rdch** that immediately follows a branch instruction
   - An asynchronous interrupt that occurs after a branch is committed
   - A correctable instruction or data ECC error after a branch is committed but before another instruction is committed
   - If the SPU is stopped through the *SPU_RunCntl* register or by a thermal monitor event after a branch is committed but before another instruction is committed
3. This signal is triggered each time a channel read or channel write instruction is issued to the affected channel and the count for the channel is zero. Note that a channel read or write can be issued speculatively following a mispredicted branch instruction.

| Bit | Type | Description |
|---|---|---|
| 17 | B | Stalled waiting on any blocking channel write (see **Note 3**). |
| 18 | B | Stalled waiting on external event status (Channel 0) (see **Note 3**). |
| 19 | B | Stalled waiting on SPU Signal Notification 1 (Channel 3) (see **Note 3**). |
| 20 | B | Stalled waiting on SPU Signal Notification 2 (Channel 4) (see **Note 3**). |
| 21 | B | Stalled waiting on DMA Command Opcode or ClassID Register (Channel 21) (see **Note 3**). |
| 22 | B | Stalled waiting on memory flow controller (MFC) Read Tag-Group Status (Channel 24) (see **Note 3**). |
| 23 | B | Stalled waiting on MFC Read List Stall-and-Notify Tag Status (Channel 25) (see **Note 3**). |
| 24 | B | Stalled waiting on SPU Write Outbound Mailbox (Channel 28) (see **Note 3**). |
| 30 | B | Stalled waiting on SPU Mailbox (Channel 29) (see **Note 3**). |

1. When SPU word 0 is enabled, the speed-converted data is routed from the SPU to debug bus lanes 0 and 1. When SPU word 2 is enabled, the speed-converted data is routed from the SPU to debug bus lanes 2 and 3.
2. Examples of false pulses on misprediction signals include the following items:
   - A blocking channel (count = 0) accessed through **wrch** or **rdch** that immediately follows a branch instruction
   - An asynchronous interrupt that occurs after a branch is committed
   - A correctable instruction or data ECC error after a branch is committed but before another instruction is committed
   - If the SPU is stopped through the *SPU_RunCntl* register or by a thermal monitor event after a branch is committed but before another instruction is committed
3. This signal is triggered each time a channel read or channel write instruction is issued to the affected channel and the count for the channel is zero. Note that a channel read or write can be issued speculatively following a mispredicted branch instruction.

### C.4.2 SPU Trigger (NClk)

| Enable | Address Offset | Register Short Name | Value |
|---|---|---|---|
| Trigger 0 | x‘40 1028’ | SPU_trig0_sel | rmw bit 57 = 1 |
| Trigger 1 | x‘40 1030’ | SPU_trig1_sel | |
| Trigger 2 | x‘40 1038’ | SPU_trig2_sel | |
| Trigger 3 | x‘40 1040’ | SPU_trig3_sel | |
| Triggers 0, 1, 2, 3 | x‘40 1070’ | SPU_trace_cntl | rmw bit 48 = 1 |
| Trigger 0 | x‘40 03C8’ | SBI_PMCR | rmw bit 39 = 1 |
| Trigger 1 | | | rmw bit 40 = 1 |
| Trigger 2 | | | rmw bit 41 = 1 |
| Trigger 3 | | | rmw bit 42 = 1 |

| Bit | Type | Description |
|---|---|---|
| 57 | B | Stalled waiting on a channel operation (see **Notes 1 and 2**). |

1. When SPU trigger 0, 1, 2, or 3 is selected, then the speed-converted data passes through the MFC on trigger bits 4, 5, 6, or 7 of the debug bus, in addition to bits 0, 1, 2, or 3.
2. This signal is triggered each time a channel read or channel write instruction is issued to the affected channel and the count for the channel is zero. Note that a channel read or write can be issued speculatively following a mispredicted branch instruction.

### C.4.3 SPU Event (NClk)

| Enable | Address Offset | Register Short Name | Value |
|---|---|---|---|
| Event 0 | x‘40 1050’ | SPU_event0_sel | rmw bit 59, 63 = ‘1 1’ |
| Event 1 | x‘40 1058’ | SPU_event1_sel | |
| Event 2 | x‘40 1060’ | SPU_event2_sel | |
| Event 3 | x‘40 1068’ | SPU_event3_sel | |
| Events 0, 1, 2, 3 | x‘40 1070’ | SPU_trace_cntl | rmw bit 48 = 1 |
| Event 0 | x‘40 03C8’ | SBI_PMCR | rmw bit 43 = 1 (see **Note 1**) |
| Event 1 | | | rmw bit 44 = 1 (see **Note 1**) |
| Event 2 | | | rmw bit 45 = 1 |
| Event 3 | | | rmw bit 46 = 1 |

| Bit | Type | Description |
|---|---|---|
| 59 | B | Instruction fetch stall |
| 63 | E | Serialized SPU address (program counter) trace (see **Note 2)** |

1. When any of the three memory interface control signal groups is enabled for performance monitoring, event 0 and event 1 cannot be used.
2. To enable tracing of the program counter or bookmarks, pairs of events (0 and 1) or (2 and 3) must be enabled together.

## C.5 Memory Flow Controller (MFC) Signal Selection

The address offset of each setup register is relative to the BE_MMIO_Base address + (x‘2000’ × Synergistic Processor Element [SPE] number). For example, the address of the *ATO_PMCR* register for synergistic processor unit (SPU) number 7 (SPU7) is at offset x‘40 E3D0’ from the BE_MMIO_Base address.

### C.5.1 MFC Atomic Unit (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘40 03D0’ | ATO_PMCR | rmw bits 34:35 = ‘10’ | rmw bits 34:35 = ‘01’ |
| x‘40 0580’ | SMM_HID | rmw bit 23 = 1 | |
| x‘40 03C8’ | SBI_PMCR | rmw bits 28:31, 36:38 = ‘1111 011’ (even SPU) | rmw bits 20:23, 36:38 = ‘1111 011’ (even SPU) |
| | | rmw bits 20:23, 36:38 = ‘1111 101’ (odd SPU) | rmw bits 28:31, 36:38 = ‘1111 101’ (odd SPU) |

| Bit | Type | Description |
|---|---|---|
| 1 | E | An atomic load was received from the direct memory access controller (DMAC). |
| 2 | E | An atomic **dclaim** was sent to the synergistic bus interface (SBI). This includes retried requests. |
| 3 | E | An atomic **rwitm** performed was sent to the SBI. This includes retried requests. |
| 4 | E | An atomic load miss caused $M_U$ cache state. |
| 5 | E | An atomic load miss caused E cache state. |
| 6 | E | An atomic load miss caused $S_L$ cache state. |
| 7 | E | An atomic load hits the cache. |
| 8 | E | An atomic load misses cache with data intervention. This is the sum of both events 4 and 6 in this group. |
| 14 | E | **putllc** or **putlluc** misses cache without data intervention. For **putllc**, this event counts only when reservation is set for the address. |
| 17 | C | The snoop machine is busy. |
| 19 | E | A snoop caused a cache state transition from the (M \| $M_U$) to the invalid (I) state. |
| 21 | E | A snoop caused a cache state transition from the (E \| S \| $S_L$) to the I state. |
| 23 | E | A snoop caused a cache state transition from the $M_U$ to the T state. |
| 27 | E | Sent modified data intervention to a destination within the same CBEA processor. |

### C.5.2 MFC Direct Memory Access Controller (NCIk/2)

| Address Offset | Register Short Name | Enable Word 0 |
|---|---|---|
| x‘40 0880’ | DMAC_PMCR | rmw bit 40 = 1 |
| x‘40 03C8’ | SBI_PMCR | rmw bits 28:31, 36:38 = ‘1111 011’ (even SPU) |
| | | rmw bits 20:23, 36:38 = ‘1111 101’ (odd SPU) |

| Bit | Type | Description |
|---|---|---|
| 0 | V | Any flavor of direct memory access (DMA) **get[]** command issued to the SBI. |
| 1 | V | Any flavor of DMA **put[]** command issued to the SBI. |
| 2 | V | DMA put (**put**) is issued to the SBI; counts only nonzero transfer size cases. |
| 17 | V | DMA get data from effective address to local storage (**get**) issued to the SBI; counts only nonzero transfer size cases. |

### C.5.3 MFC Synergistic Bus Interface (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘40 03C8’ | SBI_PMCR | rmw bits 28:31, 36:38, 47:48, 56 = ‘1111 011 00 1’ (even SPU) | rmw bits 20:23, 36:38, 47:48, 56 = ‘1111 011 01 0’ (even SPU) |
| | | rmw bits 20:23, 36:38, 47:48, 56 = ‘1111 101 00 1’ (odd SPU) | rmw bits 28:31, 36:38, 47:48, 56 = ‘1111 101 01 0’ (odd SPU) |

| Bit | Type | Description |
|---|---|---|
| 4 | V | Load request sent to the element interconnect bus (EIB). This includes **read**, **read atomic**, **rwitm, rwitm atomic,** and retried commands. |
| 5 | V | Store request sent to the EIB. This includes **wwf, wwc, wwk, dclaim, dclaim atomic,** and retried commands. |
| 6 | E | Received data from the EIB, including partial cache line data. |
| 7 | E | Sent data to the EIB, both as a master and a snooper. |
| 8 | C | Sixteen-deep SBI queue with outgoing requests not empty. This does not include atomic requests. |
| 9 | C | Sixteen-deep SBI queue with outgoing requests full. This does not include atomic requests. |
| 10 | V | Sent request to EIB. |
| 12 | E | Received data bus grant (includes data sent for MMIO operations). |
| 13 | C | Cycles between data bus request and data bus grant. |
| 14 | V | Command (read or write) for an odd-numbered memory bank. This is valid only when resource allocation is turned on. |
| 15 | V | Command (read or write) for an even-numbered memory bank. This is valid only when resource allocation is turned on. |
| 18 | V | Request gets the retry response (includes local and global requests). |
| 19 | E | Sent data bus request to the EIB. |

### C.5.4 MFC Synergistic Memory Management (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘40 0580‘ | SMM_HID | rmw bits 12, 23, 24 = ‘1 1 1’ | rmw bits 14, 23, 24 = ‘1 1 1’ |
| x‘40 03C8’ | SBI_PMCR | rmw bits 28:31, 36:38 = ‘1111 011’ (even SPU) | rmw bits 20:23, 36:38 = ‘1111 011’ (even SPU) |
| | | rmw bits 20:23, 36:38 = ‘1111 101’ (odd SPU) | rmw bits 28:31, 36:38 = ‘1111 101’ (odd SPU) |

| Bit | Type | Description |
|---|---|---|
| 0 | E | Translation lookaside buffer (TLB) miss without parity or protection errors. |
| 1 | C | TLB miss (cycles). |
| 2 | E | TLB hit. |

## C.6 Element Interconnect Bus (EIB) Signal Selection

The address offset of each setup register is relative to the BE_MMIO_Base address.

The EIB has four separate ports to the debug bus. They can be used singly or in any combination with each other. The four ports are fed by the global address concentrator (AC0), the local address concentrator (AC1), the data arbiter, and the token manager (TKM).

### C.6.1 EIB Address Concentrator 0 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 1808‘ | EIB_AC0Perf | rmw bits 0:3, 16, 18:21 = ‘1111 0 0000’ | rmw bits 8:11, 17, 18:21 = ‘1111 0 0000’ |

| Bit | Type | Description |
|---|---|---|
| **Group 1** | | |
| 0 | V | Number of **read** and **rwitm** commands (including atomic) from AC1 to AC0. |
| 1 | V | Number of **dclaim** commands (including atomic) from AC1 to AC0. |
| 2 | V | Number of **wwk**, **wwc**, and **wwf** commands from AC1 to AC0. |
| 3 | V | Number of **sync**, **tlbsync**, and **eieio** commands from AC1 to AC0. |
| 4 | V | Number of **tlbie** commands from AC1 to AC0. |
| 11 | E | Previous adjacent address match (PAAM) content addressable memory (CAM) hit. |
| 12 | E | PAAM CAM miss. |
| 14 | V | Command reflected. |
| **Group 2** | | |
| 16 | V | Number of **read** and **rwitm** commands (including atomic) from AC1 to AC0. |
| 17 | V | Number of **dclaim** commands (including atomic) from AC1 to AC0. |
| 18 | V | Number of **wwk**, **wwc**, and **wwf** commands from AC1 to AC0. |
| 19 | V | Number of **sync**, **tlbsync**, and **eieio** commands from AC1 to AC0. |
| 20 | V | Number of **tlbie** commands from AC1 to AC0. |
| 27 | E | PAAM CAM hit. |
| 28 | E | PAAM CAM miss. |
| 30 | V | Command reflected. |

### C.6.2 EIB Address Concentrator 1 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 1860‘ | EIB_DbgL | rmw bits 0:3, 8:11 = ‘0101 1111’ | |
| x‘51 1868’ | EIB_DbgR | | rmw bits 0:3, 8:11 = ‘0110 1111’ |

| Bit | Type | Description |
|---|---|---|
| 1 | V | Local command from Synergistic Processor Element (SPE) 6. |
| 2 | V | Local command from SPE 4. |
| 3 | V | Local command from SPE 2. |
| 5 | V | Local command from the PowerPC Processor Element (PPE). |
| 6 | V | Local command from SPE 1. |
| 7 | V | Local command from SPE 3. |
| 8 | V | Local command from SPE 5. |
| 9 | V | Local command from SPE 7. |
| 12 | V | AC1-to-AC0 global command from SPE 6. |
| 13 | V | AC1-to-AC0 global command from SPE 4. |
| 14 | V | AC1-to-AC0 global command from SPE 2. |
| 15 | V | AC1-to-AC0 global command from SPE 0. |
| 16 | V | AC1-to-AC0 global command from PPE. |
| 17 | V | AC1-to-AC0 global command from SPE 1. |
| 18 | V | AC1-to-AC0 global command from SPE 3. |
| 19 | V | AC1-to-AC0 global command from SPE 5. |
| 20 | V | AC1-to-AC0 global command from SPE 7. |
| 22 | V | AC1 is reflecting any local command. |
| 23 | V | AC1 sends a global command to AC0. |
| 24 | V | AC0 reflects a global command back to AC1. |
| 25 | V | AC1 reflects a command back to the bus masters. |

### C.6.3 EIB Data Ring Arbiter - Group 1 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 1860‘ | EIB_DbgL | rmw bits 0:3, 8:11 = ‘1110 1111’ | |
| x‘51 1868‘ | EIB_DbgR | | rmw bits 0:3, 8:11 = ‘1110 1111’ |

| Bit | Type | Description |
|---|---|---|
| 0 | E | Grant on data ring 0. |
| 1 | E | Grant on data ring 1. |
| 2 | E | Grant on data ring 2. |
| 3 | E | Grant on data ring 3. |
| 4 | C | Data ring 0 is in use. |
| 5 | C | Data ring 1 is in use. |
| 6 | C | Data ring 2 is in use. |
| 7 | C | Data ring 3 is in use. |
| 8 | C | All data rings are idle. |
| 9 | C | One data ring is busy. |
| 10 | C | Two or three data rings are busy. |
| 11 | C | All data rings are busy. |
| 12 | B | An IOIF0 data request is pending. |
| 13 | B | An SPE 6 data request is pending. |
| 14 | B | An SPE 4 data request is pending. |
| 15 | B | An SPE 2 data request is pending. |
| 16 | B | An SPE 0 data request is pending. |
| 17 | B | A memory interface controller (MIC) data request is pending. |
| 18 | B | A PPE data request is pending. |
| 19 | B | An SPE 1 data request is pending. |
| 20 | B | An SPE 3 data request is pending. |
| 21 | B | An SPE 5 data request is pending. |
| 22 | B | An SPE 7 data request is pending. |
| 24 | E | IOIF0 is the data destination. |
| 25 | E | SPE 6 is the data destination. |
| 26 | E | SPE 4 is the data destination. |
| 27 | E | SPE 2 is the data destination. |
| 28 | E | SPE 0 is the data destination. |
| 29 | E | MIC is the data destination. |
| 30 | E | PPE is the data destination. |
| 31 | E | SPE 1 is the data destination. |

### C.6.4 EIB Data Ring Arbiter - Group 2 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 1860‘ | EIB_DbgL | rmw bits 0:3, 8:11 = ‘1111 1111’ | |
| x‘51 1868‘ | EIB_DbgR | | rmw bits 0:3, 8:11 = ‘1111 1111’ |

| Bit | Type | Description |
|---|---|---|
| 0 | B | An IOIF0 data request is pending. |
| 1 | B | An SPE 6 data request is pending. |
| 2 | B | An SPE 4 data request is pending. |
| 3 | B | An SPE 2 data request is pending. |
| 4 | B | An SPE 0 data request is pending. |
| 5 | B | A MIC data request is pending. |
| 6 | B | A PPE data request is pending. |
| 7 | B | An SPE 1 data request is pending. |
| 8 | B | An SPE 3 data request is pending. |
| 9 | B | An SPE 5 data request is pending. |
| 10 | B | An SPE 7 data request is pending. |
| 11 | B | An IOIF1 data request is pending. |
| 12 | E | IOIF0 is the data destination. |
| 13 | E | SPE 6 is the data destination. |
| 14 | E | SPE 4 is the data destination. |
| 15 | E | SPE 2 is the data destination. |
| 16 | E | SPE 0 is the data destination. |
| 17 | E | MIC is the data destination. |
| 18 | E | PPE is the data destination. |
| 19 | E | SPE 1 is the data destination. |
| 20 | E | SPE3 is the data destination. |
| 21 | E | SPE 5 is the data destination. |
| 22 | E | SPE 7 is the data destination. |
| 23 | E | IOIF1 is the data destination. |
| 24 | E | A grant is on data ring 0. |
| 25 | E | A grant is on data ring 1. |
| 26 | E | A grant is on data ring 2. |
| 27 | E | A grant is on data ring 3. |
| 28 | C | All data rings are idle. |
| 29 | C | One data ring is busy. |
| 30 | C | Two or three data rings are busy. |
| 31 | C | All four data rings are busy. |

### C.6.5 EIB Token Manager (NClk/2)

The address offset of the *TKM_PMCR* register is relative to the BE_MMIO_Base address.

The TKM performance monitor bus is 64 bits wide. The first word is routed to the debug bus lane 0; the second word is routed to the debug bus lane 2.

The TKM performance monitor bus is divided into four halfwords (16 bits wide). All signal groups in group A are presented to debug bus 0:15, all in group B are presented to debug bus 16:31, all in group C are presented to debug bus 64:79, and all in group D are presented to debug bus 80:95.

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 36, 38:41 = ‘1 1 1111’ | rmw bits 11, 37, 46:49 = ‘1 1 1111’ |
| | | rmw bits 12:16 =<br>‘00000’: to disable signal group Ax<br>‘10000’: to select signal group A0<br>‘01000’: to select signal group A1<br>‘00100’: to select signal group A2 | rmw bits 22:26 =<br>‘00000’: to disable signal group Cx<br>‘10000’: to select signal group C0<br>‘01000’: to select signal group C1<br>‘00100’: to select signal group C2<br>‘00010’: to select signal group C3 |
| | | rmw bits 17:21 =<br>‘00000’: to disable signal group Bx<br>‘10000’: to select signal group B0<br>‘01000’: to select signal group B1<br>‘00100’: to select signal group B2 | rmw bits 27:31 =<br>‘00000’: to disable signal group Dx<br>‘10000’: to select signal group D0<br>‘01000’: to select signal group D1<br>‘00100’: to select signal group D2 |

#### C.6.5.1 *EIB Token Manager - Group A0 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 0[0:15] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 12:16, 36, 38:41 = ‘1 10000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 0 | S | An even extreme data rate input and output (XIO) token was unused by resource allocation group (RAG) 0. |
| 1 | S | An odd XIO token was unused by RAG 0. |
| 2 | S | An even bank token was unused by RAG 0. |
| 3 | S | An odd bank token was unused by RAG 0. |
| 8 | S | A token was granted for SPE 0. |
| 9 | S | A token was granted for SPE 1. |
| 10 | S | A token was granted for SPE 2. |
| 11 | S | A token was granted for SPE 3. |
| 12 | S | A token was granted for SPE 4. |
| 13 | S | A token was granted for SPE 5. |
| 14 | S | A token was granted for SPE 6. |
| 15 | S | A token was granted for SPE 7. |

### C.6.5.2 *EIB Token Manager - Group A1 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 0[0:15] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 12:16, 36, 38:41 = ‘1 01000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 0 | S | An even XIO token was wasted by RAG 0. This is valid only when Unused Enable (UE) = 1 in *TKM_CR* register. |
| 1 | S | An odd XIO token was wasted by RAG 0. This is valid only when UE = 1 in *TKM_CR* register. |
| 2 | S | An even bank token was wasted by RAG 0. This is valid only when UE = 1 in *TKM_CR* register. |
| 3 | S | An odd bank token was wasted by RAG 0. This is valid only when UE = 1 in *TKM_CR* register. |
| 12 | S | An even XIO token was wasted by RAG U. |
| 13 | S | An odd XIO token was wasted by RAG U. |
| 14 | S | An even bank token was wasted by RAG U. |
| 15 | S | An odd bank token was wasted by RAG U. |

### C.6.5.3 *EIB Token Manager - Group A2 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 0[0:15] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 12:16, 36, 38:41 = ‘1 00100 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 0 | S | An even XIO token from RAG 0 was shared with RAG 1. |
| 1 | S | An even XIO token from RAG 0 was shared with RAG 2. |
| 2 | S | An even XIO token from RAG 0 was shared with RAG 3. |
| 3 | S | An odd XIO token from RAG 0 was shared with RAG 1. |
| 4 | S | An odd XIO token from RAG 0 was shared with RAG 2. |
| 5 | S | An odd XIO token from RAG 0 was shared with RAG 3. |
| 6 | S | An even bank token from RAG 0 was shared with RAG 1. |
| 7 | S | An even bank token from RAG 0 was shared with RAG 2. |
| 8 | S | An even bank token from RAG 0 was shared with RAG 3. |
| 9 | S | An odd bank token from RAG 0 was shared with RAG 1. |
| 10 | S | An odd bank token from RAG 0 was shared with RAG 2. |
| 11 | S | An odd bank token from RAG 0 was shared with RAG 3. |

### C.6.5.4 *EIB Token Manager - Group B0 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 0[16:31] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 17:21, 36, 38:41 = ‘1 10000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 16 | S | An even XIO token was unused by RAG 1. |
| 17 | S | An odd XIO token was unused by RAG 1. |
| 18 | S | An even bank token was unused by RAG 1. |
| 19 | S | An odd bank token was unused by RAG 1. |
| 25 | S | A token was granted for IOC0. |
| 26 | S | A token was granted for IOC1. |

### C.6.5.5 *EIB Token Manager - Group B1 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 0[16:31] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 17:21, 36, 38:41 = ‘1 01000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 16 | S | An even XIO token was wasted by RAG 1. This is valid only when UE = 1 in TKM_CR. |
| 17 | S | An odd XIO token was wasted by RAG 1. This is valid only when UE = 1 in TKM_CR. |
| 18 | S | An even bank token was wasted by RAG 1. This is valid only when UE = 1 in TKM_CR. |
| 19 | S | An odd bank token was wasted by RAG 1. This is valid only when UE = 1 in TKM_CR. |

### C.6.5.6 *EIB Token Manager - Group B2 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 0[16:31] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 17:21, 36, 38:41 = ‘1 00100 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 16 | S | An even XIO token from RAG 1 was shared with RAG 0. |
| 17 | S | An even XIO token from RAG 1 was shared with RAG 2. |
| 18 | S | An even XIO token from RAG 1 was shared with RAG 3. |
| 19 | S | An odd XIO token from RAG 1 was shared with RAG 0. |
| 20 | S | An odd XIO token from RAG 1 was shared with RAG 2. |
| 21 | S | An odd XIO token from RAG 1 was shared with RAG 3. |
| 22 | S | An even bank token from RAG 1 was shared with RAG 0. |
| 23 | S | An even bank token from RAG 1 was shared with RAG 2. |
| 24 | S | An even bank token from RAG 1 was shared with RAG 3. |
| 25 | S | An odd bank token from RAG 1 was shared with RAG 0. |
| 26 | S | An odd bank token from RAG 1 was shared with RAG 2. |
| 27 | S | An odd bank token from RAG 1 was shared with RAG 3. |
| 28 | S | An even XIO token from RAG U was shared with RAG 1. |
| 29 | S | An odd XIO token from RAG U was shared with RAG 1. |
| 30 | S | An even bank token from RAG U was shared with RAG 1. |
| 31 | S | An odd bank token from RAG U was shared with RAG 1. |

### C.6.5.7 *EIB Token Manager - Group C0 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 2[64:79] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 22:26, 37, 46:49 = ‘1 10000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 64 | S | An even XIO token was unused by RAG 2. |
| 65 | S | An odd XIO token was unused by RAG 2. |
| 66 | S | An even bank token was unused by RAG 2. |
| 67 | S | An odd bank token was unused by RAG 2. |
| 68 | S | An IOIF0 in token was unused by RAG 0. |
| 69 | S | An IOIF0 out token was unused by RAG 0. |
| 70 | S | An IOIF1 in token was unused by RAG 0. |
| 71 | S | An IOIF1 out token was unused by RAG 0. |

### C.6.5.8 *EIB Token Manager - Group C1 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 2[64:79] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 22:26, 37, 46:49 = ‘1 01000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 64 | S | An even XIO token was wasted by RAG 2. This is valid only when UE = 1 in TKM_CR. |
| 65 | S | An odd XIO token was wasted by RAG 2. This is valid only when UE = 1 in TKM_CR. |
| 66 | S | An even bank token was wasted by RAG 2. This is valid only when UE = 1 in TKM_CR. |
| 67 | S | An odd bank token was wasted by RAG 2. This is valid only when UE = 1 in TKM_CR. |

### C.6.5.9 *EIB Token Manager - Group C2 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 2[64:79] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 22:26, 37, 46:49 = ‘1 00100 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 64 | S | An even XIO token from RAG 2 was shared with RAG 0. |
| 65 | S | An even XIO token from RAG 2 was shared with RAG 1. |
| 66 | S | An even XIO token from RAG 2 was shared with RAG 3. |
| 67 | S | An odd XIO token from RAG 2 was shared with RAG 0. |
| 68 | S | An odd XIO token from RAG 2 was shared with RAG 1. |
| 69 | S | An odd XIO token from RAG 2 was shared with RAG 3. |
| 70 | S | An even bank token from RAG 2 was shared with RAG 0. |
| 71 | S | An even bank token from RAG 2 was shared with RAG 1. |
| 72 | S | An even bank token from RAG 2 was shared with RAG 3. |
| 73 | S | An odd bank token from RAG 2 was shared with RAG 0. |
| 74 | S | An odd bank token from RAG 2 was shared with RAG 1. |
| 75 | S | An odd bank token from RAG 2 was shared with RAG 3. |

### C.6.5.10 *EIB Token Manager - Group C3 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 2[64:79] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 22:26, 37, 46:49 = ‘1 00010 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 64 | S | An IOIF0 in token was wasted by RAG 0. This is valid only when UE = 1 in TKM_CR. |
| 65 | S | An IOIF0 out token was wasted by RAG 0. This is valid only when UE = 1 in TKM_CR. |
| 66 | S | An IOIF1 in token was wasted by RAG 0. This is valid only when UE = 1 in TKM_CR. |
| 67 | S | An IOIF1 out token was wasted by RAG 0. This is valid only when UE = 1 in TKM_CR. |

### C.6.5.11 *EIB Token Manager - Group D0 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 2[80:95] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 27:31, 37, 46:49 = ‘1 10000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 80 | S | An even XIO token was unused by RAG 3. |
| 81 | S | An odd XIO token was unused by RAG 3. |
| 82 | S | An even bank token was unused by RAG 3. |
| 83 | S | An odd bank token was unused by RAG 3. |

### C.6.5.12 *EIB Token Manager - Group D1 (NClk/2)*

| Address Offset | Register Short Name | Enable Word 2[80:95] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 27:31, 37, 46:49 = ‘1 01000 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 80 | S | An even XIO token was wasted by RAG 3. This is valid only when UE = 1 in TKM_CR. |
| 81 | S | An odd XIO token was wasted by RAG 3. This is valid only when UE = 1 in TKM_CR. |
| 82 | S | An even bank token was wasted by RAG 3. This is valid only when UE = 1 in TKM_CR. |
| 83 | S | An odd bank token was wasted by RAG 3. This is valid only when UE = 1 in TKM_CR. |

### C.6.5.13 *EIB Token Manager - Group D2 (NCIk/2)*

| Address Offset | Register Short Name | Enable Word 2[80:95] |
|---|---|---|
| x‘50 AFF0’ | TKM_PMCR | rmw bits 11, 27:31, 37, 46:49 = ‘1 00100 1 1111’ |

| Bit | Type | Description |
|---|---|---|
| 80 | S | An even XIO token from RAG 3 was shared with RAG 0. |
| 81 | S | An even XIO token from RAG 3 was shared with RAG 1. |
| 82 | S | An even XIO token from RAG 3 was shared with RAG 2. |
| 83 | S | An odd XIO token from RAG 3 was shared with RAG 0. |
| 84 | S | An odd XIO token from RAG 3 was shared with RAG 1. |
| 85 | S | An odd XIO token from RAG 3 was shared with RAG 2. |
| 86 | S | An even bank token from RAG 3 was shared with RAG 0. |
| 87 | S | An even bank token from RAG 3 was shared with RAG 1. |
| 88 | S | An even bank token from RAG 3 was shared with RAG 2. |
| 89 | S | An odd bank token from RAG 3 was shared with RAG 0. |
| 90 | S | An odd bank token from RAG 3 was shared with RAG 1. |
| 91 | S | An odd bank token from RAG 3 was shared with RAG 2. |

## C.7 Memory Interface Controller (MIC) Signal Selection

The address offset of each setup register is relative to the BE_MMIO_Base address.

### C.7.1 MIC Group 1 (NClk/2)

| Address Offset | Register Short Name | Enable Words 0 and 1 | Enable Words 2 and 3 |
|---|---|---|---|
| x‘50 A238‘ | MIC_FIR_Debug | rmw bits 32:33 = ‘10’ | rmw bits 32:33 = ‘01’ |
| x‘50 A1C8’ | MIC_Ctl_Debug_1 | write bits 0:63 = x‘FFFF FFFF FFFF FFFF’ | |
| x‘50 A048’ | MIC_Ctl_Debug2 | rmw bits 0:31 = x‘FFFF FFFF’ | |
| x‘50 A228’ | MIC_DF_Debug_Yrac_1 | write bits 0:63 = x‘FFFF FFFF FFFF FFFF’ | |

| Bit | Type | Description |
|---|---|---|
| 9 | C | XIO1 - The read command queue is empty. |
| 10 | C | XIO1 - The write command queue is empty. |
| 12 | C | XIO1 - The read command queue is full. |
| 13 | S | XIO1 - The MIC responds with a retry for a read command because the read command queue is full. |
| 14 | C | XIO1 - The write command queue is full. |
| 15 | S | XIO1 - The MIC responds with a retry for a write command because the write command queue is full. |
| 34 | S | XIO1 - A read command is dispatched. This includes high-priority and fast-path reads (see **Note 1**). |
| 35 | S | XIO1 - A write command is dispatched (see **Note 1**). |
| 36 | S | XIO1 - A read-modify-write command (with data size < 16 bytes) is dispatched (see **Note 1**). |
| 37 | S | XIO1 - A refresh is dispatched (see **Note 1**). |
| 39 | S | XIO1 - A byte-masking write command (with data size ≥ 16 bytes) is dispatched (see **Note 1**). |
| 41 | S | XIO1 - A write command is dispatched after a read command was previously dispatched (see **Note 1**). |
| 42 | S | XIO1 - A read command is dispatched after a write command was previously dispatched (see **Note 1**). |

1. This performance monitor signal originates from the MiClk domain, which is synchronous with the extreme data rate interface. It is reliable only when the system operates at nominal clock frequencies (that is, MiClk ≤ NClk/2). In slow mode, for instance, when the NClk/2 is slower than the MiClk, the performance monitor event signal from the MiClk domain changes so frequently that it is not possible to capture and synchronize all MiClk data for the NClk representation.
2. The synergistic processor unit (SPU) event bit 0 cannot be used at the same time that this group is enabled for monitoring.

### C.7.2 MIC Group 2 (NClk/2)

| Address Offset | Register Short Name | Enable Words 0 and 1 | Enable Words 2 and 3 |
|---|---|---|---|
| x‘50 A238‘ | MIC_FIR_Debug | rmw bits 32:33 = ‘01’ | rmw bits 32:33 = ‘10’ |
| x‘50 A088’ | MIC_Ctl_Debug_0 | write bits 0:63 = x‘FFFF FFFF FFFF FFFF’ | |
| x‘50 A048’ | MIC_Ctl_Debug2 | rmw bits 0:31 = x‘FFFF FFFF’ | |
| x‘50 A220’ | MIC_DF_Debug_Yrac_0 | write bits 0:63 = x‘0000 0000 0000 0000’ | |

| Bit | Type | Description |
|---|---|---|
| 9 | C | XIO0 - The read command queue is empty. |
| 10 | C | XIO0 - The write command queue is empty. |
| 12 | C | XIO0 - The read command queue is full. |
| 13 | S | XIO0 - The MIC responds with a retry for a read command because the read command queue is full. |
| 14 | C | XIO0 - The write command queue is full. |
| 15 | S | XIO0 - The MIC responds with a retry for a write command because the write command queue is full. |
| 34 | S | XIO0 - A read command was dispatched. This includes high-priority and fast-path reads (see **Note 1**). |
| 35 | S | XIO0 - A write command was dispatched (see **Note 1**). |
| 36 | S | XIO0 - A read-modify-write command (with data size < 16 bytes) was dispatched (see **Note 1**). |
| 37 | S | XIO0 - A refresh was dispatched (see **Note 1**). |
| 41 | S | XIO0 - A write command was dispatched after a read command was previously dispatched (see **Note 1**). |
| 42 | S | XIO0 - A read command was dispatched after a write command was previously dispatched (see **Note 1**). |

1. This performance monitor signal originates from the MiClk domain. It is reliable only when the system operates at nominal clock frequencies (that is, MiClk ≤ NClk/2). In slow mode, for instance, when the NClk/2 is slower than the MiClk, the performance monitor event signal from the MiClk domain changes so frequently that it is not possible to capture and synchronize all MiClk data for the NClk representation.
2. The SPU event bit 1 cannot be used while this group is enabled for monitoring.

### C.7.3 MIC Group 3 (NClk/2)

| Address Offset | Register Short Name | Enable Words 0 and 1 | Enable Words 2 and 3 |
|---|---|---|---|
| x‘50 A238‘ | MIC_FIR_Debug | rmw bits 32:33 = ‘01’ | rmw bits 32:33 = ‘10’ |
| x‘50 A088’ | MIC_Ctl_Debug_0 | write bits 0:63 = x‘FDFF FFFF FFFF FFFF’ | |
| x‘50 A048’ | MIC_Ctl_Debug2 | rmw bits 0:31 = x‘FDFF FFFF’ | |
| x‘50 A220’ | MIC_DF_Debug_Yrac_0 | write bits 0:63 = x‘FFFF FFFF FFFF FFFF’ | |

| Bit | Type | Description |
|---|---|---|
| 35 | S | XIO0 - A write command is dispatched (see **Note 1**). |
| 36 | S | XIO0 - A read-modify-write command (of data size < 16 bytes) was dispatched (see **Note 1**). |
| 37 | S | XIO0 - A refresh was dispatched (see **Note 1**). |
| 39 | S | XIO0 - A byte-masking write command (of data size ≥ 16 bytes) was dispatched (see **Note 1**). |

1. This performance monitor signal originates from the MiClk domain. It is reliable only when the system operates at nominal clock frequencies (that is, MiClk ≤ NClk/2). In slow mode, for instance, when the NClk/2 is slower than the MiClk, the performance monitor event signal from the MiClk domain changes so frequently that it is not possible to capture and synchronize all MiClk data for the NClk representation.
2. The SPU event bit 1 cannot be used while this group is enabled for monitoring.

## C.8 Cell Broadband Engine Interface (BEI)

The address offset of each setup register is relative to the BE_MMIO_Base address.

The Cell Broadband Engine interface (BIF) controller (BIC) implements two input and output (I/O) links (IOIF0 and IOIF1), and these two instances make use of the same building blocks. Each BIC building block in IOIF0 and IOIF1 places performance monitor signals onto the same word (word 0 or word 2), but the link outputs can be swapped between word 0 and word 2 controlled by the IF0TRC0[9:10] for IOIF0 and IOIF1 as described in the following tables.

All I/O controller (IOC) signal groups, except group 3, place their signals on the same word (word 0 or word 2). Their link outputs can be swapped between word 0 and word 2, using the same scheme (shown below) as for the BIC groups. For IOC group 3, however, the selection of the output word is controlled by byte selectors in the IOC_DTB_Cfg0 register.

*IOIF0 Word Swapping*

| Address Offset | Register Short Name | **Typical**<br>IOIF0 Word 0 → BEI Word 0<br>IOIF0 Word 2 → BEI Word 2 | **Swapped**<br>IOIF0 Word 0 → BEI Word 2<br>IOIF0 Word 2 → BEI Word 0 |
|---|---|---|---|
| x‘51 1000’ | IF0TRC0 | rmw bits 9:10 = ‘00’ | rmw bits 9:10 = ‘11’ |

*IOIF1 Word Swapping*

| Address Offset | Register Short Name | **Typical**<br>IOIF1 Word 0 → BEI Word 0<br>IOIF1 Word 2 → BEI Word 2 | **Swapped**<br>IOIF1 Word 0 → BEI Word 2<br>IOIF1 Word 2 → BEI Word 0 |
|---|---|---|---|
| x‘51 1400’ | IF1TRC0 | rmw bits 9:10 = ‘00’ | rmw bits 9:10 = ‘11’ |

*IOIF0 Byte Enables* (applies to the data *after* the IOIF0 word swap)

| Address Offset | Register Short Name | BEI Word 0 | BEI Word 2 |
|---|---|---|---|
| x‘51 1000’ | IF0TRC0 | rmw bits 11:14 = ‘1111’ | rmw bits 1:4 = ‘1111’ |

*IOIF1 Byte Enables* (applies to the data *after* the IOIF1 word swap)

| Address Offset | Register Short Name | BEI Word 0 | BEI Word 2 |
|---|---|---|---|
| x‘51 1400’ | IF1TRC0 | rmw bits 1:4 = ‘1111’ | rmw bits 11:14 = ‘1111’ |

### C.8.1 BIF Controller - IOIF0 Word 0 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 2000‘ | IF0TRC2 | rmw bits 0, 24:28 = ‘1 01000’ | |
| x‘51 1040’ | IF0TRC1 | rmw bits 0, 4:7, 35 = ‘1 0100 1’ | |
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bit 0 = 1 | |
| x‘51 1400’ | IF1TRC0 | rmw bit 0 = 1 | |
| x‘51 1000’ | IF0TRC0 | rmw bits 0, 9:10, 11:14 = ‘1 00 1111’ | rmw bits 0, 1:4, 9:10 = ‘1 1111 11’ |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘51 2000‘ | IF0TRC2 | rmw bits 24:28 = ‘10000’ | |
| x‘51 1040’ | IF0TRC1 | rmw bit 4:7, 35 = ‘1000 0’ | |

| Bit | Type | Description |
|---|---|---|
| 12 | S | Type A data physical layer group (PLG). Does not include header-only or credit-only data PLGs. In I/O interface (IOIF) mode, this event counts I/O device read data. In BIF mode, this event counts all outbound data. |
| 13 | S | Type B data PLG. In IOIF mode, this event counts I/O device read data. In BIF mode, this event counts all outbound data. |
| 14 | S | Type A data PLG. Does not include header-only or credit-only PLGs. In IOIF mode, this event counts store data to I/O device. This event does not apply in BIF mode. |
| 15 | S | Type B data PLG. In IOIF mode, this event counts store data to an I/O device. This event does not apply in BIF mode. |
| 16 | S | Data PLG. This event does not include header-only or credit-only PLGs. |
| 17 | S | Command PLG (no credit-only PLG). In IOIF mode, this event counts I/O command or reply PLGs. In BIF mode, this event counts command/reflected command or snoop/combined responses. |
| 18 | S | Type A data transfer regardless of length. This event can also be used to count type A data header PLGs (but not credit-only PLGs). |
| 19 | S | Type B data transfer. |
| 20 | S | Command-credit-only command PLG in either IOIF or BIF mode. |
| 21 | S | A data-credit-only data PLG sent in either IOIF or BIF mode. |
| 22 | S | A nonnull envelope was sent (does not include long envelopes). |
| 24 | E | A null envelope was sent (see **Note 1**). |
| 25 | V | No valid data sent this cycle (see **Note 1**). |
| 26 | E | A normal envelope was sent (see **Note 1**). |
| 27 | E | A long envelope was sent (see **Note 1**). |
| 28 | V | A null PLG was inserted in an outgoing envelope. |
| 29 | C | An outbound envelope array is full. |

1. This signal comes from the BClk domain, which is synchronous with the Rambus FlexIO interface. Counting back-to-back cycles in the BClk domain might be inaccurate because the signal crosses into the NClk domain to reach the performance monitor counters. A performance monitor signal from the BClk domain is reliable only when the NClk domain operates at more than two times the BClk domain (that is, NClk is greater than 3.33 GHz).

### C.8.2 BIF Controller - IOIF1 Word 0 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 3000‘ | IF1TRC2 | rmw bits 0, 24:28 = ‘1 01000’ | |
| x‘51 1440’ | IF1TRC1 | rmw bits 0, 4:7, 35 = ‘1 0100 1’ | |
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bit 0 = 1 | |
| x‘51 1000’ | IF0TRC0 | rmw bit 0 = 1 | |
| x‘51 1400’ | IF1TRC0 | rmw bits 0, 1:4, 9:10 = ‘1 1111 00’ | rmw bits 0, 9:10, 11:14 = ‘1 11 1111’ |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘51 3000‘ | IF1TRC2 | rmw bits 24:28 = ‘10000’ | |
| x‘51 1440’ | IF1TRC1 | rmw bit 4:7, 35 = ‘1000 0’ | |

| Bit | Type | Description |
|---|---|---|
| 19 | S | Type B data transfer |

### C.8.3 BIF Controller - IOIF0 Word 2 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 2000’ | IF0TRC2 | rmw bits 0, 32:36 = ‘1 01000’ | |
| x‘51 1040’ | IF0TRC1 | rmw bits 0, 17:20, 25:27, 35 = ‘1 0100 010 1’ | |
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bit 0 = 1 | |
| x‘51 1400’ | IF1TRC0 | rmw bit 0 = 1 | |
| x‘51 1000’ | IF0TRC0 | rmw bits 0, 9:10, 11:14, 43:45 = ‘1 11 1111 010’ | rmw bits 0, 1:4, 9:10, 43:45 = ‘1 1111 00 010’ |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘51 2000’ | IF0TRC2 | rmw bits 32:36 = ‘10000’ | |
| x‘51 1040’ | IF0TRC1 | rmw bits 17:20, 25:27, 35 = ‘1000 100 0’ | |
| x‘51 1000’ | IF0TRC0 | rmw bits 43:45 = ‘100’ | |

| Bit | Type | Description |
|---|---|---|
| 1 | E | A null envelope was received (see **Note 1**). |
| 14 | S | Command PLG, but not credit-only PLG. In IOIF mode, this event counts I/O commands or reply PLGs. In BIF mode, this event counts command/reflected command or snoop/combined responses. |
| 15 | S | Command-credit-only command PLG. |
| 20 | E | A normal envelope received is good (see **Note 1**). |
| 21 | E | A long envelope received is good (see **Note 1**). |
| 22 | S | Data-credit-only data PLG in either IOIF or BIF mode (counts a maximum of one per envelope) (see **Note 1**). |
| 23 | S | Nonnull envelope (does not include long envelopes; includes retried envelopes) (see **Note 1**). |
| 24 | S | A data grant was received. |
| 28 | S | Data PLG. This does not include header-only or credit-only PLGs. |
| 29 | S | Type A data transfer regardless of length. This can also be used to count type A data header PLGs (but not credit-only PLGs). |
| 30 | S | Type B data transfer. |

1. This signal comes from the BClk domain. Counting back-to-back cycles in the BClk domain can be inaccurate because the signal crosses into the NClk domain to reach the performance monitor counters. A performance monitor signal from the BClk domain is reliable only when the NClk domain operates at more than two times the BClk domain (that is, NClk is greater than 3.33 GHz).

### C.8.4 BIF Controller - IOIF1 Word 2 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘51 3000‘ | IF1TRC2 | rmw bits 0, 32:36 = ‘1 01000’ | |
| x‘51 1440’ | IF1TRC1 | rmw bits 0, 17:20, 25:27, 35 = ‘1 0100 010 1’ | |
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bit 0 = 1 | |
| x‘51 1000’ | IF0TRC0 | rmw bit 0 = 1 | |
| x‘51 1400’ | IF1TRC0 | rmw bits 0, 1:4, 9:10, 43:45 = ‘1 1111 11 010’ | rmw bits 0, 9:10, 11:14, 43:45 = ‘1 00 1111 010’ |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘51 3000‘ | IF1TRC2 | rmw bits 32:36 = ‘10000’ | |
| x‘51 1440’ | IF1TRC1 | rmw bits 17:20, 25:27, 35 = ‘1000 100 0’ | |
| x‘51 1400’ | IF1TRC0 | rmw bits 43:45 = ‘100’ | |

| Bit | Type | Description |
|---|---|---|
| 1 | E | A null envelope was received (see **Note 1**). |
| 14 | S | Command PLG (no credit-only PLG). This counts I/O command or reply PLGs. |
| 15 | S | Command-credit-only command PLG. |
| 20 | E | A normal envelope that was received is good (see **Note 1**). |
| 21 | E | A long envelope that was received is good (see **Note 1**). |
| 22 | S | A data-credit-only data PLG was received. This event counts a maximum of one per envelope. (see **Note 1**). |
| 23 | S | A nonnull envelope was received. This does not include long envelopes, but does includes retried envelopes (see **Note 1**). |
| 24 | S | A data grant was received. |
| 28 | S | A data PLG was received. This does not include header-only or credit-only PLGs. |
| 29 | S | Type A data transfer regardless of length. This event can also be used to count type A data header PLGs (but not credit-only PLGs). |
| 30 | S | A type B data transfer was received. |

1. This signal comes from the BClk domain. Counting back-to-back cycles in the BClk domain can sometimes be inaccurate because the signal crosses into the NClk domain to reach the performance monitor counters. A performance monitor signal from the BClk domain is reliable only when the NClk domain operates at more than two times the BClk domain (that is, NClk is greater than 3.33 GHz).

### C.8.5 I/O Controller Word 0 - Group 1 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bits 0,16:19 = ‘1 1111’ | |
| x‘51 1C70‘ | IOC_TopDbgCtl | rmw bits 24:26, 54:56, 61:63 = ‘100 000 000’ | |
| x‘51 1400’ | IF1TRC0 | rmw 0, 1:4, 9:10 = ‘1 1111 00’ | rmw bits 0, 9:10, 11:14 = ‘1 11 1111’ |
| x‘51 1000’ | IF0TRC0 | rmw bit 0 = 1 | |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘50 8600‘ | IOC_DTB_Cfg0 | rmw bits 16:19 = ‘0000’ | |
| x‘51 1C70’ | IOC_TopDbgCtl | rmw bits 24:26 = ‘000’ | |

| Bit | Type | Description |
|---|---|---|
| 8 | E | Received a memory-mapped input and output (MMIO) read targeted to IOIF1. |
| 9 | E | Received an MMIO write targeted to IOIF1. |
| 10 | E | Received an MMIO read targeted to IOIF0. |
| 11 | E | Received an MMIO write targeted to IOIF0. |
| 12 | S | A command was sent to IOIF0. |
| 13 | S | A command was sent to IOIF1. |

### C.8.6 I/O Controller Word 2 - Group 2 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bits 0, 24:27 = ‘1 1111’ | |
| x‘51 1C70‘ | IOC_TopDbgCtl | rmw bits 33:35 = ‘000’ | |
| x‘51 1400’ | IF1TRC0 | rmw bits 0, 1:4, 9:10 = ‘1 1111 11’ | rmw bits 0, 9:10, 11:14 = ‘1 00 1111’ |
| x‘51 1000’ | IF0TRC0 | rmw bit 0 = 1 | |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘50 8600‘ | IOC_DTB_Cfg0 | rmw bits 24:27 = ‘0000’ | |

| Bit | Type | Description |
|---|---|---|
| 5 | B | IOIF0 dependency matrix 3 is occupied by a dependent command (see **Note 1**). |
| 6 | B | IOIF0 dependency matrix 4 is occupied by a dependent command (see **Note 1**). |
| 7 | B | IOIF0 dependency matrix 5 is occupied by a dependent command (see **Note 1**). |
| 10 | E | Received a read request from IOIF0. |
| 11 | E | Received a write request from IOIF0. |
| 14 | E | Received an interrupt from the IOIF0. |

1. The dependency matrix (DM) entries are numbered 0-15. The IOC logic generally uses the lowest-numbered available slot for new incoming requests. Therefore, the lower numbered DM entries tend to be used more often than the higher-numbered entries.

### C.8.7 I/O Controller - Group 3 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bits 0, 40:43 = ‘1 1111’ | rmw bits 0, 32:35 = ‘1 1111’ |
| x‘50 8618’ | IOC_DTB_Cfg3 | rmw bit 47 = 1 | |
| x‘51 1C00’ | IOC_IOCmd_Cfg | rmw bit 17 = 1 | |
| x‘51 1400’ | IF1TRC0 | rmw bits 0, 9:10 = ‘1 00’ | |
| x‘51 1000’ | IF0TRC0 | rmw bits 0, 9:10 = ‘1 00’ | |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bits 32:35, 40:43 = ‘0000 0000’ | |
| x‘50 8618’ | IOC_DTB_Cfg3 | rmw bit 47 = 0 | |

| Bit | Type | Description |
|---|---|---|
| 16 | S | IOIF0 request for a token for even banks (0-14). This is valid only if resource allocation is enabled. |
| 17 | S | IOIF0 request for a token for odd banks (1-15). This is valid only if resource allocation is enabled. |
| 18 | S | IOIF0 request for token 1, 3, 5, or 7. This is valid only if resource allocation is enabled. |
| 19 | S | IOIF0 request for token 9, 11, 13, or 15. This is valid only if resource allocation is enabled. |
| 24 | S | IOIF0 request for a token 16. This is valid only if resource allocation is enabled. |
| 25 | S | IOIF0 request for a token 17. This is valid only if resource allocation is enabled. |
| 26 | S | IOIF0 request for a token 18. This is valid only if resource allocation is enabled. |
| 27 | S | IOIF0 request for a token 19. This is valid only if resource allocation is enabled. |

### C.8.8 I/O Controller Word 0 - Group 4 (NClk/2)

| Address Offset | Register Short Name | Enable Word 0 | Enable Word 2 |
|---|---|---|---|
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bits 0, 48:51 = ‘1 1111’ | |
| x‘51 1400’ | IF1TRC0 | rmw bit 0 = 1 | |
| x‘51 1000’ | IF0TRC0 | rmw bits 0, 9:10, 11:14 = ‘1 00 1111’ | rmw bits 0, 1:4, 9:10 = ‘1 1111 11’ |

| Address Offset | Register Short Name | Disable Word 0 | Disable Word 2 |
|---|---|---|---|
| x‘50 8600’ | IOC_DTB_Cfg0 | rmw bits 48:51 = ‘0000’ | |

| Bit | Type | Description |
|---|---|---|
| 0 | E | An I/O page table cache hit (for commands from IOIF). |
| 1 | E | An I/O page table cache miss (for commands from IOIF). |
| 3 | E | An I/O segment table cache hit. |
| 4 | E | An I/O segment table cache miss. |
| 24 | E | An interrupt was received from any synergistic processor unit (reflected command when internal interrupt controller (IIC) decides to send an acknowledge response). |
| 25 | E | IIC generated an interrupt to PowerPC processor unit (PPU) thread 0. |
| 26 | E | IIC generated an interrupt to PPU thread 1. |
| 27 | E | Received an external interrupt (using MMIO) from PPU to PPU thread 0. |
| 28 | E | Received an external interrupt (using MMIO) from PPU to PPU thread 1. |

# Glossary

**ABI** Application binary interface. The standard that a program follows to ensure that code generated by different compilers (and perhaps linking with various, third-party libraries) will run correctly on the Cell Broadband Engine Architecture (CBEA) processors. The ABI defines data types, register usage, calling conventions, and object formats. object formats.

**AC** (1) Address Compare bit. (2) Address Concentrator.

**aligned reference** A memory reference to data that resides at an address that is an integral multiple of the data size.

**AltiVec** The Freescale SIMD instruction set that is identical to IBM's vector/SIMD multimedia extension instruction set. AltiVec is a trademark of Freescale.

**ALU** Arithmetic logic unit.

**AOS** Array of structures. A method of organizing related data values. Also called the vector-across, or vec-across, form. Compare SOA

**API** Application programming interface.

**application** (1) A program designed to perform a specific task or group of tasks. (2) A user program, as opposed to a system or supervisor program; also called a problem state program in the PowerPC Architecture.

**architecture** A detailed specification of requirements for a processor or computer system. It does not specify details of how the processor or computer system must be implemented; instead it provides a template for a family of compatible implementations.

**ARPN** Abbreviated real page number.

**asynchronous** Not coordinated in time with the execution of instructions in an instruction pipeline.

**ATO** Atomic unit. Part of an SPE's MFC. It is used to synchronize with other processor units.

**atomic access** A bus access that attempts to be part of a atomic operation.

**atomic operation** A set of operations, such as read-write, that are performed as an uninterrupted unit. An atomic operation is a form of a mutual-exclusion (mutex) lock that can implement a semaphore.

**available to software** A region of memory that is not identified as allocated in Figure 5-1 CBEA Processor Memory Map on page 122, and that are known to be populated with memory chips or external MMIO registers, are available to software for any purpose.

**AVPN** Abbreviated virtual page number.

**B** (1) Byte. (2) Blocking.

| | |
|---|---|
| **b** | Bit. |
| **barrier** | A command that ensures program-ordering of memory accesses of all preceding, nonimmediate MFC commands with respect to all commands following the barrier command within the same MFC command queue. Compare fence. |
| **basic block** | A sequence of instructions that has no branching or other transfer of control. |
| **basic-block aggregation** | A method of extracting SIMD parallelism within a basic block by packing isomorphic computation on adjacent memory accesses into vector operations. |
| **BClk** | Bus interface controller (BIC) core clock. |
| **beat** | Data sent in one L2 clock cycle, which is two Cell Broadband Engine core-clock cycles. Also called data beat. |
| **BEI** | Cell Broadband Engine interface unit. |
| **BHT** | Branch history table. |
| **BIC** | Bus interface controller. Part of the Cell Broadband Engine interface (BEI) unit to I/O. |
| **BIF** | (1) Cell Broadband Engine interface protocol. The EIB's internal communication protocol. It supports coherent interconnection to other CBEA processors and BIF-compliant I/O devices, such as memory subsystems, switches, and bridge chips. The protocol is software-selectable only at power-on reset (POR). See IOIF. (2) Branch indirect if false instruction. |
| **big endian** | An ordering of bytes and bits in which the lowest-address byte is the most-significant (high) byte and lowest-numbered bit is the most-significant (high) bit. The CBEA processors support only big-endian ordering; they do not support little-endian ordering. |
| **BIU** | Bus interface unit. Part of the PPE's interface to the EIB. |
| **block** | (1) A cache line, which is 128 bytes in size. (2) The inability for an instruction in a pipeline to proceed. Blocking occurs at the instruction-dispatch stage and, in the PPE, stops only one of the two threads. Compare stall and bubble. |
| **blocking** | A property of certain SPE channels when its count is zero. Blocking channels cause the SPE to stall when the SPE reads or writes a channel with a count of zero. See channel count for more details. |
| **boundedly undefined** | Instruction-execution results that might have occurred by executing an arbitrary sequence of instructions, starting from a given machine state. |
| **branch-and-link** | A PPE branch instruction that writes the current instruction address plus 4 into the Link Register (LR). |

**branch hint** A type of branch instruction that provides a hint of the address of the branch instruction and the address of the target instruction. Hints are coded by the programmer or inserted by the compiler. The branch is assumed taken to the target. Hints are used in place of branch prediction in the SPU.

**BRU** Branch unit.

**bubble** An unused stage in a pipeline during a cycle. Compare stall and block.

**built-ins** The subset of generic intrinsics that map to assembly-language instruction sequences. A built-in cannot be expressed as a sequence of generic intrinsics, although it can be expressed as a sequence of specific intrinsics.

**bypassing the cache** Sending a cache line obtained from the L2 cache or main storage directly to the instruction pipeline, instead of first writing it into the L1 cache and then rereading it.

**C** (1) The C programming Language. (2) The Change bit in a PPE page-table entry.

**cache** High-speed memory close to a processor. A cache typically contains recently-accessed data or instructions, but certain cache-control instructions can lock, evict, or otherwise modify the caching of data or instructions.

**cacheable** An access, such as a load, store, or instruction fetch, in which the caching-inhibited (I) bit in the PTE is cleared to ‘0’.

**cache block** 128 bytes. Same as cache line.

**cache line** 128 bytes. Same as cache block.

**cache touch** A compiler-directed method, initiated by the dcbt and dcbtst instructions, of preloading data into the cache so that the subsequent accesses hit in the cache.

**caching-inhibited** A memory update policy in which caches are bypassed, and the load or store is performed to or from main storage. Only the storage location specified by the instruction (rather than a full cache line) is accessed at a caching-inhibited location. Stores to caching inhibited pages must update the memory hierarchy to a level that is visible to all processors and devices in the system. The operating system typically implements this policy, for example, for I/O devices. The “I” bit in the “WIMG” bits.

**callee** A program or routine that is called by another program or routine (the caller).

**caller** A program or routine that calls another program or routine (the callee).

**cast out** The process of writing modified data from a cache to the next-lower-level cache (for example, L1 to L2) or to main storage, and marking the cache line invalid.

| | |
|---|---|
| **CBEA** | See Cell Broadband Engine Architecture. The Cell/B.E. processor and the PowerXCell 8i processor are both implementations of the Cell Broadband Engine Architecture. |
| **CBEA processor** | A Cell/B.E. processor or a PowerXCell 8i processor. |
| **CBEA processor task** | A task running on the PPE and SPE. Each such task has one or more main threads and some number of SPE threads, and all the main threads within the task share the task's resources, including access to the SPE threads. |
| **CEC** | Central electronics complex. |
| **Cell/B.E.** | Cell Broadband Engine. |
| **Cell/B.E. processor** | The first implementation of the Cell Broadband Engine Architecture (CBEA). |
| **Cell Broadband Engine Architecture** | Extends the PowerPC 64-bit architecture with loosely coupled cooperative off-load processors. The Cell Broadband Engine Architecture provides a basis for the development of microprocessors targeted at the game, multimedia, and real-time market segments. The Cell/B.E. processor and the PowerXCell 8i processor are both implementations of the Cell Broadband Engine Architecture. |
| **CESOF** | An extension of PPE-ELF that allows PPE executable objects to contain SPE executables. The PPU and SPUs are supported by different compilers and different tools. CESOF allows a programmer to represent and resolve dependencies among PPU and SPU programs, which do not share a symbol space. |
| **channel** | A 32-bit, unidirectional communication interface between an SPE and the system, including main storage, the PPE, and other SPEs. |
| **channel count** | The value returned by the read-channel-count instruction (rchcnt). Reading the channel count of an implemented and blocking (queue-connected) channel returns the channel's available capacity. Reading the channel count of an implemented but nonblocking channel always returns a 1. Reading the channel count of an unimplemented (reserved) channel always returns a 0. |
| **CIDR** | Class ID Register. |
| **CIU** | Core interface unit. |
| **CL** | A class-ID parameter in an MFC command. |
| **classID** | See RClassID and TClassID. |

**coherence** The correct ordering of stores to a memory address, and the enforcement of any required cache write-backs during accesses to that memory address. Cache coherence is implemented by a hardware snoop (or inquire) method, which compares the memory addresses of a load request with all cached copies of the data at that address. If a cache contains a modified copy of the requested data, the modified data is written back to memory before the pending load request is serviced.

The primary objective of a coherent memory system is to provide the same image of memory to all devices in the system.

**coherence domain** All processor elements and all interfaces to main storage.

**coherence granularity** See coherency block.

**coherency block** The block size used in managing memory coherence. In the CBEA processors, it is 128 bytes.

**combined snoop response** The snoop responses from all snoopers are combined together to form the combined snoop response that is sent to all devices in the system.

**command** A result of executing certain instructions. For example, an MFC DMA command can result from an SPU's execution of several write channel (wrch) instructions that define the parameters of a DMA transfer.

**committed** The point in the execution of an instruction at which all older instructions in the pipeline are past the flush point (the point at which they can be written back to main storage).

An instruction is said to be committed when the process of recording and writing back its result has begun and cannot be prevented by an exception.

**common subexpression elimination** The replacement of repeated, redundant expression evaluations with a single computation assigned to a temporary variable.

**complete** An instruction is said to complete when its result is both available to another instruction and can be written back to memory (retired), and past the point at which the instruction can cause an exception.

**congruence class** (1) A set (or row) of a set-associative cache, including a TLB cache. Compare way. (2) All cache entries indexed by a specified effective address. (3) A cache index.

**constant folding** A compiler's precalculation of constant expressions.

**context** (1) The environment (for example, privilege and relocation) of instruction execution that is controlled by system registers (such as the Machine State Register and Storage Description Register 1) and the relevant page table. (2) The process or task environment of instruction execution.

**context switch** A process or task switch.

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| **context synchronizing** | The halting of instruction dispatch, clearing of the fetch buffer, and completion of all instructions currently in execution (that is, past the point at which they can cause an exception). The first instruction after a context-synchronizing event is fetched and executed in the context established by that instruction. Context synchronization occurs when certain instructions are executed (such as isync or rfi) or when certain events occur (such as an exception). All context-synchronizing events are also execution synchronizing. Compare execution synchronizing. |
| **control plane** | Refers to software or hardware that manages the operation of data plane software or hardware, by allocating resources, updating tables, handling errors, and so forth. See data plane. |
| **core clock** | The processor core clock (NClk). |
| **count register** | The PPE register that holds a loop count, which can be decremented during certain branch instructions, or which provides a branch target address. |
| **CPI** | Cycles per instruction. |
| **CPL** | Current priority level. |
| **CR** | Condition Register. |
| **CSA** | Context save area. |
| **CSI** | Context-synchronizing instruction. |
| **CTR** | Count Register. |
| **cumulative ordering** | The ordering of storage accesses performed by multiple sources. |
| **cycle** | Unless otherwise specified, one tick of the processor core clock (NClk), which is the frequency at which the PPE and SPEs run.run. |
| **DAR** | Data Address Register. |
| **data** | (1) Any information in memory or on the processor's data bus, including instruction operands and instruction opcodes. (2) Instruction operands, as opposed to instruction opcodes. |
| **data beat** | See beat. |
| **data hazard** | A situation in which an instruction has a data dependence or a name dependence on a prior instruction, and they occur close enough together in the instruction sequence that the processor might generate a result inconsistent with execution in program order. |
| **data plane** | Refers to software or hardware that operates on a stream or other large body of data and is managed by control plane software or hardware. See control plane. |
| **data stream** | A sequence of contiguous data cache blocks (cache lines). |

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| **DCache** | The PPE's L1 data cache. |
| **dcbf** | PPE data-cache-block flush instruction. |
| **dcbst** | PPE data-cache-block store instruction. |
| **dcbt** | PPE data-cache-block touch instruction. |
| **dcbtst** | PPE data-cache-block touch-for-store instruction. |
| **dcbz** | PPE data-cache-block set-to-zero instruction. |
| **DDR2** | Double double rate 2 synchronous dynamic random access memory (SDRAM). Transfers are performed on both the rising and falling edges of the clock. |
| **deadlock** | A state in which two elements in a process are stalled, each waiting for the other to respond. Compare livelock. |
| **DEC** | Decrementer. |
| **decrementer** | A register that counts down each time an event occurs. Each SPU contains dedicated 32-bit decrementers for scheduling or performance monitoring, by the program or by the SPU itself. |
| **demand fetch** | A request by the L1 cache to the L2 cache for a cache line that caused a an L1 cache miss. |
| **denormalized number** | A nonzero floating-point number whose exponent is the format's minimum, but represented as all zeros, and whose implicit leading significand bit is zero. |
| **dependence** | A relationship between two instructions that requires them to execute in program order. Dependence is a property of a program. |
| **D-ERAT** | Data ERAT, or the data-cache effective-to-real-address translation table. |
| **DERR** | Data error interrupt. |
| **device memory** | Memory that has both the caching-inhibited and guarded attributes. This is typical of memory-mapped I/O devices. |
| **DFQ** | Demand instruction request queue. |
| **DIQ** | Denormalized instruction queue. |
| **displacement** | An offset or index from a base address. |
| **DLQ** | Demand data load request queue. |
| **DMA** | Direct memory access. A technique for using a special-purpose controller to generate the source and destination addresses for a memory or I/O transfer. |
| **DMAC** | Direct memory access controller. A controller that performs DMA transfers. |

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| **DMA command** | A type of MFC command that transfers or controls the transfer of a memory location containing data or instructions. See MFC command. |
| **DMA list** | A list of list elements that, together with an initiating DMA list command, specifies a sequence of DMA transfers between a single area of LS and discontinuous areas in main storage. Such lists are stored in an SPE's LS, and the sequence of transfers is initiated with a DMA list command such as getl or putl. DMA list commands can only be issued by programs running on an SPE, but the PPE or other devices can create and store the lists in an SPE's LS. DMA lists can be used to implement scatter-gather functions between main storage and the LS. |
| **DMA list command** | A type of MFC command that initiates a sequence of DMA transfers specified by a DMA list stored in an SPE's LS. See DMA list. |
| **DMA queue** | A queue for holding DMA-transfer commands. The MFC contains two DMA queues—a 16-entry SPU command queue and an 8-entry proxy command queue. |
| **DMA transfer** | Byte moves by a DMAC, without regard to the numeric significance of any byte. |
| **dominant-shift reorganization** | A reorganization of misaligned data in which the offset that is dominant (most often used) among a set of streams is chosen as the offset to shift to. |
| **doubleword** | Eight bytes. |
| **DP** | Double-precision. |
| **DPFE** | Data prefetch engine. |
| **DR** | Data relocate. |
| **DSI** | Data storage interrupt. |
| **DSISR** | Data Storage Interrupt Status Register. |
| **DTS** | Digital thermal sensor. |
| **dual-issue** | Issuing two instructions at once, under certain conditions. See fetch group. |
| **dynamic branch prediction** | Methods in which hardware records the resolution of branches and uses this information to predict the resolution of a branch when it is encountered again. |
| **dynamic linking** | Linking of a program in which library procedures are not incorporated into the load module, but are dynamically loaded from their library each time the program is loaded. |
| **EA** | See effective address. |
| **eager-shift reorganization** | A reorganization of a misaligned load stream that shifts directly to the alignment of the store. |

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| **EAH** | An effective address high parameter in an MFC command. |
| **EAL** | An effective address low parameter in an MFC command. |
| **EAR** | Effective-address reference. A software structure in the .toe section of the CESOF format. |
| **ECC** | Error-correcting code. |
| **effective address** | An address generated or used by a program to reference memory. A memory-management unit translates an effective address (EA) to a virtual address (VA), which it then translates to a real address (RA) that accesses real (physical) memory. The maximum size of the effective-address space is 264 bytes. |
| **EIB** | Element interconnect bus. The on-chip coherent bus that handles communication between the PPE, SPEs, memory, and I/O devices (or a second CBEA processor). The EIB is organized as four unidirectional data rings (two clockwise and two counterclockwise). |
| **eieio** | Enforce in-order execution of I/O. |
| **ELF** | Executable and linking format. The standard object format for many UNIX® operating systems, including Linux. Originally defined by AT&T and placed in the public domain. Compilers generate ELF files. Linkers link to files with ELF files in libraries. Systems run ELF files. |
| **ELF header** | An *ELF* object-file header that is used to find a *section header* or a *program header*. |
| **endian orientation** | A view of bits and bytes in which either the little end (least-significant or low end) or the big end (most-significant or high end) is assigned the lowest number or address. Thus, there are two types of endian orientation—little-endian and big-endian. Endian orientation applies to bits, in the context of register-value interpretation, and to bytes, in the context of memory accesses. The CBEA processors support only big-endian ordering; they do not support little-endian ordering. |
| **ERAT** | Effective-to-real address translation, or a buffer or table that contains such translations, or a table entry that contains such a translation. |
| **ESID** | Effective segment ID |
| **even pipeline** | Part of an SPE's dual-issue execution pipeline. Also referred to as pipeline 0. |
| **exception** | An error, unusual condition, or external signal that might alter a status bit and will cause a corresponding interrupt, if the interrupt is enabled. See interrupt, imprecise interrupt, and precise interrupt. |

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| **execution synchronizing** | The completion of previously issued instructions—to a point where they can no longer cause an exception, but not necessarily in the context (privilege, protection, and address translation) under which they were issued—before continuing with program execution. Execution-synchronizing does not empty (flush) the instruction prefetch queue. The sync instruction is execution-synchronizing. Compare context synchronizing. |
| **external access** | From an SPE's viewpoint, a read or write access by any DMA transfer (by the local MFC or the MFC associated with another SPE) or any processor (including the PPE or another SPE) or other device—other than the SPE being referred to. |
| **fence** | A modifier to a DMA command that ensures that this command is ordered with respect to all preceding commands in the DMA command queue within the same tag group. Compare barrier. |
| **fetch** | To read instructions (but not data) from storage. |
| **fetch group** | A doubleword-aligned instruction pair. Dual-issue occurs when a fetch group has two instructions that are ready to issue, and when the first instruction can be issued on the even pipeline and the second instruction can be issued on the odd pipeline. |
| **FIFO** | First in, first out. Refers to one way elements in a queue are processed. It is analogous to "people standing in line." |
| **FIR** | Fault Isolation Register. |
| **fixed-point value** | An integer. |
| **FlexIO** | Rambus FlexIO bus. The physical-link I/O signals on the BIF and IOIF interfaces. See IOIF. |
| **FLIH** | First-level interrupt handler. |
| **flush** | (1) Write the contents of a cache line back to main storage, and invalidate the cache line. (2) Invalidate a cache line or an instruction pipeline. |
| **flush point** | The execution-pipeline point at which all instructions older than the flush point can be written back to main storage. |
| **flush transaction** | A transaction caused by a dcbf that hits on a memory-coherent cache line and is marked shared or invalid. It is sent to other snoopers who might have a copy of the line. |
| **FP** | Floating point. |
| **FPR** | Floating-Point Register. |
| **FPSCR** | Floating-Point Status and Control Register. |
| **FPU** | Floating-point unit. |
| **fres** | Floating reciprocal estimate single A form instruction. |

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| **FXU** | Fixed-point unit. In the PPE, the fixed-point integer unit. |
| **G** | The guarded bit in a page table entry which controls the processor's accesses to cache and main storage. It is part of the “WIMG” bits. |
| **GCD** | The greatest common divisor of two integers. It is the largest integer that divides them both. This is typically denoted by GCD(a,b). |
| **general-purpose registers** | The registers used for integer and string operations. |
| **generic intrinsics** | Intrinsics that map to one of several assembly-language instructions or instruction sequences, depending on the type of operands. (See intrinsic.) All generic SPU intrinsics are prefaced by the string, spu_. For example, the generic intrinsic that implements the stop assembly instruction is named spu_stop. |
| **getllar** | Get lock line and reserve MFC command. |
| **get_ps** | An operating system call that returns a base address for the problem state area. |
| **GPR** | See general-purpose register. |
| **guarded** | Regions of memory in which prefetching and other speculative storage operations are not permitted. A data access to a guarded location is performed only if either (a) the access is caused by an instruction that is required for sequential execution, or (b) the access is a load and the storage location is already in a cache. Such regions in memory are marked by setting the G bit in the relevant page-table entry. The operating system typically implements guarding, for example, on I/O devices. |
| **H** | Hash function identifier. |
| **harvest** | To reset (clear) the state of an SPE. |
| **hashed page table** | A variable-sized, chained data structure that maps virtual page numbers (VPNs) and real page numbers (RPNs). A hash function is applied to a VPN to yield an entry into the page table. Each entry in a hashed page table is associated with a page in real memory. Also called inverted page table. |
| **hazard** | A situation in which the overlapped or out-of-order execution of a pair of instructions might generate a result inconsistent with execution of the instructions in program order. A hazard is a property of a program. |
| **hcall** | Hypervisor call, which is an sc instruction with the LEV bit set to ‘1’. |
| **HDEC** | Hypervisor decrementer. |
| **HID** | Hardware-implementation dependent. |

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| **high** | The most-significant bit or byte numbers in a field, register or memory. High bits and bytes are the lowest-numbered bits or bytes in a data structure. For example, bit 0 is the high bit of a bit field or register. |
| **hint stall** | The holding of a branch instruction in or before the stage of the pipeline in which triggering occurs. The branch is held there until an associated hint-for branch trigger is loaded. |
| **hoist** | To move an instruction to an earlier point in the program-execution order. |
| **HRMOR** | Hypervisor Real-Mode Offset Register. |
| **HTAB** | Hashed page table. |
| **hypervisor** | A control (or virtualization) layer between hardware and the operating system. It allocates resources, reserves resources, and protects resources among (for example) sets of SPEs that might be running under different operating systems.<br><br>The CBEA processors have three operating modes: user, supervisor, and hypervisor. The hypervisor performs a meta-supervisor role that allows multiple independent supervisors' software to run on the same hardware platform.<br><br>For example, the hypervisor allows both a real-time operating system and a traditional operating system to run on a single PPE. The PPE can then operate a subset of the SPEs in the CBEA processors with a real-time operating system, while the other SPEs run under the traditional operating system. |
| **hypervisor call** | An sc instruction with the LEV bit set to '1'. |
| **I** | The caching-inhibited bit in a page table entry which controls the processor's accesses to cache and main storage. It is part of the "WIMG" bits. |
| **IBuf** | Instruction buffer. |
| **ICache** | The PPE's L1 instruction cache. |
| **icbi** | PPE instruction-cache-block invalidate instruction. |
| **ID** | Instruction dispatch. |
| **IEEE 754** | The IEEE 754 floating-point standard. A standard written by the Institute of Electrical and Electronics Engineers that defines operations and representations of binary floating-point arithmetic. |
| **I-ERAT** | Instruction ERAT, or instruction-cache effective-to-real-address translation table. |
| **IFAR** | Instruction Fetch Address Register. |

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| **IGP** | Interrupt generation port. One of two registers, IIC_IGP0 and IIC_IGP1 (one for each PPE thread), to which software on the PPE or an SPU can write an Interprocessor Interrupt (IPI). |
| **IIC** | Internal interrupt controller. |
| **ILP** | Instruction-level parallelism. |
| **ILP32** | 32-bit Integers, Longs, and Pointers. The SPEs are always ILP32, with DMA support for 32-bit and 64-bit effective-address pointers |
| **immediate operand** | An operand included in an instruction. Also called immediate constant or immediate value. |
| **imprecise interrupt** | A synchronous exception that does not adhere to the precise exception model. Imprecise interrupts occur one or more instructions after the execution of the instruction causing the interrupt. They are not restartable. In the CBEA processors, single-precision floating-point operations generate imprecise exceptions. See precise interrupt. |
| **inbound access** | An access initiated externally from an I/O device (also called I/O unit) to the CBEA processors. Also called an I/O access. Compare outbound access. |
| **index** | An offset from a base address. |
| **indirect** | An access is said to be indirect if a register holds its target. For example, an indirect branch is one whose target is specified in a register. |
| **inline expansion** | An optimization in which the reference to a procedure is replaced with the code of the procedure itself to eliminate calling overhead. |
| **in-order** | In program order. The PPE and SPEs execute instructions in-order; that is, they do not rearrange them (out-of-order). |
| **instruction runout** | A situation in which there are no instructions remaining in the SPU instruction pipeline, and the SPU becomes idle. |
| **instruction-sequencing error** | An error that occurs when instructions are executed in a different order other than expected. The SPU expects instructions to be executed sequentially unless modified by hint trigger. |
| **interrupt** | A change in machine state in response to an exception. See exception, imprecise interrupt, and precise interrupt. |
| **interrupt packet** | Used to signal an interrupt, typically to a processor or to another interruptible device. |
| **interrupt vector** | The starting address of an interrupt handler or service routine. |
| **intervention** | The peer-to-peer movement of data between two caches, without involving main storage. |

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| **intrinsic** | An in-line set of assembly-language instructions in the form of function call that is built into a compiler. Intrinsics make the underlying ISA accessible from a high-level programming language. |
| **inverted page table** | See hashed page table. |
| **I/O access** | (1) An I/O operation that is initiated internally or externally to the CBEA processor. (2) As used in Section 7.7 I/O Storage Model on page 188, an I/O operation that is initiated externally to the CBEA processor. |
| **I/O address** | As used in Section 7.7 I/O Storage Model on page 188, an address passed from an I/O device to the CBEA processor when the I/O device attempts to perform an access to the CBEA processor real-address space. |
| **IOC** | I/O interface controller. |
| **I/O device** | Input/output device. From software's viewpoint, I/O devices exist as memory-mapped registers that are accessed in main-storage space by load/store instructions. The operating system typically configures access to I/O devices as caching-inhibited and guarded. |
| **IOID** | I/O identifier. It is described in the IOPT and identifies an I/O unit. |
| **IOIF** | (1) One of two I/O interfaces supported by the EIB. Also called a FlexIO interface. (2) A noncoherent communication protocol used by I/O devices attached to an IOIF, which differs from the coherent BIF protocol. The IOIF protocol can be used on either IOIF0 or IOIF1; the BIF protocol can be used only on IOIF0. The protocol is software-selectable only at power-on reset (POR). |
| **IOIF0, IOIF1** | One of two I/O (IOIF) interfaces, also called a FlexIO interfaces. |
| **IOIF device** | A device that is connected directly to the CBEA processor IOIF port. |
| **IOIF protocol** | The EIB's noncoherent protocol for interconnection to I/O devices. See BIF. |
| **I/O operation** | A storage operation that crosses a Cell Broadband Engine coherence-domain boundary. |
| **IOPT** | I/O page table. |
| **IOST** | I/O segment table. |
| **I/O unit** | One or more physical I/O devices, I/O bridges, or other functional units attached to an IOIF, in which one value of the IOID described in the IOPT. A functional unit that can initiate I/O accesses. |
| **IPC** | Instructions per cycle. |
| **IPFQ** | Instruction prefetch request queue. |

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| **IPI** | Interprocessor Interrupt. In the CBEA processors, a software interrupt written by a PPE thread or an SPE to one of the Interrupt Generation Port (IGP) registers. |
| **IPP** | Interrupt pending port. |
| **IR** | Instruction relocate. |
| **IS** | (1) Instruction issue. (2) Invalidation selector. |
| **ISEG** | Instruction segment exception. |
| **ISI** | Instruction storage interrupt. |
| **island** | A separate logic unit that the performance monitor facility can monitor (see Section C.1 Selecting Performance Monitor Signals on the CBEA Processor Debug Bus on page 795 for more information). |
| **ISRC** | Interrupt source. |
| **IU** | Instruction unit. |
| **JSRE** | Joint Software Reference Environment |
| **K** | 210 (as in KB for 1024 bytes). |
| **KB** | Kilobyte. |
| **key** | A value, ‘0’ or ‘1’, used in conjunction with the page protection (PP) bits in page-table entries to determine access rights to locations in main storage. The value of the key is equal to the value of the Ks bit or value of the Kp bit in an SLB entry, depending on the PR bit in the machine state register (MSR). |
| **Kp** | The problem-state (user-mode) storage key bit in SLB entries. In address translations, the Kp bit is used in conjunction with the PR, PP, N, and Ks bits to determine access privilege. |
| **Ks** | The supervisor-state storage key bit in SLB entries. In address translations, the Ks bit is used in conjunction with the PR, PP, N, and Ks bits to determine access privilege. |
| **L** | Large page indicator. |
| **L1** | Level-1 cache memory. The closest cache to a processor, measured in access time. |
| **L2** | Level-2 cache memory. The second-closest cache to a processor, measured in access time. An L2 cache is typically larger than an L1 cache. |
| **LAR** | Load and reserve. |
| **LARX-reserve transaction** | An address-only transaction that sets the reservation for every cache level below the level serviced by a read atomic operation. |

**latency** (1) The number of cycles between when an instruction begins execution to when the result is available. (2) The issue-to-issue latency for a dependent instruction. (3) The time between when a function is called and when it returns.

**lazy-shift reorganization** A reorganization of misaligned data that delays eager-shift reorganization as long as possible.

**ld** PPE load doubleword instruction.

**ldarx** PPE load-doubleword-and-reserve indexed instruction.

**LEAL** List-element effective-address low. A parameter in a DMA-list MFC command.

**least-recently used** An algorithm for replacing cache entries to identifies or approximates the oldest entry so that it can be cast out (evicted) to make room for a new entry.

**least-significant** The highest-numbered bits or bytes (big-endian) in a data structure.

**length conversion** A conversion between data of different lengths.

**linker** A program that resolves cross-references between separately compiled or assembled object modules and then assigns final addresses to create a single relocatable load module. If a single object module is linked, the linker simply makes it relocatable. Also called a link editor.

**link register** The PPE register used to provide a branch-target address and to hold the return address for certain branch instructions.

**list element** Also called transfer element. See DMA list.

**livelock** A state in which processors interact in a way such that no processor makes forward progress, such as in an endless loop of program execution. Compare deadlock. In a livelock, processing continues to take place. In a deadlock, no processing continues.

**lmw** PPE load multiple word instruction.

**lnop** A nop in an SPU's odd pipeline. It can be inserted in code to align for dual issue of subsequent instructions.

**load combining** The combining of multiple load accesses to storage into a single load access.

**loader** A program that reads the load module into main storage, performing all necessary dynamic linking so that the module can execute.

**local storage** The 256 KB on-chip memory associated with each SPE. It holds both instructions and data. Abbreviated as LS.

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| **lock** | (1) To use a replacement management table (RMT) entry to prevent L2-cache or TLB entries from being cast out. (2) To obtain exclusive access to a main-storage location using the lwarx and stwcx. instructions. (3) See pinned. |
| **logical partition** | A virtual machine managed by a hypervisor. |
| **logical PPE** | One of the two PPE threads. |
| **loop distribution** | The splitting off of those parts of a loop that cannot be SIMDized |
| **looped-back operation** | A command that is received by the Cell Broadband Engine from an IOIF device and routed back out the same IOIF. |
| **loop interchange** | The swapping of an inner (nested) loop with an outer loop. This is done when the inner loop appears to be the candidate for SIMDization. |
| **loop-level aggregation** | A method of extracting SIMD parallelism across loop iterations. Computations on stride-one accesses across loop iterations can be combined into vector operations. |
| **loop unrolling** | A programming optimization that increases the step of a loop, and duplicates the expressions within a loop to reflect the increase in the step. This can improve instruction scheduling and memory access time. |
| **loop versioning** | A technique for hoisting an individual array-index exception-check outside a loop by providing two copies of the loop: the safe loop and the unsafe (original) loop. |
| **low** | The least-significant bit or byte numbers in a field, register or memory. Low bits and bytes are the highest-numbered bits or bytes (big-endian) in a data structure. For example, bit 63 is the low bit of a doubleword field or register. |
| **LP** | Large page selector. |
| **LP64** | 64-bit longs and pointers. 32-bit integers. |
| **LPAR** | Logical partition. |
| **LPCR** | Logical Partition Control Register. |
| **LPES** | Logical Partitioning Environment Selector. A field in the LPCR register. |
| **LPID** | Logical partition ID. |
| **LR** | Link Register. |
| **LRU** | Least recently used. See least-recently used. |
| **LS** | See local storage. |
| **LSA** | Local storage address. |
| **LSCSA** | Local storage context save area, |

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| **LSU** | Load and store unit. |
| **lswi** | PPE load string word immediate instruction. |
| **lswx** | PPE load string word indexed instruction. |
| **LSZ** | List size. A parameter in an MFC command. |
| **LTS** | List-element transfer size. A parameter in an MFC command. |
| **lwarx** | PPE load-word-and-reserve indexed instruction. |
| **M** | (1) The memory-coherence bit in a page table entry which controls the processor's accesses to cache and main storage. It is part of the “WIMG” bits. (2) 220, as in MB for 1,048,576 bytes. |
| **m** | The number of bits in a real address. For the CBEA processors, this value is 42. |
| **mailbox** | A queue in an SPE’s MFC for exchanging 32-bit messages between the SPE and the PPE or other devices. Two mailboxes (the SPU Write Outbound Mailbox and SPU Write Outbound Interrupt Mailbox) are provided for sending messages from the SPE. One mailbox (the SPU Read Inbound Mailbox) is provided for sending messages to the SPE. |
| **main memory** | See main storage. |
| **main storage** | (1) The effective-address (EA) space. It consists physically of physical memory (whatever is external to the memory-interface controller), SPU LSs, memory-mapped registers and arrays, memory-mapped I/O devices (all I/O is memory-mapped), and pages of virtual memory that reside on disk. It does not include caches or execution-unit register files. (2) The level of storage hierarchy in which all storage state is visible to all processors and mechanisms in the system. See storage |
| **mangle** | To add an extra prefix or suffix to a global symbol that is produced by a compiler. Global symbols produced by a compiler are said to be mangled if a symbol has any extra prefix or suffix added by the tool chain. |
| **MB** | Megabyte. |
| **MC** | Microcode. |
| **memory-mapped** | Mapped into the Cell Broadband Engine’s addressable-memory space. Registers, SPE local storages (LSs), I/O devices, and other readable or writable storage can be memory-mapped. Privileged software does the mapping. |
| **memory stream** | A sequence of contiguous memory locations that are accessed by a memory reference throughout the lifetime of a loop. |
| **MERSI** | The Modified (M), Exclusive (E), Recent (R), Shared (S), and Invalid (I) cache states. The CBEA processors use the MERSI protocol, plus two additional states: Unsolicited Modified (Mu) and Tagged (T). |

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| **MFC** | Memory flow controller. It is part of an SPE and provides two main functions: moves data via DMA between the SPE's local storage (LS) and main storage, and synchronizes the SPU with the rest of the processing units in the system. |
| **MFC command** | A command issued to an SPE's MFC. The command and any related parameters are written with a series of instructions executed either on the MFC's associated SPU or on the PPE or other device. MFC commands provide the main mechanism by which software running on an SPU accesses main storage and maintains synchronization with other devices in the system. DMA commands are a primary type of MFC command. |
| **mfceieio** | MFC enforce in-order execution of I/O command. |
| **MFC pause** | MFC Pause state. One of six power-management states. |
| **MFC proxy commands** | MFC commands issued using the MMIO interface. |
| **mfcsync** | MFC synchronize command. |
| **mfspr** | PPE move from special-purpose register instruction. |
| **MIC** | Memory interface controller. The CBEA processor MIC supports two memory channels. |
| **MiClk** | MIC core clock. |
| **microarchitecture** | A microprocessor's hardware architecture. |
| **miss penalty** | The time required to fill a cache line after a cache miss. |
| **mixed-mode SIMDization** | The devirtualization of virtual vectors and their replacement by a single SIMD instruction, a sequence of SIMD instructions, a library call, or a sequence of scalar operations. |
| **MMIO** | Memory-mapped input/output. The documentation for the CBEA processors defines an "MMIO register" as any internal or external register that is accessed through the main-storage space with load and store instructions, whether or not the register is associated with an I/O device. See memory-mapped. |
| **MMU** | Memory management unit. A functional unit that translates between effective addresses (EAs) used by programs and real addresses (RAs) used by physical memory. The MMU also provides protection mechanisms and other functions. |
| **M:N thread model** | A programming model in which M user threads are mapped to N kernel threads (or virtual processors). |
| **MOD** | Modified snoop response code. |
| **MODINTV** | Modified-intervention snoop response code. |
| **most-significant** | The lowest-numbered bits or bytes (big-endian) in a data structure. |

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| **MPI** | Message passing interface. A standard for high-performance communication on massively parallel architectures and clustered distributed-memory systems. |
| **MSb** | Most-significant bit. |
| **MSR** | Machine State Register. |
| **MT** | Multithreading. See multithreading. |
| **mtmsr** | PPE move to machine state register instruction. |
| **mtspr** | PPE move to special-purpose register instruction. |
| **multithreading** | Simultaneous execution of more than one program thread. It is implemented by sharing one software process and set of execution resources but duplicating the architectural state (registers, program counter, flags, and so forth.) of each thread. |
| **mutex lock** | Mutual-exclusion lock. An atomic operation used to implement a semaphore. |
| **N** | (1) No Execute bit in SLB entries. (2) Intervention bit. (3) Nonblocking. (4) Normal run thermal-management state. |
| **n** | The number of bits in a virtual address. For the CBEA processors, this value is 65. |
| **NaN** | Not-a-number. A special string of bits encoded according to the IEEE 754 Floating-Point Standard. A NaN is the proper result for certain arithmetic operations; for example, 0/0 = NaN. There are two types of NaNs, quiet NaNs and signaling NaNs. Only the signaling NaN raises a floating-point exception when it is generated. |
| **NClk** | Core clock. The clock for the PPU and SPU. It is the highest-frequency processor clock. |
| **NCU** | Noncacheable unit. |
| **no-op** | See nop. |
| **nop** | No operation. Also called no-op. |
| **NPC** | An SPE's Next Program Counter Register. |
| **NUMA** | Nonuniform memory access. |
| **object module** | The output file of a compiler or other language translator. It includes the machine language translation and other information for symbolic binding and relocation. |
| **odd pipeline** | Part of an SPE's dual-issue execution pipeline. Also referred to as pipeline 1. |
| **offset** | An index that is added to a base address. |

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| **OGSA** | Open Grid Services Architecture. A standard that enables communication across heterogeneous, geographically dispersed environments. |
| **OpenMP** | Open Multiprocessing. An API that supports multiplatform, shared-memory parallel programming. |
| **ordered** | Said of an exception that only generates an interrupt in a prioritized order, with respect to the state of the interrupt-processing mechanism. |
| **OS** | Operating system. |
| **outbound access** | An access initiated internally from the CBEA processor to an I/O device (also called I/O unit). Compare inbound access. |
| **out of order** | Not in program order. |
| **overlay** | SPU code that is dynamically loaded and executed by a running SPU program. |
| **p** | Page size. A power-of-2 variable representing the size of a page. Three concurrent page sizes can be used: 4 KB (p = 12) and any two of the following large page sizes: 64 KB (p = 16), 1 MB (p = 20), 16 MB (p = 24). |
| **page** | A unit of main-memory storage. Each page can have independent protection and control attributes, and Change and Reference status bits can be independently recorded. See page table. |
| **page fault** | A restartable interrupt that causes the loading of a page from disk to memory. |
| **page-history recording** | Same as storage-access recording. |
| **page table** | A table that maps pages of virtual addresses (VAs) to real addresses (RA) and contains related protection parameters and other information about memory locations. See hashed page table. |
| **page-table entry group** | A 64-byte data structure in the hashed page table that contains eight page-table entries. |
| **parity** | A means of checking data reliability in which data bits are concatenated with a parity bit whose value makes the total number of ‘1’ bits even or odd. |
| **path length** | The number of instructions in an instruction sequence. |
| **PC** | (1) Program counter. The PPE maintains separate program counters for each thread. (2) Personal computer. |
| **performance simulation** | Simulation by the IBM Full System Simulator for the Cell Broadband Engine in which both the functional behavior of operations and the time required to perform the operations is simulated. Also called cycle-accurate simulation. |

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| **pervasive logic** | Logic that provides power management, thermal management, clock control, software-performance monitoring, trace analysis, and so forth. |
| **physical address** | (1) A PowerPC Architecture real address (RA). (2) An address on the CBEA processor address-bus signals. |
| **pinned** | A set of virtual-memory pages is said to be pinned or locked if the operating system has forced them resident in memory, so that they are not swapped out to disk. In addition to pages of memory being locked, the TLB entries associated with those pages can also be locked, thus preventing TLB misses. |
| **pipelined hint mode** | An SPU pipelined mode in which a hint-for branch stall is negated and no further hint stall is generated for the duration of the pipelined hint mode. Pipelined hint mode lasts for 16 cycles after the last stallable hint is issued in pipelined hint mode. |
| **pipelining** | A technique that breaks operations, such as instruction processing or bus transactions, into smaller stages so that a subsequent stage in the pipeline can begin before the previous stage has completed. |
| **PIR** | Processor Identification Register. |
| **PLG** | Physical layer group. |
| **PLL** | Phase-locked loop. |
| **plug-in** | Fully resolved code that is dynamically loaded and executed by running an SPU program. Plug-ins facilitate code overlays. |
| **POR** | Power-on reset. |
| **POSIX** | Portable Operating System Interface. |
| **POSIX threads library** | The standard UNIX threads library, available at http://www-106.ibm.com/developerworks/linux/library/l-posix1.html. See POSIX. |
| **PowerPC** | Of or relating to the PowerPC Architecture or the microprocessors that implement this architecture. |
| **PowerPC 970** | A 64-bit microprocessor from IBM in the PowerPC family. It supports both the PowerPC and vector/SIMD multimedia extension instruction sets. |
| **PowerPC Architecture** | A computer architecture that is based on the third generation of RISC processors. The PowerPC Architecture was developed jointly by Apple, Motorola, and IBM. |
| **PowerXCell 8i processor** | An implementation of the CBEA that uses a DDR2 memory interface and contains SPEs with enhanced double-precision performance and instructions. |
| **PP** | Page Protection bits. |
| **PPE** | PowerPC Processor Element. The general-purpose processor in the CBEA processors. It consists of the PPU and the PPSS. |

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| **PPE Pause (0)** | PPE Pause (0) state. One of six power-management states. |
| **PPSS** | PowerPC Processor Storage Subsystem (L2 cache, NCU, CIU, BIU). Part of the PPE. It operates at half the frequency of the PPU. |
| **PPU** | PowerPC processor unit. The part of the PPE that includes execution units, memory-management unit, and L1 cache. |
| **pragma** | A compiler directive, inserted in source code by the programmer. |
| **precise interrupt** | An interrupt in which the architecturally visible processor state (in the CBEA processors, this is the PPE state) is known at the time of the exception that caused the interrupt, so that the pipeline can be stopped, instructions that preceded the faulting instruction can complete, and subsequent instructions can be flushed and redispatched after interrupt handling has completed. A precise interrupt is caused by an exception that was generated when the instruction was fetched or executed. |
| **predicate** | A Boolean-logic term denoting a logical expression that determines the state of some variables. For example, a predicate can be an expression stating that variable A must have a value of 3. |
| **predictive commoning** | A form of common subexpression elimination that exploits data reuse among consecutive loop iterations. |
| **preferred scalar slot** | The preferred location for scalar values within a 128-bit register. The preferred location for 8-bit scalars is byte 3. The preferred location for 16-bit scalars is bytes 2 through 3. The preferred location for 32-bit scalars is bytes 0 through 3. The preferred location for 64-bit scalars is bytes 0 through 7. Compare preferred slot. See Figure 3-6 on page 73. |
| **preferred slot** | The left-most word (bytes 0, 1, 2, and 3) of a 128-bit register in an SPE. This is the SIMD word element in which scalar values are naturally maintained. Compare preferred scalar slot. See Figure 3-6 on page 73. |
| **prefetch** | To fetch instructions ahead of the processor's ability to dispatch them. |
| **primitive** | A simple procedure, such as a test-and-set loop used for atomic lock acquisition. |
| **privilege 1** | A designation given to SPE registers that are the most privileged of all SPE registers, used by the hypervisor to manage the SPE on behalf of a logical partition. |
| **privilege 2** | A designation given to SPE registers that are the second-most privileged of all SPE registers, used by the operating system in a logical partition to manage the SPE within the partition. |
| **privilege 3** | A designation given to SPE registers that are the least privileged of all SPE registers, used by Problem State (application) software, if direct access to the SPE from user space is supported by the operating system. |
| **privileged software** | Software that can be executed only in supervisor or hypervisor state. |

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| **privileged state** | Also known as supervisor state. The permission level of operating system instructions. The instructions are described in the PowerPC Operating Environment Architecture, Book III and are required of software the accesses system-critical resources. |
| **problem state** | The permission level of user instructions. The instructions are described in the PowerPC Architecture Books I and II and are required of software that implements application programs. Compare supervisor state. |
| **program header** | An ELF object-file header specifies program segments in the file for use in program loading. Compare section header. |
| **program order** | The order in which instructions occur in the program. Compare sequential order. |
| **PRV** | Pervasive logic. For details, see the Cell Broadband Engine CMOS SOI 65 nm Hardware Initialization Guide. |
| **pseudo-LRU** | A cache-replacement algorithm that approximates, and performs faster than, a true least-recently used replacement algorithm. |
| **PTE** | Page table entry. See page table. |
| **PTEG** | Page table entry group. |
| **push** | A write, triggered by a snooped transaction, that copies a modified cache line to main storage. Also called push out. |
| **putllc** | MFC put lock line conditional command. |
| **putlluc** | MFC put lock line unconditional command. |
| **QoS** | Quality of service. It typically relates to a guarantee of minimum bandwidth for streaming applications. |
| **quadword** | 16 bytes. |
| **R** | Reference bit. It is set by hardware whenever the page is accessed and can be used by the operating system to determine recently used pages. |
| **RA** | Real address. See real address. |
| **RAG** | Resource allocation group. |
| **RAG U** | Unallocated resource allocation group. |
| **RAID** | Resource allocation ID, or resource allocation group ID. It is the ID of a Resource Allocation Group (RAG). |
| **RAM** | Resource allocation management. A mechanism that allocates access to resource allocation groups (RAGs). Examples are the allocation of access to memory banks or I/O interfaces. |
| **range hit** | The satisfaction of all conditions specified in the Address Range Registers. |

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| **RAS** | Reliability, availability, serviceability. |
| **RAW** | Read-after-write dependency. |
| **RC** | (1) Read-and-claim. The L2 cache has six RC state machines that move data in and out of the L2 Cache in response to PPU or snoop requests. (2) The Reference (R) and Change (C) bits in a PTE. |
| **Rc** | See record bit. |
| **rchcnt** | SPU read channel counter instruction. |
| **RClassID** | Replacement class ID. |
| **RC machine** | Read and claim machine. |
| **RC update** | An updates to the Reference (R) or Change (C) bits in a PTE. |
| **rdch** | SPU read from channel instruction. |
| **real address** | The address of a byte in real storage or on an I/O device. Real storage includes physical memory, the PPE's L1 and L2 caches, and the SPE's LSs. The maximum size of the real-address space is 242 bytes. |
| **real addressing mode** | A PPE addressing mode in which address translation is disabled (MSR[IR] = '0' for instructions, MSR[DR] = '0' for data). Accesses in real mode bypass storage protection checks. The effective address is used as the real address; on the PPE (but not on the SPEs), such a real address may be offset by an RMOR or HRMOR base address. |
| **real mode** | Same as real addressing mode. |
| **record bit** | Bit 31 (the Rc bit) in a PPE instruction opcode. When set to 1, the instruction updates the Condition Register (CR) according to the operation's result. Instructions that have the record bit set a dot (or period) suffix on their mnemonic. |
| **recording** | The setting of bits in the Condition Register (CR) to reflect characteristics of an executed instruction's result. The recording is caused by instruction mnemonics that end in a period (.); such instructions cause the record (Rc) bit of the instruction format to be set to '1'. See record bit. |
| **register spill** | A situation in which the instantaneous number of active variables exceeds the size of the register file. |
| **register stream** | A sequence of contiguous registers that are produced by an operation over the lifetime of a loop. |
| **relocation** | Virtual-address translation. Instruction and data relocation can be enabled independently using two bits (IR and DR) in the Machine State Register (MSR). |
| **replacement management table** | A software-controlled cache-replacement facility used to lock entries in a cache, thus preventing their replacement. The CBEA processors provide an RMT for the PPE TLB and L2 cache, and for each SPE TLB. |

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| **reservation** | An exclusive right to access a main-storage location. Reservations are set and cleared with the lwarx and stwcx instructions, MFC atomic commands, and other instructions that access the reservation granule in which the reservation is set. Also called lock. |
| **reservation granule** | The storage block size corresponding to the number of low-order bits ignored when a store to a real address is compared with a reservation at that address. |
| **reserved** | Register locations marked as reserved and bit fields within registers marked as reserved not implemented operate in the same way: writes have no effect and reads return all ‘1’s. Reserved areas of the MMIO-register memory map that are not assigned to any functional unit should not be read from or written to: doing so will cause serious errors in soft-ware as follows. See Reserved Regions of Memory and Registers on page 30. |
| **retire** | To write the results of a completed instruction back to main storage. Compare complete. |
| **RI** | Recoverable Interrupt. A bit in the Machine State Register (MSR). |
| **RMI** | Real-mode caching inhibited. |
| **RMLR** | Real Mode Limit Register. |
| **RMO** | Real-mode offset. |
| **RMOR** | Real-Mode Offset Register. |
| **RMR** | Range Mask Register. |
| **RMSC** | Real-mode storage control. |
| **RMT** | See replacement management table. |
| **RPN** | Real page number. |
| **RSR** | Range Start Register. |
| **runout** | See instruction runout. |
| **runtime alignment** | Memory alignment cannot be known at compile time and is instead done at runtime. |
| **SBI** | Synergistic bus interface. The MFC’s interface to the EIB. |
| **scalar** | A single data item, as opposed to a set of data items such as a vector or array. |
| **scatter-gather** | A technique for operating on sparse data, using an index vector. A scatter-gather operation takes an vector and fetches data at an address added to that of the vector. A scatter operation stores data back to memory, using the same index vector. |

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| **scheduling** | A compiler optimization that reorders the instruction sequence subject to data-flow and control-flow restrictions so as to maximize use of a processor's hardware. |
| **SCN** | SPU control unit. A unit in SPU that handles branches and program control. |
| **sdcrf** | SL1 data-cache-range flush. |
| **sdcrt** | SL1 data-cache-range touch. Implemented as a nop on the CBEA processor. |
| **sdcrtst** | SL1 data-cache-range touch-for-store. Implemented as a nop on the CBEA processor. |
| **sdcrz** | SL1 data-cache-range set-to-zero. |
| **SDR** | Storage Descriptor Register. |
| **section header** | An ELF object-file header specifies sections in the file for use in linking. Compare program header. |
| **segment** | A fixed 256 MB unit of address space that can hold code, data, or any mixture thereof. Segments are controlled by the segment lookaside buffer (SLB), which maps EAs to VAs and provides protections. |
| **semaphore** | A software flag used to indicate the status of, and lock the availability of, a shared resource. Semaphores can be implemented with atomic operations. |
| **sequential order** | The order in which the compiler output of a program appears in main storage. Compare program order. |
| **serialization** | A hardware-enforced alteration of the processor state or execution pipeline so as to match the sequential ordering of instructions in a program. For example, an instruction is said to be serializing if it (a) causes all preceding instructions, in program order, to complete before it begins execution, and (b) completes execution before any following instructions, in program order, begin execution. Compare synchronization. See also program order, sequential order, and context synchronizing. |
| **set** | A row of a set-associative cache, including a TLB cache. Also called a congruence class. Compare way. |
| **SFP** | SPU floating-point unit. It handles single-precision and double-precision floating-point operations. |
| **SFS** | SPU odd fixed-point unit. It handles shuffle operations. |
| **SFX** | SPU even fixed-point unit. It handles arithmetic, logical, and shift operations. |
| **short-loop aggregation** | The elimination of inner loops that have short trip counts. SIMDizable short loops can be collapsed into vector operations. |

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| **signal** | Information sent on a signal-notification channel. These channels are inbound (to an SPE) registers. They can be used by the PPE or other processor to send information to an SPE. Each SPE has two 32-bit signal-notification registers, each of which has a corresponding memory-mapped I/O (MMIO) register into which the signal-notification data is written by the sending processor. Unlike mailboxes, they can be configured for either one-to-one or many-to-one signalling.<br><br>These signals are unrelated to UNIX signals. See *channel* and *mailbox*. |
| **signal notification** | See signal. |
| **SIMD** | Single instruction, multiple data. Processing in which a single instruction operates on multiple data elements that make up a vector data-type. Also known as vector processing. This style of programming implements data-level parallelism. |
| **SIMDize** | Transform scaler code to vector code. |
| **single-ported** | Single-ported memory allows only one access at a time. |
| **SL1** | A first-level cache for DMA transfers between LS and main storage. The SPEs in the CBEA processors do not implement SL1s. |
| **SLB** | Segment lookaside buffer. It is used to map an effective address (EA) to a virtual address (VA). |
| **SLBE** | SLB entry. |
| **slbia** | SLB invalidate all instruction. |
| **slbie** | SLB invalidate entry instruction. |
| **slbmfee** | SLB move from entry ESID X-form instruction. |
| **slbmfev** | SLB move from entry VSID X-form instruction. |
| **slbmte** | SLB move to entry X-form instruction. |
| **SLIH** | Second-level interrupt handler. |
| **slow mode** | Slow state. One of six power-management states. |
| **SLS** | SPU load and store unit. It handles loads, stores, and branch hints. |
| **SMD** | Slow-mode divider. A clock-divider value in PMSR1[BE_Slow] that helps specify the maximum time-base frequency. |
| **SMM** | Synergistic memory management unit. It translates EAs to RAs for in an SPE. |
| **SMP** | Symmetric Multiprocessor or Symmetric Multiprocessing. Two CBEA processors can be connected together to form an SMP system. |
| **sndsig** | Send signal command. |

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| **sndsigb** | Update signal-notification registers in an I/O device or another SPE with barrier command. |
| **sndsigf** | Update signal notification registers in an I/O device or another SPE with fence command. |
| **snoop** | To compare an address on a bus with a tag in a cache, to detect operations that violate memory coherency. |
| **snoop push** | See push. |
| **SNP** | Snoop. |
| **SOA** | Structure of arrays. A method of organizing related data values. Also called parallel-array form. See AOS. |
| **software cache** | Same as software-managed cache. |
| **software-managed cache** | An SPE's local storage (LS), which is filled from main memory using software-initiated DMA transfers. Although most processors reduce latency to memory by using hardware caches, an SPE uses its DMA-filled LS. The approach provides a high degree of control for real-time programming. However, the approach is advantageous only if the DMA transfer-size is sufficiently large and the DMA command is issued well before the data is needed, because the latency and instruction overhead associated with DMA transfers exceeds the latency of servicing a cache miss on the PPE. |
| **software-managed memory** | Same as software-managed cache. |
| **software pipelining** | A loop optimization in which the body of the loop is divided into a series of stages that are executed in parallel in a manner analogous to hardware pipelining. |
| **southbridge** | A chip that interfaces a processor to I/O buses (except, typically, graphics). |
| **SP** | (1) Single-precision. (2) Stack pointer. |
| **SPE** | Synergistic Processor Element. It includes an SPU, an MFC and an LS. In this document, the term "SPE" refers generally to functionality of any part of the processor element, including the MFC, and the term "SPU" refers to the instruction set or the unit that executes the instruction set. |
| **specific intrinsic** | A type of C and C++ language extension that maps one-to-one with a single SPU assembly instruction. All SPU specific intrinsics are named by prefacing the SPU assembly instruction with si_. |
| **SPE SRI** | SPE State Retained and Isolated (SRI) state. One of six power-management states. |

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| **SPE thread** | A thread scheduled and run on an SPE. A program can have one or more SPE threads. Each thread has its own SPU local storage (LS), register file, program counter, and MFC command queues. |
| **spin lock** | A synchronization primitive for shared-storage environments that executes a tight loop, attempting to acquire a lock at each iteration. |
| **splat** | To replicate, as when a single scalar value is replicated across all elements of a SIMD vector. |
| **SPR** | Special purpose register. |
| **SPU** | Synergistic processor unit. The part of an SPE that executes instructions from its local storage (LS). In this document, the term “SPU” refers to the instruction set or the unit that executes the instruction set, and the term “SPE” refers generally to functionality of any part of the processor element, including the MFC. |
| **SPU ISA** | SPU Instruction Set Architecture. A SIMD instruction set executed in SPEs that is similar to the vector/SIMD multimedia extension instruction set executed by the PPE. |
| **SPU Pause** | SPU Pause state. One of six power-management states. |
| **SRR0** | Save and Restore Registers 0. |
| **SRR1** | Save and Restore Registers 1. |
| **SSC** | SPU channel and DMA unit. It handles all input and output functions for an SPU. |
| **SSE** | Streaming SIMD extensions. An Intel® instruction set. |
| **stale** | A value for a parameter is said to be stale if it is not the current value for that parameter. |
| **stall** | (1) The inability for an instruction in a pipeline to proceed. Possible causes of the stall include occupation of the next pipeline stage by another instruction, waiting for operands, or serialization. Stalls occur at the instruction-issue stage and, in the PPE, stop both threads. Compare block and bubble. (2) The number of cycles, after completing the execution of an instruction of a given type, before another instruction of the same type can be issued. |
| **starvation** | A condition in which a processing element is making forward progress, but at an extremely slow rate. |
| **statically built** | Built at compile time. |
| **static branch prediction** | A method in which software (for example, a compiler) gives a hint to the processor about the direction that a branch is likely to take. |
| **static linking** | The linking of procedures at compile time, rather than at link time or at load time. |

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| **stdcx** | PPE store doubleword conditional indexed instruction. |
| **sticky bit** | A bit that is set by hardware and remains so until cleared by software. |
| **stmw** | PPE store multiple word instruction. |
| **storage** | Any type of data storage, including memory, cache, disk, diskette, and tape. Compare main storage. |
| **storage access** | An access to main storage caused by a load, a store, a DMA read, or a DMA write. |
| **storage-access recording** | The setting of the PTE Reference (R) bit to '1' whenever an attempt is made to read or write the page, and the Change (C) bit to '1' whenever the page is written. The recordings are done by hardware. |
| **store combining** | The combining of multiple store accesses to storage into a single store access. |
| **STQ** | Store queue. The CBEA processors have several. |
| **stream** | A sequence of contiguous memory locations that are accessed by a memory reference throughout the lifetime of a loop (also called a memory stream), or a sequence of contiguous register values that are produced by an operation over the lifetime of a loop (also called a register stream). |
| **stream offset** | The byte-offset of the first required value in the first register of a stream. |
| **stride** | The relationship between the layout of an array's elements in main storage and the order in which those elements are accessed. A stride of length N means that, for each array element accessed, N-1 adjacent memory elements are skipped over before the next-accessed element. |
| **stride-one memory access** | A memory access pattern in which each element in a list is accessed sequentially. |
| **stswi** | PPE store string word immediate instruction. |
| **stswx** | PPE store string word indexed x-form instruction. |
| **stwcx** | PPE store word conditional indexed instruction. |
| **subword arithmetic operations** | The conversion of integer arithmetic operations to the equivalent arithmetic operation on narrower-width data types. |
| **subword data type** | A 2-byte or 1-byte data type, both of which are smaller than a word. |
| **superscalar** | The ability to execute multiple instructions per cycle. It is accomplished with multiple, parallel execution units. |
| **supervisor state** | The permission level of privileged instructions. The instructions are described in the PowerPC Operating Environment Architecture, Book III and are required of software that implements system programs. Compare problem state. |

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| **SXU** | Synergistic execution unit. It contains the SPU odd fixed-point unit (SFS), SPU even fixed-point unit (SFX), SPU odd floating-point unit (SFP), and SPU load and store unit (SLS). |
| **symmetric multiprocessing** | See SMP. |
| **sync** | A PPE or SPE synchronize instruction. |
| **synchronization** | (1) Storage-access ordering for shared-storage environments. (2) A software-enforced alteration of processor state so as to match the sequential ordering of instructions. (3) The use of atomic operations to create semaphores, mutex locks, spin locks, and other synchronization primitives for shared-storage environments. Compare serialization. See also sequential order, program order, and context synchronizing. |
| **synchronous** | Coordinated in time, with the execution of instructions in an instruction pipeline or among tasks. |
| **system** | A combination of processors, storage, and associated mechanisms that is capable of executing programs. |
| **system memory** | See main storage. |
| **system storage** | All program-addressable memory in a system, including main storage (main memory), the PPE's L1 and L2 caches, and the SPE's local storage (LS). See main storage. |
| **table walk** | Table look-up (table search). |
| **tag group** | A group of MFC DMA commands. All DMA commands except getllar, putllc, and putlluc are associated with a Tag Group. |
| **taken** | (1) Said of an interrupt whose service routine is executed. (2) Said of a conditional branch when the condition it is testing is true. |
| **task** | A process (unit of resource ownership) in a multiprogramming (multitasking) environment. A task owns a virtual address space in which it stores processor state, and it may own other resources such as protected access to other processes, I/O devices and files. |
| **TClassID** | Transfer class ID. |
| **text segment** | A segment of programming code. |

**thread** (1) A unit of operating-system scheduling and dispatching that executes sequentially and can be interrupted. Threads are created by processes (tasks), which might own one or more of them, and threads use the resources of the creating process. A thread can be running or waiting to be run.

(2) A sequence of instructions executed within the global context (shared memory space and other global resources) of a process that has created (spawned) the thread. Multiple threads (including multiple instances of the same sequence of instructions) can run simultaneously, if each thread has its own architectural state (registers, program counter, flags, and other program-visible state).

The SPE hardware supports a single thread, per SPE. The PPE hardware supports two threads.

**throughput** (1) The number of instructions completed per cycle. A high-throughput application design seeks to keep pipelines full. (2) The maximum sustained rate at which a processor can execute an instruction of a particular type, in the absence of any dependencies and assuming infinite caches.

**time base** The facility that provides the timing functions for the processor core-clock (`NClk`) domain.

**TIS** Tool Interface Standard.

**TKM** Token management unit. Part of the element interconnect bus (EIB) that software can program to regulate the rate at which particular devices are allowed to make EIB command requests.

**TLB** Translation lookaside buffer. An on-chip cache that translates virtual addresses (VAs) to real addresses (RAs). A TLB caches page-table entries for the most recently accessed pages, thereby eliminating the necessity to access the page table from memory during load/store operations.

**TLB hit** A TLB access in which the corresponding entry is present and is valid.

**tlbie** PPE TLB invalidate-entry instruction.

**TLBIE request** A request for a TLB-entry invalidation that is generated by an SPE's atomic unit and broadcast over the EIB.

**tlbsync** PPE TLB-synchronization instruction.

**TMCU** Thermal management control unit.

**.toe** A section in a CESOF object file that contains a TOE. It is generated by the SPU linker and is not seen by application programmers.

**TOE** Table of effective-address references (EARs). It is used to resolve references from an SPU's address space to the PPU's symbols.

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| **TOE shadow** | An exact copy of a toe segment. In the PPE-ELF object, the TOE shadow is defined in a .data section. |
| **token** | A grant of access to the EIB. When the token manager function is enabled, the SBI will request a token for any memory or I/O-bus transaction. After the token is granted, the SBI will send the bus request to the EIB. |
| **touch** | To cause a cache block to be speculatively loaded. The PowerPC Architecture supports instructions that perform this function. |
| **transfer element** | Same as list element. |
| **trap instruction** | An instruction that tests for a specified set of conditions. If the conditions of a trap instruction are met, the program interrupt trap handler is invoked. |
| **trip count** | The number of iterations in a loop. |
| **TS** | The transfer-size parameter in an MFC command. |
| **unified cache** | A cache that stores both instructions and data. |
| **unified register file** | A register file in which all data types—integer, single-precision and double-precision floating-point, logicals, bytes, and so forth—use the same register file. The SPEs (but not the PPE) have unified register files. |
| **unordered** | Said of an exception that may generate an interrupt at any time, regardless of the state of the interrupt-processing mechanism. |
| **unroll** | See loop unrolling. |
| **UPC** | Unified Parallel C. A a parallel-programming extension to the ANSI C language. |
| **update** | The action, by a load or store instruction, of automatically copying the target address computed by the instruction into the base register used for the address computation. Update instructions are useful for moving repetitively through data structures. |
| **user mode** | The mode in which problem state software runs. See problem state. |
| **V** | Valid. |
| **VA** | Virtual address. |
| **valid** | A page-table entry or TLB entry that is allocated and correctly associated with the object of its reference. |
| **vector** | (1) An instruction operand that consists of a set of data elements packed into a one-dimensional array. The elements can be fixed-point or floating-point values. Most vector/SIMD multimedia extension and SPU SIMD instructions operate on vector operands. Vectors are also called SIMD or packed operands. (2) See interrupt vector. |

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| **vector/SIMD multimedia extension** | The SIMD instruction set of the PowerPC Architecture, supported on the PPE. Also known as AltiVec, which is a Freescale trademark. |
| **VIQ** | VSU issue queue. |
| **virtual address** | An address to the virtual-memory space, which is much larger than the real address space and includes pages stored on disk. It is translated from an effective address (EA) by a segmentation mechanism and used by the paging mechanism to obtain the real address (RA). The maximum size of the virtual-address space is 265 bytes. |
| **virtual memory** | The address space created using the memory management facilities of a processor. |
| **virtual mode** | The mode in which virtual-address translation is enabled. |
| **virtual vector** | A compiler construct that has no alignment constraints and can have any length. Compilers use virtual vectors to carry out preliminary SIMD transformations. |
| **volatile register** | A register designated by an ABI as unnecessary to save across procedure calls. Also called a caller-save register. |
| **VPN** | Virtual page number. The number of a page in virtual memory. |
| **VSID** | Virtual segment ID. |
| **VSU** | Vector scalar unit. In the PPE, the combination of the VXU and FPU. |
| **VXU** | Vector/SIMD multimedia extension unit. |
| **W** | The write-through bit in a page table entry which controls the processor's accesses to cache and main storage. It is part of the “WIMG” bits. |
| **WAW** | A write-after-write dependency. |
| **way** | A column of a set-associative, multi-way cache, including a TLB cache. Compare set. |
| **weakly consistent storage order** | A memory-access model in which the order of processor storage accesses, the order of those accesses with respect to another processor or mechanism, and the order of those accesses in main storage, may be different. |
| **WIMG bits** | Four bits in a page table which control the processor's accesses to cache and main storage—“W” is write-through, “I” is caching-inhibited, “M” is memory-coherence, and “G” is guarded. |
| **word** | Four bytes. |
| **workload** | A set of code samples that characterizes the performance of the architecture, algorithms, libraries, tools, and compilers. |
| **WPC** | Write-port collision. |

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| **wrch** | SPU write to channel instruction. |
| **writeback** | A data-cache property that allows modified data to be written only to the cache, rather than also to main storage. The modified data is written to main storage only when the cache line is replaced. Compare write through |
| **writeback flag** | A flag written (for example, by an SPE) to main storage that notifies another processor (for example, the PPE) of an event. |
| **write through** | A data-cache property that requires modified data to be written not only to the cache but also to main storage. Compare writeback |
| **Write With Clean** | A burst operation caused by a processor executing a dcbst instruction or a bus snoop read or clean to a modified block. It is used to tell all lower-level caches that a copy still remains in this level, while updating memory or I/O. |
| **Write With Flush** | A partial-block write to memory or a sub-block burst operation from the I/O. It is used for caching-inhibited or write-through writes from a processing element. |
| **Write With Kill** | A burst operation used to tell all snoopers to invalidate any copies of this cache line in their caches, while also storing the line to memory. |
| **X2D** | PowerXCell 8i logic block that converts XDR packets to DDR2 commands. |
| **XDR** | Rambus Extreme Data Rate DRAM technology. |
| **XIO** | Rambus XDR I/O (XIO) cell. |
| **XLAT** | Translate request queue. |
| **XU** | Execution unit. The PPE unit that contains the FXU, LSU, and MMU. |
| **zero-shift policy** | A reorganization of misaligned data in which each misaligned register stream is shifted to a stream offset of 0 immediately after it is loaded from memory. |

# Index

## Symbols



## Numerics



## A



## B

## C

## D



## E

## F



## G



## H, I, J, K

## L

## M

## N



## O



## P

## Q

## R



## S

## T



## U, V, W

## X



## Z